Formal Language (3C)

• Recursively Enumerable Language

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There are three <u>equivalent</u> <u>definitions</u> of a **recursively enumerable** language:

A **recursively enumerable language** is a **recursively enumerable** <u>subset</u> in the set of all possible words over the **alphabet** of the language.

A recursively enumerable language is a formal language for which there exists a **Turing machine** (or other computable function) which will **enumerate** <u>all valid strings</u> of the language.

Note that if the language is infinite, the enumerating algorithm provided can be chosen so that it avoids repetitions, since we can test whether the string produced for number n is "already" produced for a number which is less than n. If it already is produced, use the output for input n+1 instead (recursively), but again, test whether it is "new".

A **recursively enumerable language** is a formal language for which there exists a **Turing machine** (or other computable function) that will halt and accept when presented with any string in the language as input but may either **halt** and **reject** or **loop forever** when presented with a string not in the language.

Contrast this to recursive languages, which require that the Turing machine halts in all cases.

All regular, context-free, context-sensitive and recursive languages are recursively enumerable.

The halting problem is recursively enumerable but not recursive. Indeed one can run the Turing Machine and accept if the machine halts, hence it is recursively enumerable. On the other hand the problem is undecidable.

Some other recursively enumerable languages that are not recursive include:

Post correspondence problem Mortality (computability theory) Entscheidungsproblem

References

