

# First Order Logic – Semantics (3A)

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# Based on

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Contemporary Artificial Intelligence,  
R.E. Neapolitan & X. Jiang

Logic and Its Applications,  
Burkey & Foxley

# Model

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First specify a **signature**

Constant Symbols

Predicate Symbols

Function Symbols

Determines the **language**

Given a language

A **model** is specified

A **domain** of discourse

An **interpretation**

# Model

1. a nonempty set D of **entities** called a **domain of discourse**
  - this domain is a set
  - each element in the set : entity
  - each constant symbol : one entity in the domain
  
2. an **interpretation**
  - (a) an entity in D is assigned to each of the constant symbols.  
Normally, every entity is assigned to a constant symbol.
  - (b) for each **function**,  
an entity is assigned to each possible input of entities to the **function**
  - (c) the predicate '**True**' is always assigned **the value T**  
The predicate '**False**' is always assigned **the value F**
  - (d) for every other **predicate**,  
**the value T** or **F** is assigned  
to each possible input of entities to the **predicate**

# Model

## Constant assignment

(a) an entity → the constant symbols.

## Function assignments

(b) an entity → each possible input of entities to the **function**

## Truth value assignments

(c) the value **T** → the predicate '**True**'  
the value **F** → the predicate '**False**'

(d) for every other **predicate**,  
the value **T** or **F** is assigned → every other predicate  
to each possible input of entities to the **predicate**

# Signature Model Examples

## Signature

1. constant symbols = { Mary, Fred, Sam }
2. predicate symbols = { married, young }
  - married(x, y) : arity two
  - young(x) : arity one

## Model

1. domain of discourse D : the set of three particular *individuals*

- this domain is a set
- each element in the set : entity (= *individuals*)
- each constant symbol : one entity in the domain (= one *individual*)

2. interpretation

(a) a different *individual* is assigned to each of the **constant symbols**

(a) an entity in D is assigned to each of the constant symbols.  
Normally, every entity is assigned to a constant symbol.

# Signature Model Examples

(b) for each **function**,  
an entity is assigned to each possible input of entities to the **function**

(c) the predicate '**True**' is always assigned the value T  
The predicate '**False**' is always assigned the value F

(d) the truth value assignments

$\text{young}(\text{Mary}) = \text{F}$ ,  $\text{young}(\text{Fred}) = \text{F}$ ,  $\text{young}(\text{Sam}) = \text{T}$

$\text{married}(\text{Mary}, \text{Mary}) = \text{F}$ ,  $\text{married}(\text{Mary}, \text{Fred}) = \text{T}$ ,  $\text{married}(\text{Mary}, \text{Sam}) = \text{F}$   
 $\text{married}(\text{Fred}, \text{Mary}) = \text{T}$ ,  $\text{married}(\text{Fred}, \text{Fred}) = \text{F}$ ,  $\text{married}(\text{Fred}, \text{Sam}) = \text{F}$   
 $\text{married}(\text{Sam}, \text{Mary}) = \text{F}$ ,  $\text{married}(\text{Sam}, \text{Fred}) = \text{F}$ ,  $\text{married}(\text{Sam}, \text{Sam}) = \text{F}$

(d) for every other **predicate**,  
the value T or F is assigned  
to each possible input of entities to the **predicate**



# Signature Model Examples

## Signature

1. constant symbols = { Fred, Mary, Sam }
2. predicate symbols = { love }      love(x, y) : arity two
3. function symbols = { mother }      mother(x) : arity one

## Model

1. domain of discourse D : the set of three particular individuals
2. interpretation
  - (a) a different individual is assigned to each of the **constant symbols**
  - (b) **the truth value assignments**  
love(Fred, Fred) = F, love(Fred, Mary) = F, love(Fred, Ann) = F  
love(Mary, Fred) = T, love(Mary, Mary) = F, love(Mary, Ann) = T  
love(Ann, Fred) = T, love(Ann, Mary) = T, love(Ann, Ann) = F
  - (c) **the function assignments**  
mother(Fred) = Mary, mother(Mary) = Ann, mother(Ann) = - (no assignment)

# Signature Model Examples

## Signature

1. constant symbols = { Fred, Mary, Sam }
2. predicate symbols = { love }      love(x, y) : arity two
3. function symbols = { mother }      mother(x) : arity one

## Model

1. domain of discourse D : the set of three particular individuals
2. interpretation
  - (a) a different individual is assigned to each of the **constant symbols**
  - (b) **the truth value assignments**  
love(Fred, Fred) = F, love(Fred, Mary) = F, love(Fred, Ann) = F  
love(Mary, Fred) = T, love(Mary, Mary) = F, love(Mary, Ann) = T  
love(Ann, Fred) = T, love(Ann, Mary) = T, love(Ann, Ann) = F
  - (c) **the function assignments**  
mother(Fred) = Mary, mother(Mary) = Ann, mother(Ann) = - (no assignment)

# The truth value and sentences (1)

The truth values of **all sentences** are assigned :

1. the truth values for **sentences** developed with the symbols  $\neg$ ,  $\wedge$ ,  $\vee$ ,  $\Rightarrow$ ,  $\Leftrightarrow$  are assigned as in propositional logic.
2. the truth values for two terms connected by the  $=$  symbol is T if both terms refer to the same entity; otherwise it is F
3. the truth values for  $\forall x p(x)$  has value T if  $p(x)$  has value T for **every assignment** to  $x$  of an **entity** in the domain D; otherwise it has value F
4. the truth values for  $\exists x p(x)$  has value T if  $p(x)$  has value T for **at least one assignment** to  $x$  of an **entity** in the domain D; otherwise it has value F

# The truth value and sentences (1)

5. the operator **precedence** is as follows  $\neg$ ,  $=$ ,  $\wedge$ ,  $\vee$ ,  $\Rightarrow$ ,  $\Leftrightarrow$
6. the **quantifiers** have precedence over the operators
7. **parentheses** change the order of the precedence

# Formulas and Sentences

An **formula**

- A **atomic formula**
- The operator  $\neg$  followed by a **formula**
- Two formulas separated by  $\wedge$ ,  $\vee$ ,  $\Rightarrow$ ,  $\Leftrightarrow$
- A **quantifier** following by a variable followed by a formula

A **sentence**

- A **formula** with **no free variables**

$\forall x \text{ love}(x,y)$	: free variable $y$	: <b>not</b> a sentence
$\forall x \text{ tall}(x)$	: no free variable	: a sentence

# Finding the truth value

Find the truth values of **all sentences**

1.  $\neg$ ,  $\wedge$ ,  $\vee$ ,  $\Rightarrow$ ,  $\Leftrightarrow$

2. = symbol

3.  $\forall x p(x)$

4.  $\exists x p(x)$

5. the **operator precedence** is as follows  $\neg$ , =,  $\wedge$ ,  $\vee$ ,  $\Rightarrow$ ,  $\Leftrightarrow$

6. the **quantifiers** ( $\forall$ ,  $\exists$ ) have precedence over the **operators**

7. **parentheses** change the order of the precedence

# Formal Language

## Signature

Constant Symbols = {Socrates, Plato, Zeus, Fido}

Predicate Symbols = {human, mortal, legs} all arity one

$\forall \exists$

## Model

D: the set of these four particular individuals

## Interpretation

(a) a different individual is assigned to each of the constant symbols

(b) the truth value assignment

human(Socrates)=T, human(Plato)=T, human(Zeus)=F, human(Fido)=F

mortal(Socrates)=T, mortal(Plato)=T, mortal(Zeus)=F, mortal(Fido)=T

legs(Socrates)=T, legs(Plato)=T, legs(Zeus)=T, legs(Fido)=T

# Formal Language

Sentence 1:  $\text{human}(\text{Zeus}) \wedge \text{human}(\text{Fido}) \vee \text{human}(\text{Socrates}) = \text{T}$   
F             $\wedge$             F             $\vee$             T

Sentence 2:  $\text{human}(\text{Zeus}) \wedge (\text{human}(\text{Fido}) \vee \text{human}(\text{Socrates})) = \text{F}$   
F             $\wedge$ (            F             $\vee$             T            )

Sentence 3:  $\forall x \text{human}(x) = \text{F}$   
                  $\text{human}(\text{Zeus})=\text{F}, \text{human}(\text{Fido})=\text{F}$

Sentence 4:  $\forall x \text{mortal}(x) = \text{F}$   
                  $\text{mortal}(\text{Zeus})=\text{F}$

Sentence 5:  $\forall x \text{legs}(x) = \text{T}$   
                  $\text{legs}(\text{Socrates})=\text{T}, \text{legs}(\text{Plato})=\text{T}, \text{legs}(\text{Zeus})=\text{T}, \text{legs}(\text{Fido})=\text{T}$

Sentence 6:  $\exists x \text{human}(x) = \text{T}$   
                  $\text{human}(\text{Socrates})=\text{T}, \text{human}(\text{Plato})=\text{T}$

Sentence 7:  $\forall x (\text{human}(x) \Rightarrow \text{mortal}(x)) = \text{T}$



# Formal Language

Sentence 7:  $\forall x (\text{human}(x) \Rightarrow \text{mortal}(x)) = T$

$\text{human}(\text{Socrates})=T$ ,  $\text{human}(\text{Plato})=T$ ,  $\text{human}(\text{Zeus})=F$ ,  $\text{human}(\text{Fido})=F$   
 $\text{mortal}(\text{Socrates})=T$ ,  $\text{mortal}(\text{Plato})=T$ ,  $\text{mortal}(\text{Zeus})=F$ ,  $\text{mortal}(\text{Fido})=T$   
 $T \Rightarrow T = T$        $T \Rightarrow T = T$        $F \Rightarrow F = T$        $F \Rightarrow T = T$

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