

Applications of Arrays (1A)

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- Viewing an **array** as a **pointer**
- Viewing a **pointer** as an **array**

- Viewing an **array** as a **pointer**

`int a[4] ;`

an array **a**



`int (*a)`

view **a** as a pointer

generalization

- virtual pointer
- no real memory location
- constraints :
 $\text{value}(\&a) = \text{value}(a)$

- Viewing a **pointer** as an **array**

`int (*a) ;`

a pointer **a**



`int a[N]`

view **a** as an array

a specific instance

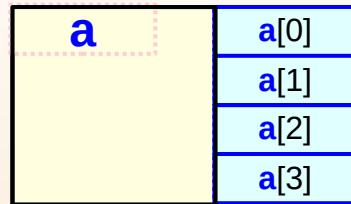
N is not fixed

`sizeof(a)` is
not the size of the array
but of a pointer variable

Array **a** and pointer **a**

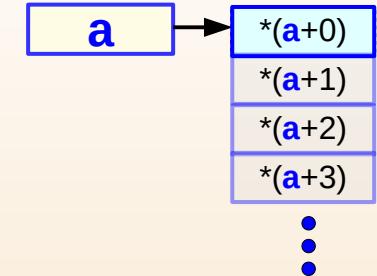
`int a[4] ;`

an array **a**



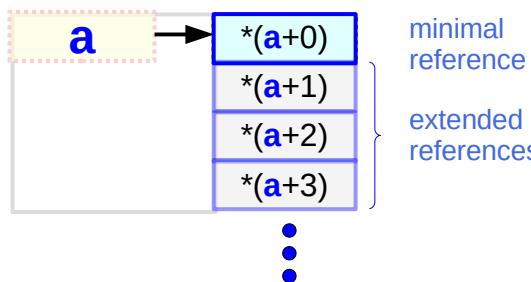
`int (*a) ;`

a pointer **a**



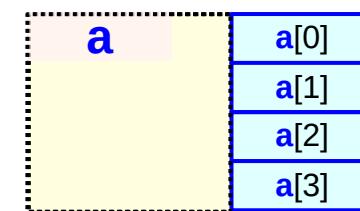
`int (*a)`

a as a pointer



`int a[N]`

a as an array



Array **a** and pointer **a**

`int a[4] ;`

an array **a**

- `type(a)` = `int [4]`
- `sizeof(a)` = an array size (16 bytes)
- `value(&a)` = `value(a)`
- fixed number of elements

`int (*a) ;`

a pointer **a**

- `type(a)` = `int (*)`
- `sizeof(a)` = a pointer size (4 bytes)
- `value(&a)` ≠ `value(a)`
- variable number of elements

`int (*a)`

a as a pointer

a is not a real pointer

- `sizeof(a)` = an array size
- `value(&a)` = `value(a)`

`int a[N]`

a as an array

a is not a real array

- `sizeof(a)` ≠ an array size
= a pointer size
- `value(&a)` ≠ `value(a)`
= assigned address

Relationship between array and array pointer types

`int b[4][2] ;` declare a **2-d array b**



generalization

`int (*b) [2]` **b as a 1-d array pointer**

`int a[4] ;` declare a **1-d array a**



generalization

`int (*a)` **a as a 0-d array pointer**

`int (*b)[2] ;` declare a **1-d array pointer b**



a specific instance

`int b[N][2]` **b as a 2-d array**

`int (*a) ;` declare a **0-d array pointer a**



a specific instance

`int a[N]` **a as a 1-d array**

Array **b** and array pointer **b**

`int b[4][2] ;`

2-d array b

- `type(b)` = `int [4]`
- `sizeof(b)` = an array size (32 bytes)
- `value(&b)` = `value(b)`
- fixed number of elements

`int (*b) [2] ;`

1-d array pointer b

- `type(b)` = `int (*)`
- `sizeof(b)` = a pointer size (4 bytes)
- `value(&b)` ≠ `value(b)`
- variable number of elements

`int (*) [2]`

b as a 1-d array pointer

b is not a real pointer

- `sizeof(b)` = an array size
- `value(&b)` = `value(b)`

`int [N][2]`

b as a 2-d array

b is not a real array

- `sizeof(b)` ≠ an array size
= a pointer size
- `value(&b)` ≠ `value(b)`
= assigned address

Array **b** and array pointer **b**

int b[4][2] ;

2-d array b

b	b[0]	b[0][0] b[0][1]
	b[1]	b[1][0] b[1][1]
	b[2]	b[2][0] b[2][1]
	b[3]	b[3][0] b[3][1]

int (*b) [2] ;

1-d array pointer b

b	→	*(b+0)	(*b+0)[0] (*b+0)[1]
		*(b+1)	(*b+1)[0] (*b+1)[1]
		*(b+2)	(*b+2)[0] (*b+2)[1]
		*(b+3)	(*b+3)[0] (*b+3)[1]
		...	

minimal reference

extended references

int (*) [2]

b as a 1-d array pointer

b	→	*(b+0)	(*b+0)[0] (*b+0)[1]
		*(b+1)	(*b+1)[0] (*b+1)[1]
		*(b+2)	(*b+2)[0] (*b+2)[1]
		*(b+3)	(*b+3)[0] (*b+3)[1]
		...	

- virtual pointer
- no real memory location
- constraints : $\&b = b$

minimal reference

extended references

int [N][2]

b as a 2-d array

b	b[0]	b[0][0] b[0][1]
	b[1]	b[1][0] b[1][1]
	b[2]	b[2][0] b[2][1]
	b[3]	b[3][0] b[3][1]

N is not fixed to 4

`sizeof(b)` is not the size of the array but the size of a pointer variable

Dual type - relaxing the 1st dimension of an array

`int [4][2]` **2-d array**

more constrained type

relaxing the
1st dimension
generalization



a specific instance

`int (*)[2]` **1-d array pointer**

more general type

`int [4]` **1-d array**

more constrained type

relaxing the
1st dimension
generalization



a specific instance

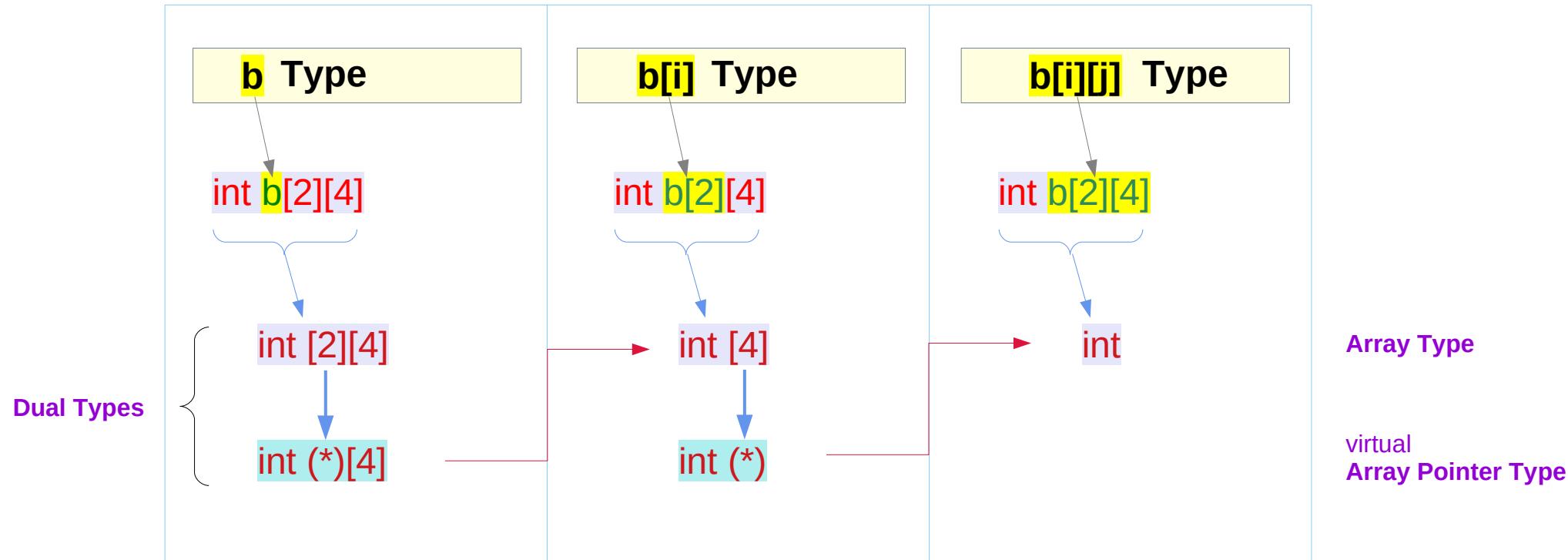
`int (*)` **0-d array pointer**

more general type

Subarray types in a 2-d array

```
int b[2][4];
```

2-d array b

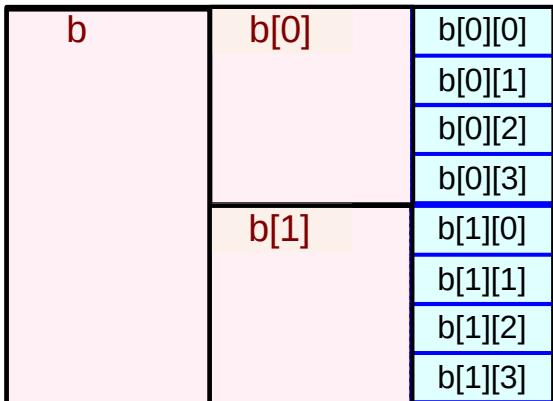


Dual types in a 2-d array

int b[2][4];

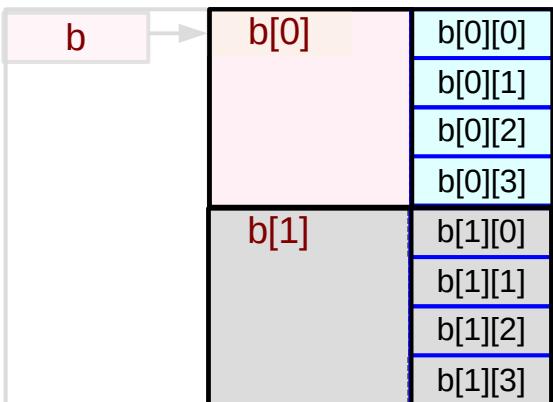
2-d array b

int [2][4]



int (*[4]

int [4]



Dual Types

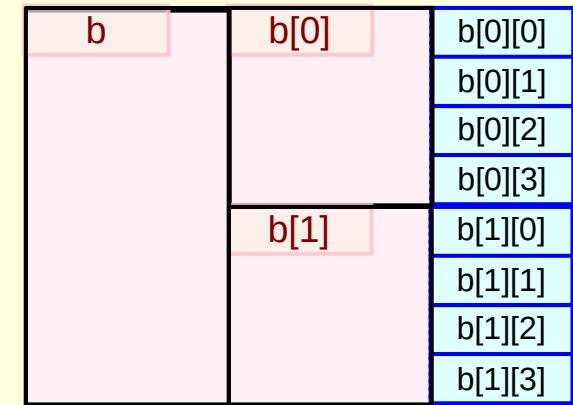
int [2][4]

int (*[4]

int [4]

int (*)

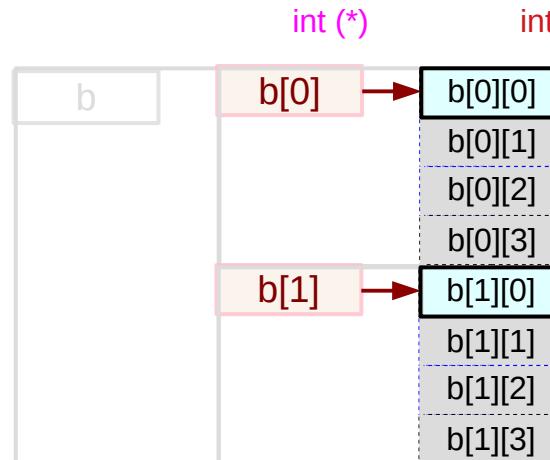
int



int (*)

int

int



Subarray type examples

```
int a[4];
```

			relaxed type	virtual
a	int [4]	1-d array type	int (*)	0-d array pointer type
a[i]	int	0-d array type		

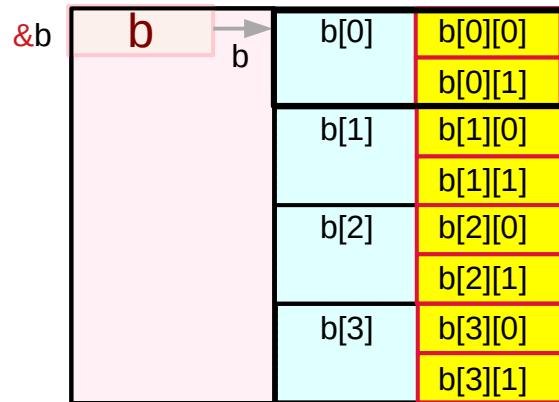
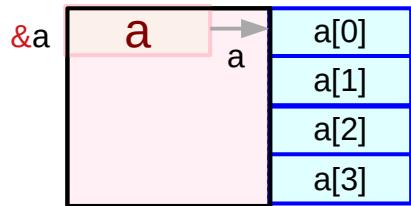
```
int b[2][4];
```

			relaxed type	virtual
b	int [2][4]	2-d array type	int (*)[4]	1-d array pointer type
b[i]	int [4]	1-d array type	int (*)	0-d array pointer type
b[i][j]	int	0-d array type		

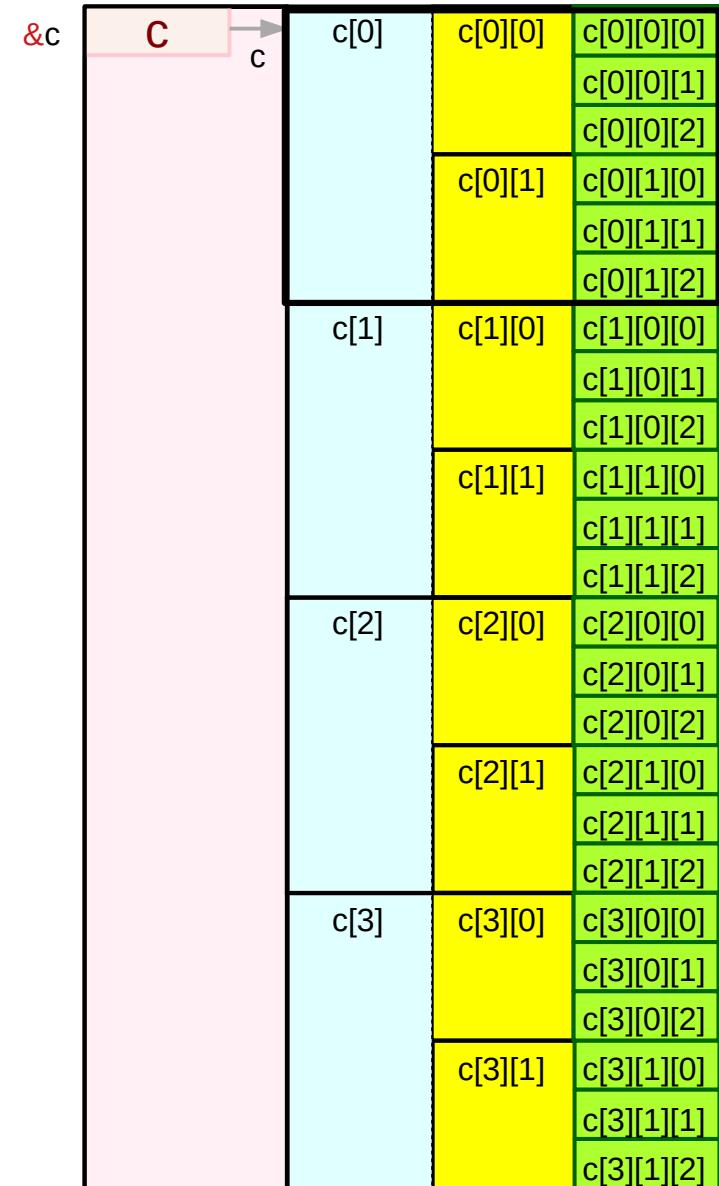
```
int c[4][2][3];
```

			relaxed type	virtual
c	int [4][2][3]	3-d array type	int (*)[2][3]	2-d array pointer type
c[i]	int [4][2]	2-d array type	int (*)[2]	1-d array pointer type
c[i][j]	int [4]	1-d array type	int (*)	0-d array pointer type
c[i][j][k]	int	0-d array type		

Types of a, b, c arrays



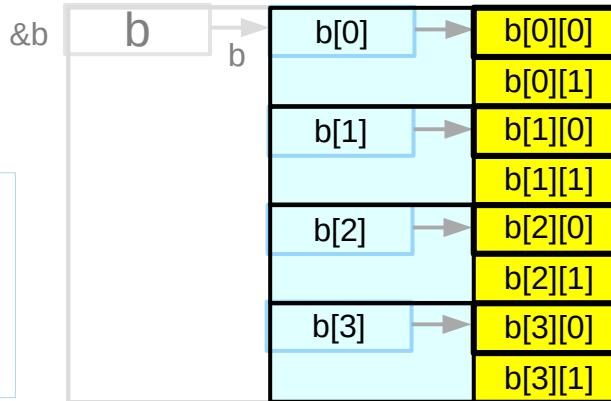
```
int a[4];  
int b[2][4];  
int c[4][2][3];
```



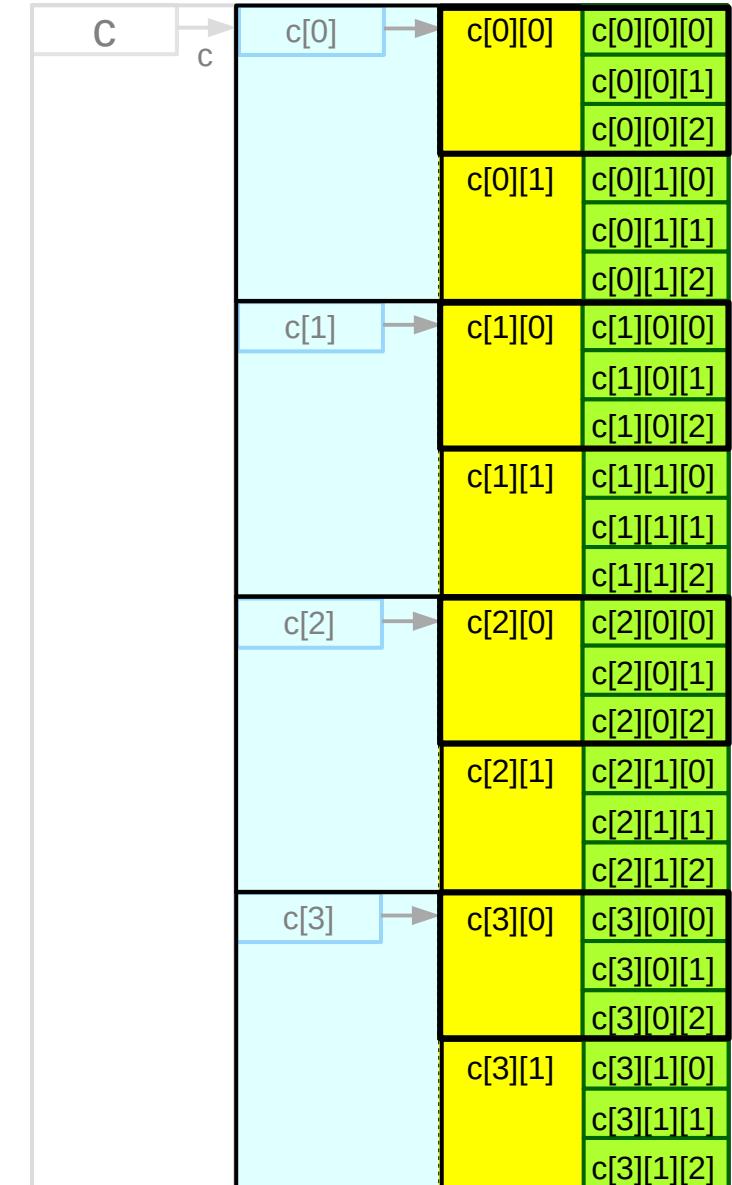
dual types

int [4]	1-d array a	a[i]
int (*)	0-d array pointer a (virtual)	*(a+i)
int [4][2];	2-d array b	b[i]
int (*)[2];	1-d array pointer b (virtual)	*(b+i)
int [4][2][3];	3-d array c	c[i]
int (*)[2][3];	2-d array pointer c (virtual)	*(c+i)

Types of $b[i]$, $c[i]$ subarrays



```
int a[4];  
int b[2][4];  
int c[4][2][3];
```



dual types

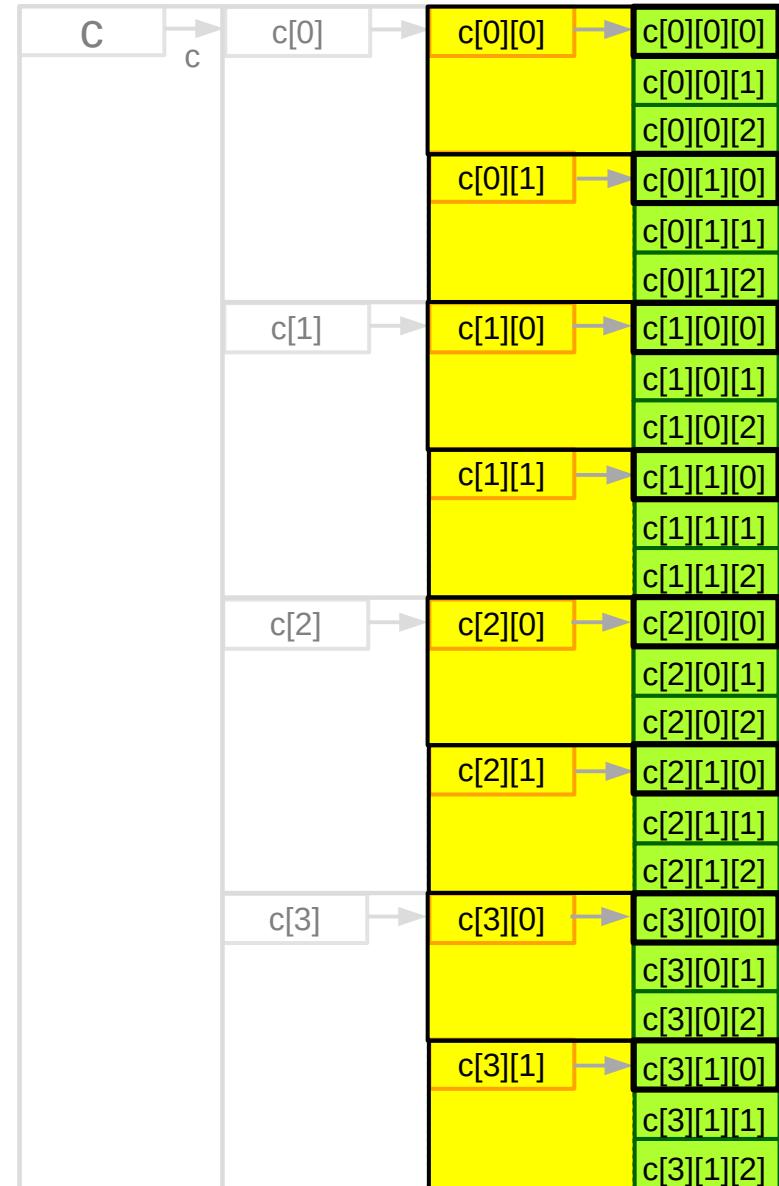
int [2]	1-d array $b[i]$	$b[i][j]$
int (*)	0-d array pointer $b[i]$ (virtual)	$*(b[i]+j)$
int [2][3];	2-d array $c[i]$	$c[i][j]$
int (*)[3];	1-d array pointer $c[i]$ (virtual)	$*(c[i]+j)$

Types of $c[i][j]$ subarrays

```
int a[4];
int b[2][4];
int c[4][2][3];
```

dual types

int [3]	1-d array $c[i][j]$	$c[i][j][k]$
int (*)	0-d array pointer $c[i][j]$ (virtual)	$*(c[i][j]+k)$



Types of a 4-d array and its subarrays

`int d[4][2][3][4];`

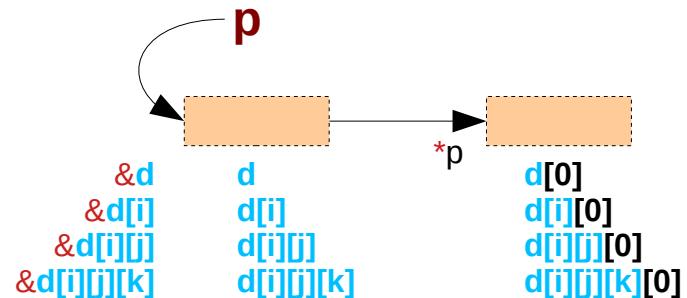
types

<code>d</code>	consider <code>d[4][2][3][4]</code> relax the 1 st dimension	→ int [4][2][3][4] → int (*)[2][3][4]	→ 4-d array → 3-d array pointer (virtual)
<code>d[i]</code>	consider <code>d[i][2][3][4]</code> relax the 1 st dimension	→ int [2][3][4] → int (*)[3][4]	→ 3-d array → 2-d array pointer (virtual)
<code>d[i][j]</code>	consider <code>d[i][j][3][4]</code> relax the 1 st dimension	→ int [3][4] → int (*)[4]	→ 2-d array → 1-d array pointer (virtual)
<code>d[i][j][k]</code>	consider <code>d[i][j][k][4]</code> relax the 1 st dimension	→ int [4] → int (*)	→ 1-d array → 0-d array pointer (virtual)

`i,j,k` are specific index values $i=[0..3]$, $j=[0..1]$, $k=[0..2]$

Initializing n -d array pointers with n -d subarrays

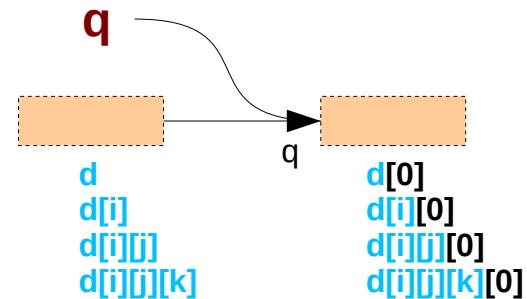
`int d[4][2][3][4];`



<code>d</code>	4-d array	<code>d[4][2][3][4]</code>	<code>p = &d</code>	abstract data
<code>p</code>	4-d array pointer	<code>(*p)[4][2][3][4]</code>	<code>int (*p)[4][2][3][4] = &d;</code> <code>(*p)[i][j][k][l] ≡ d[i][j][k][l]</code>	
<code>d[i]</code>	3-d array	<code>d[i][2][3][4]</code>	<code>p = &d[i]</code>	abstract data
<code>p</code>	3-d array pointer	<code>(*p)[2][3][4]</code>	<code>int (*p)[3][4] = &d[i];</code> <code>(*p)[j][k][l] ≡ d[i][j][k][l] given i</code>	
<code>d[i][j]</code>	2-d array	<code>d[i][j][3][4]</code>	<code>p = &d[i][j]</code>	abstract data
<code>p</code>	2-d array pointer	<code>(*p)[3][4]</code>	<code>int (*p)[4] = &d[i][j];</code> <code>(*p)[k][l] ≡ d[i][j][k][l] given i, j</code>	
<code>d[i][j][k]</code>	1-d array	<code>d[i][j][k][4]</code>	<code>p = &d[i][j][k]</code>	abstract data
<code>p</code>	1-d array pointer	<code>(*p)[4]</code>	<code>int (*p) = &d[i][j][k];</code> <code>(*p)[l] ≡ d[i][j][k][l] given i, j, k</code>	

Initializing $(n-1)$ -d array pointers with n -d subarrays

`int d[4][2][3][4];`



<code>d</code>	4-d array	<code>d[4][2][3][4]</code>	<code>q = d</code>	virtual pointer
<code>q</code>	3-d array pointer	<code>(*q)[2][3][4]</code>	<code>int (*q)[2][3][4] = d;</code> $q[i][j][k][l] \equiv d[i][j][k][l]$	
<code>d[i]</code>	3-d array	<code>d[i][2][3][4]</code>	<code>q = d[i]</code>	virtual pointer
<code>q</code>	2-d array pointer	<code>(*q)[3][4]</code>	<code>int (*q)[3][4] = d[i];</code> $q[j][k][l] \equiv d[i][j][k][l]$ given i	
<code>d[i][j]</code>	2-d array	<code>d[i][j][3][4]</code>	<code>q = d[i][j]</code>	virtual pointer
<code>q</code>	1-d array pointer	<code>(*q)[4]</code>	<code>int (*q)[4] = d[i][j];</code> $q[k][l] \equiv d[i][j][k][l]$ given i, j	
<code>d[i][j][k]</code>	1-d array	<code>d[i][j][k][4]</code>	<code>q = d[i][j][k]</code>	virtual pointer
<code>q</code>	0-d array pointer	<code>(*q)</code>	<code>int (*q) = d[i][j][k];</code> $q[l] \equiv d[i][j][k][l]$ given i, j, k	

Aggregate Data Types

Abstract Data Types

Virtual Array Pointers

Aggregate data type

an aggregate type

consists of **N** elements

each element

- starting address
- size

int [4][4]

dual type

a virtual pointer

the address and value of a virtual pointer are the same

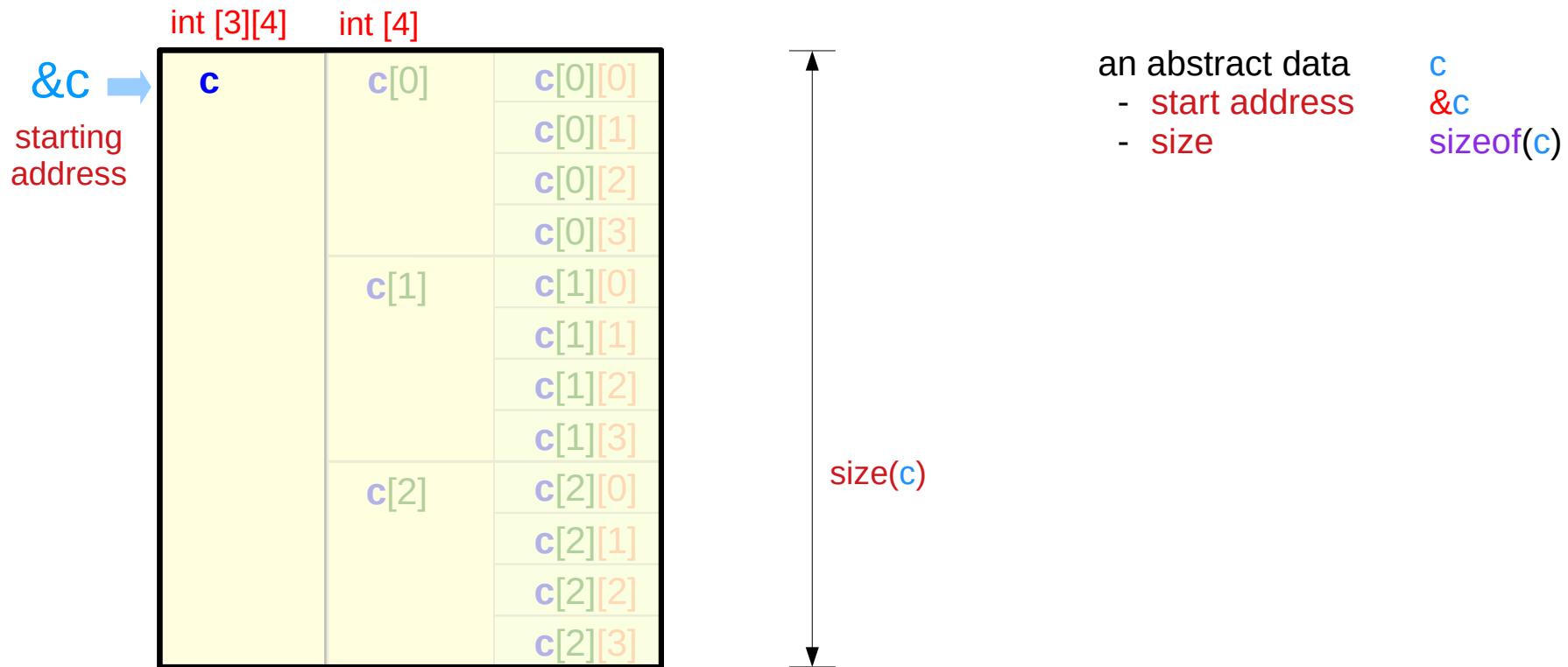
int (*[4]

an abstract type

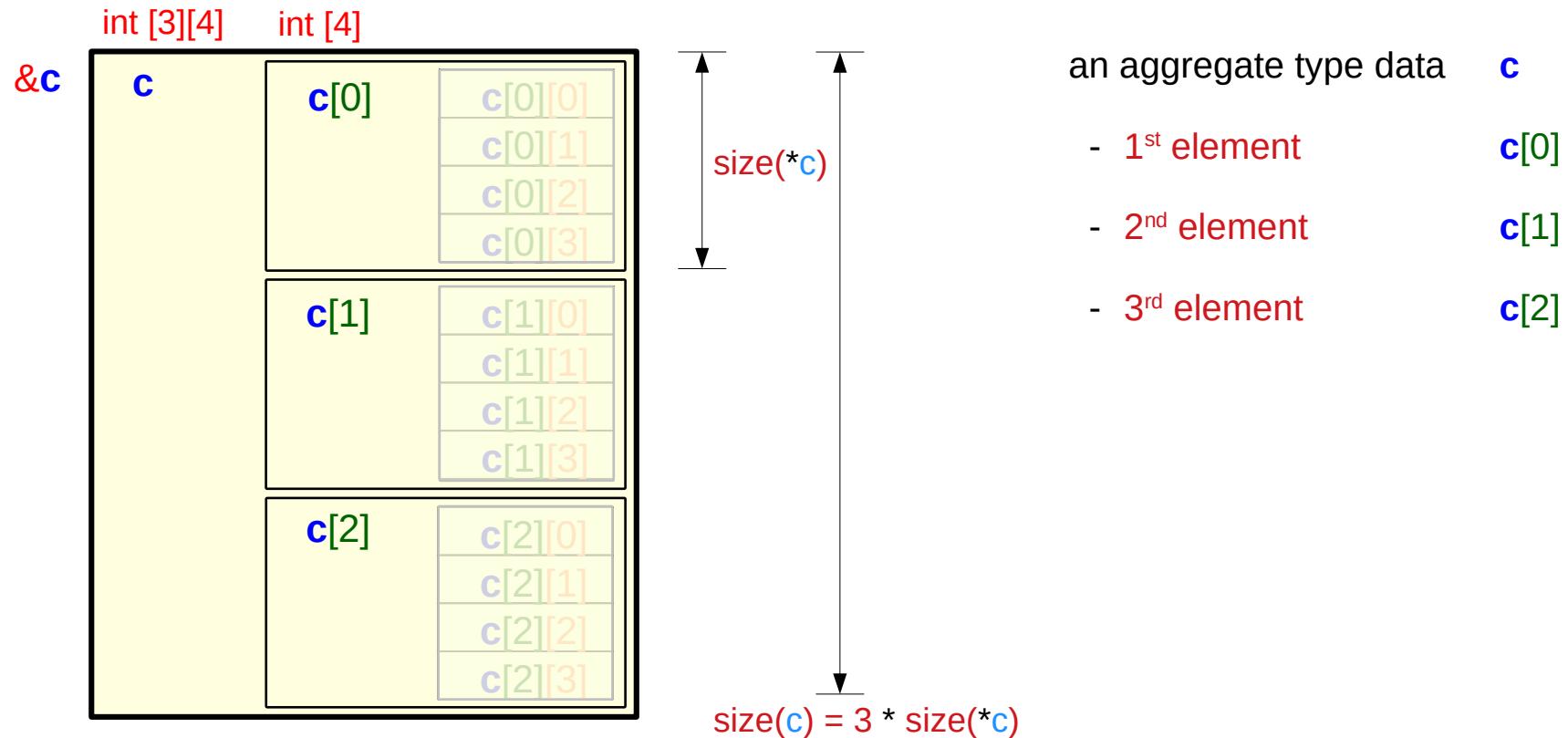
is considered as one unit

- starting address
- size

Abstract data c

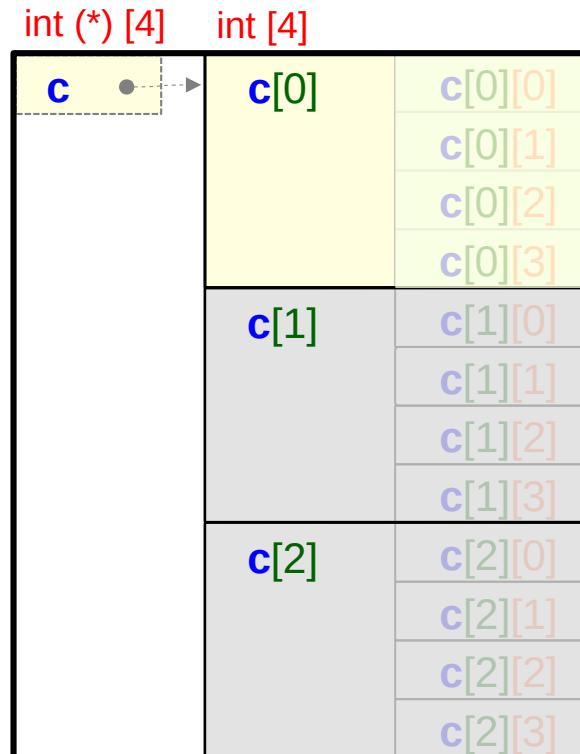


Aggregate data c



Virtual pointer **c**

$\&c = c = \&c[0]$



↑
size(*c)
↓

a virtual pointer **c**
- pointer address **&c**
- pointer value **c = &c[0]**

with the constraint
c = &c

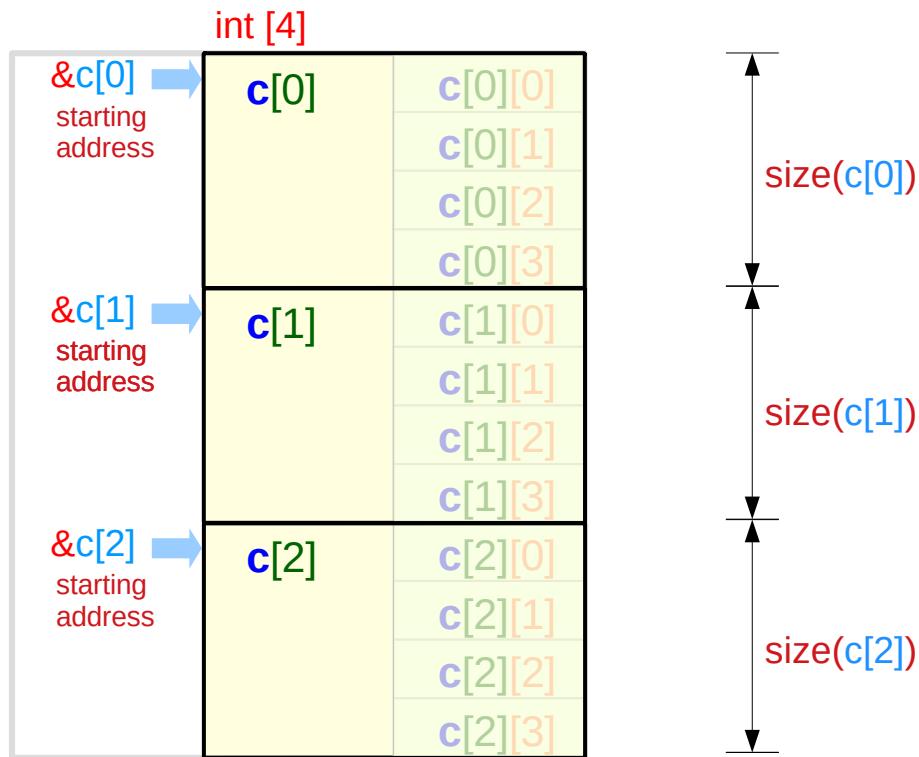
an abstract data **c[0] = *c**
- start address **&c[0] = c**
- size **sizeof(c[0])**

virtual pointer **c** points
to abstract data **c[0]**

virtual pointers

- no physical memory locations are allocated
- address and data have the same value

Abstract data $c[i]$



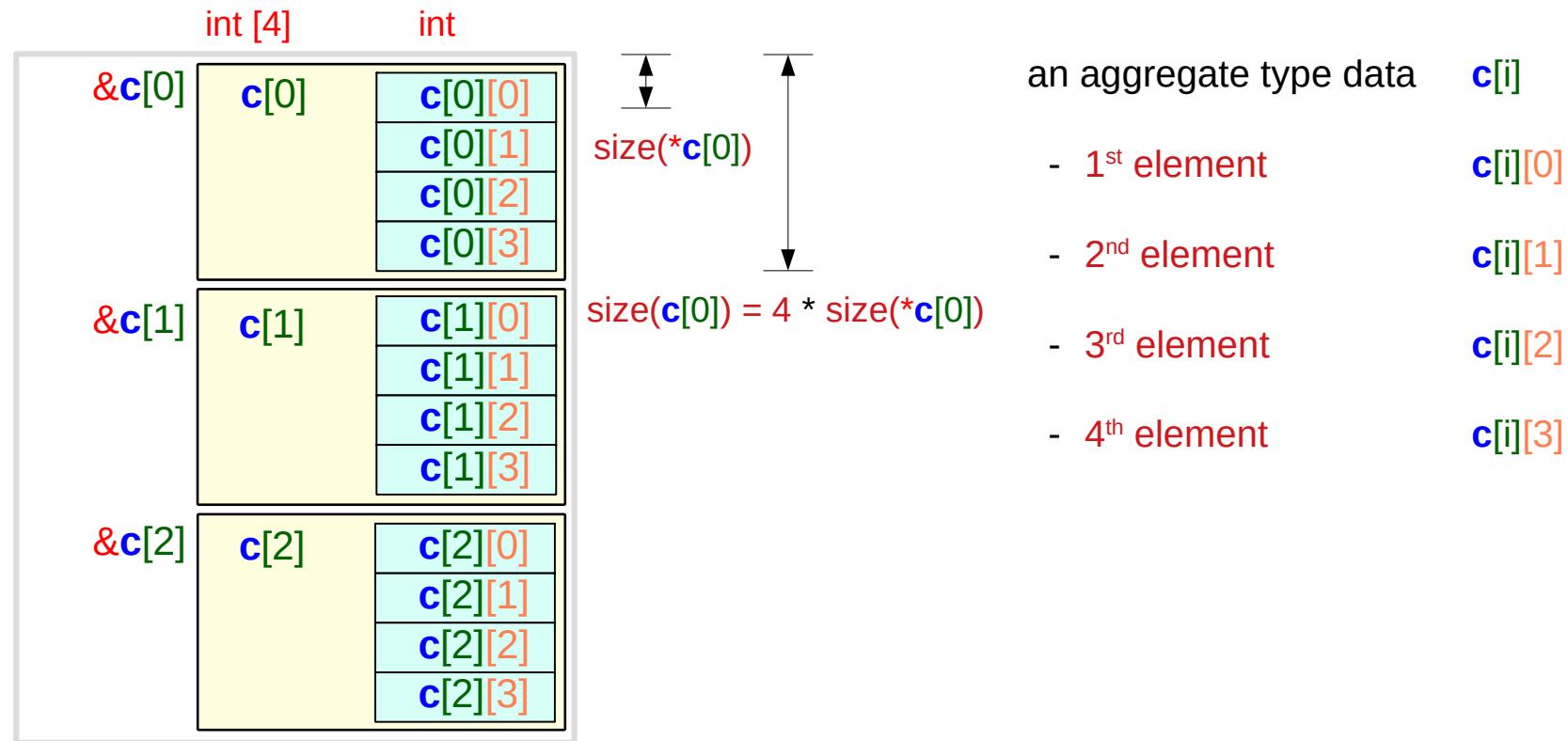
- an abstract data
 - start address
 - size
- an abstract data
 - start address
 - size
- an abstract data
 - start address
 - size

$c[0]$
 $\&c[0]$
 $\text{sizeof}(c[0])$

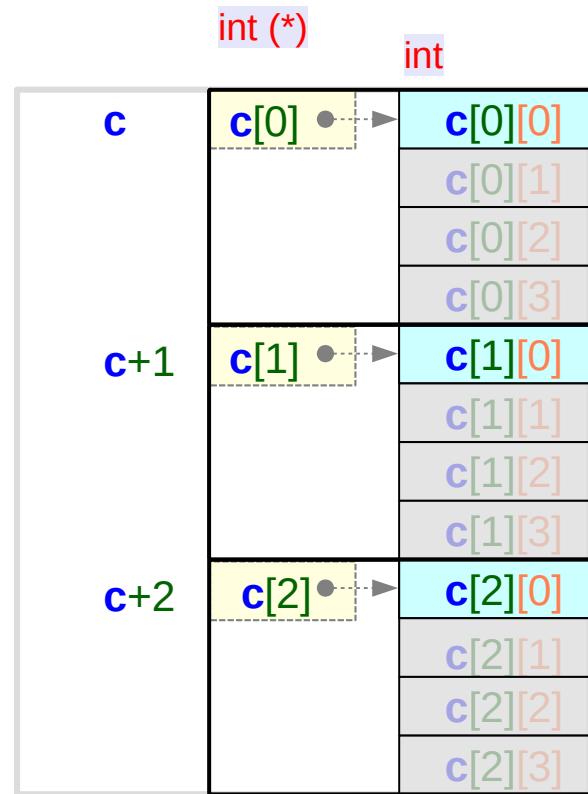
$c[1]$
 $\&c[1]$
 $\text{sizeof}(c[1])$

$c[2]$
 $\&c[2]$
 $\text{sizeof}(c[2])$

Aggregate data $c[i]$



Virtual pointer $c[i]$



a virtual pointer $c[i]$
- pointer address $\&c[i]$
- pointer value $c+i = \&c[i]$

with the constraint
 $c[i] = \&c[i]$

an primitive data $c[i][0] = *c[i]$
- start address $\&c[i][0] = c[i]$
- size $\text{sizeof}(c[i][0])$

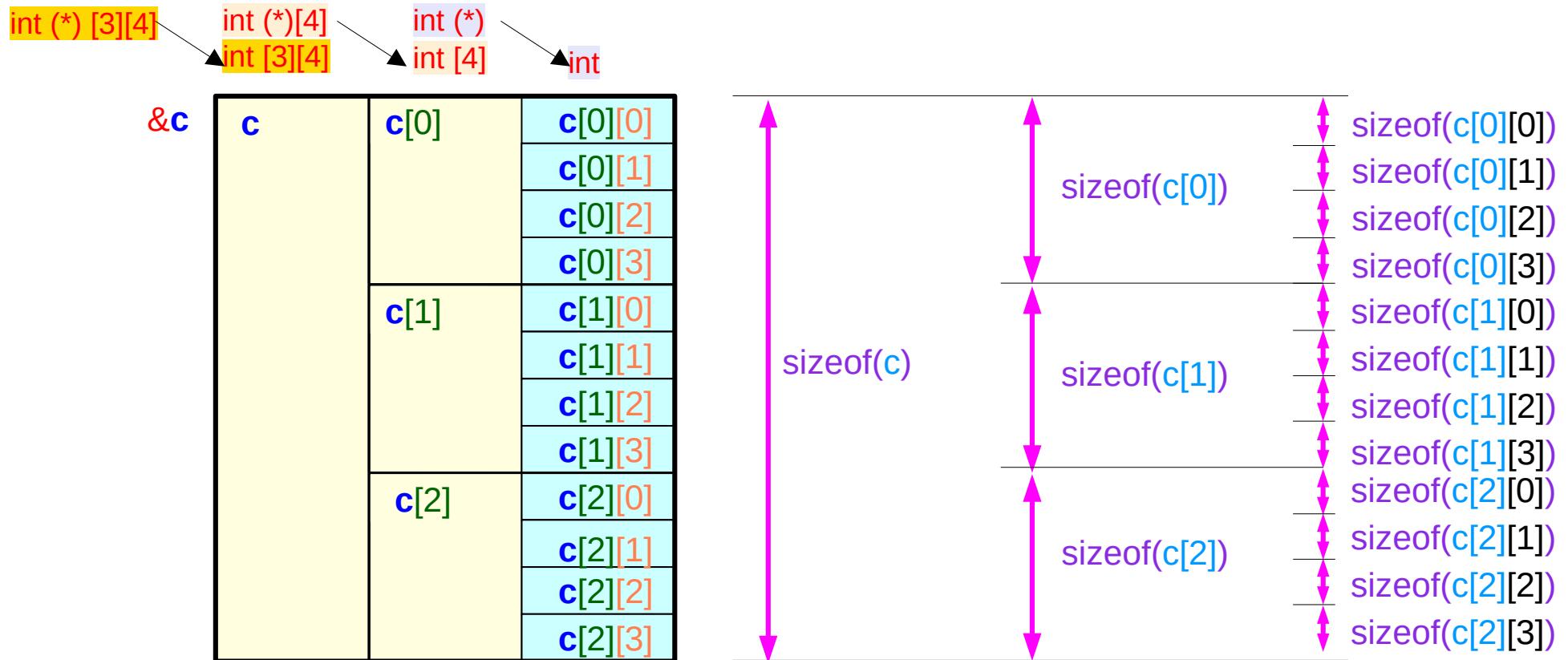
virtual pointer $c[i]$
points to primitive data $c[i][0]$

virtual pointers

- no physical memory locations are allocated
- address and data have the same value

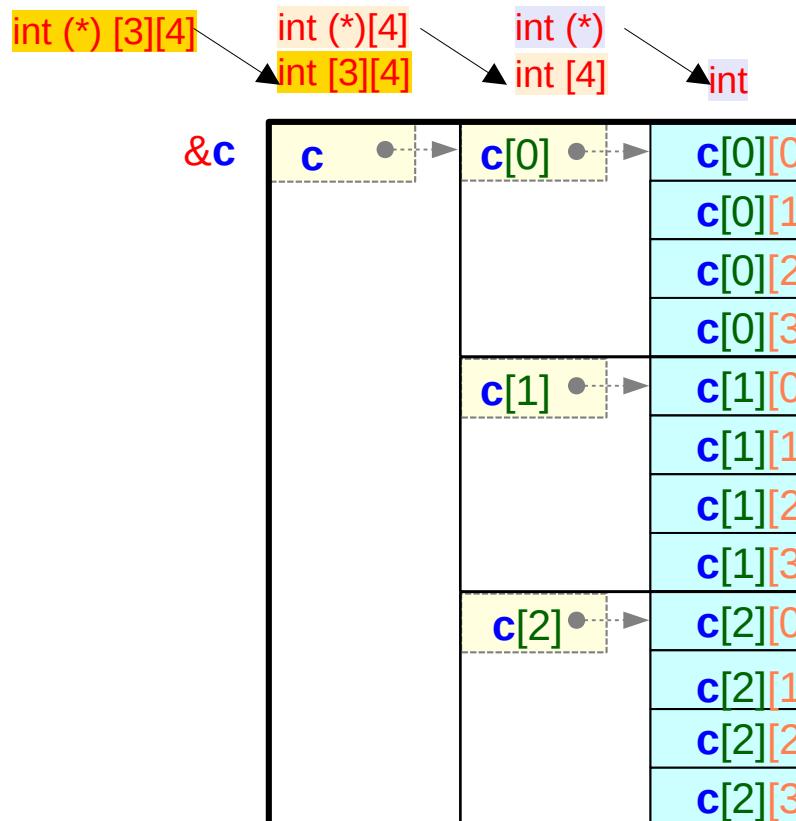
A 2-d array and its 1-d sub-arrays – a size view

```
int c[3][4];
```



A 2-d array and its 1-d sub-arrays – a virtual pointer view

```
int c[3][4];
```



$$\begin{aligned} \text{value}(c) &= \text{value}(c[0]) & = \text{value}(&c[0][0]) \\ \text{value}(&c) &= \text{value}(&c[0]) & = \text{value}(&c[0][0]) \end{aligned}$$

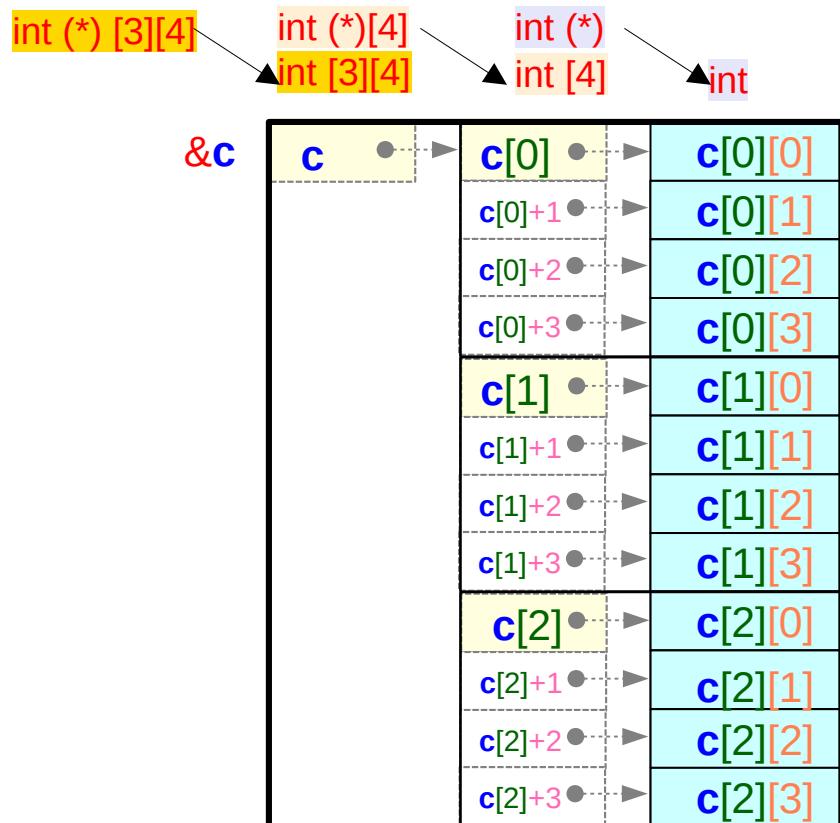
$$\begin{aligned} \text{value}(c[1]) &= \text{value}(&c[1][0]) \\ \text{value}(&c[1]) &= \text{value}(&c[1][0]) \end{aligned}$$

$$\begin{aligned} \text{value}(c[2]) &= \text{value}(&c[2][0]) \\ \text{value}(&c[2]) &= \text{value}(&c[2][0]) \end{aligned}$$

$$\begin{aligned} \text{address}(c) &= \text{address}(c[0]) & = \text{address}(c[0][0]) \\ \text{address}(c[1]) &= \text{address}(c[1][0]) & = \text{address}(c[1][0]) \\ \text{address}(c[2]) &= \text{address}(c[2][0]) & = \text{address}(c[2][0]) \end{aligned}$$

A 2-d array and its 1-d sub-arrays – size relation

```
int c[3][4];
```



`sizeof(c) = sizeof(c[0]) * 3 ... leading element`
`sizeof(c+1) = pointer size (4/8 bytes)`
`sizeof(c+2) = pointer size (4/8 bytes)`

`sizeof(c[0]) = sizeof(c[0][0]) * 4 ... leading element`
`sizeof(c[0]+1) = pointer size (4/8 bytes)`
`sizeof(c[0]+2) = pointer size (4/8 bytes)`
`sizeof(c[0]+3) = pointer size (4/8 bytes)`

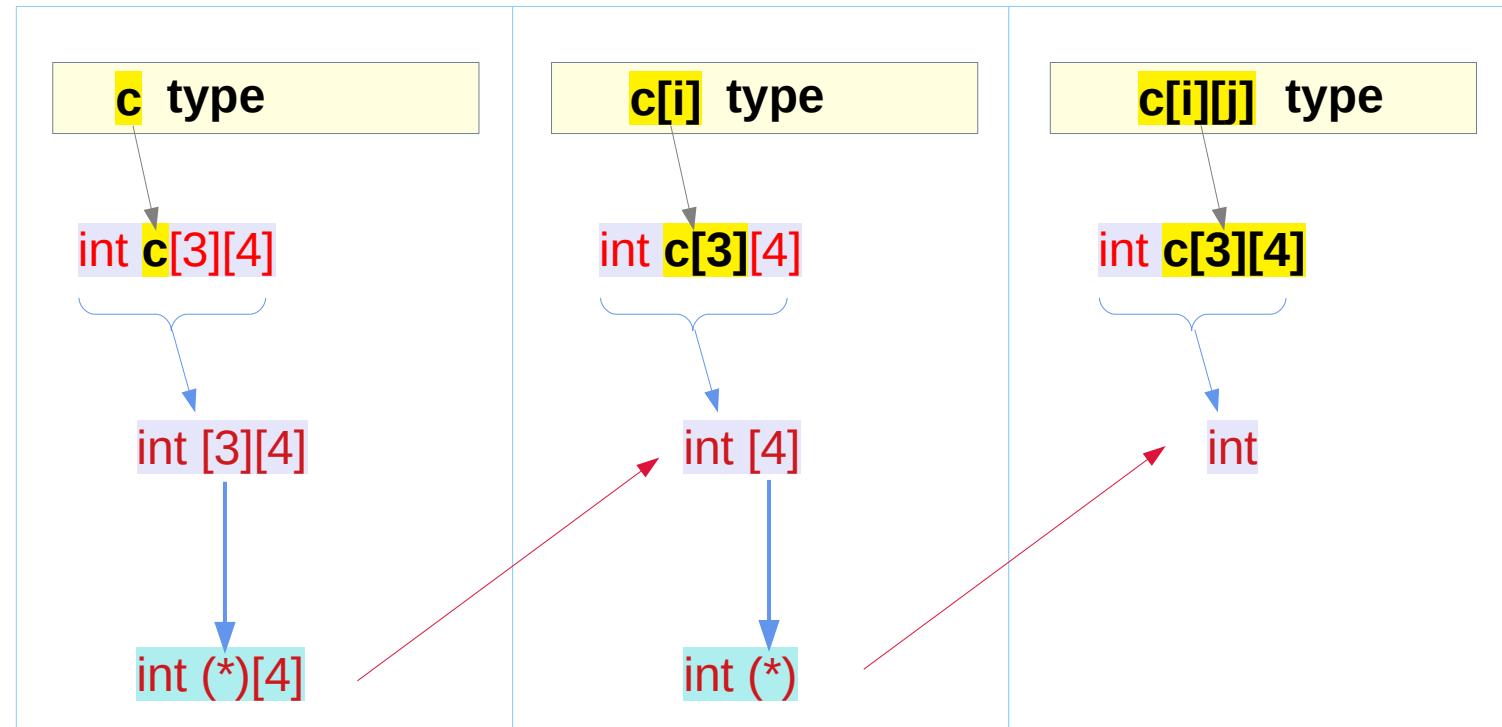
`sizeof(c[1]) = sizeof(c[1][0]) * 4 ... leading element`
`sizeof(c[1]+1) = pointer size (4/8 bytes)`
`sizeof(c[1]+2) = pointer size (4/8 bytes)`
`sizeof(c[1]+3) = pointer size (4/8 bytes)`

`sizeof(c[2]) = sizeof(c[2][0]) * 4 ... leading element`
`sizeof(c[2]+1) = pointer size (4/8 bytes)`
`sizeof(c[2]+2) = pointer size (4/8 bytes)`
`sizeof(c[2]+3) = pointer size (4/8 bytes)`

Sub-array types in a 2-d array

```
int c[3][4];
```

2-d array **c**



Dual Types

- Identifying nested arrays

in a 2-d array declaration

Nested arrays in a 2-d array declaration

int c[3] [4] ;

int c[3] [4] ;

c : a 3 element array
c[i] : each element

int c[3] [4] ;

c[i]'s type 1 : int [4]
an array of 4 integers

int c[3] [4] ;
relaxed dimension

c[i]'s type 2: int (*)
a pointer to an integer

Nested arrays

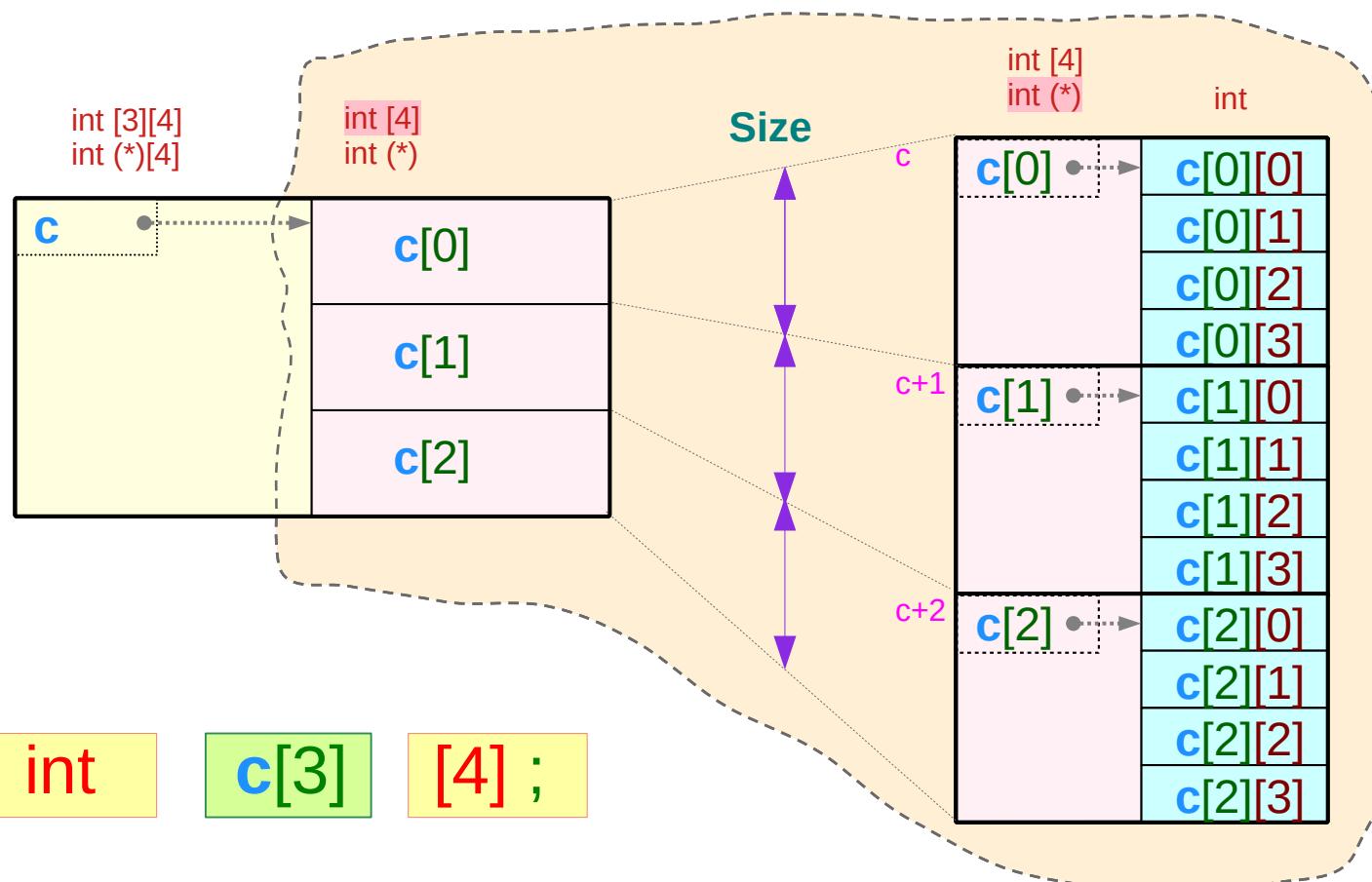
c[3]

c : a 3 element array
c[i] : each element

int

[4] ;

c[i]'s type 1 : int [4]
c[i]'s type 2 : int (*)



Address

&c[0][0] → c[0] → c

&c[1][0] → c[1]

&c[2][0] → c[2]

c : 3-element array

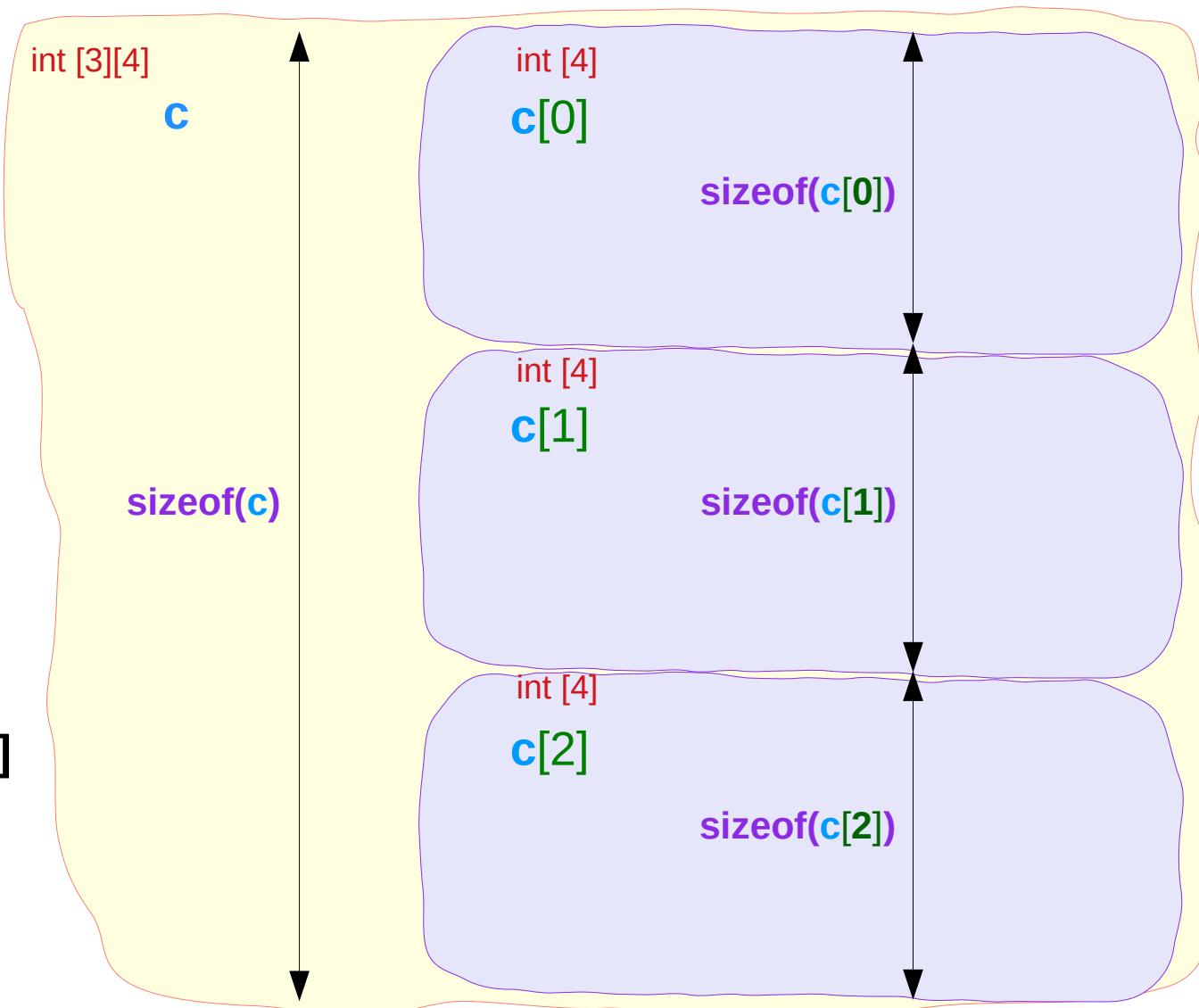
c	2-d array	int [3][4]
c[i]	1-d array	int [4]

```
int c[3][4];
```

3-element array c

abstract data element **c[i]**

each element **c[i]** has the
1-d array type **int [4]**



c : pointer to a 4-element array

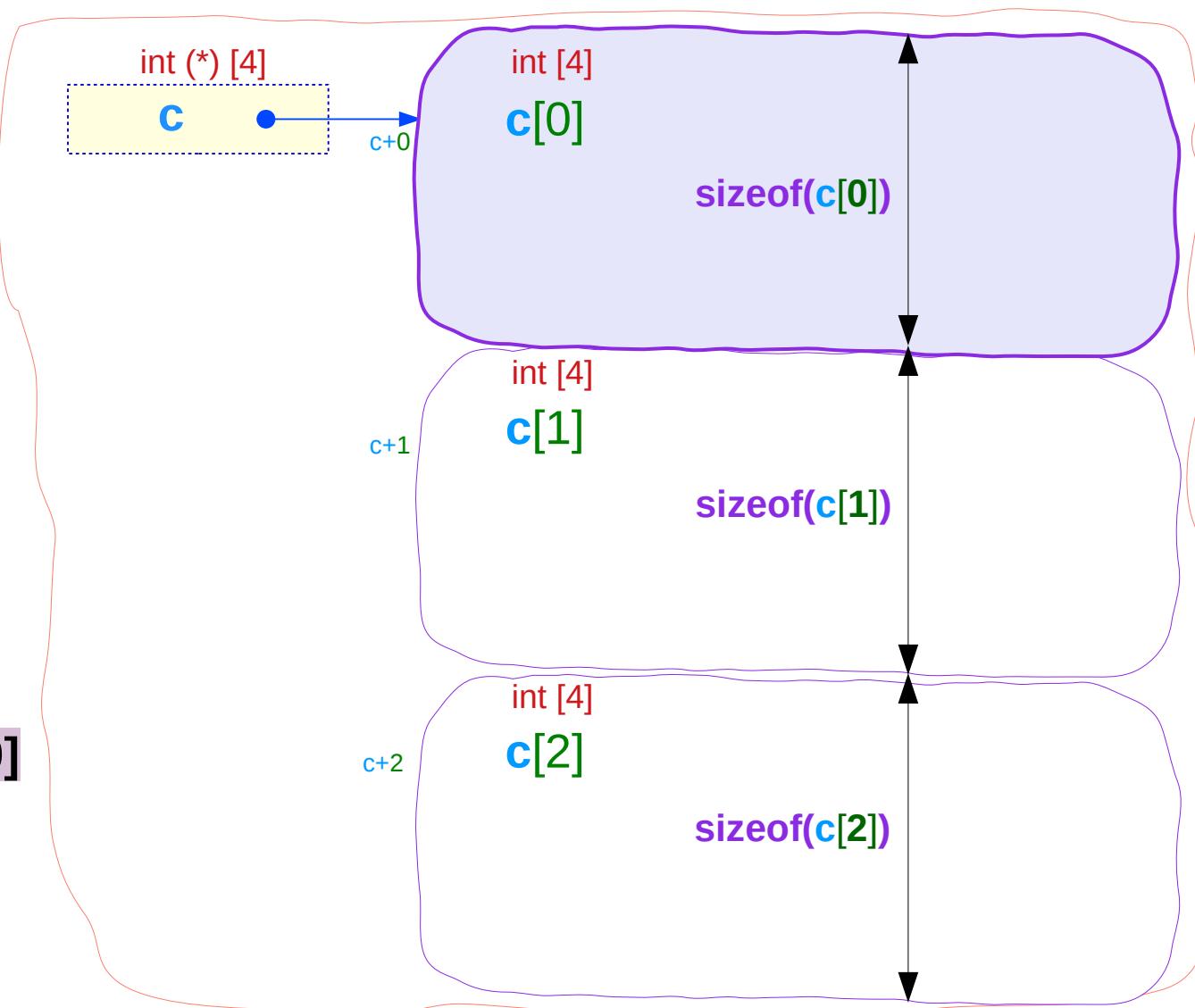
c	1-d array pointer	$\text{int } (*)[4]$
c[i]	1-d array	$\text{int } [4]$

`int` **c** [3] [4] ;
relaxed dimension

pointer c

abstract data element **c[0]**

each element **c[i]** has the
1-d array type $\text{int } [4]$



c[i] : 4-element array

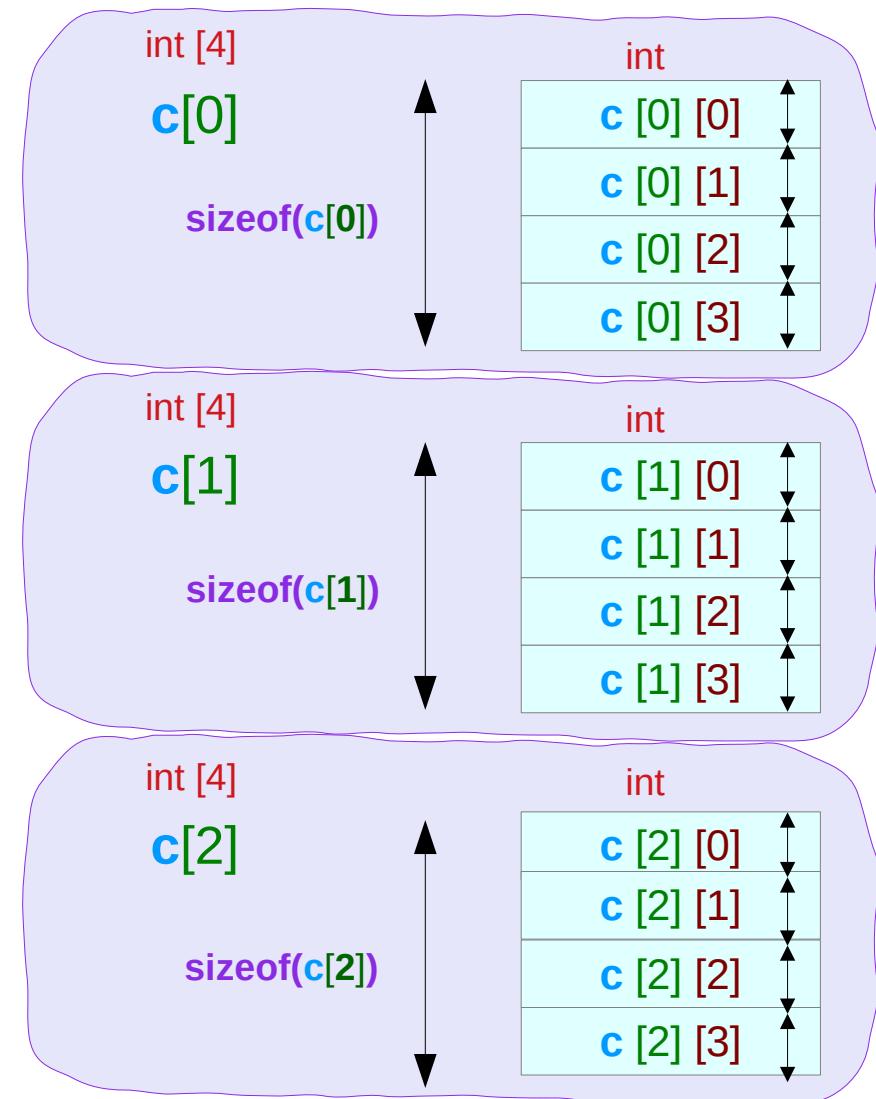
c[i]	1-d array	int [4]
c[i][j]	0-d array	int

```
int [ ] c [3] [4] ;
```

4-element array **c[i]**

primitive data element **c[i][j]**

each element **c[i][j]** has
the primitive type **int**



c[i] : pointer to a primitive data

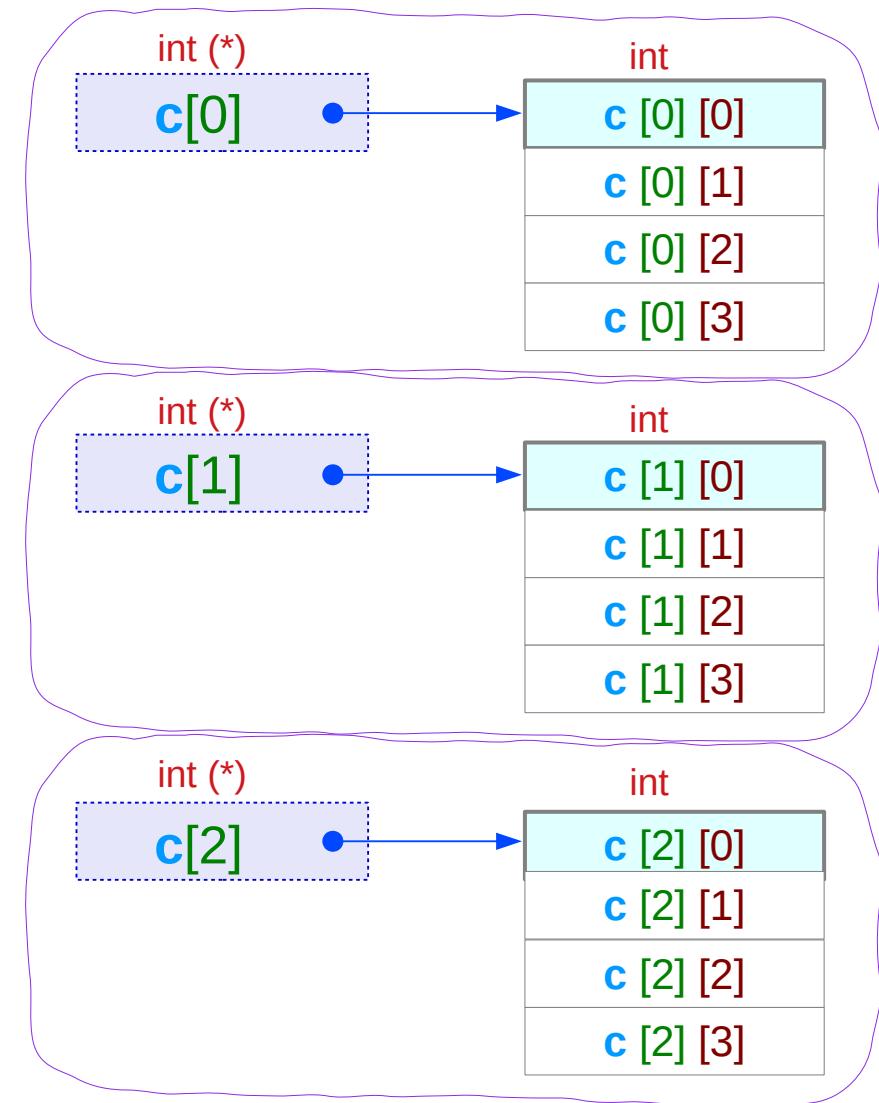
c[i]	0-d array pointer	int (*)
c[i][j]	0-d array	int

int c [3] [4];
 relaxed dimension

pointer **c[i]**

primitive data element **c[i][0]**

each element **c[i][j]** has
the primitive type **int**

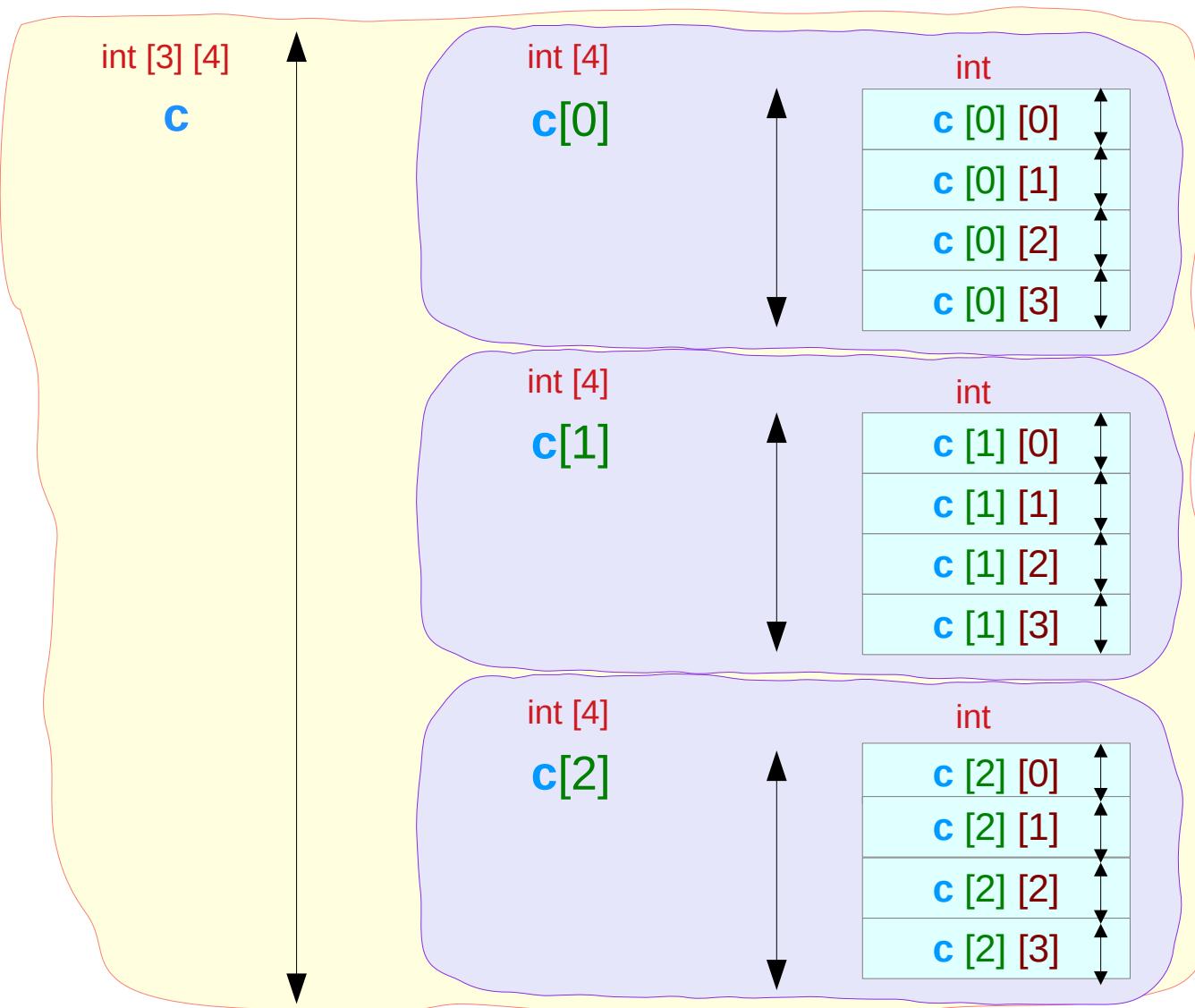


Recursive data view

<code>c</code>	2-d array	<code>int [3][4]</code>
<code>c</code>	1-d array pointer	<code>int (*)[4]</code>
<code>c[i]</code>	1-d array	<code>int [4]</code>
<code>c[i]</code>	0-d array pointer	<code>int (*)</code>
<code>c[i][j]</code>	0-d array	<code>int</code>

`int` `c[3]` `[4] ;`

3-element array `c`
4-element array `c[i]`



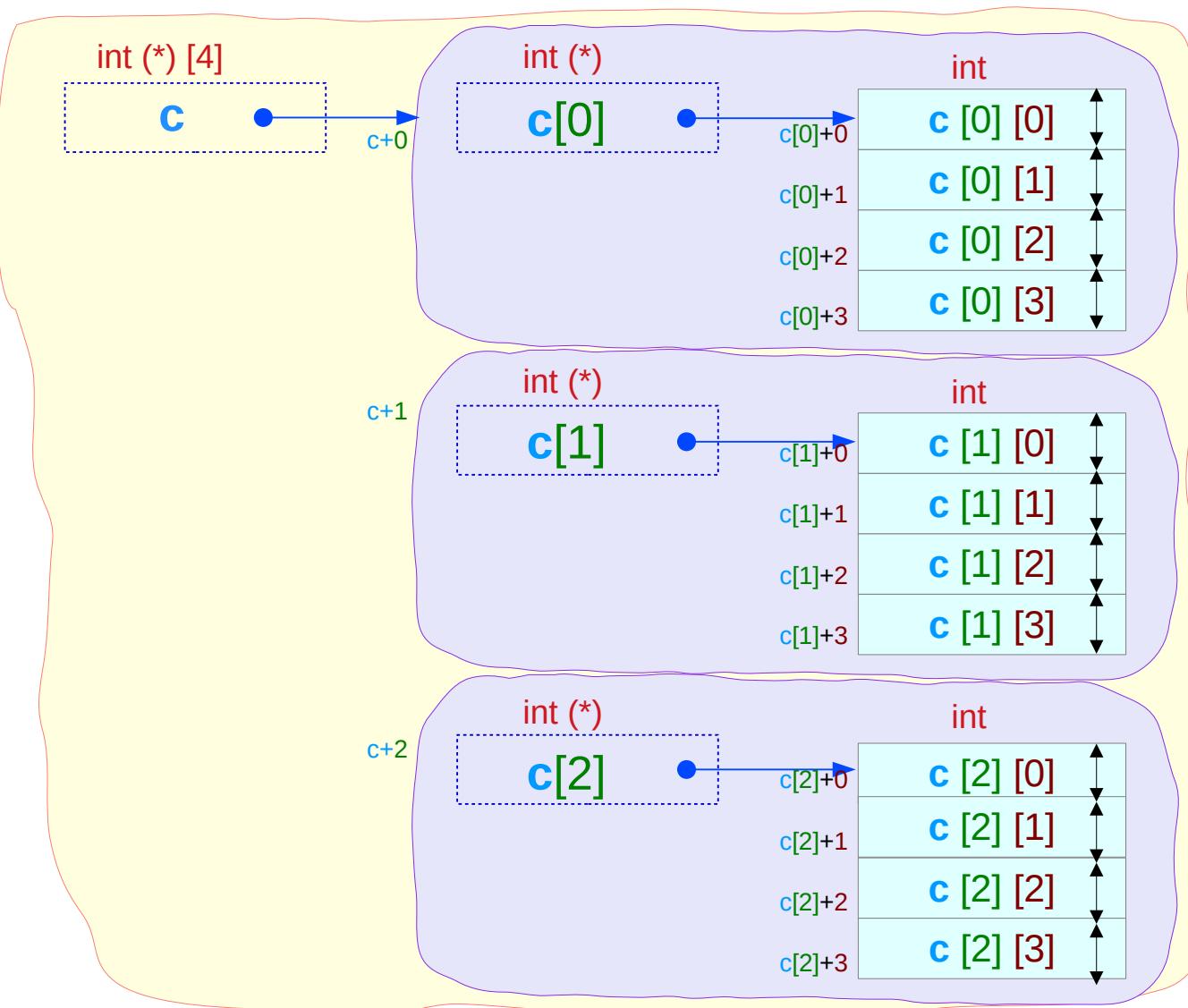
Pointer view

c	2-d array	int [3][4]
c	1-d array pointer	int (*)[4]
c[i]	1-d array	int [4]
c[i]	0-d array pointer	int (*)
c[i][j]	0-d array	int

int c[3] [4] ;

$$\begin{aligned} v(c) &= v(c[0]) = v(\&c[0][0]) \\ v(c[1]) &= v(\&c[1][0]) \\ v(c[2]) &= v(\&c[2][0]) \end{aligned}$$

v ≡ value



1-d array pointer

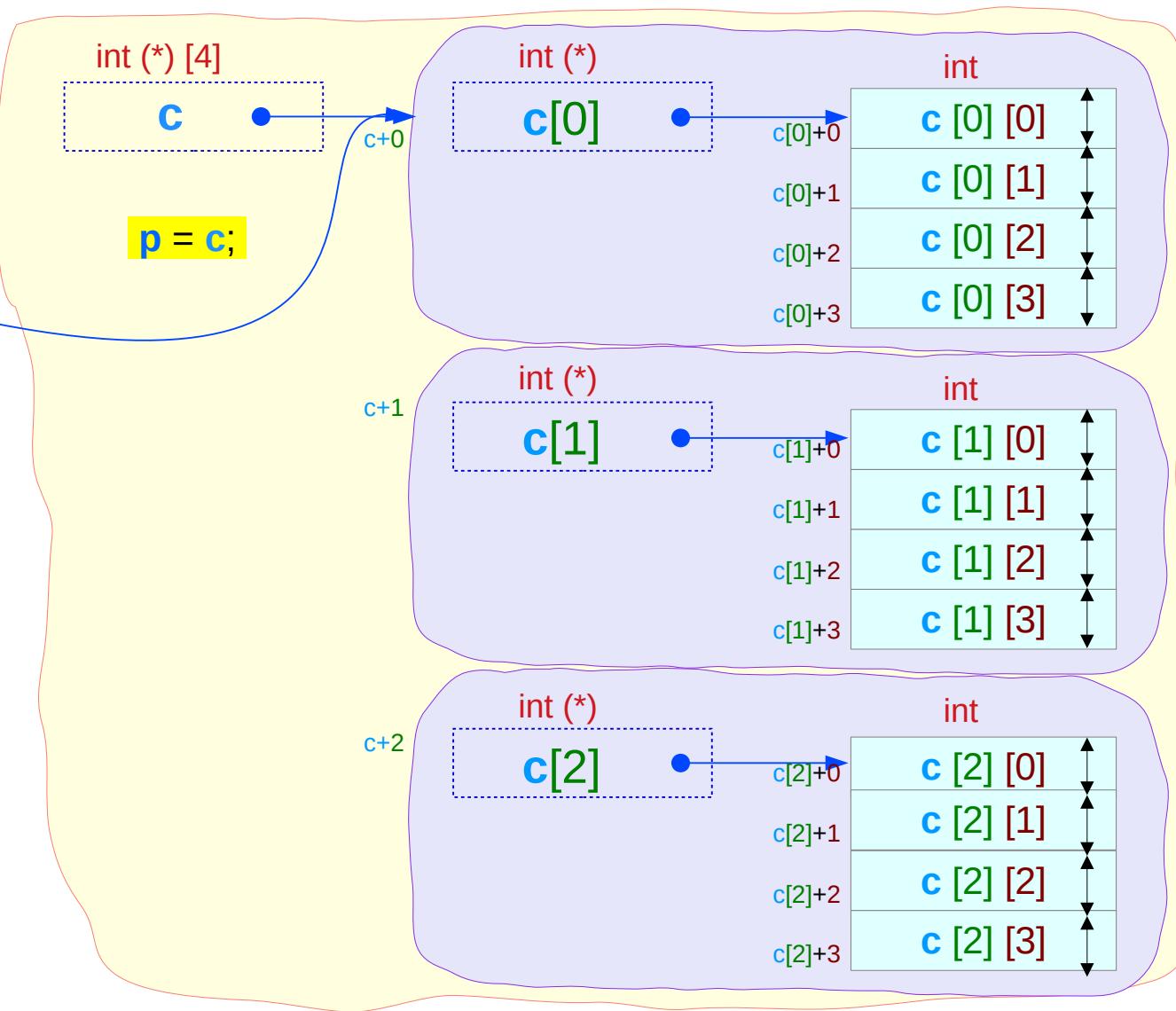
```
int (*p) [4];
```

```
int (*) [4] p;
```

```
int c[3] [4];
```

$v(c) = v(c[0]) = v(\&c[0][0])$
 $v(c[1]) = v(\&c[1][0])$
 $v(c[2]) = v(\&c[2][0])$

$v \equiv \text{value}$

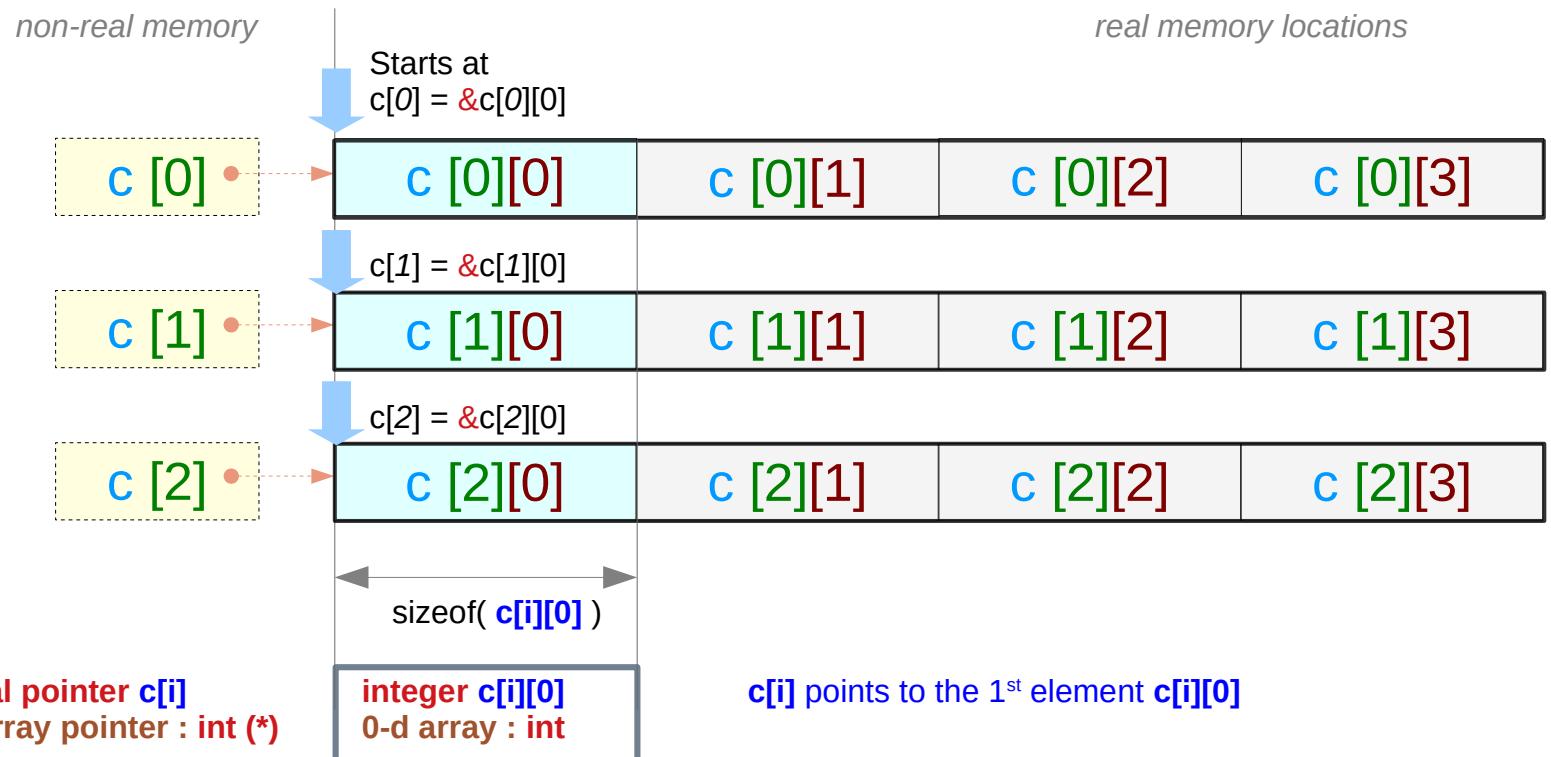


Pointer **c[i]** and integer **c[i][0]**

```
int c[3][4];
```

non-real pointer **c[i]** : value(**c[i]**) = &**c[i][0]**

0-d array pointer

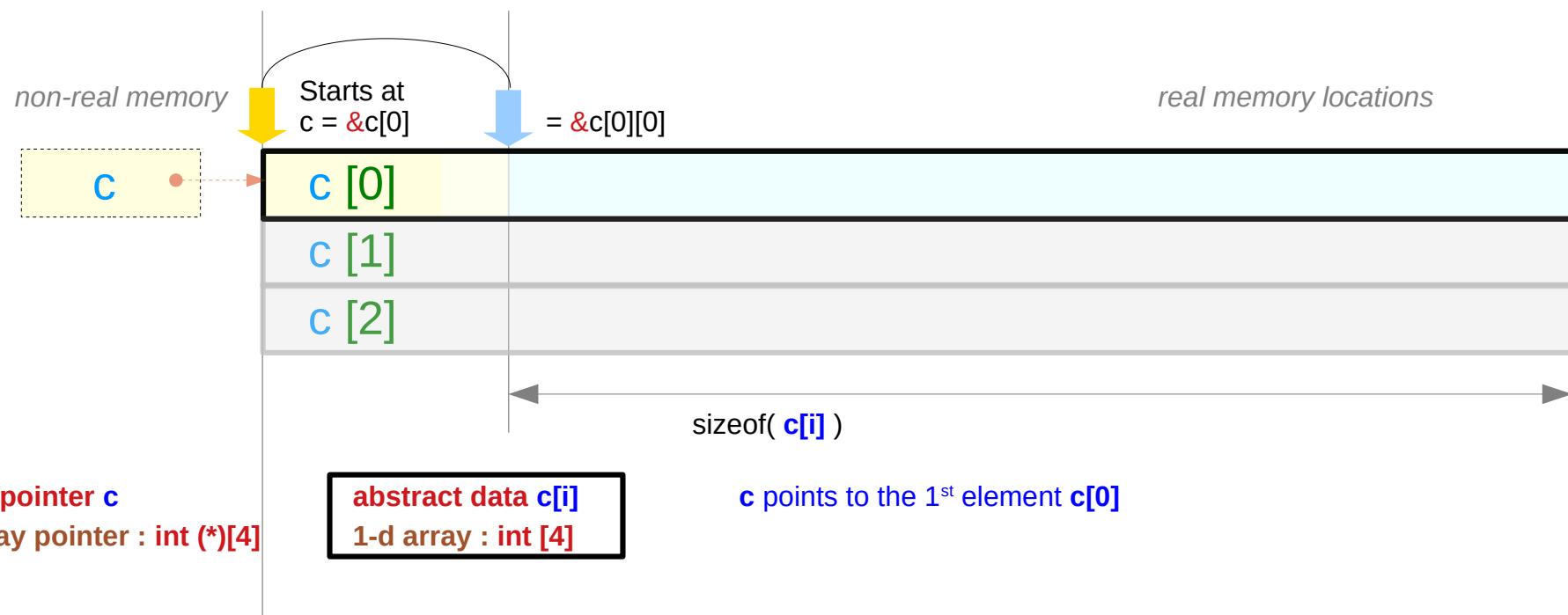


Pointer **c** and abstract data **c[i]**

```
int c [3] [4];
```

non-real pointer **c** : value(**c**) = $\&c[0] = \&c[0][0]$
abstract data **c[i]** : sizeof(**c[i]**) = $4 * \text{sizeof}(\text{int})$

1-d array pointer
1-d array



Abstract data c

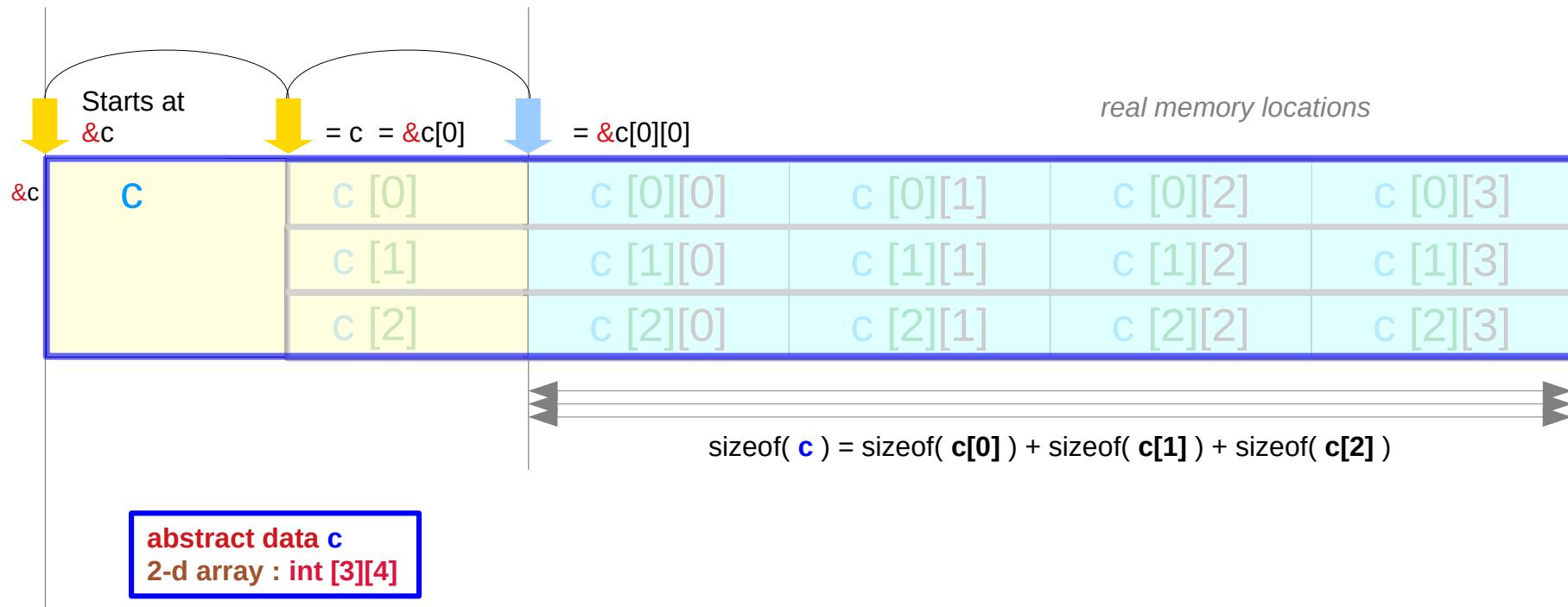
```
int c [3] [4];
```

abstract data

c:

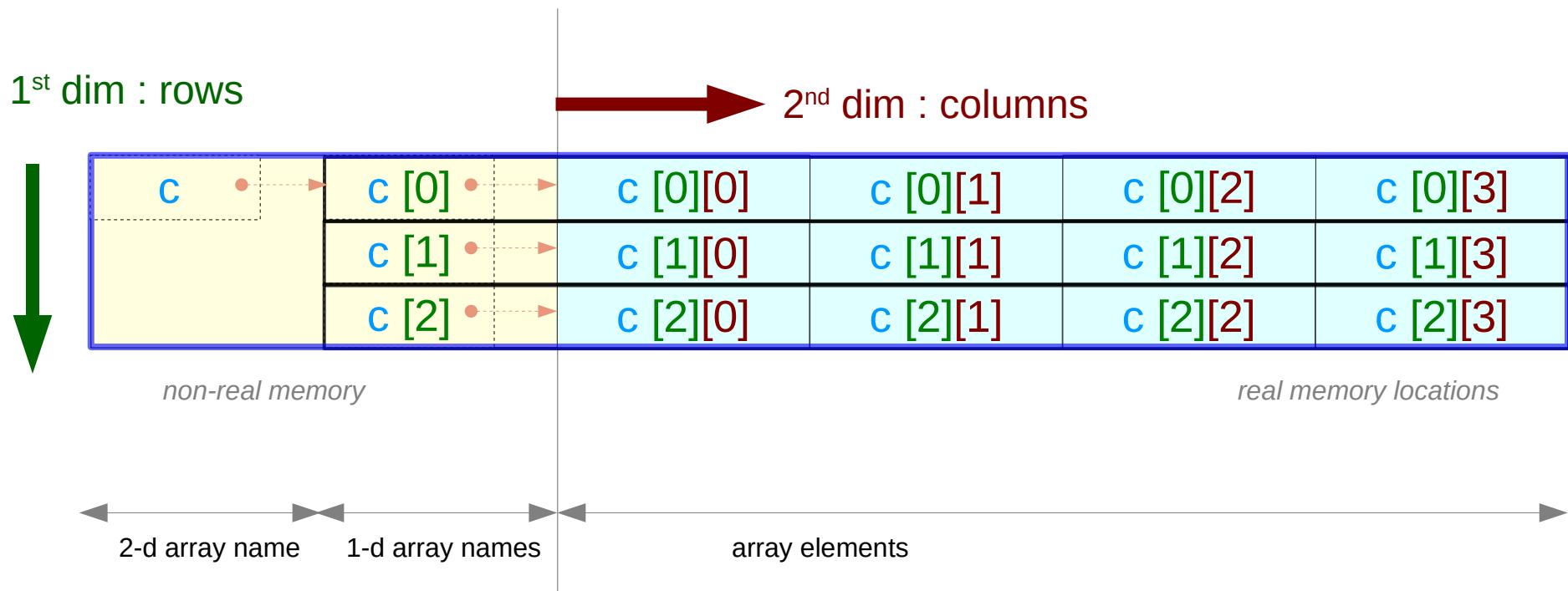
$\text{sizeof}(c) = 3 * \text{sizeof}(c[i])$

2-d array



Rows and columns of a 2-d array c

```
int c[3][4];
```



Determining types of sub-arrays from the declaration of an array

Types of array names

```
int [ ] a [4] ;
```

a is the name of the 1-d array

int [4]

$$\text{sizeof}(a) = 4 * 4$$

[3] is declared;
[0], [1], [2] are used

```
int [ ] c [3] [4] ;
```

c[i] is the name of the 1-d subarray

int [4]

$$\text{sizeof}(c[i]) = 4 * 4$$

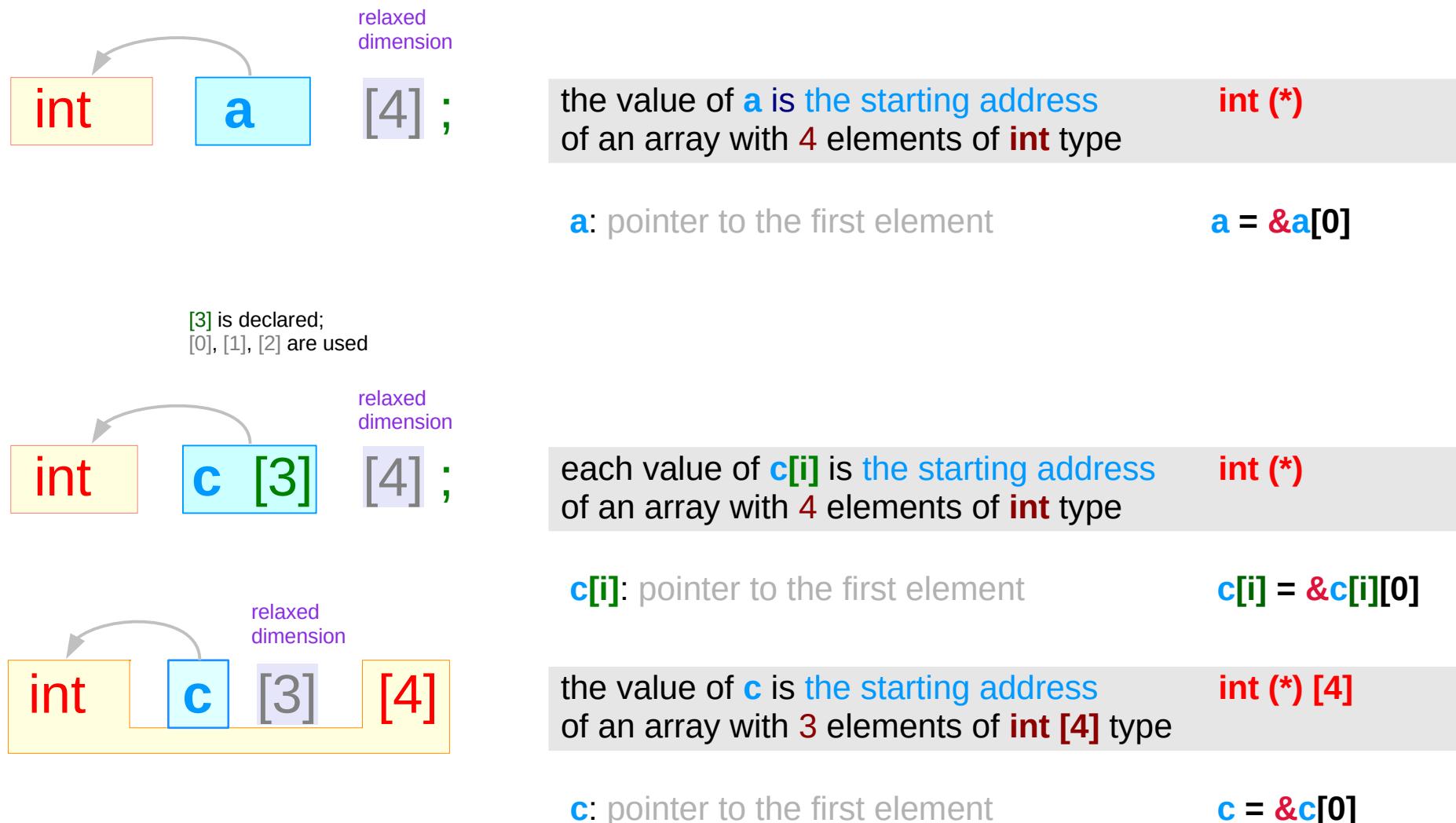
```
int [ ] c [3] [4]
```

c is the name of the 2-d array

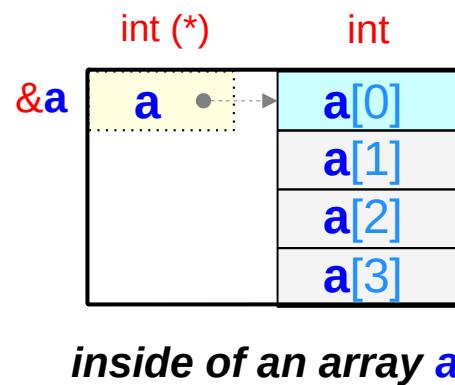
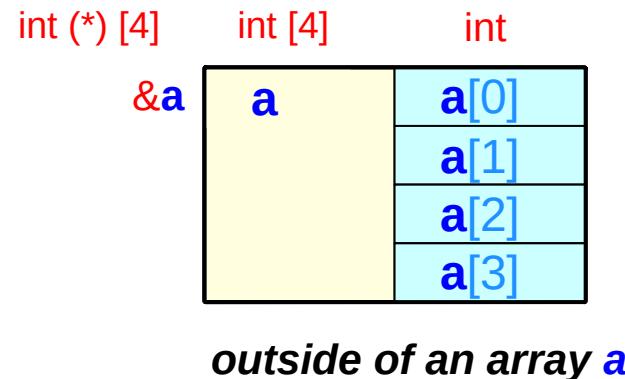
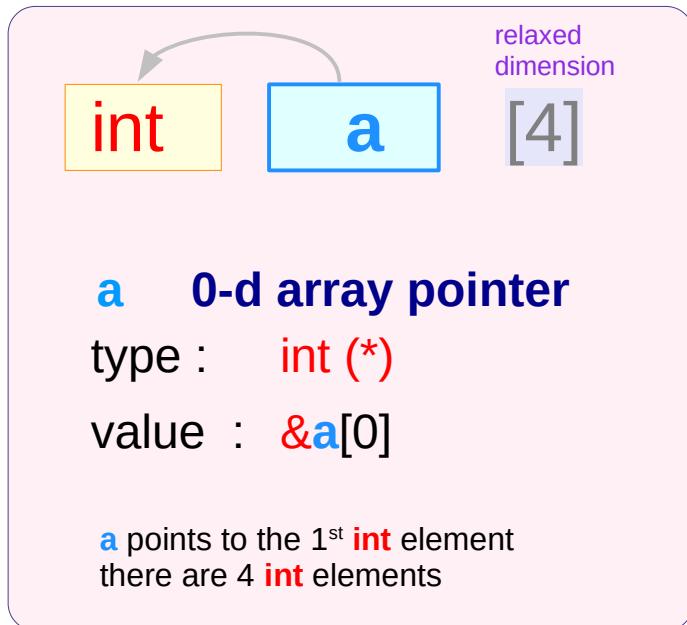
int [3][4]

$$\text{sizeof}(c) = 3 * 4 * 4$$

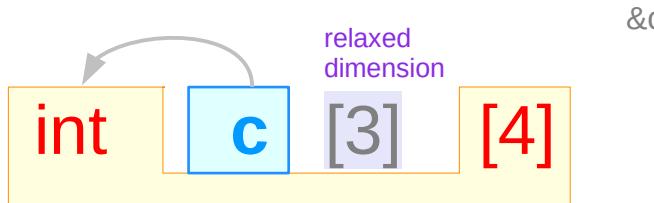
Values of array names



Array and pointer types in a 1-d array



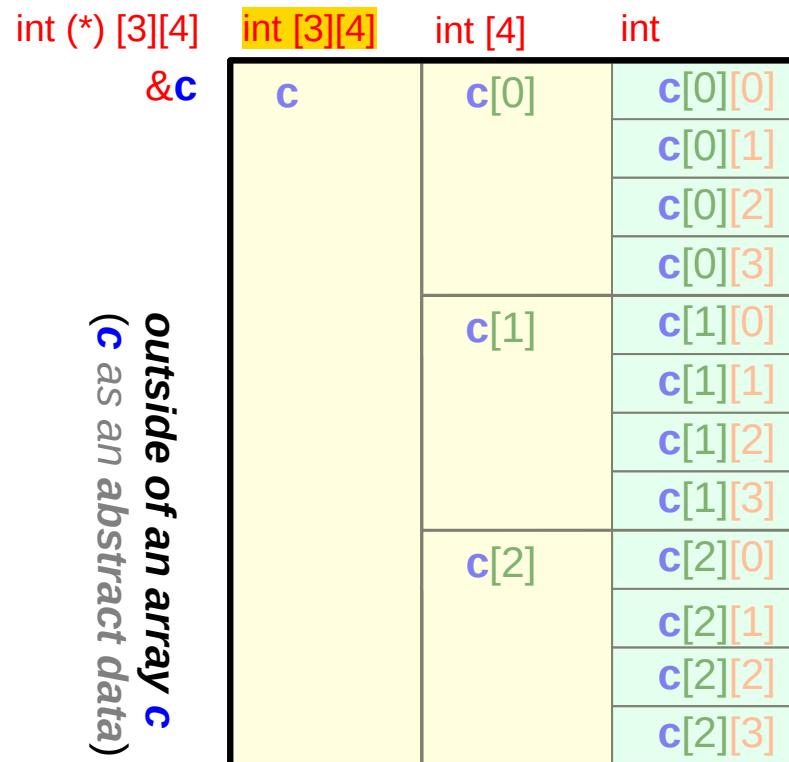
2-d array type



c 1-d array pointer

type : int (*) [4]

value : **c** = &**c**[0][0]



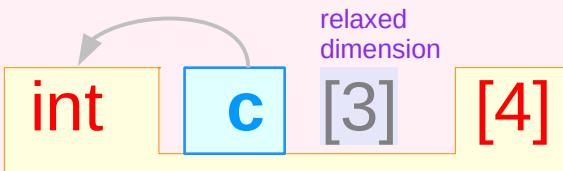
1-d array pointer type



c 2-d array

type : int [3][4]

size : 3 * 4 * 4

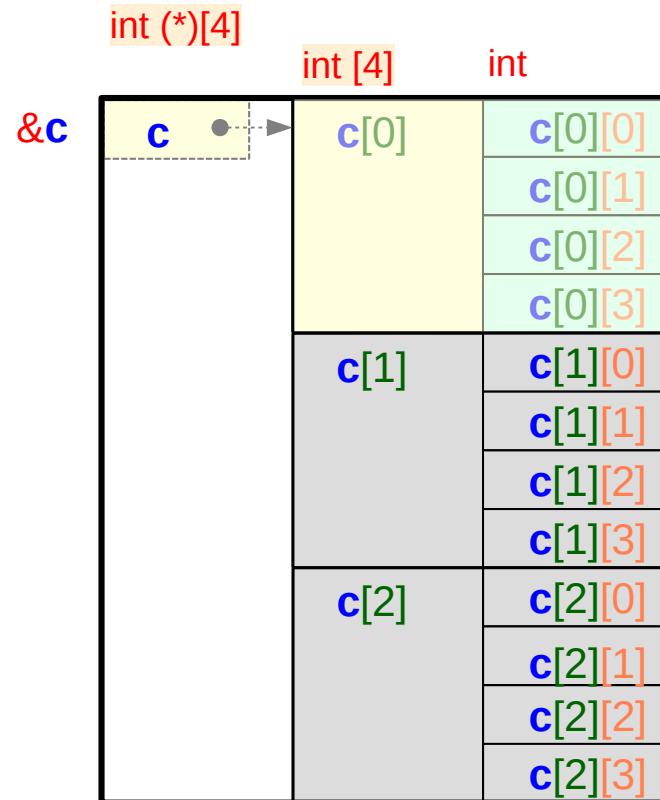


c 1-d array pointer

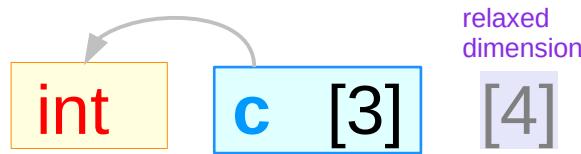
type : int (*) [4]

value : c = &c[0][0]

c points to the 1st int [4] element
There are 3 int [4] elements



1-d array type

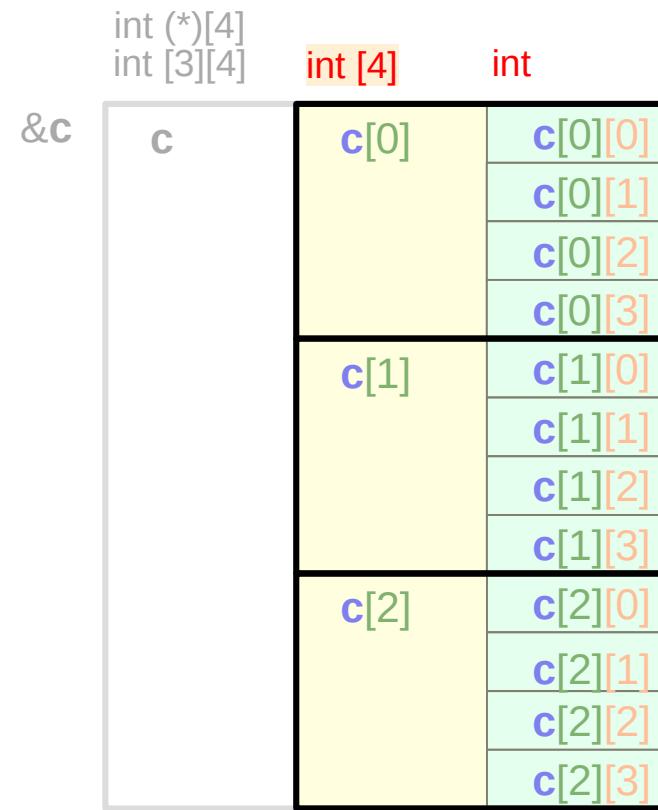


`c[i]` 0-d array pointer

type : `int (*)`

value : `c[i] = &c[i][0]`

`c[i]` points to the 1st `int` element
There are 4 `int` elements



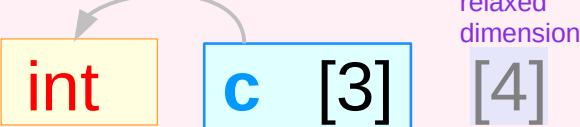
0-d array pointer type



`c[i]` 1-d array

type : `int [4]`

size : `4 * 4`

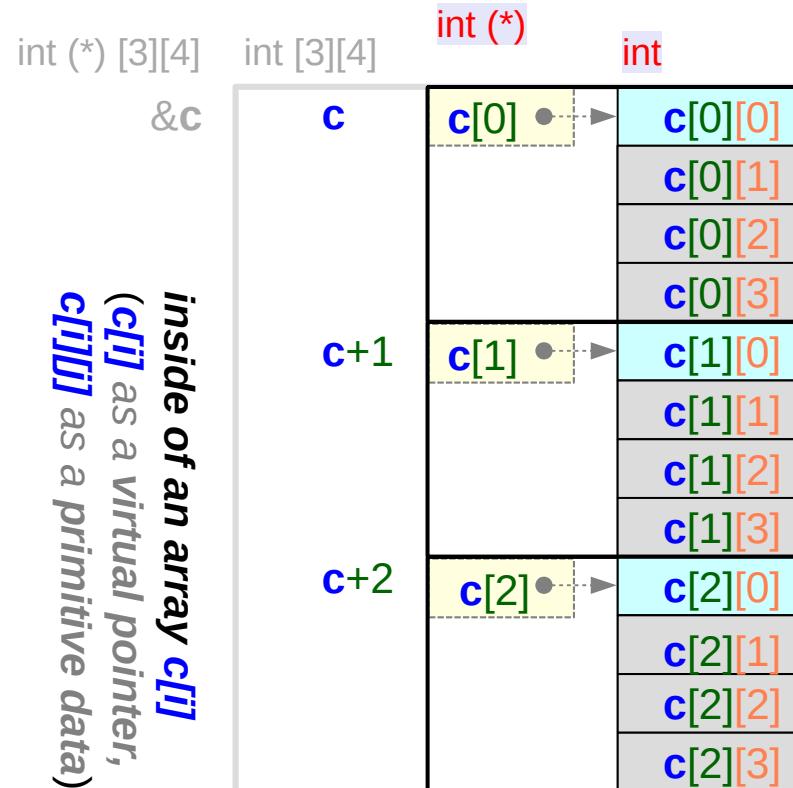


`c[i]` 0-d array pointer

type : `int (*)`

value : `c[i] = &c[i][0]`

`c[i]` points to the 1st `int` element
There are 4 `int` elements



*inside of an array `c[i]`
(`c[i]` as a virtual pointer,
`c[i][j]` as a primitive data)*

Types in a 2-d array



c 2-d array

type : int [3][4]
size : $3 * 4 * 4$



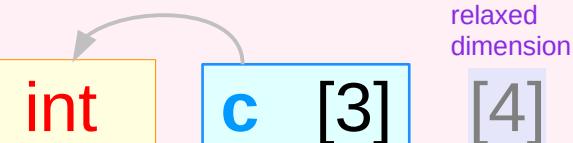
c 1-d array pointer (virtual)

type : int (*)[4]
value : &c[0][0]



c[i] 1-d array

type : int [4]
size : $4 * 4$



c[i] 0-d array pointer (virtual)

type : int (*)
value : &c[i][0]

The name of a 2-d array

```
int      a [4];
```

```
int      c [4] [4];
```

1. the name of the nested array (recursive definition)

2. a double pointer

3. a pointer to an array

2-d array **c** and 1-d array **q**

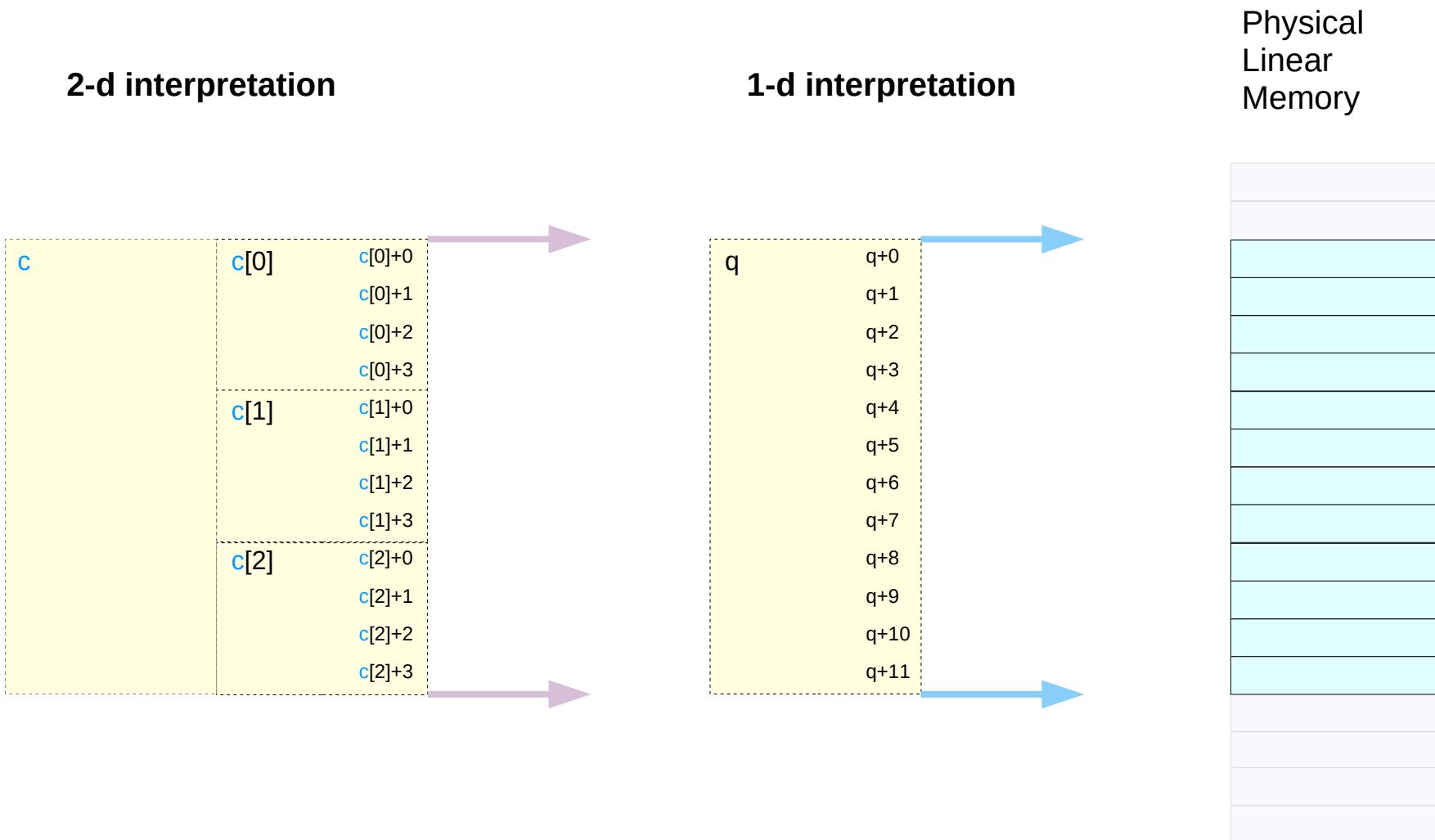
int c [3] [4];

c	c[0]	c[0]+0	c[0][0]
		c[0]+1	c[0][1]
		c[0]+2	c[0][2]
		c[0]+3	c[0][3]
	c[1]	c[1]+0	c[1][0]
		c[1]+1	c[1][1]
		c[1]+2	c[1][2]
		c[1]+3	c[1][3]
	c[2]	c[2]+0	c[2][0]
		c[2]+1	c[2][1]
		c[2]+2	c[2][2]
		c[2]+3	c[2][3]

int q [3*4];

q	q+0	q[0*4+0]
	q+1	q[0*4+1]
	q+2	q[0*4+2]
	q+3	q[0*4+3]
	q+4	q[1*4+0]
	q+5	q[1*4+1]
	q+6	q[1*4+2]
	q+7	q[1*4+3]
	q+8	q[2*4+0]
	q+9	q[2*4+1]
	q+10	q[2*4+2]
	q+11	q[2*4+3]

2-d and 1-d interpretations of linear memories



A 2-d array stored as a 1-d array (row major order)

int c [3] [4];

c[i][j]

[i*4+j]

[k]

c	c[0]	c[0]+0	c[0][0]
		c[0]+1	c[0][1]
		c[0]+2	c[0][2]
		c[0]+3	c[0][3]
	c[1]	c[1]+0	c[1][0]
		c[1]+1	c[1][1]
		c[1]+2	c[1][2]
		c[1]+3	c[1][3]
	c[2]	c[2]+0	c[2][0]
		c[2]+1	c[2][1]
		c[2]+2	c[2][2]
		c[2]+3	c[2][3]

index values

0	=[0*4+0]
1	=[0*4+1]
2	=[0*4+2]
3	=[0*4+3]
4	=[1*4+0]
5	=[1*4+1]
6	=[1*4+2]
7	=[1*4+3]
8	=[2*4+0]
9	=[2*4+1]
10	=[2*4+2]
11	=[2*4+3]

q	q+0	q[0]
	q+1	q[1]
	q+2	q[2]
	q+3	q[3]
	q+4	q[4]
	q+5	q[5]
	q+6	q[6]
	q+7	q[7]
	q+8	q[8]
	q+9	q[9]
	q+10	q[10]
	q+11	q[11]

2-d array access via a single pointer

```
int *p = c[0];
```

```
int c [3][4];
```

```
p[ i*4 + j ]
```

```
c[ i ][ j ]
```

```
*(p+ i*4 + j)
```

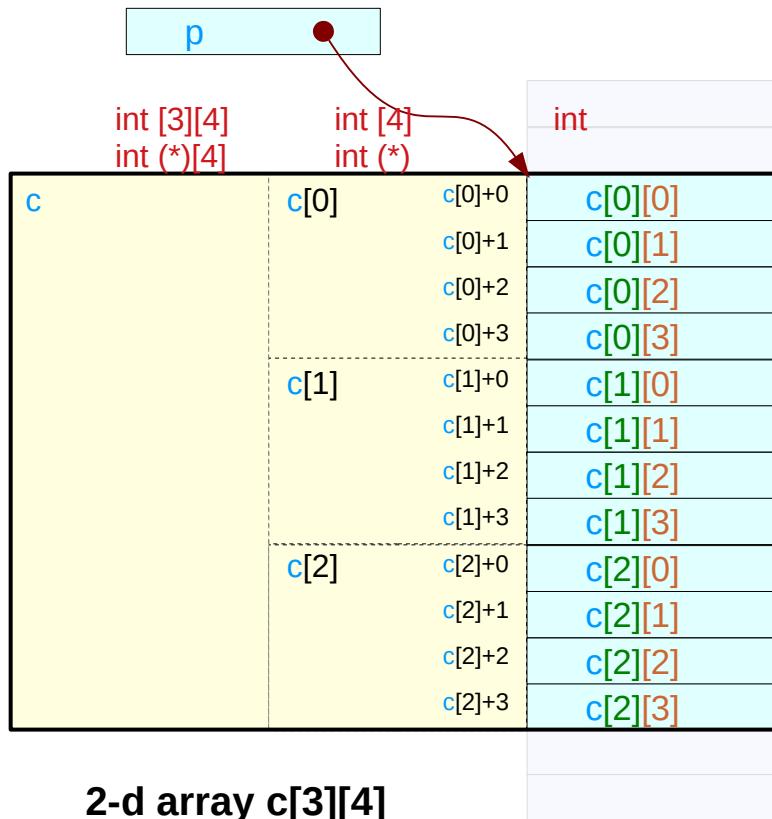
```
*(*(c+i)+ j)
```

```
*(p+k)      i = k / 4;  
              j = k % 4;
```

View a 2-d array as a 1-d array

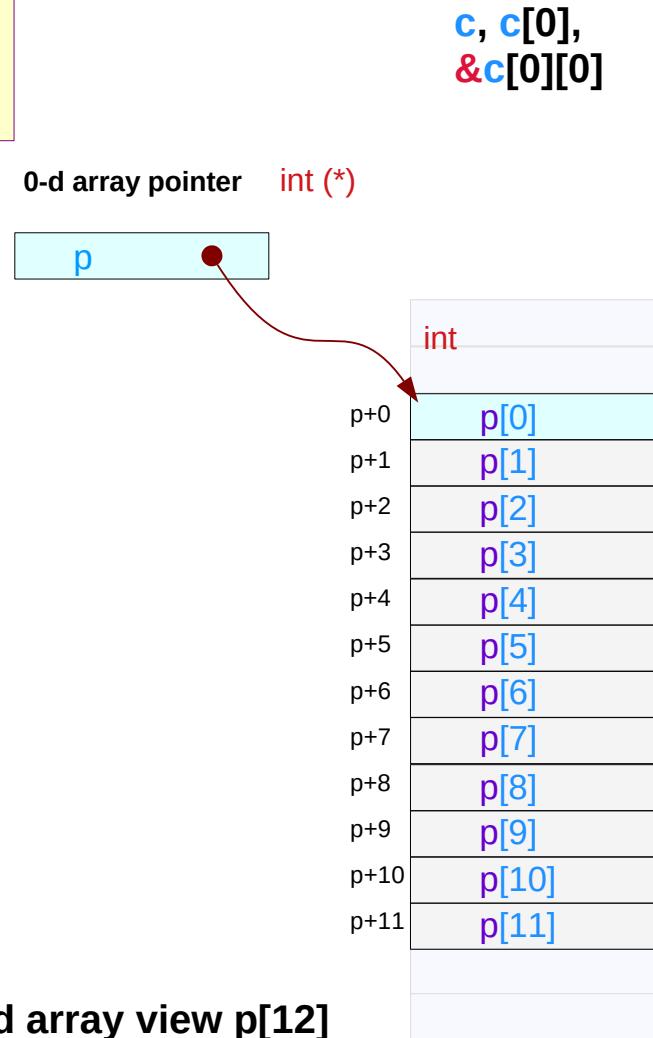
`int c [3][4];`

0-d array pointer `int (*)`



`int *p = c[0];`

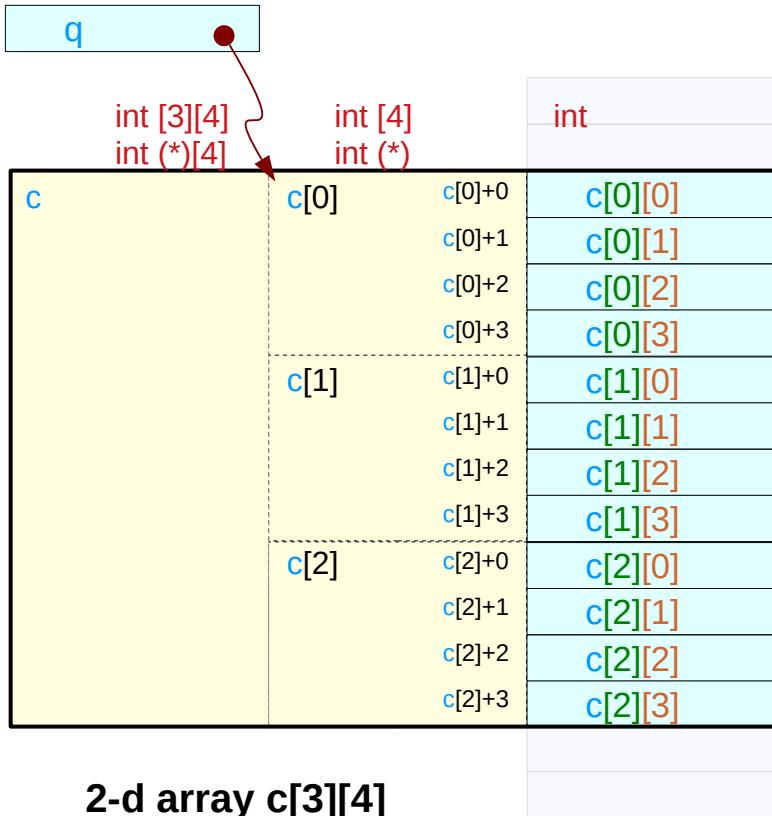
0-d array pointer `int (*)`



View a 2-d array as another 2-d array

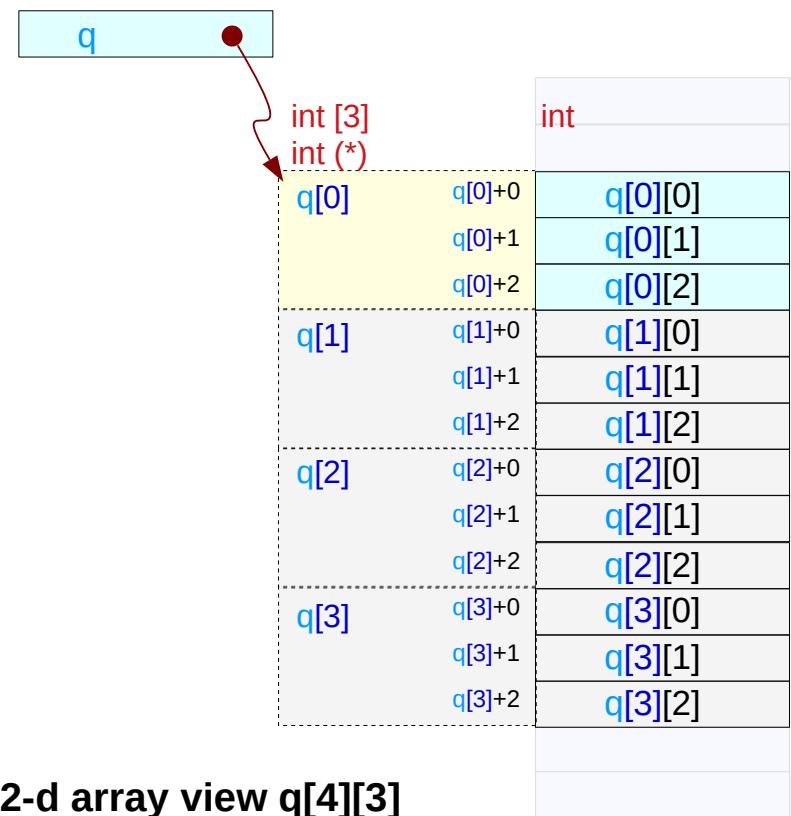
`int c [3][4];`

1-d array pointer `int (*) [3]`



`int (*q) [3] = (int (*) [3]) c;`

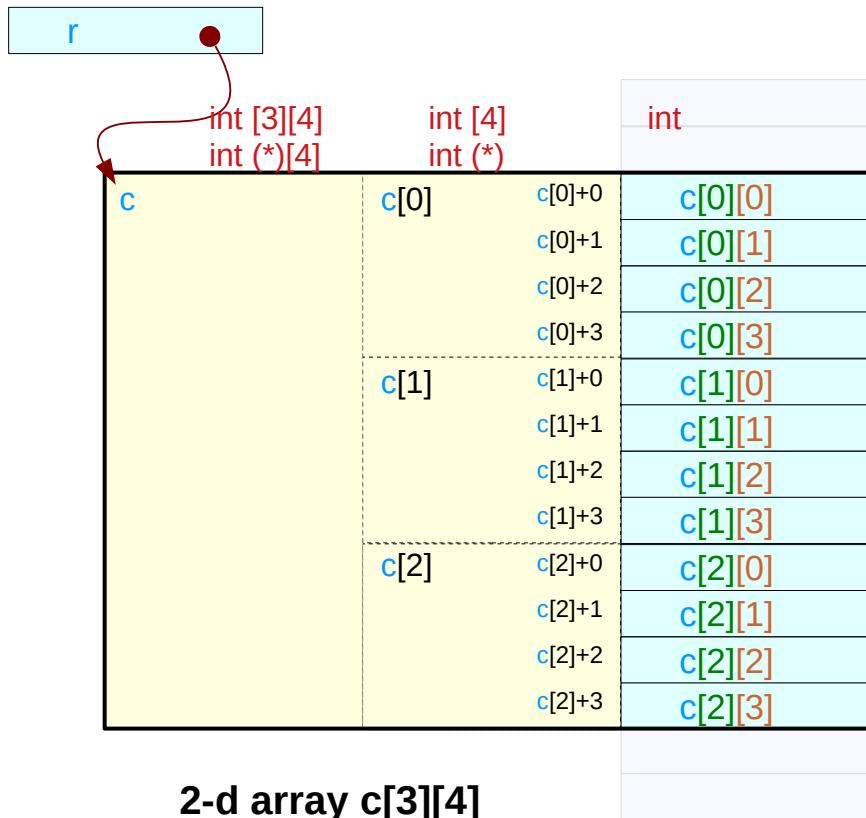
1-d array pointer `int (*) [3]`



A 2-d array stored as a 1-d array (row major order)

`int c [3] [4];`

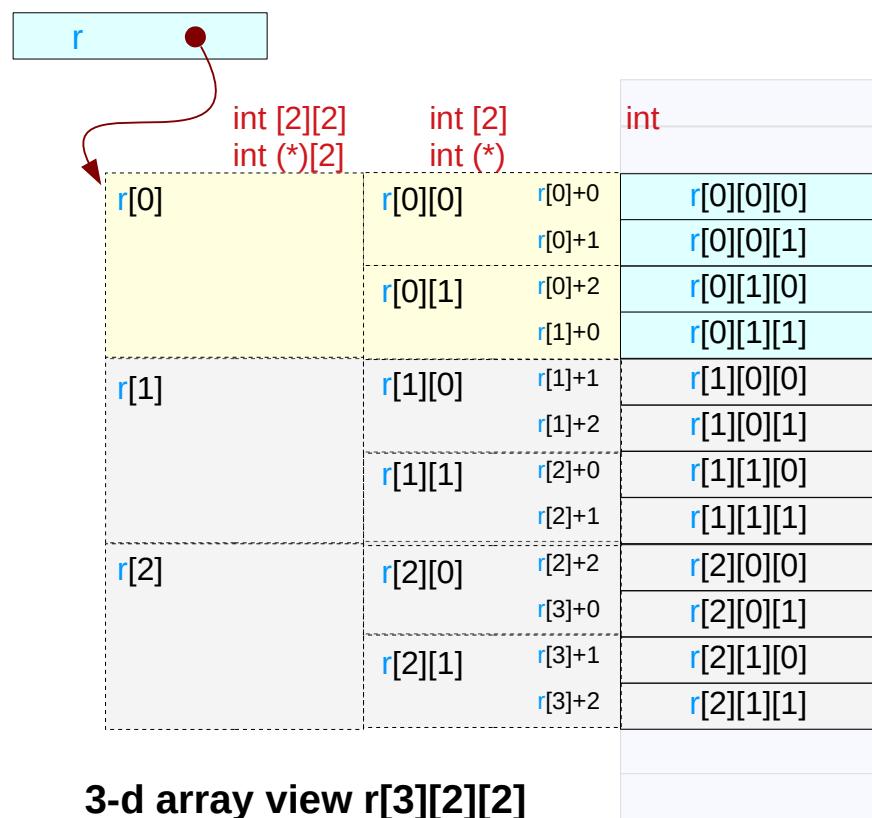
2-d array pointer `int (*) [2][2]`



`int (*r) [2][2] = (int (*) [2][2]) c;`

`c, c[0],
&c[0][0]`

2-d array pointer `int (*) [2][2]`



2-d array access via pointers

```
int c [3][4];
```

```
int *p = c[0] ;
```

1. recursive pointers

c [i][j]

(*c+i)[j]

*(c[i]+ j)

*(*c+i)+ j)

int (*p)[4];

p[i*4 + j]

*(p+ i*4 + j)

2. linear array pointers

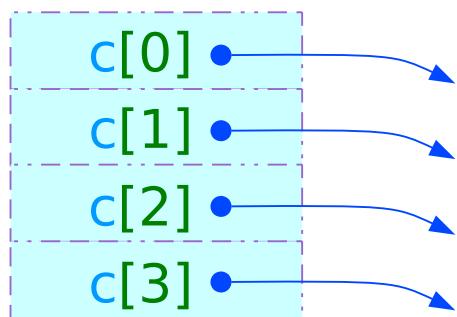
Static Allocation of a 2-d Array

int A [3][4];

A in %eax,
i in %edx,
j in %ecx

```
sall $2, %ecx  
leal (%edx, %edx, 2), %edx  
leal (%ecx, %edx, 4), %edx  
movl (%eax, %edx), %eax
```

```
; j * 4  
; i * 3  
; j * 4 + i * 12  
; read M[ X_A+4(3i +j) ]
```

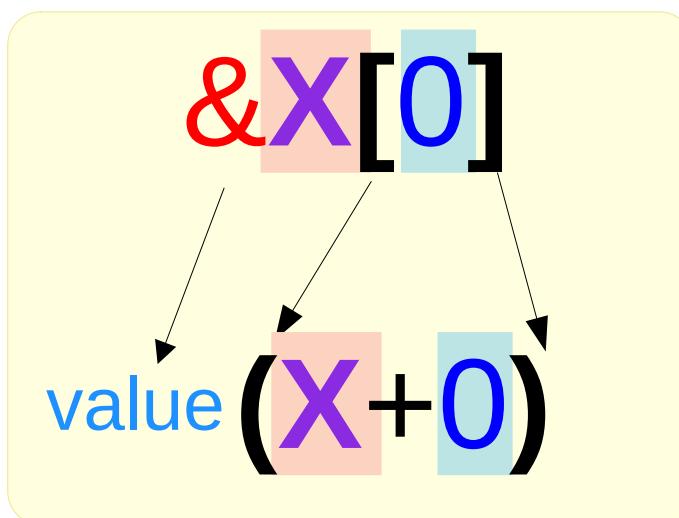
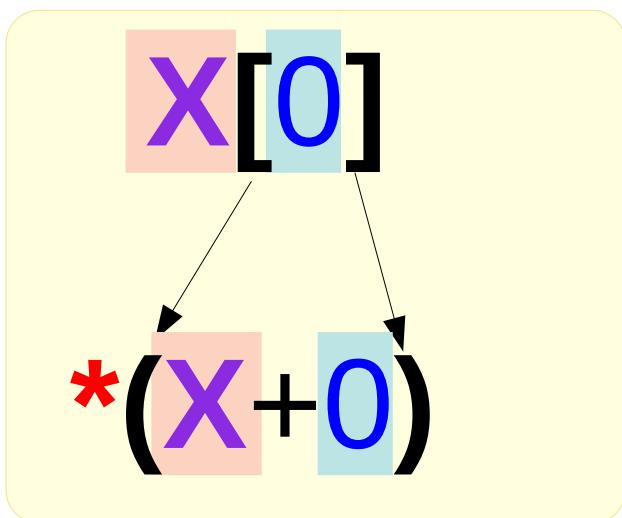
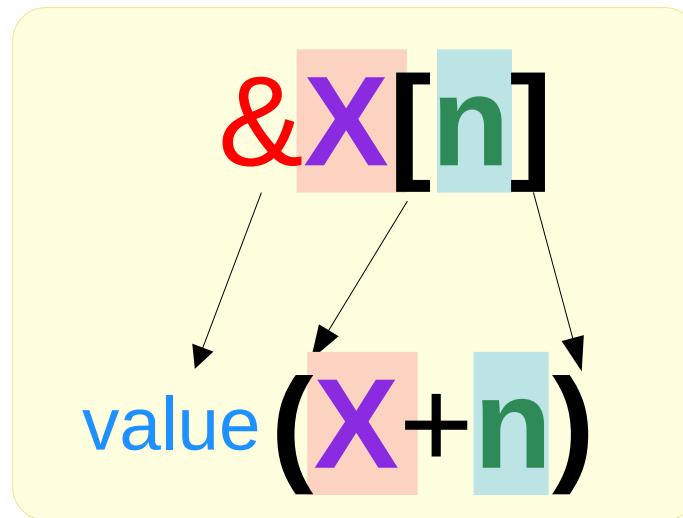
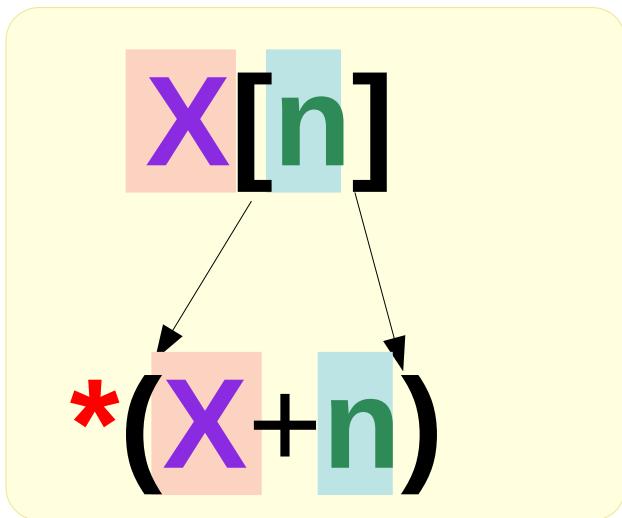


The pointer array :
not allocated
in the memory

<code>c[0]+0</code>	<code>*(c [0]+0)</code>
<code>c[0]+1</code>	<code>*(c [0]+1)</code>
<code>c[0]+2</code>	<code>*(c [0]+2)</code>
<code>c[0]+3</code>	<code>*(c [0]+3)</code>
<code>c[1]+0</code>	<code>*(c [1]+0)</code>
<code>c[1]+1</code>	<code>*(c [1]+1)</code>
<code>c[1]+2</code>	<code>*(c [1]+2)</code>
<code>c[1]+3</code>	<code>*(c [1]+3)</code>
<code>c[2]+0</code>	<code>*(c [2]+0)</code>
<code>c[2]+1</code>	<code>*(c [2]+1)</code>
<code>c[2]+2</code>	<code>*(c [2]+2)</code>
<code>c[2]+3</code>	<code>*(c [2]+3)</code>

Pointers, arrays, and operator precedence

Equivalences between *, &, and [] operators



Operator Precedence of * and []

$$*\mathbf{x}[m] \equiv *(\mathbf{x}[m])$$

$$\mathbf{x}[m][n] \equiv (\mathbf{x}[m])[n]$$

$$\mathbf{**x} \equiv *(*\mathbf{x})$$

[] has a **higher** priority than *

[] has **left-to-right** associativity

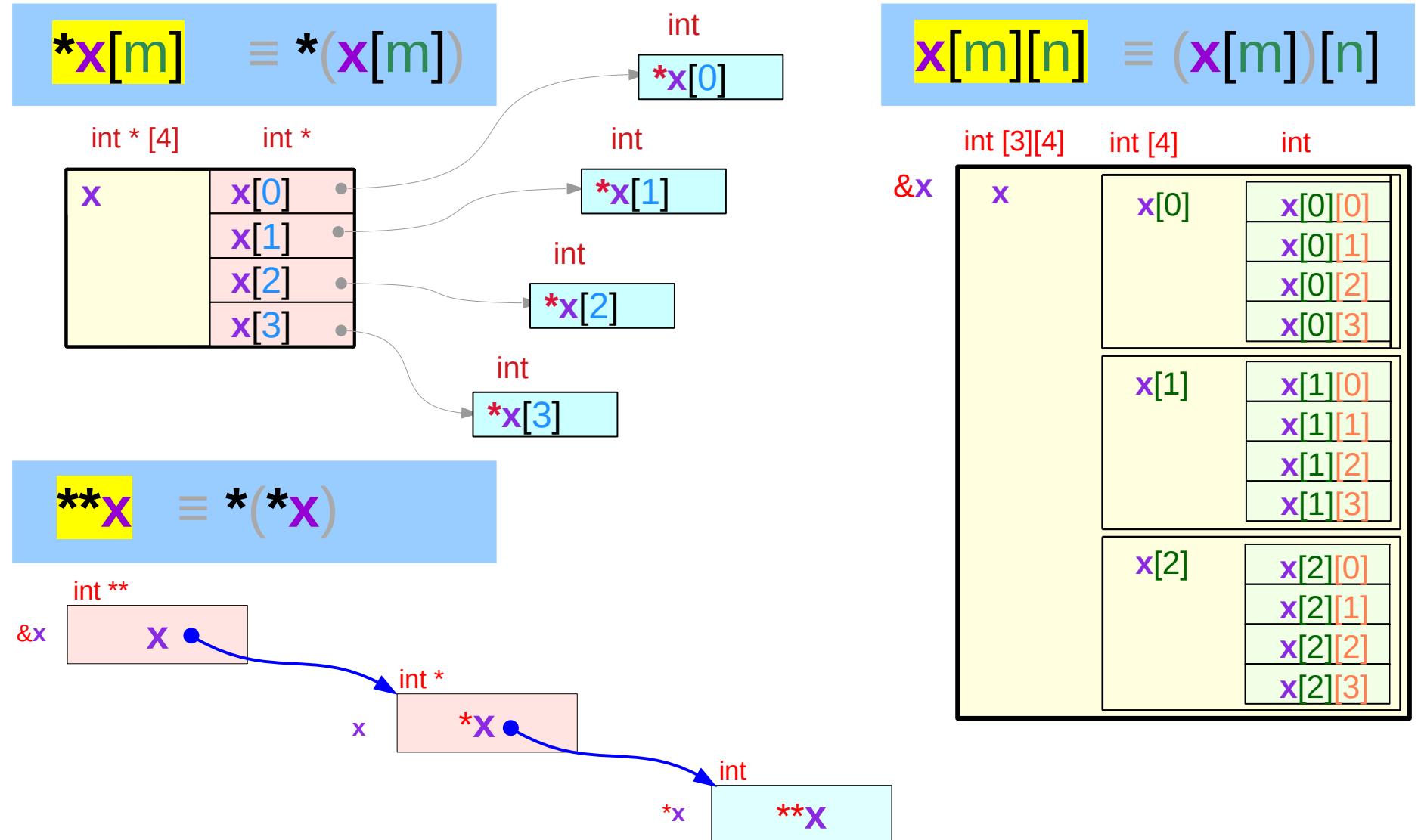
* has **right-to-left** associativity

$$(*\mathbf{x})[m][n] \leftrightarrow ((*\mathbf{x})[m])[n]$$

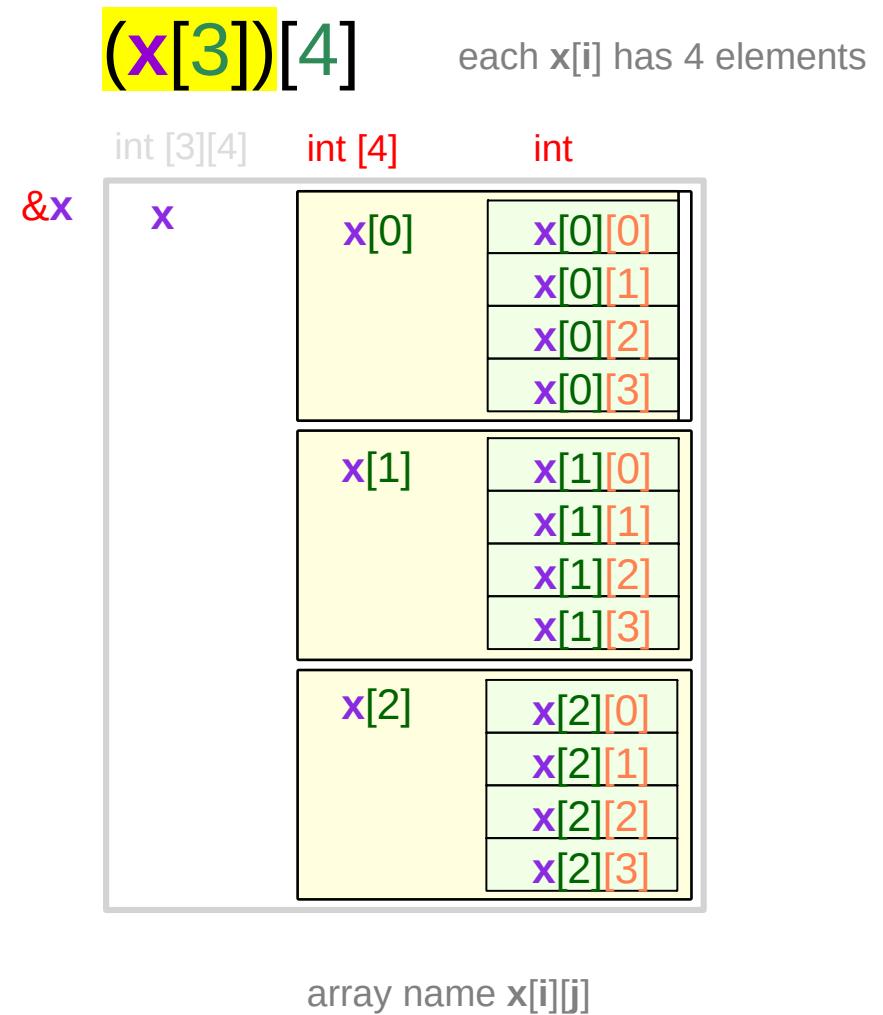
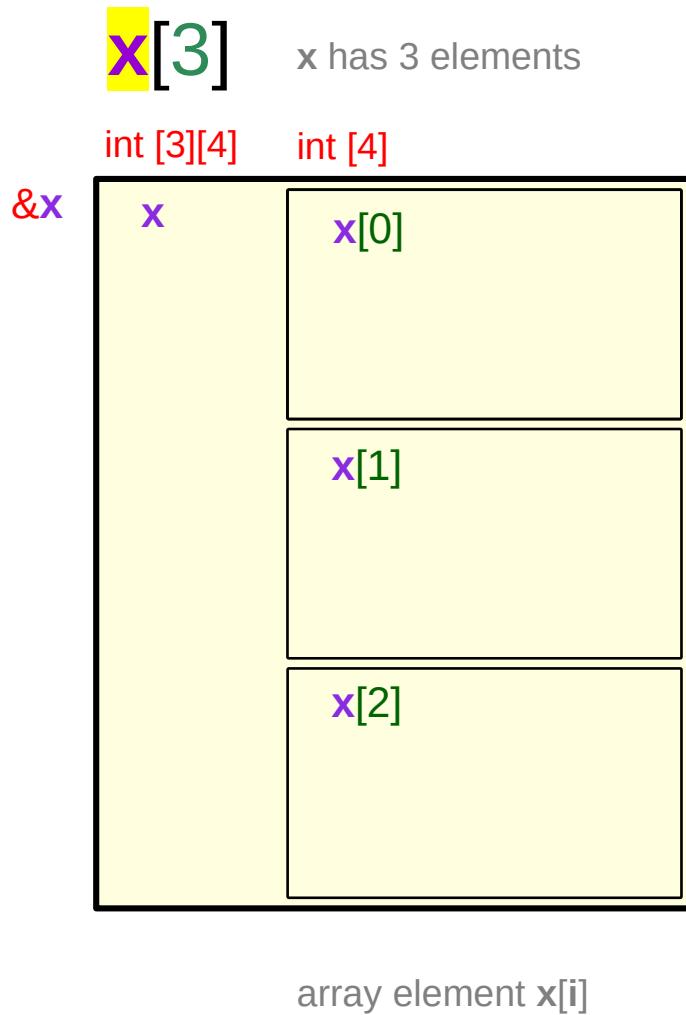
red parentheses () must not be removed
gray parentheses () can be removed

$$(*\mathbf{x}[m])[n] \leftrightarrow (*(\mathbf{x}[m]))[n]$$

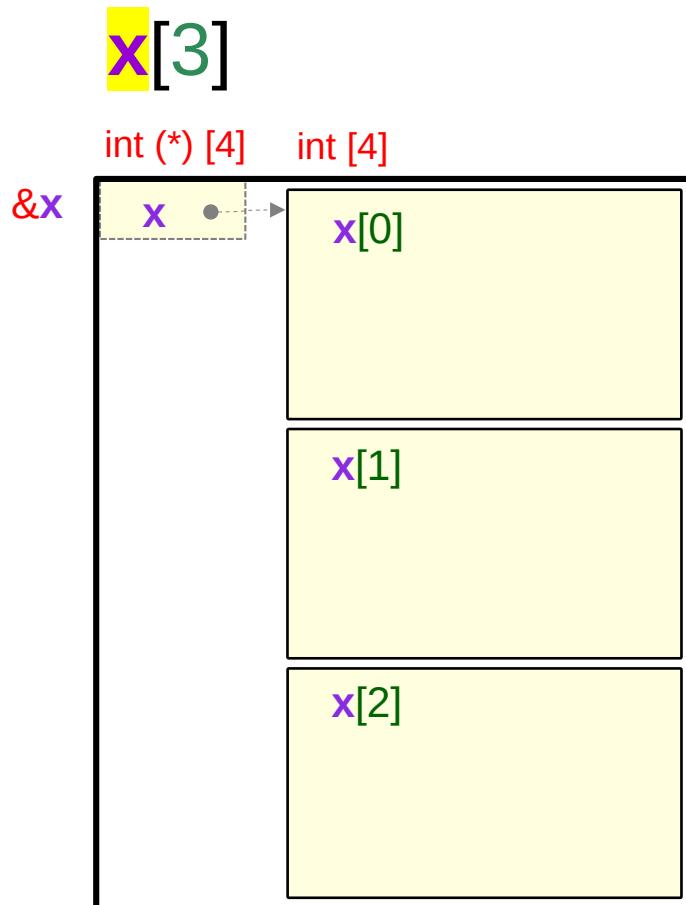
Operator Precedence of * and []



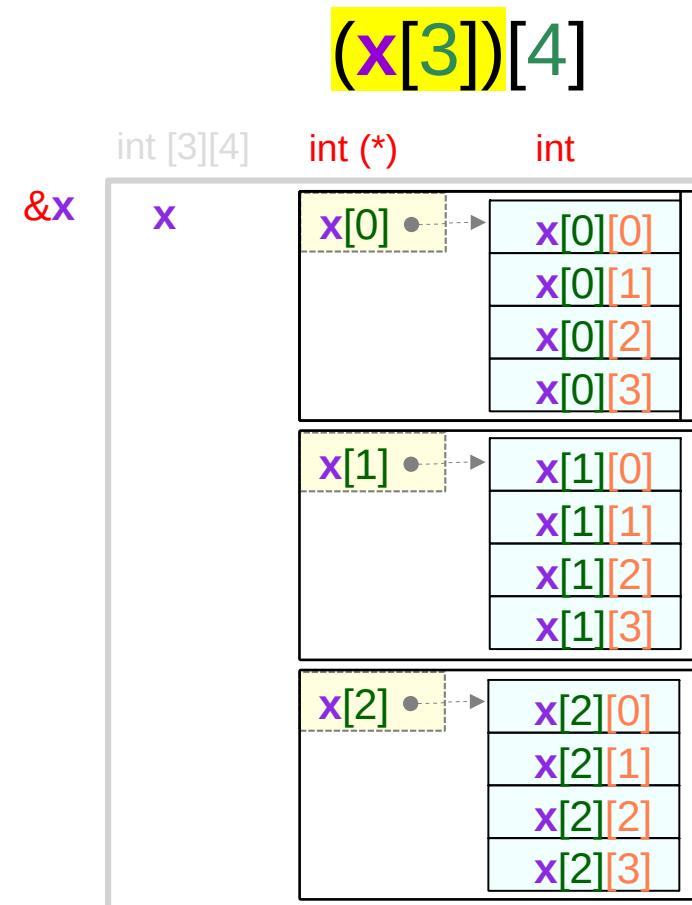
Abstract Data x and $x[i]$



Virtual Pointers x and $x[i]$



array name x virtual pointer
array element $x[i]$ abstract data



array name $x[i]$ virtual pointer
array element $x[i][j]$ primitive data

Left-to-right and right-to-left associative operators

$$\begin{aligned} p[i] &\equiv p[\cancel{i}] \\ p[i][j] &\equiv (\cancel{p[i]}[j])[\cancel{j}] \\ p[i][j][k] &\equiv ((\cancel{p[i]}[\cancel{j}])[k])[\cancel{k}] \end{aligned}$$

$$\begin{aligned} &\rightarrow *(\cancel{p+i}) \\ &\rightarrow *\cancel{(*(p+i)+j)} \\ &\rightarrow *\cancel{(*(*(\cancel{p+i})+j)+k)} \end{aligned}$$

$$\begin{aligned} *p &\equiv *(p) \\ **p &\equiv *\cancel{(*(p))} \\ ***p &\equiv *\cancel{(*(*(\cancel{p})))} \end{aligned}$$

$$\begin{aligned} &\rightarrow p[\cancel{0}] \\ &\rightarrow (\cancel{p[0]}[\cancel{0}])[\cancel{0}] \\ &\rightarrow ((\cancel{p[0]}[\cancel{0}])[\cancel{0}])[\cancel{0}] \end{aligned}$$

Equivalence on relaxing outer most dimension

$$p[i] \equiv *(\&p + i)$$

$$p[i][j] \equiv *(p[i] + j)$$

$$p[i][j][k] \equiv *(p[i][j] + k)$$

$$\&p[i] \equiv \text{value}(\&p + i)$$

$$\&p[i][j] \equiv \text{value}(p[i] + j)$$

$$\&p[i][j][k] \equiv \text{value}(p[i][j] + k)$$

$$p[0] \equiv *p$$

$$p[i][0] \equiv *p[i]$$

$$p[i][j][0] \equiv *p[i][j]$$

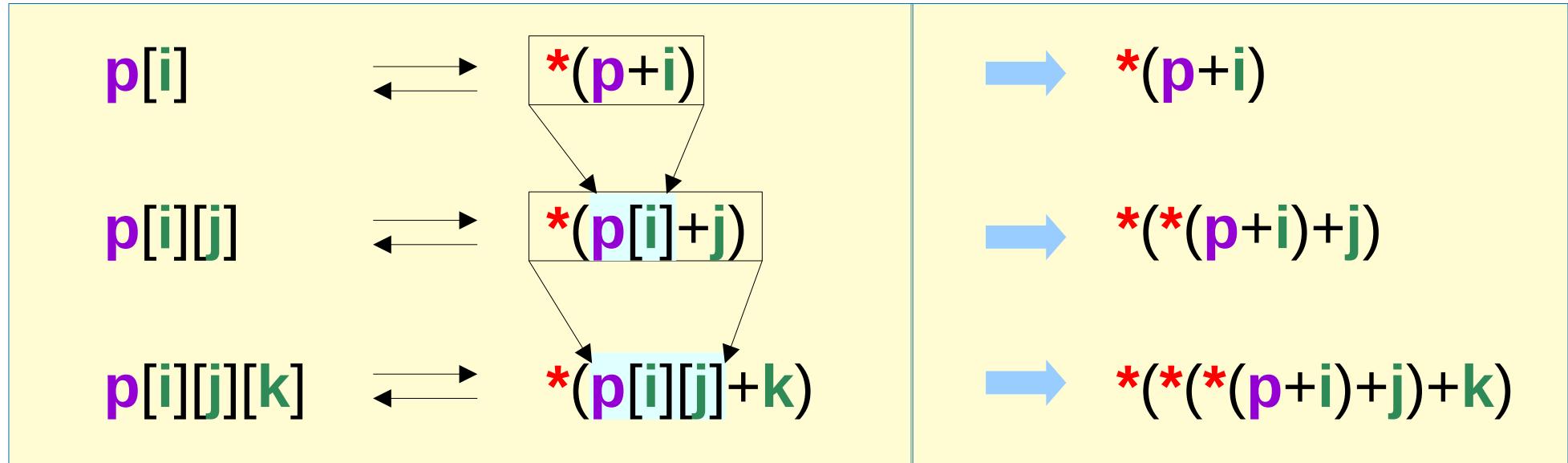
$$\&p[0] \equiv \text{value}(p)$$

$$\&p[i][0] \equiv \text{value}(p[i])$$

$$\&p[i][j][0] \equiv \text{value}(p[i][j])$$

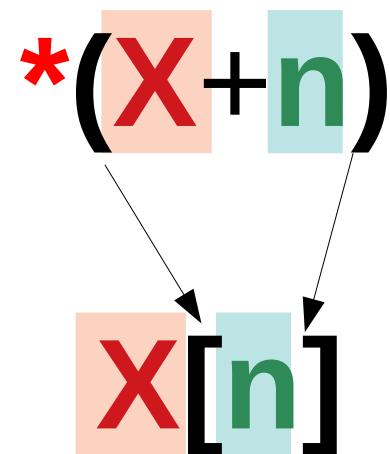
valid for proper i, j, k values

Relaxing all the dimensions



$$\ast(X+n) \equiv X[n]$$

valid for proper i, j, k values



Equivalences on relaxing all the dimensions

$$p[i] \equiv *(\mathbf{p} + i)$$

$$p[i][j] \equiv *(*(\mathbf{p} + i) + j)$$

$$p[i][j][k] \equiv *(*(*(\mathbf{p} + i) + j) + k)$$

$$\&p[i] \equiv \text{value}(\mathbf{p} + i)$$

$$\&p[i][j] \equiv \text{value}(*(\mathbf{p} + i) + j)$$

$$\&p[i][j][k] \equiv \text{value}(*(*(\mathbf{p} + i) + j) + k)$$

$$p[0] \equiv *p$$

$$p[0][0] \equiv **p$$

$$p[0][0][0] \equiv ***p$$

$$\&p[0] \equiv \text{value}(p)$$

$$\&p[0][0] \equiv \text{value}(*p)$$

$$\&p[0][0][0] \equiv \text{value}(**p)$$

valid for proper i, j, k values

Address Calculation (1) Array Pointer Approach

int **c** [2][3][4] ;

$$\begin{aligned} \mathbf{c}[i] &\equiv *(\mathbf{c} + i) \\ \mathbf{c}[i][j] &\equiv *(\mathbf{c}[i] + j) \\ \mathbf{c}[i][j][k] &\equiv *(\mathbf{c}[i][j] + k) \end{aligned}$$

$$\begin{aligned} \&\mathbf{c}[i] &\equiv \text{value}(\mathbf{c} + i) \\ \&\mathbf{c}[i][j] &\equiv \text{value}(\mathbf{c}[i] + j) \\ \&\mathbf{c}[i][j][k] &\equiv \text{value}(\mathbf{c}[i][j] + k) \end{aligned}$$

address replication

$\text{value}(\mathbf{c}[i][j][k]) \neq \text{value}(\&\mathbf{c}[i][j][k])$ ← primitive data & address

$\text{value}(\mathbf{c}[i][j]) = \text{value}(\&\mathbf{c}[i][j]) = \text{value}(\&\mathbf{c}[i][j][0])$

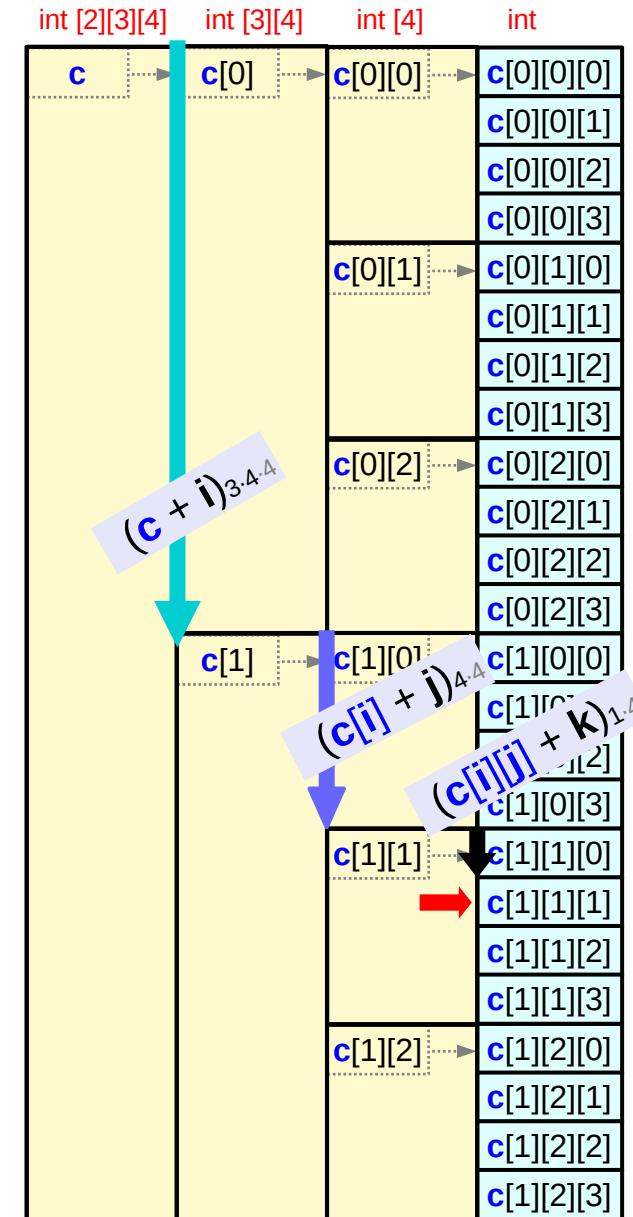
$\text{value}(\mathbf{c}[i]) = \text{value}(\&\mathbf{c}[i]) = \text{value}(\&\mathbf{c}[i][0][0])$

$\text{value}(\mathbf{c}) = \text{value}(\&\mathbf{c}) = \text{value}(\&\mathbf{c}[0][0][0])$

skip i elements of \mathbf{c}
skip j elements of $\mathbf{c}[i]$
skip k elements of $\mathbf{c}[i][j]$



skip $i*3*4$ primitive elements of \mathbf{c}
skip $j*4$ primitive elements of \mathbf{c}
skip k primitive elements of \mathbf{c}



Address Calculation (2) Pointer Array Approach

```
int ** c [2] ;
int * b [2*3] ;
int a [2*3*4] ;
```

$$\begin{aligned} b[j] &\equiv (a + j \cdot 4) \\ *(b[j]+k) &= *(a + j \cdot 4 + k); \\ b[j][k] &\equiv a[j \cdot 4 + k] \end{aligned}$$

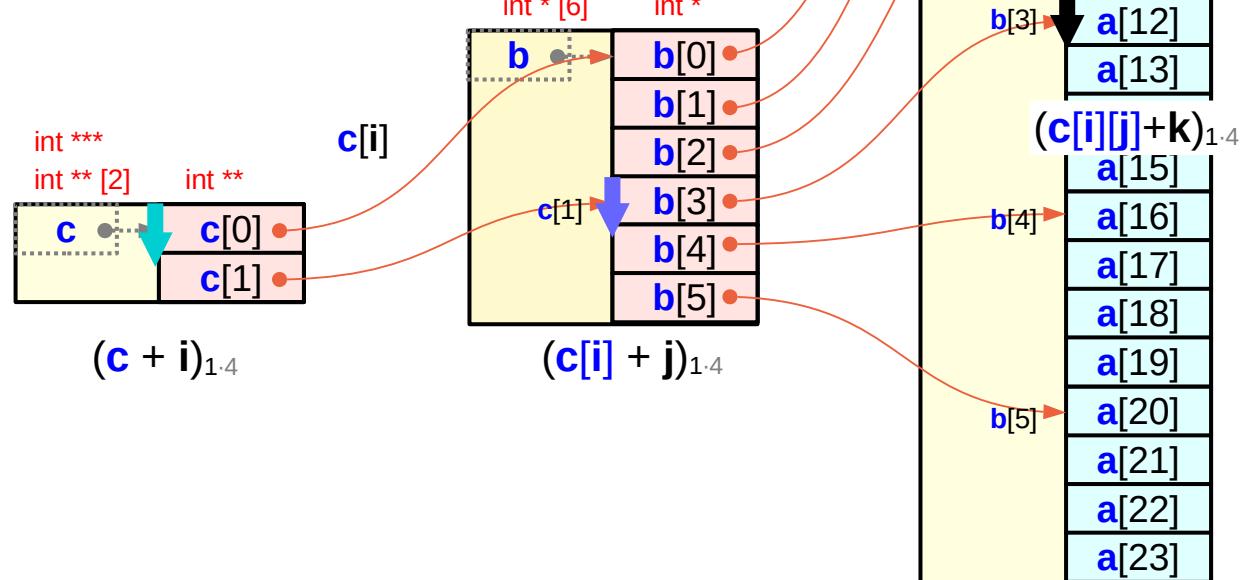
$$\begin{aligned} c[i] &\equiv (b + i \cdot 3) \\ *(c[i]+j) &= *(b + i \cdot 3 + j); \\ c[i][j] &\equiv b[i \cdot 3 + j] \end{aligned}$$

$$\begin{aligned} c[i][j] &\equiv (a + (i \cdot 3 + j) \cdot 4) \\ *(c[i][j]+k) &= *(a + (i \cdot 3 + j) \cdot 4 + k); \\ c[i][j][k] &\equiv a[(i \cdot 3 + j) \cdot 4 + k] \end{aligned}$$

$$\begin{aligned} c[i] &= \&b[3*i] \quad (= b + 3*i) \\ b[j] &= \&a[4*j] \quad (= a + 4*j) \end{aligned}$$

skip i elements of c
 skip j elements of b
 skip k elements of a

skip $i \cdot 3 \cdot 4$ elements of a
 skip $j \cdot 4$ elements of a
 skip k elements of a

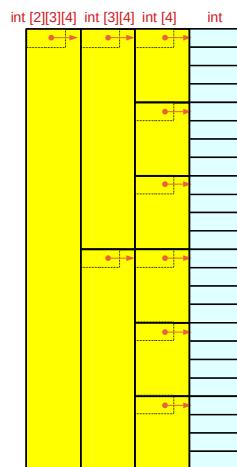


Address Calculation (3)

$$\begin{aligned}
 \text{value}(\mathbf{c} + \mathbf{i}) &= \text{value}(\mathbf{c}) + \mathbf{i} * 3 * 4 * 4 \\
 \text{value}(\mathbf{c[i]} + \mathbf{j}) &= \text{value}(\mathbf{c[i]}) + \mathbf{j} * 4 * 4 \\
 \text{value}(\mathbf{c[i][j]} + \mathbf{k}) &= \text{value}(\mathbf{c[i][j]}) + \mathbf{k} * 4
 \end{aligned}$$

$$\begin{aligned}
 \text{value}(\mathbf{c} + \mathbf{i}) &= \text{value}(\mathbf{c}) + \mathbf{i} * \text{sizeof}(*\mathbf{c}) \\
 \text{value}(\mathbf{c[i]} + \mathbf{j}) &= \text{value}(\mathbf{c[i]}) + \mathbf{j} * \text{sizeof}(*\mathbf{c[i]}) \\
 \text{value}(\mathbf{c[i][j]} + \mathbf{k}) &= \text{value}(\mathbf{c[i][j]}) + \mathbf{k} * \text{sizeof}(*\mathbf{c[i][j]})
 \end{aligned}$$

Array Pointer Approach



$$\begin{aligned}
 \mathbf{c[i]} &\equiv *(\mathbf{c} + \mathbf{i}) \\
 \mathbf{c[i][j]} &\equiv *(\mathbf{c[i]} + \mathbf{j}) \\
 \mathbf{c[i][j][k]} &\equiv *(\mathbf{c[i][j]} + \mathbf{k})
 \end{aligned}$$

$$\begin{aligned}
 \&\mathbf{c[i]} &\equiv \text{value}(\mathbf{c} + \mathbf{i}) \\
 \&\mathbf{c[i][j]} &\equiv \text{value}(\mathbf{c[i]} + \mathbf{j}) \\
 \&\mathbf{c[i][j][k]} &\equiv \text{value}(\mathbf{c[i][j]} + \mathbf{k})
 \end{aligned}$$

$$\begin{aligned}
 \text{value}(\mathbf{c} + \mathbf{i}) &= \text{value}(\mathbf{c}) + \mathbf{i} * 4 \\
 \text{value}(\mathbf{c[i]} + \mathbf{j}) &= \text{value}(\mathbf{c[i]}) + \mathbf{j} * 4 \\
 \text{value}(\mathbf{c[i][j]} + \mathbf{k}) &= \text{value}(\mathbf{c[i][j]}) + \mathbf{k} * 4
 \end{aligned}$$

$(c + i)_{1 \cdot 4}$ $(c[i] + j)_{1 \cdot 4}$ $(c[i][j] + k)_{1 \cdot 4}$

Pointer Array Approach

Subscript [] and dereference * notations (2a)

$$p[i] \leftrightarrow *(p+i)$$

from p , skip
 $i \cdot M \cdot N$ integers

$$\&p[i] = \text{value}((p + i)_{M \cdot N \cdot 4}) \\ = \text{value}(p) + i * M \cdot N \cdot 4$$

$$p[i][j] \leftrightarrow *(*(p+i)+j)$$

from $p[i]$, skip
 $j \cdot N$ integers

$$\&p[i][j] = \text{value}((p[i] + j)_{N \cdot 4}) \\ = \text{value}(p[i]) + j * N \cdot 4$$

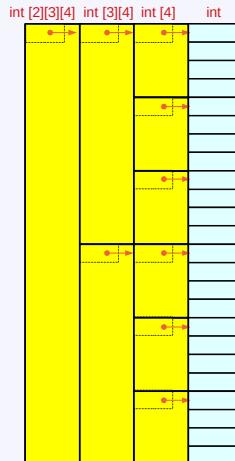
$$p[i][j][k] \leftrightarrow *(*(*(p+i)+j)+k)$$

from $p[i][j]$, skip
 k integers

$$\&p[i][j][k] = \text{value}((p[i][j]+k)_{1 \cdot 4}) \\ = \text{value}(p[i][j]) + k * 1 \cdot 4$$

int $p[L][M][N]$

**Array Pointer
Approach**



address replication

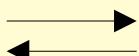
$$\text{value}(p[i]) = \text{value}(p + i)$$

$$\text{value}(p[i][j]) = \text{value}(p[i] + j)$$

$$\&p[i][j][k] = \text{value}(p + i * M \cdot N \cdot 4 + j * N \cdot 4 + k * 4)$$

Subscript [] and dereference * notations (2b)

p[i]



***(p+i)**

skip i pointers
from p

$$\&p[i] = \text{value}(p + i)_{1..4} \\ = \text{value}(p) + i * 4$$

p[i][j]

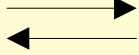


***(*(p+i)+j)**

skip j pointers
from p[i]

$$\&p[i][j] = \text{value}(p[i] + j)_{1..4} \\ = \text{value}(p[i]) + j * 4$$

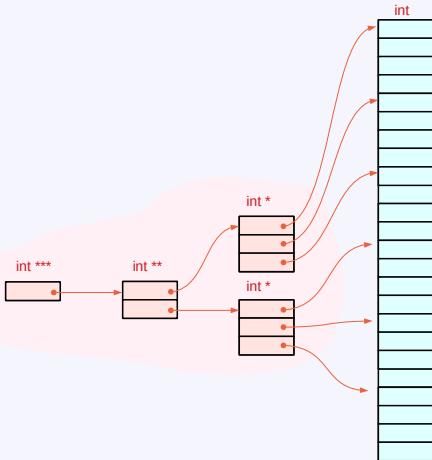
p[i][j][k]



***(*(*p+i)+j)+k)**

skip k integers
from p[i][j]

$$\&p[i][j][k] = \text{value}(p[i][j] + k)_{1..4} \\ = \text{value}(p[i][j]) + k * 4$$



**int ** p [L] ;
int * q [L·M] ;
int r [L·M·N] ;**

**Pointer Array
Approach**

$$\&p[i][j][k] \\ = \text{value}(\text{*value}(\text{*value}(p + i * 4) + j * 4) + k * 4)$$

Equivalences (3)

$p[i]$	\equiv	$*(p+i)$
$p[i][j]$	\equiv	$*(*(p+i)+j)$
$p[i][j][k]$	\equiv	$*(*(*(p+i)+j)+k)$

C Expressions

$\&p[i]$	\equiv	$value(p+i)$
$\&p[i][j]$	\equiv	$value(*(p+i)+j)$
$\&p[i][j][k]$	\equiv	$value(*(*(p+i)+j)+k)$

C Expressions

int $p[L][M][N]$;

$value(*X) = value(X)$ (address replication)

$p[i]$	\longrightarrow	$*(p+i)_{M \cdot N \cdot 4}$
$p[i][j]$	\longrightarrow	$*(*(p+i)_{M \cdot N \cdot 4}+j)_{N \cdot 4}$
$p[i][j][k]$	\longrightarrow	$*(*(*(p+i)_{M \cdot N \cdot 4}+j)_{N \cdot 4}+k)_{1 \cdot 4}$

Math Expressions

$\&p[i]$	\longrightarrow	$value(p+i)_{M \cdot N \cdot 4}$
$\&p[i][j]$	\longrightarrow	$value((p+i)_{M \cdot N \cdot 4}+j)_{N \cdot 4}$
$\&p[i][j][k]$	\longrightarrow	$value(((p+i)_{M \cdot N \cdot 4}+j)_{N \cdot 4}+k)_{1 \cdot 4}$

Math Expressions

int $** p[L], * q[L \cdot M], r[L \cdot M \cdot N]$;

$p[i]$	\longrightarrow	$*(p+i)_{1 \cdot 4}$
$p[i][j]$	\longrightarrow	$*(*(p+i)_{1 \cdot 4}+j)_{1 \cdot 4}$
$p[i][j][k]$	\longrightarrow	$*(*(*(p+i)_{1 \cdot 4}+j)_{1 \cdot 4}+k)_{1 \cdot 4}$

Math Expressions

$\&p[i]$	\longrightarrow	$value(p+i)_{1 \cdot 4}$
$\&p[i][j]$	\longrightarrow	$value(*(p+i)_{1 \cdot 4}+j)_{1 \cdot 4}$
$\&p[i][j][k]$	\longrightarrow	$value(*(*(p+i)_{1 \cdot 4}+j)_{1 \cdot 4}+k)_{1 \cdot 4}$

Math Expressions

Equivalences (4)

int p [L][M][N] ;

$\text{value}(*\mathbf{x}) = \text{value}(\mathbf{x})$ (address replication)

$$\begin{aligned}\&p[i] &= \text{value}((p + i)_{M \cdot N \cdot 4}) = \boxed{\text{value}(p + i * M \cdot N \cdot 4)} \\ \&p[i][j] &= \text{value}((p[i] + j)_{N \cdot 4}) = \boxed{\text{value}(p[i] + j * N \cdot 4)} \\ \&p[i][j][k] &= \text{value}((p[i][j] + k)_{1 \cdot 4}) = \boxed{\text{value}(p[i][j] + k * 1 \cdot 4)} \\ \&p[i][j][k] &= \text{value}(p + i * M \cdot N \cdot 4 + j * N \cdot 4 + k * 4)\end{aligned}$$

$\&p[i]$	\longrightarrow	$\text{value}(p + i)_{M \cdot N \cdot 4}$
$\&p[i][j]$	\longrightarrow	$\text{value}((p + i)_{M \cdot N \cdot 4} + j)_{N \cdot 4}$
$\&p[i][j][k]$	\longrightarrow	$\text{value}(((p + i)_{M \cdot N \cdot 4} + j)_{N \cdot 4} + k)_{1 \cdot 4}$

Math Expressions

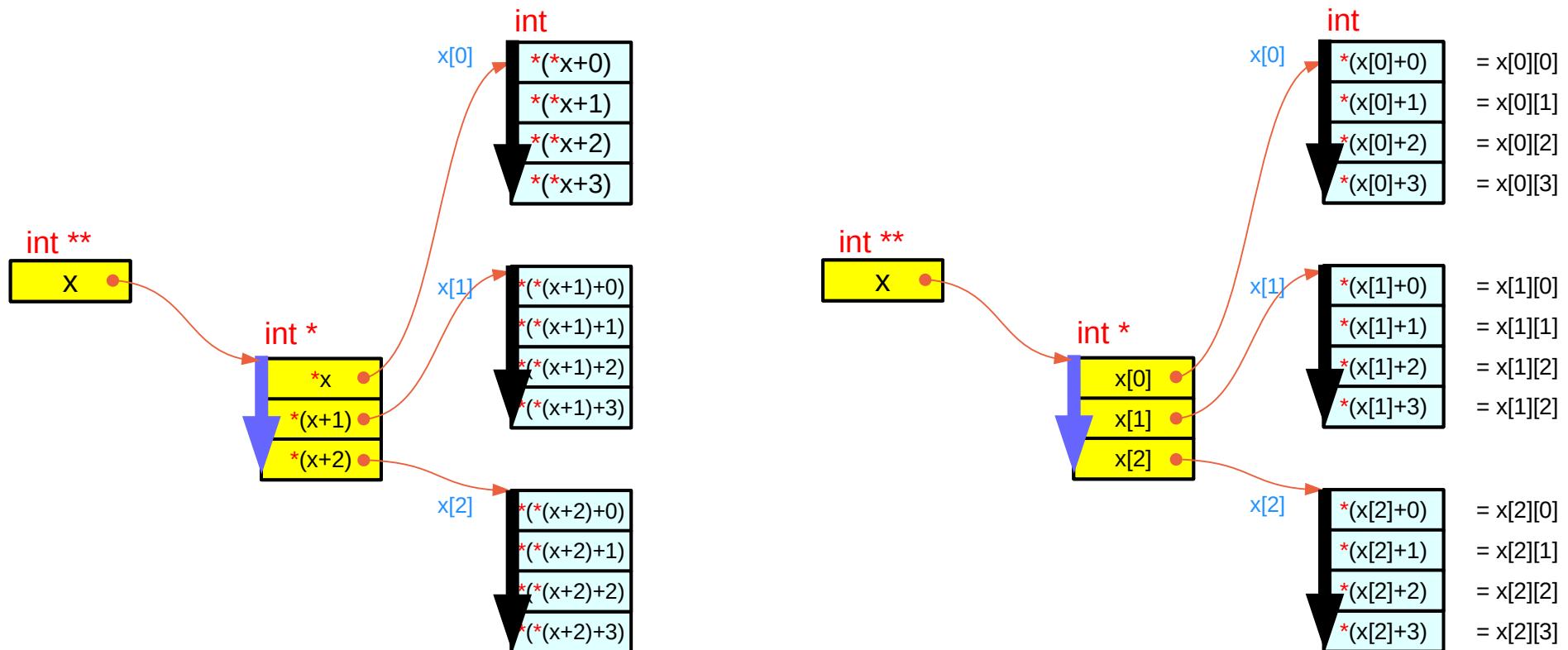
int ** p[L], * q[L·M], r[L·M·N] ;

$$\begin{aligned}\&p[i] &= \text{value}(p + i)_{1 \cdot 4} = \boxed{\text{value}(p + i * 1 \cdot 4)} \\ \&p[i][j] &= \text{value}(p[i] + j)_{1 \cdot 4} = \boxed{\text{value}(p[i] + j * 1 \cdot 4)} \\ \&p[i][j][k] &= \text{value}(p[i][j] + k)_{1 \cdot 4} = \boxed{\text{value}(p[i][j] + k * 1 \cdot 4)} \\ \&p[i][j][k] &= \text{value}(*\text{value}(*\text{value}(p + i * 4) + j * 4) + k * 4)\end{aligned}$$

$\&p[i]$	\longrightarrow	$\text{value}(p + i)_{1 \cdot 4}$
$\&p[i][j]$	\longrightarrow	$\text{value}(*(p + i)_{1 \cdot 4} + j)_{1 \cdot 4}$
$\&p[i][j][k]$	\longrightarrow	$\text{value}(*(*(p + i)_{1 \cdot 4} + j)_{1 \cdot 4} + k)_{1 \cdot 4}$

Math Expressions

* and [] notations (1)



C expression

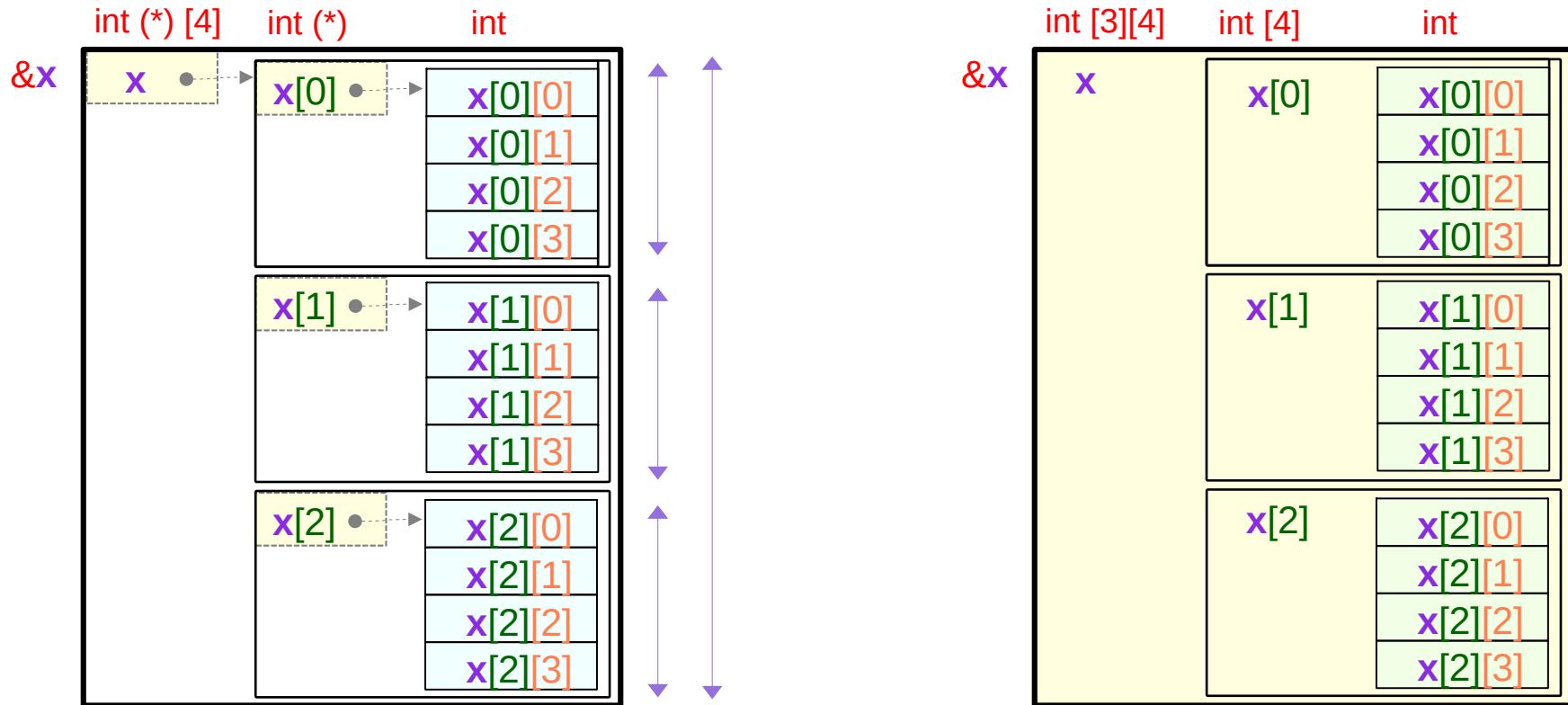
 $\ast(\ast(x+i)+j)$

 $x[i][j]$

Math expression

 $\ast(\ast(x+i)_{1..4}+j)_{1..4}$

* and [] notations (2)



C expression

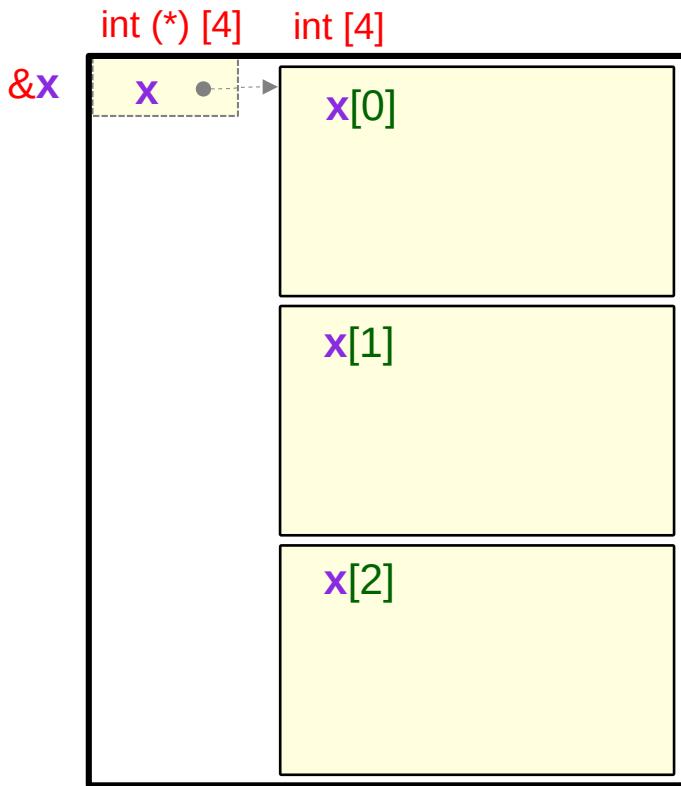
$$*(\ast(x+i)+j)$$


$$x[i][j]$$

Math expression

$$*(\ast(\ast(x+i)_{4 \cdot 4} + j)_{1 \cdot 4})$$

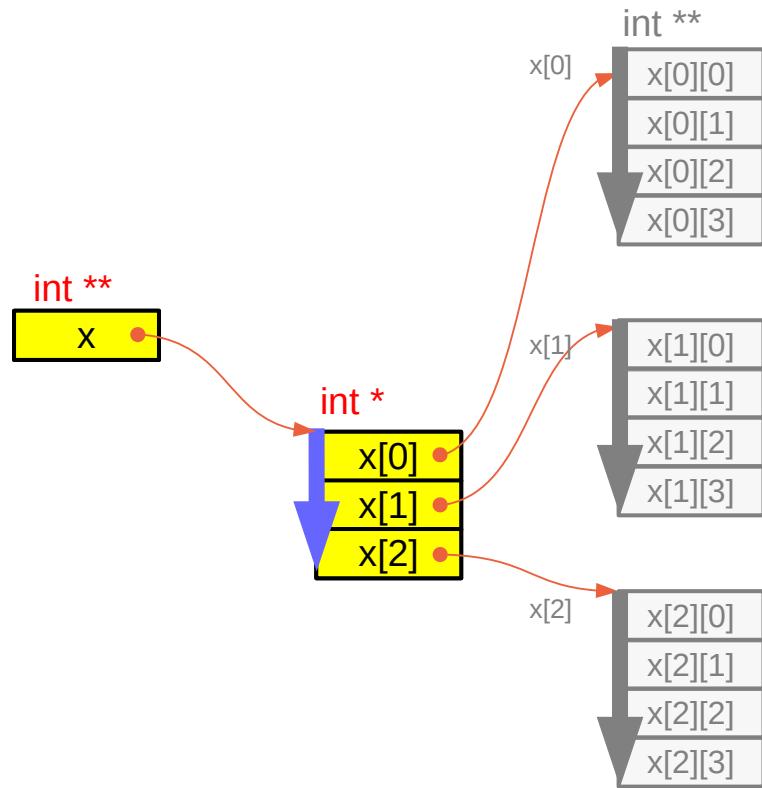
Virtual pointers vs. real pointers (1)



$$\&x = x \quad \text{sizeof}(x) = 3 * \text{sizeof}(*x)$$

$$x[i] = *(x + i)_{4..4}$$

$$\begin{aligned} \text{value}(x+i) &= \text{value}(x) + i * \text{sizeof}(*x) \\ &= \text{value}(x) + i * 4 * 4 \end{aligned}$$

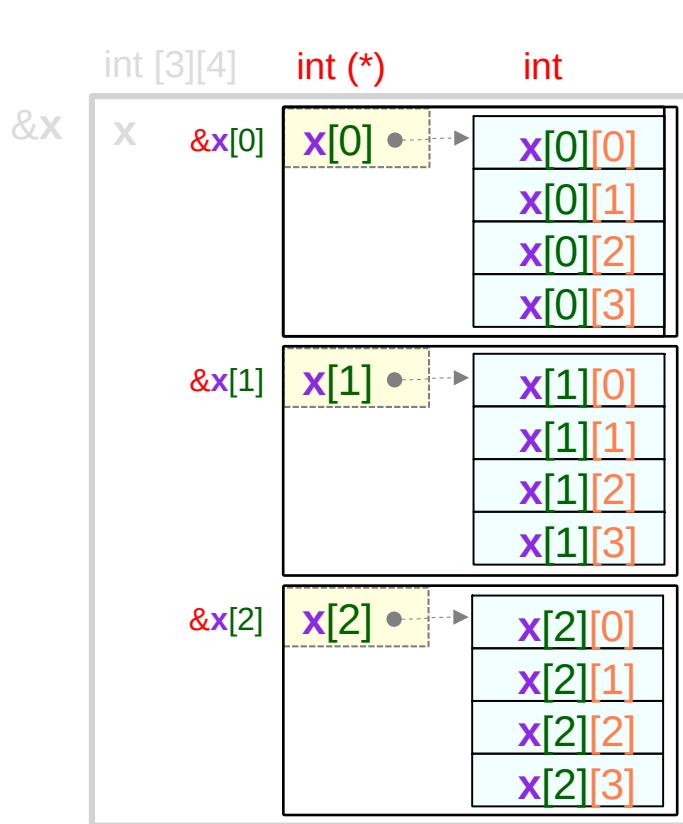


$$\&x \neq x \quad \text{sizeof}(x) = \text{sizeof}(*x) = 4$$

$$x[i] = *(x + i)_{1..4}$$

$$\begin{aligned} \text{value}(x+i) &= \text{value}(x) + i * \text{sizeof}(*x) \\ &= \text{value}(x) + i * 4 \end{aligned}$$

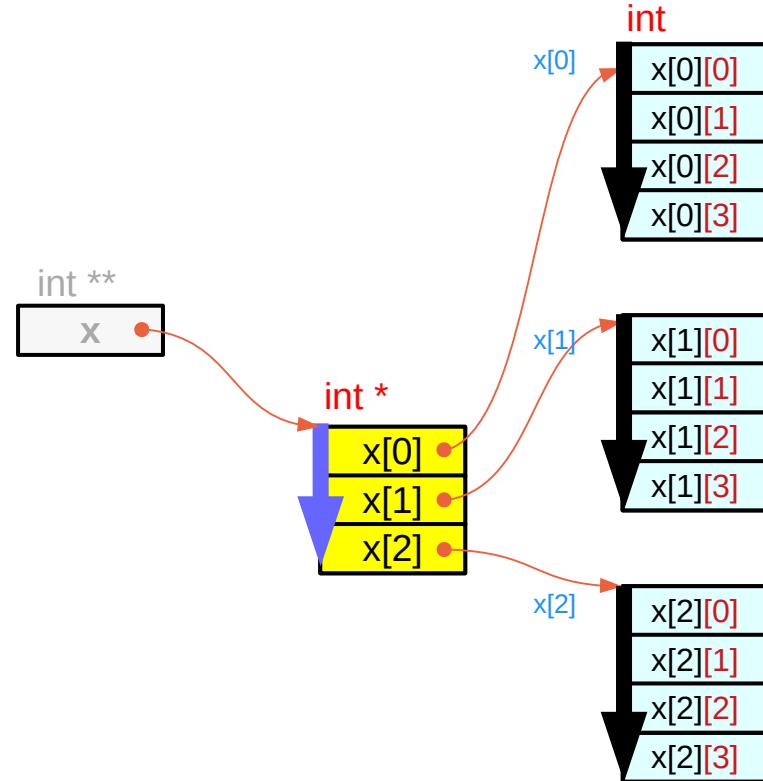
Virtual pointers vs. real pointers (2)



$$\&x[i] = x[i] \quad \text{sizeof}(x[i]) = 4 * \text{sizeof}(*x[i])$$

$$x[i] = *(x + i)_{1..4}$$

$$\begin{aligned} \text{value}(x[i]+j) &= \text{value}(x[i]) + j * \text{sizeof}(*x[i]) \\ &= \text{value}(x[i]) + j * 4 \end{aligned}$$

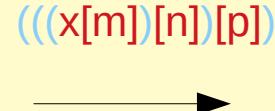


$$\&x[i] = x[i] \quad \text{sizeof}(x[i]) = \text{sizeof}(*x[i]) = 4$$

$$x[i][j] = *(x[i] + j)_{1..4}$$

$$\begin{aligned} \text{value}(x[i]+j) &= \text{value}(x[i]) + j * \text{sizeof}(*x[i]) \\ &= \text{value}(x[i]) + j * 4 \end{aligned}$$

Operator Precedence

Precedence	Operator	Description	Associativity
1	<code>++ --</code> <code>()</code> <code>[]</code> <code>.</code> <code>-></code> <code>(type){list}</code>	Suffix/postfix increment and decrement Function call Array subscripting Structure and union member access member access through pointer Compound literal(C99)	Left-to-right 
2	<code>++ --</code> <code>+ -</code> <code>! ~</code> <code>(type)</code> <code>*</code> <code>&</code> <code>sizeof</code> <code>_Alignof</code>	Prefix increment and decrement Unary plus and minus Logical NOT and bitwise NOT Type cast Indirection (dereference) Address-of Size-of Alignment requirement(C11)	Right-to-left 

https://en.cppreference.com/w/c/language/operator_precedence

Limitations

No index Range Checking

Array Size must be a constant expression

Variable Array Size

Arrays cannot be Copied or Compared

Aggregate Initialization and Global Arrays

Precedence Rule

Index Type Must be Integral

References

- [1] Essential C, Nick Parlante
- [2] Efficient C Programming, Mark A. Weiss
- [3] C A Reference Manual, Samuel P. Harbison & Guy L. Steele Jr.
- [4] C Language Express, I. K. Chun
- [5] <https://pdos.csail.mit.edu/6.828/2008/readings/pointers.pdf>