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NUMERICAL OPTIMIZATION ALGORITHM FOR ENGINEERING PROBLEMS USING MICROCOMPUTER

by

Dong Soo, Kim

September 1984

Thesis Advisor:

G. N. Vanderplaats

21

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	A general purpose computer p	· · ·	oped to perform						
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	The program is developed especia	ally for use on 1	microcomputers and						
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Typical applications of MSCOP program are in the design of machine components and simple beam and truss structures. Solutions to three sample problems are given. Approved for public release; distribution unlimited.

Numerical Optimization Algorithm for Engineering Problems Using Micro-computer

by

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Submitted in partial fulfillment of the requirements for the degree of

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ABSTRACT

A general purpose computer program is developed to perform nonlinear constrained optimization of engineering design problems. The program is developed especially for use on microcomputers and is called Microcomputer Software for Constrained Optimization Problems (MSCOP). It will accept a nonlinear objective function and up to 50 inequality constraint functions and up to 20 hounded design variables.

MSCOP employs the method of feasible directions. Although developed for microcomputers, for speed of development, the MSCOP was implemented on an IBM 3033 using standard basic language, Waterloc BASIC Version 2.0. It is directly transportable to a variety of microcomputers.

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I. <u>INTRODUCTION</u>

A. PURPOSE

This thesis describes the development of a microcomputer oriented program called MSCOP (Microcomputer Software for Constrained Optimization Problems) for constrained optimization of engineering design problems. Problems which can be solved by the MSCOP are nonlinear programming problems arising in several areas of machine and structural design, such as the minimum weight design of structures subject to stress and displacement constraints [Ref. 1].

In recent years, several powerful general purpose optimization programs have become available for engineering design problems, e.g., COPES/CONMIN [Ref. 2], and ADS-1 These programs can handle a wide range of design [Ref. 3]. problems and contain a variety of solution techniques. Also, several programs are available that include optimization in an integrated analysis / design code, e.g., ACCESS, ASOP, EAI, PARS, SAVES, SPAR, STARS and TSO [Ref. 4]. All of the above optimization programs are written in FORTRAN, and are built for use on a mainframe computer. Their use can be cumbersome, especially for the occasional user. Since many engineers are now using microcomputers, there is a need to develop an optimization program contained in a microcomputer software package for use on microcomputers. This thesis fills that need by developing a compact program written in a standard BASIC language suitable for a wide range of microcomputers.

B. IPPLEMENTATION

The nature of an optimization program depends on the computer and programming method available. The MSCOP software is designed for use on a microcomputer. However, for the speed of development and testing, MSCOP was developed on the IEM 3033 computer at the W. R. Church Computer Center in Naval Postgraduate School, and was written in WEASIC (Waterloo Basic) Version 2.0.

To make sure that the program is easily portable to a microcomputer, only standard BASIC commands and functions are used. For example, FOR I = 1 TO NDB ... NEXT I, GOSUB etc., were used. The commands and functions not available in all variations of EASIC are avoided, for example, TRN(A), MAT(A), etc.

MSCOF provides design engineers with a convenient tool for optimization of engineering design problems with up to 20 bounded design variables and as many as 50 inequality constraints.

C. GENERAL OPTIMIZATION MODEL

The general optimization problem to be solved is of the form : Find the set of design variables X that will

Minimize $F(\underline{X})$ (1.1)

Subject to $G_{i}(\underline{X}) < 0$ $j = 1, \dots, m$ (1.2)

$$X_{i}^{\perp} < X_{i} < X_{i}^{\perp}$$
 $i = 1, ..., n$ (1.3)

where X is referred to as the vector of design variables. $F(\underline{X})$ is the objective function which is to be minimized. $G(\underline{X})$ are inequality constraint functions, and X_{i}^{2} and X_{i}^{2} are lower and upper bounds, respectively, on the design

variables. Although these bounds or "side constraints" could be included in the inequality constraint set given by Eq(1.2), it is convenient to treat them separately because of their special structure. The objective function and constraint functions may be nonlinear, explicit or implicit in X. However, they must be continuous and should have continuous first derivatives.

In general engineering optimization problems, the objective to be minimized is usually the weight or volume of a structure being designed while the constraints gives limits on compressive stress, tensile stress, Euler buckling, displacement, frequencies (eigenvalues), etc. [Ref. 5 : p.264]. Equality constraints are not included because their inclusion complicates the solution techniques and because in engineering situations, equality constraints are rare.

Most optimization algorithms require that an initial value of design variables X^o be specified. Beginning from these starting values, the design is iteratively improved. The iterative procedure is given by

 $\underline{\mathbf{x}}^{\mathbf{q}+1} = \underline{\mathbf{x}}^{\mathbf{q}} + \mathbf{a}^{\mathbf{x}} \underline{\mathbf{s}}^{\mathbf{q}}$ (1.4)

where q is the iteration number, S is a search direction vector in the design space, and a* is a scalar parameter which defines the amcunt of change in \underline{X} . At iteration q, it is desirable to determine a direction \underline{S} which will reduce the objective function (usable direction) without violating the constraints (feasible direction). After determining the search direction, the design variables, \underline{X} , are updated by Eq (1.4) so that the minimum objective value is found in this direction. [Ref. 6].

Thus, it is seen that nonlinear optimization algorithms for the general optimization problem based on Eq(1.4) can be separated into two parts, determination of search direction and determination of scalar parameter a*.

D. ORGANIZATION OF THIS THESIS

This chapter has stated the purpose of the thesis and has put the general concept of engineering optimization into a preliminary perspective. Chapter 2 will describe the essential aspects of the optimization algorithm used in MSCOP such as finding a search direction, the onedimensional search and convergence criteria. Chapter 3 describes program usage. In chapter 4, there are three examples which are sclved by the MSCOP. Summary and conclusions are given in chapter 5. The program is listed in the appendix.

II. CPTIMIZATION ALGORITHM

A. INTRODUCTION

There are many optimization algorithms for constrained nonlinear problems such as generalized reduced gradient method, feasible direction method, penalty function methods, Augmented Lagrangian multiplier method, and sequential linear programming. The feasible direction method is chosen for development in this thesis for three main reasons. First it progresses rapidly to a near optimum design. Second it only requires gradients of objective and constraint functions that are active at any given point in the optimization process [Ref. 7]. Third, because it maintains a feasible design, engineer cannot fail to meet safety requirements as defined by the contraints. However, the method does have several disadvantages in that it is prone to "zig-zag" between constraint boundaries and that it is usually does not achieve a precise optimum. This method solves the nonlinear programming problem by moving from a feasible point (can be initially infeasible) to another feasible point with an improved value of the objective value.

The following strategy is typical of feasible direction method : Assuming that an initial feasible point X° is known, first find a usable-feasible direction S. The algorithm for this is similar to linear programming and complementary pivoting algorithms. Having found the search direction, a move is made in this direction to update the X vector according to Eq(1.4). The scalar a* is found by a one-dimensional search to reduce the objective function as much as possible subject to constraints. That is MIN

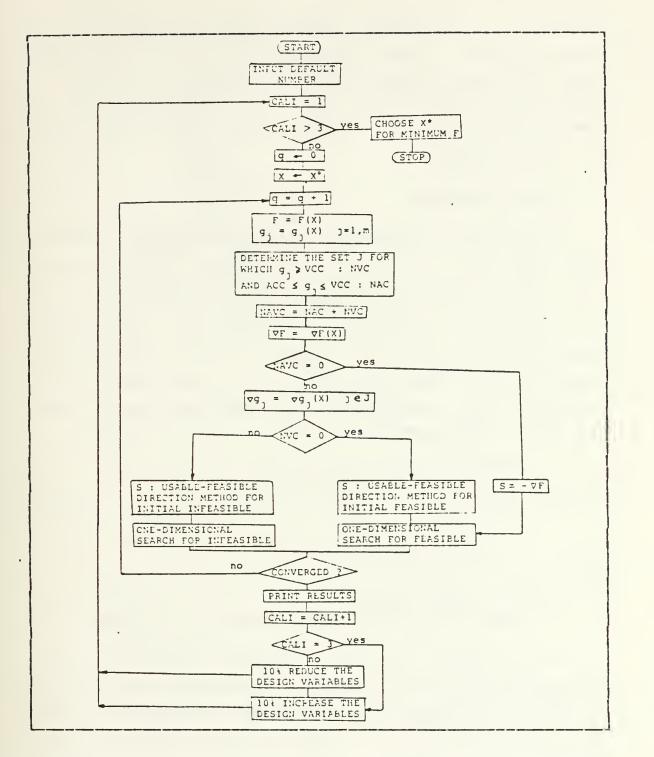


Figure 2.1 Algorithm for the Feasible Direction Method.

F(X+a*S) subject to $G(X+a*S) \leq 0$. It is assumed that the initial design X° is feasible, but if it is not, a search

direction is found which will direct the design to the feasible region. After updating the X° vector, the convergence test must be performed in the iterative algorithm. A convergence criteria used in this is implementation are described in section D. The general algorithm used in MSCOP is given in Figure 2.1

B. SEARCH DIRECTION

In the feasible direction algorithm, a usable - feasible search direction S is found which will reduce the objective function without violating any constraints for some finite move. It is assumed that at any point in the design space (at any \underline{X}) the value of the objective and constraint functions as well as the gradients of these functions with respect to the design variables can be calculated. Since these gradients cannot usually be calculated analytically, the finite difference method Eq(2.1) is used in MSCOF.

$$\frac{\partial F(\underline{X})}{\partial \underline{X}_{i}} = \frac{F(\underline{X} + \varepsilon e_{i}) - F(\underline{X})}{\varepsilon}$$
(2.1)

where e is the ith unit vector

Eis a small scalar.

In MSCOP, ε is 0.1% of the ith design variable

In the feasible direction algorithm, there are usually one or more "active" constraints. A constraint $G(\underline{X}) \leq 0$ is "active" at \underline{X} if $g(\underline{X}) \approx 0$. As shown in Figure 2.1, if no constraints are active the standard steepest descent direction $\underline{S} = -\nabla F$ is used.

1. Usable-Feasible Direction

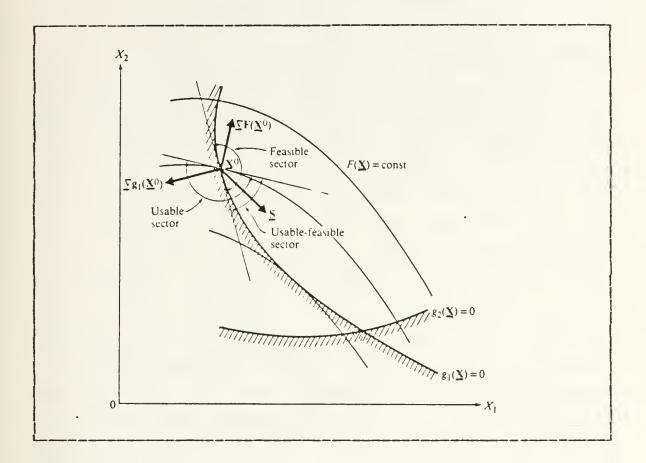


Figure 2.2 Usable-Feasible Direction.

Assume there are NAC active constraints at \underline{X} . The direction \underline{S} is "usable" if it reduces the objective function, i.e.,

$$\nabla F \cdot S < 0 \tag{2.2}$$

Similarly the direction is feasible if for a small movement in this direction, no constraint will be violated, i.e.,

$$\nabla G \cdot S < 0 \tag{2.3}$$

This is shown geometrically in Figure 2.2

2. Active Constraints

It is necessary to determine if a constraint is active or violated in the feasible direction algorithm. A constraint $G(\underline{X}) \leq 0$ is "active" at \underline{X}° if $G(\underline{X}^{\circ}) \approx 0$. In order to avoid the zigzagging effect between one OT more constraint boundaries, a tolerance band about zero is used for determining whether or not a constraint is active. From the engineering point of view, a constraint $G(\underline{X}) \leq 0$ is active near the boundary $G(\underline{X}) = 0$ whenever ACC $\leq G(\underline{X}) \leq VCC$. ACC is the active constraint criterion and VCC is the violated constraint criterion in MSCOP. Assuming the feasible constraints are normalized so that G(X) ranges between -1 and 0 for reasonable values of X, the constraint $G(X) \leq 0$ is considered active if $G(X) \geq -0.1$. The constraint is considered to be violated if G(X) > 0.004. This is an algorithmic trick which improves efficiency and reliability of the algorithm. However, since in the one dimensional search, all interpolations for constraint $G(\underline{X})$ are done for zeros of a linear or quadratic approximation to $G(\underline{X})$ in order to find a*, at the optimum the value of active constraints are very near zero, but may be as large as 0.004 [Ref. 6]. From an engineering point of view, a 0.4 % constraint violation is considered to be acceptable.

3. Suboptimization Problem and Push-Off Factors

Zoutendijk [Ref. 8] has shown that a usable feasible direction S may be found as follows :

> Maximize β (2.4) Subject to ;

> > $\mathbf{\nabla}\mathbf{F}(\underline{X})\cdot\mathbf{S} + \mathbf{\beta} < 0 \tag{2.5}$

$$\nabla G(X) \cdot S + \Theta_{j} \beta < 0 \quad j \in J \quad (2.6)$$

Where scalar β is a measure of the satisfaction of the usability and feasibility requirements. The scalar θ_j in Eq (2.6) is referred to as the "push-off" factor which effectively pushes the search direction away from the active

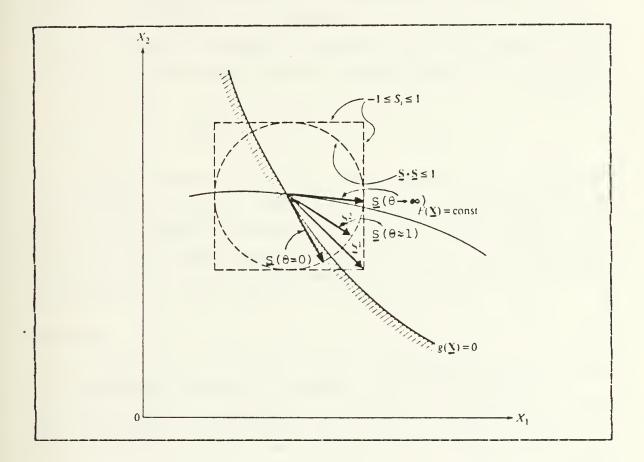


Figure 2.3 Push-Off Factor and Bounding of the S-Vector.

constraints. In Eq (2.6), if the push-off factor is zero, the search direction is tangent to the active constraints, and if it is infinite, then the search direction is tangent to the objective function. It has been found that a push-off factor is defined as follows gives good results
[Ref. 5: p.167]:

$$\Theta_{j} = \left[1 - \frac{G_{j}(X)}{ACC} \right]^{2} \Theta_{o}$$
 (2.8)

where $\theta_{1} = 1$.

To avoid an unbounded solution when seeking a usable - feasible direction it is necessary to impose bounds on the search direction §. Cne method of imposing bounds on search direction is to impose bounds on the components of S-vector of form :

$$-1 < s_i < 1$$
 (2.9)

This choice of bounding the S-vector actually biases the search direction. This is undesirable since we wish to use the push-off factors as our means of controlling the search direction. A method which avoids this bias in search direction is the circle as shown Figure 2.3. The norm here is

$$\underline{S} \cdot \underline{S} < 1$$
 (2.9.1)

4. Simple Simplex-like Method for Search Direction

Vanderplaats [Ref. 5: pp.168-169] provides the matrix formulation which solves the above sub-optimization problem by using the Zoutendijk method.

Maximize $\underline{P} \cdot \underline{y}$ (2.10) Subject to ; $\underline{A} \cdot \underline{y} < 0$ (2.11)

$$\underline{y} \cdot \underline{y} < 1 \tag{2.12}$$

Where

$$\underline{\mathbf{y}} = \begin{bmatrix} \underline{\mathbf{S}}_{1} \\ \underline{\mathbf{S}}_{2} \\ \vdots \\ \underline{\mathbf{S}}_{n} \\ \mathbf{g} \end{bmatrix} \qquad \underline{\mathbf{P}} = \begin{bmatrix} \mathbf{0} \\ \mathbf{0} \\ \vdots \\ \mathbf{0} \\ \mathbf{1} \end{bmatrix} \qquad (2.13)$$

$$\underline{A} = \begin{pmatrix} \nabla^{T} G_{1}(X), & \theta_{1} \\ \nabla^{T} G_{2}(X), & \theta_{2} \\ \vdots & \vdots \\ \nabla^{T} G_{j}(X), & \theta_{j} \\ \nabla^{T} G_{j}(X), & \theta_{j} \\ \nabla^{T} F(X), & 1 \end{pmatrix}$$
(2.14)

and where j is the number of active constraints (NAC)

When the solution to Eq(2.10) through (2.12) is found, S may be normalized to some value other than unity, but the form of the normalization is the same. A solution to the above problem may be obtained by solving the following system derived from the Kuhn-Tucker conditions for that problem :

$$\begin{bmatrix} B & I \end{bmatrix} \begin{bmatrix} u \\ v \end{bmatrix} = \underline{c}$$
 (2.15)

Where

₿ =		(2.17)
I =	Identity matrix	(2.18)
<u>c</u> =	- <u>A</u> ·₽	(2.19)

Above system can be solved using a complimentary pivot algorithm. Choose an initial basic solution to Eq(2.15) is to be

$$\underline{\mathbf{v}} = \underline{\mathbf{c}}, \qquad \underline{\mathbf{u}} = \mathbf{0} \tag{2.20}$$

where \underline{v} is the set of basic variables and \underline{u} is the set of nonbasic variables. If all $v_i > 0$, Eq(2.16) is also satisfied and problem is solved. If some $v_i < 0$, the solution procedure is as follows :

Let B_{ii} be the diagonal element of the i-th nonbasic variable.

- Given the condition that some c is less then zero, we find max (c;/B;:) which is the incoming row to the basis.
- The incoming column is changed to a tasic column, the tableau is updated by a standard simplex pivot on B_{ii}.
- 3. Until all $c_i > 0$, repeat steps 1. and 2.
- 4. When all $c_i > 0$, the iteration is complete. The value of u is now the desired solution.
- 5. By using $\underline{y} = \underline{p} \underline{A}^T \cdot \underline{u}$, we get the usable-feasible search direction S which is first NDV components of y.
 - 5. Initially Infeasible Designs

The method of feasible directions assumes that we begin with a feasible design and feasibility is maintained throughout the optimization process. If the initial design is infeasible, then a search direction pointing toward the feasible region can be found by a simple modification to direction finding problem.

A design situation can exist in which the violated constraints are strongly dependent on part of the design variables, while the objective function is primarily dependent on the other design variables. This suggests a method for finding a search direction which will' simultaneously minimize the objective while overcoming the constraint violations. These considerations lead to the following statement of the direction finding problem [Ref. 5 : pp.171-172]:

Maximize $- \nabla F(\underline{X}) \cdot \underline{S} + \underline{\Phi} \underline{\beta}$ (2.21) Subject to ; $\nabla G(\underline{X}) \cdot \underline{S} + \underline{\Theta}_{j} \underline{\beta} \leq 0$ $j \in J$ (2.22)

 $5 \cdot 5 < 1 \tag{2.23}$

where J is the set of active and violated constraints, and where the scalar \mathbf{F} in Eq(2.21) is a weighting factor determining the relative importance of the objective and the constraints. Usually a value of $\mathbf{F} > 10000$ will ensure that the resulting S-vector will point toward the feasible region. Incorporating Eq(2.21) and Eq(2.22) into the direction finding algorithm requires only that we modify the p-vector given in Eq(2.24) and the A-matrix of Eq(2.25).

$$P = \begin{pmatrix} - \nabla F(\underline{X}) \\ \overline{\Phi} \end{pmatrix}$$
(2.24)

$$\underline{A} = \begin{pmatrix} \underline{\nabla}^{\mathrm{T}} \mathbf{G}_{1} (\mathbf{X}), & \boldsymbol{\theta}_{1} \\ \underline{\nabla}^{\mathrm{T}} \mathbf{G}_{2} (\mathbf{X}), & \boldsymbol{\theta}_{2} \\ \vdots & \vdots \\ \underline{\nabla}^{\mathrm{T}} \mathbf{G}_{j} (\mathbf{X}), & \boldsymbol{\theta}_{j} \end{pmatrix}$$

$$(2.25)$$

$$\boldsymbol{\theta}_{j} \leq 50$$

$$(2.26)$$

We use the simple simplex-like method to find the search direction toward the feasible region.

C. ONE-DIMENSIONAL SFARCH

1. No Violated Constraints

If no constraints are violated, we find the largest a* in Eq(1.4) from all possible values that will minimize the objective on S without violating any constraints, active or inactive.

The procedure in MSCOF is as follows :

- Let a0, a1, a2, a3 be the scalar in Eq(1.4) ccrresponding to pcints X0, X1, X2, X3, X4.
- 2. aC = 0 at given point <u>XO</u>.
- 3. In order to get a1, we can calculate the a1 to reduce the objective by at most 10% or to change each of the design variable $\frac{X}{2}$ by at most 10%.
- 4. Update the design variables to X1 using Eq(1.4).
- 5. Evaluate the objective for $\underline{X1}$, and check the feasibility. If one or more constraints is violated, then a1 is reduced to a1/2, and we go to step 4.
- 6. In order to estimate a2, we can use the quadratic approximation with 2 points X, $\underline{X1}$ and the \underline{VF} .

- Update the design variables to <u>X2</u> by Eq(1.4) and check the side constraints.
- 8. Evaluate the objective and constraints.
- 9. Now having 3 a's, and values of objectives and constraints for design variables <u>X0</u>, <u>X1</u>, <u>X2</u> are known, so by using 3-point quadratic approximation, a value of a3 is found.
- Update the new optimal point in search direction by Eq(1.4).
- 11. Evaluate the objective and constraints.
- 12. Now choose last 3 values, a1, a2, a3 and find a new a3 using 3-points Quadratic approximation
- 13. Choose the a* among the 5 points which corresponds to the minimum objective function value with no-viclated constraints.
 - 2. One or More Constraints Violated

If one or more constraints are initially violated, a modified usable-feasible direction is found. It is then necessary to find the scalar a* in Eq(1.4) which will minimize the maximum constraint violation, using the most violated constraint j, a good initial estimate for a* is

$$a^{*} = \frac{\int_{j}^{-G} (\underline{X})}{\nabla G_{j} (\underline{X}) \cdot \underline{S}}$$
(2.27)

Since the gradients of the violated constraints are known, the scalar which is required to obtain a feasible design with respect to violated constraint in the search direction, is given to a first approximation by Eq(2.27).

The more detail procedure in MSCOP is as follow ;

- 1. Choose the most violated constraint j.
- Calculate a* for violated constraint j using Eq(2.27).

- Update the design variables for a* and check the side constraints.
- 4. If one or more violated constraints still exist, then calculate the derivative of objective, violated and active constraints and find a new search direction and then go to step 1. Otherwise proceed with the optimization in the normal fashion.

D. CCNVERGENCE CRITERIA

A desired property of an algorithm for solving a nonlinear problem is that it should generate a sequence of points converging to a global optimal point. In many cases, however, we may have to be satisfied with less faverable outcomes. In fact, as a result of non-convexity, problem size, and other difficulties, we may stop the iterative procedure if a point belongs to a described set, which is defined in MSCOP as follows ;

1. $Q_1 = \{ \underline{X} \mid | \underline{X}^\circ - \underline{X} | < \mathcal{E}_{\mathbf{X}} \mid | \underline{X}^\circ | \}$ 2. $Q_2 = \{ \underline{X} \mid | \mathbf{F} (\underline{X}^\circ) - \mathbf{F} (\underline{X}) | < \mathcal{E}_{\mathbf{F}} \cdot | \mathbf{F} (\underline{X}^\circ) | \}$

In MSCOP, the algorithm is terminated if a point \underline{X} is reached such that $\underline{X} \in Q_1 \cap Q_2$. \mathcal{E}_x is 0.001 and \mathcal{E}_f is approximatly 0.001. Since in engineering design problems it is not necessary to find solutions with more than three significant digits.

III. MSCOP USAGE

A. INTRODUCTION

Since this MSCOP is written in WATERLOO BASIC Version 2.0, it is very convenient to use. The user must first formulate the design problem with the classical machine Given the formulation design criteria. of the design problem as a nonlinear program, the user then enters the problem as a part of a BASIC program. The user defines the objective function and constraint functions using EASIC statements. Other parameters are input as data : the number of design variables NDV, the number of inequality constraints NIQC, variable bounds an initial design value and a print control number.

B. PRCBIEM FORMULATION

Generally, the experienced design engineer will be able to choose the appropriate objective for optimization depending on the requirements of the particular application. The physical phenomena of significance should first be summarized for the device to be designed. The appropriate objective can then be selected and constraints can be imposed on the remaining phenomena to assure an acceptable design from all standpoints. However, the initial formulation for the optimization problem should not be more complicated then necessary and this often requires the making of some simplifying assumptions. [Ref. 9].

After completing the formulation of the design problem, the design engineer should be able to answer the following questions :

1. What are the design variables ?

- 2. What is the objective function ?
- 3. What are the inequality constraints ?
- 4. What are the bounds on the variables ?

The engineer is then ready to input the program to the MSCOP. However, additional study and preparation of the problem may be useful. In particular, redundant constraints should be avoided if possible. MSCOP will operate with redundant constraints but it will operate faster without Selection of an initial design point from which to them. start this program is important, since it affects performance and running time. The user should use any available information which gives a good initial approximation. If side constraints exist, the user must be sure the initial values of the design variables do not violate the side constraints. This program will automatically handle an initial design point which is infeasible with respect to the G(X) < 0 constraints. However, if the initial point does not violate these constraints, convergence will likely be more rapid.

C. PROBLEM ENTRY

Problem entry is accomplished by editing the main program directly. As an example, consider the following simple NLP with two design variables, and three constraint functions.

Minimize $F(\underline{X}) = \frac{2}{1} + 3 + 3 + 2 + 2 + 2 + 1 + 1$

subject to ;

$$\begin{array}{c} x_{1} + x_{2} - 3 < 0 \\ 1 + \frac{1}{2} - 2 < 0 \\ 1 + \frac{1}{x_{2}} - 2 < 0 \end{array}$$

$$x_{1}^{2} + x_{1} - x_{2} - 2 < 0$$

 $x_{1}^{2} \ge 0.1$

With the MSCOP loaded into memory and listed on the CRT, modifications are made on the program lines as follows to input this example :

Line 100

Just after the word "data", three integers are added, separated by a comma. The first number is NDV which is the number of design variables, the second is NIQC which is the number of inequality constraints, and the third is IPRT which is print control number (0; only final results, 1; given data and final results, 2; given data and iterative subcrtimal results)

for example : 100 data 2,3,2

Lines 201-220

Each line here corresponds to a separate design variable, beginning with X(1) and continuing in order to input X(NDV). On each line, three values are separated by commas. After the word "data", these values are the initial values of the design variable, the lower bound on the variable and the upper bound on the variable. If no bound is to be specified, the entry is filled by "no".

For the sample problem, the input is :

201 data 3.,0.1,no 202 data 3.,0.1,no

Lines 400 - 450

These lines are available for defining the objective function. The objective function must be defined in terms of subscripted design variables X(1), X(2), etc.

For the sample problem, the input is :

 $400 \text{ fn}_f = x(1) **2 + x(1) * x(2) + 2 \cdot * x(2) **2 - x(1) - x(2) + 1.$ Lines 500-650

These lines are available for defining the inequality constraint functions, which must be expressed using the format :

601 if i = k then $fn_g = G_i(x) - b_i$

For the sample problem, the input is :

00601 if i = 1 then fn_g = x(1) + x(2) - 3. 00602 if i = 2 then fn_g = 1./x(1) + 1./x(2) - 2. 00603 if i = 3 then fn_g = x(1) * 2 + x(1) - x(2) - 2.

If there are many constant values in the constraint functions, the user may input data for these functions on lines 501-600 in order to simplify their statements.

IV. EXAMPLE PROBLEMS

A. DESIGN OF CANTILEVERED BEAM

1. Uniform Cantilevered Beam

Assume a cantilevered beam as shown in Figure 4.1 must be designed. The objective is to find the minimum

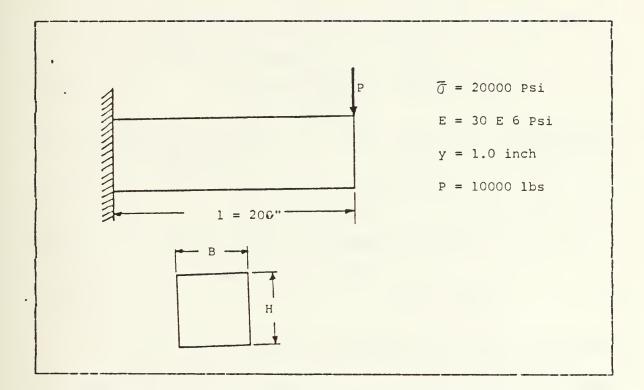


Figure 4.1 Design of a Uniform Cantilevered Beam.

volume of material which will support the load P.

The design variables are the width B and height H in the team. The design task is as follows : Find B and H to minimize volume V = B H l (4.1) we wish to design the beam subject to limit on bending stress, shear stress, deflection and geometric conditions. The bending stress in the beam must not exceed 20,000 psi.

$$\sigma_{\rm b} = \frac{M_{\rm c}}{I} = \frac{6 \, P \, I}{\frac{2}{B \, H}} < 20,000 \tag{4.2}$$

The shear stress must not exceed 10,000 psi.

$$\sigma_{\rm h} = \frac{3 P}{2 A} = \frac{3 P}{2 B H} < 10,000$$
(4.3)

and the deflection under the load must not exceed 1 inch.

$$\delta = \frac{P1^{3}}{3EI} = \frac{4P1^{3}}{EBH^{3}} \le 1.0$$
(4.4)

Additionally, geometric limits are imposed on the heam size.

$$H/b < 10.$$
 (4.7)

Now we can input this problem to MSCOP.

Input NDV, NIQC, IPRT

00100 data 2,4,2

Initial starting points

00210 data 3.5,0.5,5.0 00220 data 16.0,1.0,20.0

Evaluation of objective

 $00400 \text{ fn}_f = tl * x(1) * x(2)$

32

Evaluation of constraints 00500 tl = 200. 00501 be = 30.e+6 00502 bp = 10000. 00503 if i = 1 then fn_g = 6.*bp*t1/(20000.*b*h**2)-1. 00503 if i = 2 then fn_g = 3.*bp/(10000.*2.*b*h)-1. 00503 if i = 3 then fn_g = 4.*bp*t1**3/(be*b*h**3)-1. 00503 if i = 4 then fn_g = h/b-10.

TABLE I

The Solution of a Uniform Cantilevered Beam

objective ; 6664.0

design variable ;

X	(1)	=	1.852
Х	(2)	=	17.99

constraint ;

g(1) = 0.000902 g(2) = -0.9549 g(3) = -0.0109g(4) = -0.0286

As a result of this problem are in Table 4.1.

2. Variable Cantilevered Beam

The cantilevered beam shown in Figure 4.2 is to be designed for minimum material volume. The design variables and height h at each of are the width b 5 segments. We wish to design the beam subject to limits on stress(calculated at left end of each segment), deflection under the load, and the geometric requirement that the height of any segment does not exceed 20 times the width.

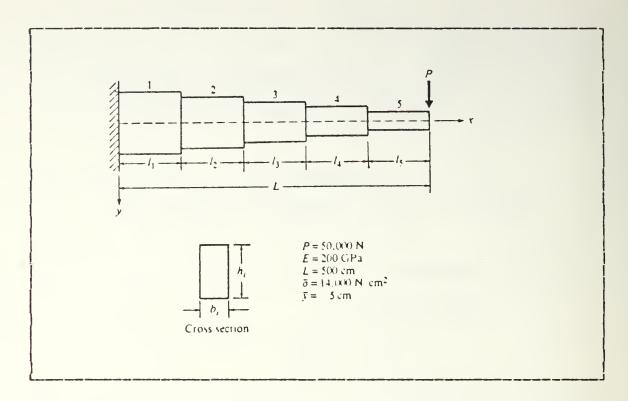


Figure 4.2 Design of a Variable Cantilevered Beam.

The deflection y at the right end of segment i is calculated by the following recursion formulas :

$$y_{0} = y_{0} = 0$$
 (4.8)

$$y' = \frac{P_{i}}{E_{i}} \left[L + \frac{1_{i}}{2} + \sum_{j=1}^{i} 1_{j} \right] + y'_{i-1}$$
(4.9)

$$y = \frac{P_{i}}{2 E_{i}} \left[L - \sum_{j=1}^{i} 1_{j} + \frac{21_{i}}{3} \right] + y_{i-1} 1_{i} + y_{i-1}$$
(4.10)

where the deflection y is defined as positive downward, y' is the derivative of y with respect to the X, and l; is the length of of segment i. Young's modulus E is the same for all segments, and the moment of inertia for segment i is

$$I_{i} = \frac{\frac{b_{i} h_{i}^{3}}{12}}{12}$$
 (4.11)

The bending moment at the left end of segment i is calculated as

$$M_{i} = P \left[L + l_{j=1} - \sum_{j=1}^{i} l_{j} \right]$$

$$(4.12)$$

and the corresponding maximum bending stress is

$$\sigma_{i} = \frac{M_{i} h_{i}}{2 I_{i}}$$
(4.13)

The design task is now defined as

Minimize:
$$V = \sum_{i=1}^{N} b h l$$
 (4.14)

Subject to :

$$\frac{\sigma_i}{\overline{\sigma}} - 1 < 0 \qquad i = 1, \dots, N \qquad (4.16)$$

(4.15)

$$\frac{y_{N}}{y} - 1 < 0$$
 (4.17)

$$h_{i} = 20 \ b_{i} < 0 \qquad i = 1, \dots, N$$
 (4.18)

 $h_{i} > 5.0$ i = 1, ..., N> 1.0 b i (4.19)

where $\overline{\sigma}$ is the allowable bending stress and \overline{y} is the allowable displacement. This is a design problem in 10 variables. There are 6 nonlinear constraints defined by Eq (4.16) and Eq(4.17), and 5 linear constraints defined by Eq(4.18), and 10 side constraints on the design variables defined by Eq (4.19).

Now we can input this problem to MSCOP.

Input NDV, NIQC, IPRI

00100 data 10,11,2

00210 data

Initial starting points

5.,1.,no

data 5.,1.,no data 5.,1.,no data 5.,1.,no data 5.,1.,no data 5.,1.,no data 40.,5.,no data 40.,5.,no data 40.,5.,no data 40.,5.,no data 40.,5.,no $\begin{array}{c} 00210\\ 00220\\ 00230\\ 00240\\ 00250\\ 00260\\ 002270\\ 00280\\ 00290\\ 00290\\ 00230\\ 00230\\ 00230\\ 00230\\ 00230\\ 00230\\ 00230\\ 00230\\ 00000\\ 00000\\ 0000\\$ Evaluation of objective * (00400 fn f = 100.x(1) * x(6) + x(2) * x(7) + x(3) * x(8) $x(4) * x(9)^{-} + x(5) * x(10)$) Evaluation of constraints. 00490 def fn g(x,i) 00498 dim bm(10),bi(10),sigi(10),ypb(10),yb(10) 00500 pcb = 50000. 00501 be = 200.e+5 00502 tl = 200. sigb = 14000.for m = 1 to km = m+5 - 5 bi(m) = x(m) * x(km) **3/12. sigi(m) = bm(m) * x(km)/(2.*bi(m)) next m $y_{zo} = 0.$ $y_{pzo} = 0.$ for m = 1 to 5

00517 00518 00519	<pre>ypb(m) = (pcb*sl*(tl+sl/2m*sl))/(be*bi dumm = pcb*sl**2*(tl-m*sl+2.*sl/3.) yb(m) = dumm/(2.*be*bi(m))+ypzo*sl+yzo</pre>	(m))-ypzo
00520 00521 00522	$ypzo = ypb(\pi)$ $yzo = yb(\pi)$ next m	
00550 00560 00570	rom constraint function	
00580	if i = 1 then fn_g = sigi(1)/sigb-1. if i = 2 then fn_g = sigi(2)/sigb-1. if i = 3 then fn_g = sigi(3)/sigb-1. if i = 4 then fn_g = sigi(4)/sigb-1. if i = 5 then fn_g = sigi(5)/sigb-1. if i = 6 then fn_g = yb(5)/ybb-1. if i = 7 then fn_g = x(6)-20.*x(1) if i = 8 then fn_g = x(7)-20.*x(2) if i = 9 then fn_g = x(8)-20.*x(3) if i = 10 then fn_g = x(9)-20.*x(4)	
00610 00620 00630	if $i = 5$ then $fn^{-}g = sigi(5)/sigb-1$. if $i = 6$ then $fn^{-}g = yb(5)/ybb-1$. if $i = 7$ then $fn^{-}g = x(6)-20.*x(1)$ if $i = 8$ then $fn^{-}g = x(7)-20.*x(2)$ if $i = 9$ then $fn^{-}g = x(8)-20.*x(3)$	
00632	if i = 1 then fn_g = sigi (1) /sigb-1. if i = 2 then fn_g = sigi (2) /sigb-1. if i = 3 then fn_g = sigi (3) /sigb-1. if i = 4 then fn_g = sigi (4) /sigb-1. if i = 5 then fn_g = sigi (5) /sigb-1. if i = 6 then fn_g = x (6) -20.*x (1) if i = 8 then fn_g = x (7) -20.*x (2) if i = 9 then fn_g = x (9) -20.*x (4) if i = 10 then fn_g = x (10) -20.*x (5)	

TABLE II

The Solution of a Variable Cantilevered Beam objective : 62133.35 design variables constraints X(1) = 2.994G(1) = -0.00219X(2) = 2.782G(2) = -0.00415X(3) = 2.528G(3) = -0.00508X(4) = 2.208G(4) = -0.00406X(5) = 1.761G(5) = -0.0177X(6) = 59.88G(6) = -0.4401X(7) = 55.62G(7) = -0.0101X(8) = 50.56G(8) = -0.0231X(9) = 44.14G(9) = 0.0000

X(10) = 35.19

- G(10) = -0.0248
 - G(11) = -0.0278

B. SIMPLE TRUSS

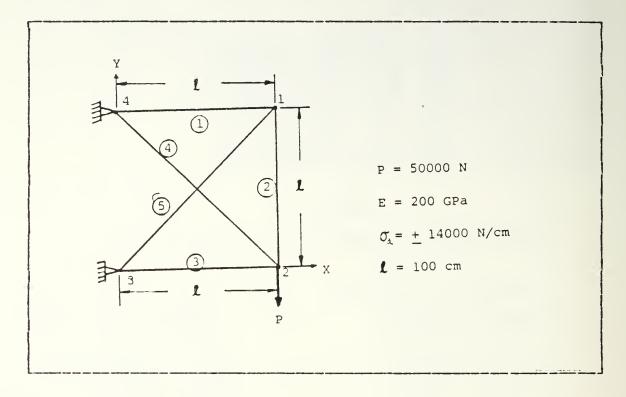


Figure 4.3 Design of a 5-Bar Truss.

A simple truss with 5 members as shown in Figure 4.3 is designed for the minimum volume. The design variables are the sectional areas of the members. The constraints are formed for the stresses of the members not to exceed the given allowable stress. The lower bound for each design variable is also considered. The stresses are obtained by the displacement method of the finite element analysis. The equation to be solved is given by

$$\underline{K} \cdot \underline{u} = \underline{P} \tag{4.20}$$

where \underline{K} is the stiffness matrix, \underline{u} is the displacement vector and \underline{P} is the load vector as follows :

$$\underline{\underline{U}} = \begin{bmatrix} u \\ 1 \\ v \\ 1 \\ u \\ 2 \\ v \\ 2 \\ v \\ 2 \end{bmatrix} \qquad \underline{\underline{p}} = \begin{bmatrix} 0 \\ 0 \\ 0 \\ 0 \\ -5000 \end{bmatrix} \qquad (4.21)$$

$$K = E \begin{bmatrix} \frac{A_{1}}{2} + \frac{A_{5}}{\sqrt{2}2} & \frac{A_{5}}{\sqrt{2}2} & 0 & 0 \\ \frac{A_{5}}{\sqrt{2}2} & \frac{A_{2}}{2} + \frac{A_{5}}{\sqrt{2}2} & 0 & -\frac{A_{2}}{2} \\ 0 & 0 & \frac{A_{3}}{2} + \frac{A_{4}}{\sqrt{2}2} & -\frac{A_{4}}{2} \\ 0 & 0 & \frac{A_{3}}{2} + \frac{A_{4}}{\sqrt{2}2} & -\frac{A_{4}}{2} \\ 0 & -\frac{A_{2}}{2} & -\frac{A_{4}}{2} & \frac{A_{2}}{2} + \frac{A_{4}}{\sqrt{2}2} \end{bmatrix}$$
(4.22)

From Eq. (4.20) the displacements are solved by

$$\underline{U} = \underline{K} \cdot \underline{P}$$
(4.23)

Having displacements at all nodes, we can calculate the stress for each element.

$$\sigma_{i} = E \cdot \varepsilon = \frac{1}{1}$$
(4.24)

where

$$\Delta l_{1} = \sqrt{(l_{1} + u_{1})^{2} + v_{1}^{2}} - l_{1}$$

$$\Delta l_{2} = \sqrt{(l_{2} + v_{1} - v_{2})^{2} + (u_{1} - u_{2})^{2} - l_{2}}$$

$$\Delta l_{3} = \sqrt{(l_{3} + u_{2})^{2} + v_{2}^{2}} - l_{3}$$
(4.25)

$$\Delta l_{4} = \sqrt{(l_{3} + u_{2})^{2} + (l_{2} - v_{2})^{2}} - l_{4}$$

$$\Delta l_{5} = \sqrt{(l_{3} + u_{1})^{2} + (l_{2} + v_{1})^{2}} - l_{5}$$

The design problem is given by

minimize
$$V = \sum_{i=1}^{5} A_i l$$
 (4.26)

Subject to

$$G_{i} = \frac{|O_{i}|}{\sigma_{a}} - 1.0 \le 0 \quad i = 1,...,5 \quad (4.27)$$

$$A > 0.1$$

i = 1,...,5 (4.28)

The MSCOF input for this problem is given as follows : Input NDV, NIQC, IPRT 00100 data 5,5,2 Initial starting point $\begin{array}{c} 00200 \\ 00202 \\ data \\ 3...1,no \\ 00204 \\ data \\ 3...1,no \\ 00208 \\ data \\ 3...1,no \\ 00008 \\ data \\ 0...1,no \\ 0...$

```
0504  cs = 2.*sqr(2.)
0505  ct = te/t1
0506  k11 = (x(1)+x(5)/cs)*ct
0508  k21 = k12
0509  k22 = (x(2)+x(5)/cs)*ct
0510  k24 = -x(2)*ct
0511  k33 = (x(3)+x(4)/cs)*ct
0511  k33 = (x(3)+x(4)/cs)*ct
0513  k42 = k24
0514  k43 = k34
0515  k44 = (x(2)+x(4)/cs)*ct
0516  dk1 = -k11/k12
0517  dk2 = -(k12+k22*dk1)/k24
0518  dk3 = -k33*dk2/k34
0519  pp = -50000.
0520  vv(1) = pp/(k42*dk1+k43*dk3+k44*dk2)
0522  vv(3) = dk3*vv(1)
0522  vv(3) = dk3*vv(1)
0522  vv(4) = dk2*vv(1)
0592  dl1 = sgr((t1+vv(1))**2+vv(2)**2)-t1
0592  dl2 = sgr((t1+vv(3))**2+vv(4)**2)-t1
0595  h1 = sgr(2.)*t1
0596  dl4 = sgr((h1+vv(3))**2+(h1-vv(4))**2)-h1
0598  dl5 = sgr((h1+vv(1))**2+(h1-vv(4))**2)-h1
0600  rem constraint.
0602  if i = 1 then fn_g = te*d11/(t1*sigb)-1.
0604  if i = 3 then fn_g = te*d12/(t1*sigb)-1.
0608  if i = 3 then fn_g = te*d13/(t1*sigb)-1.
0608  if i = 3 then fn_g = te*d14/(h1*sigb)-1.
0608  if i = 3 then fn_g = te*d13/(t1*sigb)-1.
0600  if i = 5 then fn_g = te*d13/(t1*sigb)-1.
0610  if i = 5 then fn_g = te*d13/(t1*sigb)-1.
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0610  if i = 5 then fn_g = te*d13/(t1*sigb)-1.
0610  if i = 5 then fn_g = te*d13/(t1*sigb)-1.
```

TABLE III

The Sclution of a 5-Bar Truss

objective ; 108.52

design variables

constraint

=	0.1	G (1)	=	-1.9988
=	0.1	G(2)	=	-2.0030
=	3.514	G(3)	=	-0.0030
=	4.948	G (4)	=	-0.1203
Ξ	0.1	G(5)	=	-1.8797
		= 0.1 = 0.1 = 3.514 = 4.948 = 0.1	= 0.1 G(2) = 3.514 G(3) = 4.948 G(4)	= 0.1 G(2) = = 3.514 G(3) = = 4.948 G(4) = = 151

V. SUMMARY AND CONCLUSION

Numerical optimization is a powerful technique for those confronted with practical engineering design problems. It is also a useful tool for obtaining reasonable solutions to the classical engineering design problems. Since many engineers are now using microcomputers for solving design problems, the development of microcomputer software which can be easily used is needed.

In this thesis, an algorithm for constrained optimization problems is programmed in standard BASIC language (WBASIC version 2.0) on an IBM 3033. The users can easily convert this to other microcomputers.

MSCOF (Microcomputer Software for Constrained Optimization Problems) employs the method of feasible directions and specific modifications of a one-dimensional search for constrained optimization. MSCOP has been validated by tests on three constrained optimization problems. Its performance is good and could be made better through refinement of the algorithm.

Since microcomputers are available with reasonable memory size and computational speed, their capabilities will continue to improve as more engineering software becomes available. MSCOP is considered to be a first step toward more widespread use cf optimization techniques on microcomputers.

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<u>APPENDIX A</u>

MSCOP PROGRAM LISTING

cption base 1 dim x(21),x0(21),gcv(51),ngcv(51),df(21),dg(51,21) dim thta(21),wrky(51,51) dim a(51,21),b(51,51),p(21),y(21),s(21),u(51),c(51) dim iwrk(51),jwrk(51),wrk1(51),wrk2(51),wrk3(51) dim wrku(51),wrk1(51),lowb(21),uprb(21),lo3(6),upf rem input data 0010 $\begin{array}{c} 0010\\ 0020\\ 0021\\ 0030\\ 0040\\ 0050\\ 0060\\ 0070\\ 0080\\$ 1),u(51),c(51) 1),wrk3(51) 1),lo3(6),up‡(6) 10000 gosub gosub 10000
rem input number of design variables and constraints.
read ndv,niqc,iprt
data 2,4,2
for i = 1 to ndv
rem input initial value of design variables
 read x(i)
 x0(i) = x(i)
 if niqc = 0 then 160
 read Io\$,up\$
 if lo\$ = 'no' then lowb(i) = bnlo else lowb(i) =
 value(lo\$) 0115 0120 0125 0130 0135 0140 value(lo\$) if up\$ = 'no' then uprb(i) = bnup else uprb(i) 0150 = value (up\$) 0160 0200 0360 0370 0375 0380 0390 0400 0420 0420 0430 next i data 3.5,0.5,10. data 16.,1.0,20. rem evalute the objective-function obj = fn_f (x) itri = 1 $f_{rem}^{1trl} = 1$ rem objective function def fn f(x) fn_f = 200.*x(1)*x(2) fnend tor i = 1 to nigc
 gcv(i) = fn_g(x,i)
next i rem evaluate the constraints 0440 0450 0450 0480 0490 0500 0510 0520 0520 0530 rem constraint functions fn g (x,i) = 200. = 30.e+6 = 10000. def tl = be bp iI = = (6.*bp*t1)/ (20000.*x(1)*x(2)**2)-1. (3.*bp)/(2)000.*x(1)*x(2))-1. (4.*bp*t1**3)/ (be*x(1)*x(2)**3)-1. x(2)/(10.*x(1))-1. i fn_g = 1 then = 0540 if if i 2 3 fn_g fn_g then = = then = = 0560 0650 0700 0710 0720 if i = 4 then $fn_g =$ rem initial counting number input ical = 1 if ical > 3 then stor ical = 1
if ical > 3 then stop
rem call the optimization code.
gosub 2000
rem print results.
rem 0720 0730 0740 0750 0760 0770 0780 rem rem re-counting number input. ical = ical+1 if ical = 3 then 850 rem 10% reduce the design variables. ŎŹŠ 079 0 0800 for i = 1 to ndv

x(i) = 0.9 * x(i)x0(i) = x(i) 0810 0820 0830 next i goto 720 rem 10% increase design variables. for i = 1 to ndv x(i) = 1.1*x(i) x0(i) = x(i) pert i 0840 0850 0850 0860 0870 0880 0890 0900 next i goto 720 2000 2001 2002 2003 rem calculate the obj. constraint for. obj = fn_f(x)
for i = T to nigc
 gcv(i) = fn_g(x,i)
next i 2004 2008 2010 2020 itrq = 1 itrg = itrg+1 rem calculate the number of active and violate constraints. $\begin{array}{c} 2030\\ 2040 \end{array}$ gosub 3500 active or violated constraints. gosub 3800 if navc = 0 then 2190 gosub 3900 rem calculate the gradient of objective and 2050 2060 2070 2080 2090 2100 2110 rem calculate the push-off factors gosub 4000 rem making the matrix c rem normalized the df(i) gosub 4100 rem normalized the DG(i) gosub 4200 if nvc > 0 then gosub 4400 else gosub 4600 rem calaulate the usable-feasible direction s(i) gosub 5000 goto 2230 rem normalize the df(i)
for i = 1 to ndv
 s(i) = -(df(i))
next i rem normalize the s(i) gosub 5700 rem one-dimensional search if nvc = 0 then gosub 6000 else gosub 9000 rem update x for alph gosub 7000 gosub 7100 ğosub rem calculate new point value.
nobj = fn_f(x)
rem convergence test 2320 2330 2340 2350 2360 2370 gosub 6780 walp <= accx and delf <= dabf then 2470 itri = itri+1 if itri > mxit then print 'check the problem' obj = nobj r i = 1 to ndv x0(i) = x(i) xt i íf 2380 2390 for next i
for i = 1 to nicc gcv(i) = fn_g(x,i) if iprt = 2 then 2460 goto 2010 řem print final results print '**** final resu final results ***** gosub 9200 2500 2500 retúrn 3000 rem initialize the integer working array

```
for i = 1 to nigm
    iwrk(i) = 0
3005
3010
3010
3015
3020
3050
3055
3060
        next i
         return
         rem initialize the integer working array
for i = 1 to nigm
              jwrk(i)
t i
                           = 0
nexť
         return
         rem initialize the one-dimension working array
for i = 1 to nigm
    wrk1(i) = 0.
         next i
         return
         rem initialize the one-dimension working array for i = 1 to nigm
        wrk2(i)
next i
                           = 0.
         return
         rem initialize the one-dimension working array for i = 1 to nigc
              wrk3(i) = gc\vec{v}(i)
         next i
         return
        rem initialize the two-dimension working array
for i = 1 to nigm
for j = 1 tc ndvm
                  wrky(i,j) = 0.
kt j
              next
        next i
        return
        rem initialize the derivative of objective DF(i)
for i = 1 to ndvm
    df(i) = 0.
next i
         return
        rem initialize the a(i,j),p(i),y(i),c(i)
for i = 1 to ndvm
            i =

p(i) = 0.

for j = 1

a(j,i)

+
                              tc niqm
= 0.
        next i
for j = 1
    c(j) =

                         to
                             niqm
                         0.
                      =
        next
         return
        rem initialize the derivative of constraints DG(i,j)
for i = 1 to nigm
   for j = 1 to ndvm
             1 = 1
for j = 1
    dg(i,j)
next j
                              tondvm
                                =
                                    Ο.
         next i
         return
        rem initialize the b(i,j)
for i = 1 to nigm
    for j = 1 to nigm
        b(i,j) = 0.
        next j

        next i
         return
rem Calculate the number of active and violate
3500
        constraints.
gosub 3000
gosub 3100
3502
3504
3510
3520
3530
                   0
        ňac
              =
              =
i
        nvc
                   0
         for
                  = 1 to nigc
```

```
if gcv(i) \ge vcc then 3580
if gcv(i) < acc then 3590
nac = nac+1
goto 3590
        ñvç
                 =
                     nvc+1
next
            i
navc
if n
            =
                nac+nvc
      navc = 0 then 3790
        ii =
jj =
for
                    1
                     1
               if gcv(i) >= vcc
if gcv(i) >= vcc
if gcv(i) < acc
wrk(nvc+ii)
wrk1(nvc+ii)
ii = ii+1
goto 3750
wrkk(jj) = i
wrk1(jj) = gcv(i)
jj = jj+1
               if
if
                                                           then 3720
hen 3750
                                                        then
= i
                                                         =
                                                              qcv(i)
                                            gcv(i)
        jj=
next i
return
rem calculate the gradient of f(x)
gosub 3300
for i = 1 to ndv
qosub
for i
dxi =
if
       sub 3300
si = 1 to ndv
. = fdm*abs(x(i))
if dxi <= mfds then dxi = mfds
    x(i) = x(i) +dxi
    dobj = fn f(x)
    df(i) = (dobj-obj)/dxi
    x(i) = x0(i)
t i
</pre>
next
return
        calculate the DG(i,j)
rem
gosub 3400
for i = 1 to ndv
    dxi = fdm*x(i)
    if dxi < mfds then dxi = mfds</pre>
       if dx1 < mids then dx1 = mids
x(i) = x(i) + dxi
for j = 1 tc navc
    k = iwrk(j)
    dcon = fn g(x,k)
    dg(j,i) = (dcon-wrk1(j))/dxi
next j
x(i) = x0(i)</pre>
       next
x(i)
next
             i
return
rem calcilate the push-off factor
for i = 1 to navc
thta(i) = tht0*(1.-wrk1(i)/acc
                           = tht0*(1.-wrk1(i)/acc)**2
(i) > thtm then thta(i) = thtm
        if thta (i)
t i
next
return
rem normalize the DF(i)
gosub 3200
fsg = 0.
for i = 1 to ndv
                     fsq+df(i) **2
        fsg =
next i
fsq = sqr(fsq)
if fsg = 0. then fsq = zro
for i = 1 to ndv
    wrk3(i) = (1./fsg)*df(i)
next i

return
rem normalize the DG(i)
gosub 3250
for i = 1 to navc
        gsq = 0.
```

```
for j
gsq
next j
                   = 1 \text{ tc ndv}
                     = gsg+dg(i,j) **2
       gsq = sqr(gsq)
if gsq = 0. then gsq = zro
for j = 1 to ndv
wrky(i,j) = (1./gsq)*dg
next j
                                        (1./gsq)*dg(i,j)
 next
next i
 return
 rem exist the
gosub 3350
for i = 1 to
for j = 1
                     the violate constraints
                            navc
             = j = 1
a(i,j)
                            tc ndv
                            = wrky(i,j)
       next j
a(i,ndv+1)
                             = thta(i)
 next
for i
                 1 to
            =
                            ndv
 p(i)
next i
                  = -wrk3(i)
 p (ndv+1)
for i =
yy =
for j
                   = phid
1 to navc
                  Ö
                        1 tc ndv+1
a(i,j)*p(j)
yy+xx
              j
xx
                    =
xx

next

c(i) =

ndt =

ret
                    Ξ
                  j
                        (-1.)*yy
ndt-
return
rem only exis
gosub 3350
for i = 1 to
for j = 1
a(i,j)
-+ j
-+1
 ndb = navc
                    exist active constraints
                            navc
                            tc ndv
                            = wrky(i,j)
       next j
a(i,ndv+1) = thta(i)
 next i
for j =
       j = 1 to ndv
a(navc+1,j) = wrk3(j)
 next j
a(navc+1,ndv+1) =
p(ndv+1) = 1.
for i = 1 to navc;
                                      1.
                  1 to navc+1
       cc = a(i, ndv+1) * p(ndv+1)

c(i) = (-1.) * cc
 next`i
ndb = navc+1
 return
 return
rem calculate the usable-feasible direction
gosub 3000
gosub 3250
gosub 3450
for i = 1 to ndt
    for j = 1 to ndv+1
        wrky(j,i) = a(i,j)
    next j
    next j
      ti

i = 1

for j =

ff = 0.

for k =

tf =

tf =

tf =
 next i
for i =
                   1 \text{ to ndb} = 1 \text{ to ndb}
                               1 to ndv+1
a(i,k)*wrky(k,j)
ff+tf
              next k
b(i,j)
t j
                             =
                                  (-1.) *ff
        next
```

next iter i 0 5*ndb = nmax = rem begin iteration iter = iter+1 cbmx = 0. cbmx = 0. ichk = 0 for i = 1 to ndb ci = c(i) bii = b(i i) if bii = 0. then if ci > 0. then then 5340 cb = ci/bii if cb <= cbmx then 5340 ichk = i Ξ cbmx cb next i if cbm if ich cbmx ichk jj = jj = if b zro or iter > nmax then 5550
0 then 5550 < = 0 then iwrk(ichk) = ichk else iwr b(ichk,ichk) = 0. then b(ichk,ichk) if bb = 1./b(ichk,ichk) if bb = 0. then bb = zro for i = 1 to ndb b(ichk,i) = bb*b(ichk,i) next i C(ichk) if = ichk else iwrk (ichk) = 0 = zro b (lCn... next i c(ichk) = cbmx for i = 1 to ndb if i = ichk then 5530 bbi = b(i,ichk) b(i,ichk) = 0. for j = 1 to ndb if j = ichk then 5520 b(i,j) = b(i,j)-bbi*b(ichk,j) pext j vext j vext j 5220 0 goto ner for = i for i = 1 to ndb
 u(i) = 0.
 j = iwrk(i)
if j > 0 then u(i) = c(j)
for i í ff for = 1 to ndb 0. ff for 1 to ndb ff+wrky(i,j)*u(j) = = $\begin{array}{c} next \\ y \\ s \\ i \\ t \\ i \end{array}$ j p{i y(i) -ff = next return rem normalized the s(i) ssq = 0. for i = 1 to ndyssg = ssg+s(i)**2 next i ssq = sqr(ssq)
if fslp = 0. then fslp = zro
for i = 1 to ndv s(i) t i = (1./ssq)*s(i) next return rem one-dimensional search for initial feasible point. rem calculate for slope of f(x)fslp = 0. for i = 1 to ndv

6020 6025 6035 6040 fslp = fslp+df(i) *s(i) next i rem idenfy the initial point. next alow = 0. flow = obj for i = 1 to nicc wrkl(i) = gcv(i) next i rem update x for a1st. alph = a1st gosub 7000 gosub 7100 fem calculate the f1st and wrk1(i) f1st = fn f(x) for i = 1 to nigc 1205050 1205050 13305050 145050 $wrk1(i) = fn_g(x,i)$ next next i
rem check the feasibility.
ncv1 = 0
for i = 1 to nigc
 if wrk1(i) < vcc then 6170
 ncv1 = ncv1+1</pre> next i if ncv1 = 0 the a1st = 0.5*a1st goto 6105 = 0 then 6200 rem find a2nd ; the 2nd mid-point. rem 2-points quadratic fit interpolation for minimum f(alpa). a0 = flow a1 = fslp a2 = (flst-a1*a1st-a0)/(a1st**2) if a2 <= 0. then a2 = zro a2nd = -a1/(2.*a2) rem 2-points linear interpolation for g(alpa)=0. for i = 1 to nigc a0 = wrkl(i) if a1st = 0. then a1st = zro a1 = (wrk1(i)-a0)/a1st if a1 <= 0. then a1 = zro walp = -a0/a1 if walp <= 0. then walp = 1000. if walp >= a2nd then 6265 a2nd = walp flow a0 = next i rem update x for a2nd. alph = a2nd gosub 7000 gosub 7100 fem calculate f2nd and wrk2(i) f2nd = fn f(x) for i = 1 to nigc $wrk2(i) = fn_g(x,i)$ next i rem find final roint aupr by using 3-points guadratic fit. f1 = flow 6320 6321 6325 6326 alp1 = alowf2 = f1ctalp2 = a1st

6330 6331 6335 6340 f3 = f2nd alp3 = a2nd rem 3-points quadratic fit interpolation. gosub 6600 if a2 = 0 then a2 = zro alp: gosub 6600 if a2 = 0. then a2 a3rd = -a1/(2.*a2) if a3rd <= 0. then for i = 1 to nigc f1 = wrk1(i) f2 = wrk1(i) f3 = wrk2(i) gosub 6600 gosub 6630 if alps > a3rd the a3rd = alps = zro then a3rd = 1000. then 6380 i next rem update x for aupr alph = a3rd gosub 7000 gosub 7100 fem calculate the fupr and wrku(i)
fupr = fn_f(x)
for i = 1 to nigc next i rem find 4th new point. f1 = f1st f2 = f2nd f3 = f3rd alp1 alp2 alp3 rem = a1st= a2nd = a3rd 3-points quadratic fit. gosub if 2 = 0. then a2 = -a1/(2.*a2)1 to nigc wrk1(i) wrk2(i) wrk3(i) = a1ct a2 then a2 = zro aupr i = f1 = for f2 f3 = = alp1 a1st = a2nd alp3 = alp3 = a3rdgosub 6600 gosub 6630 alps > aurr then 6540 pr = alps íf aupr next i rem update x for aupr = aupr 7000 7100 alph. gosub qosub rem evaluate furr and wrku(i) furr = fn_f(x) for i = 1 to nigc wrku(i) = $f\bar{n}_g(x,i)$ next find optimum alpa for not violating constraints. b 14300 rer gosub return rem 3-points quadratic if alp1 = alp2 cr alp2 fit = a alp3 or alp1 = alp3 return
((f3-f1)/(alp3-alp1)(f2-f1)/(alp2-alp1))/(alp3-alp2)
(f2-f1)/(alp2-alp1)-a2*(alp1+alp2)
f1-a1*alp1-a2*alp1**2 then a2 = 6605 6610 6615 6620 6630 a 1 = a0 = return rem zero of polynomial for q(alpa)

```
dd
if
if
        = a1 * * 2 - 4 * a2 * a0
                   = 0.
                             then 6715
        dd
                <
<=
        a2 <
a2 =
a1pb
a1pc
if a
if a
                           then a2 = zro
then a2 = zro
(-a1+sqr(dd))/(2.*a2)
(-a1-sqr(dd))/(2.*a2)
<= 0 and alpc <= 0. t
>= 0. and alpc >= 0.
>= 0. and alpc < 0. t
= alpc</pre>
                   0.
                    =
                    =
               alph
                                                                                then 6715
             alpb
alpb
alps
6720
                                                                                                6695
                                                                                  then
                                                                                then 6685
                           = alpc
 goto
        alps = alpb
goto 6720
if alpb >= alpc then 6710
alps = alpb
                alps = 0.6720
                             = alpb
goto 6720
alps =
goto 6720
alps = 1000.
return
                                  alpc
 rem update aboj and alpx
delf = abs(obj-nobj)
diff = abs(delf/obj)
abcj = (aboj+diff)/2.
walp = 0.
abcj =
walp =
welx =
for i =
                 Õ
                    ).
1
                          to ndv
        delx = abs(x0(i)-x(i)
difx = abs(delx/x0(i))
if delx >= welx then
if difx <= walp then
walp = difx
                                                             welx
                                                                          =
                                                                                delx
                                                              6880
             i
 next
 alpx
dabf
                  (alpx+walp) /2.
accf*abs(obj)
             =
             =
 return
 rem update the x(i)
for i = 1 to ndv
        x(i)
t i
                    = x0(i)+alph*s(i)
 next
 return
 rem check the side-constraints.
for i = 1 to ndv
    if x(i) < lcwb(i) then x(i)
    if x(i) > uprb(i) then x(i)
                                                                               = lowb(i)
= uprb(i)
           'i
 next
 return
 rem estimate the alpa
fstr = flow
alpa = alow
nvc1 = 0
for i = 1 to niqc
    if wrk1(i) < vcc then 8070
    nvc1 = nvc1+1
next i
    if nvc1 > 0 th cn 8120
next
if n
if f
        nvc1
f1st
                    >
                          0 th
fstr
                               then 8120
tr then 8120
        alpa
                          a1st
f1st
                    =
        fstr
                   =
nvc1 = 0
for i = 1 to nigc
    if wrk2(i) < vcc then 8160
    nvc1 = nvc1+1
    if i = nvc1+1</pre>
next
if n
if f
        nvc1
f2nd
alpa
                          0 th
fstr
                    >
                                       then 8210
                          a2nd
f2nd
                    =
         fstr
                     =
                 0
nvc1
             =
```

8220 8230 8240 8250 8260 8260 for i = 1 to nigc
 if wrk3(i) < vcc then 3250
 nvc1 = nvc1+1</pre> next i if nvc if f3r nvc1 > 0 the f3rd > fstr 0 then 8300 then 8300 8280 8290 8300 8310 alpa = a3rd fstr = f3rd fstr - ____ nvc1 = 0 for i = 1 to nigc if wrku(i) < vcc then 8340 nvc1 = nvc1+1 = 1
nvc1 =
if nvc1
if fr 8320 8320 8330 8340 8350 8350 8370 nvc1 > 0 th fupr > fstr alpa = aupr fstr = fupr > 0 then 8390 > fstr then 8390 8380 8390 8400 alph = alpa return 9000 rem one-dimensional search for initial infeasible point. 9002 ii = 1qcvm = wrk1(1)for i = 1 to navc if wrk1(i) <= gcvm then 9014 9004 9006 9010 ii =`i 9**01**2 gcym = wrk1(i)9014 next i 9016 rem calculate the slope of badly violation. gslp = 0. for i = 1 to ndv gslp = gslp+dg(ii,i)*s(i) next i rem calculate the alph. if gslp = 0. then gslp = zro alph = -gcvm/gslp rem update X fcr alph. gosub 7000 gosub 7100 rem evalute the objective and constraint. obj = fn_f(x) for i = T to nigc 9040 9042 $gcy(i) = fn_{\bar{g}}(x,i)$ 9044 next i 9046 rem calculate the NVC. gosub 3500 9048 9048 9050 9052 9054 9056 9058 if nvc = 0 then return rem update initial value. for i = 1 to ndv x0(i) = x(i) next i 9060 9062 9064 rem calculate df(i),dg(i,j) and push-off factor. gosub 3800 gosub 3900 gosub gosub 4000 9066 9068 9070 9072 9074 rem normalize the df(i), dg(i, j) gosub 4200 rem find the search direction. 9076 9078 gosub 5000 9078 9030 9200 9205 9210 9215 9225 9225 9230 goto 9000 rem print the results print '' print '**************** da 1.1 print The number of design variables = 'ndv The number of inequality constraints = ', niqc print print 11 print

```
print
next i
print
                          'the number of active constraints = ';nac
             print
                           1.1
            print
             print
                          'the number of vionate constraints = ';nvc
                          11
            print
print
           print '' constr
for i = 1 to nigc
    print 'g(';i;')
next i
return
                          ***** constraint value *****
                                                        = ';gcv(i)
            return
            rem default number
           rem default number
mxit = 50 ! maximum iteration number
fdm = .01 ! finite difference step
mfds = .001 ! maximum absolute finite difference step
vcc = .004 ! violated constraint criteria (thickness)
acc = -.1 ! active constraints criteria (thickness)
tht0 = 1. ! push-off factor multiplier (theta zero)
thtm = 50. ! maximum value of push-off factor
phid = 100000. ! weighting-factor used in direction
when infeasible
accf = .001 ! absolute convergence criteria
                                                                                                                   (thickness)
(thickness)
            accf = .001
accx = 0.001
zro = .0001
espl = .005 !
bnlo = -1.e+70
bnup = 1.e+70
dalp = .01 !
9590
                                               ! absolute convergence criteria
9600
9600
9620
9620
9630
9650
                                               ! absolute convergence criteria.
                                               ! defined zero
                                              used to prevent division by zero
! the value of low boundary
! the value of upper boundary
step size of alpa in one-dimensional
                                               search
                           0.1
21
51
9660
9670
9680
                                          ! step size for reduce objective
! reduce the design variable factor
the number of maximum design variable
! the number of maximum inequality
            abcj
                        =
            alpx
ndvm
                       =
                       =
                                       1
9690
                        =
            niqm
                                              ccnstraints
9700 return
9800 end
```

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