

State Monad (3E)

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Based on

[Haskell in 5 steps](https://wiki.haskell.org/Haskell_in_5_steps)

https://wiki.haskell.org/Haskell_in_5_steps

Maybe Monad

```
class Monad m where
  return :: a -> m a
  (>>=) :: m a -> (a -> m b) -> m b
```

```
instance Monad Maybe where
  -- return    :: a -> Maybe a
  return x    = Just x

  -- (>>=)     :: Maybe a -> (a -> Maybe b) -> Maybe b
  Nothing >>= _ = Nothing
  (Just x) >>= f = f x
```

```
f :: a -> m b
```

<https://cseweb.ucsd.edu/classes/wi13/cse230-a/lectures/monads2.html>

Maybe Monad

a monad is a **parameterized type** **m**
that supports **return** and **>>=** functions of the specified types

m must be a parameterized type,
rather than just a type (not a concrete type)

It is because of this declaration
that the **do** notation can be used to sequence **Maybe** values.

More generally, Haskell supports the use of this notation
with any monadic type.

examples of types that are monadic,
the benefits that result from recognizing and exploiting this fact.

<https://cseweb.ucsd.edu/classes/wi13/cse230-a/lectures/monads2.html>

Monad, Monoid

monad (plural monads)

- An ultimate atom, or simple, unextended point; something ultimate and **indivisible**.
- (mathematics, computing) A monoid in the category of endofunctors.
- (botany) A single **individual** (such as a pollen grain) that is free from others, not united in a group.

monoid (plural monoids)

- (mathematics) A **set** which is closed under an **associative** binary operation, and which contains an element which is an **identity** for the operation.

<https://en.wiktionary.org/wiki/monad>, [monoid](https://en.wiktionary.org/wiki/monoid)

List Monad

The **Maybe** monad provides a simple model of computations that can fail,

a value of type `Maybe a` is either `Nothing` (**failure**)
the form `Just x` for some `x` of type `a` (**success**)

The **list** monad generalises this notion,
by permitting multiple results in the case of success.

More precisely, a value of `[a]` is
either the empty list `[]` (**failure**)
or the form of a non-empty list `[x1,x2,...,xn]` (**success**)
for some `xi` of type `a`

<https://cseweb.ucsd.edu/classes/wi13/cse230-a/lectures/monads2.html>

List Monad

```
instance Monad [] where
  -- return :: a -> [a]
  return x = [x]

  -- (>>=) :: [a] -> (a -> [b]) -> [b]
  xs >>= f = concat (map f xs)
```

return converts a value into a *successful* result containing that value

>>= provides a means of *sequencing* computations that may produce *multiple results*:

xs >>= **f** applies the function **f** to each of the *results* in the list **xs** to give a *nested list of results*, which is then concatenated to give a *single list of results*.

```
xs :: [a]
f :: a -> [b]
(>>=) :: [a] -> (a -> [b]) -> [b]
```

(Aside: in this context, [] denotes the list type [a] without its parameter.)

<https://cseweb.ucsd.edu/classes/wi13/cse230-a/lectures/monads2.html>

List Monad

instance Monad [] where

-- return :: a -> [a]

return x = [x]

-- (>>=) :: [a] -> (a -> [b]) -> [b]

xs >>= **f** = concat (map **f** **xs**)

instance Monad ST where

-- return :: a -> ST a

return x = \s -> (x,s)

-- (>>=) :: ST a -> (a -> ST b) -> ST b

st >>= **f** = \s -> let (x,s') = **st** s in **f** x s'

<https://cseweb.ucsd.edu/classes/wi13/cse230-a/lectures/monads2.html>

A State Transformer

```
type State = ...
```

```
type ST = State -> State
```

the problem of writing **functions** that manipulate some kind of **state**, represented by a type, whose detail is not our concern now.

a **state transformer (ST)**, which takes the **current state** as its argument, and produces a **modified state** as its result, which reflects any **side effects** performed by the function:

<https://cseweb.ucsd.edu/classes/wi13/cse230-a/lectures/monads2.html>

A Generalized State Transformer

```
type State = ...
```

```
type ST = State -> State
```

```
type ST a = State -> (a, State)
```

want to return a result value in addition to the modified state
generalized state transformers also return a result value,
as a parameter of the **ST** type

```
State -> (a, State)
```

```
  s -> (v, s')
```

s: input state, v: the result value, s': output state

<https://cseweb.ucsd.edu/classes/wi13/cse230-a/lectures/monads2.html>

A Curried Generalized State Transformer

type `ST a = State -> (a, State)` generalized ST

type `ST' a b = b -> State -> (a, State)` further generalized ST

`b -> ST a = b -> State -> (a, State)` think currying

also may need to take argument values
no need to use more generalized ST type
can be exploiting currying.

a state transformer that takes a character and returns an integer
would have type `Char -> ST Int`

`Char -> State -> (Int, State)` **curried form**

<https://cseweb.ucsd.edu/classes/wi13/cse230-a/lectures/monads2.html>

ST Monad

```
instance Monad ST where
```

```
-- return :: a -> ST a
```

```
return x = \s -> (x,s)
```

```
-- (>>=) :: ST a -> (a -> ST b) -> ST b
```

```
st >>= f = \s -> let (x,s') = st s in f x s'
```

>>= provides a means of sequencing **state transformers**:

st >>= f applies the **state transformer st** to an initial state **s**,

then applies the function **f** to the resulting value **x**

to give a second **state transformer (f x)**,

which is then applied to the modified state **s'** to give the final result:

```
st :: ST a
```

```
f :: a -> ST b
```

```
(>>=) :: ST a -> (a -> ST b) -> ST b
```

```
st :: State -> (a, State)
```

```
f :: a -> State -> (b, State)
```

```
(>>=) :: State -> (a, State) -> (a -> ST b)  
-> ST b
```

```
(x,s') = st s
```

```
f x s'
```

<https://cseweb.ucsd.edu/classes/wi13/cse230-a/lectures/monads2.html>

ST Monad

instance Monad **ST** where

-- return :: a -> ST a

return x = \s -> (x,s)

-- (>>=) :: ST a -> (a -> ST b) -> ST b

st >>= f = \s -> let (x,s') = **st s** in **f x s'**

st :: ST a

f :: a -> ST b

(>>=) :: ST a -> (a -> ST b) -> ST b

st :: State -> (a, State) (x,s') = st s s → (x,s')

f :: a -> State -> (b, State) f x s'

(>>=) :: State -> (a, State) -> (a -> State -> (b, State)) -> State -> (b, State)

<https://cseweb.ucsd.edu/classes/wi13/cse230-a/lectures/monads2.html>

let ... in ...

```
cylinder :: (RealFloat a) => a -> a -> a
```

```
cylinder r h =
```

```
  let sideArea = 2 * pi * r * h
```

```
      topArea = pi * r ^2
```

```
  in sideArea + 2 * topArea
```

The form is **let** <bindings> **in** <expression>.

The names that you define in the **let** part are accessible to the expression after the **in** part.

Notice that the names are also aligned in a single column.

For now it just seems that **let** puts the bindings first and the expression that uses them later **whereas** where is the other way around.

<http://learnyouahaskell.com/syntax-in-functions>

List Monad

```
instance Monad [] where
  -- return :: a -> [a]
  return x = [x]

  -- (>>=) :: [a] -> (a -> [b]) -> [b]
  xs >>= f = concat (map f xs)
```

```
instance Monad Maybe where
  -- return :: a -> Maybe a
  return x = Just x

  -- (>>=) :: Maybe a -> (a -> Maybe b) -> Maybe b
  Nothing >>= _ = Nothing
  (Just x) >>= f = f x
```

```
instance Monad ST where
  -- return :: a -> ST a
  return x = \s -> (x,s)

  -- (>>=) :: ST a -> (a -> ST b) -> ST b
  st >>= f = \s -> let (x,s') = st s in f x s'
```

<https://cseweb.ucsd.edu/classes/wi13/cse230-a/lectures/monads2.html>

Dummy Constructor S0

```
type ST a = State -> (a, State)
```

generalized ST

```
data ST0 a = S0 (State -> (a, State))
```

types defined using the **type** mechanism
cannot be made into **instances** of classes.

types defined using the **data** mechanism
can be made into **instances** of classes.
but requires a **dummy constructor** (S0)

<https://cseweb.ucsd.edu/classes/wi13/cse230-a/lectures/monads2.html>

Removing Data Constructor

```
type ST a = State -> (a, State)           generalized ST
```

```
data ST0 a = S0 (State -> (a, State))
```

In order to remove the dummy constructor,
define our own application function (**apply0**) for this type

```
apply0 :: ST0 a -> State -> (a, State)
```

```
apply0 (S0 f) x = f x
```

```
f :: State -> (a, State)
```

* Curried Function

f x y

f :: a -> b -> c

(f x) y

f :: a -> (b -> c)

f x returns a function of type **b -> c**

g y

g :: b -> c

<https://cseweb.ucsd.edu/classes/wi13/cse230-a/lectures/monads2.html>

ST Monad

```
type ST a = State -> (a, State)           generalized ST
```

```
data ST0 a = S0 (State -> (a, State))
```

```
apply0 :: ST0 a -> State -> (a, State)
```

```
apply0 (S0 f) x = f x
```

```
apply0 (S0 g) x = g x
```

```
instance Monad ST0 where
```

```
  -- return :: a -> ST a
```

```
  return x = S0 (\s -> (x,s))
```

```
  -- (>>=) :: ST a -> (a -> ST b) -> ST b
```

```
  st >>= f = S0 (\s -> let (x, s') = apply0 st s in apply0 (f x) s')
```

```
instance Monad ST where
```

```
  -- return :: a -> ST a
```

```
  return x = \s -> (x,s)
```

```
  -- (>>=) :: ST a -> (a -> ST b) -> ST b
```

```
  st >>= f = \s -> let (x,s') = st s in f x s'
```

<https://cseweb.ucsd.edu/classes/wi13/cse230-a/lectures/monads2.html>

Data Constructor

```
data Colour = Red | Green | Blue
```

```
data Colour = RGB Int Int Int
```

```
RGB :: Int -> Int -> Int -> Colour
```

<https://stackoverflow.com/questions/18204308/haskell-type-vs-data-constructor>

Examples (1)

```
pairs :: [a] -> [b] -> [(a,b)]           do  
pairs xs ys = do x <- xs  
                y <- ys  
                return (x, y)
```

this function returns all possible ways of pairing elements from two lists

each possible value x from the list xs, and
each value y from the list ys, and
return the pair (x,y).

<https://cseweb.ucsd.edu/classes/wi13/cse230-a/lectures/monads2.html>

Examples (1)

```
pairs :: [a] -> [b] -> [(a,b)]           do  
pairs xs ys = do x <- xs  
                y <- ys  
                return (x, y)
```

```
pairs xs ys = [(x,y) | x <- xs, y <- ys]   comprehension
```

In fact, there is a formal connection
between the do notation and the comprehension notation.
Both are simply different shorthands
for repeated use of the `>>=` operator for lists.

<https://cseweb.ucsd.edu/classes/wi13/cse230-a/lectures/monads2.html>

Simple Examples (1)

```
(>>) :: Monad m => m a -> m b -> m b;
```

`a1 >> a2` takes the actions `a1` and `a2` and returns the mega action which is `a1-then-a2-returning-the-value-returned-by-a2`.

<https://cseweb.ucsd.edu/classes/wi13/cse230-a/lectures/monads2.html>

Simple Examples (1)

```
type State = Int

fresh :: ST0 Int
fresh = S0 (\n -> (n, n+1))

wtf1 = fresh >>
      fresh >>
      fresh >>
      fresh

ghci> apply0 wtf1 0
```

<https://cseweb.ucsd.edu/classes/wi13/cse230-a/lectures/monads2.html>

Simple Examples (2)

```
return :: a -> ST0 a
```

```
> wtf2 = fresh >>= \n1 ->
```

```
>   fresh >>= \n2 ->
```

```
>   fresh >>
```

```
>   fresh >>
```

```
>   return [n1, n2]
```

```
> wtf2' = do { n1 <- fresh;
```

```
>           n2 <- fresh;
```

```
>           fresh ;
```

```
>           fresh ;
```

```
>           return [n1, n2];
```

```
>           }
```

<https://cseweb.ucsd.edu/classes/wi13/cse230-a/lectures/monads2.html>

Simple Examples (3)

```
ghci> apply0 wtf2 0
```

```
> wtf3 = do n1 <- fresh
```

```
>     fresh
```

```
>     fresh
```

```
>     fresh
```

```
>     return n1
```

<https://cseweb.ucsd.edu/classes/wi13/cse230-a/lectures/monads2.html>

Dice Examples

to generate `Int` dice - result : a number between 1 and 6
throw results from a pseudo-random generator of type `StdGen`.

the type of the **state processors** will be

```
State StdGen Int
```

```
StdGen -> (Int, StdGen)
```

https://en.wikibooks.org/wiki/Haskell/Understanding_monads/State

randomR

the StdGen type : an instance of **RandomGen**

randomR :: (Random a, RandomGen g) => (a, a) -> g -> (a, g)

assume a is Int and g is StdGen

the type of **randomR**

randomR (1, 6) :: StdGen -> (Int, StdGen)

already have a **state processing function**

https://en.wikibooks.org/wiki/Haskell/Understanding_monads/State

randomR

```
randomR (1, 6) :: StdGen -> (Int, StdGen)
```

```
rollDie :: State StdGen Int
```

```
rollDie = state $ randomR (1, 6)
```

https://en.wikibooks.org/wiki/Haskell/Understanding_monads/State

Some Examples (1)

```
module StateGame where
```

```
import Control.Monad.State
```

```
-- Example use of State monad  
-- Passes a string of dictionary {a,b,c}  
-- Game is to produce a number from the string.  
-- By default the game is off, a C toggles the  
-- game on and off. A 'a' gives +1 and a b gives -1.  
-- E.g  
-- 'ab'   = 0  
-- 'ca'   = 1  
-- 'cabca' = 0  
-- State = game is on or off & current score  
--       = (Bool, Int)
```

https://wiki.haskell.org/State_Monad

Some Examples (2)

```
type GameValue = Int
type GameState = (Bool, Int)

playGame :: String -> State GameState GameValue
playGame [] = do
  (_, score) <- get
  return score
```

https://wiki.haskell.org/State_Monad

Some Examples (3)

```
playGame (x:xs) = do
  (on, score) <- get
  case x of
    'a' | on -> put (on, score + 1)
    'b' | on -> put (on, score - 1)
    'c'   -> put (not on, score)
    _     -> put (on, score)
  playGame xs

startState = (False, 0)

main = print $ evalState (playGame "abcaaacbbcabbab") startState
```

https://wiki.haskell.org/State_Monad

References

- [1] <ftp://ftp.geoinfo.tuwien.ac.at/navratil/HaskellTutorial.pdf>
- [2] <https://www.umiacs.umd.edu/~hal/docs/daume02yaht.pdf>