Side Effects (1A)

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Please send corrections (or suggestions) to youngwlim@hotmail.com. This document was produced by using LibreOffice. Haskell in 5 steps https://wiki.haskell.org/Haskell_in_5_steps

Variables

Imperative programming:

- variables as changeable locations in a computer's memory
- **imperative programs** explicitly commands (instructs) the computer what to do

functional programming

- a way to think in higher-level mathematical terms
- defining how variables relate to one another
- the **compiler** will **translate** these **functions** and **variables**

to **instructions** so that the computer can process.

Δ

Haskell Language Features (I)

Haskell Functional Programming (I)

- Immutability
- Recursive Definition : only in functions
- No Data Dependency

Redefinition : not allowed

imperative programming:	r = 5
after setting r = 5 and <u>then</u> changing it to r = 2 .	r = 2
Hakell programming:	
an error: "multiple declarations of r ".	
within a given scope, a variable in Haskell	
are defined only once and cannot change,	
like variables in mathematics.	

r = 5 r = 2 no mutation

in Haskell

Variables in a file

Immutable:

they can <u>change</u> only based on the data we <u>enter</u> to run the program.

We <u>cannot</u> define \mathbf{r} two ways <u>in the same code</u>, but we could <u>change</u> the value by <u>changing the file</u> Vars.hs

a = 100 r = 5 pi = 3.14159 e = 2.7818

No Mutation



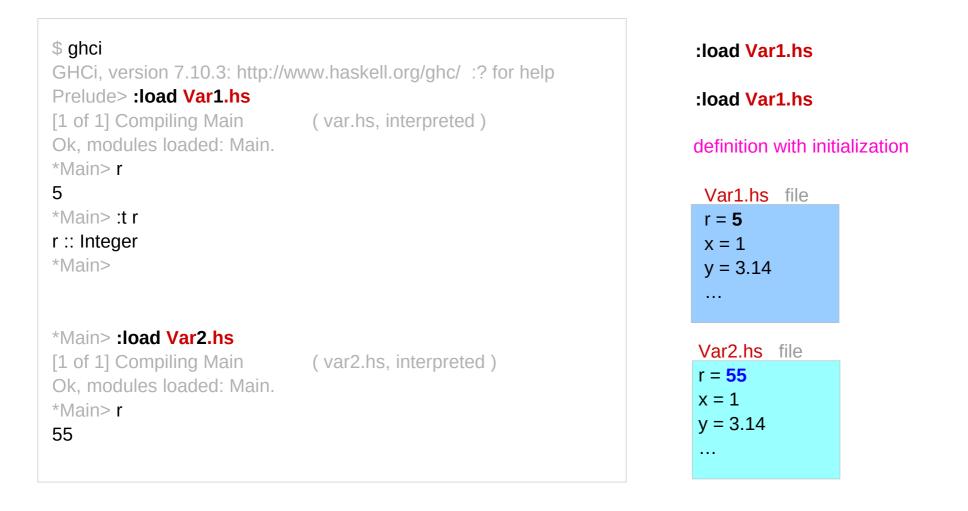
let r = 33

https://en.wikibooks.org/wiki/Haskell/Variables_and_functions

Side Effects (1A)

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Loading a variable definition file



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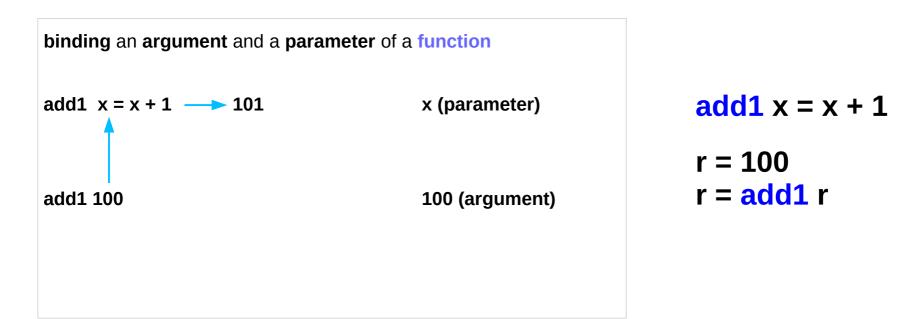
https://en.wikibooks.org/wiki/Haskell/Variables_and_functions

Side Effects (1A)

Incrementing by one

imperative programming:	r = r + 1
incrementing the variable r	
(updating the value in memory)	
Hakell programming:	
No compound assignment like operations	r = 3 r = r + 1
	r=r+1
if r had been defined with any value beforehand,	
then $\mathbf{r} = \mathbf{r} + 1$ in Haskell would bring an error message.	
multiple definition not allowed	as an event and a never star of a function in simple scars
	as an argument and a parameter of a function in simple cases add1 x = x + 1
Use a function	add1 xatd x140 1 x (parameter) = 100 (argument)
	r = 3
	r = add1 r

Arguments and parameters of a function



Recursive Definition

Hakell programming:			
a recursive definition of r (defining it in terms of itself)	a += b a -= b a *= b a /= b	(a = a + b) (a = a - b) (a = a * b) (a = a / b)	r f f
			r

No compound assignment like operations are allowed

if \mathbf{a} had been defined with any value beforehand,

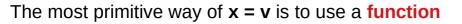
then **a** = **a** + **b** in Haskell would **multiply defined**

recursive function factorial 0 = 1factorial n = n * factorial (n - 1)

non-recursive function add1 x = x + 1

recursive definitions are allowed only in function definition

Simulating imperative codes



taking \mathbf{x} as a **parameter**, and pass the **argument** \mathbf{v} to that function.

i = <mark>s</mark> = 0;	// sum 0100	
while (i <= 100) {		
<mark>s = s+i</mark> ;		
i++;		
}		
return <mark>s</mark> ;		

 sum = f 0 0
 -- the initial values

 where
 -- increment i, augment s

 f i s | i <=100</td>
 = f (i+1) (s+i)

 otherwise
 -- return s at the end

 s = (s+i)

This code is not pretty functional programing code, but it is simulating imperative code

https://stackoverflow.com/questions/43525193/how-can-i-re-assign-a-variable-in-a-function-in-haskell

 $\mathbf{x} = \mathbf{v}$

No Data Dependency

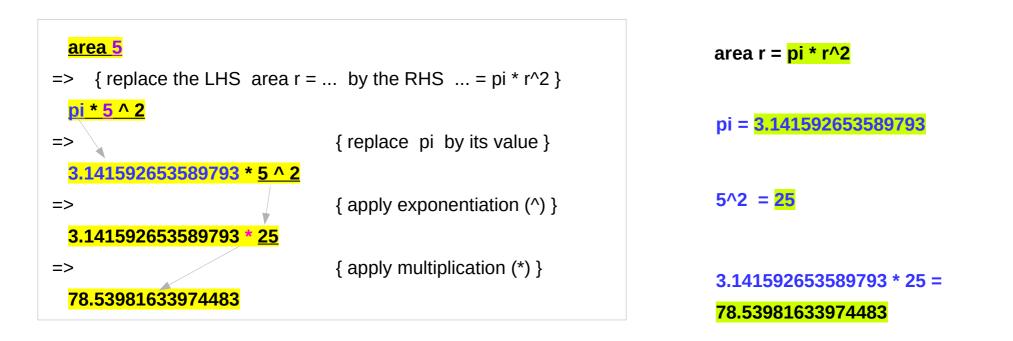


Hakell programming:

because the values of variables do not change variables can be defined <u>in any order</u>

no mandatory : "x being declared before y"

Evaluation examples



Translation to instructions

functional programming • leaving the compiler to translate functions and variables to the step-by-step instructions that the computer can process. replace each function and variable with its definition repeatedly replace the results until a single value remains. to apply or call a function means to replace the LHS of its definition by its RHS.

Scope

Scope rules define the **visibility rules** for **names** in a programming language.

What if you have references to a **variable** named **k** in <u>different parts</u> of the program?

Do these refer to the same variable or to different ones?

https://courses.cs.washington.edu/courses/cse341/03wi/imperative/scoping.html

Haskell Scope

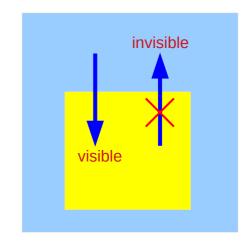
Most languages, including Haskell, are **statically scoped**.

- A **block** defines a new **scope**.
- Variables can be declared in that scope,

and are <u>not visible</u> from the outside.

- However, variables <u>outside the scope</u> (in enclosing scopes) are <u>visible</u> unless they are overridden.
- In Haskell, these scope rules also apply
 - to the names of **functions**.

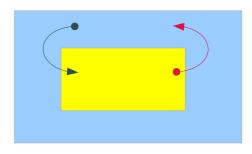
Static scoping is also sometimes called **lexical scoping**.



https://courses.cs.washington.edu/courses/cse341/03wi/imperative/scoping.html

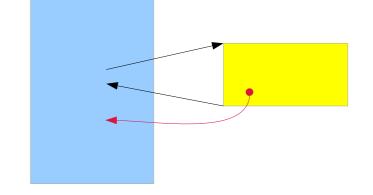
Side Effects Definition

a function or expression is said to have a side effect if it modifies some state outside its scope or has an observable interaction with its calling functions or the outside world besides returning a value.



a particular function might

- modify a **global** variable or **static** variable
- modify one of its arguments
- raise an **exception**
- write data to a display or file
- read data from a keyboard or file
- call other side-effecting functions



https://en.wikipedia.org/wiki/Side_effect_(computer_science)

Side Effects (1A)

Some Monad types to handle side effects

State monad

manages global variables

Error monad

enables exceptions

IO monad

handles interactions with the **file system**, and other **resources** <u>outside</u> the program

the **program** itself has <u>no side effects</u> the **action** in monads does have <u>side effects</u>

the functional nature of the **program** is maintained (**pure**, **no side effects**)

https://blog.osteele.com/2007/12/overloading-semicolon/

actions in State, Error, IO monad have side effects

Side Effects (1A)



History, Order, and Context

In the presence of side effects,

a program's behaviour may depend on history;

the order of evaluation matters.

the context and histories

<u>imperative</u> programming : <u>frequent</u> utilization of **side effects**. **<u>functional</u>** programming : **side effects** are <u>rarely</u> used.

The lack of side effects makes it easier to do **formal verifications** of a program

https://en.wikipedia.org/wiki/Side_effect_(computer_science)

Side Effects Examples in C

// The assignment function returns 10// which automatically casts to "<u>true</u>"// so the loop conditional always evaluates to true

while (**b** = 10) { }

https://en.wikipedia.org/wiki/Side_effect_(computer_science)

Side Effects (1A)

Haskell Language Features (II)

Haskell Functional Programming (II)

- Pure Function
- Simple IO
- Laziness
- Sequencing

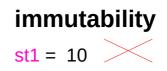
Pure Languages

Haskell is a <mark>pure</mark> language	no side effects	
programs are made of functions	pure functions	
that <u>cannot</u> change		
any global state or variables ,		
they <u>can</u> only		
do some computations and return their results.		
not modify arguments of a function		
every variable's value does <u>not change</u> in time		
However, some problems are <u>inherently</u> stateful		

in that they <u>rely</u> on some <u>state</u> that <u>changes</u> <u>over time</u>.

a bit tedious to model Haskell has the **state monad** features

http://learnyouahaskell.com/for-a-few-monads-more

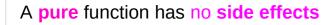


use a function for stateful computations

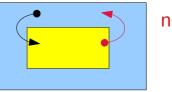
s -> (x,s) st1 (v,10)

Side Effects (1A)

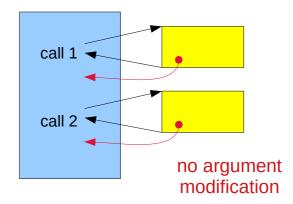
Pure Function



- **no state** nor **no** access to **external states** (global variables)
 - → the function call starts from the scratch (no memory)
 - every invocation with the <u>same set</u> of <u>arguments</u>
 <u>returns</u> always the <u>same result</u>
- no argument modifications
 - → calling a pure function is the <u>same</u> as
 - → calling it twice and <u>discarding</u> the <u>result</u> of the <u>first call</u>.



no global variables

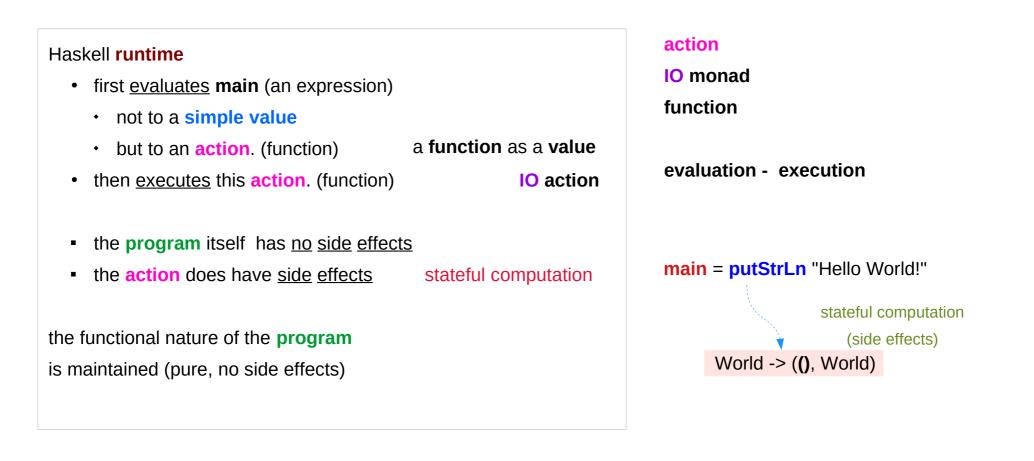


easily parallelizeable

no side effect means no data races



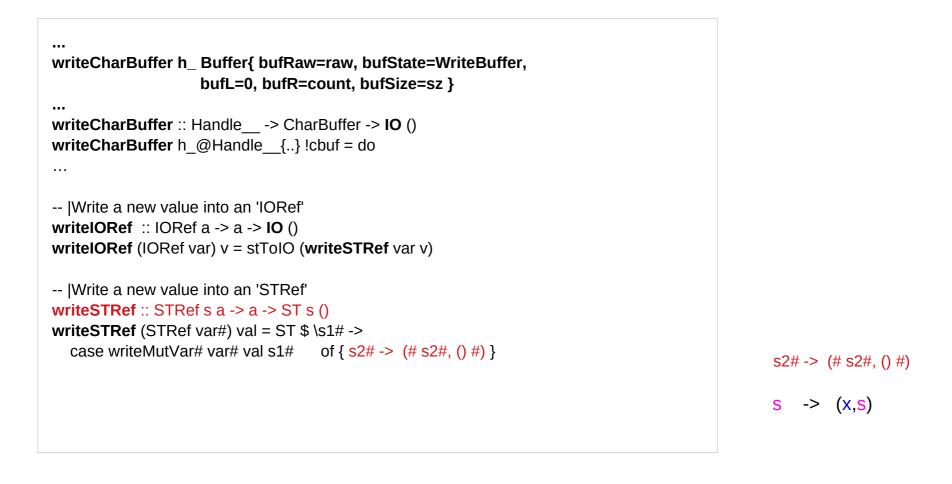
Actions



Simple IO



Primitives in PutStrLn

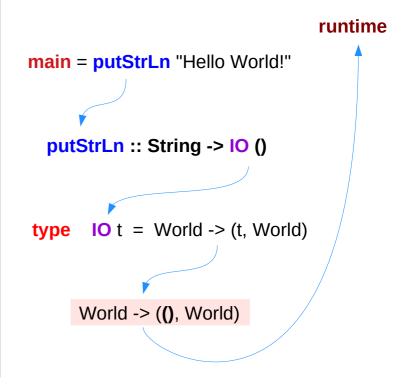


http://hackage.haskell.org/package/base-4.11.1.0/docs/src/GHC.IO.Handle.Text.html#local-6989586621679303176

IO actions in main

IO action is invoked, after the Haskell program has run

- an IO action can <u>never</u> be <u>executed</u> inside the program in order to <u>execute</u> a function of the type World -> (t, World) must <u>supply</u> a value of the type World
- once <u>created</u>, an IO action keeps percolating up until it ends up in main and is <u>executed</u> by the runtime.
- IO action can be also <u>discarded</u>, but that means it will <u>never</u> be <u>evaluated</u>





Laziness

Haskell will <u>not</u> <u>calculate</u> anything unless it's strictly <u>necessary</u> or is <u>forced</u> by the programmer

Haskell will not even evaluate

<u>arguments</u> to a function <u>before</u> <u>calling</u> it

Haskell <u>assumes</u> that the <u>arguments</u> will <u>not</u> be <u>used</u>, so it <u>procrastinates</u> as long as possible. unless proven otherwise

Laziness and Pure Functions

A **pure function** has <u>no</u> **side effects**. Calling a **function once** is the same as calling it **twice** and <u>discarding</u> the **result** of the <u>first</u> **call**.

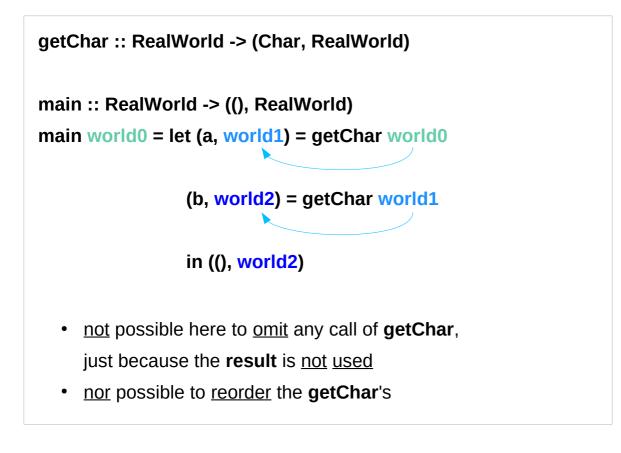
not modifying its **arguments** but modifying only the **result**

furthermore, if the **result** of any function call is <u>not used</u>, Haskell will spare itself the trouble and will <u>never call</u> the **function**.

exception IO ()

-- () non-significant result

Laziness and Pure Functions



world2 requires world1
world1 requires world0

the result () is not used

https://wiki.haskell.org/IO_inside#Welcome_to_the_RealWorld.2C_baby

Laziness Example 1

Division by zero : undefined - never be evaluated.

main = print \$ undefined + 1

no compile time error

but a runtime error

because of an attempt to evaluate undefined.

foo x = **1**

```
main = print $ (foo undefined) + 1
```

Haskell calls **foo** but <u>never</u> <u>evaluates</u> its argument **undefined** (just returns 1)

Laziness Example 2

this does not come from <u>optimization</u> : from the definition of foo , the compiler figures out that its argument is <u>unnecessary</u> .	
but the result is the same if the definition of foo is hidden from view in another module.	
<pre>{-# START_FILE Foo.hs #-} show module Foo (foo) where foo x = 1</pre>	{-# START_FILE Main.hs #-} show import Foo main = print \$ (foo undefined) + 1

Laziness with infinity

laziness allows it to deal with

- infinity (like an infinite list)
- the **future** that hasn't materialized yet

Laziness and IO action

Laziness or not, a program will be executed at some time.

why an expression should be evaluated? among many reasons, the fundamental one is to <u>display</u> its <u>result</u>.

without I/O, nothing would ever be evaluated

Do Notation

Larger IO actions are composed of smaller IO actions.

- the order of composition matters
- sequence IO actions

special syntax for sequencing : the **do** notation.

Do Notation Example

main = do

putStrLn "The answer is: "

print 43

sequencing two IO actions

- one IO action returned by putStrLn
- another IO action returned by print

inside a **do** block proper **indentation**.

Do Notation – input action (1)

whatever you receive from the <u>user</u> or from a <u>file</u> you assign to a <u>variable</u> and use it later.

main = do

str <- getLine putStrLn str

when <u>executed</u>, <u>creates</u> an <u>action</u> that will take the <u>input</u> from the user. then <u>pass</u> this input to the <u>rest</u> of **actions** of the <u>do</u> block under the **name str** when the rest is <u>executed</u>. (not ordinary variable, but a **binding**) immutable variable just a binding

 x <- monadic value
 (only the <u>result</u> of the monadic value execution)

getLine str binded name

Do Notation – input action (2)

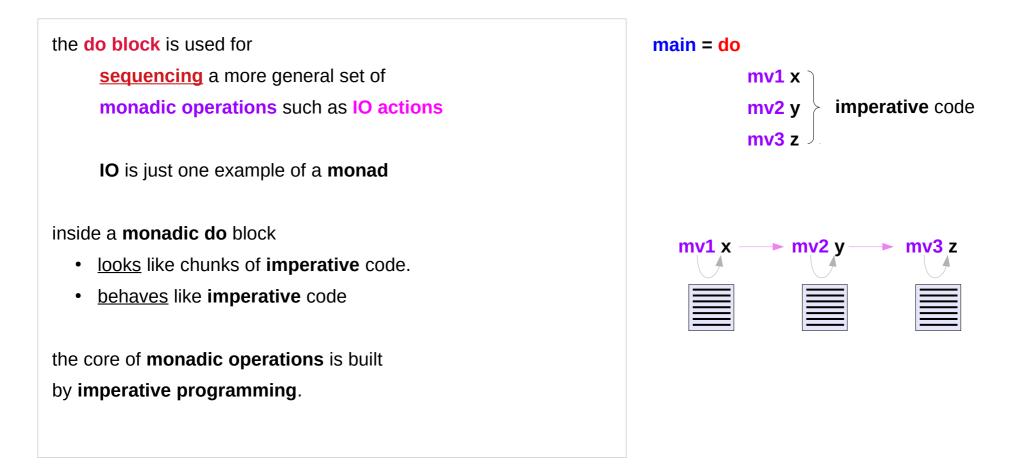


only the returned **result** is passed

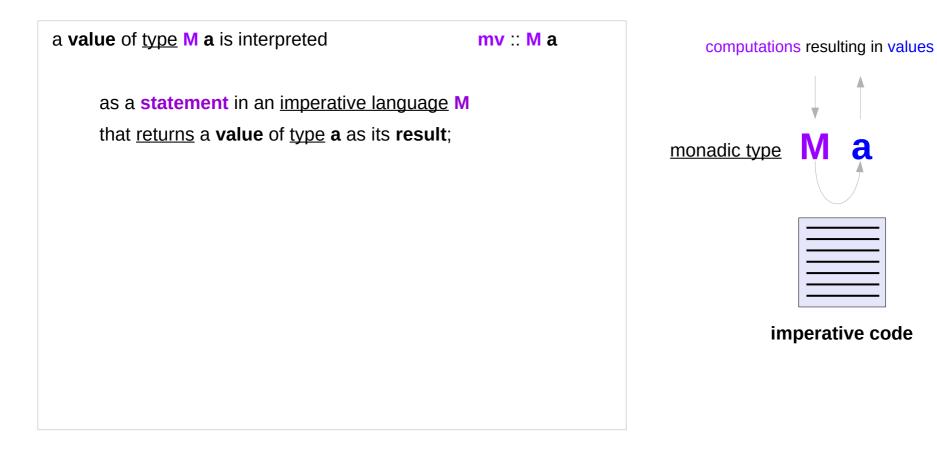
- str is not really a variable
- <- is not really an assignment
- <- creates an action (execution)
- <- <u>binds</u> the <u>name</u> str to the value (String) that will be <u>returned</u> by <u>executing</u> the <u>action</u> of <u>getLine</u>.

In Haskell you <u>never</u> <u>assign</u> to a variable, (immutable) instead you <u>bind</u> a name to a value. getLine creates an action that, when the action executed will take the input from the user. It will then pass that input to the rest of the do block (which is also an action) under the name str when it (the rest) is executed.

do block operations



Monadic value



Semicolon Overloading

The way the **actions** are <u>glued</u> together is the essence of the **Monad**.

Since the <u>glueing</u> happens between the <u>lines</u>, the <u>Monad</u> is sometimes described as an "**overloading of the semicolon**."

Different **monads** overload it differently.

main = do
putStrLn "The answer is: ";
print 43

main =
 putStrLn "The answer is: " >>
 print 43

Semicolon Overloading Examples

can define your own sequencing rule

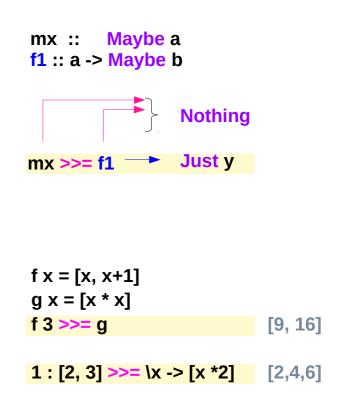
- execute the <u>first</u> statement <u>once</u>, and <u>then</u> execute the <u>next</u> statement
- the <u>first</u> statement computes a <u>value</u>, which the <u>next</u> statement can <u>use</u>

the Maybe monad

 <u>execute</u> the <u>first</u> statement, but <u>only execute</u> the <u>next</u> statement if the value so far <u>isn't null</u>

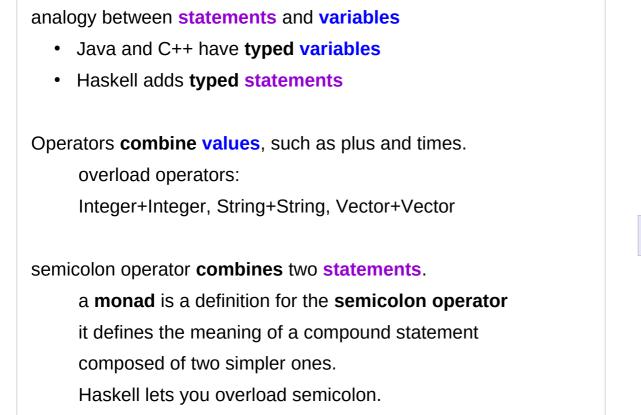
the List monad

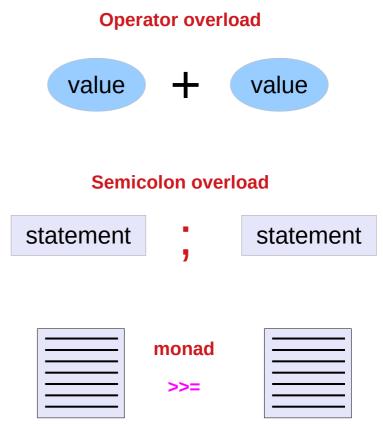
 the <u>first</u> statement computes a <u>list</u> of values, and the <u>second</u> statement <u>runs</u> once using <u>each</u> of them



https://blog.osteele.com/2007/12/overloading-semicolon/

Combining two statements





https://blog.osteele.com/2007/12/overloading-semicolon/

Stateful Computations & IO: Side Effects in Haskell

The functional language Haskell expresses **side effects** such as **I/O** and other **stateful computations**

using monadic actions

IO monad

State monad

https://en.wikipedia.org/wiki/Side_effect_(computer_science)

Stateful Computation

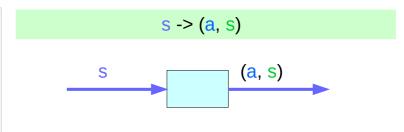
a stateful computation is a function that

takes some **state** and returns a **value** along with some **new state**.

That function would have the following type:

<mark>s -> (a,s</mark>)

- s is the type of the state and
- a the result of the stateful computation.



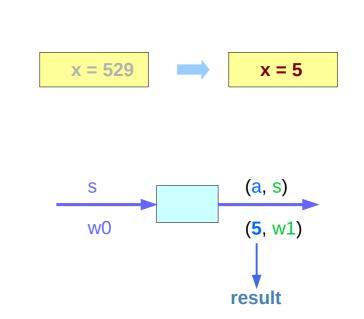
a **function** is an <u>executable</u> <u>data</u> when <u>executed</u>, a <u>result</u> is produced

action (an executable function) **result** is produced if executed

http://learnyouahaskell.com/for-a-few-monads-more

Assignment in the Haskell runtime

Assignment in an imperative language : will assign the value 5 to the variable x will have the value 5 as an expression
Assignment in a functional language as a function that takes a state and returns a result and a new state



http://learnyouahaskell.com/for-a-few-monads-more

Assignment as a stateful computation

Assignment in a <u>functional</u> language	s -> (a, s)
as a function that	s (a, s)
takes a state and	
returns a result and a new state	
	all the variables that have been assignedall the previous mapped variablepreviouslyplus the newly assigned variable
an input state :	
all the variables that have been assigned previously	a result : 5
a result : 5	
a new state :	–
all the previous variable mappings plus	x = 5
the newly assigned variable.	

http://learnyouahaskell.com/for-a-few-monads-more

A value with a **context**

The stateful computation:

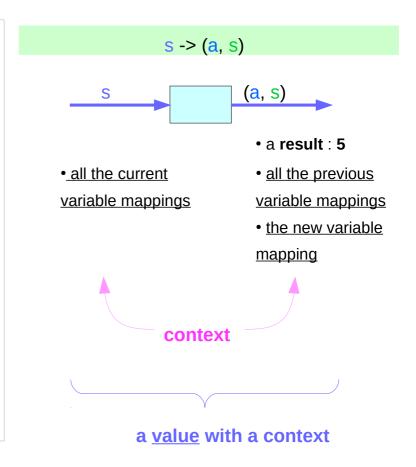
- a function that
 - → takes a state and
 - → returns a **result** and a **new state**
- can be considered as a value with a context

the actual value is the result

the **context** is

an **initial state** that must be provided to get the **result** not only the **result**, but also a **new state** is obtained through the **execution** of the function

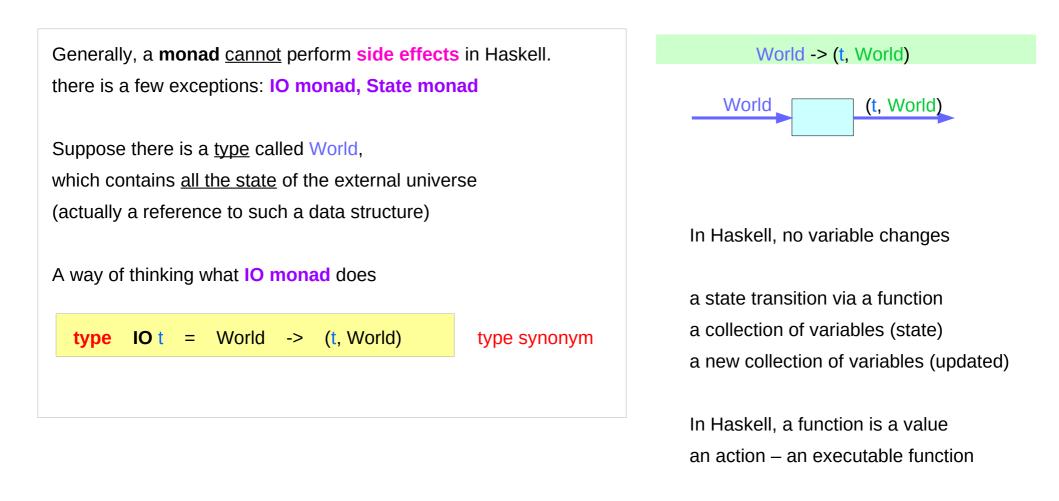
the **result** is determined <u>based</u> on the **initial state** the **result** and the **new state** <u>depend</u> on the **initial state**



http://learnyouahaskell.com/for-a-few-monads-more

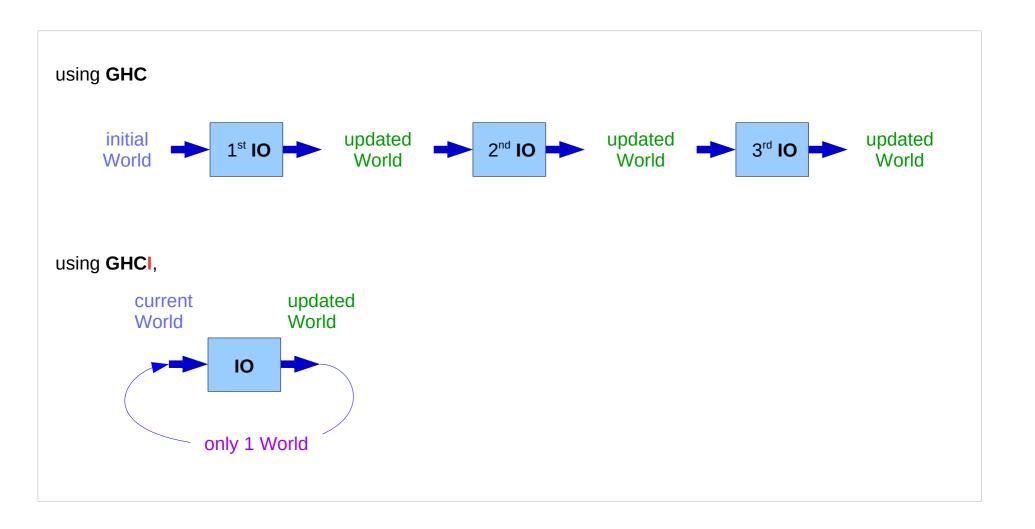
Side Effects (1A)

Stateful computations of IO Monad



https://www.cs.hmc.edu/~adavidso/monads.pdf

Stateful computation models of IO monad



https://www.cs.hmc.edu/~adavidso/monads.pdf





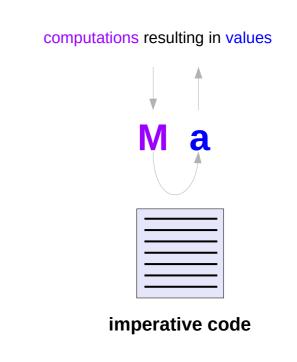
Pure subset of a language

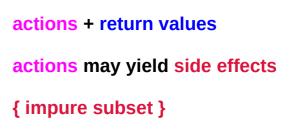
Some functional languages allow **expressions** to <u>yield</u> actions in addition to return values.

These actions are called side effects to emphasize that the return value is the most important result of a function

pure languages prohibit side effects but, pure <u>subsets</u> is still useful

beneficial to write a significant part of a code as **pure** and the remaining error prone **impure** part as small as possible





https://wiki.haskell.org/Functional_programming#Purity

Side Effects (1A)

Pure language features

Immutable Data Referential Transparency Lazy Evaluation Purity and Effects

<u>altered copies</u> are used the <u>same result</u> on each invocation <u>defer</u> until needed <u>mutable</u> array and IO

Pure functional programs typically operate on **immutable data**.

Instead of <u>altering</u> existing values, <u>altered</u> <u>copies</u> are <u>created</u> and the <u>original</u> is <u>preserved</u>.

Since the **unchanged parts** of the structure <u>cannot</u> be <u>modified</u>, they can often be <u>shared</u> between the old and new copies, which <u>saves memory</u>.

Referential Transparency

Pure computations yield the <u>same value</u> each time they are <u>invoked</u>.

This property is called **referential transparency** and makes possible to conduct **equational reasoning** on the code.

no argument modification no global variable access : no side effects

Referential Transparency Examples

y = f x

g = h y y

then we should be able to replace the definition of **g** with

g = h(f x)(f x)

and get the <u>same result;</u> only the **efficiency** might change.

Lazy Evaluation

Since **pure** computations are **referentially transparent** they can be performed <u>at any time</u> and still yield the <u>same result</u>.

This makes it possible to <u>defer</u> the computation of values <u>until</u> they are <u>needed</u>, that is, to **compute** them **lazily**.

Lazy evaluation <u>avoids</u> unnecessary computations and <u>allows</u> infinite data structures to be defined and used.

Purity and Effects

Even though **purely functional programming** is very <u>beneficial</u>, the programmer might want to use **features** that are not available in pure programs, like efficient **mutable arrays** or **convenient I/O**.

There are 2 **approaches** to this problem.

extended impure function
 simulating monads

Using impure functions

Some functional languages **extend** their purely functional core **with side effects**.

The programmer must be <u>careful</u> <u>not</u> to use **impure functions** in places *where only pure functions are expected*.

Using monads

Another way of **introducing side effects** to a pure language is to **simulate** them using **monads**.

While the **language** <u>remains</u> **pure** and **referentially transparent**, **monads** can provide **implicit state** by threading it inside them.

The **compiler** does <u>not</u> care about the **imperative features** because the **language** itself <u>remains</u> **pure**,

however usually the **implementations** do care about them due to the **efficiency reasons**, for instance to provide **O(1) mutable arrays**. stateful computation

Monads enable lazy evaluation

Allowing side effects <u>only through monads</u> and keeping the language **pure** makes it possible to have **lazy evaluation** that does <u>not conflict</u> with the **effects** of **impure code**.

Even though **lazy expressions** can be evaluated **in any order**, the **monad structure** <u>forces</u> the effects to be executed **in the correct order**.

But now, when you look at a do block, it looks very much like imperative code with hidden side effects. The Either monadic code looks like using functions that can throw exceptions. State monad code looks as if the state were a global mutable variable. You access it using get with no arguments, and you modify it by calling put that returns no value. So what have we gained in comparison to C?

We might not see the hidden effects, but the compiler does. It desugars every do block and typechecks it. The state might look like a global variable but it's not. Monadic bind makes sure that the state is threaded from function to function. It's never shared. If you make your Haskell code concurrent, there will be no data races.

https://www.schoolofhaskell.com/school/starting-with-haskell/basics-of-haskell/12-State-Monad

If you have a global environment, which various functions read from (and you might, for example, initialise from a configuration file) then you should thread that as a parameter to your functions (after having, very likely, set it up in your 'main' action). If the explicit parameter passing annoys you, then you can 'hide' it with a Monad.

https://wiki.haskell.org/Global_variables

References

- [1] ftp://ftp.geoinfo.tuwien.ac.at/navratil/HaskellTutorial.pdf
- [2] https://www.umiacs.umd.edu/~hal/docs/daume02yaht.pdf