

Monad P3 : Haskel Expressions (1E)

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Haskell Expressions

Expressions and values

Because Haskell is a **purely functional language**,
all **computations** are done via the **evaluation** of
expressions (syntactic terms) to yield **values**

Every **value** has an associated **type**.
(Intuitively, we can think of **types** as **sets of values**.)

Examples of **expressions** include **atomic values**
such as the **integer 5**, the **character 'a'**,
and the **function $\lambda x \rightarrow x+1$** ,
as well as **structured values**
such as the **list [1,2,3]** and the **pair ('b',4)**.

Expressions



Value

Type



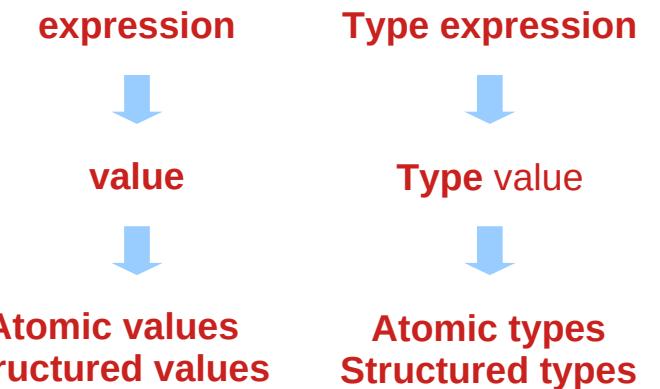
Atomic values
Structured values

<https://www.haskell.org/tutorial/goodies.html>

Type expressions and types

Just as **expressions** denote **values**,
type expressions are **syntactic terms**
that denote **type values** (or just **types**).

Examples of **type expressions** include the **atomic types**
Integer (infinite-precision integers),
Char (characters),
Integer->Integer (functions mapping Integer to Integer),
as well as the **structured types**
[Integer] (homogeneous lists of integers) and
(Char,Integer) (character, integer pairs).



<https://www.haskell.org/tutorial/goodies.html>

First class values

All Haskell values are "first-class"

- they may be passed as arguments to functions,
- returned as results,
- placed in data structures, etc.

Haskell types, on the other hand, are not first-class.

<https://www.haskell.org/tutorial/goodies.html>

Typing

Types in a sense describe values, and
the association of a value with its type is called a **typing**.

Using the examples of values and types above,
we write **typing** as follows: (the ":" can be read "has type.")

```
5 :: Integer
'a' :: Char
inc :: Integer -> Integer
[1,2,3] :: [Integer]
('b',4) :: (Char,Integer)
```

<https://www.haskell.org/tutorial/goodies.html>

Function definition and declaration

Functions in Haskell are normally defined by a **series of equations**.

For example, the function **inc** can be defined by the single equation:

```
inc n      = n+1
```

An **equation** is an example of a **declaration**.

Another kind of **declaration** is a **type signature** declaration,
with which we can declare an **explicit typing** for inc:

```
inc      :: Integer -> Integer
```

<https://www.haskell.org/tutorial/goodies.html>

Expression evaluation =>

when we wish to indicate that an **expression e1 evaluates**, or
"reduces," to *another expression* or **value e2**, we will write:

e1 => e2

For example, note that:

inc (inc 3) => 5

<https://www.haskell.org/tutorial/goodies.html>

Statements vs Expressions

Many programming languages differentiate statements from expressions.

Statement: What code does

Expression: What code is

can think the term "**statement**" very broadly to refer to anything that is not an expression or type declaration.

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Imperative vs functional languages

statements vs. **expressions** closely parallels
imperative languages vs. **functional languages**:

Imperative: A language that *emphasizes statements*

Functional: A language that *emphasizes expressions*

C lies at one end of the spectrum (imperative),
relying heavily on **statements** to accomplish everything.

Haskell lies at the exact opposite extreme (functional),
using **expressions** heavily:

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Statement examples in the imperative language C

```
#include <stdio.h>

int main(int argc, char *argv[]) {
    int elems[5] = {1, 2, 3, 4, 5};          // statement

    int total = 0;
    int i;

    for (i = 0; i < 5; i++) {                // statement
        total += elems[i];                  // statement
    }
    printf("%d\n", total);                  // statement

    return 0;
}
```

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Expression examples in the functional language Haskell (1)

everything in Haskell is an **expression**,
and even **statements** are **expressions**.

```
main = print (sum [1..5])          -- Expression
```

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Expression examples in the functional language Haskell (2)

For example, the following code might appear to be
a traditional imperative-style sequence of statements:

```
main = do
    putStrLn "Enter a number:"          -- Statement?
    str <- getLine                   -- Statement?
    putStrLn ("You entered: " ++ str)  -- Statement?
```

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Expression examples in the functional language Haskell (3)

but **do** notation is merely **syntactic sugar**

for nested applications of (**>>=**), which is itself nothing more than
an infix higher-order function:

main =

```
putStrLn "Enter a number:"  >>= (\_  ->      -- Expression
          getLine                  >>= (str ->      -- Sub-expression
          putStrLn ("You entered: " ++ str) ))    -- Sub-expression
```

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Statement-as-expression

In Haskell, "statements" are actually **nested expressions**,
and **sequencing statements** just builds larger and larger **expressions**.

This **statement-as-expression** paradigm promotes consistency
and prevents arbitrary language limitations,
such as Python's restriction of lambdas to single statements.

In Haskell, you cannot limit
the number of statements a **term** uses
any more than you can limit the number of sub-expressions.

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Monads

do notation works for more than just **IO**.

Any **type** that implements the **Monad class**
can be "sequenced" in **statement form**,
as long as it supports the following two operations:

class Monad m where

(>>=) :: m a -> (a -> m b) -> m b

return :: a -> m a

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Statement-like syntax using monads

This provides a uniform interface for translating imperative **statement-like** syntax into **expressions** under the hood.

For example, the **Maybe** type implements the Monad class:

```
data Maybe a = Nothing | Just a
```

```
instance Monad Maybe where
```

```
  m >>= f = case m of
```

```
    Nothing -> Nothing
```

```
    Just a -> f a
```

```
  return = Just
```

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

do notation using monads

This lets you assemble **Maybe-based** computations using **do** notation

example :: Maybe Int

example = do

x <- Just 1

y <- Nothing

return (x + y)

example =

Just 1 >>= (\x ->

Nothing >>= (\y ->

return (x + y)))

The above code desugars to nested calls to (**>>=**):

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Substitute `>>=` and `return`

The compiler then **substitutes** in our definition of (`>>=`) and **return**

```
example = case (Just 1) of  
    Nothing -> Nothing  
    Just x -> case Nothing of  
        Nothing -> Nothing  
        Just y -> Just (x + y)
```

```
example =  
  Just 1  >>= (\x ->  
    Nothing >>= (\y ->  
      return (x + y) ))
```

```
instance Monad Maybe where  
  m >>= f = case m of  
    Nothing -> Nothing  
    Just a -> f a  
  return = Just
```

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Evaluate the outer and inner case expression

We can then hand-evaluate this expression to prove that it short-circuits when it encounters Nothing:

-- Evaluate the outer `case`
example = case Nothing of
 Nothing -> Nothing
 Just y -> Just (1 + y)

-- Evaluate the remaining `case`
example = Nothing

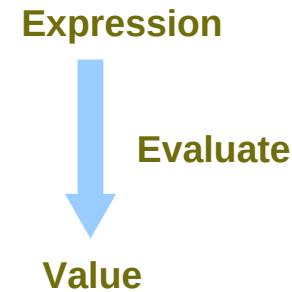
example = case (Just 1) of
 Nothing -> Nothing
 Just x -> case Nothing of
 Nothing -> Nothing
 Just y -> Just (x + y)

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Everything is an expression to be evaluated

Notice that we can evaluate these **Maybe** "statements"
without invoking any sort of **abstract machine**.

When everything is an **expression**,
everything is simple to **evaluate**
and does not require *understanding* or
invoking an execution model.



FSM not needed
for sequencing

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Semantics

Semantics

In fact, the distinction between **statements** and **expressions** also closely parallels another important divide: the difference between **operational semantics** and **denotational semantics**.

Operational semantics:

Translates code to **abstract machine statements**

Denotational semantics:

Translates code to **mathematical expressions**

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Expressions and their meaning

Haskell teaches you
to **think denotationally** in terms of **expressions** and their **meanings**
instead of **statements** and an **abstract machine**.

This is why Haskell makes you a better programmer:
you *separate* your mental model
from the underlying execution model, ... abstract machine
so you can more easily identify common patterns
between diverse programming languages and problem domains.

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Haskell expression

the distinction between **statements** and **expressions**
in imperative languages

`x = 2 + 2;`

the `x = ...;` part being a **statement**

the `2 + 2` part being an **expression**.

The **body** of a Haskell function is

always one single expression

although you can split that one expression apart for convenience

<https://stackoverflow.com/questions/63144227/what-is-an-expression-in-haskell>

Haskell expression

So if you want to "do more than one thing",
which is an **imperative** notion of a **function**
being able to change **global state**,
you solve this with **monads**, like so:

<https://stackoverflow.com/questions/63144227/what-is-an-expression-in-haskell>

Web service examples

Scotty is a [web framework](#) written in Haskell,
which is similar to **Ruby's Sinatra**.

You can install it using the following commands:

```
$ sudo apt-get install cabal-install  
$ cabal update  
$ cabal install scotty
```

You can compile and start the server from the terminal

```
$ runghc hello-world.hs  
Setting phasers to stun... (port 3000) (ctrl-c to quit)
```

<http://shakthimaan.com/posts/2016/01/27/haskell-web-programming/news.html>

hello-world.hs

```
$ runghc hello-world.hs
```

The service will run on port 3000, and you can open localhost:3000 in a browser to see the 'Hello, World!' text.

You can also use **Curl** to make a query to the server.

```
$ sudo apt-get install curl
```

```
$ curl localhost:3000
```

Hello, World!

```
-- hello-world.hs
```

```
{-# LANGUAGE OverloadedStrings #-}
```

```
import Web.Scotty
```

```
main :: IO ()
```

```
main = scotty 3000 $ do
```

```
    get "/" $ do
```

```
        html "Hello, World!"
```

<http://shakthimaan.com/posts/2016/01/27/haskell-web-programming/news.html>

Web service requests and responses

```
{-# LANGUAGE OverloadedStrings #-}
```

```
import Web.Scotty  
import Network.HTTP.Types
```

```
main = scotty 3000 $ do  
  
  get "/" $ do  
    text "This was a GET request!"  
  
  delete "/" $ do  
    html "This was a DELETE request!"  
  
  post "/" $ do  
    text "This was a POST request!"  
  
  put "/" $ do  
    text "This was a PUT request!"  
  
  -- handle GET request on "/" URL  
  -- send 'text/plain' response  
  
  -- handle DELETE request on "/" URL  
  -- send 'text/html' response  
  
  -- handle POST request on "/" URL  
  -- send 'text/plain' response  
  
  -- handle PUT request on "/" URL  
  -- send 'text/plain' response
```

<https://dev.to/parambirs/how-to-write-a-haskell-web-servicefrom-scratch---part-3-5en6>

Overloaded Strings

{-# LANGUAGE OverloadedStrings #-}

is called a **language pragma** and
extends the language with nice features.

In this case, **OverloadedStrings** allows us to write a string and
it gets automatically converted to the **string type** we need
(**String**, **ByteString**, or **Text**).

{-# LANGUAGE OverloadedStrings #-}

<https://www.stackbuilders.com/blog/getting-started-with-haskell-projects-using-scotty/>

Entry function **scotty**

scotty is the entry function
that **Scotty** defines for running an application.

The first **parameter** is the **port** that we want it to run in, and
the rest is the **application**,
which looks like a **list of routes** and **handlers**.

For now, we only have one route (the root) and a **handler**,
which is a **GET** and returns an **HTML string** with a **title**.

```
scotty 3000 $  
get "/" $  
html "<h1>Shortener</h1>"
```

<https://www.stackbuilders.com/blog/getting-started-with-haskell-projects-using-scotty/>

Named and unnamed parameters

-- named parameters:

```
get "/askfor/:word" $ do
  w <- param "word"
  html $ mconcat ["<h1>You asked for ", w, ", you got it!</h1>" ]
```

-- unnamed parameters from a query string or a form:

```
post "/submit" $ do
  -- e.g. http://server.com/submit?name=somename
  name <- param "name"
  text name
```

<https://dev.to/parambirshow-to-write-a-haskell-web-servicefrom-scratch---part-3-5en6>

Haskell expression in scotty examples (1)

```
{-# LANGUAGE OverloadedStrings #-}  
module Main (main) where  
import Web.Scotty  
  
main :: IO ()  
main = scotty 3000 $  
    get("/:who") $ do  
        who <- param "who"  
        text ("Beam " <> who <> " up, Scotty!")
```

```
Ghci> [1,2,3] <> [4,5,6]          -- concatenation  
[1,2,3,4,5,6]
```

<https://stackoverflow.com/questions/63144227/what-is-an-expression-in-haskell>

Haskell expression in scotty examples (2)

Here, **main's body** (a monadic action, not a function) is a single expression, **scotty 3000 (...)**.

While the **linebreak1** after **scotty 3000 \$** doesn't carry meaning and only makes the code look nicer,

the **linebreak2** in the **do** block actually
reduces multiple actions into one expression
via **syntactic sugar**.

```
main :: IO ()  
main = scotty 3000 $ -- linebreak1  
    get "/:who" $ do -- linebreak2  
        who <- param "who"  
        Text ("..." <> who <> "...")
```

<https://stackoverflow.com/questions/63144227/what-is-an-expression-in-haskell>

Haskell expression in scotty examples (3)

So while it may seem that this **event handler** does two things things:

- (1) **param "who"**
- (2) **text (...)**

it is still one expression equivalent to this:

```
{-# LANGUAGE OverloadedStrings #-}

module Main (main) where
import Web.Scotty

main :: IO ()
main = scotty 3000 $ 
    get "/:who" $ do
        who <- param "who"
        text ("Beam " <> who <> " up, Scotty!")
```

<https://stackoverflow.com/questions/63144227/what-is-an-expression-in-haskell>

Haskell expression in scotty examples (4)

```
main =  
  scotty 3000  
    (get "/:who"  
      (param "who" >>=  
        (who -> text ("Beam " <> who <> " up, Scotty!"))))
```

with **>>=** being the invisible operator between the **do-block lines**.

When expressions begin to grow, this becomes very inconvenient,
so you split parts of them into **sub-expressions**
and give those **names**, e.g. like:

<https://stackoverflow.com/questions/63144227/what-is-an-expression-in-haskell>

Haskell expression in scotty examples (5)

```
main = scotty 3000 handler
```

```
where
```

```
  handler = do
```

```
    get "/:who" getWho
```

```
    post "/" postWho
```

```
  getWho = do
```

```
  ...
```

```
  postWho = do
```

```
  ...
```

But it is essentially equivalent to one big expression.

<https://stackoverflow.com/questions/63144227/what-is-an-expression-in-haskell>

Haskell expression in scotty examples (6)

There are many things in the language beyond function bodies that are not expressions; in the example above, the following are not expressions:

- `{-# LANGUAGE OverloadedStrings #-}` (a language pragma)
- `module Main (main) where` (a module, export list)
- `import Web.Scotty` (an import declaration)
- `main :: IO ()` (a type signature)
- `main =` (a top declaration, or a value binding)

<https://stackoverflow.com/questions/63144227/what-is-an-expression-in-haskell>

Haskell expression in scotty examples (7)

import Web.Scotty could be called a kind of **statement**,
since *grammatically* it's in **imperative form**,
but if we're going to be imprecise,
It would be ok to call them all **declarations**.

More interestingly, in Haskell you have
both an **expression language**
at the **value level** and one at the **type level**.

So **IO ()** isn't a **value expression**, but it's a **type expression**.
If you had the ability to mix those two expression languages up,
you'd have **dependent types**.

- **{-# LANGUAGE OverloadedStrings #-}**
(a language pragma)
- **module Main (main) where**
(a module, export list)
- **import Web.Scotty**
(an import declaration)
- **main :: IO ()**
(a type signature)
- **main =**
(a top declaration, or a value binding)

<https://www.haskell.org/tutorial/goodies.html>

Lazy evaluation

Operational semantics

Operational semantics (1)

It is one of the key properties of
purely functional languages like Haskell
that a **direct mathematical interpretation** like "1+9 denotes 10"
carries over to **functions**, too:

in essence, the **denotation** of a program of type **Integer -> Integer**
is a **mathematical function $Z \rightarrow Z$** between integers.

https://en.wikibooks.org/wiki/Haskell/Denotational_semantics

Operational semantics (2)

While we will see that this expression needs refinement generally, to include non-termination,

the situation for **imperative languages** is clearly worse:
a **procedure** with that type denotes something
that changes the **state** of a machine in possibly unintended ways.

Imperative languages are tightly tied to **operational semantics**
which describes their way of execution on a machine.

https://en.wikibooks.org/wiki/Haskell/Denotational_semantics

Operational semantics (3)

It is possible to define a **denotational semantics** for **imperative programs** and to use it to reason about such programs, but the semantics often has **operational nature** and sometimes must be extended in comparison to the **denotational semantics** for **functional languages**.[

https://en.wikibooks.org/wiki/Haskell/Denotational_semantics

Operational semantics (4)

In contrast, the meaning of **purely functional languages** is by **default** completely **independent** from their **way of execution**.

The Haskell98 standard even goes as far as to specify only Haskell's **non-strict denotational semantics**, leaving open how to implement them.

https://en.wikibooks.org/wiki/Haskell/Denotational_semantics

Operational semantics (5)

The real quantity we're interested in formally describing is **expressions** in programming languages.

A programming language **semantics** is described by the **operational semantics** of the language.

The **operational semantics** can be thought of as a description of an **abstract machine** which operates over the **abstract terms** of the programming language in the same way that a virtual machine might operate over instructions.

http://dev.stephendiehl.com/fun/004_type_systems.html

Operational semantics (6)

Denotational semantics for a language provides a **function** that translates from **program syntax** into **mathematical objects** like sets, functions, lists or even some other programming language

- a denotational semantics acts like a **compiler**

Operational semantics works by rewriting or executing programs **step-by-step**

- it uses only one program syntax to explain how a program runs

<https://www.cs.princeton.edu/~dpw/cos441-11/notes/slides13-lambda-calc.pdf>

Operational semantics (7)

As languages become more complicated, it is often easier to define operational semantics than **denotational semantics**

- it requires less math to do so
- but you might not be able to prove particularly strong theorems using the semantics

<https://www.cs.princeton.edu/~dpw/cos441-11/notes/slides13-lambda-calc.pdf>

Operational semantics (8)

The **operational library** makes it easy to implement monads with tricky **control flow**.

This is very useful for:

- writing **web applications** in a **sequential** style,
- programming **games** with a **uniform interface** for human and AI players and easy replay,
- implementing fast **parser monads**,
- designing **monadic DSLs**, etc.

Embedded Domain Specific Language means that you embed a **Domain specific language** in a language like Haskell.

<https://apfelmus.nfshost.com/articles/operational-monad.html>

Operational semantics (9)

For instance, to write a [web application](#)

where the user is guided through a [sequence of tasks](#) ("wizard").

To [structure](#) your application, you can use a [custom monad](#)

that supports an instruction [askUserInput :: CustomMonad UserInput](#).

This command [sends a web form](#) to the user
and [returns a result](#) when he submits the form.

However, you [don't want](#) your server to [block](#)
[while waiting](#) for the user, so you have to [suspend](#) the computation
and [resume](#) it at some later point.

tricky to implement

This library makes it easy.

<https://apfelmus.nfshost.com/articles/operational-monad.html>

Operational semantics (10)

The idea is to identify a set of primitive instructions and to specify their **operational semantics**.

Then, the library makes sure that the monad laws hold automatically.

In the web application example,
the primitive instruction would be **AskUserInput**.

Any monad can be implemented in this way.

Ditto for monad transformers.

<https://apfelmus.nfshost.com/articles/operational-monad.html>

Sharing (1)

Sharing means that **temporary data** is physically stored,
if it is used multiple times.

```
let x = sin 2  
in x*x
```

x is used twice as factor in the product x*x.

Due to **referential transparency**, it does not play a role,
whether sin 2 is computed twice or
whether it is computed once and the result is stored and reused.

https://wiki.haskell.org/Lazy_evaluation

Sharing (2)

However, when you write **let** expression,
the **Haskell compiler** will certainly decide to store the result.

This can be **the wrong way**,
if a computation is cheap but its result is huge.

[0..1000000] ++ [0..1000000]

where it is much cheaper to compute the list of numbers
than to store it with full length.

https://wiki.haskell.org/Lazy_evaluation

Sharing (3)

Because the **sharing** property **cannot** be observed in Haskell,
it is **hard** to transfer the sharing property to foreign programs
when you use Haskell as an **Embedded** domain specific language.

You must design a **monad** or
use **unsafePerformIO** hacks, which should be avoided.

https://wiki.haskell.org/Lazy_evaluation

References

- [1] <ftp://ftp.geoinfo.tuwien.ac.at/navratil/HaskellTutorial.pdf>
- [2] <https://www.umiacs.umd.edu/~hal/docs/daume02yaht.pdf>