

Monad P3 : Haskell Expressions (1E)

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Haskell Expressions

Expressions and values

Because Haskell is a **purely functional language**, all **computations** are done via the **evaluation** of **expressions** (**syntactic terms**) to yield **values**

Every **value** has an associated **type**.
(Intuitively, we can think of **types** as **sets of values**.)

Examples of **expressions** include **atomic values**
such as the **integer 5**, the **character 'a'**,
and the **function $\lambda x \rightarrow x+1$** ,
as well as **structured values**
such as the **list $[1,2,3]$** and the **pair $('b',4)$** .

Expressions



Value

Type



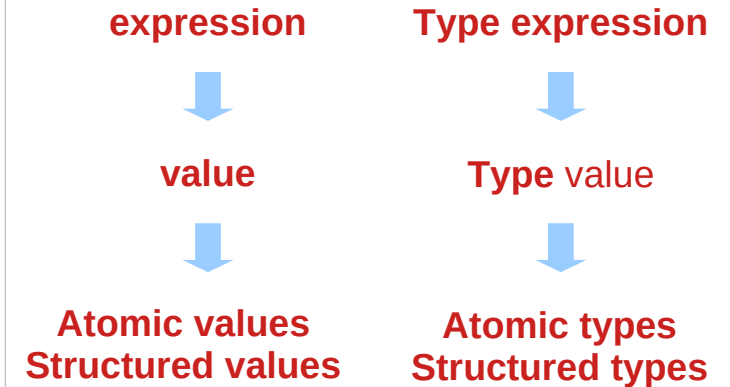
Atomic values
Structured values

<https://www.haskell.org/tutorial/goodies.html>

Type expressions and types

Just as **expressions** denote **values**,
type expressions are **syntactic terms**
that denote **type values** (or just **types**).

Examples of **type expressions** include the **atomic types**
Integer (infinite-precision integers),
Char (characters),
Integer->Integer (functions mapping Integer to Integer),
as well as the **structured types**
[Integer] (homogeneous lists of integers) and
(Char,Integer) (character, integer pairs).



<https://www.haskell.org/tutorial/goodies.html>

First class values

All Haskell values are "**first-class**"

- they may be passed as arguments to functions,
- returned as results,
- placed in data structures, etc.

Haskell types, on the other hand, are not first-class.

<https://www.haskell.org/tutorial/goodies.html>

Typing

Types in a sense describe values, and the association of a **value** with its **type** is called a **typing**.

Using the examples of values and types above, we write **typing** as follows: (the "::" can be read "has type.")

```
5 :: Integer
'a' :: Char
inc :: Integer -> Integer
[1,2,3] :: [Integer]
('b',4) :: (Char,Integer)
```

<https://www.haskell.org/tutorial/goodies.html>

Function definition and declaration

Functions in Haskell are normally defined by a **series of equations**.
For example, the **function `inc`** can be defined by the single equation:

```
inc n    = n+1
```

An **equation** is an example of a **declaration**.

Another kind of **declaration** is a **type signature declaration**,
with which we can declare an **explicit typing** for `inc`:

```
inc      :: Integer -> Integer
```

<https://www.haskell.org/tutorial/goodies.html>

Expression evaluation =>

when we wish to indicate that an **expression e1 evaluates**, or "**reduces**," to *another expression* or **value e2**, we will write:

e1 => e2

For example, note that:

inc (inc 3) => 5

<https://www.haskell.org/tutorial/goodies.html>

Statements vs Expressions

Many programming languages differentiate **statements** from **expressions**.

Statement: What code does

Expression: What code is

can think the term "**statement**" very broadly to refer to anything that is not an **expression** or **type declaration**.

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Imperative vs functional languages

statements vs. **expressions** closely parallels

imperative languages vs. **functional languages**:

Imperative: A language that *emphasizes* **statements**

Functional: A language that *emphasizes* **expressions**

C lies at one end of the spectrum (imperative),
relying heavily on **statements** to accomplish everything.

Haskell lies at the exact opposite extreme (functional),
using **expressions** heavily:

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Statement examples in the imperative language C

```
#include <stdio.h>

int main(int argc, char *argv[]) {
    int elems[5] = {1, 2, 3, 4, 5};    // statement

    int total = 0;
    int i;

    for (i = 0; i < 5; i++) {         // statement
        total += elems[i];           // statement
    }
    printf("%d\n", total);           // statement

    return 0;
}
```

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Expression examples in the functional language Haskell (1)

everything in Haskell is an **expression**,
and even **statements** are **expressions**.

```
main = print (sum [1..5])      -- Expression
```

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Expression examples in the functional language Haskell (2)

For example, the following code might appear to be a traditional imperative-style sequence of statements:

```
main = do
  putStrLn "Enter a number:"      -- Statement?
  str <- getLine                  -- Statement?
  putStrLn ("You entered: " ++ str) -- Statement?
```

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Expression examples in the functional language Haskell (3)

but **do** notation is merely **syntactic sugar**
for nested applications of (**>>=**), which is itself nothing more than
an infix higher-order function:

```
main =  
  putStrLn "Enter a number:" >>= (\_ ->      -- Expression  
    getLine                          >>= (\str -> -- Sub-expression  
      putStrLn ("You entered: " ++ str) )) -- Sub-expression
```

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Statement-as-expression

In Haskell, "**statements**" are actually **nested expressions**, and **sequencing statements** just builds larger and larger **expressions**.

This **statement-as-expression** paradigm promotes consistency and prevents arbitrary language limitations, such as Python's restriction of lambdas to single statements.

In Haskell, you cannot limit the number of statements a **term** uses any more than you can limit the number of **sub-expressions**.

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Monads

do notation works for more than just **IO**.

Any **type** that implements the **Monad class** can be "sequenced" in **statement form**, as long as it supports the following two operations:

class Monad m where

(>>=) :: m a -> (a -> m b) -> m b

return :: a -> m a

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Statement-like syntax using monads

This provides a uniform interface for translating imperative **statement-like** syntax into **expressions** under the hood.

For example, the **Maybe** type implements the Monad class:

```
data Maybe a = Nothing | Just a
```

```
instance Monad Maybe where
```

```
  m >= f = case m of
```

```
    Nothing -> Nothing
```

```
    Just a -> f a
```

```
  return = Just
```

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

do notation using monads

This lets you assemble **Maybe-based** computations using **do** notation

example :: Maybe Int

example = **do**

x <- **Just 1**

y <- **Nothing**

return (x + y)

example =

Just 1 >>= (\x ->

Nothing >>= (\y ->

return (x + y))

The above code desugars to nested calls to (>>=):

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Substitute >>= and return

The compiler then **substitutes** in our definition of (**>>=**) and **return**

```
example = case (Just 1) of
  Nothing -> Nothing
  Just x   -> case Nothing of
    Nothing -> Nothing
    Just y   -> Just (x + y)
```

```
example =
  Just 1 >>= (\x ->
  Nothing >>= (\y ->
  return (x + y) ))
```

instance Monad Maybe where

```
m >>= f = case m of
  Nothing -> Nothing
  Just a   -> f a
```

```
return = Just
```

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Evaluate the outer and inner **case** expression

We can then hand-evaluate this expression to prove that it short-circuits when it encounters `Nothing`:

-- Evaluate the outer `case`

example = case `Nothing` of

`Nothing` -> `Nothing`

`Just y` -> `Just (1 + y)`

-- Evaluate the remaining `case`

example = `Nothing`

example = case (`Just 1`) of

`Nothing` -> `Nothing`

`Just x` -> case `Nothing` of

`Nothing` -> `Nothing`

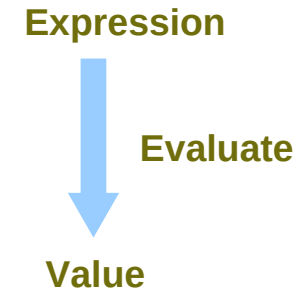
`Just y` -> `Just (x + y)`

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Everything is an expression to be evaluated

Notice that we can evaluate these **Maybe** "statements" without invoking any sort of **abstract machine**.

When everything is an **expression**, **everything** is simple to **evaluate** and does not require understanding or *invoking an execution model*.



FSM not needed for sequencing

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Semantics

Semantics

In fact, the distinction between **statements** and **expressions** also closely parallels another important divide: the difference between **operational semantics** and **denotational semantics**.

Operational semantics:

Translates code to **abstract machine statements**

Denotational semantics:

Translates code to **mathematical expressions**

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Expressions and their meaning

Haskell teaches you
to **think denotationally** in terms of **expressions** and their **meanings**
instead of **statements** and an **abstract machine**.

This is why Haskell makes you a better programmer:
you *separate* your mental model
from the underlying execution model, ... **abstract machine**
so you can more easily identify *common patterns*
between diverse **programming languages** and **problem domains**.

<https://www.haskellforall.com/2013/07/statements-vs-expressions.html>

Haskell expression

the distinction between **statements** and **expressions**
in **imperative languages**

```
x = 2 + 2;
```

the `x = ...;` part being a **statement**

the `2 + 2` part being an **expression**.

The **body** of a **Haskell function** is

always one single expression

although you can split that one expression apart for convenience

<https://stackoverflow.com/questions/63144227/what-is-an-expression-in-haskell>

Haskell expression

So if you want to "do more than one thing", which is an **imperative** notion of a **function** being able to change **global state**, you solve this with **monads**, like so:

<https://stackoverflow.com/questions/63144227/what-is-an-expression-in-haskell>

Web service examples

Scotty is a **web framework** written in Haskell,
which is similar to **Ruby's Sinatra**.

You can install it using the following commands:

```
$ sudo apt-get install cabal-install
```

```
$ cabal update
```

```
$ cabal install scotty
```

You can compile and start the server from the terminal

```
$ runghc hello-world.hs
```

```
Setting phasers to stun... (port 3000) (ctrl-c to quit)
```

<http://shakthimaan.com/posts/2016/01/27/haskell-web-programming/news.html>

hello-world.hs

```
$ runghc hello-world.hs
```

The service will run on [port 3000](#), and you can [open localhost:3000](#) in a browser to see the `Hello, World!` text.

You can also use **Curl** to make a [query](#) to the server.

```
$ sudo apt-get install curl
```

```
$ curl localhost:3000
```

```
Hello, World!
```

```
-- hello-world.hs
{-# LANGUAGE OverloadedStrings #-}

import Web.Scotty

main :: IO ()
main = scotty 3000 $ do
  get "/" $ do
    html "Hello, World!"
```

<http://shakthimaan.com/posts/2016/01/27/haskell-web-programming/news.html>

Web service requests and responses

```
{-# LANGUAGE OverloadedStrings #-}
import Web.Scotty
import Network.HTTP.Types

main = scotty 3000 $ do
  get "/" $ do
    text "This was a GET request!"
  delete "/" $ do
    html "This was a DELETE request!"
  post "/" $ do
    text "This was a POST request!"
  put "/" $ do
    text "This was a PUT request!"
  -- handle GET request on "/" URL
  --   send 'text/plain' response
  -- handle DELETE request on "/" URL
  --   send 'text/html' response
  -- handle POST request on "/" URL
  --   send 'text/plain' response
  -- handle PUT request on "/" URL
  --   send 'text/plain' response
```

<https://dev.to/parambirs/how-to-write-a-haskell-web-servicefrom-scratch---part-3-5en6>

Overloaded Strings

```
{-# LANGUAGE OverloadedStrings #-}
```

is called a **language pragma** and extends the language with nice features.

In this case, **OverloadedStrings** allows us to write a string and it gets automatically converted to the **string type** we need (**String**, **ByteString**, or **Text**).

```
{-# LANGUAGE OverloadedStrings #-}
```

<https://www.stackbuilders.com/blog/getting-started-with-haskell-projects-using-scotty/>

Entry function **scotty**

scotty is the entry function that **Scotty** defines for running an application.

The first **parameter** is the **port** that we want it to run in, and the rest is the **application**, which looks like a **list** of **routes** and **handlers**.

For now, we only have one **route** (the root) and a **handler**, which is a **GET** and returns an **HTML string** with a **title**.

```
scotty 3000 $  
  get "/" $  
    html "<h1>Shortener</h1>"
```

<https://www.stackbuilders.com/blog/getting-started-with-haskell-projects-using-scotty/>

Named and unnamed parameters

-- named parameters:

```
get "/askfor/:word" $ do
  w <- param "word"
  html $ mconcat ["<h1>You asked for ", w, ", you got it!</h1>"]
```



-- unnamed parameters from a query string or a form:

```
post "/submit" $ do      -- e.g. http://server.com/submit?name=somename
  name <- param "name"
  text name
```


<https://dev.to/parambirs/how-to-write-a-haskell-web-servicefrom-scratch---part-3-5en6>

Haskell expression in scotty examples (1)

```
{-# LANGUAGE OverloadedStrings #-}
module Main (main) where
import Web.Scotty

main :: IO ()
main = scotty 3000 $
  get "/:who" $ do
    who <- param "who"
    text ("Beam " <> who <> " up, Scotty!")
```

```
Ghci> [1,2,3] <> [4,5,6]      -- concatenation
[1,2,3,4,5,6]
```

<https://stackoverflow.com/questions/63144227/what-is-an-expression-in-haskell>

Haskell expression in scotty examples (2)

Here, **main's body** (a **monadic action**, not a function) is a single expression, **scotty 3000** (...).

While the **linebreak1** after **scotty 3000 \$** doesn't carry meaning and only makes the code look nicer,

the **linebreak2** in the **do** block actually reduces multiple actions into one expression via **syntactic sugar**.

```
main :: IO ()
main = scotty 3000 $      -- linebreak1
      get "/:who" $ do   -- linebreak2
        who <- param "who"
        Text ("..." <> who <> " ...")
```

<https://stackoverflow.com/questions/63144227/what-is-an-expression-in-haskell>

Haskell expression in scotty examples (3)

So while it may seem that this **event handler** does two things things:

- (1) **param** "who"
- (2) **text** (...)

it is still one expression equivalent to this:

```
{-# LANGUAGE OverloadedStrings #-}
module Main (main) where
import Web.Scotty

main :: IO ()
main = scotty 3000 $
  get "/:who" $ do
    who <- param "who"
    text ("Beam " <> who <> " up, Scotty!")
```

<https://stackoverflow.com/questions/63144227/what-is-an-expression-in-haskell>

Haskell expression in scotty examples (4)

```
main =  
  scotty 3000  
    (get "/:who"  
      (param "who" >>=  
        (\who -> text ("Beam " <> who <> " up, Scotty!")))))
```

with `>>=` being the invisible operator between the `do-block` lines.

When expressions begin to grow, this becomes very inconvenient,
so you split parts of them into `sub-expressions`
and give those `names`, e.g. like:

<https://stackoverflow.com/questions/63144227/what-is-an-expression-in-haskell>

Haskell expression in scotty examples (5)

```
main = scotty 3000 handler
  where
    handler = do
      get "/:who" getWho
      post "/" postWho

    getWho = do
      ...
    postWho = do
      ...
```

But it is essentially equivalent to one big expression.

<https://stackoverflow.com/questions/63144227/what-is-an-expression-in-haskell>

Haskell expression in scotty examples (6)

There are many things in the language beyond function bodies that are not expressions; in the example above, the following are not expressions:

- `{-# LANGUAGE OverloadedStrings #-}` (a language pragma)
- `module Main (main) where` (a module, export list)
- `import Web.Scotty` (an import declaration)
- `main :: IO ()` (a type signature)
- `main =` (a top declaration, or a value binding)

<https://stackoverflow.com/questions/63144227/what-is-an-expression-in-haskell>

Haskell expression in scotty examples (7)

`import Web.Scotty` could be called a kind of **statement**, since *grammatically* it's in **imperative form**, but if we're going to be imprecise, it would be ok to call them all **declarations**.

More interestingly, in Haskell you have both an **expression language** at the **value level** and one at the **type level**.

So `IO ()` isn't a **value expression**, but it's a **type expression**. If you had the ability to mix those two expression languages up, you'd have **dependent types**.

- `{-# LANGUAGE OverloadedStrings #-}`
(a language pragma)
- `module Main (main) where`
(a module, export list)
- `import Web.Scotty`
(an import declaration)
- `main :: IO ()`
(a type signature)
- `main =`
(a top declaration, or a value binding)

<https://www.haskell.org/tutorial/goodies.html>

Lazy evaluation

Operational semantics

Operational semantics (1)

It is one of the key properties of **purely functional languages** like Haskell that a **direct mathematical interpretation** like "1+9 denotes 10" carries over to **functions**, too:

in essence, the **denotation** of a program of type **Integer -> Integer** is a **mathematical function** $\mathbf{Z} \rightarrow \mathbf{Z}$ between integers.

https://en.wikibooks.org/wiki/Haskell/Denotational_semantics

Operational semantics (2)

While we will see that this expression needs **refinement** generally, to include non-termination,

the situation for **imperative languages** is clearly worse: a **procedure** with that type denotes something that changes the **state** of a machine in possibly unintended ways.

Imperative languages are tightly tied to **operational semantics** which describes their way of execution on a machine.

https://en.wikibooks.org/wiki/Haskell/Denotational_semantics

Operational semantics (3)

It is possible to define a **denotational semantics** for **imperative programs** and to use it to reason about such programs, but the semantics often has **operational nature** and sometimes must be extended in comparison to the **denotational semantics** for **functional languages**.]

https://en.wikibooks.org/wiki/Haskell/Denotational_semantics

Operational semantics (4)

In contrast, the meaning of **purely functional languages** is by **default** completely independent from their way of execution.

The Haskell98 standard even goes as far as to specify only Haskell's **non-strict denotational semantics**, leaving open how to implement them.

https://en.wikibooks.org/wiki/Haskell/Denotational_semantics

Operational semantics (5)

The real quantity we're interested in formally describing is **expressions** in programming languages.

A programming language **semantics** is described by the **operational semantics** of the language.

The **operational semantics** can be thought of as a description of an **abstract machine** which operates over the **abstract terms** of the programming language in the same way that a virtual machine might operate over instructions.

http://dev.stephendiehl.com/fun/004_type_systems.html

Operational semantics (6)

Denotational semantics for a language provides a **function** that translates from **program syntax** into **mathematical objects** like sets, functions, lists or even some other programming language

– a denotational semantics acts like a **compiler**

Operational semantics works by rewriting or executing programs **step-by-step**

– it uses only one program syntax to explain how a program runs

<https://www.cs.princeton.edu/~dpw/cos441-11/notes/slides13-lambda-calc.pdf>

Operational semantics (7)

As languages become more complicated, it is often easier to define **operational semantics** than **denotational semantics**

- it requires less math to do so
- but you might not be able to prove particularly strong theorems using the semantics

<https://www.cs.princeton.edu/~dpw/cos441-11/notes/slides13-lambda-calc.pdf>

Operational semantics (8)

The **operational library** makes it easy to implement monads with tricky **control flow**.

This is very useful for:

writing **web applications** in a **sequential** style,
programming **games** with a **uniform interface**
for human and AI players and easy replay,
implementing fast **parser monads**,
designing **monadic DSLs**, etc.

Embedded Domain Specific Language means
that you embed a **Domain specific language** in a language like Haskell.

<https://apfelmus.nfshost.com/articles/operational-monad.html>

Operational semantics (9)

For instance, to write a [web application](#) where the user is guided through a [sequence of tasks](#) ("wizard"). To [structure](#) your application, you can use a [custom monad](#) that supports an instruction `askUserInput :: CustomMonad UserInput`.

This command sends a web form to the user and returns a result when he submits the form.

However, you don't want your server to block while waiting for the user, so you have to suspend the computation and resume it at some later point.

tricky to implement

This library makes it easy.

<https://apfelmus.nfshost.com/articles/operational-monad.html>

Operational semantics (10)

The idea is to identify a set of primitive instructions and to specify their **operational semantics**.

Then, the library makes sure that the monad laws hold automatically.

In the web application example, the primitive instruction would be **AskUserInput**.

Any monad can be implemented in this way.

Ditto for monad transformers.

<https://apfelmus.nfshost.com/articles/operational-monad.html>

Sharing (1)

Sharing means that **temporary data** is physically stored,
if it is used multiple times.

```
let x = sin 2  
in x*x
```

x is used twice as factor in the product $x*x$.

Due to **referential transparency**, it does not play a role,
whether **sin 2** is computed twice or
whether it is computed once and the result is stored and reused.

https://wiki.haskell.org/Lazy_evaluation

Sharing (2)

However, when you write **let** expression,
the **Haskell compiler** will certainly decide to store the result.

This can be **the wrong way**,
if a computation is cheap but its result is huge.

```
[0..1000000] ++ [0..1000000]
```

where it is much cheaper to compute the list of numbers
than to store it with full length.

https://wiki.haskell.org/Lazy_evaluation

Sharing (3)

Because the **sharing** property **cannot** be observed in Haskell, it is **hard** to transfer the sharing property to foreign programs when you use Haskell as an **Embedded** domain specific language.

You must design a **monad** or use **unsafePerformIO** hacks, which should be avoided.

https://wiki.haskell.org/Lazy_evaluation

References

- [1] <ftp://ftp.geoinfo.tuwien.ac.at/navratil/HaskellTutorial.pdf>
- [2] <https://www.umiacs.umd.edu/~hal/docs/daume02yaht.pdf>