

Critical Section

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Useful Links

Uoc.gr

<http://www.csd.uoc.gr/~hy586/material/lectures/cs586-Section2.pdf>

Until Peterson's algorithm

UWisc, EECE6354,

<http://pages.cs.wisc.edu/~remzi/OSTEP/threads-locks.pdf>

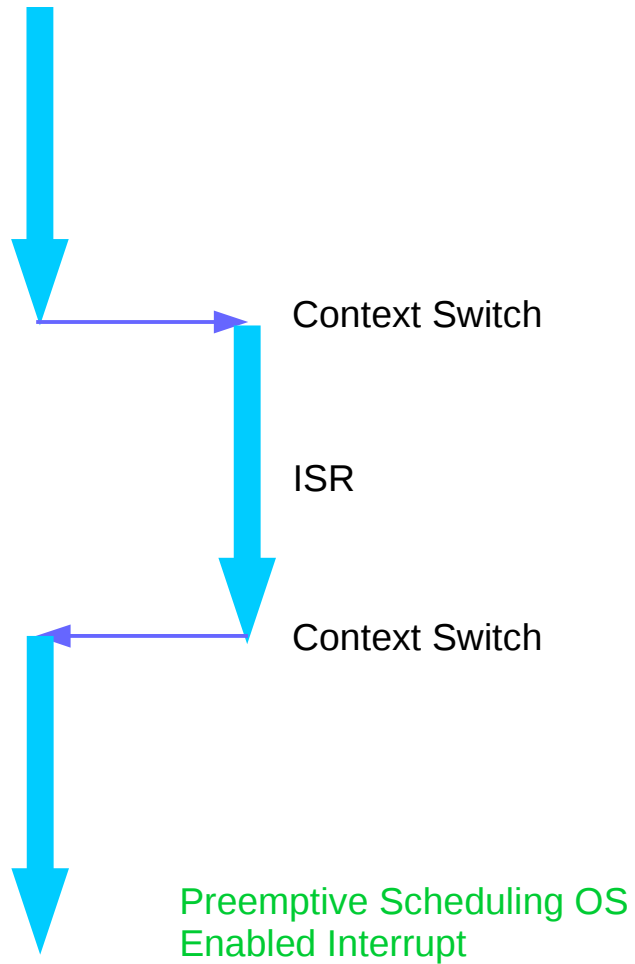
Multitasking, RTK,

Uregina,

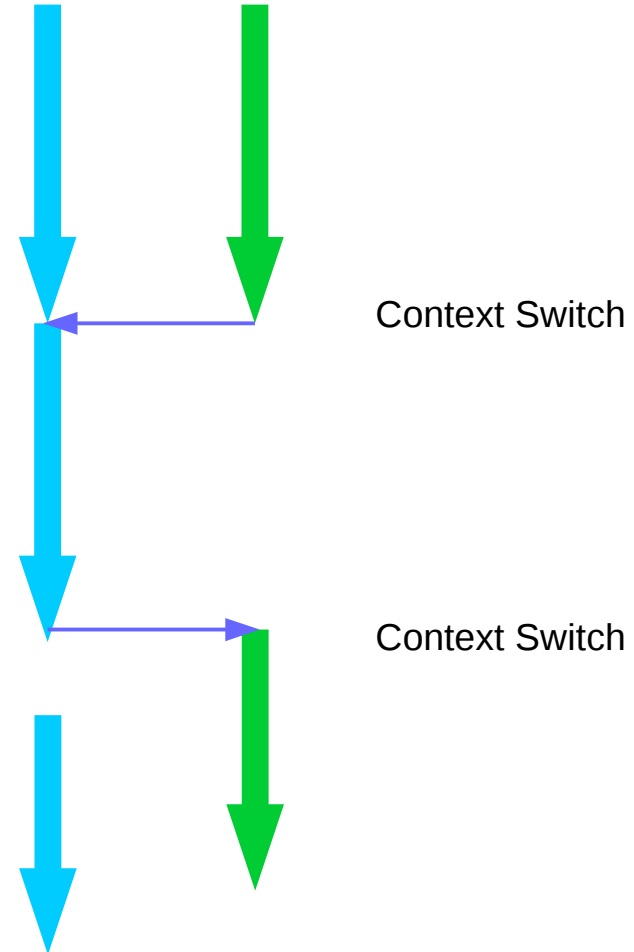
<http://www2.cs.uregina.ca/~hamilton/courses/330/notes/synchro/node3.html>

Single Processor v.s. Multi-processor

Single Processor



Multi-Processor



Try 1 – working examples

lock = false;

```
P0:  
while (true) {  
  while (lock) { ; }  
  lock = true;  
  critical section  
  lock = false;  
  non-critical section  
}
```

```
P1:  
while (true) {  
  while (lock) { ; }  
  lock = true;  
  critical section  
  lock = false;  
  non-critical section  
}
```

Wait for the releasing
of the **lock**
Set the **lock** before
entering the CS.

P0 sets / clears **lock**
P1 sets / clears **lock**

```
P0:  
while (true) {  
  while (lock) { ; }  
  lock = true;  
  critical section  
  lock = false;  
  non-critical section  
}
```

```
P1:  
while (true) {  
  while (lock) { ; }  
  lock = true;  
  critical section  
  lock = false;  
  non-critical section  
}
```

Try 2 – working examples

```
want[2] = {false, false};
```

```
P0:  
while (true) {  
  while (want[1]) { ; }  
  want[0] = true;  
  critical section  
  want[0] = false;  
  non-critical section  
}
```

```
P1:  
while (true) {  
  while (want[0]) { ; }  
  want[1] = true;  
  critical section  
  want[1] = false;  
  non-critical section  
}
```

```
P0:  
while (true) {  
  while (want[1]) { ; }  
  want[0] = true;  
  critical section  
  want[0] = false;  
  non-critical section  
}
```

```
P1:  
while (true) {  
  while (want[0]) { ; }  
  want[1] = true;  
  critical section  
  want[1] = false;  
  non-critical section  
}
```

Wait for the releasing of the `want[J]`
Set the `want[I]` before entering the CS.

P0 sets / clears `want[0]`
P1 sets / clears `want[1]`

If you want, I'll wait

Try 3 – working examples

turn = 1;

```
P0:
while (true) {
  while (turn != 0) { ; }
  critical section
  turn = 1;
  non-critical section
}
```

```
P1:
while (true) {
  while (turn != 1) { ; }
  critical section
  turn = 0;
  non-critical section
}
```

Wait for my **turn**
Switch the **turn** after
exiting the CS.

P0 sets **turn** to 1
P1 sets **turn** to 0

I'll wait for my turn

```
P0:
while (true) {
  while (turn != 0) { ; }
  critical section
  turn = 1;
  non-critical section
}
```

```
P1:
while (true) {
  while (turn != 1) { ; }
  critical section
  turn = 0;
  non-critical section
}
```

Try 4 – working examples

```
want[2] = {false, false};
```

```
P0:  
while (true) { want[0] = true;  
  while (want[1]) {  
    want[0] = false;  
    while (want[1]) {;}  
    want[0] = true;  
  } critical section  
  want[0] = false;  
  non-critical section  
}
```

```
P1:  
while (true) { want[1] = true;  
  while (want[0]) {  
    want[1] = false;  
    while (want[0]) {;}  
    want[1] = true;  
  } critical section  
  want[1] = false;  
  non-critical section  
}
```

```
P1:  
while (true) { want[1] = true;  
  while (want[0]) {  
    want[1] = false;  
    while (want[0]) {;}  
    want[1] = true;  
  } critical section  
  want[1] = false;  
  non-critical section  
}
```

```
P0:  
while (true) { want[0] = true;  
  while (want[1]) {  
    want[0] = false;  
    while (want[1]) {;}  
    want[0] = true;  
  } critical section  
  want[0] = false;  
  non-critical section  
}
```

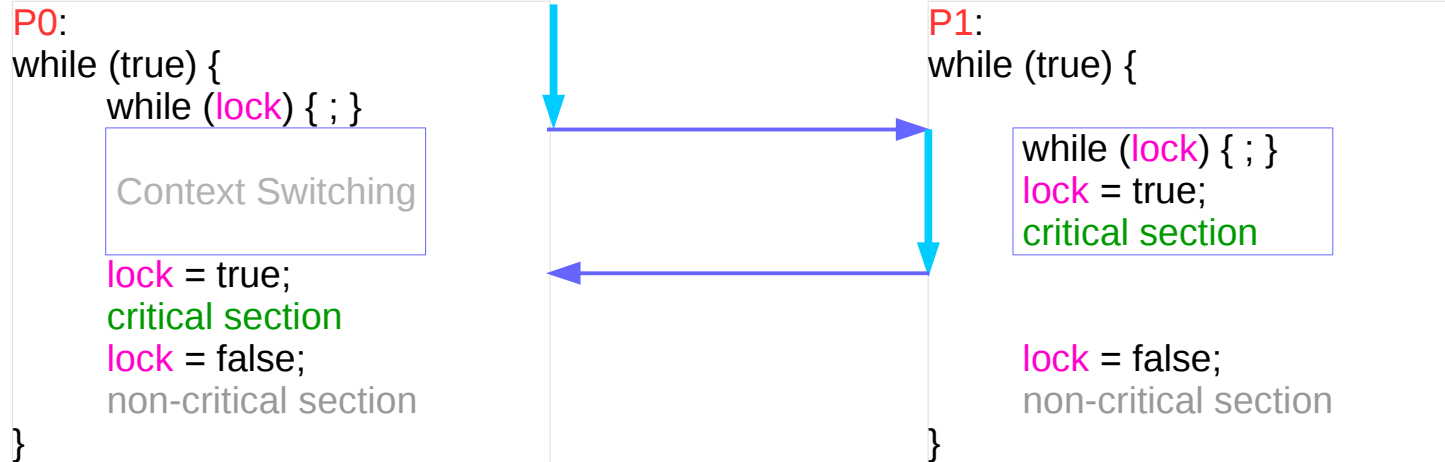
Wait for the releasing of the `want[J]`
While waiting, temporarily clears `want[I]`
Set the `want[I]` before entering the CS.

P0 sets / clears `want[0]`
P1 sets / clears `want[1]`

*If you want, go ahead
I'll wait and do*

Try 1 – Not working examples

```
lock = false;
```



after seeing false `lock`,
a context switch occurs
plan to enter its `critical section`

after seeing false `lock`,
entered its `critical section`
but before releasing the `lock`
a context switch occurs

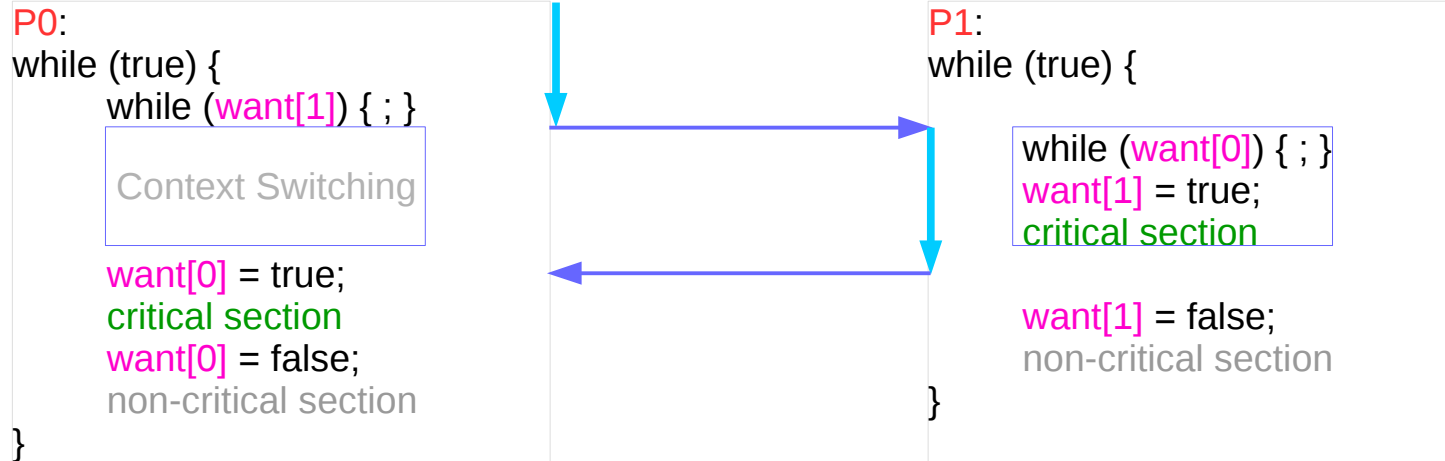
Both P0 & P1 are in the `critical section`

Wait for the releasing
of the `lock`
Set the `lock` before
entering the CS.

P0 sets / clears `lock`
P1 sets / clears `lock`

Try 2 – Not working examples

```
want[2] = {false, false};
```



after seeing false **want[1]**,
a context switch occurs
plan to enter its **critical section**

after seeing false **want[0]**,
entered its **critical section**
but before releasing the **lock**
a context switch occurs

Both **P0** & **P1** are in the **critical section**

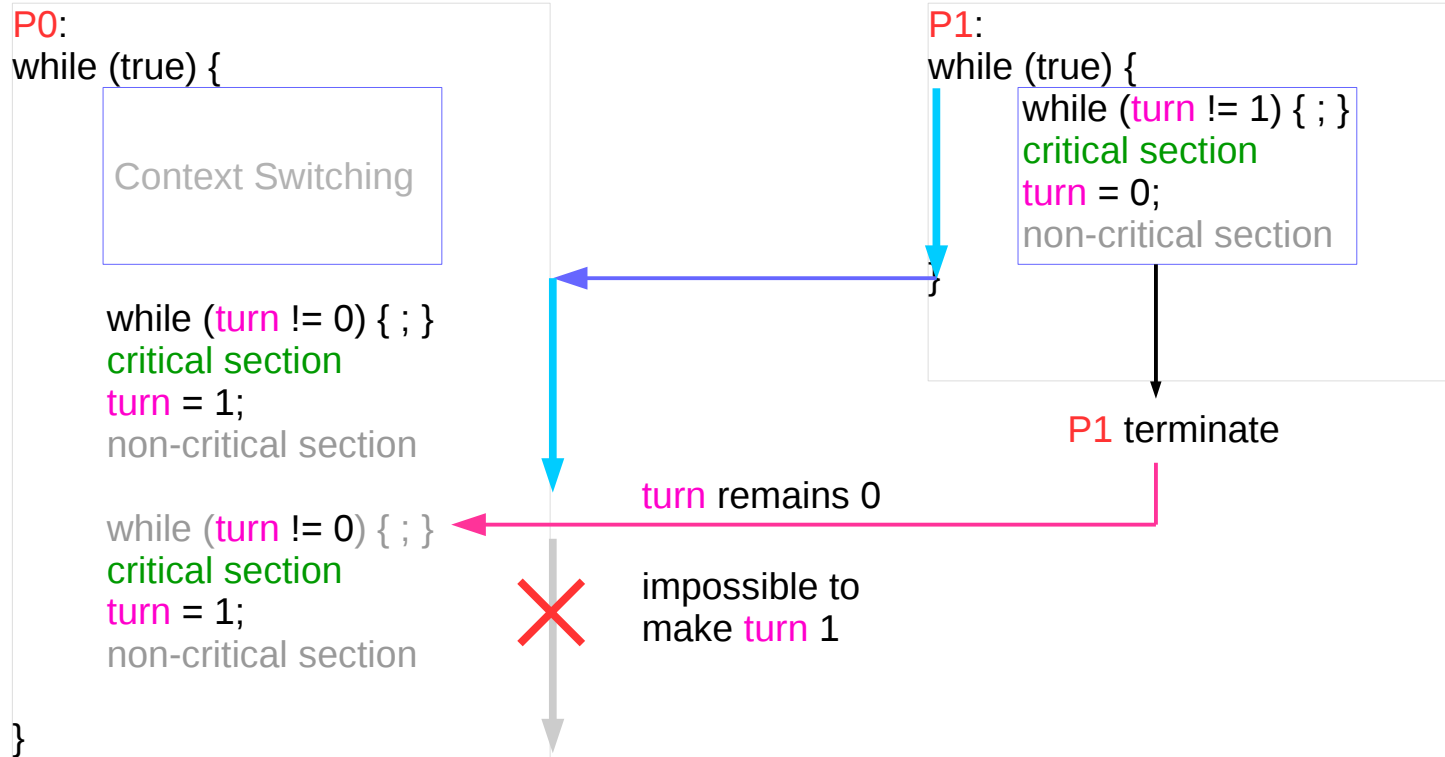
Wait for the releasing of
the **want[J]**
Set the **want[I]** before
entering the CS.

P0 sets / clears **want[0]**
P1 sets / clears **want[1]**

If you want, I'll wait

Try 3 – Not working examples

turn = 1;



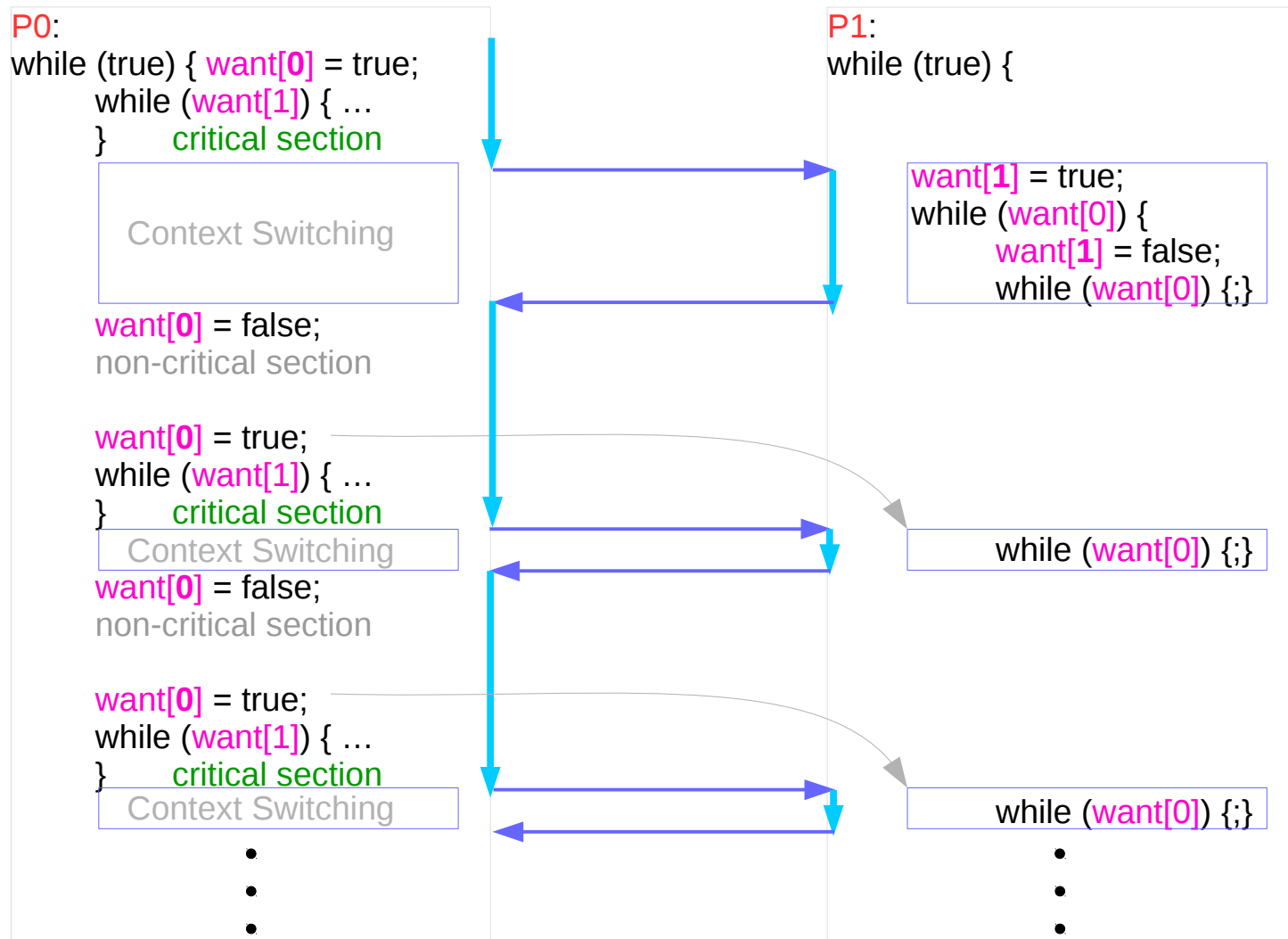
Wait for my **turn**
Switch the **turn** after
exiting the CS.

P0 sets **turn** to 1
P1 sets **turn** to 0

I'll wait for my turn

Try 4 – Not working examples

```
want[2] = {false, false};
```



Wait for the releasing of the want[J]
While waiting, temporarily clears want[I]
Set the want[I] before entering the CS.

P0 sets / clears want[0]
P1 sets / clears want[1]

*If you want, go ahead
I'll wait and do*

Paterson's Algorithm

```
want[2] = {false, false};    turn;
```

```
P0:  
while (true) {  
    want[0] = true;  
    turn = 1;  
    while  
        (want[1] && turn==1) { ; }  
    critical section  
    want[0] = false;  
    non-critical section  
}
```

```
P1:  
while (true) {  
    want[1] = true;  
    turn = 0;  
    while  
        (want[0] && turn==0) { ; }  
    critical section  
    want[1] = false;  
    non-critical section  
}
```

busy waiting until context sw

Wait for the releasing of either **want[J]** or **J turn**. Before waiting, first yield to **J turn**

P0 sets / clears **want[0]**
P1 sets / clears **want[1]**

P0 sets **turn** to 1
P1 sets **turn** to 0

<http://www2.cs.uregina.ca/~hamilton/courses/330/notes/synchro/node3.html>

Paterson's Algorithm

```
want[2] = {false, false};    turn;
```

```
P0:  
while (true) {  
    want[0] = true;  
    turn = 1;  
    while  
        (want[1] && turn==1) { ; }  
    critical section  
    want[0] = false;  
    non-critical section  
}
```

```
P1:  
while (true) {  
    want[1] = true;  
    turn = 0;  
    while  
        (want[0] && turn==0) { ; }  
    critical section  
    want[1] = false;  
    non-critical section  
}
```

busy waiting until context sw

Wait for the releasing of either **want[J]** or **J turn**. Before waiting, first yield to **J turn**

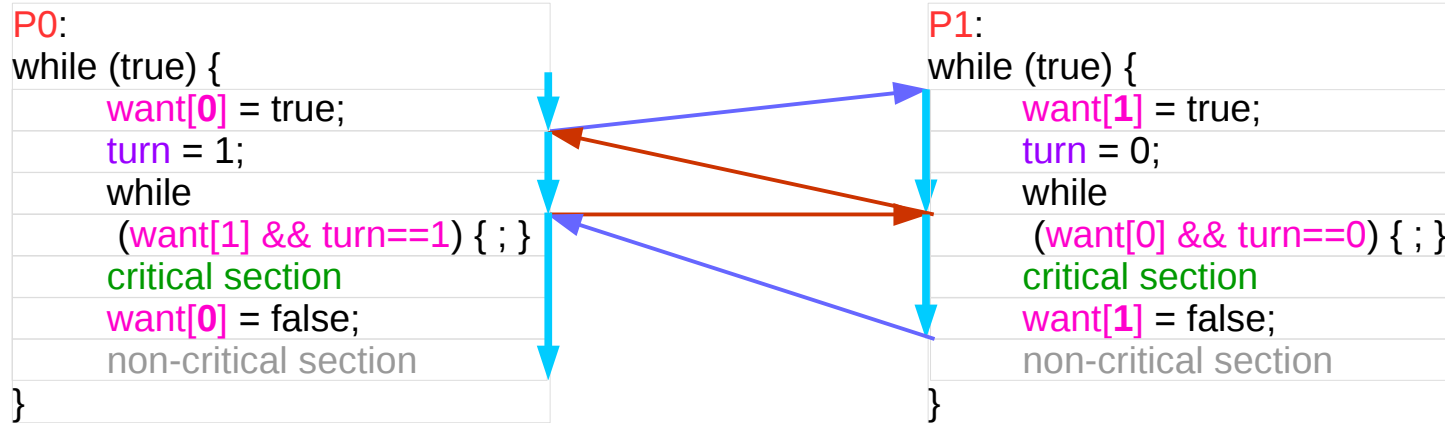
P0 sets / clears **want[0]**
P1 sets / clears **want[1]**

P0 sets **turn** to 1
P1 sets **turn** to 0

<http://www2.cs.uregina.ca/~hamilton/courses/330/notes/synchro/node3.html>

Paterson's Algorithm

```
want[2] = {false, false};    turn;
```



busy waiting until context sw

busy waiting until context sw

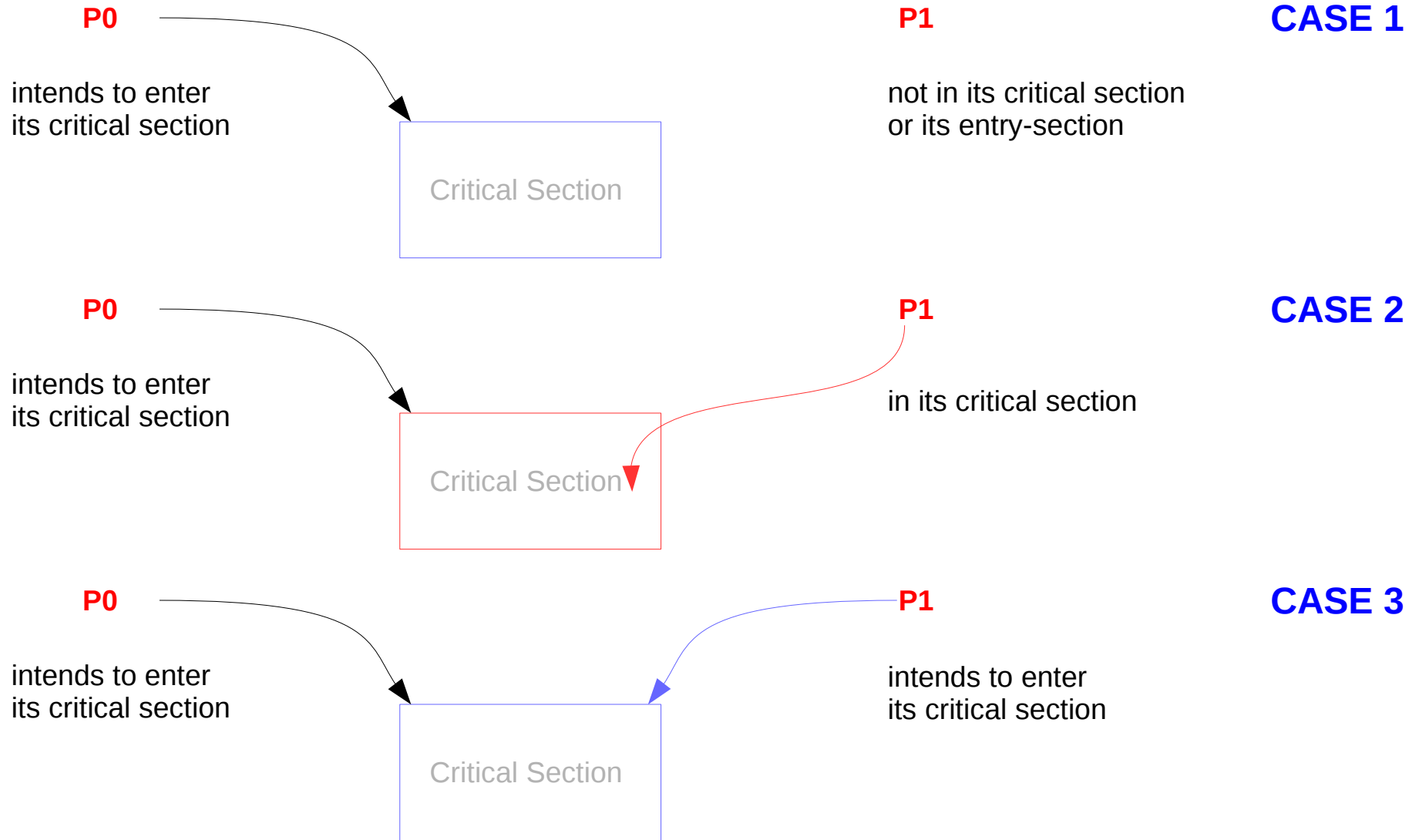
Wait for the releasing of either $want[J]$ or J turn. Before waiting, first yield to J turn

P0 sets / clears $want[0]$
P1 sets / clears $want[1]$

P0 sets $turn$ to 1
P1 sets $turn$ to 0

<http://www2.cs.uregina.ca/~hamilton/courses/330/notes/synchro/node3.html>

Test Cases



Case 1

```
want[2] = {false, false};    turn;
```

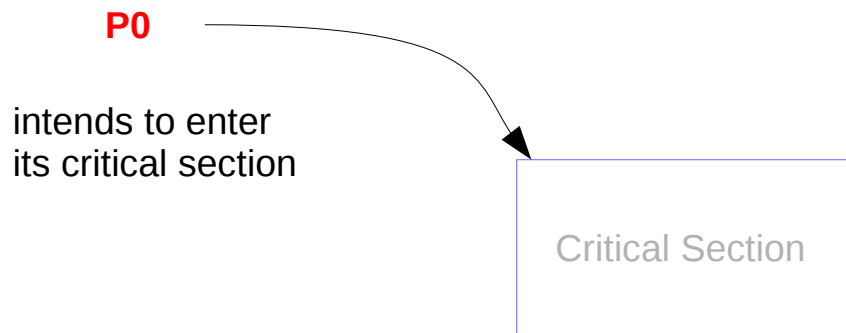
```
P0:  
while (true) {  
    want[0] = true;  
    turn = 1;  
    while  
        (want[1] && turn==1) { ; }  
    critical section  
    want[0] = false;  
    non-critical section  
}
```

```
P1:
```

Wait for the releasing of either **want[J]** or **J turn**. Before waiting, first yield to **J turn**

P0 sets / clears **want[0]**
P1 sets / clears **want[1]**

P0 sets **turn** to 1
P1 sets **turn** to 0



P1

not in its critical section or its entry-section

CASE 1

Case 2

```
want[2] = {false, false};    turn;
```

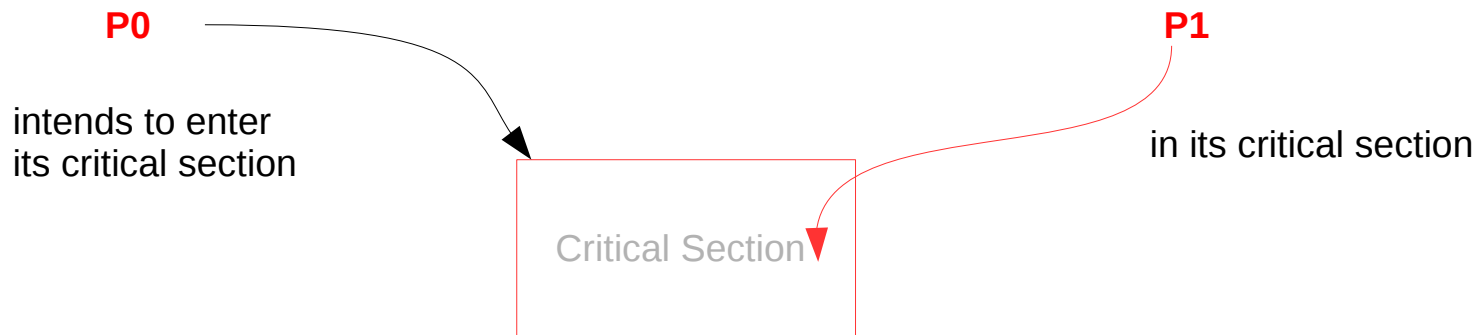
```
P0:  
while (true) {  
    want[0] = true;  
    turn = 1;  
  
    while  
        (want[1] && turn==1) { ; }  
    critical section  
    want[0] = false;  
    non-critical section  
}
```

```
P1:  
while (true) {  
    want[1] = true;  
    turn = 0;  
    while  
        (want[0] && turn==0) { ; }  
    critical section  
    want[1] = false;  
    non-critical section  
}
```

Wait for the releasing of either **want[J]** or **J turn**. Before waiting, first yield to **J turn**

P0 sets / clears **want[0]**
P1 sets / clears **want[1]**

P0 sets **turn** to 1
P1 sets **turn** to 0



Case 3

```
want[2] = {false, false};    turn;
```

```
P0:  
while (true) {  
    want[0] = true;  
    turn = 1;  
    while  
        (want[1] && turn==1) { ; }  
    critical section  
    want[0] = false;  
    non-critical section  
}
```

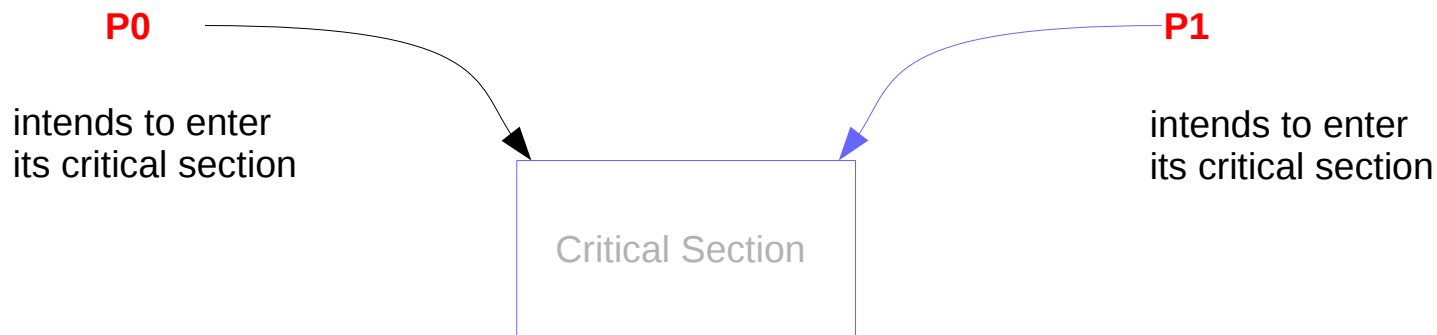
```
P1:  
while (true) {  
    want[1] = true;  
  
    turn = 0;  
    while  
        (want[0] && turn==0) { ; }  
    critical section  
    want[1] = false;  
    non-critical section  
}
```

Wait for the releasing of either **want[J]** or **J turn**. Before waiting, first yield to **J turn**

P0 sets / clears **want[0]**
P1 sets / clears **want[1]**

P0 sets **turn** to 1
P1 sets **turn** to 0

P1 first executes **turn = J**; waits until **PJ** executes **turn = I** and then enters its critical section. The **PJ** will be the next



CASE 3

References

- [1] <http://crystal.uta.edu/~ylei/cse6324/data/critical-section.pdf>