

Booth Encoding

20161005

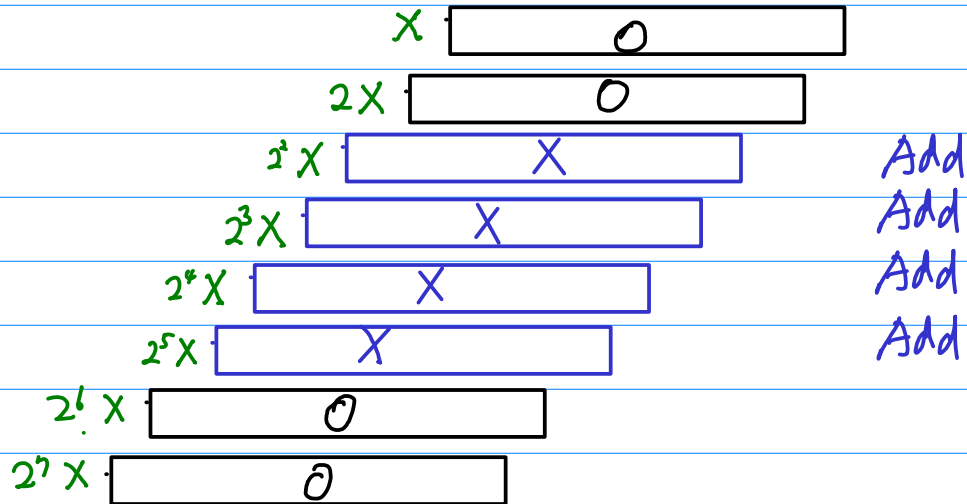
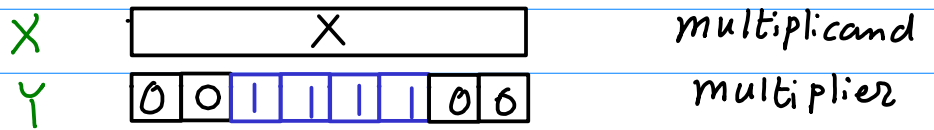
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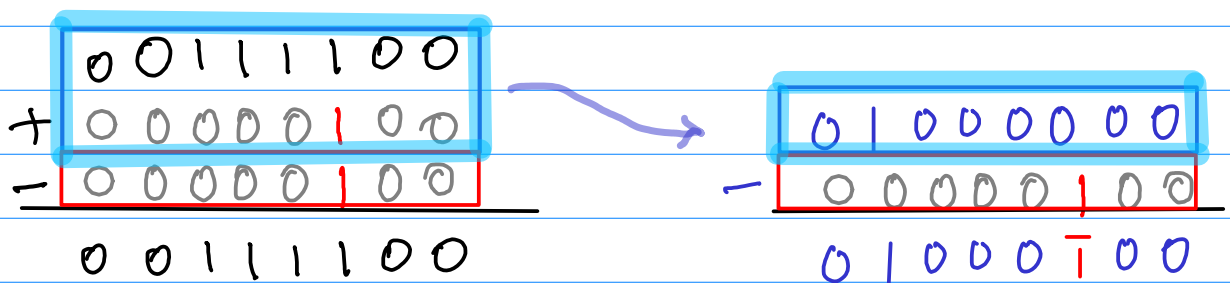
Based on

MJ Flynn's EE 486 Lecture 7 : Integer Multiplication, Stanford University

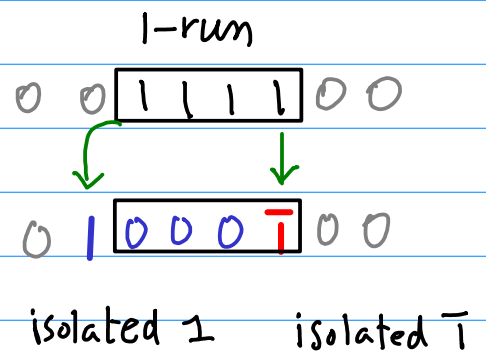
of Partial Products to be added



Booth Encoding

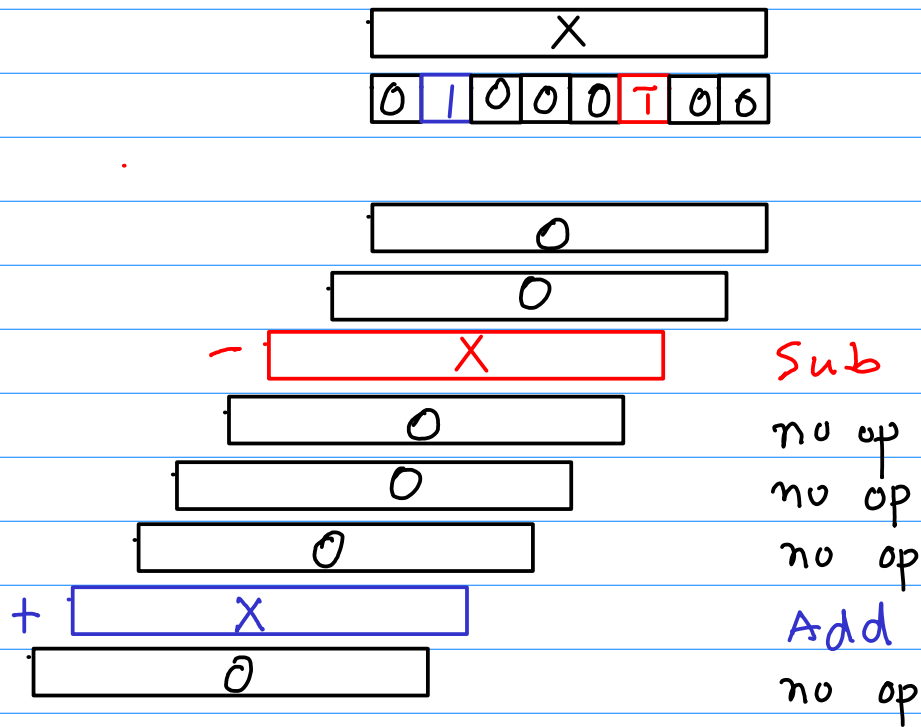


$+1 \rightarrow 0$

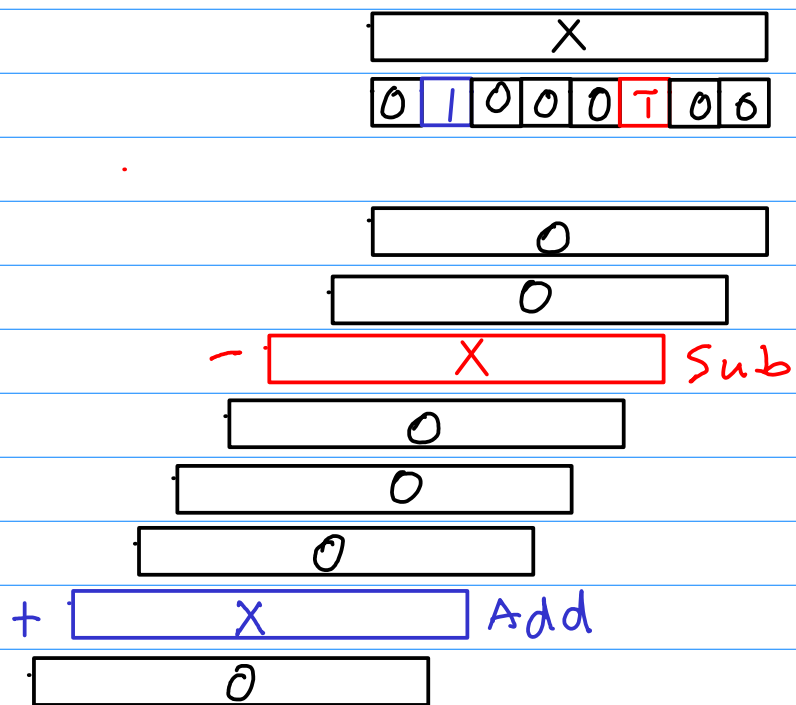
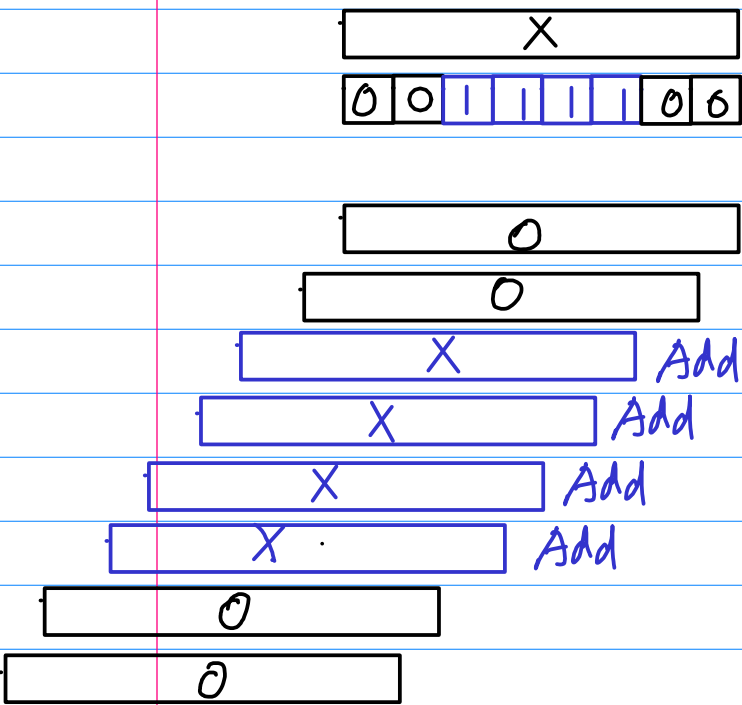


of Partial Products to be added

Original Booth Encoded



{ Best case ✓
 { Average case ✓
 { Worst case



worst case

\textcircled{n} add

=

worst case

\textcircled{n} add/sub

Booth encoding :

the worst case

still

the number of partial products $\Rightarrow \textcircled{n}$

Not much gain

reduces the number of partial products to be added
delay reduction

works well for serial multiplication
Variable latency

⊙ Worst case : alternating 1's and 1's
1 1 1 1 1 1 1 1

⊙ Booth encoding does not significantly improve
the worst case

⇒ use modified booth encoding

Modified Booth 2

2-bit encoding + 1-bit overlapping

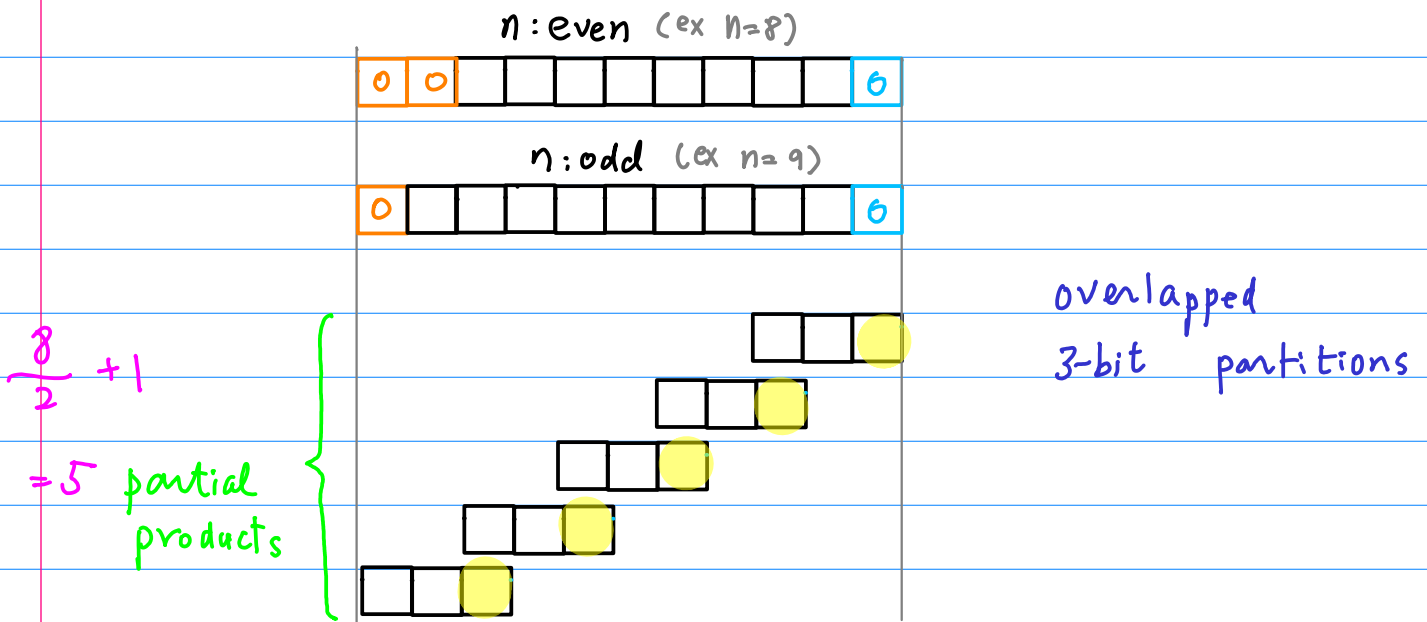
worst case # of p.p.'s = $\frac{N}{2} + 1$

modified Booth 3

3-bit encoding + 1-bit overlapping

worst case # of p.p.'s = $\frac{N}{3} + 1$

Modified Booth 2 (unsigned case)



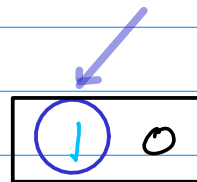
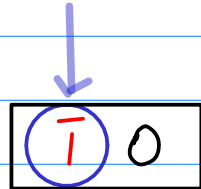
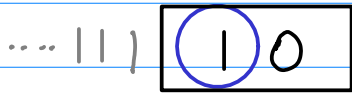
2-bit encoding		2^1 2^0	scale factor	
all zero's	$000 \rightarrow 000$		$0 \cdot 2 + 0 = +0$	$+0$
end of 1's	$001 \rightarrow 010$		$0 \cdot 2 + 1 = +1$	$+X$
isolated 1	$010 \rightarrow 1T0$		$1 \cdot 2 + T = +1$	$+X$
end of 1's	$011 \rightarrow 100$		$1 \cdot 2 + 0 = +2$	$+2X$
start of 1's	$100 \rightarrow T00$		$T \cdot 2 + 0 = -2$	$-2X$
isolated 0	$101 \rightarrow T10$		$T \cdot 2 + 1 = -1$	$-X$
start of 1's	$110 \rightarrow 0T0$		$0 \cdot 2 + T = -1$	$-X$
all 1's	$111 \rightarrow 000$		$0 \cdot 2 + 0 = 0$	$+0$

Scale factor $\{0, \pm 1, \pm 2\}$

0-run & 1-run

start of a possible 1-run

end of a possible 1-run

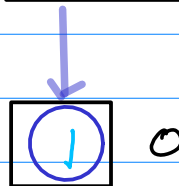
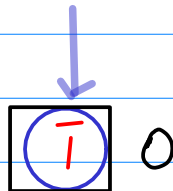
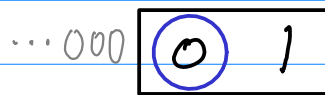
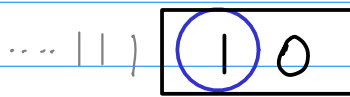


start of a possible 1's

end of a possible 1's

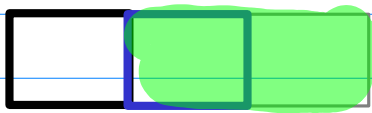
→ start of 1-run

→ start of 0-run



to encode a 1-bit, we need 2-bit info.

Needs 3-bits for encoding a 2-bit



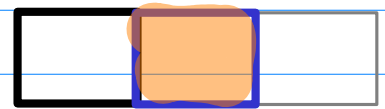
0	0
0	1
1	0
1	1

0-run

0-run starts

1-run starts

1-run



0
1
1
0



0	0
0	1
1	0
1	1

0-run

0-run starts

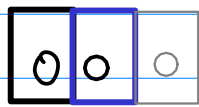
1-run starts

1-run

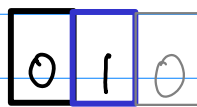


0
1
1
0

Encoding 2-bit : 2 choices

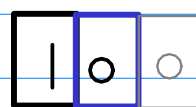


0 0 0

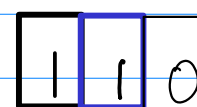


0 T 0
 1 0 0

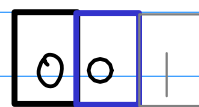
 1 T 0



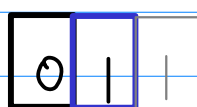
T 0 0



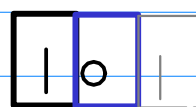
0 T 0



0 1 0

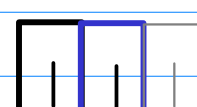


1 0 0

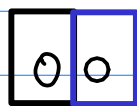


0 1 0
 T 0 0

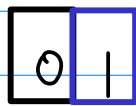
 T 1 0



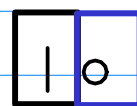
0 0 0



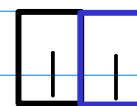
{ 0 0
 0 1



{ 1 T
 1 0



{ T 0
 T 1



{ 0 T
 0 0



+0



1 +X



2-1 +X



2 +2X



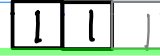
-2 -2X



-2+1 -X



-1 -X



+0

Scale Factor for 2-bit encoding

all zero's	→	middle of 0-run	000 → 000	+0
end of 1's	→	start of 0-run at (j)	001 → 010	+X
isolated 1	→	isolated 1 at (j)	010 → 1T0	+X
end of 1's	→	start of 0-run at (j+1)	011 → 100	+2X
start of 1's	→	start of 1-run at (j+1)	100 → T00	-2X
isolated 0	→	isolated 0 at (j)	101 → T10	-X
start of 1's	→	start of 1-run at (j)	110 → 0T0	-X
all 1's	→	middle of 1-run	111 → 000	+0

$$\begin{array}{r}
 X = \boxed{} \\
 \times Y = \boxed{10111011} \\
 \hline
 \end{array}$$

multiplier

multiplier

↓
modified booth encode 2

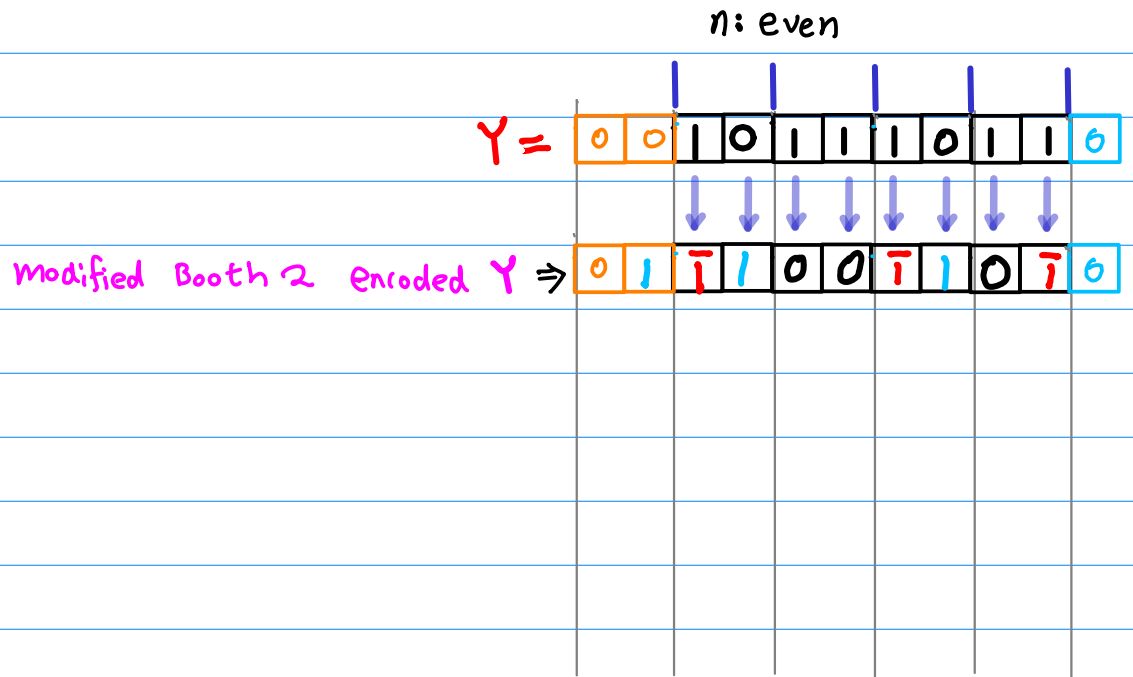
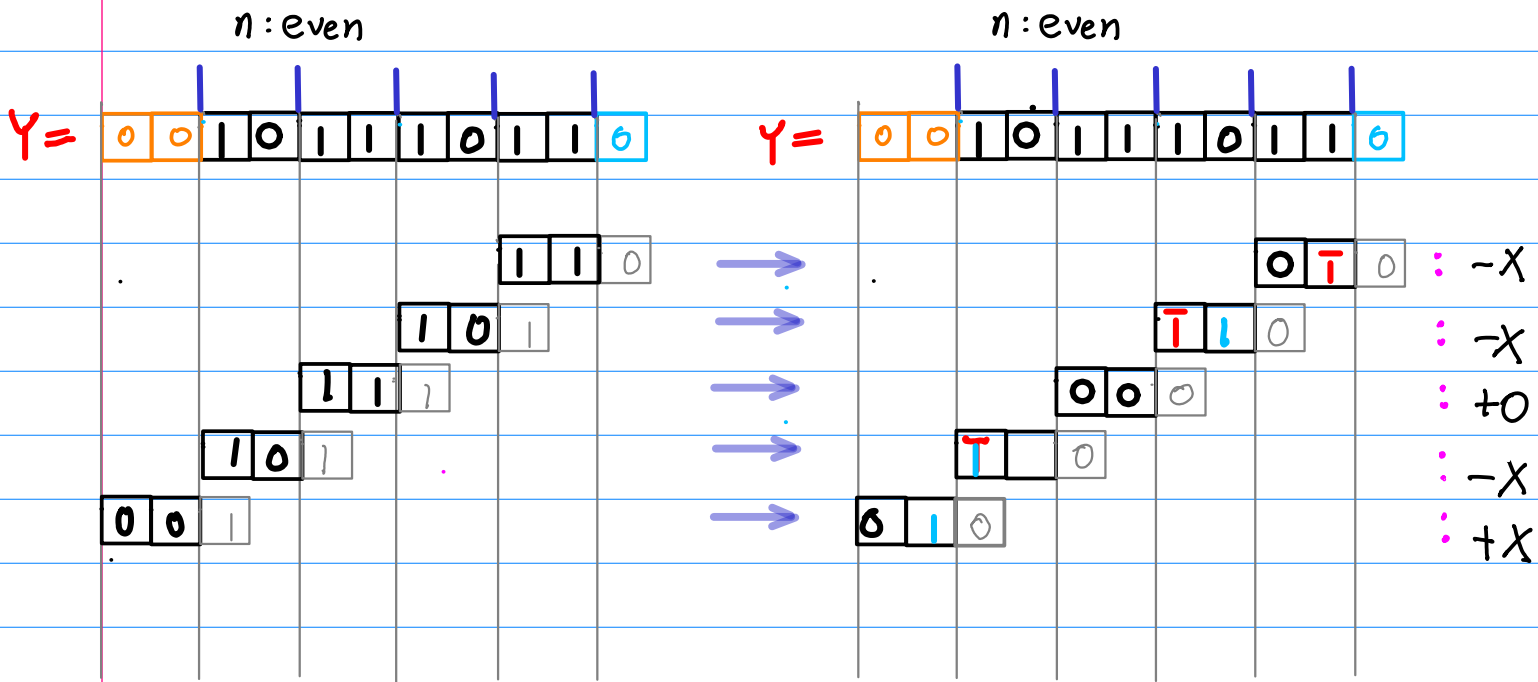
3

↓
scale factor

↓
shift multiplicand = X

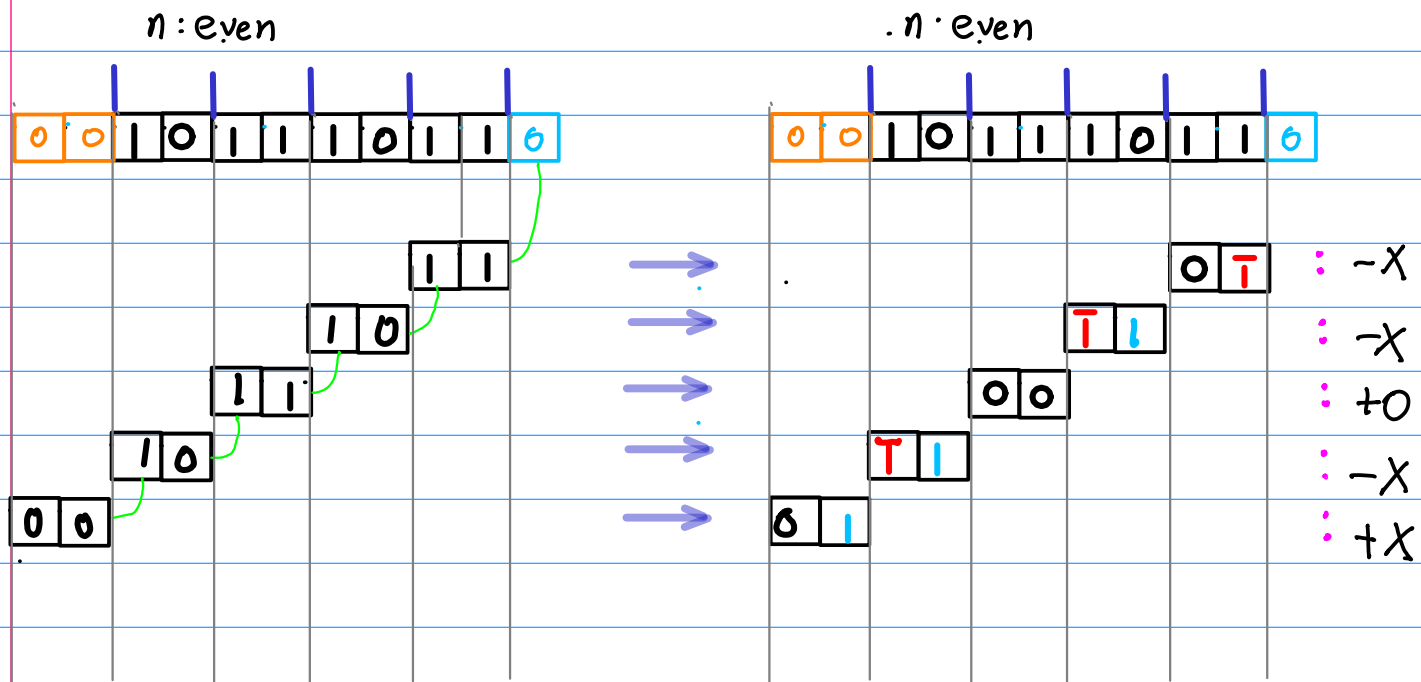
↓
reduced number of
partial products
for multiplication

Multiplier Encoding



$$\begin{array}{r}
 X = \boxed{} \\
 \times Y = \boxed{10111011}
 \end{array}$$

Basically, 2-bit encoding scheme



Context
encoding

increasing overlapping bit \rightarrow no. need

0 0 0 0	middle of 0-run	0 0 0 0
0 0 0 1	middle of 0-run	0 0 0 0
0 0 1 0	start of 0-run at (j)	0 1 0 0
0 0 1 1	start of 0-run at (j)	0 1 0 0
0 1 0 0	isolated 1 at (j)	1 1 0 0
0 1 0 1	isolated 1 at (j)	1 1 0 0
0 1 1 0	start of 0-run at (j+1)	1 0 0 0
0 1 1 1	start of 0-run at (j+1)	1 0 0 0
1 0 0 0	start of 1-run at (j+1)	1 0 0 0
1 0 0 1	start of 1-run at (j+1)	1 0 0 0
1 0 1 0	isolated 0 at (j)	1 1 0 0
1 0 1 1	isolated 0 at (j)	1 1 0 0
1 1 0 0	start of 1-run at (j)	0 1 0 0
1 1 0 1	start of 1-run at (j)	0 1 0 0
1 1 1 0	middle of 1-run	0 0 0 0
1 1 1 1	middle of 1-run	0 0 0 0

Overlapping bit at both ends \rightarrow no need

0 0 0 0	middle of 0-run	0 0 0 0
0 0 0 1	start of 0-run at (j)	0 0 1 0
0 0 1 0	isolated 1 at (j)	0 1 1 0
0 0 1 1	start of 0-run at (j+1)	0 1 0 0
0 1 0 0	start of 1-run at (j+1)	0 1 0 0
0 1 0 1	isolated 0 at (j)	0 1 1 0
0 1 1 0	start of 1-run at (j)	0 0 1 0
0 1 1 1	middle of 1-run	0 0 0 0
1 0 0 0	middle of 0-run	0 0 0 0
1 0 0 1	start of 0-run at (j)	0 0 1 0
1 0 1 0	isolated 1 at (j)	0 1 1 0
1 0 1 1	start of 0-run at (j+1)	0 1 0 0
1 1 0 0	start of 1-run at (j+1)	0 1 0 0
1 1 0 1	isolated 0 at (j)	0 1 1 0
1 1 1 0	start of 1-run at (j)	0 0 1 0
1 1 1 1	middle of 1-run	0 0 0 0

2-bit context encoding



	middle of 0-run	
	middle of 1-run	
	start of 0-run at (j)	
	start of 0-run at (j+1)	
	start of 1-run at (j)	
	start of 1-run at (j+1)	
	isolated 0 at (j)	
	isolated 1 at (j)	

⌈		middle of 0-run	
		start of 0-run at (j)	
⌈		isolated 1 at (j)	
		start of 0-run at (j+1)	
⌈		start of 1-run at (j+1)	
		isolated 0 at (j)	
⌈		start of 1-run at (j)	
		middle of 1-run	



context encoding

5 Types of Partial Products

$000 \rightarrow 000$		$+0$	} 5 type of scaled multiplicand
$001 \rightarrow 010$	1	$+X$	
$010 \rightarrow 1T0$	$2-1$	$+X$	
$011 \rightarrow 100$	2	$+2X$	
$100 \rightarrow T00$	-2	$-2X$	
$101 \rightarrow T10$	$-2+1$	$-X$	
$110 \rightarrow 0T0$	-1	$-X$	
$111 \rightarrow 000$		-0	

$+0$	0	} 2's complement
$+X$	x_i	
$+2X \leftarrow$	x_{i-1}	
$-X$	$\overline{x_i}$	
$-2X \leftarrow$	$\overline{x_{i-1}}$	

\uparrow i -th bit of the scaled partial products

$$1 * X \longrightarrow (i)$$

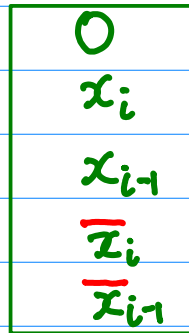
$$2 * X \longrightarrow (i-1)$$

$$(+) \longrightarrow \square$$

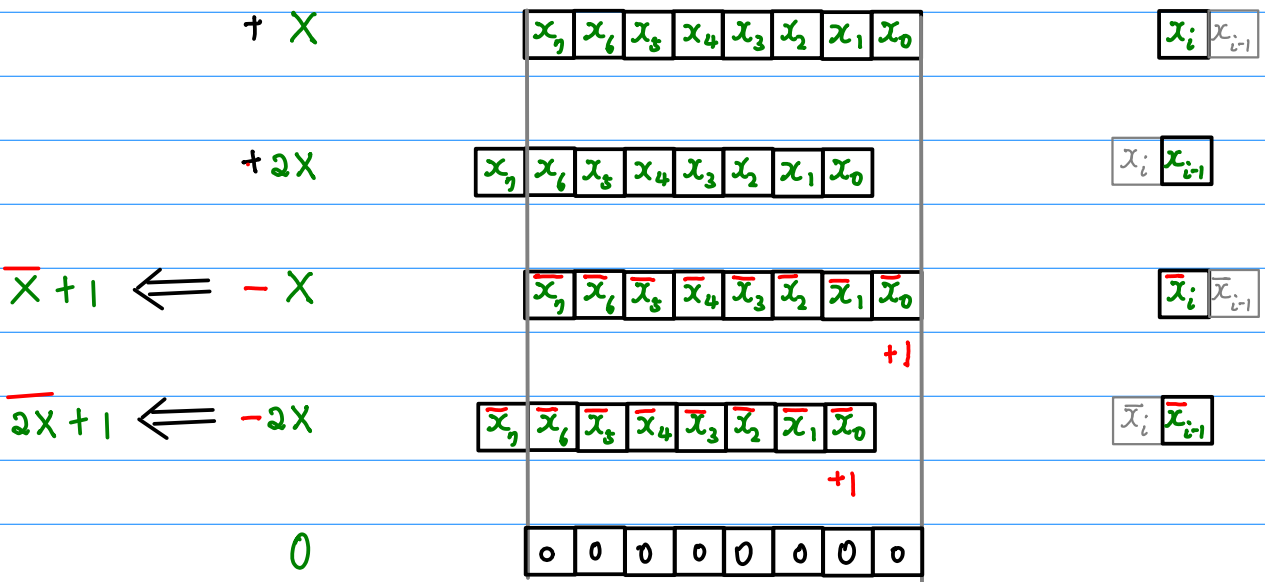
$$(-) \longrightarrow \square \quad \text{1's complement}$$

5 types of scaling \times (=multiplier)

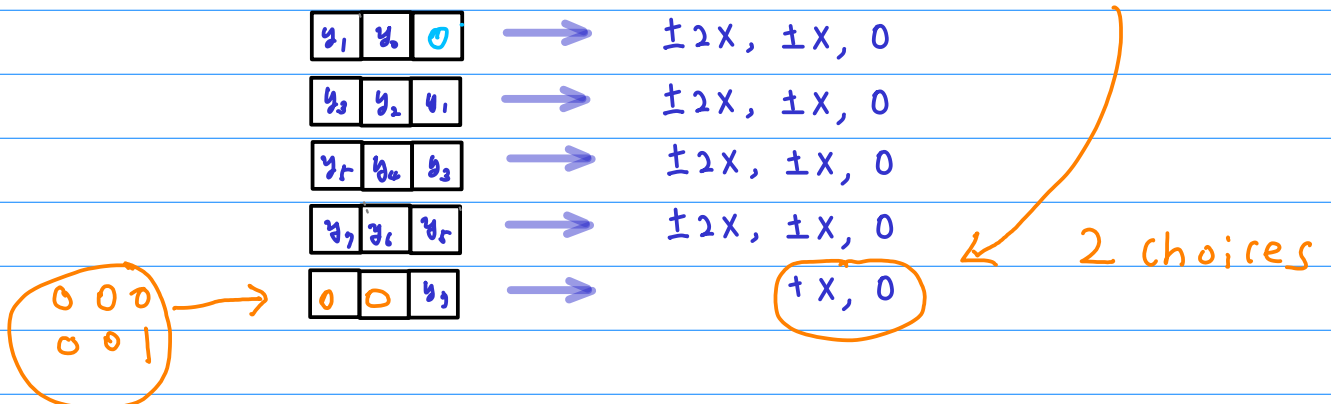
$+0$
 $+X$
 $+2X \leftarrow$
 $-X$
 $-2X \leftarrow$



i -th bit of the scaled partial products

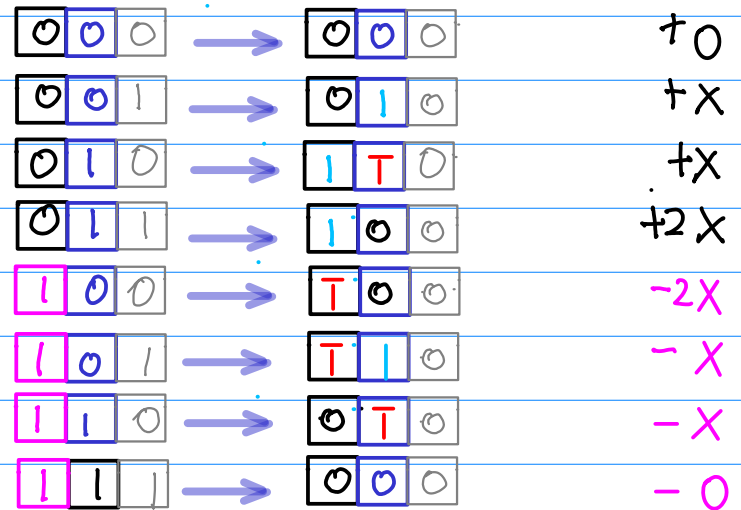
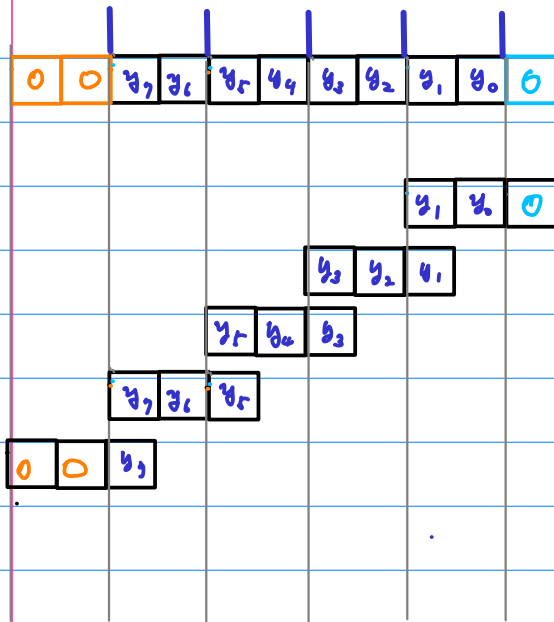


* the last partial product (sign-bit p.p.)

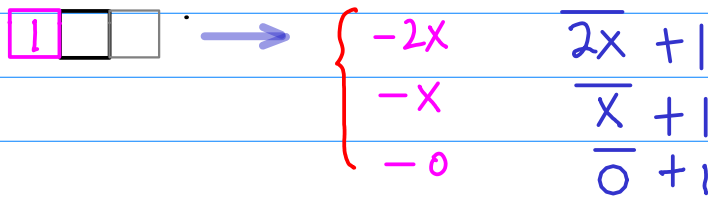


$y_1, y_3, y_5, y_7 = 1 \Rightarrow$ negative scaling

n : even



3-bit partition



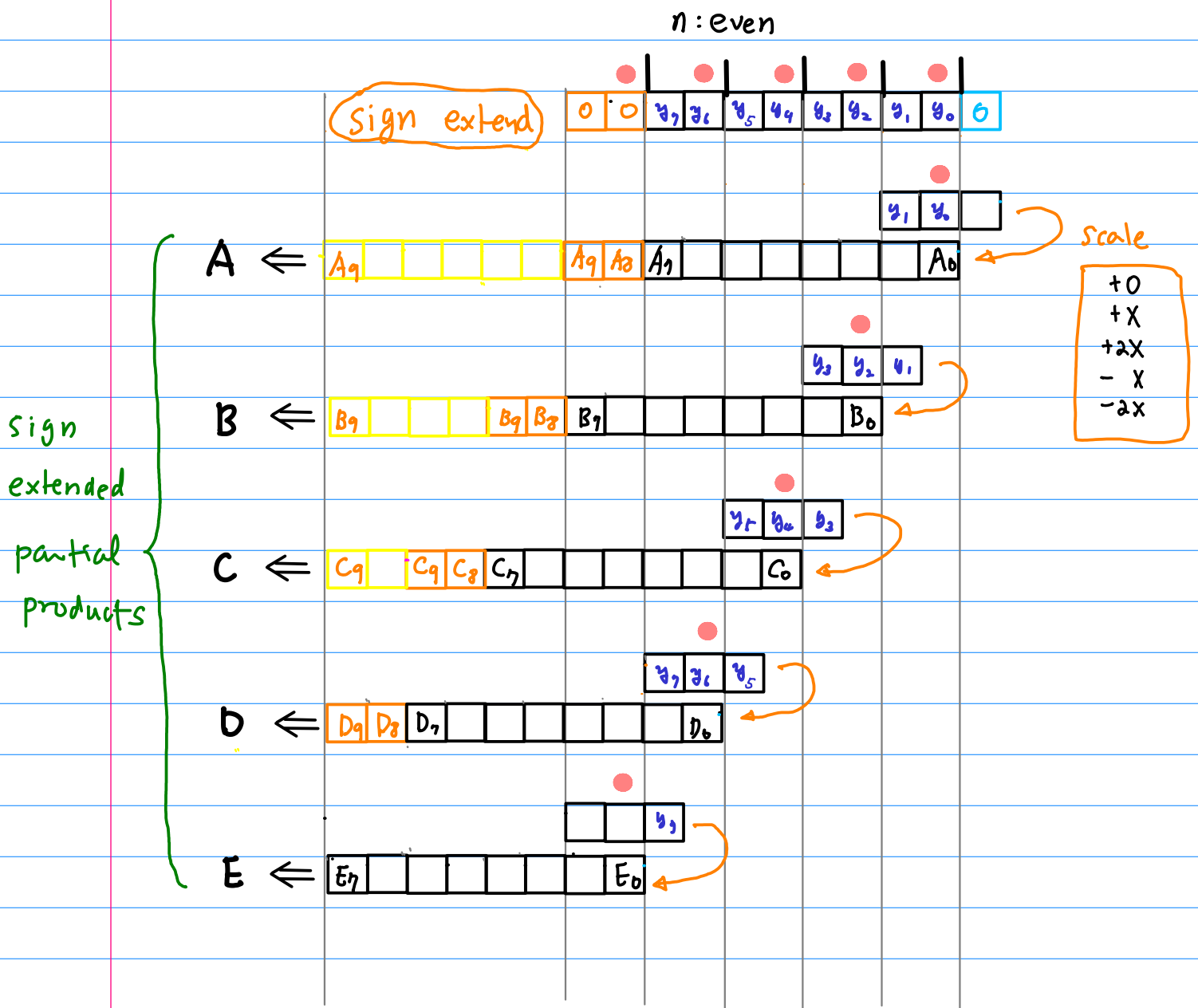
Complement $\begin{pmatrix} 2X \\ X \\ 0 \end{pmatrix} + 1$

Scaled X's are aligned at y_0, y_2, y_4, y_6, y_8

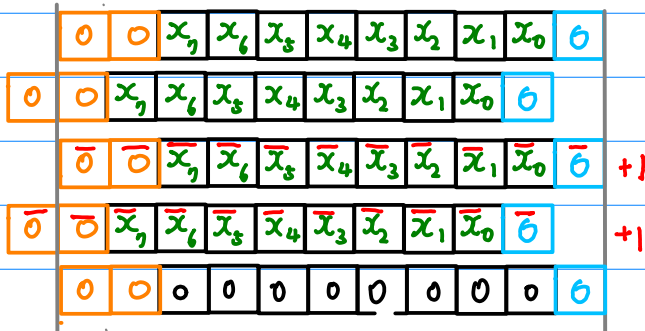
Sign extended partial product A B C D E

$$X = \boxed{x_7 x_6 x_5 x_4 x_3 x_2 x_1 x_0} \quad \text{multiplicand}$$

$$Y = \overset{\bullet}{y_7} \overset{\bullet}{y_6} \overset{\bullet}{y_5} \overset{\bullet}{y_4} \overset{\bullet}{y_3} \overset{\bullet}{y_2} \overset{\bullet}{y_1} \overset{\bullet}{y_0} \quad \text{multiplier}$$

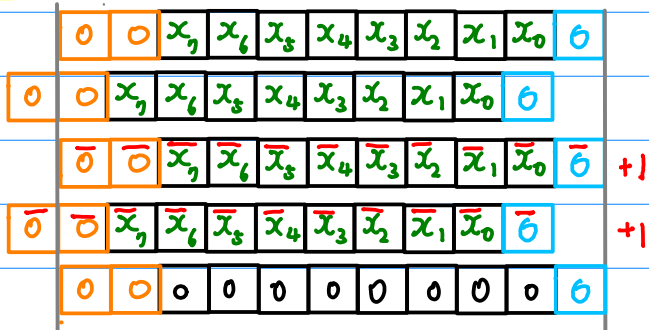


A ← A_9 A_8 A_7 A_6 A_5 A_4 A_3 A_2 A_1 A_0



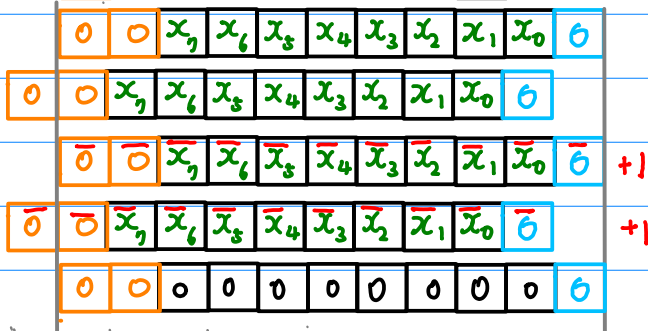
5 choices

B ← B_9 B_8 B_7 B_6 B_5 B_4 B_3 B_2 B_1 B_0



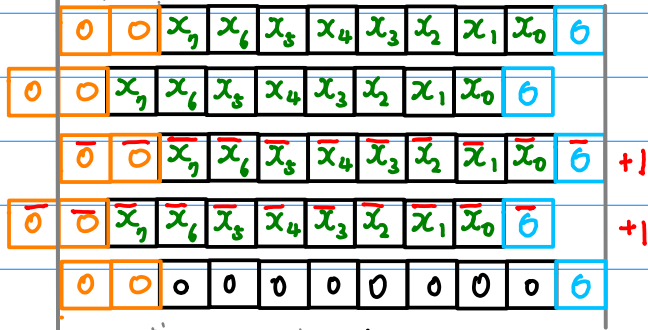
5 choices

C ← C_9 C_8 C_7 C_6 C_5 C_4 C_3 C_2 C_1 C_0



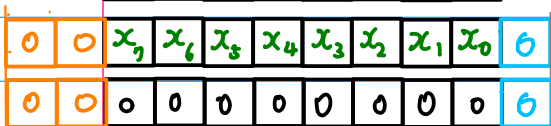
5 choices

D ← D_9 D_8 D_7 D_6 D_5 D_4 D_3 D_2 D_1 D_0



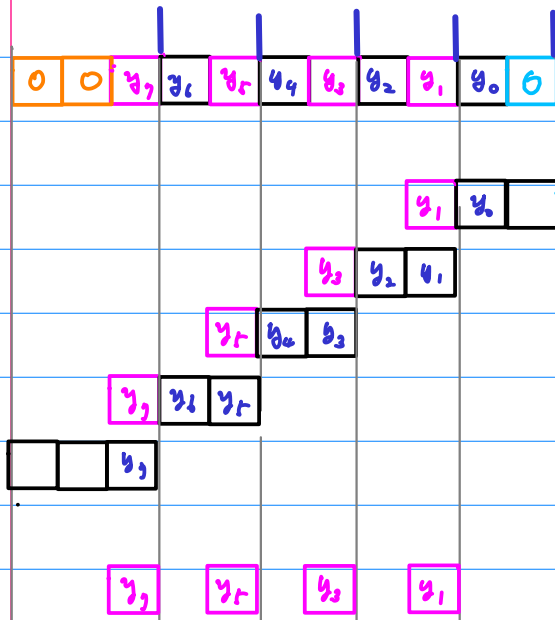
5 choices

E ← E_7 E_6 E_5 E_4 E_3 E_2 E_1 E_0



2 choices

n : Even



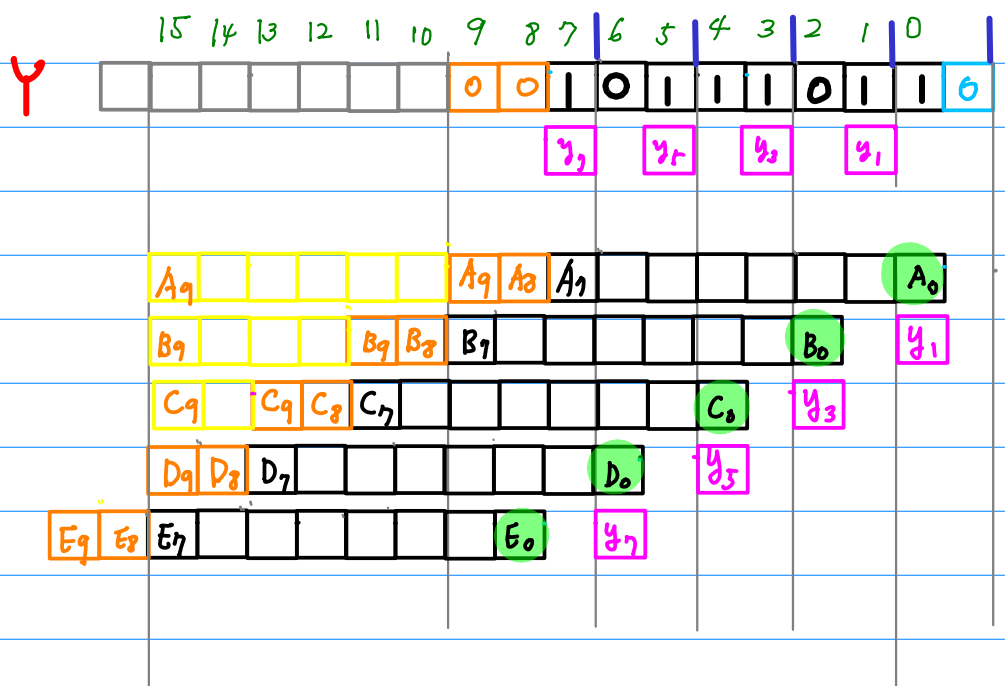
if $y_1=1$, complement partial product + y_1

if $y_3=1$, complement partial product + y_3

if $y_5=1$, complement partial product + y_5

if $y_7=1$, complement partial product + y_7

n : Even



if $y_1=1$, complement A

if $y_3=1$, complement B

if $y_5=1$, complement C

if $y_7=1$, complement D

E

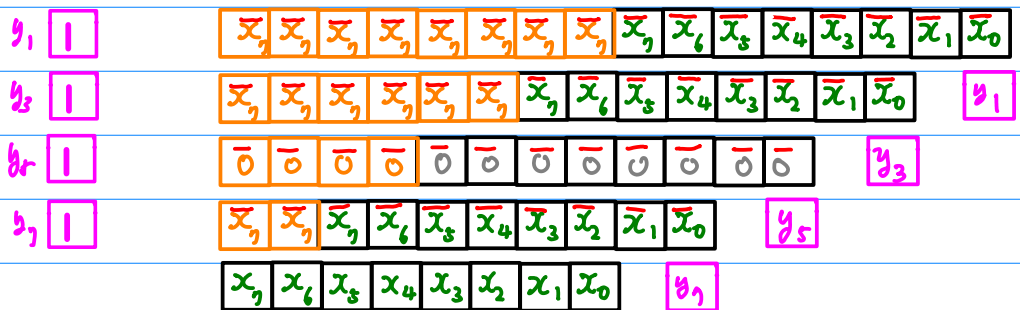
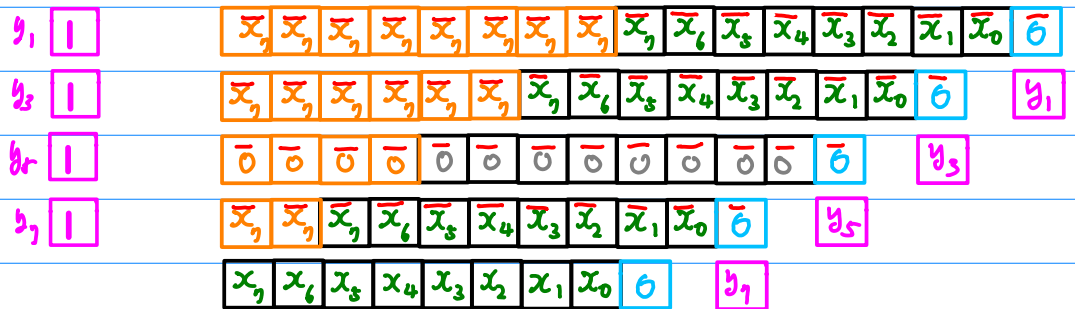
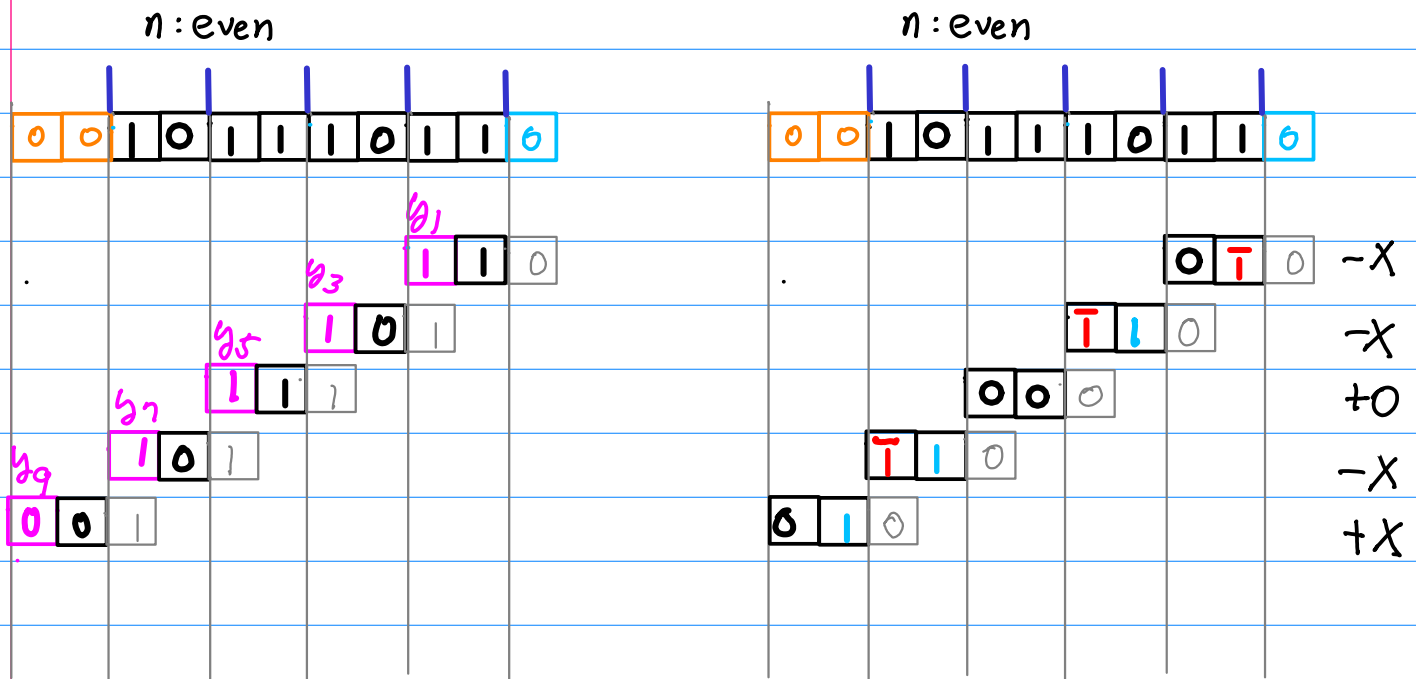
if $y_1=1$, add 1 else add 0 ($y_1=0$)

if $y_3=1$, add 1 else add 0 ($y_3=0$)

if $y_5=1$, add 1 else add 0 ($y_5=0$)

if $y_7=1$, add 1 else add 0 ($y_7=0$)

Handling Negative Scales

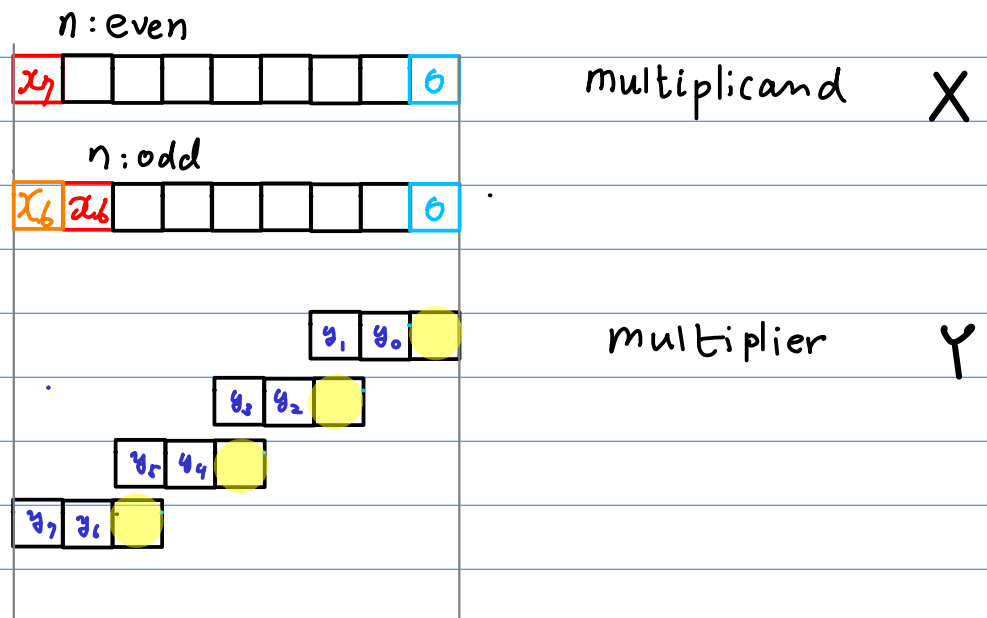


Modified Booth 2 (signed case)

the worst case $\frac{N}{2}$ partial products

$$\frac{8}{2} = 4$$

make
4 partial products



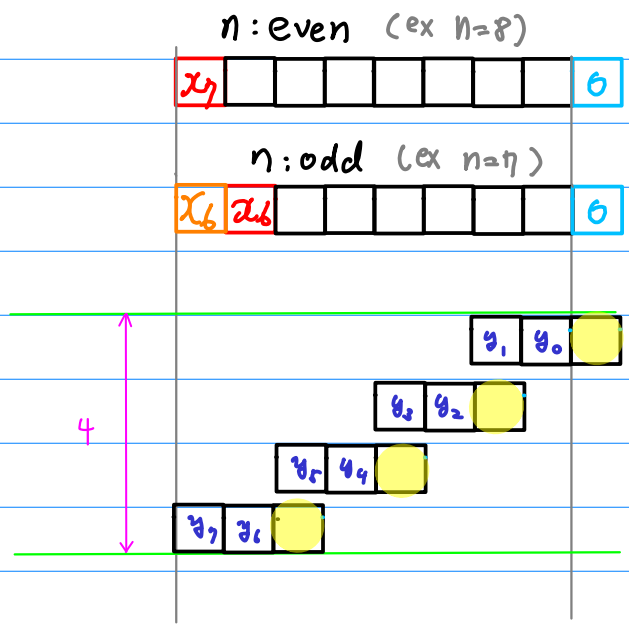
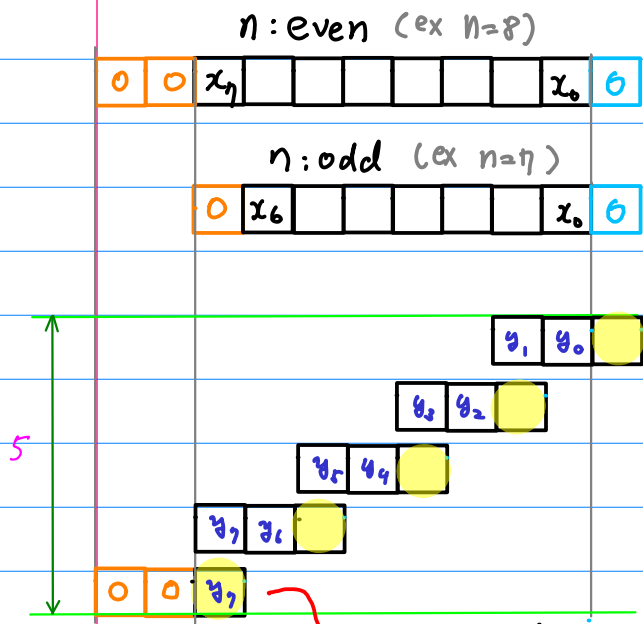
x_7, y_7 Sign bit

Compared to unsigned multiplication

One less partial products

Booth 2 Unsigned Case

Booth 2 Signed Case



$$\frac{8}{2} + 1 = 5$$

0	0	0	($y_7 = 0$)	+0
0	0	1	($y_7 = 1$)	+X

Necessary to make
all positive numbers
(unsigned)

$$\frac{8}{2} = 4$$

Booth 2 code :

Signed Digit Number
→ no special treatment
for the Sign bit (y_7)

4-bit unsigned number & Booth 2 Code

0	0 0 0 0 0	0 0 0 0 0	0
1	0 0 0 0 1	0 0 0 1 1	$2 - 1 = 1$
2	0 0 0 1 0	0 0 1 1 0	$4 - 2 = 2$
3	0 0 0 1 1	0 0 1 0 1	$4 - 1 = 3$
4	0 0 1 0 0	0 1 1 0 0	$8 - 4 = 4$
5	0 0 1 0 1	0 1 1 1 1	$8 - 4 + 2 - 1 = 5$
6	0 0 1 1 0	0 1 0 1 0	$8 - 2 = 6$
7	0 0 1 1 1	0 1 0 0 1	$8 - 1 = 7$
8	0 1 0 0 0	1 1 0 0 0	$16 - 8 = 8$
9	0 1 0 0 1	1 1 0 1 1	$16 - 8 + 2 - 1 = 9$
10	0 1 0 1 0	1 1 1 1 0	$16 - 8 + 4 - 2 = 10$
11	0 1 0 1 1	1 1 1 0 1	$16 - 8 + 4 - 1 = 11$
12	0 1 1 0 0	1 0 1 0 0	$16 - 4 = 12$
13	0 1 1 0 1	1 0 1 1 1	$16 - 4 + 2 - 1 = 13$
14	0 1 1 1 0	1 0 0 1 0	$16 - 2 = 14$
15	0 1 1 1 1	1 0 0 0 1	$16 - 1 = 15$

Booth 2 code :

Signed Digit Number

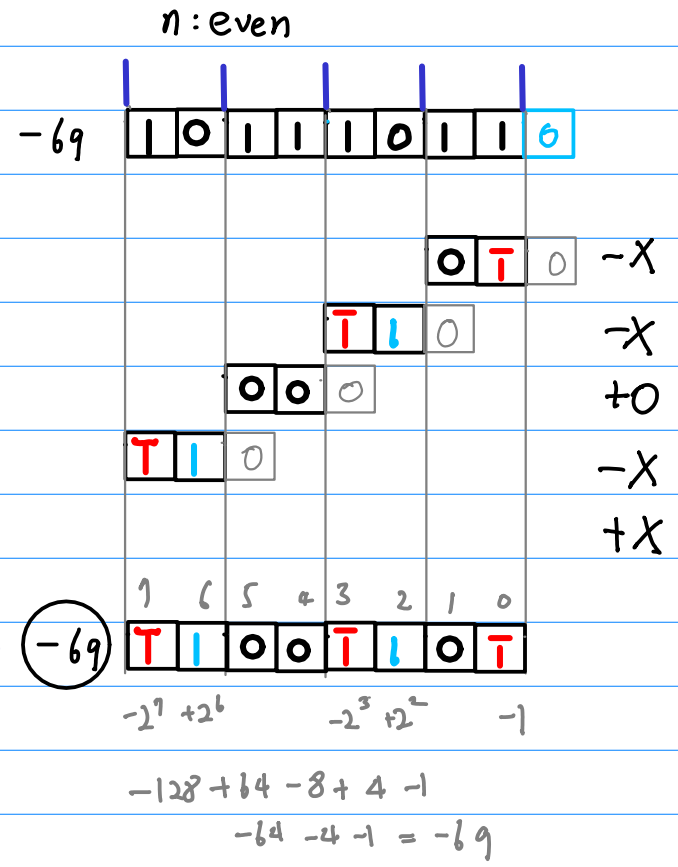
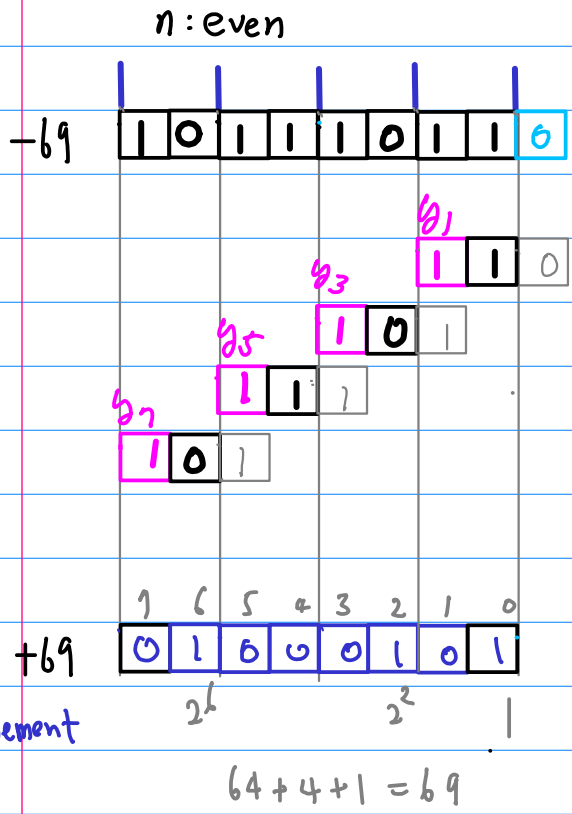
→ Need special treatment

to make all positive numbers

4-bit 2's complement number & Booth 2 Code (Signed)

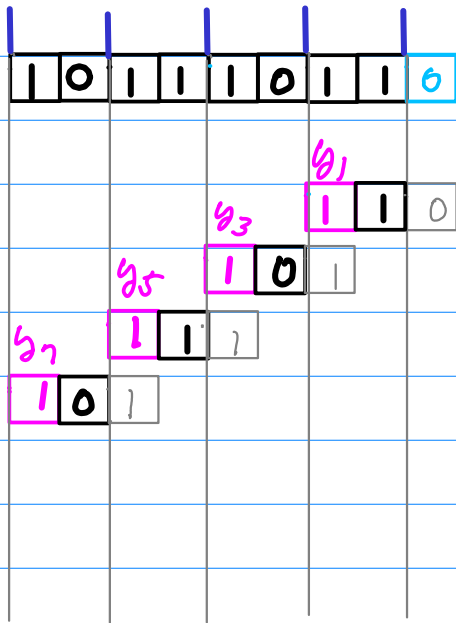
0	0000	0000	0
1	0001	0011	$2 - 1 = 1$
2	0010	0110	$4 - 2 = 2$
3	0011	0101	$4 - 1 = 3$
4	0100	1100	$8 - 4 = 4$
5	0101	1111	$8 - 4 + 2 - 1 = 5$
6	0110	1010	$8 - 2 = 6$
7	0111	1001	$8 - 1 = 7$
-8	1000	1000	-8
-7	1001	1011	$-8 + 2 - 1 = -7$
-6	1010	1110	$-8 + 4 - 2 = -6$
-5	1011	1101	$-8 + 4 - 1 = -5$
-4	1100	0100	-4
-3	1101	0111	$-4 + 2 - 1 = -3$
-2	1110	0010	-2
-1	1111	0001	-1

Booth 2 code : Signed Digit Number
 → no special treatment
 for the sign bit (y₃)

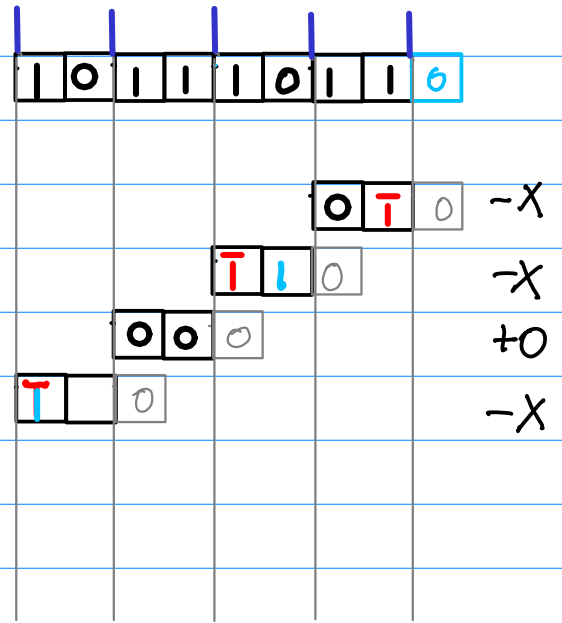


no problem in Booth encoding
signed numbers!

n : even



n : even



q_1

q_3

q_5

q_7

$\bar{x}_7 \bar{x}_6 \bar{x}_5 \bar{x}_4 \bar{x}_3 \bar{x}_2 \bar{x}_1 \bar{x}_0$

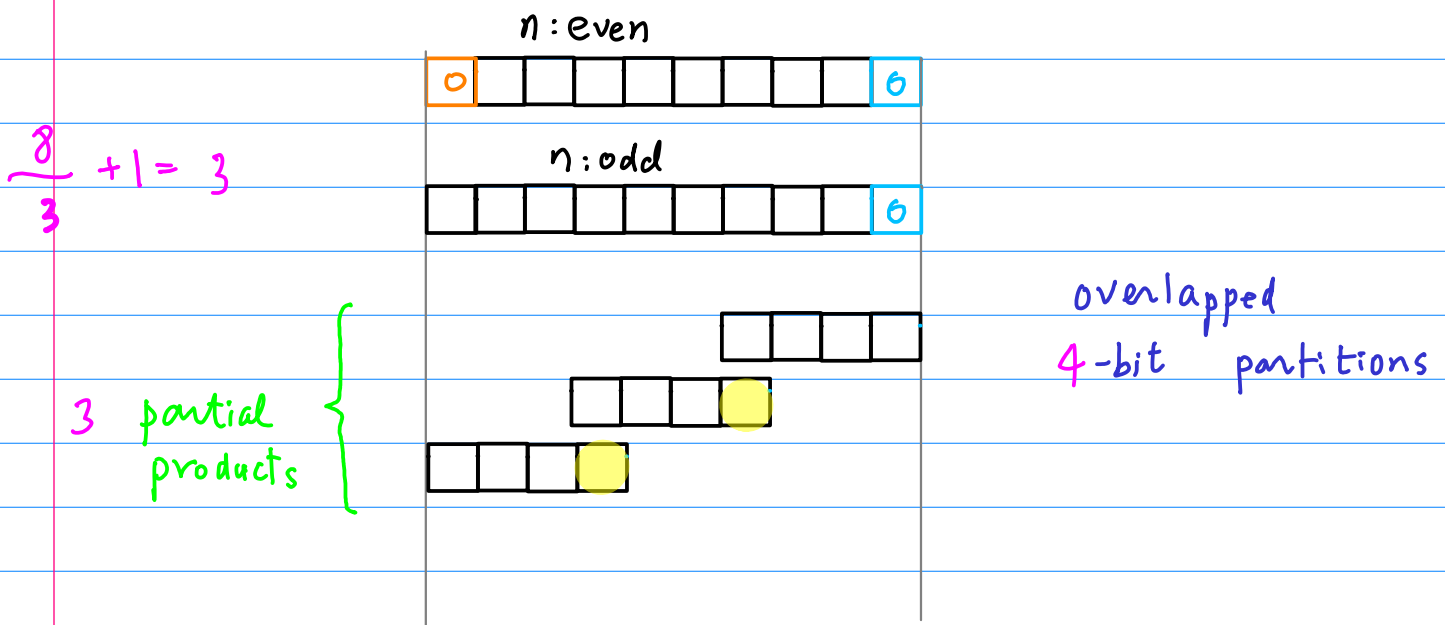
$\bar{x}_7 \bar{x}_6 \bar{x}_5 \bar{x}_4 \bar{x}_3 \bar{x}_2 \bar{x}_1 \bar{x}_0$

$\bar{0} \bar{0} \bar{0} \bar{0} \bar{0} \bar{0} \bar{0} \bar{0}$

$\bar{x}_7 \bar{x}_6 \bar{x}_5 \bar{x}_4 \bar{x}_3 \bar{x}_2 \bar{x}_1 \bar{x}_0$

Modified Booth 3 (unsigned)

the worst case $\frac{N}{3} + 1$ partial products



Scale factor

0000	→	0000	+0	1000	→	1000	-4X
0001	→	0010	+X	1001	→	1010	-3X
0010	→	0110	+X	1010	→	1110	-3X
0011	→	0100	+2X	1011	→	1100	-2X
0100	→	1100	+2X	1100	→	0100	-2X
0101	→	1110	+3X	1101	→	0110	-X
0110	→	1010	+3X	1110	→	0010	-X
0111	→	1000	+4X	1111	→	0000	-0



