

Overflow Flag

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1 Based on

2 Overflow flag

- TOC: Overflow flag
- Overflow flag in unsigned and signed computations
- Rules for the overflow flag
- Method 1 for computing the overflow flag
- Method 2 for computing the overflow flag
- More examples of the overflow flag

- 1 "Self-service Linux: Mastering the Art of Problem Determination",

Mark Wilding

- 1 "Computer Architecture: A Programmer's Perspective", Bryant & O'Hallaron

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Compiling 32-bit program on 64-bit gcc

- `gcc -v`
- `gcc -m32 t.c`
- `sudo apt-get install gcc-multilib`
- `sudo apt-get install g++-multilib`
- `gcc-multilib`
- `g++-multilib`
- `gcc -m32`
- `objdump -m i386`

TOC: Overflow flag

- Overflow flag in unsigned and signed computations
- Rules for the overflow flag
- Method 1 for computing the overflow flag
- Method 2 for computing the overflow flag
- More examples of the overflow flag

- Overflow flag

Overflow flag (1)

- **overflow** flag is based on **signed** arithmetic
- to decide if the **overflow** flag is turned on or off, only need to look at the **sign bits** (leftmost) of the three numbers

| | | | | |
|---------|---|-----------|---|------------|
| augend | + | addend | = | sum |
| minuend | - | subrahend | = | difference |

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt

Overflow flag (2)

- in **signed** arithmetic,
 - watch the **overflow** flag to detect errors
 - **overflow** flag on means the result is wrong
 - errors can be detected by examining the sign of the result, in the 2's complement arithmetic ($P + P \rightarrow N$ or $N + N \rightarrow P$)
- in **unsigned** arithmetic,
 - the **overflow** flag tells you nothing interesting

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Overflow flag (3)

- when two positive numbers are added
 - if the result is a positive, ($P + P \rightarrow P$), then no overflow
 - if the result is a negative, ($P + P \rightarrow N$), then overflow
- when two negative numbers are added
 - if the result is a negative, ($N + N \rightarrow N$), then no overflow
 - if the result is a positive, ($N + N \rightarrow P$), then overflow

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Overflow flag (4)

- adding negative (**N**) and positive (**P**) numbers cannot be wrong, because the sum is between the addends ($[\mathbf{N}, \mathbf{P}]$) .
 - if opposite signed numbers are added, then no overflow
 - both of the addends lies in the allowable range of numbers, their sum is between the opposite signed addends, therefore the sum lies also in the allowable range
 - $(P + N \rightarrow P \text{ or } N)$ no overflow always
 - $(N + P \rightarrow P \text{ or } N)$ no overflow always

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TOC Rules for the overflow flag

- the 1st rule for setting OF
- the 2nd rule for setting OF
- cases for clearing OF (1 ~ 6)

Overflow flag setting and clearing conditions

| | | Method 1 | | Method 2 |
|---|------|-----------------------|-----------------------|--------------------------|
| | | ADD conditions | SUB conditions | |
| 1 | OF=1 | $P + P \rightarrow N$ | $P - N \rightarrow N$ | $c_n \oplus c_{n-1} = 1$ |
| 2 | OF=1 | $N + N \rightarrow P$ | $N - P \rightarrow P$ | $c_n \oplus c_{n-1} = 1$ |
| 3 | OF=0 | $P + P \rightarrow P$ | $P - N \rightarrow P$ | $c_n \oplus c_{n-1} = 0$ |
| 4 | OF=0 | $N + N \rightarrow N$ | $N - P \rightarrow N$ | $c_n \oplus c_{n-1} = 0$ |
| 5 | OF=0 | $P + N \rightarrow P$ | $P - N \rightarrow P$ | $c_n \oplus c_{n-1} = 0$ |
| 6 | OF=0 | $P + N \rightarrow N$ | $P - P \rightarrow N$ | $c_n \oplus c_{n-1} = 0$ |
| 7 | OF=0 | $N + P \rightarrow P$ | $N - N \rightarrow P$ | $c_n \oplus c_{n-1} = 0$ |
| 8 | OF=0 | $N + P \rightarrow N$ | $N - P \rightarrow N$ | $c_n \oplus c_{n-1} = 0$ |

$$+P = -(-P) = -N$$

$$+N = -(-N) = -P$$

The 1st case - setting the overflow flag

① $P + P \rightarrow N$ ($P - N \rightarrow N$)

If the **sum** of two **signed** numbers with the sign bits off (0, 0) yields a result number with the sign bit on (1),

the **overflow flag** is turned on

signed addition

0100 carries

0100 (+4)

+0100 (+4)

01000 (-8)

signed subtraction

0100 (+4)

-1100 -(-4)

01000 (-8)

unsigned addition

0100 (4)

+0100 +(4)

01000 (8)

Method 1 $OF = 1$

when $\overline{a_{n-1}} \cdot \overline{b_{n-1}} \cdot s_{n-1}$ for addition
and $\overline{a_{n-1}} \cdot b_{n-1} \cdot s_{n-1}$ for subtraction

Method 2 $OF = c_n \oplus c_{n-1}$

$c_4 \oplus c_3 = 0 \oplus 1 = 1$

The 2nd case - setting the overflow flag

2 $N + N \rightarrow P$ ($N - P \rightarrow P$)

If the **sum** of two numbers with the sign bits on (1, 1) yields a result number with the sign bit off (0)

the **overflow flag** is turned on.

signed addition

```
1001 carries
 1001 (-7)
+1001 +(-7)
-----
10010 ( 2)
```

signed subtraction

```
1001 (-7)
-0111 -(+7)
-----
10010 ( 2)
```

unsigned addition

```
1001 ( 9)
+1001 +( 9)
-----
10010 (18)
```

Method 1 $OF = 1$

when $a_{n-1} \cdot b_{n-1} \cdot \overline{s_{n-1}}$ for addition
and $a_{n-1} \cdot \overline{b_{n-1}} \cdot \overline{s_{n-1}}$ for subtraction

Method 2 $OF = c_n \oplus c_{n-1}$

$c_4 \oplus c_3 = 1 \oplus 0 = 1$

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The 3rd case - clearing the overflow flag

③ $P + P \rightarrow P$ ($P - N \rightarrow P$)

If the **sum** of two **signed** numbers with the sign bits off (0, 0) yields a result number with the sign bit off (0), the **overflow flag** is turned off

signed addition

```
0011  carries
0011  (+3)
+0011  (+3)
-----
00110  (+6)
```

signed subtraction

```
0011  (+3)
-1101  -(-3)
-----
00110  (+6)
```

unsigned addition

```
0011  ( 3)
+0011  +( 3)
-----
00110  ( 6)
```

Method 1 $OF = 0$

when $\overline{a_{n-1}} \cdot \overline{b_{n-1}} \cdot \overline{s_{n-1}}$ for addition
and $\overline{a_{n-1}} \cdot b_{n-1} \cdot \overline{s_{n-1}}$ for subtraction

Method 2 $OF = c_n \oplus c_{n-1}$

$c_4 \oplus c_3 = 0 \oplus 0 = 0$

The 4th case - clearing the overflow flag

④ $N + N \rightarrow N$ ($N - P \rightarrow N$)

If the **sum** of two **signed** numbers with the sign bits on (1, 1) yields a result number with the sign bit on (1), the **overflow flag** is turned off

signed addition

```
1101 carries
 1101 (-3)
+1101 +(-3)
-----
11010 (-6)
```

signed subtraction

```
1101 (-3)
-0011 -(+3)
-----
11010 (-6)
```

unsigned addition

```
1101 (13)
+1101 +(13)
-----
11010 (26)
```

Method 1 $OF = 0$

when $a_{n-1} \cdot b_{n-1} \cdot s_{n-1}$ for addition
and $a_{n-1} \cdot \overline{b_{n-1}} \cdot s_{n-1}$ for subtraction

Method 2 $OF = c_n \oplus c_{n-1}$

$c_4 \oplus c_3 = 1 \oplus 1 = 0$

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The 5th case - clearing the overflow flag

5 $P + N \rightarrow P, (P - P \rightarrow P)$

If the **sum** of two **signed** numbers with the sign bits off and on (0, 1) yields a result number with the sign bit off (0), the **overflow flag** is turned off

signed addition

```
1100 carries
 0100 (+4)
+1101 +(-3)
-----
10001 (+1)
```

signed subtraction

```
0100 (+4)
-0011 -(+3)
-----
10001 (+1)
```

unsigned addition

```
0100 ( 4)
+1101 +(13)
-----
10001 (17)
```

Method 1 $OF = 0$

when $\overline{a_{n-1}} \cdot b_{n-1} \cdot \overline{s_{n-1}}$ for addition
and $\overline{a_{n-1}} \cdot \overline{b_{n-1}} \cdot \overline{s_{n-1}}$ for subtraction

Method 2 $OF = c_n \oplus c_{n-1}$

$c_4 \oplus c_3 = 1 \oplus 1 = 0$

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The 6th case - clearing the overflow flag

6 $P + N \rightarrow N$ ($P - P \rightarrow N$)

If the **sum** of two **signed** numbers with the sign bits off and on (0, 1) yields a result number with the sign bit on (1), the **overflow flag** is turned off

signed addition

```
0000 carries
 0011 (+3)
+1100 +(-4)
-----
01111 (-1)
```

signed subtraction

```
0011 (+3)
-0100 -(+4)
-----
01111 (-1)
```

unsigned addition

```
0011 ( 3)
+1100 +(12)
-----
01111 (15)
```

Method 1 $OF = 0$

when $\overline{a_{n-1}} \cdot b_{n-1} \cdot s_{n-1}$ for addition
and $\overline{a_{n-1}} \cdot b_{n-1} \cdot s_{n-1}$ for subtraction

Method 2 $OF = c_n \oplus c_{n-1}$

$c_4 \oplus c_3 = 0 \oplus 0 = 0$

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The 7th case - clearing the overflow flag

$$\textcircled{7} \quad N + P \rightarrow P, (N - N \rightarrow P)$$

If the **sum** of two **signed** numbers with the sign bits on and off (1, 0) yields a result number with the sign bit off (0), the **overflow flag** is turned off

signed addition

```
1100 carries
 1101 (-3)
+0100 (+4)
-----
10001 (+1)
```

signed subtraction

```
0011 (-3)
-1100 -(-4)
-----
10001 (+1)
```

unsigned addition

```
1101 (13)
+0100 +( 4)
-----
10001 (17)
```

Method 1 $OF = 0$

when $a_{n-1} \cdot \overline{b_{n-1}} \cdot \overline{s_{n-1}}$ for addition
and $a_{n-1} \cdot b_{n-1} \cdot \overline{s_{n-1}}$ for subtraction

Method 2 $OF = c_n \oplus c_{n-1}$

$c_4 \oplus c_3 = 1 \oplus 1 = 0$

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The 8th case - clearing the overflow flag

8 $N + P \rightarrow N, (N + P \rightarrow N)$

If the **sum** of two **signed** numbers with the sign bits on and off (1, 0) yields a result number with the sign bit on (1),

the **overflow flag** is turned off

signed addition

```
0000 carries
 1100 (-4)
+0011 +(3)
-----
01111 (-1)
```

signed subtraction

```
0100 (-4)
-1101 -(-3)
-----
01111 (-1)
```

unsigned addition

```
1100 (12)
+0011 +( 3)
-----
01111 (15)
```

Method 1 $OF = 0$ when $a_{n-1} \cdot \overline{b_{n-1}} \cdot s_{n-1}$
and $a_{n-1} \cdot b_{n-1} \cdot s_{n-1}$

Method 2 $OF = c_n \oplus c_{n-1}$ $c_4 \oplus c_3 = 0 \oplus 0 = 0$

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TOC Method 1 for computing the overflow flag

- Adding two numbers with the same sign
- Overflow conditions for additions and subtractions
- Overflow condition for an addition
- Overflow conditions for a subtraction
- Overflow in signed computations

Adding two numbers with the same sign

- **overflow** can only happen when adding two numbers of the same sign results in a different sign ($P + P \rightarrow N$, $N + N \rightarrow P$)

- n -bit **signed** binary arithmetic $A + B = S$

$$A = (a_{n-1}, \dots, a_1, a_0)$$

$$B = (b_{n-1}, \dots, b_1, b_0)$$

$$S = (s_{n-1}, \dots, s_1, s_0)$$

- to detect overflow
 - only the **sign** bits are considered
 - **msb** (most significant bit) $a_{n-1}, b_{n-1}, s_{n-1}$
 - the other bits are ignored

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Overflow conditions for additions and subtractions

- with two operands (A and B) and one result (S), three sign bits ($a_{n-1}, b_{n-1}, s_{n-1}$) are considered
→ $2^3 = 8$ possible combinations
- only two cases result in **overflow** for an addition
 - 0 0 1 ($p + p \rightarrow n$)
 - 1 1 0 ($n + n \rightarrow p$)
- only two cases are considered as **overflow** for an subtraction
 - 0 1 1 ($p - n \rightarrow n$)
 - 1 0 0 ($n - p \rightarrow p$)

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Overflow condition for an addition

- Overflow in an addition ($A + B$)

| | a_{n-1} | b_{n-1} | c_{n-1} | |
|------|-----------|-----------|-----------|-----------------------|
| | 0 | 0 | 0 | $p + p \rightarrow p$ |
| OVER | 0 | 0 | 1 | $p + p \rightarrow n$ |
| | 0 | 1 | 0 | $p + n \rightarrow p$ |
| | 0 | 1 | 1 | $p + n \rightarrow n$ |
| | 1 | 0 | 0 | $n + p \rightarrow p$ |
| | 1 | 0 | 1 | $n + p \rightarrow n$ |
| OVER | 1 | 1 | 0 | $n + n \rightarrow p$ |
| | 1 | 1 | 1 | $n + n \rightarrow n$ |

- adding two positives should be positive
- adding two negatives should be negative

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Overflow conditions for a subtraction

- Overflow in a subtraction ($A - B$)

| | a_{n-1} | b_{n-1} | c_{n-1} | |
|------|-----------|-----------|-----------|-----------------------|
| | 0 | 0 | 0 | $p - p \rightarrow p$ |
| | 0 | 0 | 1 | $p - p \rightarrow n$ |
| | 0 | 1 | 0 | $p - n \rightarrow p$ |
| OVER | 0 | 1 | 1 | $p - n \rightarrow n$ |
| OVER | 1 | 0 | 0 | $n - p \rightarrow p$ |
| | 1 | 0 | 1 | $n - p \rightarrow n$ |
| | 1 | 1 | 0 | $n - n \rightarrow p$ |
| | 1 | 1 | 1 | $n - n \rightarrow n$ |

- subtracting a negative is the same as adding a positive
- subtracting a positive is the same as adding a negative

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Overflow in signed computations

- ALU might contain a small logic that sets the **overflow** flag to "1" if and only if any one of the above four **OV conditions** is met.
- in **signed** computations, adding two numbers of the same sign must produce a result of the same sign, otherwise overflow happened.

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TOC Method 2 for computing the overflow flag

- Carry into and carry out of the sign bit
- Overflow in 2's complement arithmetic
- Overflow flag = $c_n \oplus c_{n-1}$
- Examples of 4-bit signed additions
- c_n and c_{n-1} in a n -bit addition
- Overflow flag computation
- Examples of computing overflow flag
- Hexadecimal carry, octal carry, decimal carry
- No carry into the sign bit

Carry into and carry out of the sign bit

- When adding two n -bit binary values, consider
 - the **carry** *coming into* the most significant bit (msb)
 c_{n-1} : **carry** into the **sign** bit
 - the **carry** *going out of* the most significant bit (msb)
 c_n : **carry** out of the **sign** bit
this is the **carry** flag (**CF**) in the processor

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Overflow in 2's complement arithmetic

- **overflow** in 2's complement happens (**OF=1**) when
 - there is a **carry** *into* the **sign** bit ($c_{n-1} = 1$)
but no carry *out of* the **sign** bit ($c_n = 0$)
 - there is no carry *into* the **sign** bit ($c_{n-1} = 0$)
but a **carry** *out of* the **sign** bit ($c_n = 1$)

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Overflow flag = $c_n \oplus c_{n-1}$

- the **overflow** flag is the **XOR** ($c_n \oplus c_{n-1}$) of
 - of the **carry coming into** the **sign** bit (c_{n-1})
 - with the **carry going out of** the **sign** bit (c_n)
- **overflow** happens when the **carry in** (c_{n-1}) does not equal to the **carry out** (c_n)

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Examples of 4-bit signed additions (1)

- 4-bit 2's complement addition examples

```
0000
 0100 (+4) (pos sign 0)
+1000 (-8) (neg sign 1)
=====
01100 (-4) (neg sign 1)
```

```
C4 carry out 0 (1+0+0)
C3 carry in 0 (0+1+0)
0 XOR 0 = NO OVERFLOW
```

```
1100
 1100 (-4) (neg sign 1)
+0100 (+4) (pos sign 0)
=====
10000 ( 0) (pos sign 0)
```

```
C4 carry out 1 (1+0+1)
C3 carry in 1 (1+1+0)
1 XOR 1 = NO OVERFLOW
```

```
0100
 0100 (+4) (pos sign 0)
+0100 (+4) (pos sign 0)
=====
01000 (-8) (neg sign 1)
```

```
C4 carry out 0 (0+0+1)
C3 carry in 1 (1+1+0)
0 XOR 1 = OVERFLOW!
```

```
1000
 1100 (-4) (neg sign 1)
+1000 (-8) (neg sign 1)
=====
10100 (+4) (pos sign 0)
```

```
C4 carry out 1 (1+1+0)
C3 carry in 0 (1+0+0)
1 XOR 0 = OVERFLOW!
```

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Examples of 4-bit signed additions (2)

- same sign addition → possible overflow

| | | | |
|--------------------------|--------------------------|--------------------------|--------------------------|
| ----- + +, - ----- | ----- - -, + ----- | ----- + +, + ----- | ----- - -, - ----- |
| +5 +5 | -5 -5 | +5 +1 | -5 -1 |
| ----- -6(OF) | ----- +6(OF) | ----- +6 | ----- -6 |

| | | | |
|----------------------|----------------------|----------------------|----------------------|
| 0101 0101 0101 | 1011 1011 1011 | 0001 0101 0001 | 1111 1011 1111 |
| ----- 01010 | ----- 10110 | ----- 00110 | ----- 11010 |
| ----- C4 = 0 | ----- C4 = 1 | ----- C4 = 0 | ----- C4 = 1 |
| ----- C3 = 1 | ----- C3 = 0 | ----- C3 = 0 | ----- C3 = 1 |
| ----- OF = 1 | ----- OF = 1 | ----- OF = 0 | ----- OF = 0 |

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Examples of 4-bit signed additions (3)

- mixed sign addition → no overflow

| | | | |
|--|--|--|--|
| ----- + -, + ----- +5 -1 ----- +4 | ----- + -, - ----- +5 -6 ----- -1 | ----- - +, + ----- -5 +6 ----- +1 | ----- - +, - ----- -5 +1 ----- -4 |
| 1111 0101 1111 ----- 10100 ----- C4 = 1 C3 = 1 ----- OF = 0 | 0000 0101 1010 ----- 01111 ----- C4 = 0 C3 = 0 ----- OF = 0 | 1110 1011 0110 ----- 10001 ----- C4 = 1 C3 = 1 ----- OF = 0 | 0011 1011 0001 ----- 01100 ----- C4 = 0 C3 = 0 ----- OF = 0 |

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c_n and c_{n-1} in a n -bit addition

$(n-1)^{th}$ bit – MSB

- adding operations at the $(n-1)$ bit position
- $\{c_n, s_{n-1}\} = a_{n-1} + b_{n-1} + c_{n-1}$

$$\begin{array}{r} \text{msb} \\ a_{n-1} \\ b_{n-1} \\ \hline c_{n-1} \\ \hline c_n \quad s_{n-1} \end{array}$$

- c_n : carry coming out of the msb

$(n-2)^{th}$ bit

- adding operations at the $(n-2)$ bit position
- $\{c_{n-1}, s_{n-2}\} = a_{n-2} + b_{n-2} + c_{n-2}$

$$\begin{array}{r} \text{msb} \\ a_{n-2} \\ b_{n-2} \\ \hline c_{n-2} \\ \hline c_{n-1} \quad s_{n-2} \end{array}$$

- c_{n-1} : carry coming into the msb

Overflow flag computation

ADD (addition)

$$OF = c_n \oplus c_{n-1}$$

a 2's complement addition

$$A + B = A + B + \mathbf{0} \quad (c_0 = 0)$$

$$\begin{aligned} &\{c_n, s_{n-1}\} \\ &= a_{n-1} + b_{n-1} + c_{n-1} \end{aligned}$$

$$\begin{aligned} &\{c_{n-1}, s_{n-2}\} \\ &= a_{n-2} + b_{n-2} + c_{n-2} \end{aligned}$$

SUB (subtraction)

$$OF = c_n \oplus c_{n-1}$$

the transformed addition

$$A - B = A + \overline{B} + \mathbf{1} \quad (c_0 = 1)$$

$$\begin{aligned} &\{c_n, s_{n-1}\} \\ &= a_{n-1} + \overline{b_{n-1}} + c_{n-1} \end{aligned}$$

$$\begin{aligned} &\{c_{n-1}, s_{n-2}\} \\ &= a_{n-2} + \overline{b_{n-2}} + c_{n-2} \end{aligned}$$

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Hexadecimal carry, octal carry, decimal carry

- Note that this XOR method only works with the **binary** carry that goes into the sign **bit**.
- not works with **hexadecimal carry**
decimal carry, **octal carry**
 - the carry doesn't go into the sign **bit**
 - can't XOR that non-binary carry with the outgoing carry.

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No carry into the sign bit

- Hexadecimal addition example
(showing that XOR doesn't work for hex carry):

```
8Ah
+8Ah
====
114h
```

- The hexadecimal carry of 1 resulting from A+A does not affect the sign bit.
- If you do the math in binary, you'll see that there is **no** carry **into** the sign bit; but, there is carry out of the sign bit. Therefore, the above example sets OVERFLOW on. (The example adds two negative numbers and gets a positive number.)

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Summary I

| unsigned add/sub | | | signed addition | | | signed subtraction | | | CF | OF |
|------------------|------------|-----|-----------------|-------|-----|--------------------|-------|-----|----|----|
| 1101 | (13) | | 1101 | (-3) | | 1101 | (-3) | | | |
| +1110 | +(14) | ADD | +1110 | +(-2) | ADD | -0010 | -(-2) | | | |
| ----- | ----- | | ----- | ----- | | ----- | ----- | | | |
| 11011 | (11) (+16) | | 11011 | (-5) | | 11011 | (-5) | | 1 | 0 |
| | | | | | | | | | | |
| 0011 | (3) | | 0011 | (+3) | | 0011 | (+3) | | | |
| -1110 | -(14) | SUB | +0010 | +(+2) | | -1110 | -(-2) | SUB | | |
| ----- | ----- | | ----- | ----- | | ----- | ----- | | | |
| 10101 | (5) (-16) | | 00101 | (+5) | | 00101 | (+5) | | 1 | 0 |
| | | | | | | | | | | |
| 0011 | (3) | | 0011 | (+3) | | 0011 | (+3) | | | |
| +0010 | +(2) | ADD | +0010 | +(+2) | ADD | -1110 | -(-2) | | | |
| ----- | ----- | | ----- | ----- | | ----- | ----- | | | |
| 00101 | (5) (+ 0) | | 00101 | (+5) | | 00101 | (+5) | | 0 | 0 |
| | | | | | | | | | | |
| 1101 | (13) | | 1101 | (-3) | | 1101 | (-3) | | | |
| -0010 | -(2) | SUB | +1110 | +(-2) | | -0010 | -(-2) | SUB | | |
| ----- | ----- | | ----- | ----- | | ----- | ----- | | | |
| 11011 | (11) (-16) | | 11011 | (-5) | | 11011 | (-5) | | 0 | 0 |

Summary II

| unsigned add/sub | | | signed addition | | | signed subtraction | | | CF | OF |
|------------------|------------|-----|-----------------|-------|-----|--------------------|-------|-----|----|----|
| 1011 | (11) | | 1011 | (-5) | | 1011 | (-5) | | | |
| +1100 | +(12) | ADD | +1100 | +(-4) | ADD | -0100 | -(+4) | | | |
| ----- | ----- | | ----- | ----- | | ----- | ----- | | | |
| 10111 | (7) (+16) | | 10111 | (+7) | | 10111 | (+7) | | 1 | 1 |
| | | | | | | | | | | |
| 0101 | (5) | | 0101 | (+5) | | 0101 | (+5) | | | |
| -1100 | -(12) | SUB | +0100 | +(+4) | | -1100 | -(-4) | SUB | | |
| ----- | ----- | | ----- | ----- | | ----- | ----- | | | |
| 11001 | (9) (-16) | | 01001 | (-7) | | 01001 | (-7) | | 1 | 1 |
| | | | | | | | | | | |
| 0101 | (5) | | 0101 | (+5) | | 0101 | (+5) | | | |
| +0100 | +(4) | ADD | +0100 | +(+4) | ADD | -1100 | -(-4) | | | |
| ----- | ----- | | ----- | ----- | | ----- | ----- | | | |
| 01001 | (9) (+ 0) | | 01001 | (-7) | | 01001 | (-7) | | 0 | 1 |
| | | | | | | | | | | |
| 1011 | (11) | | 1011 | (-5) | | 1011 | (-5) | | | |
| -0100 | -(4) | SUB | +1100 | +(-4) | | -0100 | -(+4) | SUB | | |
| ----- | ----- | | ----- | ----- | | ----- | ----- | | | |
| 00111 | (7) (0) | | 10111 | (+7) | | 10111 | (+7) | | 0 | 1 |

Cases for setting the overflow flag (1) CF=1, OF=1

- signed integer overflow (OF=1 means incorrect S)

| * unsigned addition | | * signed addition | signed subtraction |
|---------------------|--|---------------------------------|--------------------|
| 1011 (11) | | 1000 | |
| +1100 +(12) ADD | | 1011 (-5) | 1011 (-5) |
| ----- | | +1100 +(-4) ADD | -0100 -(+4) |
| 10111 (7) (+16) | | ----- | ----- |
| | | 10111 (+7) | 10111 (+7) |
| OF=1 | | n + n -> p (OF=1) | n - p -> p (OF=1) |
| OF meaningless | | -> incorrect S | -> incorrect S |
| S = 0111 | | S = 0111 | S = 0111 |
| * think hand | | * OF <- C4 XOR C3 = 1 XOR 0 = 1 | |
| addition | | of signed addition | |

* CF=1, S=0111, OF=1 for all three interpretations

Cases for setting the overflow flag (2) CF=1, OF=1

- signed integer overflow (OF=1 means incorrect S)

| * unsigned subtraction | | signed addition | | * signed subtraction |
|--------------------------|--|---------------------------------|--|----------------------|
| 0101 (5) | | 0100 | | 0101 (+5) |
| -1100 -(12) SUB | | +0100 +(4) | | -1100 -(-4) SUB |
| ----- | | ----- | | ----- |
| 11001 (9) (-16) | | 01001 (-7) | | 01001 (-7) |
| OF=1 | | p + p -> n (OF=1) | | p - n -> n (OF=1) |
| OF meaningless | | -> incorrect S | | -> incorrect S |
| S = 1001 | | S = 1001 | | S = 1001 |
| ----- | | | | |
| * think hand subtraction | | * OF <- C4 XOR C3 = 0 XOR 1 = 1 | | of signed addition |
| ----- | | | | |

* CF=1, S=1001, OF=1 for all three interpretations

Cases for setting the overflow flag (3) CF=0, OF=1

- signed integer overflow (OF=1 means incorrect S)

| * unsigned addition | | * signed addition | signed subtraction |
|---------------------|--|---------------------------------|--------------------|
| 0101 (5) | | 0100 | |
| +0100 +(4) ADD | | 0101 (+5) | 0101 (+5) |
| ----- | | +0100 +(4) ADD | -1100 -(-4) |
| 01001 (9) (+ 0) | | ----- | ----- |
| | | 01001 (-7) | 01001 (-7) |
| OF=1 | | p + p -> n (OF=1) | p - n -> n (OF=1) |
| OF meaningless | | -> incorrect S | -> incorrect S |
| S = 1001 | | S = 1001 | S = 1001 |
| * think hand | | * OF <- C4 XOR C3 = 0 XOR 1 = 1 | |
| addition | | of signed addition | |

* CF=0, S=1001, OF=1 for all three interpretations

Cases for setting the overflow flag (4) CF=0, OF=1

- signed integer overflow (OF=1 means incorrect S)

| * unsigned subtraction | | signed addition | | * signed subtraction |
|--------------------------|--|---------------------------------|--|----------------------|
| 1011 (11) | | 1000 | | 1011 (-5) |
| -0100 -(4) SUB | | +1100 +(-4) | | -0100 -(+4) SUB |
| ----- | | ----- | | ----- |
| 00111 (7) (0) | | 10111 (+7) | | 10111 (+7) |
| OF=1 | | n + n -> p (OF=1) | | n - p -> p (OF=1) |
| OF meaningless | | -> incorrect S | | -> incorrect S |
| S = 0111 | | S = 0111 | | S = 0111 |
| * think hand subtraction | | * OF <- C4 XOR C3 = 1 XOR 0 = 1 | | of signed addition |

* CF=0, S=0111, OF=1 for all three interpretations

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt

Cases for clearing the overflow flag (1) CF=1, OF=0

- no signed integer overflow (CF=0 means correct S)

| * unsigned addition | | * signed addition | | signed subtraction |
|---------------------|--|---------------------------------|--|--------------------|
| 1101 (13) | | 1100 | | |
| +1110 +(14) ADD | | 1101 (-3) | | 1101 (-3) |
| ----- | | +1110 +(-2) ADD | | -0010 -(+2) |
| ----- | | ----- | | ----- |
| 11011 (11) (+16) | | 11011 (-5) | | 11011 (-5) |
| | | | | |
| OF=0 | | n + n -> n (OF=0) | | n - p -> n (OF=0) |
| | | | | |
| OF meaningless | | -> correct S | | -> correct S |
| S = 0000 | | S = 0000 | | S = 0000 |
| | | | | |
| * think hand | | * OF <- C4 XOR C3 = 1 XOR 1 = 0 | | |
| addition | | of signed addition | | |

* CF=1, S=1011, OF=0 for all three interpretations

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt

Cases for clearing the overflow flag (2) CF=1, OF=0

- no signed integer overflow (CF=0 means correct S)

| * unsigned subtraction | | signed addition | | * signed subtraction |
|--------------------------|--|---|--|----------------------|
| 0011 (3) | | 0010 | | 0011 (+3) |
| -1110 -(14) SUB | | +0010 +(2) | | -1110 -(-2) SUB |
| ----- | | ----- | | ----- |
| 10101 (5) (-16) | | 00101 (+5) | | 00101 (+5) |
| CF=1 | | p + p -> p (OF=0) | | p - n -> p (OF=0) |
| OF meaningless | | -> correct S | | -> correct S |
| S = 0101 | | S = 0101 | | S = 0101 |
| * think hand subtraction | | * OF <- C4 XOR C3 = 0 XOR 0 = 0 of signed addition | | |

* CF=1, S=0101, OF=0 for all three interpretations

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt

Cases for clearing the overflow flag (3) CF=0, OF=0

- no signed integer overflow (CF=0 means correct S)

| * unsigned addition | | * signed addition | signed subtraction |
|---------------------|--|---------------------------------|--------------------|
| 0011 (3) | | 0010 | |
| +0010 +(2) ADD | | 0011 (+3) | 0011 (+3) |
| ----- | | +0010 +(2) ADD | -1110 -(-2) |
| 00101 (5) (+0) | | ----- | ----- |
| | | 00101 (+5) | 00101 (+5) |
| OF=0 | | | |
| | | p + p -> p (OF=0) | p - n -> p (OF=0) |
| OF meaningless | | -> correct S | -> correct S |
| S = 0101 | | S = 0101 | S = 0101 |
| * think hand | | * OF <- C4 XOR C3 = 0 XOR 0 = 0 | |
| addition | | of signed addition | |

* CF=0, S=0101, OF=0 for all three interpretations

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt

Cases for clearing the overflow flag (4) CF=0, OF=0

- no signed integer overflow (CF=0 means correct S)

| * unsigned addition | * signed addition | signed subtraction |
|---------------------|---------------------------------|--------------------|
| 1101 (13) | 1100 | |
| -0010 -(2) SUB | +1110 +(-2) | 1101 (-3) |
| ----- | ----- | -0010 -(+2) SUB |
| 11011 (11) (-16) | 11011 (-5) | ----- |
| | | 11011 (-5) |
| OF=0 | n + n -> n (OF=0) | n - p -> n (OF=0) |
| | | |
| OF meaningless | -> correct S | -> correct S |
| S = 1011 | S = 1011 | S = 1011 |
| | | |
| * think hand | * OF <- C4 XOR C3 = 1 XOR 1 = 0 | |
| subtraction | of signed addition | |

* CF=0, S=1011, OF=0 for all three interpretations

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt