

FPGA Carry Chain Adder (1A)

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FPGA Carry Chain Cell



$$s_i = (a_i \oplus b_i) \oplus c_i = p_i \oplus c_i$$

$$c_{i+1} = (a_i \cdot b_i) + (a_i \oplus b_i) c_i = \bar{p}_i \cdot g_i + p_i \cdot c_i = \bar{p}_i \cdot a_i + p_i \cdot c_i = \bar{p}_i \cdot b_i + p_i \cdot c_i$$

when $\bar{p}_i = 1$, then $a_i = b_i$

when $g_i = 1$, then $a_i = b_i = 1$

$p(i)$	0	1
0	0	1
1	1	0

$g(i)$	0	1
0	0	0
1	0	1

FPGA Carry Chain Cell



Synthesis of Arithmetic Circuits: FPGA, ASIC and Ebedded Systems, J-P Deschamps et al

FPGA Carry Chain

FPGAs generally contain dedicated computation resources for generating fast adders

The Virtex family programmable arrays include logic gates (**XOR**) and **multiplexers** that along with the general purpose **lookup tables** allow one to build effective carry-chain adders

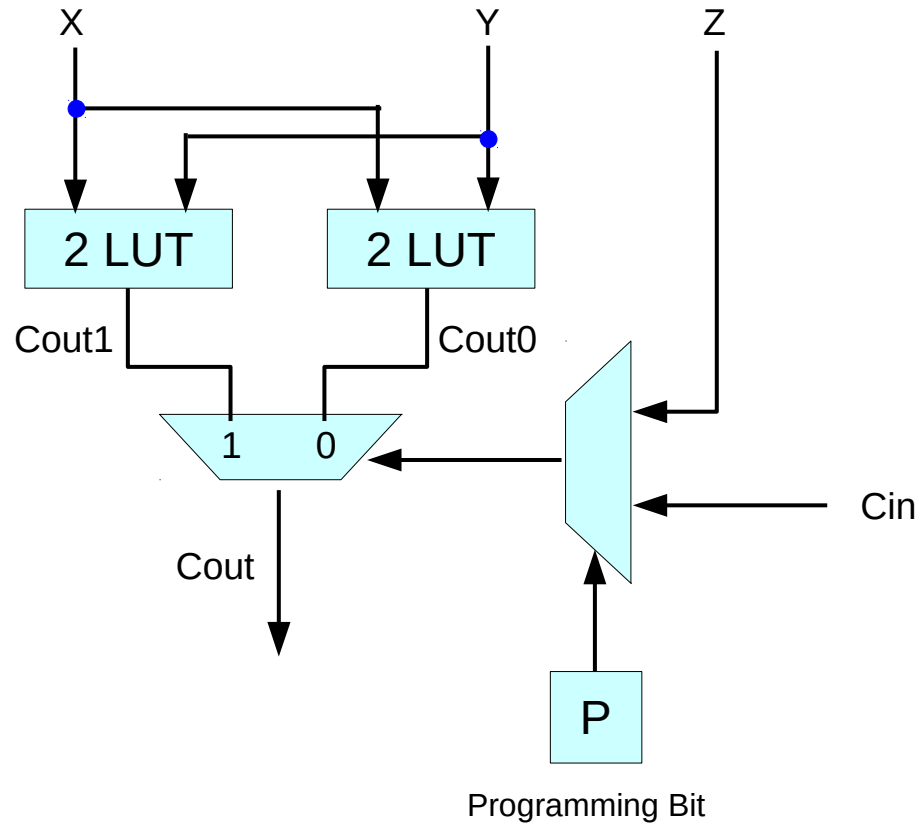
The carry chain is made up of multiplexers belonging to adjacent configurable blocks

the lookup table is used for implementing the exclusive or function

$$p(i) = x(i) \text{ xor } y(i)$$

https://en.wikipedia.org/wiki/Carry-lookahead_adder

FPGA Carry Chain Cell



Cout1, Cout2 : functions of X, Y, Cin

Cout1 = X+Y when Cin=1

Cout0 = X Y when Cin=0

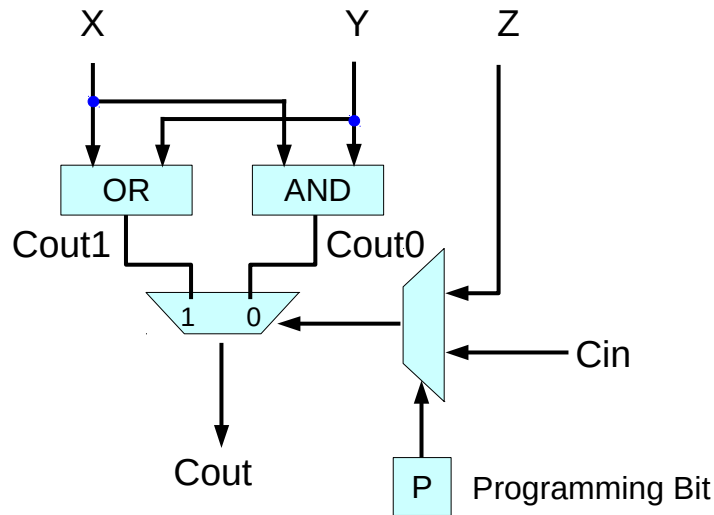
$Cout = (X + Y) Cin + X Y \overline{Cin}$

$Cout = P' Cin + G \overline{Cin} \dots P' = \text{relaxed } P$

Cout1	Cout0	Cout	Name
0	0	0	Kill
0	1	\overline{Cin}	Inverse Propagate
1	0	Cin	Propagate
1	1	1	Generate

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FPGA Carry Chain Cell



X	Y	Cin	$\overline{\text{Cin}}$	$\overline{X} \overline{Y}$
		Cout1	Cout0	
0	0	0	0	$\overline{X} \overline{Y}$
0	1	1	0	$\overline{X} Y$
1	0	1	0	$X \overline{Y}$
1	1	1	1	$X Y$

Cout : functions of X, Y, Cin

$$\text{Cout}(X, Y, 1) = \text{Cout1} = X + Y$$

$$\text{Cout}(X, Y, 0) = \text{Cout0} = X Y$$

$$\text{Cout1} = X + Y \text{ when Cin}=1$$

$$\text{Cout0} = X Y \text{ when Cin}=0$$

$$\text{Cout1} = P' \text{Cin} \dots P' = \text{relaxed } P$$

$$\text{Cout0} = G \overline{\text{Cin}}$$

If $\overline{\text{Cin}}$, then $\text{Cout} = (\overline{X} Y + X \overline{Y} + X Y)$
 If Cin , then $\text{Cout} = X Y$

$$\text{Cin} (X + Y) + \overline{\text{Cin}} X Y$$

$$\text{Cin} (\overline{X} Y + X \overline{Y} + X Y) + \overline{\text{Cin}} X Y$$

$$\text{Cin} (\overline{X} Y + X \overline{Y}) + (\text{Cin} + \overline{\text{Cin}}) X Y$$

$$P \text{Cin} + G$$

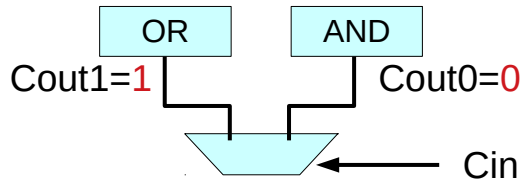
$$\text{Cin} (X + Y) + \overline{\text{Cin}} X Y$$

$$\text{Cin } P' + \overline{\text{Cin}} G$$

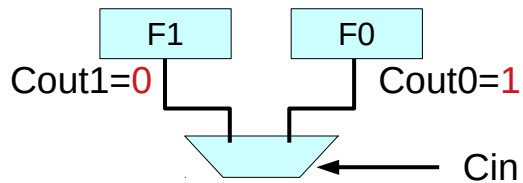
... P' : relaxed P

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FPGA Carry Chain Cell



Cout1=1 when Cin=1
 Cout0=0 when Cin=0
 Cout = Cin



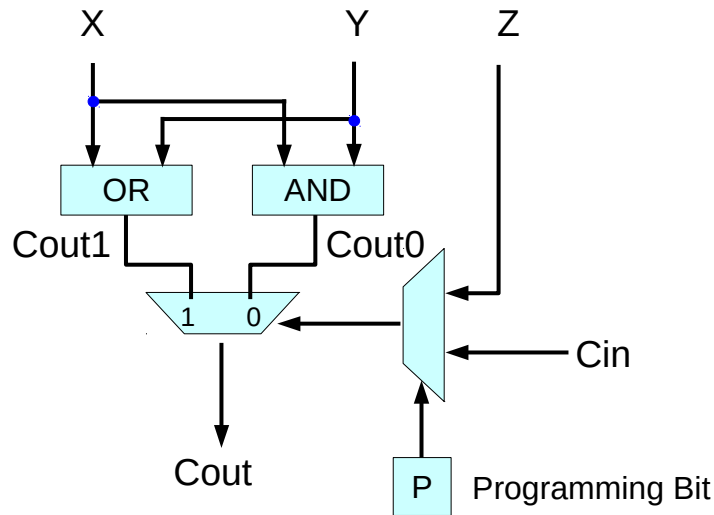
Cout1=0 when Cin=1
 Cout0=1 when Cin=0
 Cout = $\overline{\text{Cin}}$

Cout0	Cout1	Cout	Name
0	0	0	Kill
0	1	$\overline{\text{Cin}}$	Propagate
1	0	$\overline{\text{Cin}}$	Inverse Propagate
1	1	1	Generate

Cout1	Cout0	Cout	Name
0	0	0	Kill
0	1	$\overline{\text{Cin}}$	Inverse Propagate
1	0	$\overline{\text{Cin}}$	Propagate
1	1	1	Generate

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Carry Chain



X	Y	Cin	$\overline{\text{Cin}}$	
		Cout1	Cout0	
0	0	0	0	$\overline{X} \overline{Y}$
0	1	1	0	$\overline{X} Y$
1	0	1	0	$X \overline{Y}$
1	1	1	1	$X Y$

Cout1	Cout0	Cout	Name
0	0	0	Kill
0	1	$\overline{\text{Cin}}$	Inverse Propagate
1	0	Cin	Propagate
1	1	1	Generate

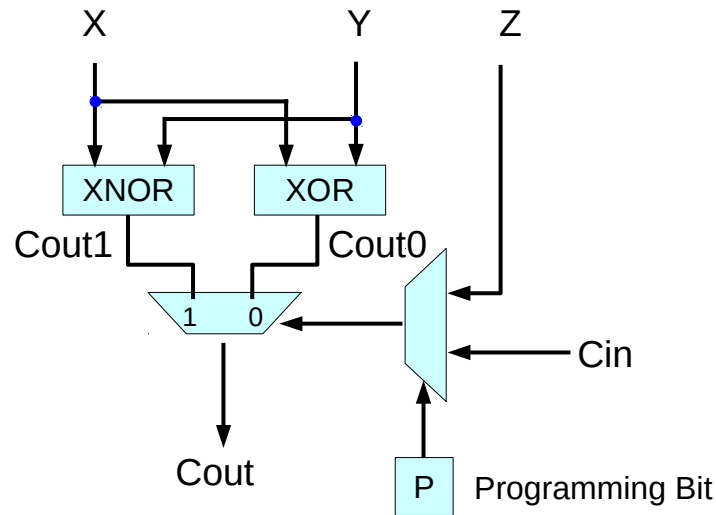
Carry Out

X	Y	Cin	Cout
0	0	Cin	$\overline{\text{Cin}}$
0	1	Cin	$\overline{\text{Cin}}$
1	0	Cin	$\overline{\text{Cin}}$
1	1	Cin	Cin

Cout1=1 when Cin=1
 Cout0=0 when Cin=0
 Cout = Cin propagate

Cout1=0 when Cin=1
 Cout0=1 when Cin=0
 Cout = $\overline{\text{Cin}}$ inverse propagate

Parity Checker



X	Y	Cin	$\overline{\text{Cin}}$	
		Cout1	Cout0	
0	0	1	0	$\overline{X} \overline{Y}$
0	1	0	1	$\overline{X} Y$
1	0	0	1	$X \overline{Y}$
1	1	1	0	$X Y$

Cout1	Cout0	Cout	Name
0	0	0	Kill
0	1	$\overline{\text{Cin}}$	Inverse Propagate
1	0	Cin	Propagate
1	1	1	Generate

Computing Parity

$X \oplus Y \oplus \text{Cin}$	
$0 \oplus 0 \oplus \text{Cin}$	$\overline{\text{Cin}}$
$0 \oplus 1 \oplus \text{Cin}$	$\overline{\text{Cin}}$
$1 \oplus 0 \oplus \text{Cin}$	$\overline{\text{Cin}}$
$1 \oplus 1 \oplus \text{Cin}$	Cin

Cout1=1 when Cin=1

Cout0=0 when Cin=0

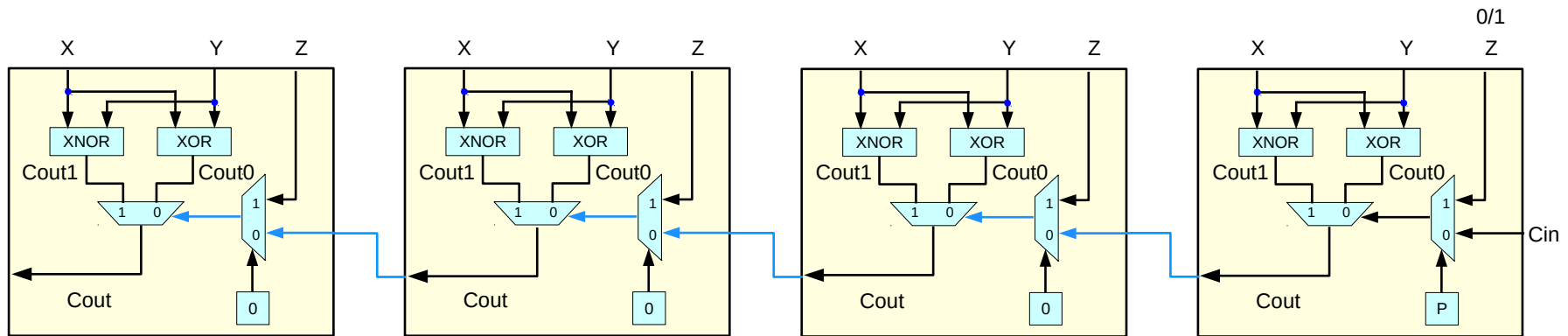
Cout = Cin propagate

Cout1=0 when Cin=1

Cout0=1 when Cin=0

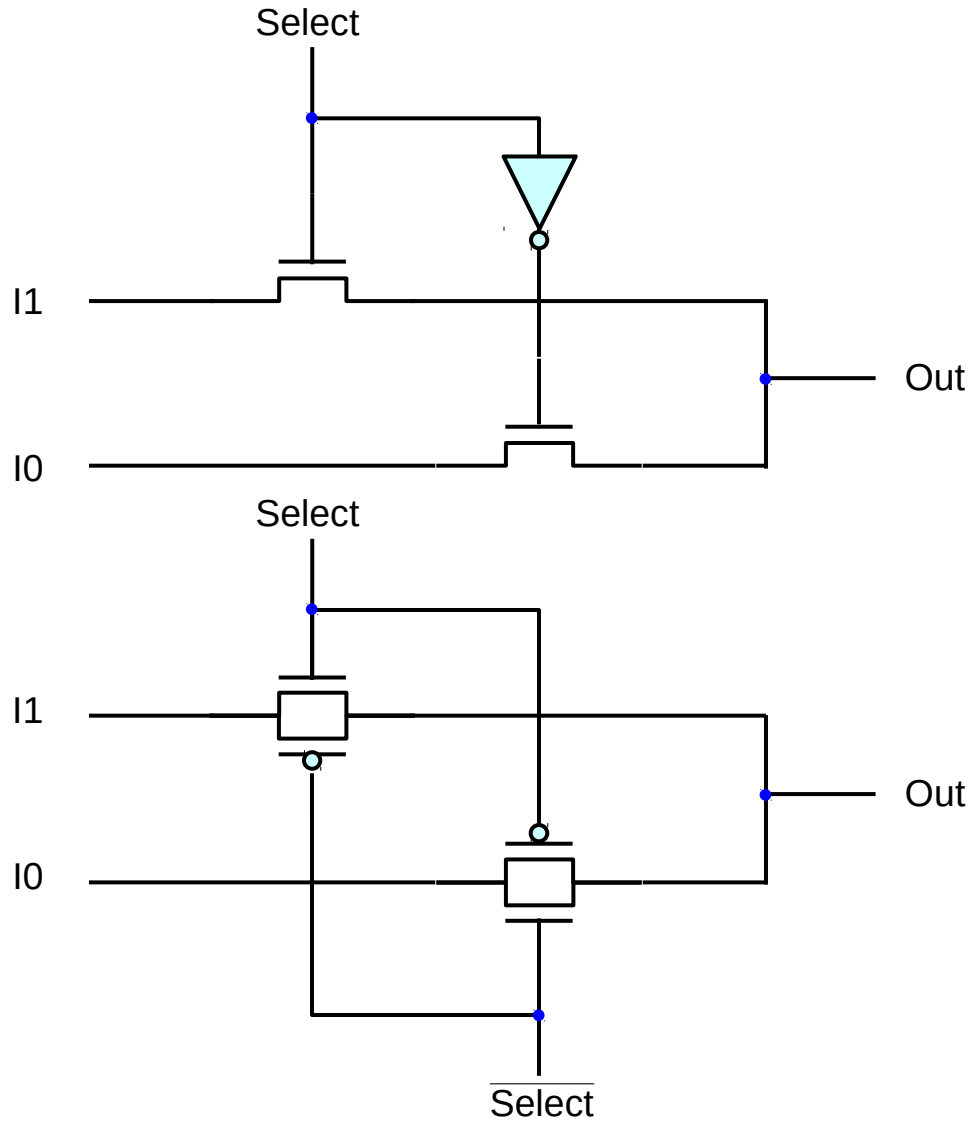
Cout = $\overline{\text{Cin}}$ inverse propagate

Ripple Carry Chain



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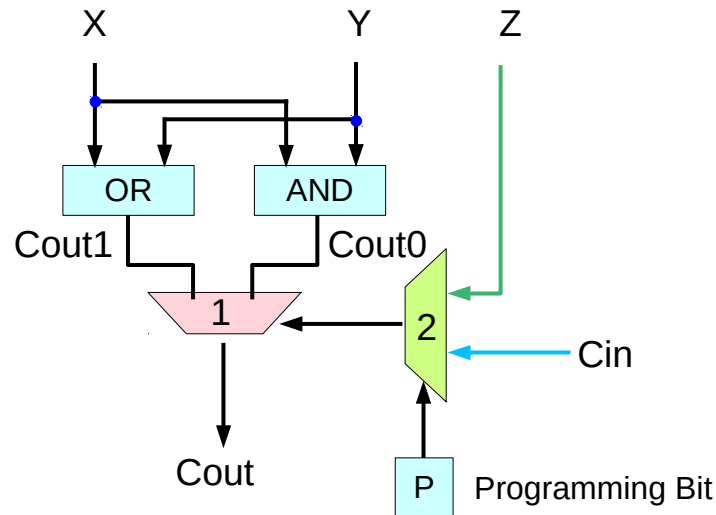
Ripple Carry Chain



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FPGA Carry Chain Cell

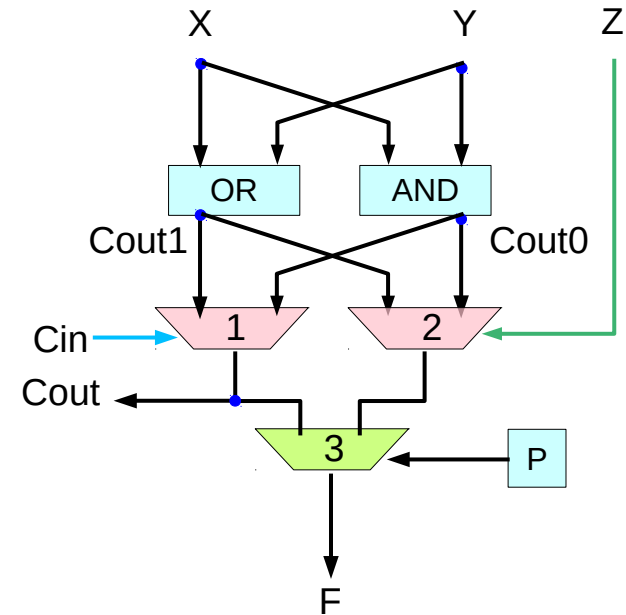
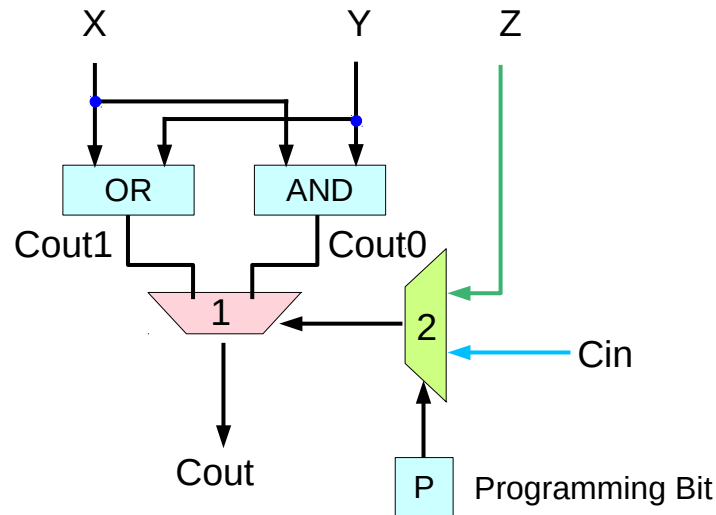


Significantly slower
Two muxes on the carry chain in each cell
Delay 1 for first cell
Delay 3 for each additional cell in the carry chain
1 delay for mux 2 and 2 delays for mux 1
Overall $2n-2$ for an n -cell carry chain

The critical path comes from the 2-LUTs and not input Z
Since the delay through the 2-LUTs will be larger than
Through mux 2 in the first cell

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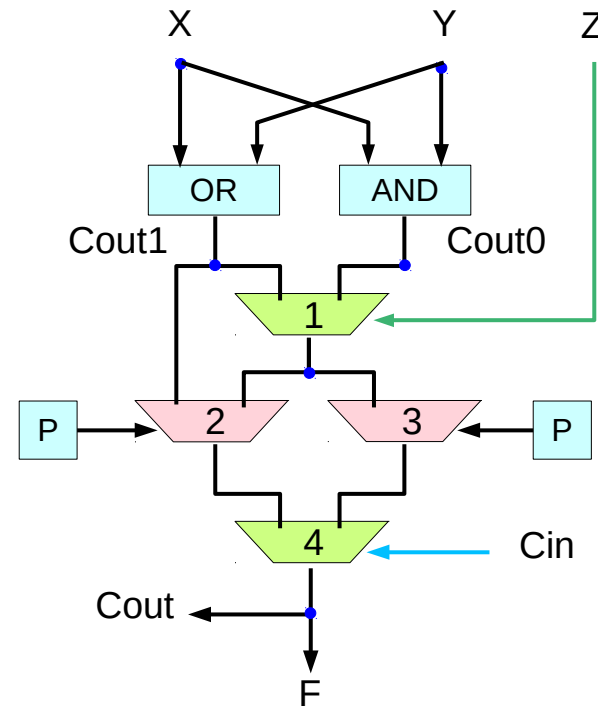
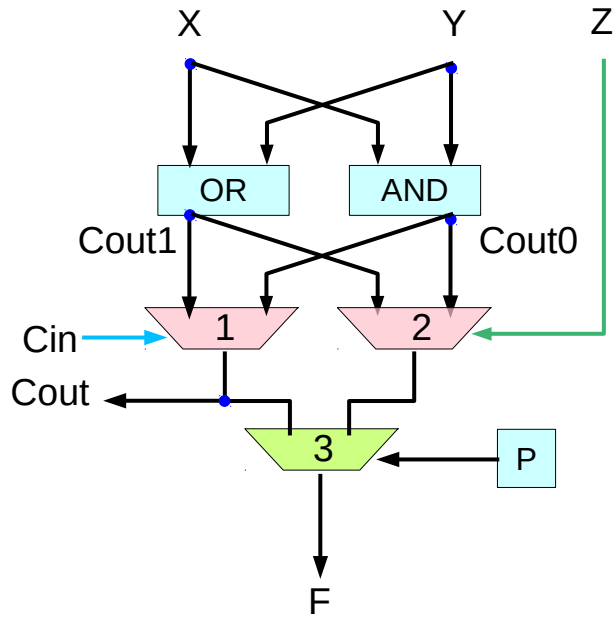
FPGA Carry Chain Cell



Not logically equivalent
No longer use the Z input in the first cell
Since Z is only attached to mux2
And mux 2 does not lead to the carry cells

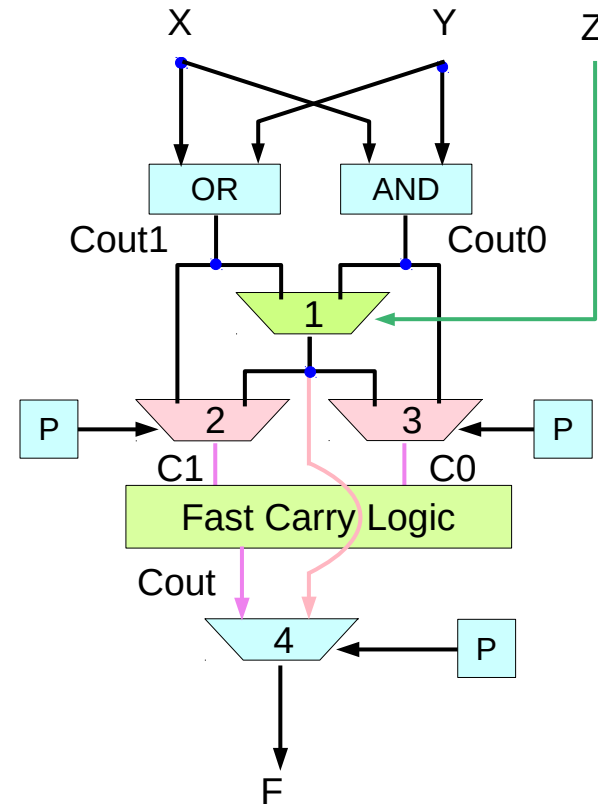
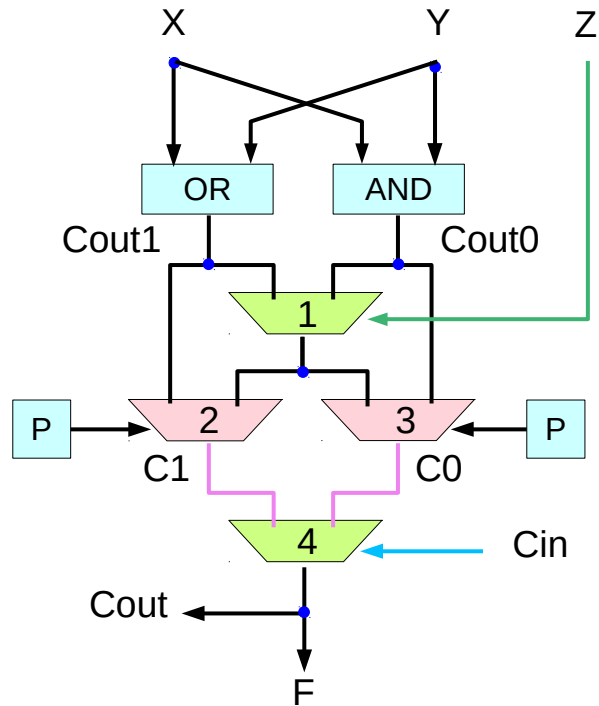
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FPGA Carry Chain Cell



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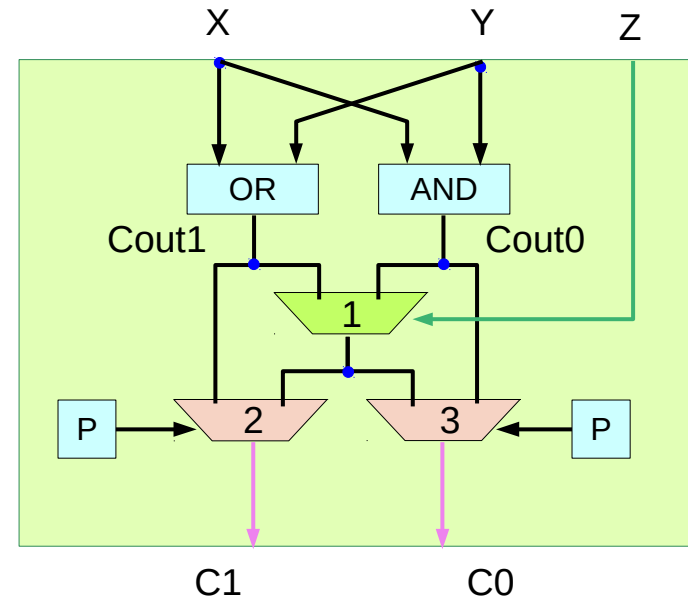
FPGA Carry Chain Cell



$$Cout_i = (Cout_{i-1} \cdot C1_i) + (\overline{Cout_{i-1}} \cdot C0_i)$$

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FPGA Carry Chain Cell



$$Cout_i = (Cout_{i-1} \cdot C1_i) + (\overline{Cout_{i-1}} \cdot C0_i)$$

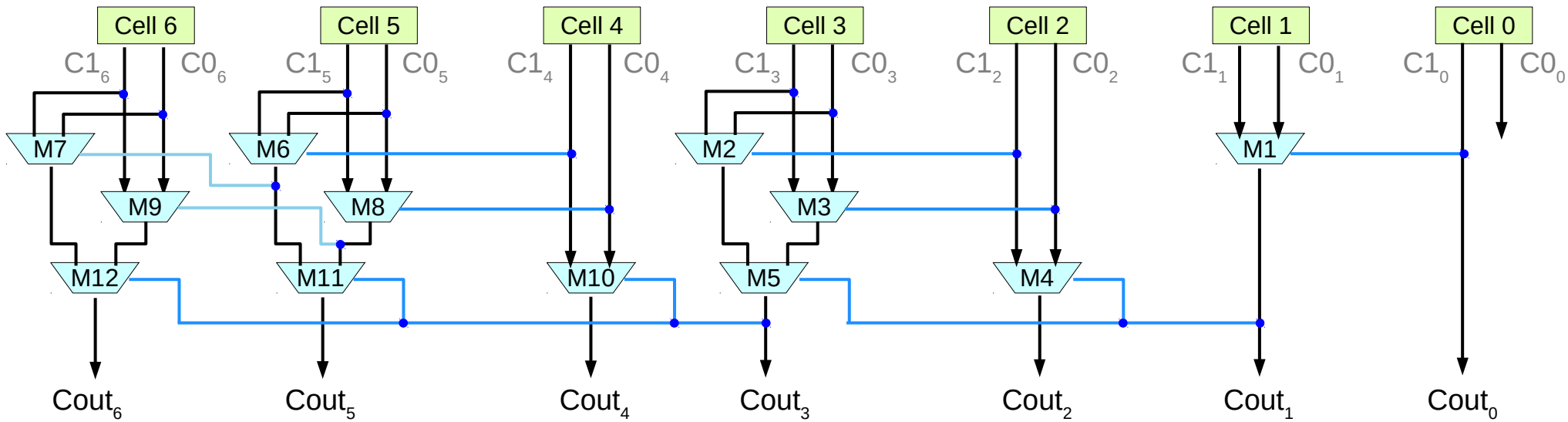
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Fast Carry Logc

Carry Select Adder
Carry Lookahead Adder
 Brent-Kung
Variable Block
Ripple Carry Adder

https://en.wikipedia.org/wiki/Carry-lookahead_adder

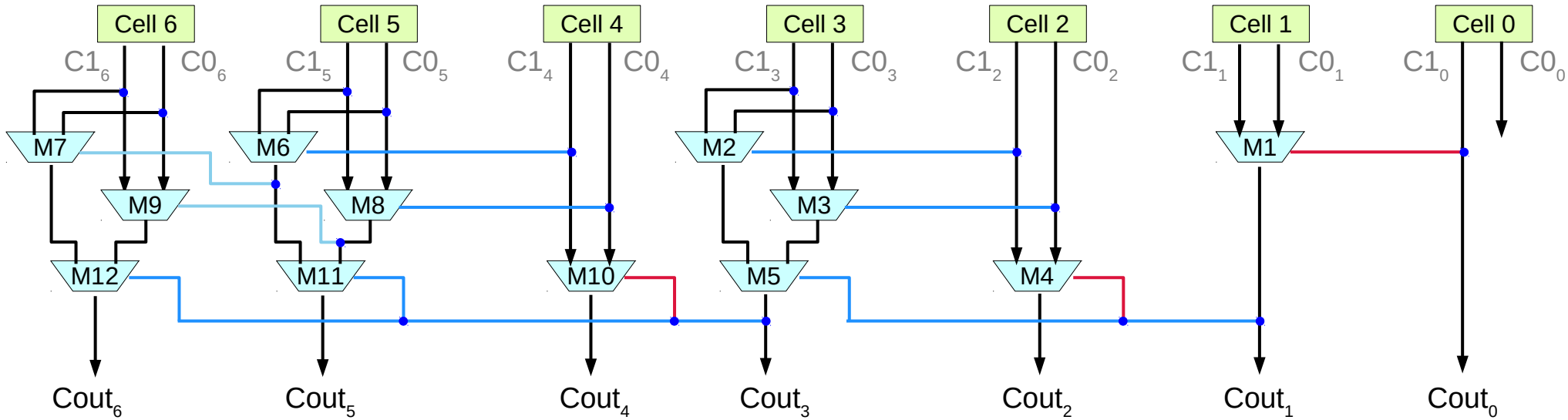
FPGA Carry Chain Cell



$$Cout_i = (Cout_{i-1} \cdot C1_i) + (\overline{Cout_{i-1}} \cdot C0_i)$$

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FPGA Carry Chain Cell



$$Cout_i = (Cout_{i-1} \cdot C1_i) + (\overline{Cout_{i-1}} \cdot C0_i)$$

$$Cout_1 = (Cout_0 \cdot C1_1) + (\overline{Cout_0} \cdot C0_1)$$

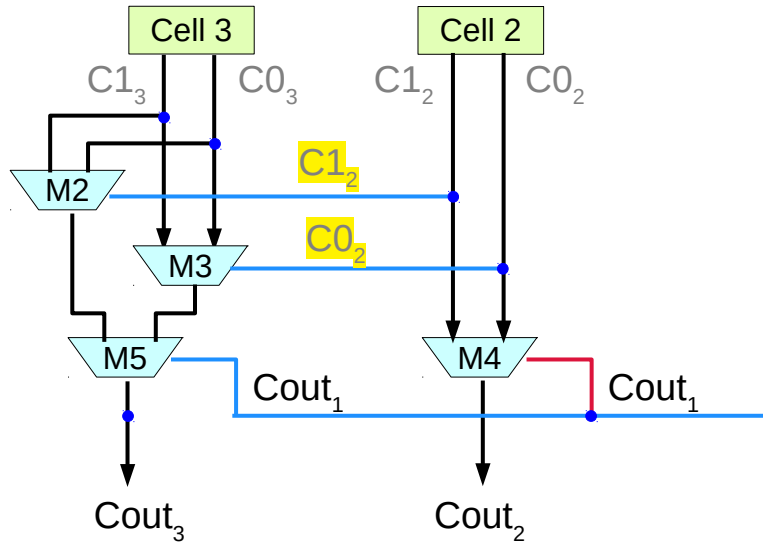
$$Cout_{i+1} = (Cout_i \cdot C1_{i+1}) + (\overline{Cout_i} \cdot C0_{i+1})$$

$$Cout_1 = (C1_0 \cdot C1_1) + (\overline{C1_0} \cdot C0_1)$$

$$Cout_{i+1} = (((Cout_{i-1} \cdot C1_i) + (\overline{Cout_{i-1}} \cdot C0_i)) \cdot C1_{i+1}) + (\overline{((Cout_{i-1} \cdot C1_i) + (\overline{Cout_{i-1}} \cdot C0_i))} \cdot C0_{i+1})$$

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FPGA Carry Chain Cell



$$Cout_i = (Cout_{i-1} \cdot C1_i) + (\overline{Cout_{i-1}} \cdot C0_i)$$

$$Cout_{i+1} = (Cout_i \cdot C1_{i+1}) + (\overline{Cout_i} \cdot C0_{i+1})$$

$$Cout_2 = (Cout_1 \cdot C1_2) + (\overline{Cout_1} \cdot C0_2)$$

$$Cout_3 = (Cout_2 \cdot C1_3) + (\overline{Cout_2} \cdot C0_3)$$

$$= (((Cout_1 \cdot C1_2) + (\overline{Cout_1} \cdot C0_2)) \cdot C1_3)$$

$$+ (((Cout_1 \cdot C1_2) + (\overline{Cout_1} \cdot C0_2)) \cdot C0_3)$$

$$(((Cout_1 \cdot C1_2) + (\overline{Cout_1} \cdot C0_2)) \cdot C1_3)$$

$$= (C1_3 C1_2 Cout_1 + C1_3 C0_2 \overline{Cout_1})$$

$$(((\overline{Cout_1} \cdot C1_2) + (\overline{Cout_1} \cdot C0_2)) \cdot C0_3)$$

$$= (((\overline{Cout_1} + \overline{C1_2}) \cdot (\overline{Cout_1} + \overline{C0_2})) \cdot C0_3)$$

$$= (\overline{Cout_1} \overline{Cout_1} + \overline{C1_2} \overline{Cout_1} + \overline{Cout_1} \overline{C0_2} + \overline{C1_2} \overline{C0_2}) \cdot C0_3$$

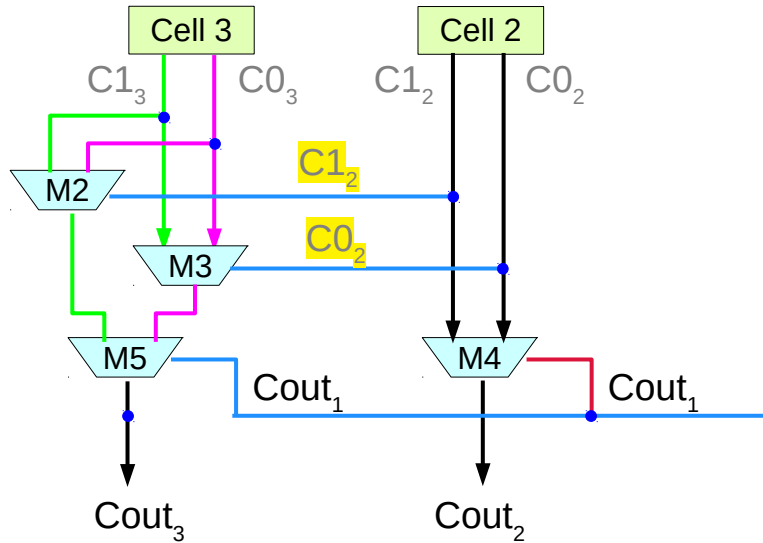
$$= (\overline{C1_2} \overline{Cout_1} + \overline{C0_2} \overline{Cout_1}) \cdot C0_3$$

$$= (C0_3 \overline{C1_2} \overline{Cout_1} + C0_3 \overline{C0_2} \overline{Cout_1})$$

$$(C1_3 C1_2 + C0_3 \overline{C1_2}) Cout_1 + (C1_3 C0_2 + C0_3 \overline{C0_2}) \overline{Cout_1}$$

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FPGA Carry Chain Cell

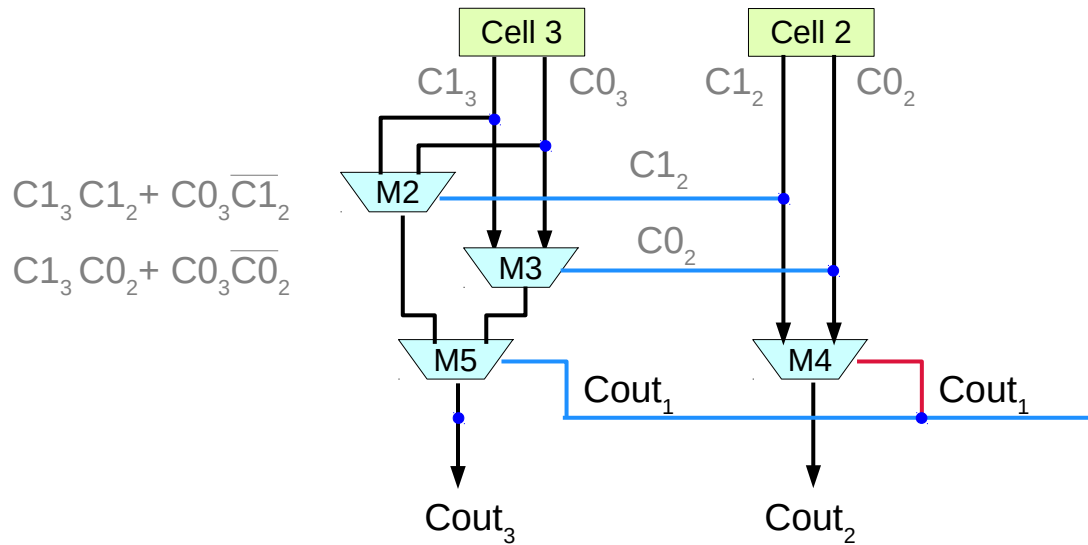


$$\begin{aligned}
 &= (\overline{Cout_1}Cout_1 + \overline{C1_2}Cout_1 + \overline{Cout_1}C0_2 + \overline{C1_2}C0_2) \cdot C0_3 \\
 &= (\overline{C1_2}Cout_1 + \overline{C0_2}Cout_1) \cdot C0_3 \\
 &= (C0_3\overline{C1_2}Cout_1 + C0_3\overline{C0_2}Cout_1)
 \end{aligned}$$

$$(C1_3 C1_2 + C0_3 \overline{C1_2})Cout_1 + (C1_3 C0_2 + C0_3 \overline{C0_2})\overline{Cout_1}$$

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FPGA Carry Chain Cell



$$C1_3 C1_2 + C0_3 \overline{C1_2}$$

$$C1_3 C0_2 + C0_3 \overline{C0_2}$$

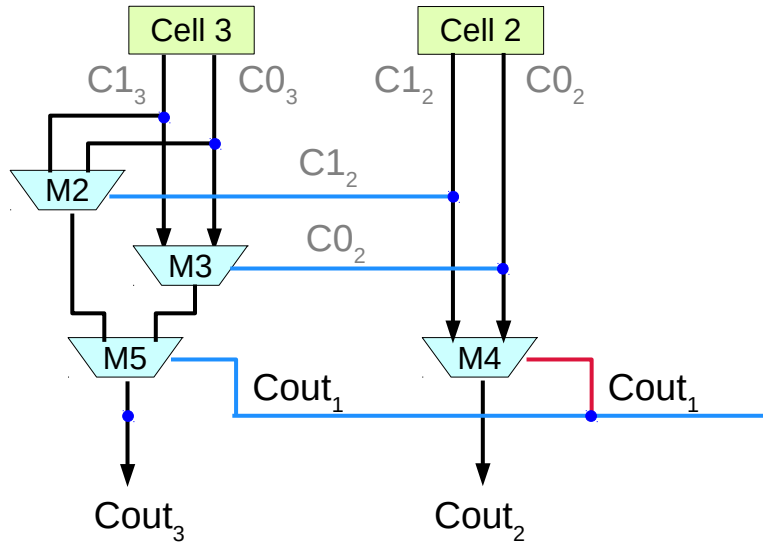
$$(C1_3 C1_2 + C0_3 \overline{C1_2}) Cout_1 + (C1_3 C0_2 + C0_3 \overline{C0_2}) \overline{Cout_1}$$

$$= C1_3 \cdot (C1_2 Cout_1 + C0_2 \overline{Cout_1})$$

$$+ C0_3 \cdot (\overline{C1_2} Cout_1 + \overline{C0_2} \overline{Cout_1})$$

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FPGA Carry Chain Cell



$$Cout_i = (Cout_{i-1} \cdot C1_i) + (\overline{Cout_{i-1}} \cdot C0_i)$$

$$Cout_{i+1} = (Cout_i \cdot C1_{i+1}) + (\overline{Cout_i} \cdot C0_{i+1})$$

$$Cout_{i+1} = \left(\left[(Cout_{i-1} \cdot C1_i) + (\overline{Cout_{i-1}} \cdot C0_i) \right] \cdot C1_{i+1} \right) + \left(\left[(Cout_{i-1} \cdot C1_i) + (\overline{Cout_{i-1}} \cdot C0_i) \right] \cdot C0_{i+1} \right)$$

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References

[1] <http://en.wikipedia.org/>

[2] J-P Deschamps, et. al., “Sunthesis of Arithmetic Circuits”, 2006