

# Stack Frames (12A)

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# Based on

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ARM System-on-Chip Architecture, 2<sup>nd</sup> ed, Steve Furber

Introduction to ARM Cortex-M Microcontrollers  
– Embedded Systems, Jonathan W. Valvano

Digital Design and Computer Architecture,  
D. M. Harris and S. L. Harris

ARM assembler in Raspberry Pi  
Roger Ferrer Ibáñez

<https://thinkingeek.com/arm-assembler-raspberry-pi/>

# Nested and recursive function calls

## Nested function call

```
int main(void) {  
    f1( ... );  
}  
  
void f1 ( ... ) {  
    f2 ( ... );  
}  
  
void f2 ( ... ) {  
}
```

## Recursive function call

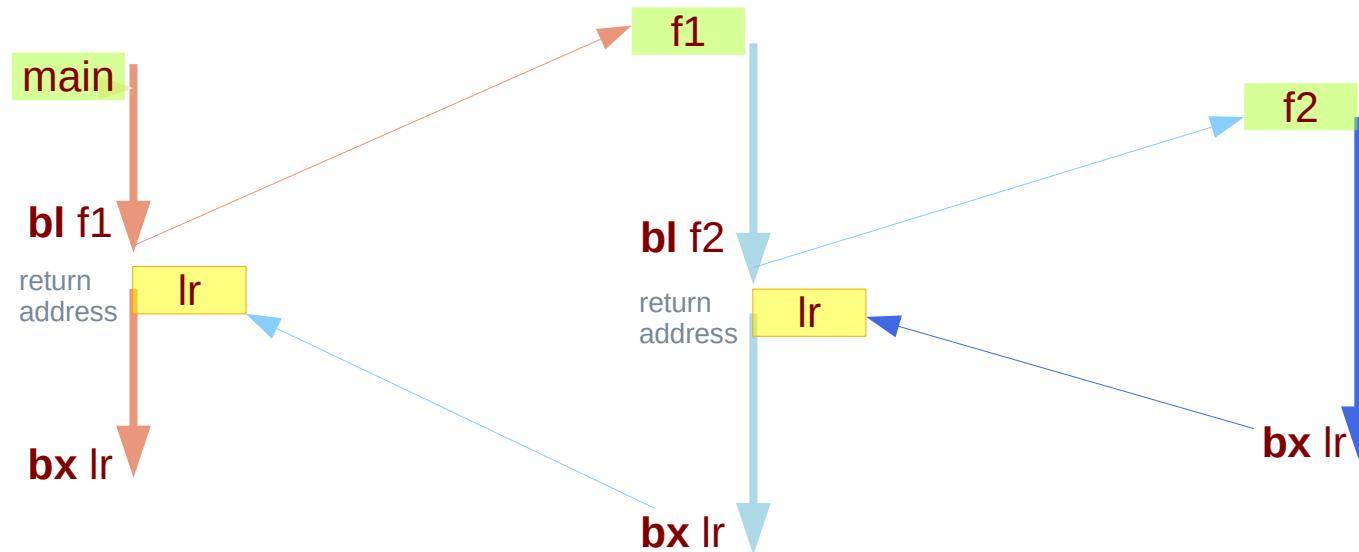
```
int fact (int n)  
{  
    if (n < 1)  
        return (1);  
    else  
        return (n * fact(n-1));  
}  
  
int fact (int n)  
{  
    if (n < 1)  
        return (1);  
    else  
        return (n * fact(n-1));  
}  
  
int fact (int n)  
{  
    if (n < 1)  
        return (1);  
    else  
        return (n * fact(n-1));  
}
```

# Nested and recursive function calls

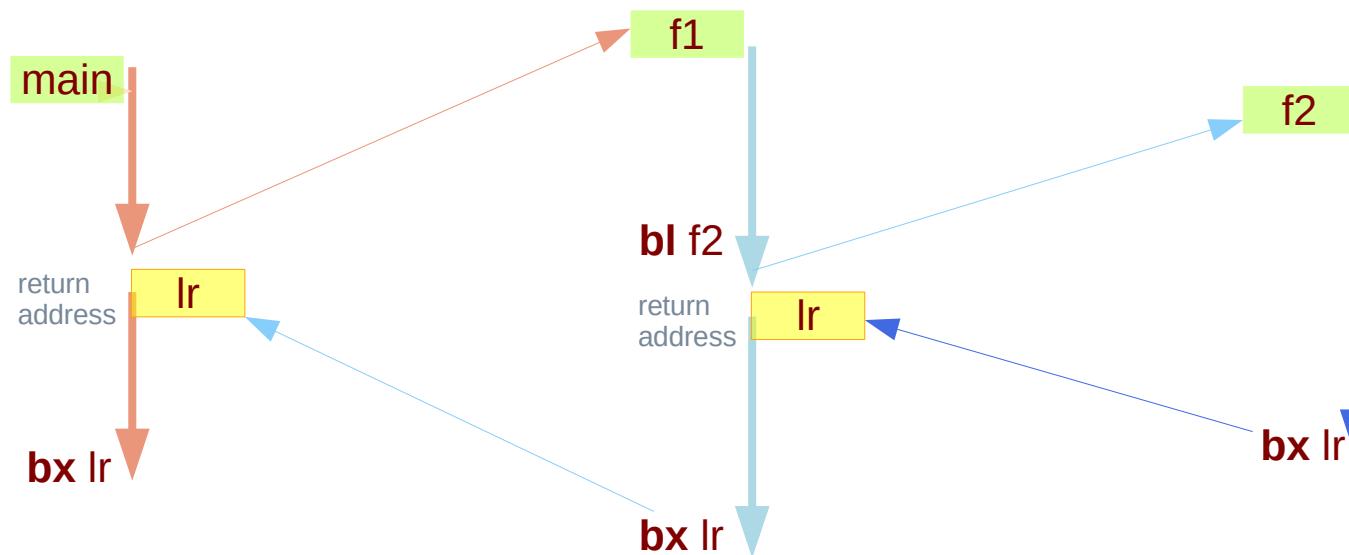
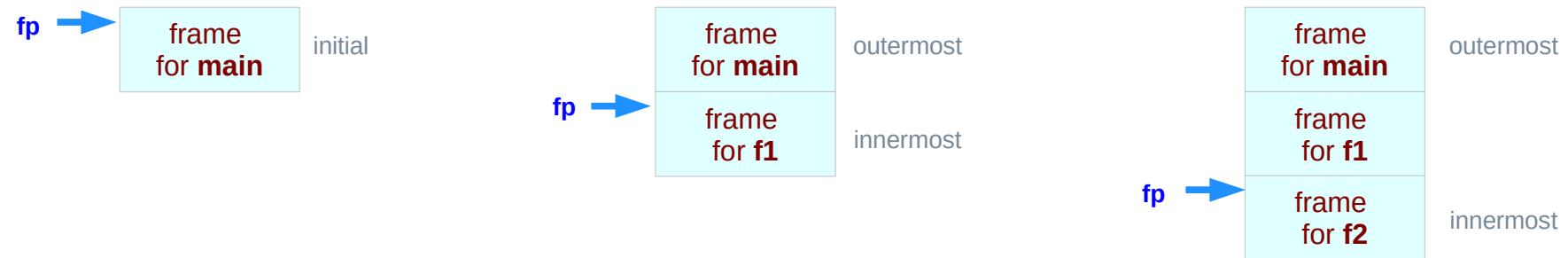
At least, **LR** must not be overwritten

save the followings in a **frame**

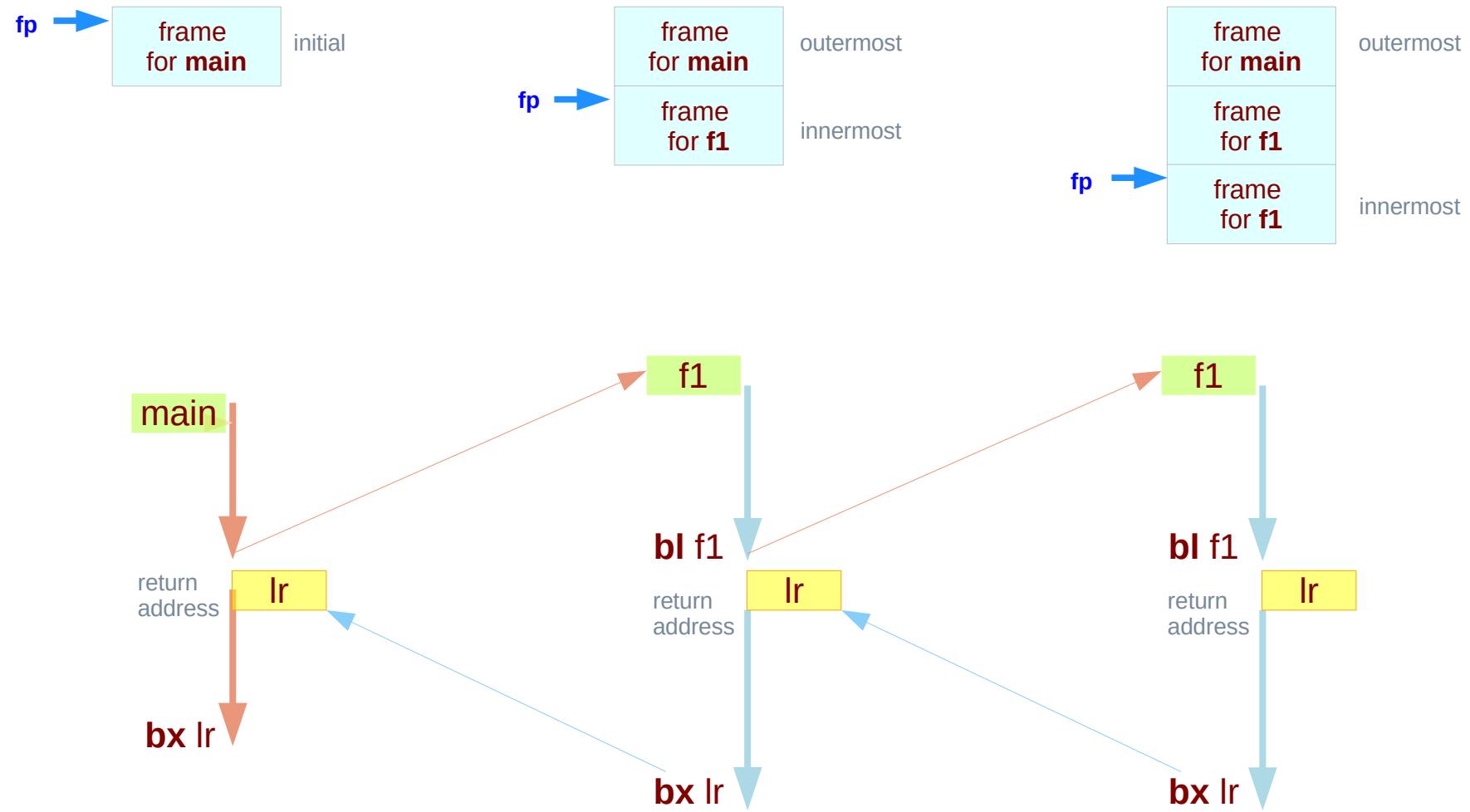
- return address
- arguments
- local variables.



# Nested and recursive function calls



# Recursive function calls



# Activation records (1)

## local variables

- created upon entry to **function**.
- destroyed when function returns.

each **invocation** of a function has  
its own instantiation of **local variables**.

- recursive and nest calls to a function require several instantiations to exist simultaneously.
- functions return only after all functions it calls have returned last-in-first-out(**LIFO**) behavior.
- a **LIFO** structure called a **stack** is used to hold each instantiation.

the portion of the stack used for an **invocation** of a function is called the function's **stack frame** or **activation record**

<https://www.cs.princeton.edu/courses/archive/spring03/cs320/notes/7-1.pdf>

# Activation records (2)

## a stack frame

a frame of data that gets pushed onto the stack.

## a call stack

divided up into contiguous pieces called **stack frames**  
which represent a **function call** and its **argument** data.

- return address
- arguments
- local variables.

architecture-dependent.

processor knows the size of each frame  
and moves the **stack pointer** accordingly  
as **frames** are pushed and popped off the stack.

<https://stackoverflow.com/questions/10057443/explain-the-concept-of-a-stack-frame-in-a-nutshell>

# Activation records (3)

when your program is started,  
the **call stack** has only one frame,  
that of the function **main()**.  
the **initial frame** or the **outermost frame**.

each time a function is called,  
a new frame is added.  
each time a function returns,  
the frame for that function call is eliminated.

for a recursive function,  
there can be many frames for the same function.

the frame for the currently executing function  
is called the **innermost frame**.  
the most recently created frame

[http://www.qnx.com/developers/docs/qnxcar2/index.jsp?topic=%2Fcom.qnx.doc.neutrino.prog%2Ftopic%2Fusing\\_gdb\\_StackFrames.html](http://www.qnx.com/developers/docs/qnxcar2/index.jsp?topic=%2Fcom.qnx.doc.neutrino.prog%2Ftopic%2Fusing_gdb_StackFrames.html)

# Activation records (4)

A **stack frame** consists of many bytes

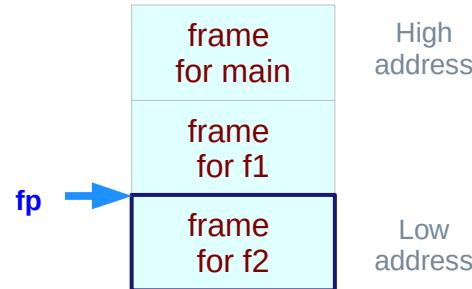
**stack frames** are identified by their addresses.

**the address of the frame** depends on architectures

Usually this address is kept in a register

called the **frame pointer register fp**

while execution is going on in that frame.

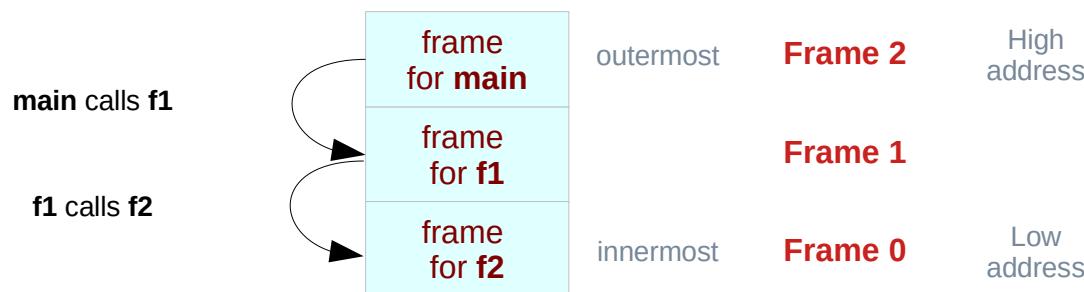


[http://www.qnx.com/developers/docs/qnxcar2/index.jsp?topic=%2Fcom.qnx.doc.neutrino.prog%2Ftopic%2Fusing\\_gdb\\_StackFrames.html](http://www.qnx.com/developers/docs/qnxcar2/index.jsp?topic=%2Fcom.qnx.doc.neutrino.prog%2Ftopic%2Fusing_gdb_StackFrames.html)

# Activation records (5)

**GDB** assigns numbers to all existing stack frames,  
starting with **0** for the **innermost** frame,  
**1** for the frame that called it, and so on upward.

These numbers don't really exist in your program;  
they're assigned by GDB to give you  
a way of designating stack frames in GDB commands.



[http://www.qnx.com/developers/docs/qnxcar2/index.jsp?topic=%2Fcom.qnx.doc.neutrino.prog%2Ftopic%2Fusing\\_gdb\\_StackFrames.html](http://www.qnx.com/developers/docs/qnxcar2/index.jsp?topic=%2Fcom.qnx.doc.neutrino.prog%2Ftopic%2Fusing_gdb_StackFrames.html)

# Activation records (6)

## a call stack

a stack data structure that stores information about the **active subroutines** of a computer program.

Although maintenance of the **call stack** is important for the proper functioning of most software, the details are normally **hidden** and **automatic** in high-level programming languages.

Many computer instruction sets provide **special instructions** for manipulating stacks.

also known as an

- execution stack
- program stack
- control stack
- run-time stack
- machine stack

[https://en.wikipedia.org/wiki/Call\\_stack](https://en.wikipedia.org/wiki/Call_stack)

# Activation records (7)

A **call stack** is used for several related purposes, but the main reason for having one is to keep track of the point to which each **active subroutine** should return control when it finishes executing.

An **active subroutine** is one that has been called, but is yet to complete execution, after which control should be handed back to the point of call.

Such **activations** of subroutines may be nested to any level (recursive as a special case), hence the **stack structure**.

[https://en.wikipedia.org/wiki/Call\\_stack](https://en.wikipedia.org/wiki/Call_stack)

# Argument, scratch, variable, return result registers

**R0 – R3, R12 :**

argument or scratch registers

that are not preserved by the **callee** on a procedure call

**R4 – R11**

8 variable registers that must be preserved on a procedure call  
(if used, the **callee** must save and restore them)

**R0, R1 :**

return result registers

The called performs the calculations,  
places the result (if any) in **R0** and **R1**  
and returns control to the caller using **MOV PC, LR**

# Argument, scratch, variable, return result registers

Registers that is preserved across a procedure

variable registers **R4 – R11**

stack pointer register **sp**

link register **lr**

stack above the stack pointer

Registers that is not preserved across a procedure

argument registers **R0 – R3**

intra procedure call scratch register **r12**

stack below the stack pointer

# Frame pointer and stack pointer registers (1)

**LR (R14, link register, )**

where you were

**PC (R15, program counter)**

where you are

**FP (R11, frame pointer)**

where the stack was

**SP (R13, stack pointer)**

where the stack is

<https://stackoverflow.com/questions/15752188/arm-link-register-and-frame-pointer>

# APCS Register Use Convention

High Address ↑	R15	PC	Program counter
	R14	LR	Link address / scratch register
	R13	SP	Lower end of current stack frame
	R12	IP	Scratch register / specialist use by linker
Low Address	R11	FP	Frame Pointer

# Frame pointer and stack pointer registers (4)

The basic frame layout is,

- fp[-0] saved pc, where we stored this frame.
- fp[-1] saved lr, the return address for this function.
- fp[-2] previous sp, before this function eats stack.
- fp[-3] previous fp, the last stack frame.
- many optional registers...

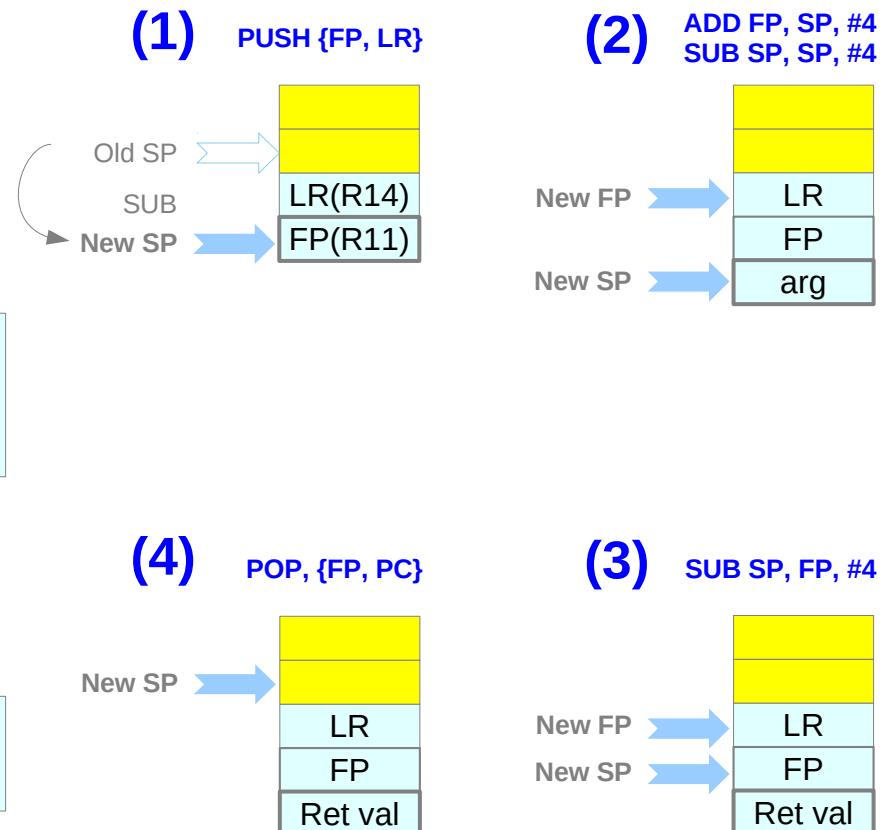
<https://stackoverflow.com/questions/15752188/arm-link-register-and-frame-pointer>

# Stack frame example A

```
Int add(int a, int b) {  
    int c;  
    c = a + b;  
    some_func(a,b);  
    return c;  
}
```

```
0x00010414 <+0>  
0x00010418 <+4>  
0x0001041c <+8>  
0x00010420 <+12>  
0x00010424 <+16>  
0x00010428 <+20>  
0x0001042c <+24>  
0x00010430 <+28>  
0x00010434 <+32>
```

push	{fp, lr}
add	fp, sp, #4
sub	sp, sp, #4
add	r3, r0, r1
str	r3, [fp-#8]
<b>bl</b>	<b>some_func</b>
str	r0, [fp-#8]
sub	sp, fp, #4
pop	{fp, pc}



<https://lloydrochester.com/post/c/stack-of-frames-arm/>

# Stack frame example B

```
int one(int, int);
int two(int, int);
int three(int, int);

Int main(int argc, char *argv[])
{
    int ia, ib, ic;

    ia = 1;
    ib = 2;
    ic = one(ia, ib);

    return ic;
}
```

```
Int one(int a, int b)
{
    int c;
    c = two(a,b);
    return c;
}

Int two(int a, int b)
{
    int c;
    c = three(a,b);
    return c;
}

Int three(int a, int b)
{
    int c;
    c = a+b;
    return c;
}
```

<https://lloydrochester.com/post/c/stack-of-frames-arm/>

# Disassembly of main

(gdb) disassemble **main**

Dump of assembler code for function main:

0x000103d0 <+0>:	push	{r11, lr}	; lr=0xbfe84718 r11 at lowest address
0x000103d4 <+4>:	add	r11, sp, #4	; r11=fp=0x0
0x000103d8 <+8>:	sub	sp, sp, #24	; sp=0xbefff4d8, frame is size 28=24+4
0x000103dc <+12>:	str	r0, [r11, #-24]	; 0xffffffffe8
0x000103e0 <+16>:	str	r1, [r11, #-28]	; 0xffffffffe4
0x000103e4 <+20>:	mov	r3, #1	
0x000103e8 <+24>:	str	r3, [r11, #-8]	
0x000103ec <+28>:	mov	r3, #2	
0x000103f0 <+32>:	str	r3, [r11, #-12]	
0x000103f4 <+36>:	ldr	r1, [r11, #-12]	
0x000103f8 <+40>:	ldr	r0, [r11, #-8]	
0x000103fc <+44>:	<b>bl</b>	0x10414 <one>	; here the lr will be set to 0x00010400
0x00010400 <+48>:	str	r0, [r11, #-16]	; r0 has the return value from function one
0x00010404 <+52>:	ldr	r3, [r11, #-16]	
0x00010408 <+56>:	mov	r0, r3	; r0 will return with the value of int ic
0x0001040c <+60>:	sub	sp, r11, #4	; point sp one word above fp
0x00010410 <+64>:	pop	{r11, pc}	; pc will be restored to 0xbfe84718

End of assembler dump.

<https://lloydrochester.com/post/c/stack-of-frames-arm/>

# Disassembly of one

(gdb) disassemble **one**

Dump of assembler code for function one:

0x00010414 <+0>:	push {r11, lr}	; lr=0x00010400 r11=fp=0xbefff4d0
0x00010418 <+4>:	add r11, sp, #4	; r11=fp=0xbefff4d4
0x0001041c <+8>:	sub sp, sp, #16	; sp=0xbefff4c0 frame is size 20=16+4
0x00010420 <+12>:	str r0, [r11, #-16]	
0x00010424 <+16>:	str r1, [r11, #-20]	; 0xfffffec
0x00010428 <+20>:	ldr r1, [r11, #-20]	; 0xfffffec
0x0001042c <+24>:	ldr r0, [r11, #-16]	
0x00010430 <+28>:	<b>bl</b> 0x10448 < <b>two</b> >	; lr will be 0x00010434
0x00010434 <+32>:	str r0, [r11, #-8]	
0x00010438 <+36>:	ldr r3, [r11, #-8]	
0x0001043c <+40>:	mov r0, r3	
0x00010440 <+44>:	sub sp, r11, #4	; point sp one word above fp
0x00010444 <+48>:	pop {r11, pc}	; fp=0xbefff4f4, lr=0x00010400

End of assembler dump.

<https://lloydrochester.com/post/c/stack-of-frames-arm/>

# Disassembly of two

(gdb) disassemble **two**

Dump of assembler code for function two:

0x00010448 <+0>:	push {r11, lr}	; lr=0x00010434, r11=fp=0xbefff4d4
0x0001044c <+4>:	add r11, sp, #4	; fp=0xbefff4bc
0x00010450 <+8>:	sub sp, sp, #16	; sp=0xbefff4a8 frame is 20=16+4 words
0x00010454 <+12>:	str r0, [r11, #-16]	
0x00010458 <+16>:	str r1, [r11, #-20]	; 0xfffffec
0x0001045c <+20>:	ldr r1, [r11, #-20]	; 0xfffffec
0x00010460 <+24>:	ldr r0, [r11, #-16]	
0x00010464 <+28>:	<b>bl</b> 0x1047c <three>	; lr will be set to 0x00010468
0x00010468 <+32>:	str r0, [r11, #-8]	
0x0001046c <+36>:	ldr r3, [r11, #-8]	
0x00010470 <+40>:	mov r0, r3	
0x00010474 <+44>:	sub sp, r11, #4	
0x00010478 <+48>:	pop {r11, pc}	

End of assembler dump.

<https://lloydrochester.com/post/c/stack-of-frames-arm/>

# Disassembly of three

(gdb) disassemble **three**

Dump of assembler code for function three:

0x0001047c <+0>:	push	{r11}	
0x00010480 <+4>:	add	r11, sp, #0	; (str r11, [sp, #-4]!) NOTICE no lr!!
0x00010484 <+8>:	sub	sp, sp, #20	; dont add #4 here since no fp=0xbefff4a4
0x00010488 <+12>:	str	r0, [r11, #-16]	; stack is size 20 sp=0xbfff490
0x0001048c <+16>:	str	r1, [r11, #-20]	; 0xfffffec
0x00010490 <+20>:	ldr	r2, [r11, #-16]	
0x00010494 <+24>:	ldr	r3, [r11, #-20]	; 0xfffffec
0x00010498 <+28>:	add	r3, r2, r3	
0x0001049c <+32>:	str	r3, [r11, #-8]	
0x000104a0 <+36>:	ldr	r3, [r11, #-8]	
0x000104a4 <+40>:	mov	r0, r3	
0x000104a8 <+44>:	add	sp, r11, #0	
0x000104ac <+48>:	pop	{r11}	; (ldr r11, [sp], #4)
0x000104b0 <+52>:	<b>bx</b>	lr	; lr=0x10468

End of assembler dump.

<https://lloydrochester.com/post/c/stack-of-frames-arm/>

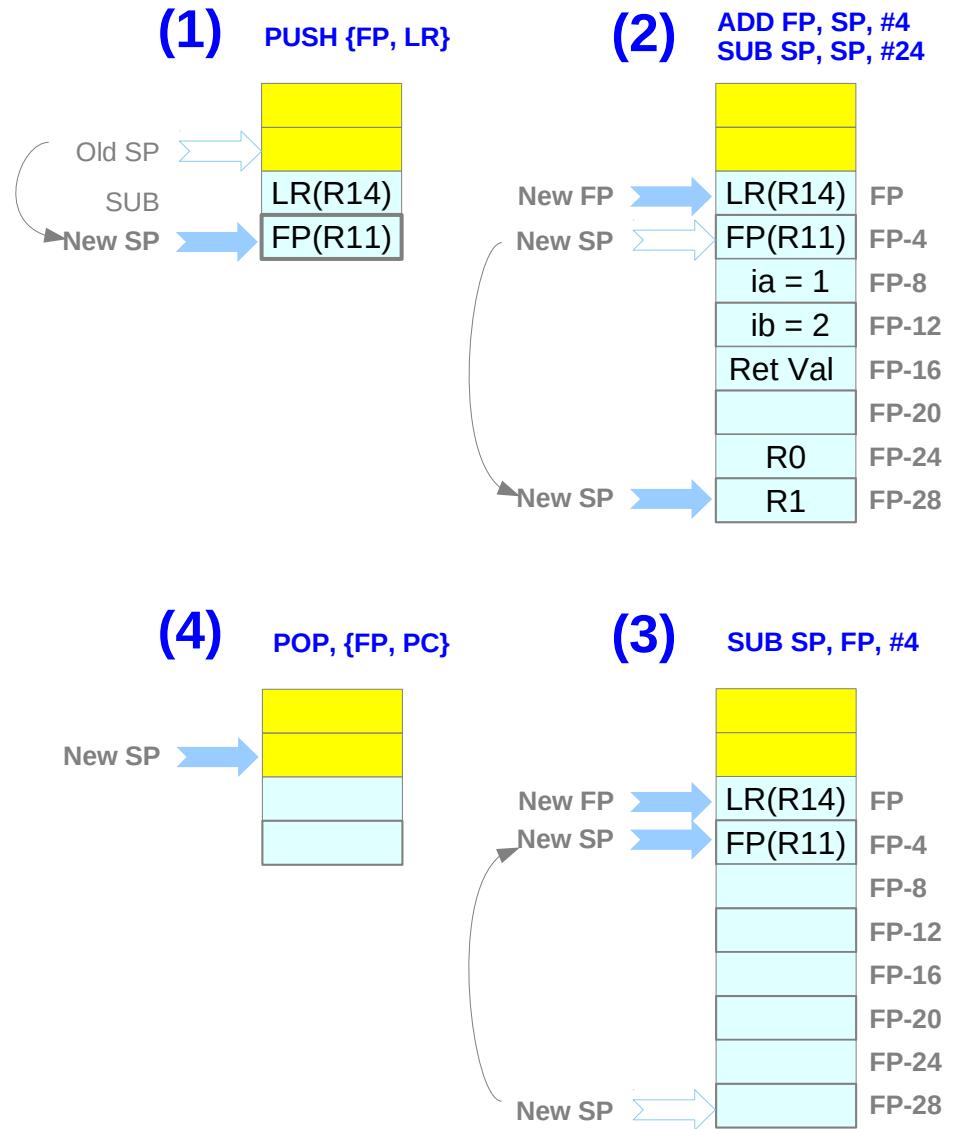
# Stack frame for main

main

```

push {r11, lr}
add r11, sp, #4
sub sp, sp, #24
str r0, [r11, #-24]
str r1, [r11, #-28]
mov r3, #1
str r3, [r11, #-8]
mov r3, #2
str r3, [r11, #-12]
ldr r1, [r11, #-12]
ldr r0, [r11, #-8]
bl 0x10414 <one>
str r0, [r11, #-16]
ldr r3, [r11, #-16]
mov r0, r3
sub sp, r11, #4
pop {r11, pc}

```

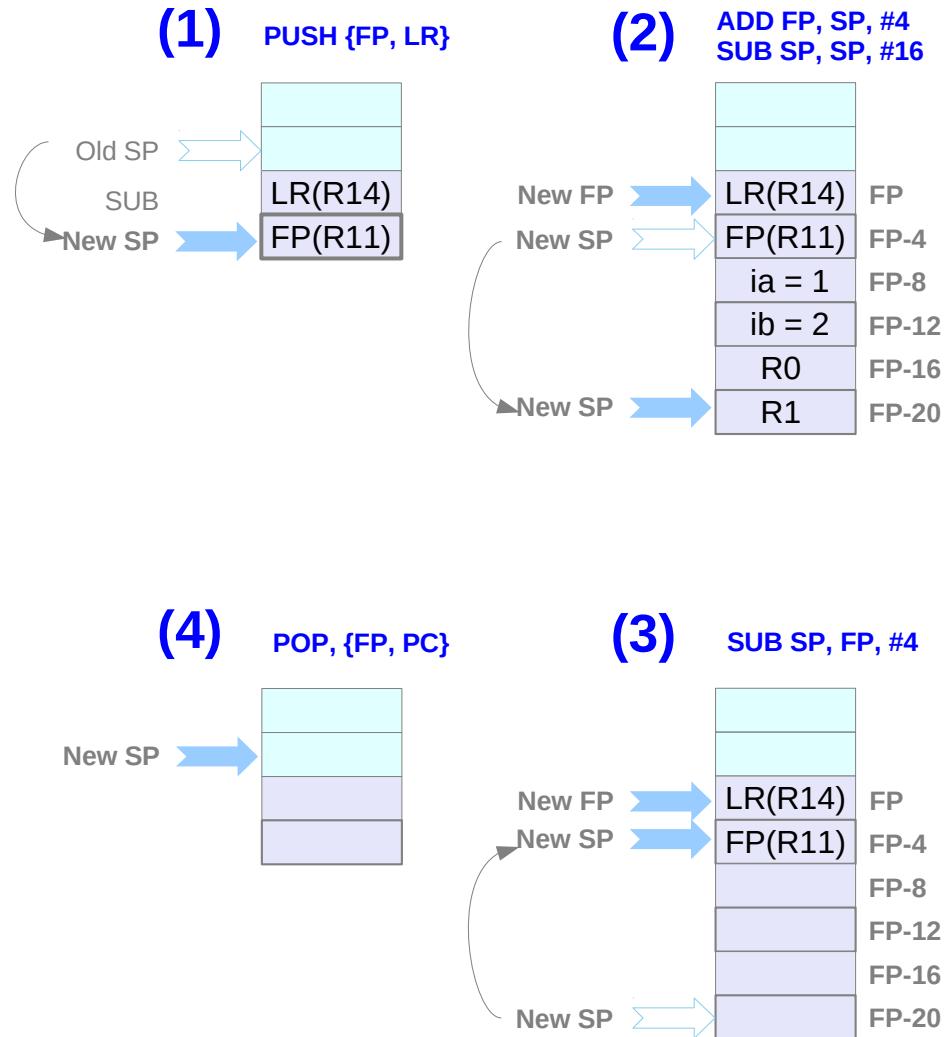


<https://lloydrochester.com/post/c/stack-of-frames-arm/>

# Stack frame for one

one

```
push {r11, lr}
add r11, sp, #4
sub sp, sp, #16
str r0, [r11, #-16]
str r1, [r11, #-20]
ldr r1, [r11, #-20]
ldr r0, [r11, #-16]
bl 0x10448 <two>
str r0, [r11, #-8]
ldr r3, [r11, #-8]
mov r0, r3
sub sp, r11, #4
pop {r11, pc}
```

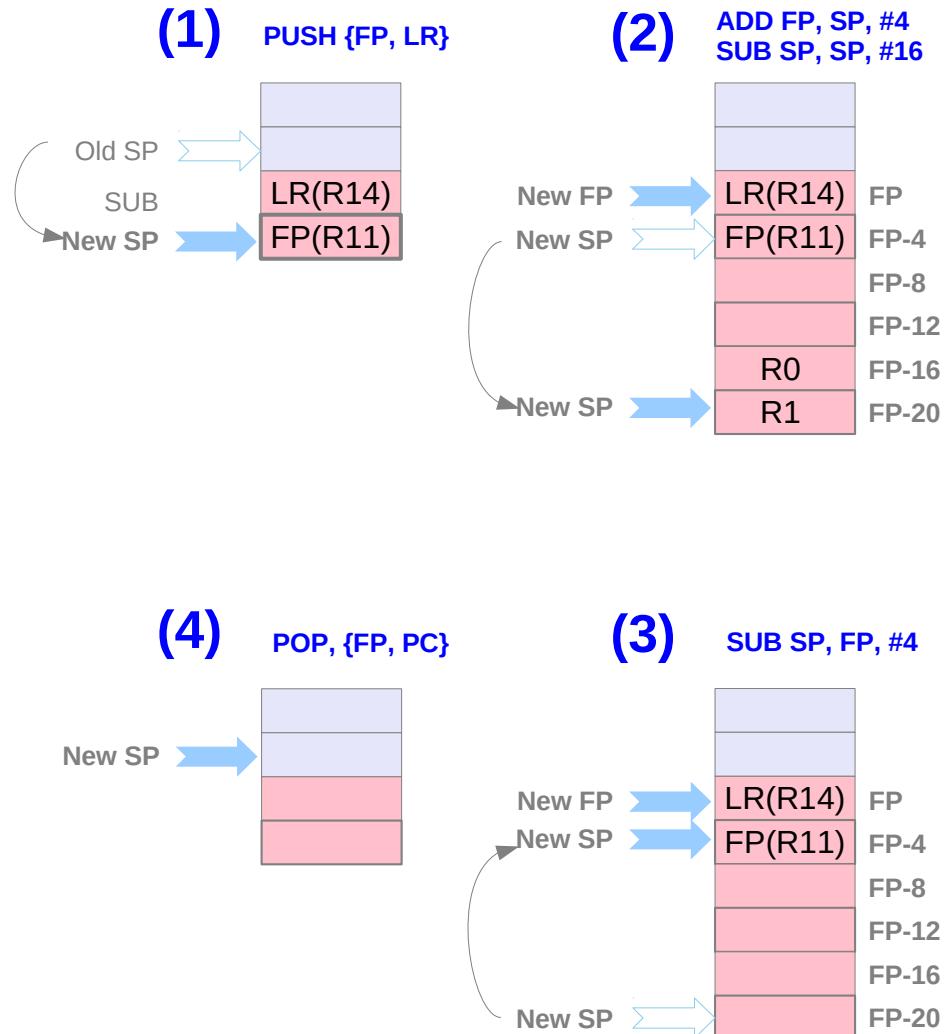


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# Stack frame for two

two

```
push {r11, lr}
add r11, sp, #4
sub sp, sp, #16
str r0, [r11, #-16]
str r1, [r11, #-20]
ldr r1, [r11, #-20]
bl 0x1047c <three>
str r0, [r11, #-8]
ldr r3, [r11, #-8]
mov r0, r3
sub sp, r11, #4
pop {r11, pc}
```

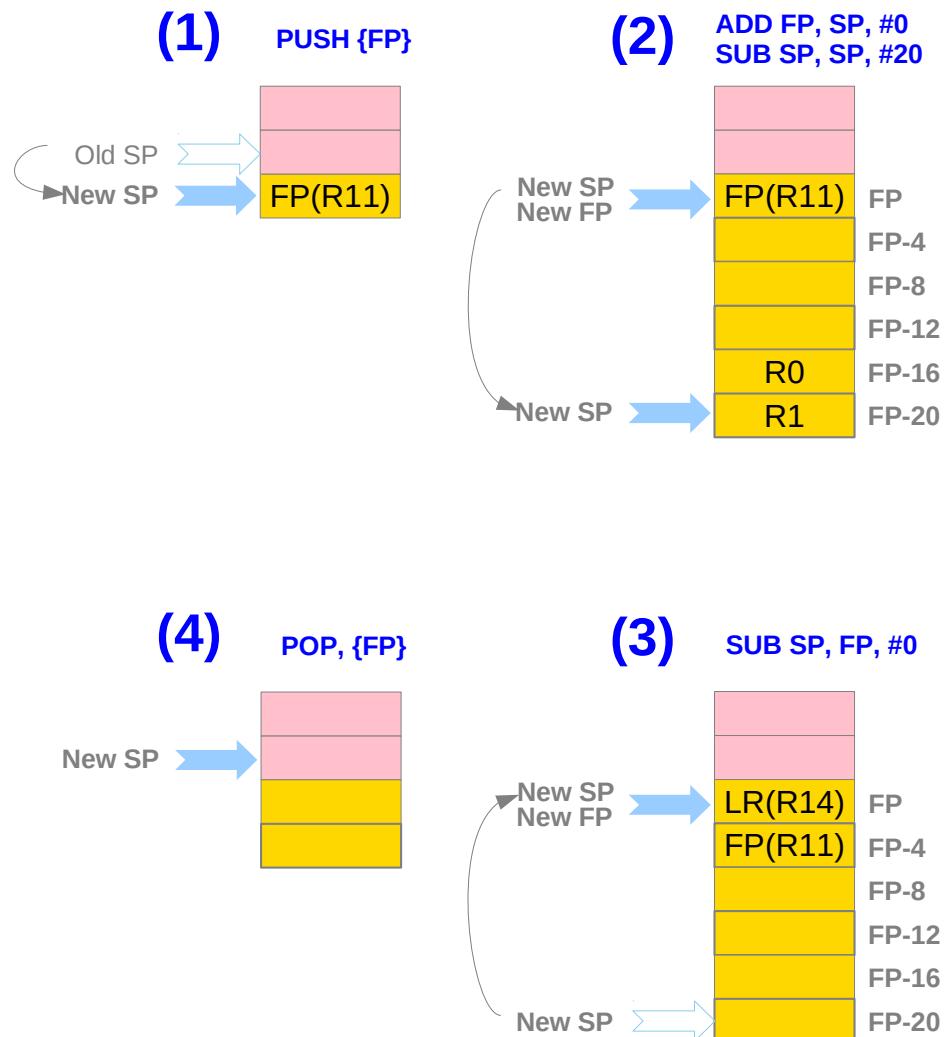


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# Stack frame for three

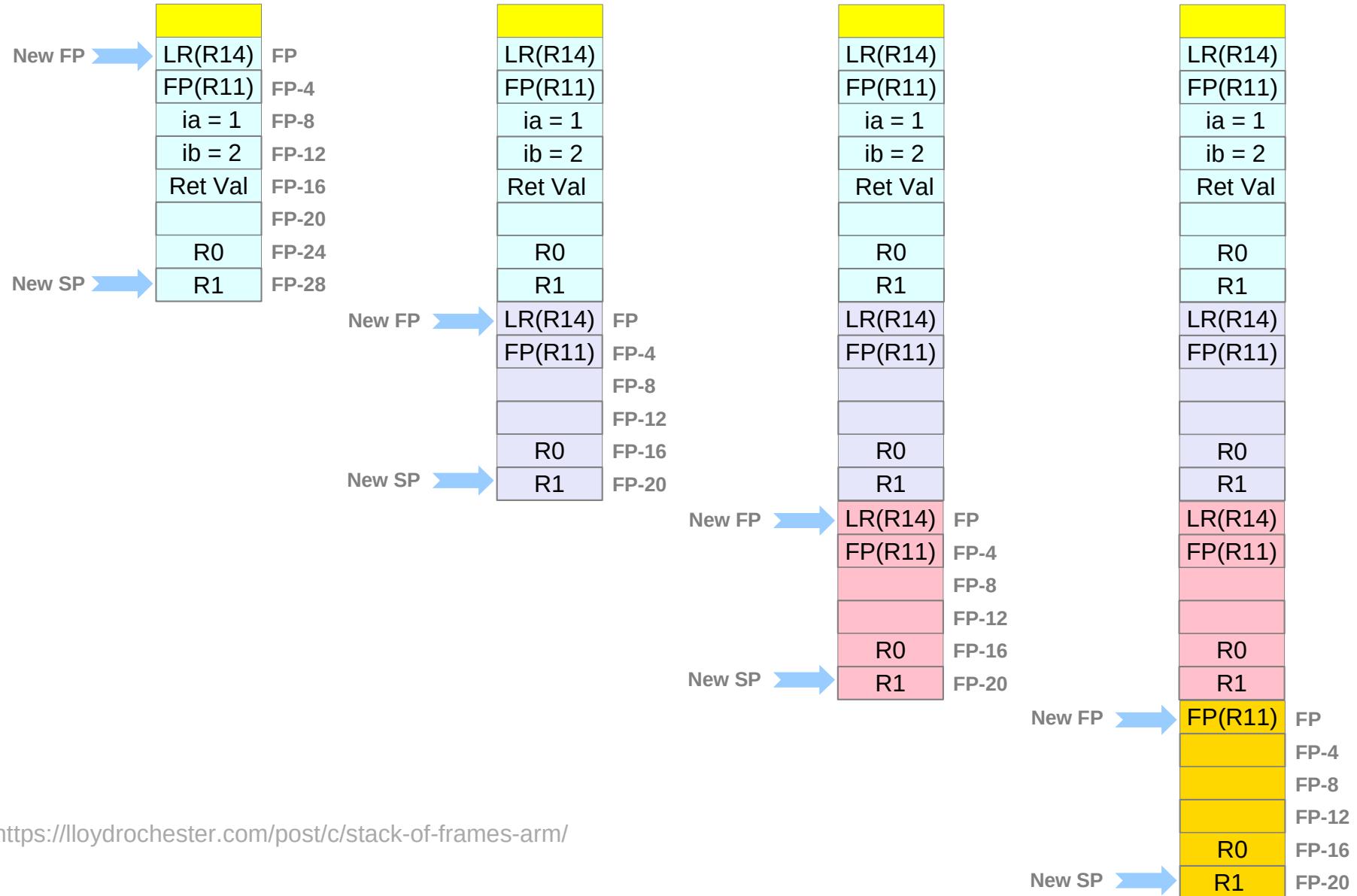
three

```
push {r11}
add r11, sp, #0
sub sp, sp, #20
str r0, [r11, #-16]
str r1, [r11, #-20]
ldr r2, [r11, #-16]
ldr r3, [r11, #-20]
add r3, r2, r3
str r3, [r11, #-8]
ldr r3, [r11, #-8]
mov r0, r3
add sp, r11, #0
pop {r11}
bx lr
```



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# Stack frames



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# Activation records (1)

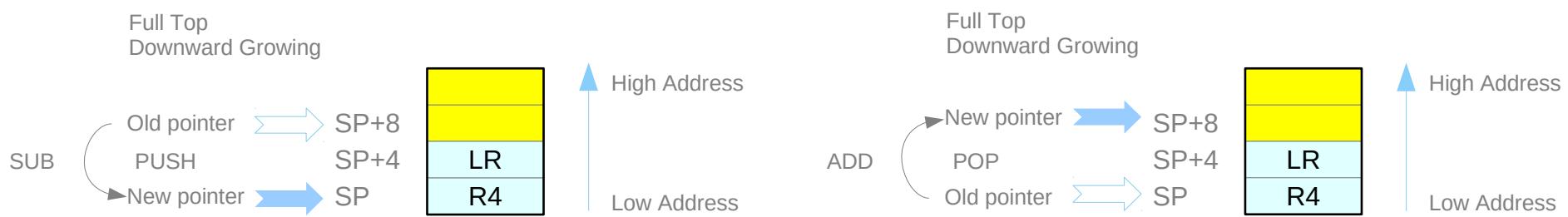
**function:** ; keep callee-saved registers

**push {r4, lr}** ; keep the callee saved registers

...

**pop {r4, lr}** ; restore the callee saved registers

**bx lr** ; return from the function



<https://thinkingeek.com/2013/02/07/arm-assembler-raspberry-pi-chapter-10/>

# Activation records (2)

function:

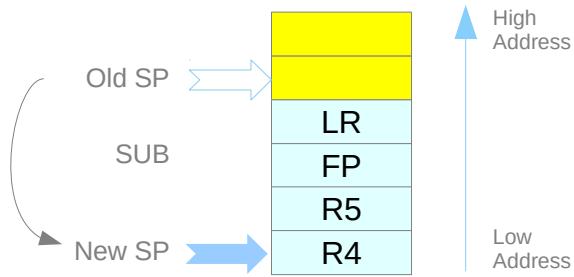
```
push {r4, r5, fp, lr}  
mov fp, sp  
...  
mov sp, fp  
pop {r4, r5, fp, lr}  
  
bx lr
```

; keep callee-saved registers  
; keep the callee saved registers.  
; we added **r5** to keep the stack 8-byte aligned  
; but the important thing here is **fp**  
; **fp** ← **sp**. Keep dynamic link in **fp**  
; code of the function  
; **sp** ← **fp**. Restore dynamic link in **fp**  
; restore the callee saved registers.  
; this will restore **fp** as well  
; return from the function

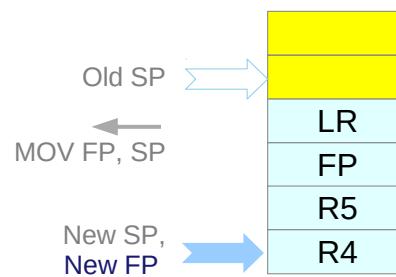
<https://thinkingeek.com/2013/02/07/arm-assembler-raspberry-pi-chapter-10/>

# Activation records (2)

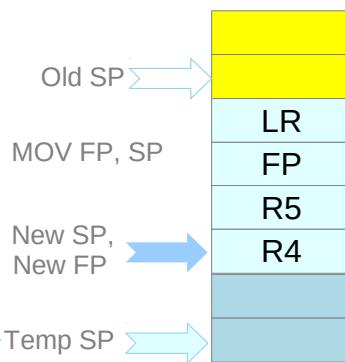
1. **push {r4, r5, fp, lr}**



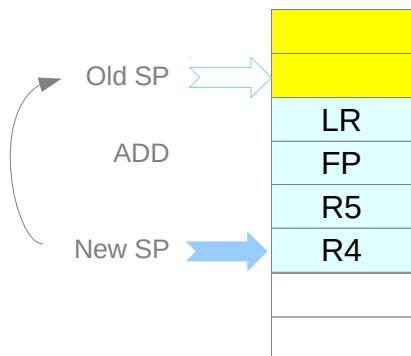
2. **mov fp, sp**



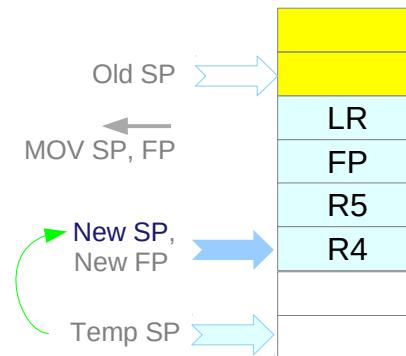
3. **function code**



5. **pop {r4, r5, fp, lr}**



4. **mov sp, fp**



<https://thinkingeek.com/2013/02/07/arm-assembler-raspberry-pi-chapter-10/>

# Activation records (3)

function:

```
push{r4, r5, fp, lr}  
mov fp, sp  
sub sp, sp, #8
```

...

```
mov sp, fp  
pop {r4, r5, fp, lr}
```

```
bx lr
```

; keep callee-saved registers  
; keep the callee saved registers.  
; w added r5 to keep the stack 8-byte aligned  
; but the important thing here is fp  
; fp ← sp. Keep dynamic link in fp  
; enlarge the stack by 8 bytes  
; code of the function  
; sp ← fp. restore dynamic link in fp  
; restore the callee saved registers.  
; this will restore fp as well  
; return from the function

<https://thinkingeek.com/2013/02/07/arm-assembler-raspberry-pi-chapter-10/>

# Local Data Generating Examples

```
void sq(int *c)
{
    (*c) = (*c) * (*c);
}
```

```
int sq_sum5(int a, int b, int c, int d, int e)
{
    sq(&a);
    sq(&b);
    sq(&c);
    sq(&d);
    sq(&e);
    return a + b + c + d + e;
}
```

```
...
    sq_sum5(1, 2, 3, 4, 5);
...
```

callee  
function

- **sq** received a reference
- registers do not have an address
- allocate temporary local storage

caller  
function

# Callee Function Code

```
sq_sum5:  
push { fp, lr }  
mov fp, sp  
sub sp , sp , #16  
  
str r0, [ fp, #-16 ] *( fp - 16 ) <- r0  
str r1, [ fp, #-12 ] *( fp - 12 ) <- r1  
str r2, [ fp, #-8 ] *( fp - 8 ) <- r2  
str r3, [ fp, #-4 ] *( fp - 4 ) <- r3
```

```
mov sp , fp  
pop { fp, lr }  
bx lr
```

```
sq:  
ldr r1, [ r0 ] r1 <- (*r0 )  
mul r1, r1, r1 r1 <- r1 * r1  
str r1, [ r0 ] (*r0 ) <- r1  
bx lr
```

```
sub r0, fp, #16 r0 <- fp - 16  
bl sq call sq ( &a )  
sub r0, fp, #12 r0 <- fp - 12  
bl sq call sq ( &b )  
sub r0, fp, #8 r0 <- fp - 8  
bl sq call sq ( &c )  
sub r0, fp, #4 r0 <- fp - 4  
bl sq call sq ( &d )  
add r0, fp, #8 r0 <- fp + 8  
bl sq call sq ( &e )
```

```
ldr r0, [ fp, #-16 ] r0 <- *( fp - 16 ) :a  
ldr r1, [ fp, #-12 ] r1 <- *( fp - 12 ) :b  
add r0, r0, r1 r0 <- r0 + r1  
ldr r1, [ fp, #-8 ] r1 <- *( fp - 8 ) :c  
add r0, r0, r1 r0 <- r0 + r1  
ldr r1, [ fp, #-4 ] r1 <- *( fp - 4 ) :d  
add r0, r0, r1 r0 <- r0 + r1  
ldr r1, [ fp, #8 ] r1 <- *( fp + 8 ) :e  
add r0, r0, r1 r0 <- r0 + r1
```

# Caller Function Code

```
.data  
.align 4  
  
message:  
.asciz "Sum of 1^2 + 2^2 + 3^2 + 4^2 +  
5^2 is %d\n"
```

```
.text  
  
sq: <<defined above>>  
sq_sum5: <<defined above>>
```

```
.globl main  
main:
```

```
push { r4, lr }
```

```
pop { r4, lr }
```

```
bx lr
```

```
mov r0, #1      a ← 1  
mov r1, #2      b ← 2  
mov r2, #3      c ← 3  
mov r3, #4      d ← 4  
  
mov r4, #5      r4 ← 5  
  
sub sp , sp , #8  
str r4, [sp]     e ← 5  
  
bl sq_sum5     sq_sum5 ( 1, 2, 3, 4, 5 )  
  
add sp , sp , #8  
  
mov r1, r0  
ldr r0, address_of_message  
  
bl printf  
  
address_of_message: . word message
```

```
void sq(int *c) {  
    (*c) = (*c) * (*c);  
}
```

sq:

```
ldr r1, [r0]      ; r1 ← (*r0)      ; r0 : argument register  
mul r1, r1, r1   ; r1 ← r1 * r1  
str r1, [r0]      ; (*r0) ← r1  
bx lr             ; return from the function
```

<https://thinkingeek.com/2013/02/07/arm-assembler-raspberry-pi-chapter-10/>

# sq\_sum5 (1)

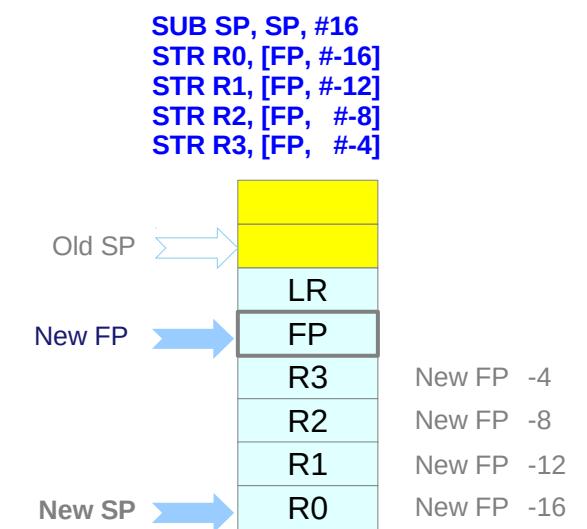
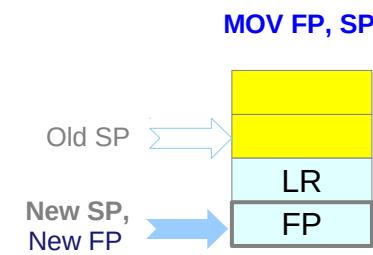
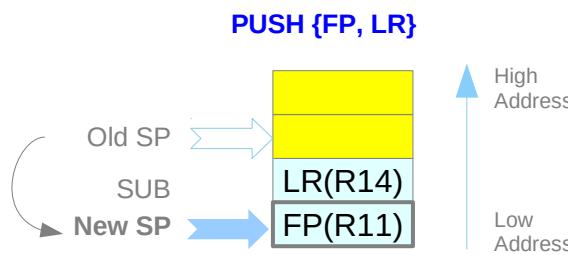
```
int sq_sum5(int a, int b, int c, int d, int e) {  
    sq(&a);  
    sq(&b);  
    sq(&c);  
    sq(&d);  
    sq(&e);  
    return a + b + c + d + e;  
}
```

<https://thinkingeek.com/2013/02/07/arm-assembler-raspberry-pi-chapter-10/>

# sq\_sum5 (2)

sq\_sum5:

```
push {fp, lr}          ; keep fp and all callee-saved registers.  
mov fp, sp            ; set the dynamic link  
                      ; allocate space for 4 integers in the stack  
                      ; keep parameters in the stack  
sub sp, sp, #16        ; sp ← sp - 16.  
str r0, [fp, #-16]      ; *(fp - 16) ← r0  
str r1, [fp, #-12]      ; *(fp - 12) ← r1  
str r2, [fp, #-8]       ; *(fp - 8) ← r2  
str r3, [fp, #-4]       ; *(fp - 4) ← r3
```



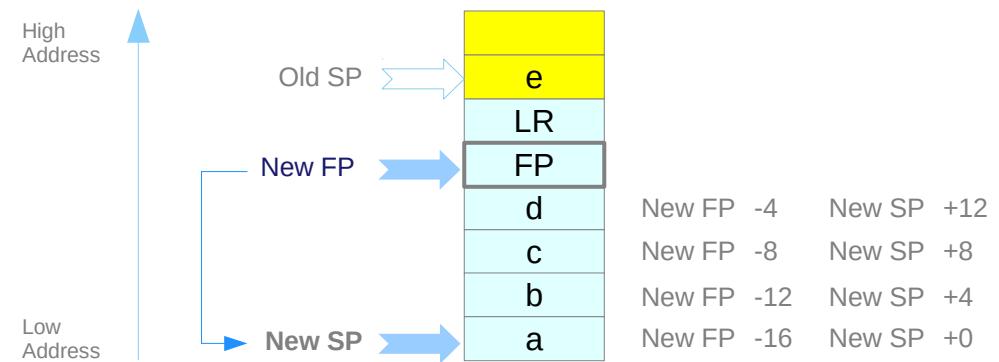
<https://thinkingeek.com/2013/02/07/arm-assembler-raspberry-pi-chapter-10/>

# sq\_sum5 (3)

Value	Address(es)	
a	[fp, #-16]	[sp]
b	[fp, #-12]	[sp, #4]
c	[fp, #-8]	[sp, #8]
d	[fp, #-4]	[sp, #12]
fp(r11)	[fp]	[sp, #16]
lr(r14)	[fp, #4]	[sp, #20]
e	[fp, #8]	[sp, #24]

High Address

fp[-0] saved pc  
fp[-1] saved lr  
fp[-2] previous sp  
fp[-3] previous fp



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# sq\_sum5 (4)

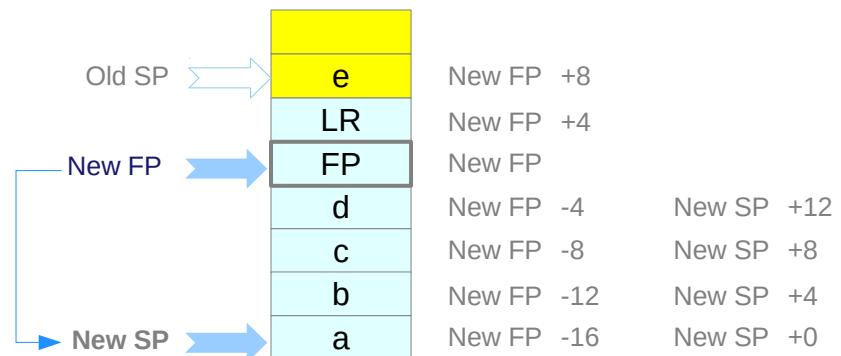
```
sub r0, fp, #16      ; r0 ← fp - 16
bl  sq               ; call sq(&a);
sub r0, fp, #12      ; r0 ← fp - 12
bl  sq               ; call sq(&b);
sub r0, fp, #8       ; r0 ← fp - 8
bl  sq               ; call sq(&c);
sub r0, fp, #4       ; r0 ← fp - 4
bl  sq               ; call sq(&d)
add r0, fp, #8      ; r0 ← fp + 8
bl  sq               ; call sq(&e)
```



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# sq\_sum5 (5)

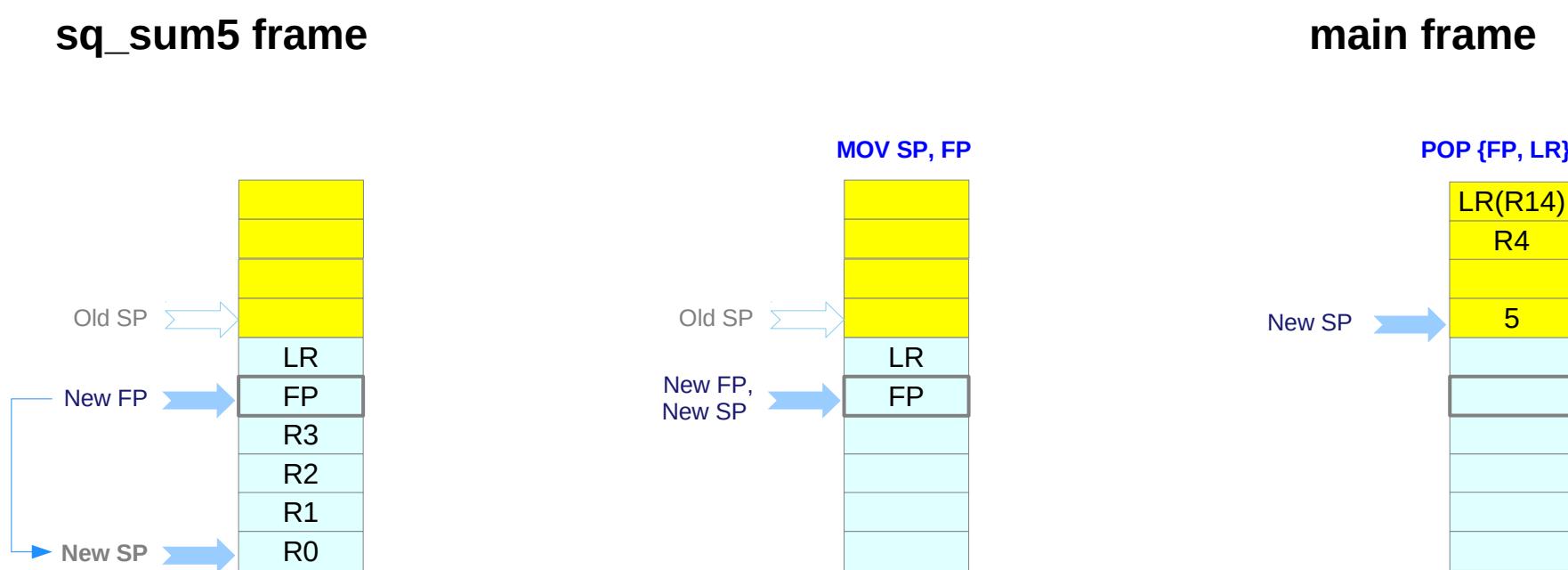
```
ldr r0, [fp, #-16] ; r0 ← *(fp - 16). ; Loads a into r0
ldr r1, [fp, #-12] ; r1 ← *(fp - 12). ; Loads b into r1
add r0, r0, r1      ; r0 ← r0 + r1      ; (a +b)
ldr r1, [fp, #-8]   ; r1 ← *(fp - 8). ; Loads c into r1
add r0, r0, r1      ; r0 ← r0 + r1      ; (a +b +c)
ldr r1, [fp, #-4]   ; r1 ← *(fp - 4). ; Loads d into r1
add r0, r0, r1      ; r0 ← r0 + r1      ; (a +b +c +d)
ldr r1, [fp, #8]    ; r1 ← *(fp + 8). ; Loads e into r1
add r0, r0, r1      ; r0 ← r0 + r1      ; (a +b +c +d +e)
```



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# sq\_sum5 (6)

```
mov sp, fp          ; Undo the dynamic link  
pop {fp, lr}        ; Restore fp and callee-saved registers  
bx lr              ; Return from the function
```



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# main (1)

```
/* squares.s */
.data

.align 4
.message: .asciz "Sum of 1^2 + 2^2 + 3^2 + 4^2 + 5^2 is %d\n"

.text

sq:
<<defined above>>

sq_sum5:
<<defined above>>

.globl main
```

<https://thinkingeek.com/2013/02/07/arm-assembler-raspberry-pi-chapter-10/>

# main (2)

main:

```
push {r4, lr} ; Keep callee-saved registers
```

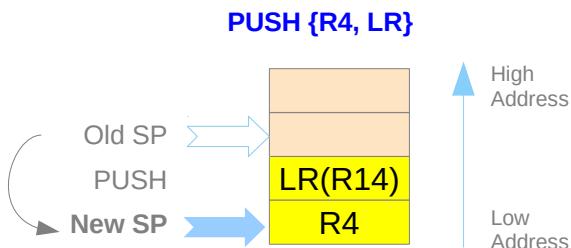
```
; Prepare the call to sq_sum5
```

```
mov r0, #1 ; Parameter r0 ← a=1
```

```
mov r1, #2 ; Parameter r1 ← b=2
```

```
mov r2, #3 ; Parameter r2 ← c=3
```

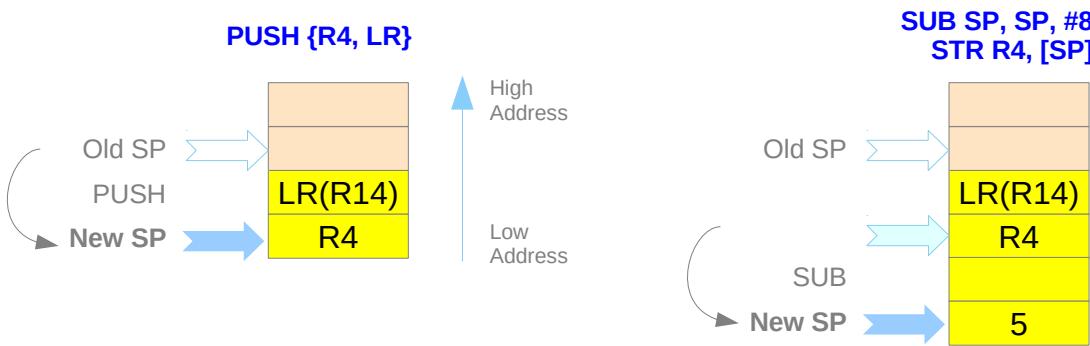
```
mov r3, #4 ; Parameter r3 ← d=4
```



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# main (3)

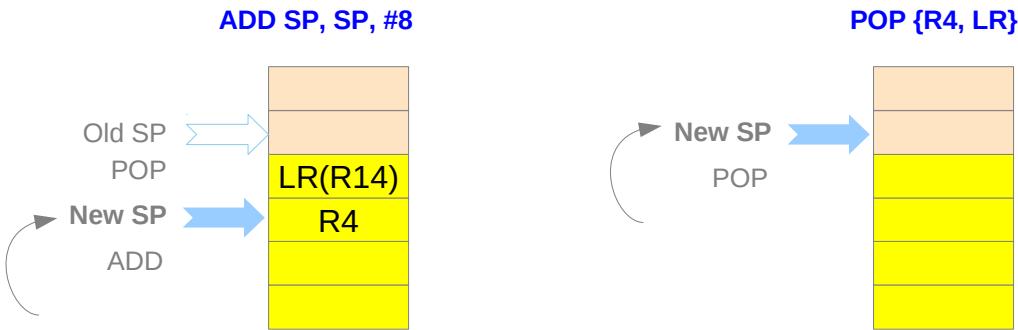
```
; Parameter e goes through the stack,  
; so it requires enlarging the stack  
mov r4, #5           ; r4 ← 5  
sub sp, sp, #8       ; Enlarge the stack 8 bytes,  
                     ; we will use only the  
                     ; topmost 4 bytes  
str r4, [sp]         ; Parameter e ← 5  
bl sq_sum5          ; call sq_sum5(1, 2, 3, 4, 5)
```



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# main (4)

```
add sp, sp, #8      ; Shrink back the stack  
  
; Prepare the call to printf  
mov r1, r0          ; The result of sq_sum5  
ldr r0, address_of_message  
bl printf           ; Call printf  
  
pop {r4, lr}         ; Restore callee-saved registers  
bx lr
```



<https://thinkingeek.com/2013/02/07/arm-assembler-raspberry-pi-chapter-10/>

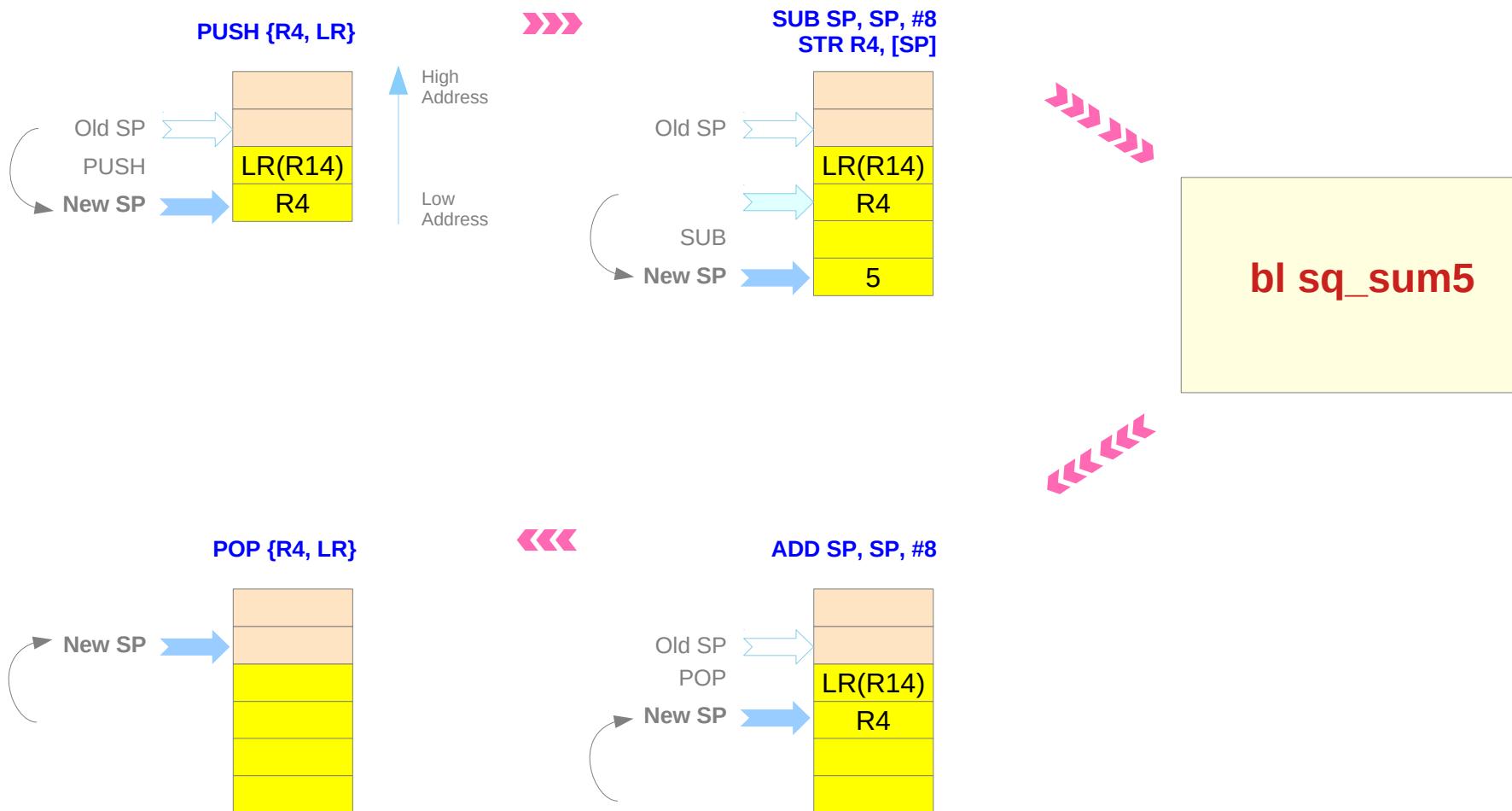
# main (6)

```
address_of_message:    .word    message
message:                .asciz    "Sum of 1^2 + 2^2 + 3^2 + 4^2 + 5^2 is %d\n"
```

```
$ ./square
Sum of 1^2 + 2^2 + 3^2 + 4^2 + 5^2 is 55
```

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# main's stack frame

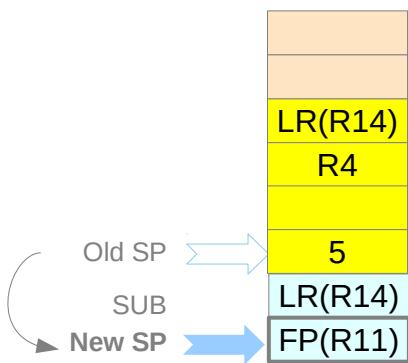


<https://thinkingeek.com/2013/02/07/arm-assembler-raspberry-pi-chapter-10/>

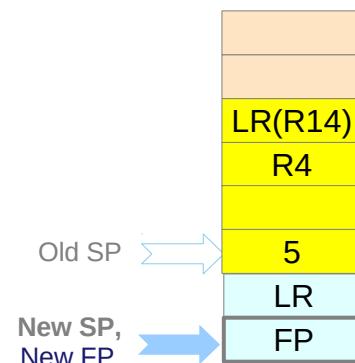
# sq\_sum5's stack frame (1)

sq_sum5:	sub	r0, fp, #16	ldr	r0, [fp, #-16]	mov	sp, fp
	bl	sq	ldr	r1, [fp, #-12]	pop	{fp, lr}
	sub	r0, fp, #12	add	r0, r0, r1	bx	lr
	bl	sq	ldr	r1, [fp, #-8]		
	sub	r0, fp, #8	add	r0, r0, r1		
	bl	sq	ldr	r1, [fp, #-4]		
	sub	r0, fp, #4	add	r0, r0, r1		
	bl	sq	ldr	r1, [fp, #8]		
	add	r0, fp, #8	add	r0, r0, r1		
	bl	sq				

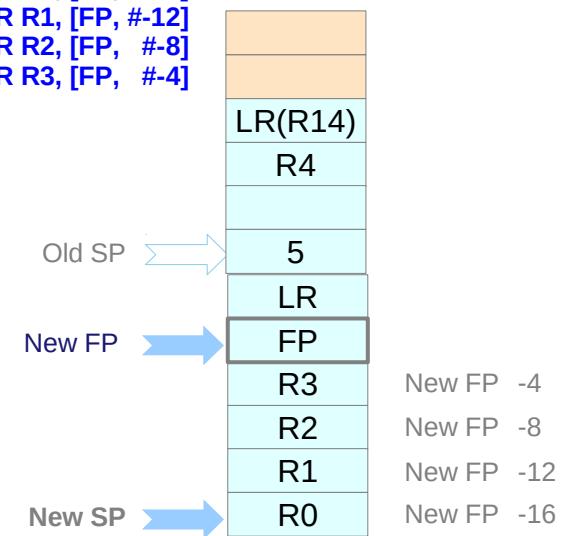
PUSH {FP, LR}



MOV FP, SP



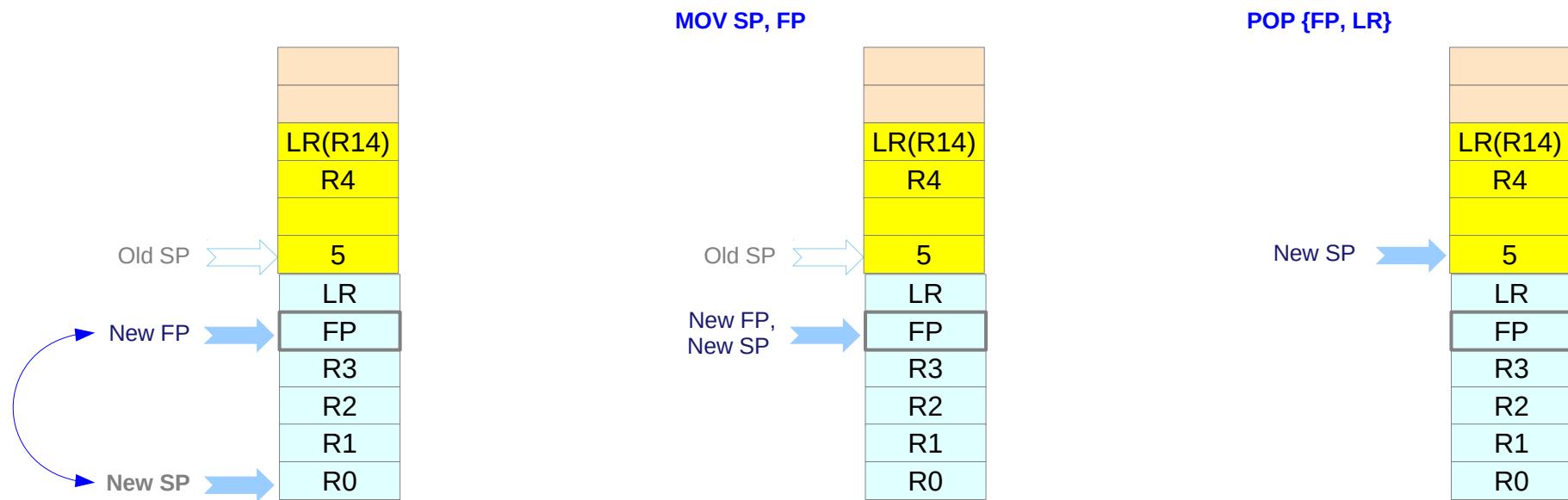
SUB SP, SP, #16  
STR R0, [FP, #-16]  
STR R1, [FP, #-12]  
STR R2, [FP, #-8]  
STR R3, [FP, #-4]



<https://thinkingeek.com/2013/02/07/arm-assembler-raspberry-pi-chapter-10/>

# sq\_sum5's stack frame (2)

sq_sum5:	sub	r0, fp, #16	ldr	r0, [fp, #-16]	mov	sp, fp
push {fp, lr}	bl	sq	ldr	r1, [fp, #-12]	pop	{fp, lr}
mov fp, sp	sub	r0, fp, #12	add	r0, r0, r1	bx	lr
sub sp, sp, #16	bl	sq	ldr	r1, [fp, #-8]		
str r0, [fp, #-16]	sub	r0, fp, #8	add	r0, r0, r1		
str r1, [fp, #-12]	bl	sq	ldr	r1, [fp, #-4]		
str r2, [fp, #-8]	sub	r0, fp, #4	add	r0, r0, r1		
str r3, [fp, #-4]	bl	sq	ldr	r1, [fp, #8]		
	add	r0, fp, #8	add	r0, r0, r1		
	bl	sq				

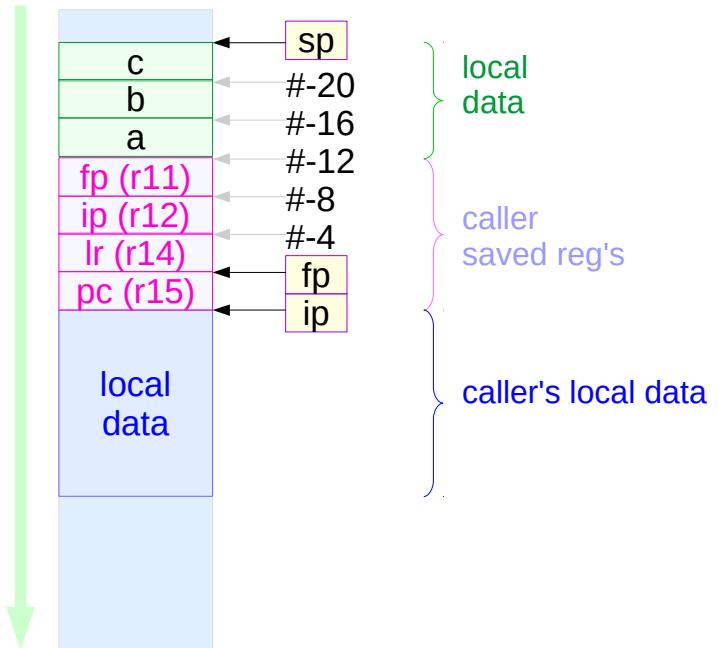


<https://thinkingeek.com/2013/02/07/arm-assembler-raspberry-pi-chapter-10/>

# -fno-omit-frame-pointer

```
main:  
mov    ip, sp  
stmfd  sp!, { fp, ip, lr, pc }  
sub    fp, ip, #4  
sub    sp, sp, #12  
ldr    r2, [fp, #-16]  
ldr    r3, [fp, #-20]  
add    r3, r3, r2  
str    r3, [fp, #-24]  
sub    sp, fp, #12  
ldmfd  sp, {fp, sp, pc}
```

```
main()  
{  
    volatile int a, b, c;  
    c = a + b;  
}
```

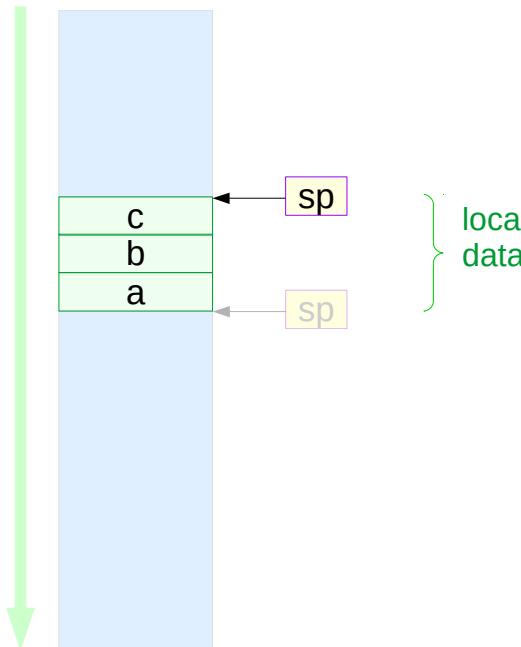


<https://community.arm.com/thread/7092>

# -fomit-frame-pointer

```
main:  
sub    sp, sp, #12  
ldr    r2, [sp, #8]  
ldr    r3, [fp, #4]  
add    r3, r3, r2  
str    r3, [sp, #0]  
sub    sp, sp, #12
```

```
main()  
{  
    volatile int a, b, c;  
    c = a + b;  
}
```



<https://community.arm.com/thread/7092>

## References

- [1] [http://wiki.osdev.org/ARM\\_RaspberryPi\\_Tutorial\\_C](http://wiki.osdev.org/ARM_RaspberryPi_Tutorial_C)
- [2] <http://blog.bobuhir011.net/2014/01-13-baremetal.html>
- [3] <http://www.valvers.com/open-software/raspberry-pi/>
- [4] <https://www.cl.cam.ac.uk/projects/raspberrypi/tutorials/os/downloads.html>