

Applications of Array Pointers (1A)

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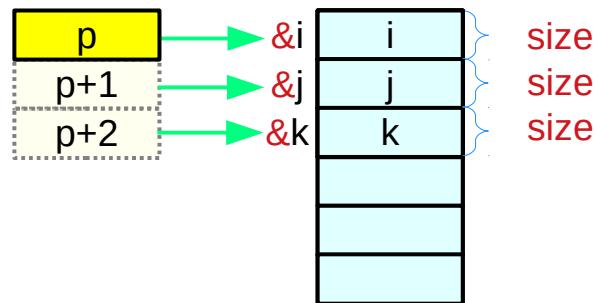
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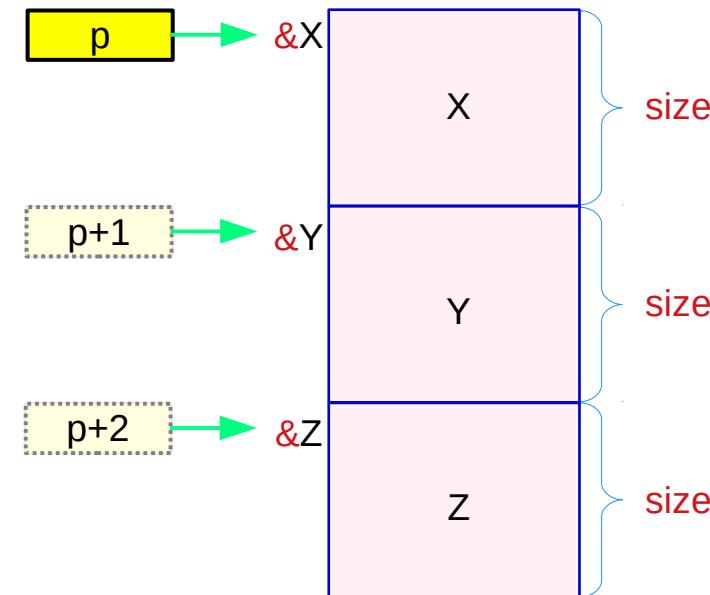
Virtual Array Pointers in Multi-dimensional Arrays

Pointers to various data types

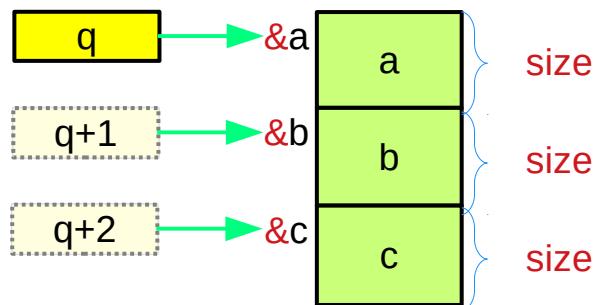
int *p; **int i, j, k;**



T *p; **T X, Y, Z;**



double *q; **double a, b, c;**



pointer

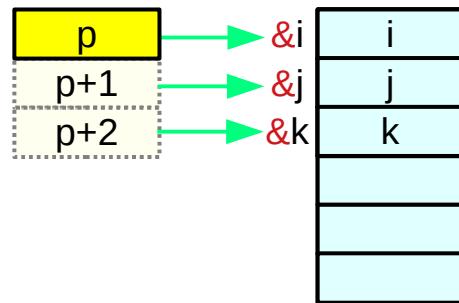
abstract data

Pointers to primitive data

int *p;

int i, j, k;

sizeof(int) = 4 bytes



size
size
size

= sizeof(i)
= sizeof(j)
= sizeof(k)

= sizeof(*p)
= sizeof(*p+1)
= sizeof(*p+2)

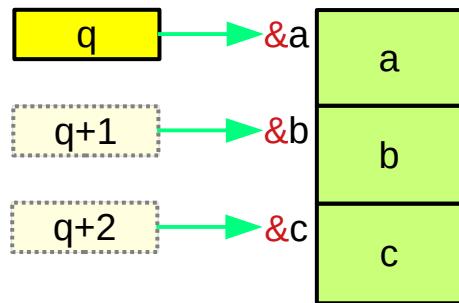
≠ sizeof(p)
≠ sizeof(p+1)
≠ sizeof(p+2)

pointer size
4 or 8 bytes

double *q;

double a, b, c;

sizeof(double) = 8 bytes



size
size
size

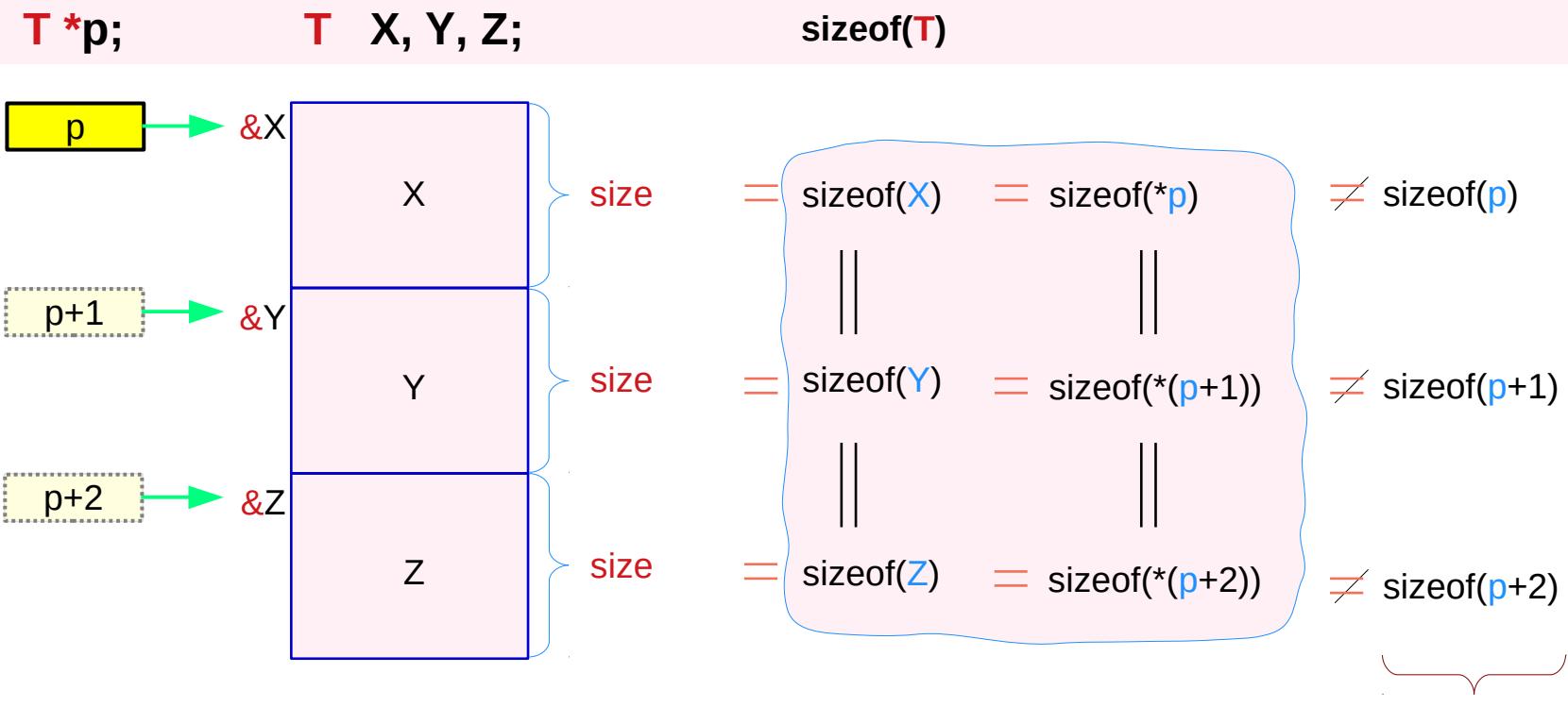
= sizeof(a)
= sizeof(b)
= sizeof(c)

= sizeof(*q)
= sizeof(*q+1)
= sizeof(*q+2)

≠ sizeof(q)
≠ sizeof(q+1)
≠ sizeof(q+2)

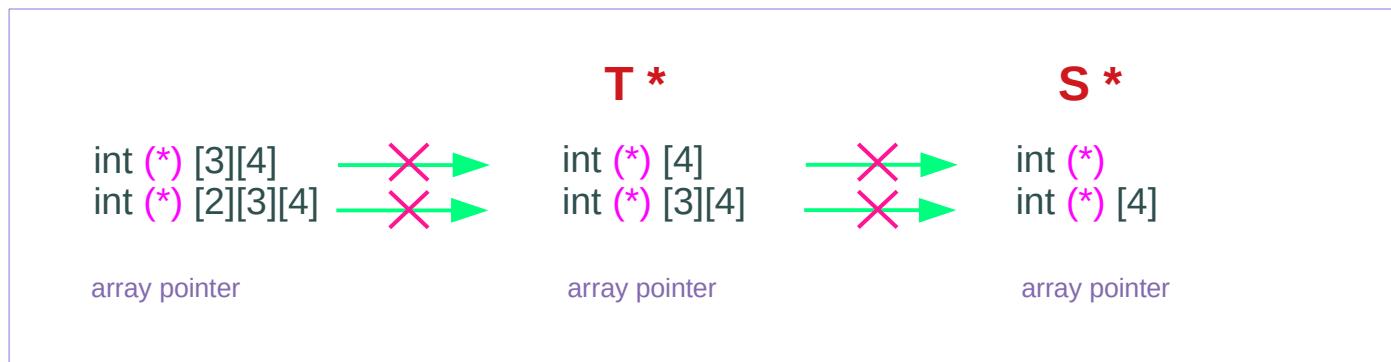
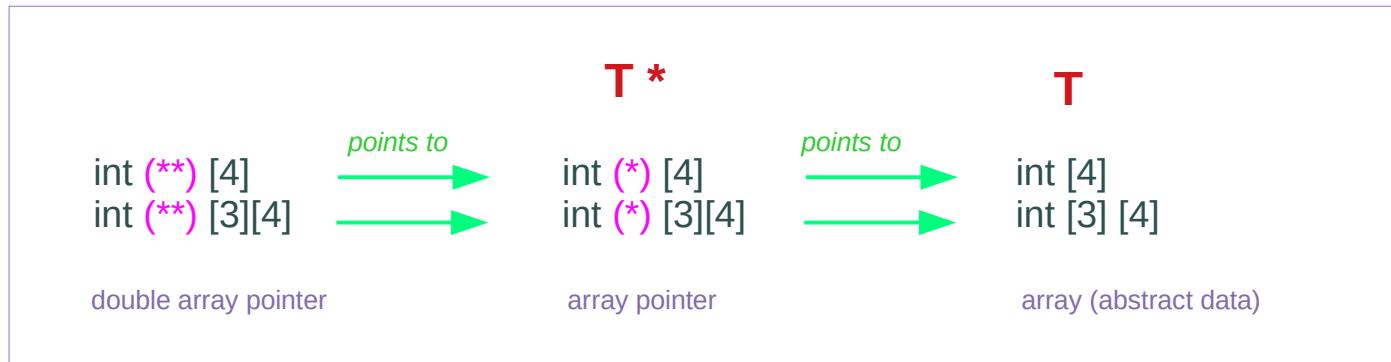
pointer size
4 or 8 bytes

Pointers to abstract data

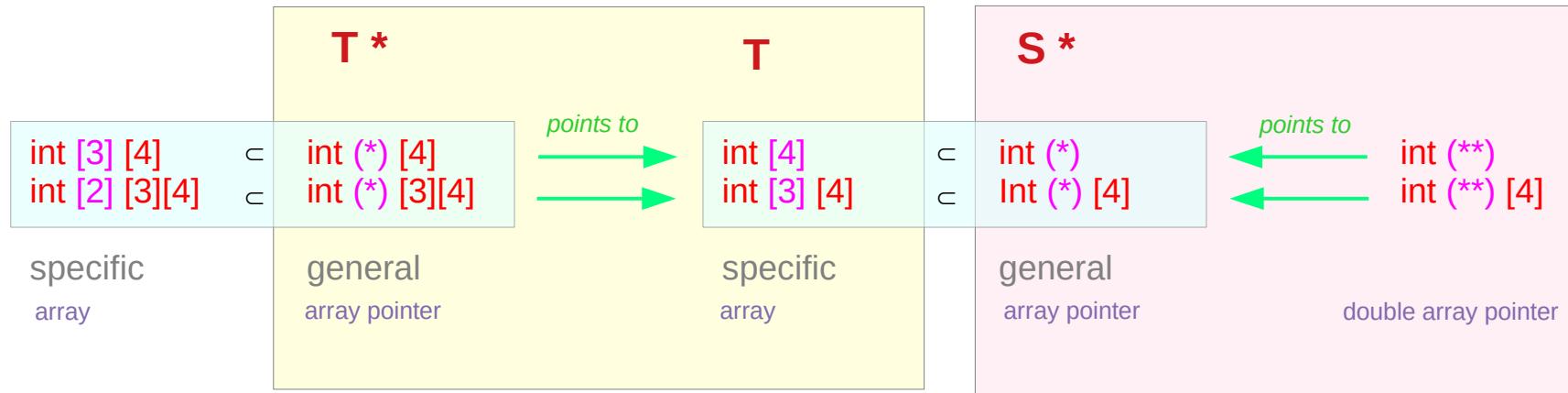


type ----- array
value ----- start address
increment size ----- size

Array pointer and array types



Specific array types v.s. general array pointer types



Array pointers have augmented dimensions

typedef int (*T1) [4];	typedef int T2[4];
typedef int (*T1) [3][4];	typedef int T2[3][4];
int (*) [4]	int [4]
int (*) [3][4]	int [3] [4]
general	specific

T1 a;
T2 b;

T1 is a pointer type
T2 is an array type
T1 has one more dimension than T2

a = &b;
***a = b;**

a references **b**
b is the dereference of **a**

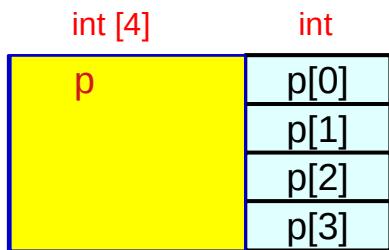
(a+1) = ?
***(a+1) = ?**

(a+2) = ?
***(a+2) = ?**

Virtual pointers in an array of integers

`int p[3];`

p is an abstract data (array)

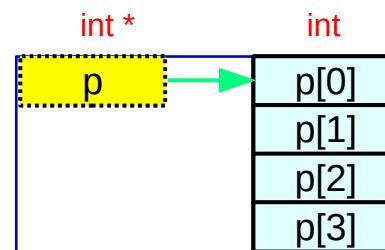


p is the name of an array

p has the size of the whole array

p has an array type (abstract data)

p can also be viewed as a pointer



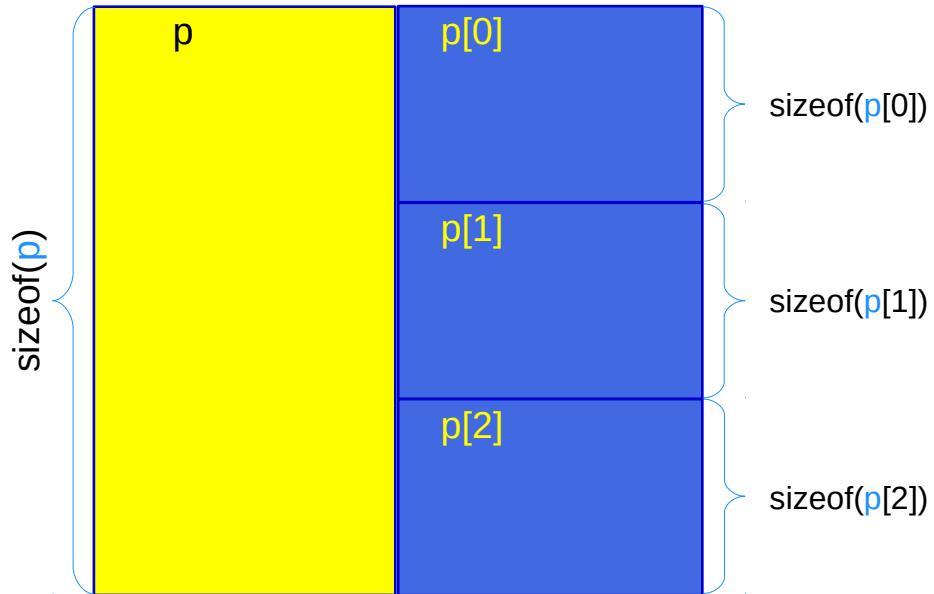
p also has pointer characteristics

p has the value of the starting address

p is a virtual pointer

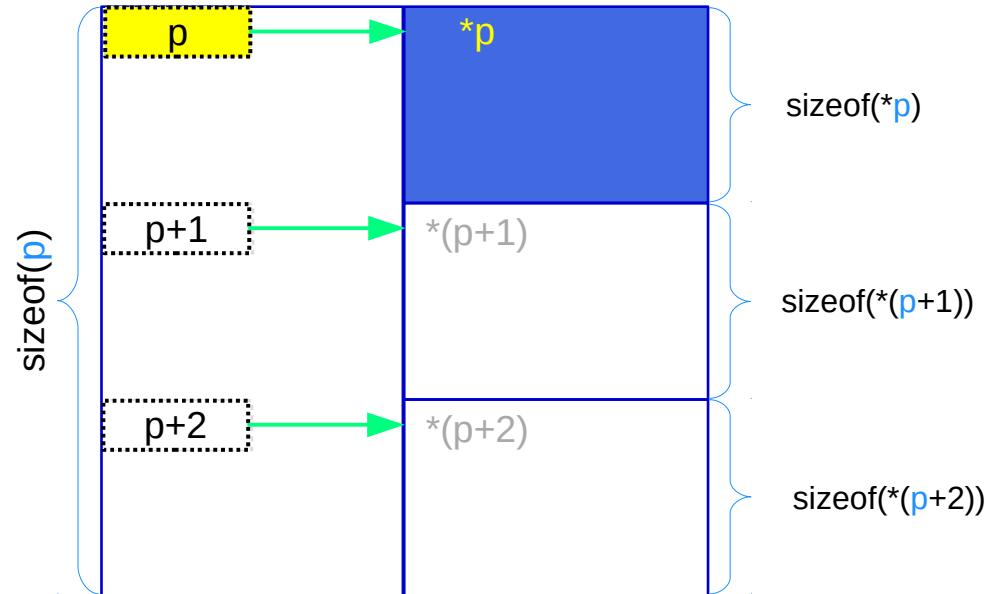
Virtual pointers in an array of abstract data

Abstract data array p



p has an array type (abstract data element)
p is the name of an array
p has the size of the whole array

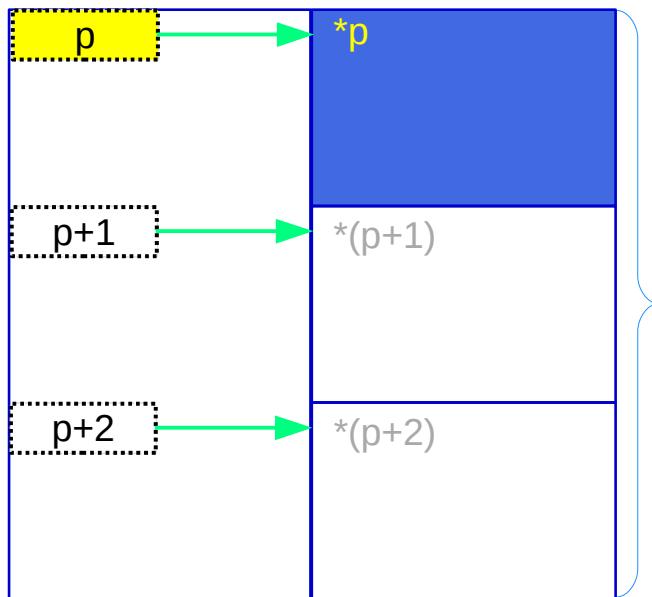
Virtual pointer p



p also has a pointer type
p has the value of the starting address
p is a virtual array pointer

Virtual pointer to abstract data

virtual pointer p abstract data *p



whole array size

$\text{sizeof}(p)$

\equiv

$\text{sizeof}(*p) * 3$

~~$\text{sizeof}(p+1)$~~

~~\equiv~~

$\text{sizeof}(*p+1) * 3$

~~$\text{sizeof}(p+2)$~~

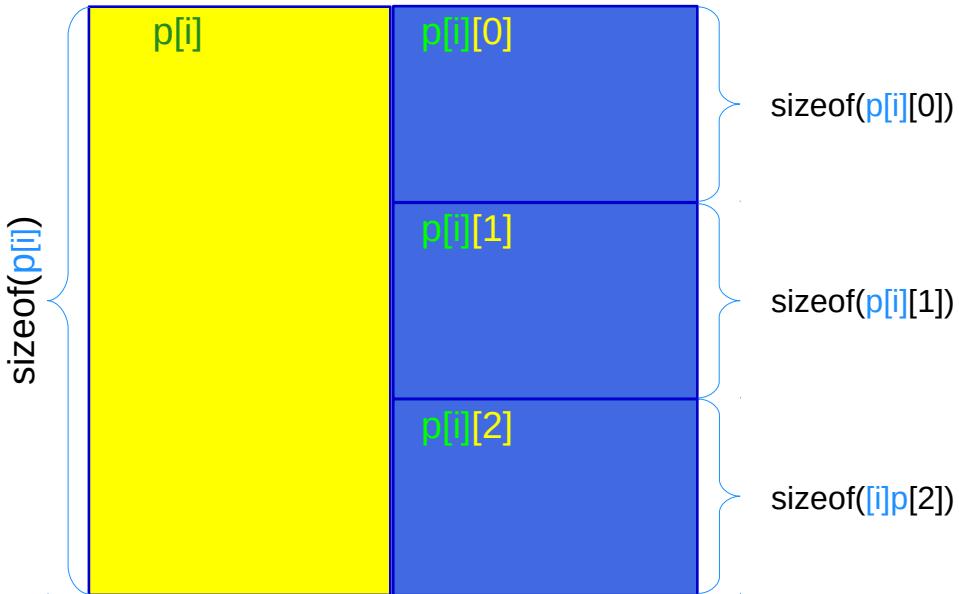
~~\equiv~~

$\text{sizeof}(*p+2) * 3$

pointer size
4 / 8 bytes

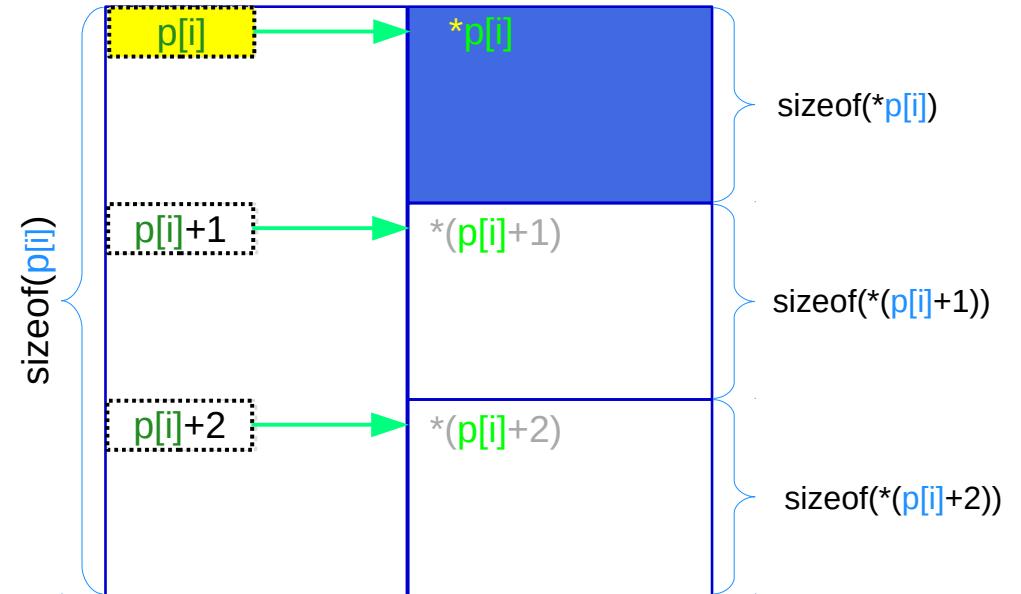
Virtual array pointers in a multi-dimensional array

Abstract data (array) $p[i]$



$p[i]$ has an array type (abstract data)
 $p[i]$ is the name of an array
 $p[i]$ has the size of the whole array

Virtual array pointer $p[i]$



$p[i]$ also has an array pointer type
 $p[i]$ has the value of the starting address
 $p[i]$ is a virtual array pointer

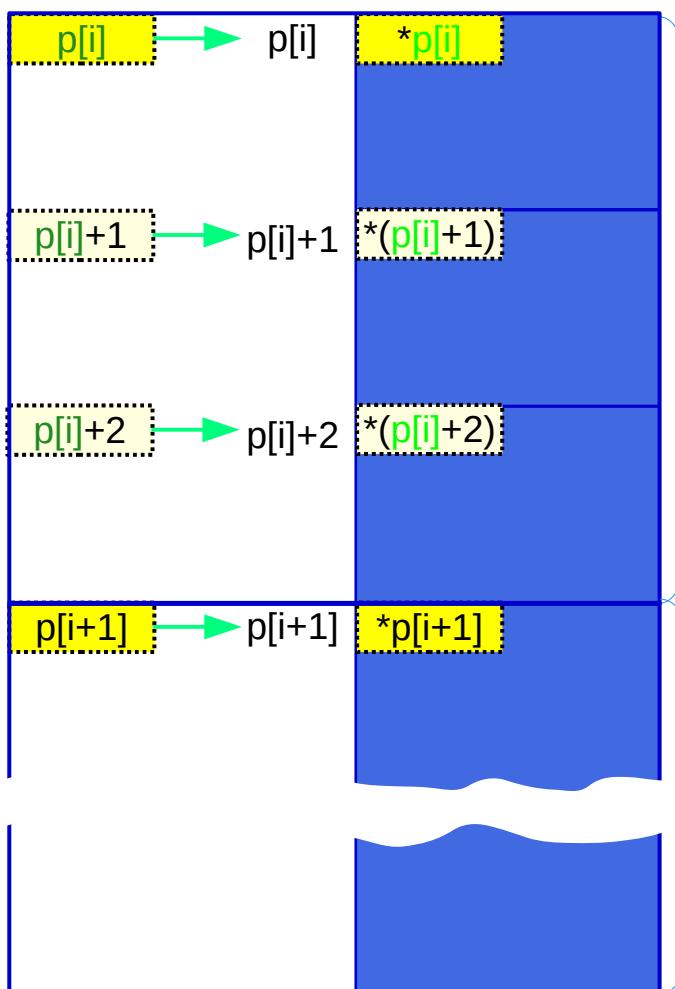
Virtual pointers to a sub array

$p[i] :: T *$

array pointer type

$*p[i], *p[i+1] :: T$

array type



$\text{sizeof}(p[i])$

$$= \text{sizeof}(*p[i]) * N$$
$$= \text{sizeof}(p[i][0]) * N$$

$\text{size} = \text{sizeof}(*p[i]) =$
 $= \text{sizeof}(p[i][0])$

$\text{size} = \text{sizeof}(*p[i]+1) =$
 $= \text{sizeof}(p[i][1])$

$\text{size} = \text{sizeof}(*p[i]+2) =$
 $= \text{sizeof}(p[i][2])$

$\text{sizeof}(p[i+1])$

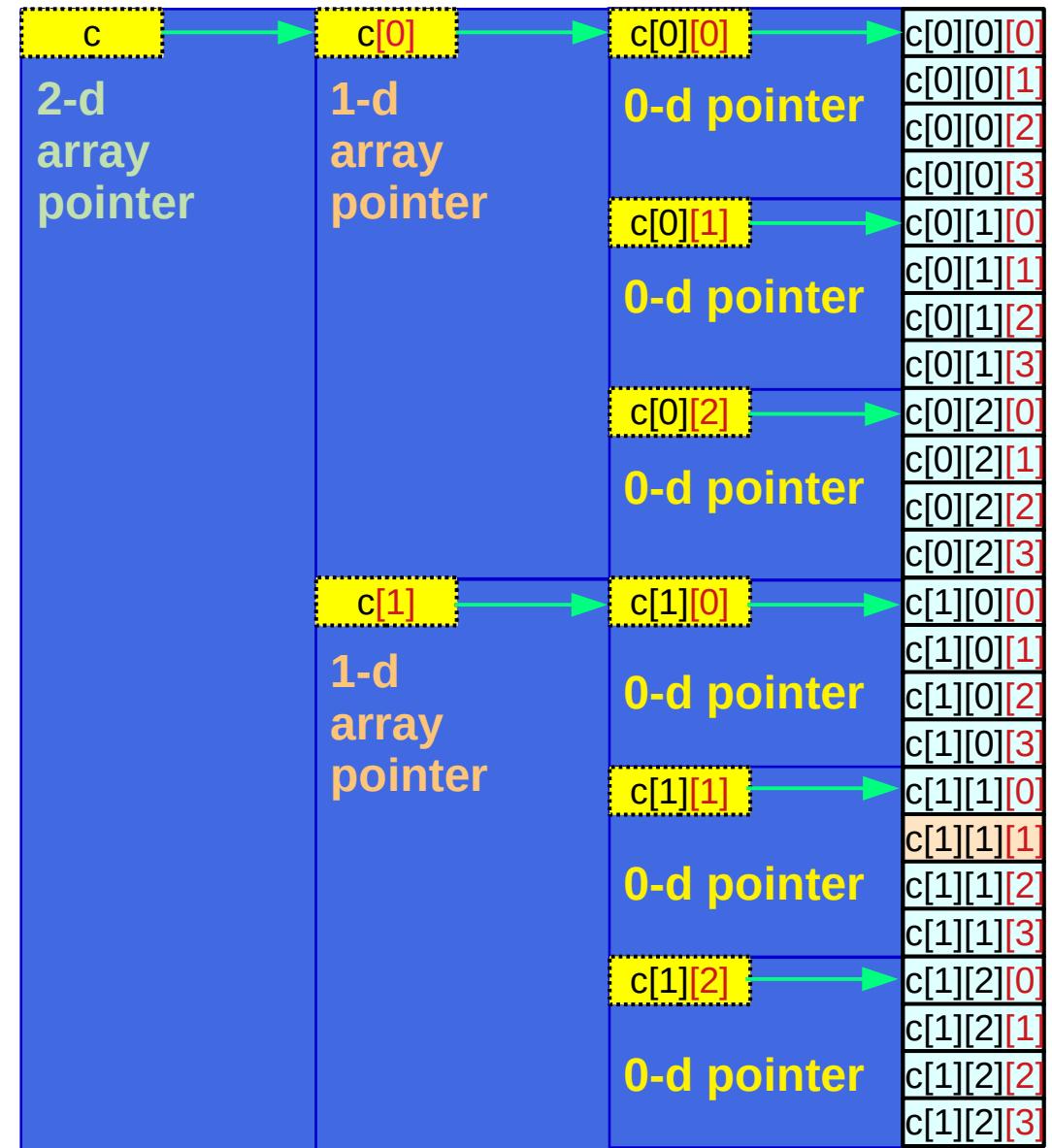
$$= \text{sizeof}(*p[i+1]) * N$$
$$= \text{sizeof}(p[i+1][0]) * N$$

3-d array structure – pointer representation

```
int c[2][3][4];
```

```
*(*(*c + i) + j) + k)
```

- Hierarchical
- Nested Structure
- Virtual Array Pointers to abstract data (subarrays)
- Contiguous and Linear Data Layout
- Row Major Order



3-d array structure – abstract data representation

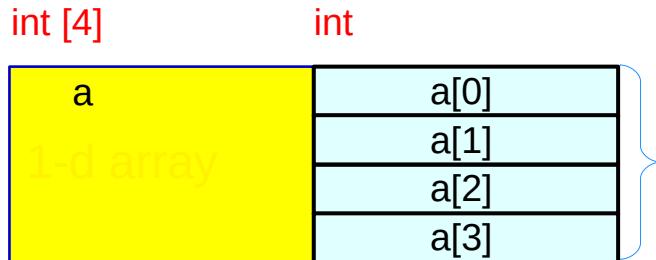
```
int c[2][3][4];
```

```
((c [i])[j])[k]
```

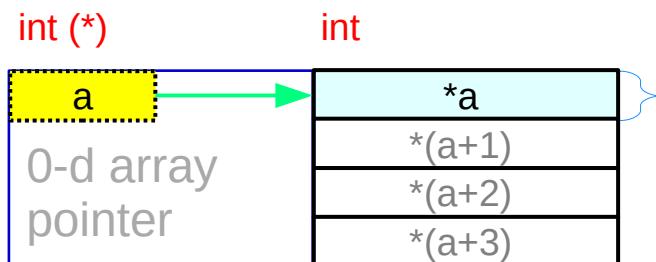
- Hierarchical
- Nested Structure
- Virtual Array Pointers to abstract data (subarrays)
- Contiguous and Linear Data Layout
- Row Major Order

c 3-d array name	c[0] 2-d array name	c[0][0] 1-d array name	c[0][0][0] c[0][0][1] c[0][0][2] c[0][0][3]
		c[0][1] 1-d array name	c[0][1][0] c[0][1][1] c[0][1][2] c[0][1][3]
		c[0][2] 1-d array name	c[0][2][0] c[0][2][1] c[0][2][2] c[0][2][3]
	c[1] 2-d array name	c[1][0] 1-d array name	c[1][0][0] c[1][0][1] c[1][0][2] c[1][0][3]
		c[1][1] 1-d array name	c[1][1][0] c[1][1][1] c[1][1][2] c[1][1][3]
		c[1][2] 1-d array name	c[1][2][0] c[1][2][1] c[1][2][2] c[1][2][3]

Array a and pointer a



1-d array a specific array type
sizeof(a)



pointer a general pointer type
sizeof(a) = sizeof(*a) * 4

- a is the name of a 1-d array
 - a also has a pointer type
 - a has the size of the array
 - a has the value of the starting address
- a is a virtual array pointer

Array **b** and pointer **b**

2-d array b specific array type

`sizeof(b)`

`int [3] [4]` `int [4]`

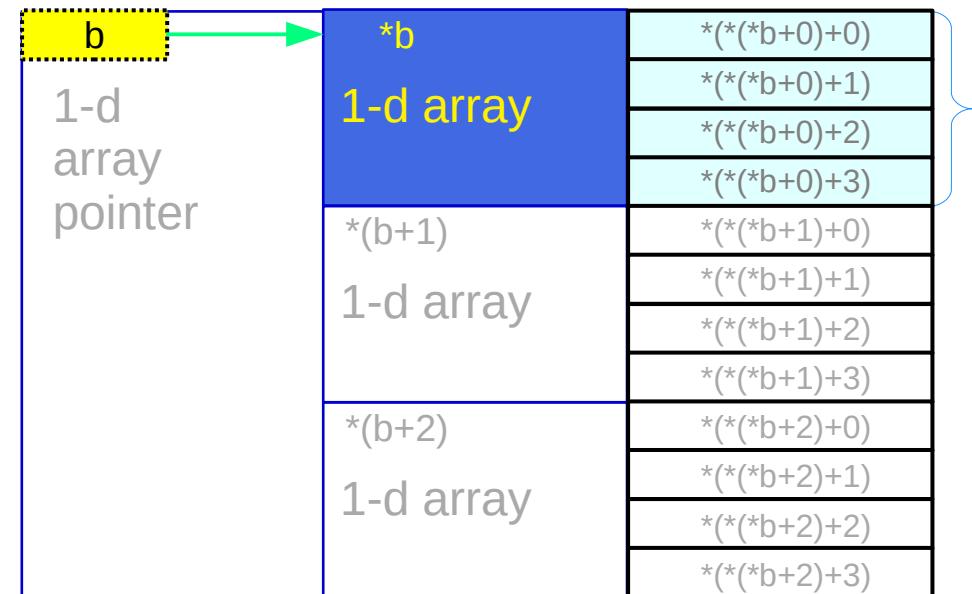


b is the name of a 2-d array
b has the size of the array

1-d array pointer b general pointer type

`sizeof(b) = sizeof(*b) * 3`

`int (*) [4]` `int [4]`



b also has a 1-d array pointer type
b has the value of the starting address

b is a virtual array pointer

Array c

3-d array c

specific array type

`sizeof(c)`

int [2][3][4]	int [3][4]	int [4]	
c 3-d array	c[0] 2-d array	c[0][0] 1-d array	c[0][0][0] c[0][0][1] c[0][0][2] c[0][0][3]
		c[0][1] 1-d array	c[0][1][0] c[0][1][1] c[0][1][2] c[0][1][3]
		c[0][2] 1-d array	c[0][2][0] c[0][2][1] c[0][2][2] c[0][2][3]
	c[1] 2-d array	c[1][0] 1-d array	c[1][0][0] c[1][0][1] c[1][0][2] c[1][0][3]
		c[1][1] 1-d array	c[1][1][0] c[1][1][1] c[1][1][2] c[1][1][3]
		c[1][2] 1-d array	c[1][2][0] c[1][2][1] c[1][2][2] c[1][2][3]

c is the name of a 3-d array

c has the size of the array

Pointer c

2-d array pointer c

general pointer type

`sizeof(c) = sizeof(*c) * 2`

c also has a 2-d array pointer type
c has the value of the starting address

c is a virtual array pointer

<code>int (*) [3][4]</code>	<code>int [3][4]</code>	<code>int [4]</code>	
<code>c</code>	<code>*c</code>	<code>*(c+0)</code>	<code>*(*(c+0)+0)</code>
2-d array pointer	2-d array	1-d array	<code>*(*(c+0)+1)</code>
			<code>*(*(c+0)+2)</code>
			<code>*(*(c+0)+3)</code>
	<code>*(c+1)</code>	<code>*(c+1)</code>	<code>*(*(c+1)+0)</code>
	2-d array	1-d array	<code>*(*(c+1)+1)</code>
			<code>*(*(c+1)+2)</code>
			<code>*(*(c+1)+3)</code>
	<code>*(c+1)</code>	<code>*(c+1)</code>	<code>*(*(c+1)+0)</code>
	2-d array	1-d array	<code>*(*(c+1)+1)</code>
			<code>*(*(c+1)+2)</code>
			<code>*(*(c+1)+3)</code>
	<code>*(c+1)</code>	<code>*(c+1)</code>	<code>*(*(c+1)+0)</code>
	2-d array	1-d array	<code>*(*(c+1)+1)</code>
			<code>*(*(c+1)+2)</code>
			<code>*(*(c+1)+3)</code>
	<code>*(c+1)</code>	<code>*(c+1)</code>	<code>*(*(c+1)+0)</code>
	2-d array	1-d array	<code>*(*(c+1)+1)</code>
			<code>*(*(c+1)+2)</code>
			<code>*(*(c+1)+3)</code>

Types of virtual array pointers in a 3-d array

```
int c[2][3][4];
```

c[i][j][k]

c[i][j]
[k]

c[i]
[j] [k]

c
[i] [j] [k]

int

int [4]
[k]

int [3][4]
[j] [k]

int [2][3][4]
[i] [j] [k]

int

int (*)
[k]

int (*)[4]
[j] [k]

int (*)[3][4]
[i] [j] [k]

array type (name)

array pointer type

Sizes of virtual array pointers in a 3-d array

```
int c[2][3][4];
```

$$\text{sizeof}(c[i][j][k]) = \text{sizeof(int)}$$

$$\text{sizeof}(c[i][j])_{[k]} = \text{sizeof(int)} * 4_{[k]}$$

$$\text{sizeof}(c[i])_{[j] [k]} = \text{sizeof(int)} * 3 * 4_{[j] [k]}$$

$$\text{sizeof}(c)_{[i] [j] [k]} = \text{sizeof(int)} * 2 * 3 * 4_{[i] [j] [k]}$$

Address values of virtual array pointers in a 3-d array

```
int c[2][3][4];
```

$$c[i][j][k] = \&c[i][j][k]$$

$$c[i][j]+k = \&c[i][j][0] + k * \text{sizeof}(c[i][j][k])$$

$\text{sizeof}(*c[i][j])$
 $= \text{sizeof}(c[i][j][0]) = \text{sizeof(int)}$

$$c[i]+j = \&c[i][0][0] + j * \text{sizeof}(c[i][j])$$

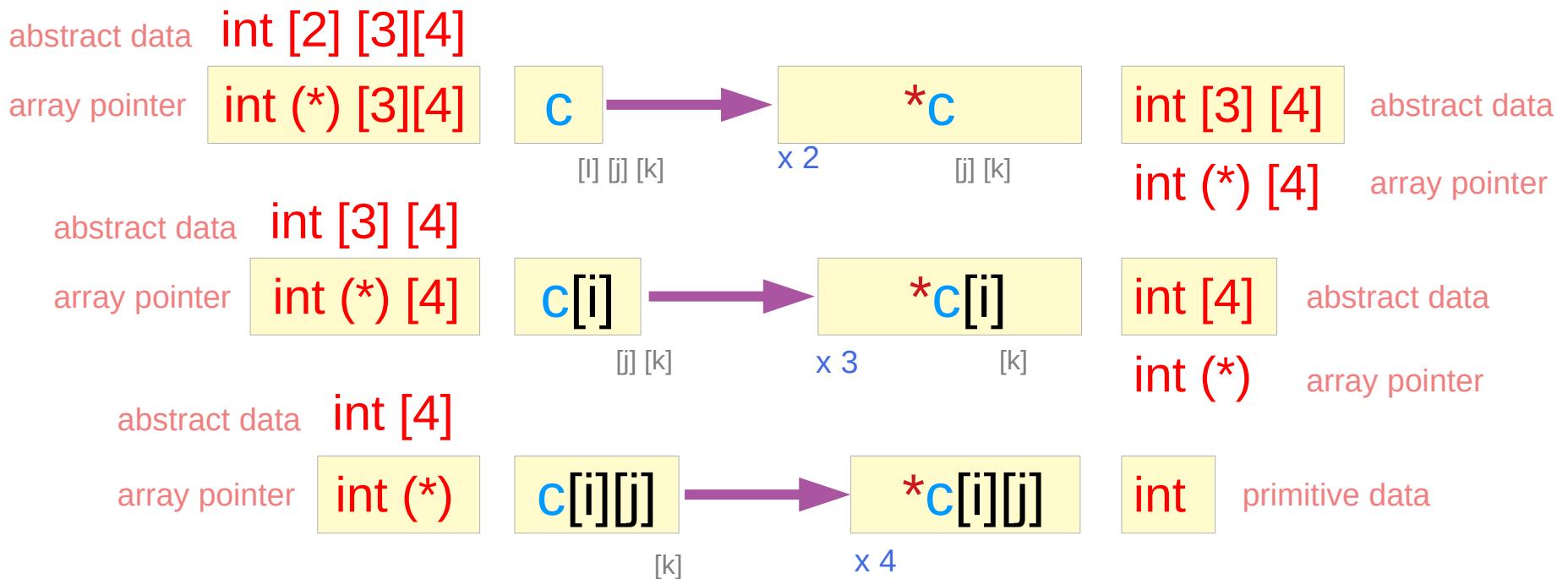
$\text{sizeof}(*c[i])$
 $= \text{sizeof}(c[i][0]) = \text{sizeof(int)} * 4$
[k]

$$c+i = \&c[0][0][0] + i * \text{sizeof}(c[i])$$

$\text{sizeof}(*c)$
 $= \text{sizeof}(c[0]) = \text{sizeof(int)} * 3 * 4$
[j] [k]

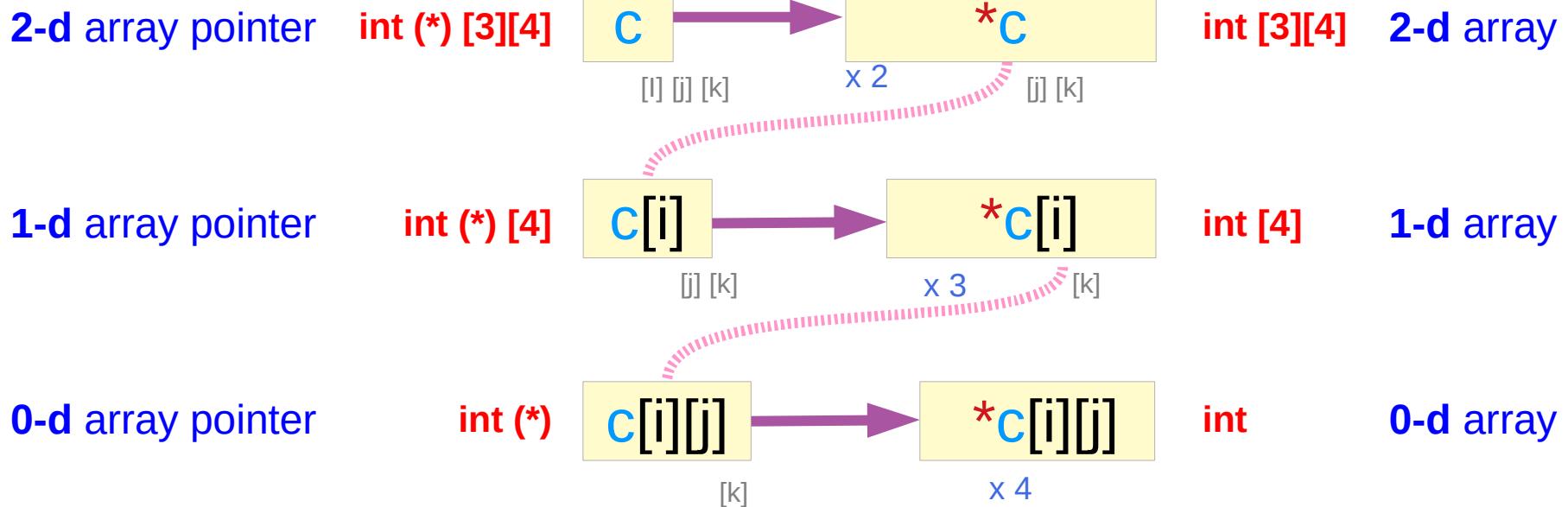
Types in a multi-dimensional 3-d array

```
int c [2][3][4];
```



Virtual array pointers and abstract data

```
int c [2][3][4];
```



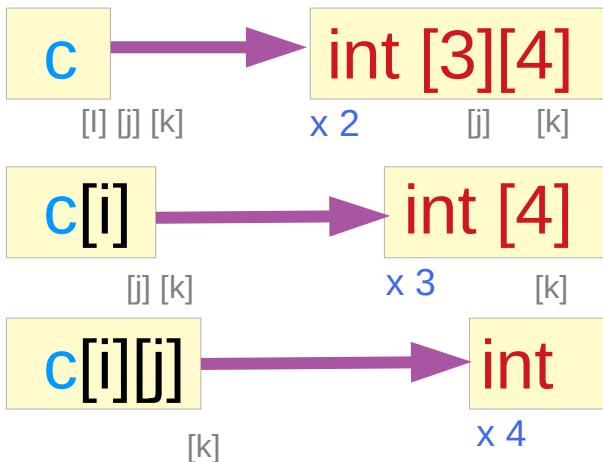
all these pointers are virtual, and
take no actual memory locations

exploiting the **contiguity** of
allocated memory locations

Abstract data sizes

```
int c [2][3][4];
```

the size of a pointer type is fixed
Here, the sizes of virtual pointers are shown
i.e, the sizes of different abstract data types



sizeof(c)	= sizeof(int [2][3][4])
sizeof(*c)	= sizeof(int [3][4])
sizeof(c[i])	= sizeof(int [3][4])
sizeof(*c[i])	= sizeof(int [4])
sizeof(c[i][j])	= sizeof(int [4])
sizeof(*c[i][j])	= sizeof(int)

all are sizes of arrays

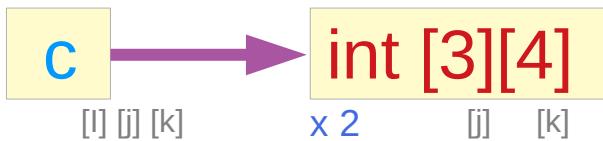
c, c[i], c[i][j] are virtual array pointers
and they are also abstract data (arrays)

when sizes are considered,
view them as abstract data (arrays)

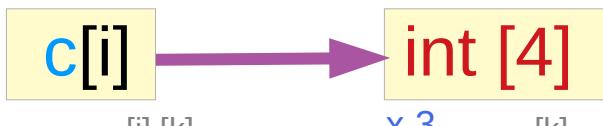
Virtual array pointer sizes and abstract data sizes

```
int c [2][3][4];
```

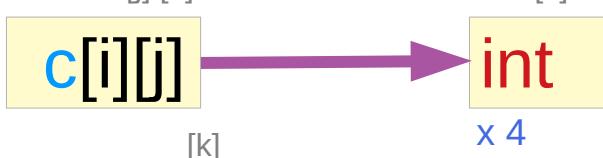
size of a virtual array pointer = size of the pointed abstract data type * the number of such types



$$\text{sizeof}(\text{c}) = \text{sizeof}(*c) * 2$$



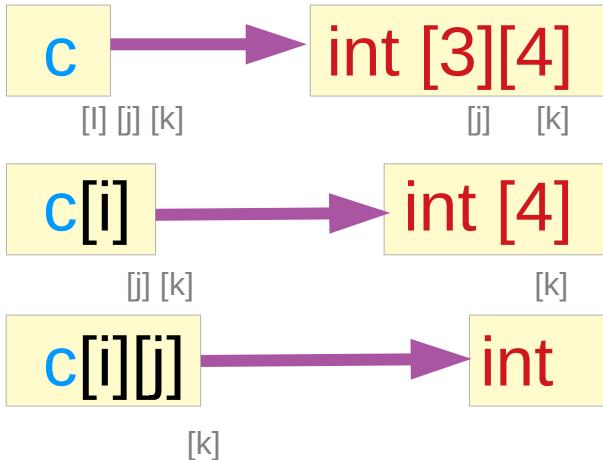
$$\text{sizeof}(\text{c}[i]) = \text{sizeof}(*c[i]) * 3$$



$$\text{sizeof}(\text{c}[i][j]) = \text{sizeof}(*c[i][j]) * 4$$

Sizes of array pointer types

```
int c [2][3][4];
```



not real array pointers
virtual array pointers

`c` `int (*)[3][4]` = `sizeof(c)`
`sizeof(int (*) [3][4])` = pointer size $\neq \text{sizeof}(c)$

`c[i]` `int (*) [4]` = `sizeof(c[i])`
`sizeof(int (*) [4])` = pointer size $\neq \text{sizeof}(c[i])$

`c[i][j]` `int [4]` = `sizeof(c[i][j])`
`sizeof(int [4])` = pointer size $\neq \text{sizeof}(c[i][j])$

4 bytes for 32-bit machines
8 bytes for 64-bit machines

Virtual array pointer increment size

```
int c [2][3][4];
```

c points to a **2-d** array
increment size: `sizeof(int[2][3][4])`

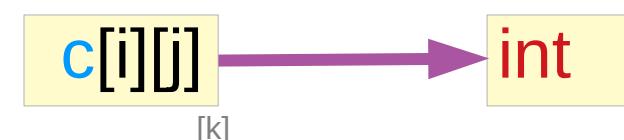
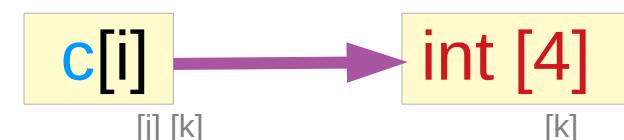
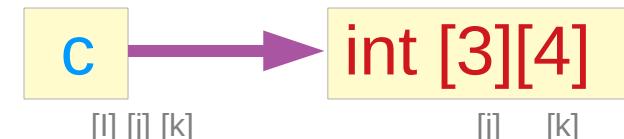
`c[i]` points to an **1-d** array
increment size: `sizeof(int[3][4])`

`c[i][j]` points to an integer
increment size: `sizeof(int[4])`

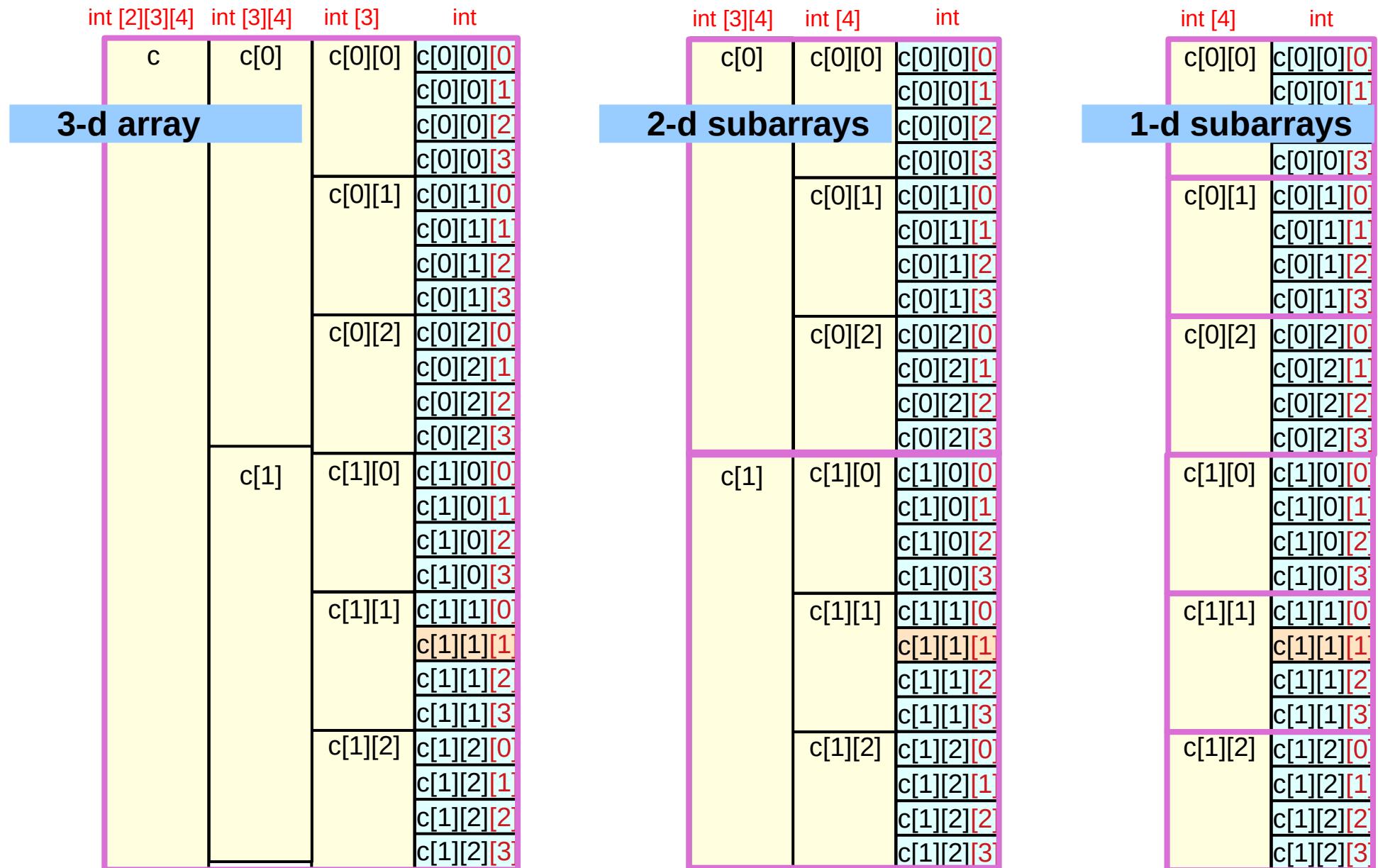
`int (*) [3][4]`

`int (*) [4]`

`int (*)`

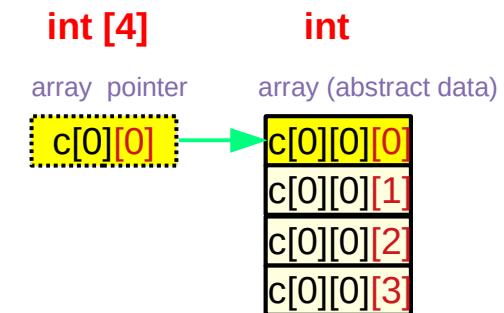
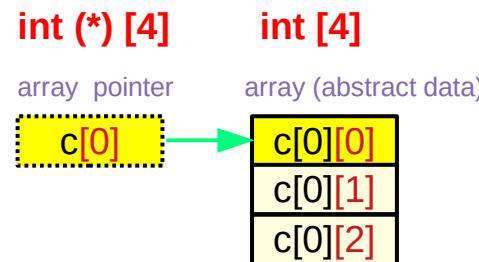
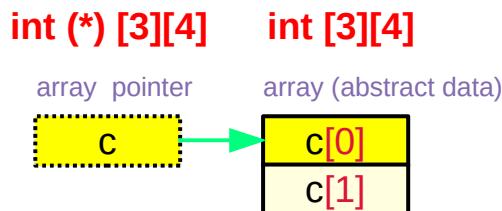


Subarrays in a 3-d array



Virtual array pointer c, c[0], c[0][0] – types and sizes

Types – array pointers



Sizes – abstract data

`sizeof(c)`
`sizeof(int [2][3][4])`
`sizeof(int) * 2 * 3 * 4`

`sizeof(c[0])`
`sizeof(int [3][4])`
`sizeof(int) * 3 * 4`

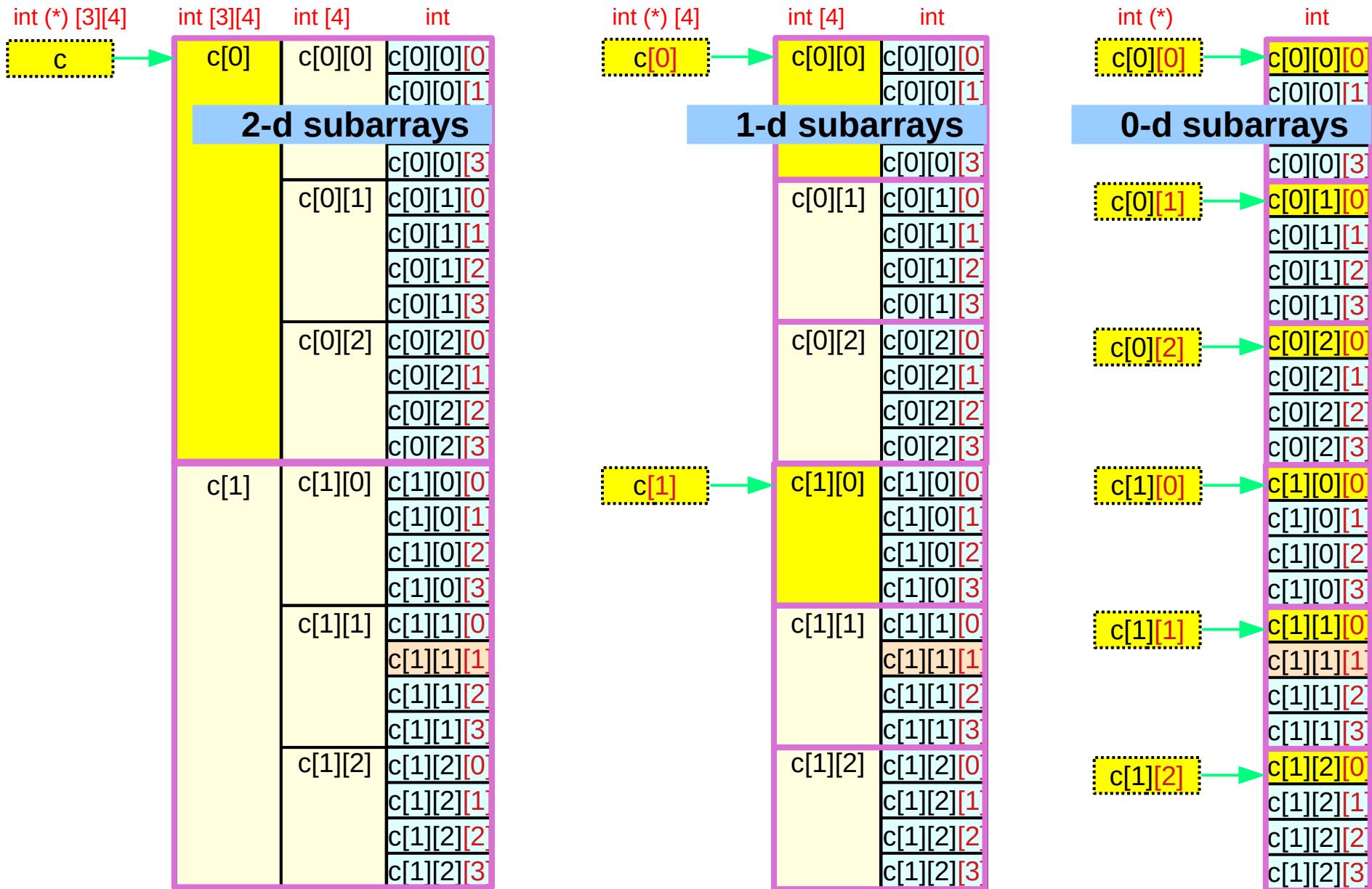
`sizeof(c[0][0])`
`sizeof(int [4])`
`sizeof(int) * 4`

`sizeof(int [2][3][4]) = 96`
`sizeof(int (*)[3][4]) = 4 / 8`

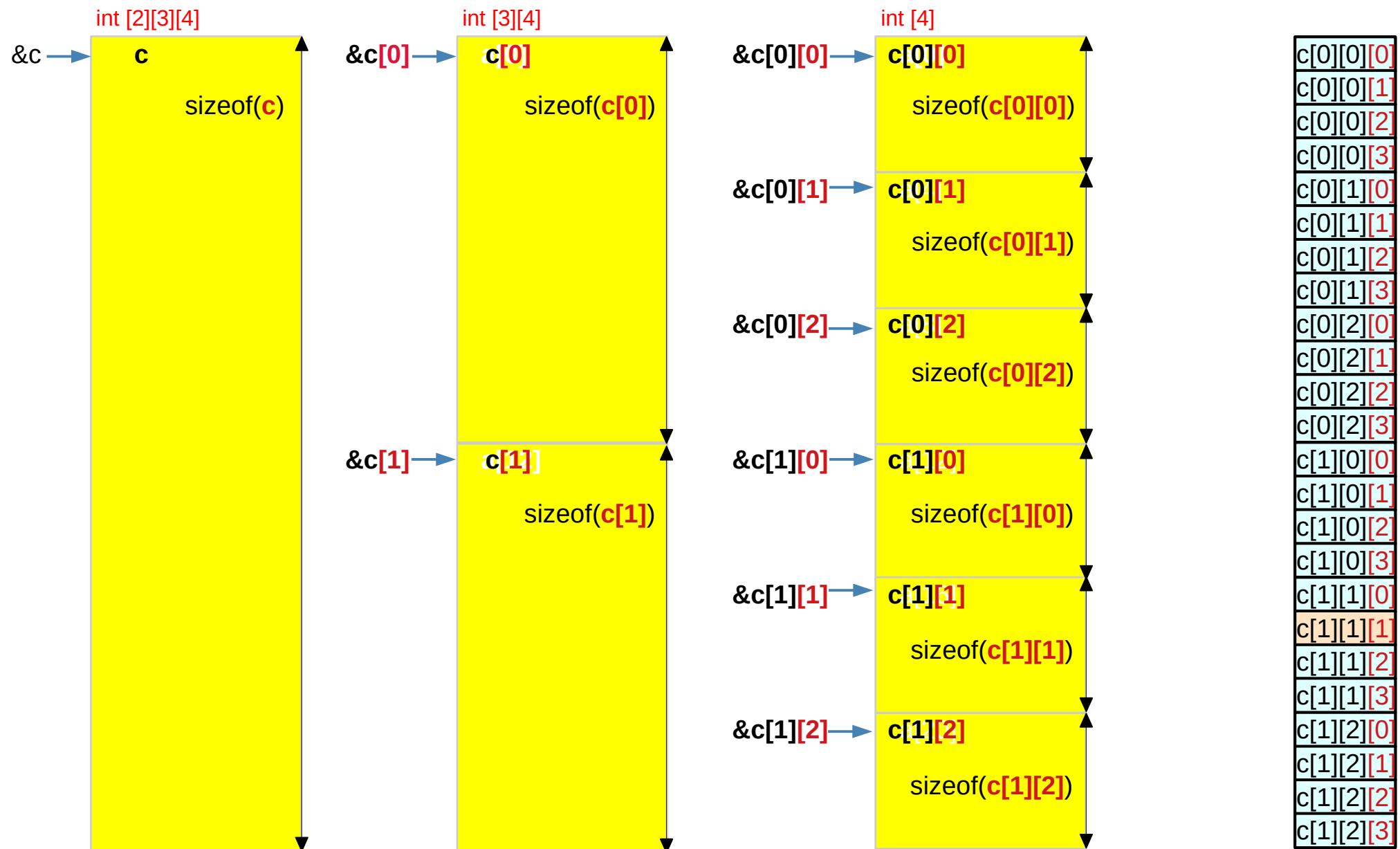
`sizeof(int [3][4]) = 48`
`sizeof(int (*)[4]) = 4 / 8`

`sizeof(int [4]) = 16`
`sizeof(int (*)) = 4 / 8`

Pointers to subarrays in a 3-d array



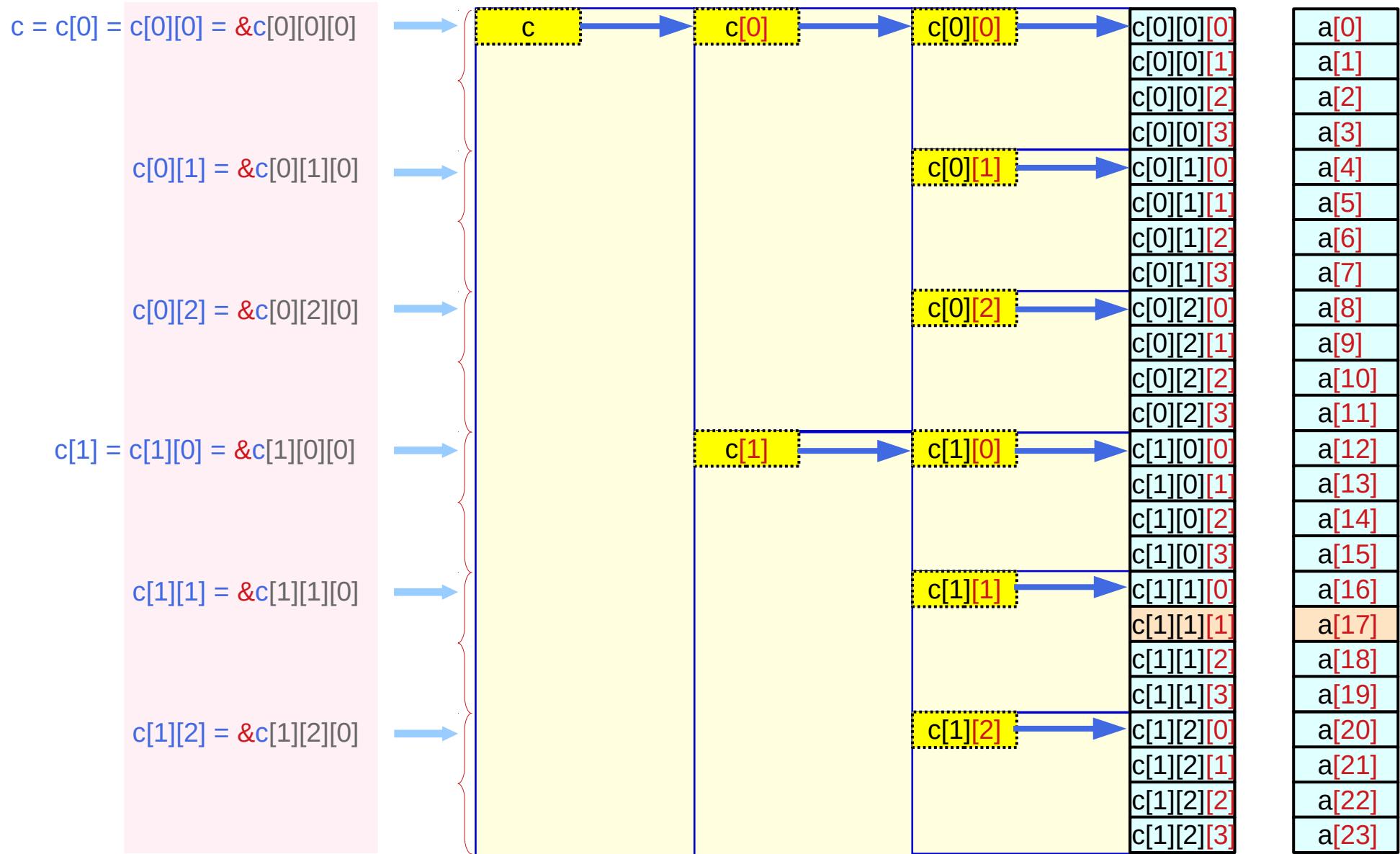
Abstract Data c , $c[i]$, $c[i][j]$ – start addresses and sizes



Virtual array pointers – types, sizes, and values

<code>int c[2][3][4];</code>	c[i][j]	c[i][j][0]	
type	<code>int [4] int (*)</code>	<code>int int</code>	<ul style="list-style-type: none">• abstract data type• array pointer type
size	<code>sizeof(c[i][j]) =</code>	<code>sizeof(c[i][j][0]) * 4</code>	<code>= sizeof(int) * 4</code>
value (address)	<code>c[i][j] =</code>	<code>&c[i][j][0]</code>	
<code>int c[2][3][4];</code>	c[i]	c[i][0]	
type	<code>int [3][4] int (*)[4]</code>	<code>int [4] int (*)</code>	<ul style="list-style-type: none">• abstract data type• array pointer type
size	<code>sizeof(c[i]) =</code>	<code>sizeof(c[i][0]) * 3</code>	<code>= sizeof(int) * 4 * 3</code>
value (address)	<code>c[i] =</code>	<code>&c[i][0][0]</code>	
<code>int c[2][3][4];</code>	c	c[0]	
type	<code>int [2][3][4] int (*)[3][4]</code>	<code>int [3][4] int (*)[4]</code>	<ul style="list-style-type: none">• abstract data type• array pointer type
size	<code>sizeof(c) =</code>	<code>sizeof(c[0]) * 2</code>	<code>= sizeof(int) * 4 * 3 * 2</code>
value (address)	<code>c =</code>	<code>&c[0][0][0]</code>	

Virtual array pointer c, c[i], c[i][j] – values (addresses)



Virtual array pointer c, c[i], c[i][j] – vertical displacement

$c = c[0] = c[0][0] = \&c[0][0][0]$



For address values

$c[0][1] = \&c[0][1][0]$



Horizontal displacements
are not counted

$c[0][2] = \&c[0][2][0]$



Only vertical displacements
are considered

$c[1] = c[1][0] = \&c[1][0][0]$



$c[0][0][0]$	a[0]
$c[0][0][1]$	a[1]
$c[0][0][2]$	a[2]
$c[0][0][3]$	a[3]
$c[0][1][0]$	a[4]
$c[0][1][1]$	a[5]
$c[0][1][2]$	a[6]
$c[0][1][3]$	a[7]
$c[0][2][0]$	a[8]
$c[0][2][1]$	a[9]
$c[0][2][2]$	a[10]
$c[0][2][3]$	a[11]
$c[1][0][0]$	a[12]
$c[1][0][1]$	a[13]
$c[1][0][2]$	a[14]
$c[1][0][3]$	a[15]
$c[1][1][0]$	a[16]
$c[1][1][1]$	a[17]
$c[1][1][2]$	a[18]
$c[1][1][3]$	a[19]
$c[1][2][0]$	a[20]
$c[1][2][1]$	a[21]
$c[1][2][2]$	a[22]
$c[1][2][3]$	a[23]

$c[1][1] = \&c[1][1][0]$



$c[1][2] = \&c[1][2][0]$



Virtual array pointer c, c[i], c[i][j] – values and types

$c = c[0] = c[0][0] = \&c[0][0][0]$

means
→

$c[0][1] = \&c[0][1][0]$

means
→

$c[0][2] = \&c[0][2][0]$

means
→

$c[1] = c[1][0] = \&c[1][0][0]$

means
→

$c[1][1] = \&c[1][1][0]$

means
→

$c[1][2] = \&c[1][2][0]$

means
→

$\text{value}(c) = \text{value}(c[0]) = \text{value}(c[0][0]) = \text{value}(\&c[0][0][0])$ $\text{type}(c) \neq \text{type}(c[0]) \neq \text{type}(c[0][0]) = \text{type}(\&c[0][0][0])$ $\text{int } (*) [3][4] \quad \text{int } (*) [4] \quad \text{int } * \quad \text{int } *$	
--	--

	$\text{value}(c[0][1]) = \text{value}(\&c[0][1][0])$ $\text{type}(c[0][1]) = \text{type}(\&c[0][1][0])$ $\text{int } * \quad \text{int } *$
--	---

	$\text{value}(c[0][2]) = \text{value}(\&c[0][2][0])$ $\text{type}(c[0][2]) = \text{type}(\&c[0][2][0])$ $\text{int } * \quad \text{int } *$
--	---

	$\text{value}(c[1]) = \text{value}(c[1][0]) = \text{value}(\&c[1][0][0])$ $\text{type}(c[1]) \neq \text{type}(c[1][0]) = \text{type}(\&c[1][0][0])$ $\text{int } (*) [4] \quad \text{int } * \quad \text{int } *$
--	---

	$\text{value}(c[1][1]) = \text{value}(\&c[1][1][0])$ $\text{type}(c[1][1]) = \text{type}(\&c[1][1][0])$ $\text{int } * \quad \text{int } *$
--	---

	$\text{value}(c[1][2]) = \text{value}(\&c[1][2][0])$ $\text{type}(c[1][2]) = \text{type}(\&c[1][2][0])$ $\text{int } * \quad \text{int } *$
--	---

Summary of virtual array pointers in a 3-d array

$$c[i] \equiv *(c + i)$$

$$c[i][j] \equiv *(c[i] + j)$$

$$c[i][j][k] \equiv *(c[i][j] + k)$$

int (*) [3][4] 2-d array pointer `c`
int [2] [3][4] 3-d array name `c`

int (*) [4] 1-d array pointers `c[i]`
Int [3] [4] 2-d array names `c[i]`

int (*) 0-d array pointers `c[i][j]`
int [4] 1-d array names `c[i][j]`

address value $c + i$

`&c[0][0][0] + i * sizeof(*c)`
`&c[0][0][0] + i * sizeof(c[0])`
`&c[0][0][0] + i * 4 * 3 * 4`

address value $c[i] + j$

`&c[i][0][0] + j * sizeof(*c[i])`
`&c[i][0][0] + j * sizeof(c[i][0])`
`&c[i][0][0] + j * 4 * 4`

address value $c[i][j] + k$

`&c[i][j][0] + k * sizeof(*c[i][j])`
`&c[i][j][0] + k * sizeof(c[i][j][0])`
`&c[i][j][0] + k * 4`

leading elements
`c[0][0][0]`

leading elements
`c[0][0][0]`

`c[1][0][0]`

leading elements
`c[0][0][0]`
`c[0][1][0]`
`c[0][2][0]`
`c[1][0][0]`
`c[1][1][0]`
`c[1][2][0]`

Dual type constraints in multi-dimensional arrays

Virtual pointers to subarrays in a 3-d array

virtual 2-d array pointer

`sizeof(c) =
sizeof(c[0]) * 2`

`int c [2][3][4];`

`int (*) [3][4]`

`c`

`int [3][4]`

`c[0]`

`int [4]`

`c[0][0]`

`int`

`c[0][0][0]`

`c[0][0][1]`

`c[0][0][2]`

`c[0][0][3]`

`c[0][1][0]`

`c[0][1][1]`

`c[0][1][2]`

`c[0][1][3]`

`c[0][2][0]`

`c[0][2][1]`

`c[0][2][2]`

`c[0][2][3]`

the first 2-d subarray

`sizeof(c[0]) =
sizeof(int [3][4])`

virtual 1-d array pointer

`sizeof(c[0]) =
sizeof(c[0][0]) * 3`

`int (*) [4]`

`c[0]`

`int [4]`

`c[0][0]`

`int`

`c[0][0][0]`

`c[0][0][1]`

`c[0][0][2]`

`c[0][0][3]`

the first 1-d subarray

`sizeof(c[0][0]) =
sizeof(int [4])`

virtual 0-d array pointer

`sizeof(c[0][0]) =
sizeof(c[0][0][0]) * 4`

`int (*)`

`c[0][0]`

`int`

`c[0][0][0]`

`c[0][0][1]`

`c[0][0][2]`

`c[0][0][3]`

the first 0-d subarray

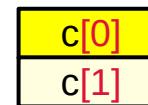
`sizeof(c[0][0][0]) =
sizeof(int)`

Contiguous subarrays

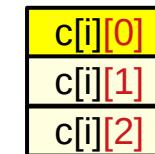
array pointer
c
int [2][3][4] Contiguous 2*3*4 integers

int c [2][3][4];

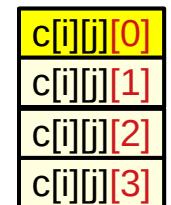
array pointer array (abstract data)
c → **c[0]** $i=0,1$ 2 $c[i]$'s, contiguous
int (*) [3][4] **int [3][4]** Contiguous 3*4 integers



for a given i array pointer array (abstract data)
c[i] → **c[i][0]** $j=0,1,2$ 3 $c[i][j]$'s, contiguous
int (*) [4] **int [4]** Contiguous 4 integers



for a given i, j array pointer array (abstract data)
c[i][j] → **c[i][j][0]** $k=0,1,2,3$ 4 $c[i][j][k]$'s, contiguous
int (*) **int**



Virtual assignments

```
int c [2][3][4];
```

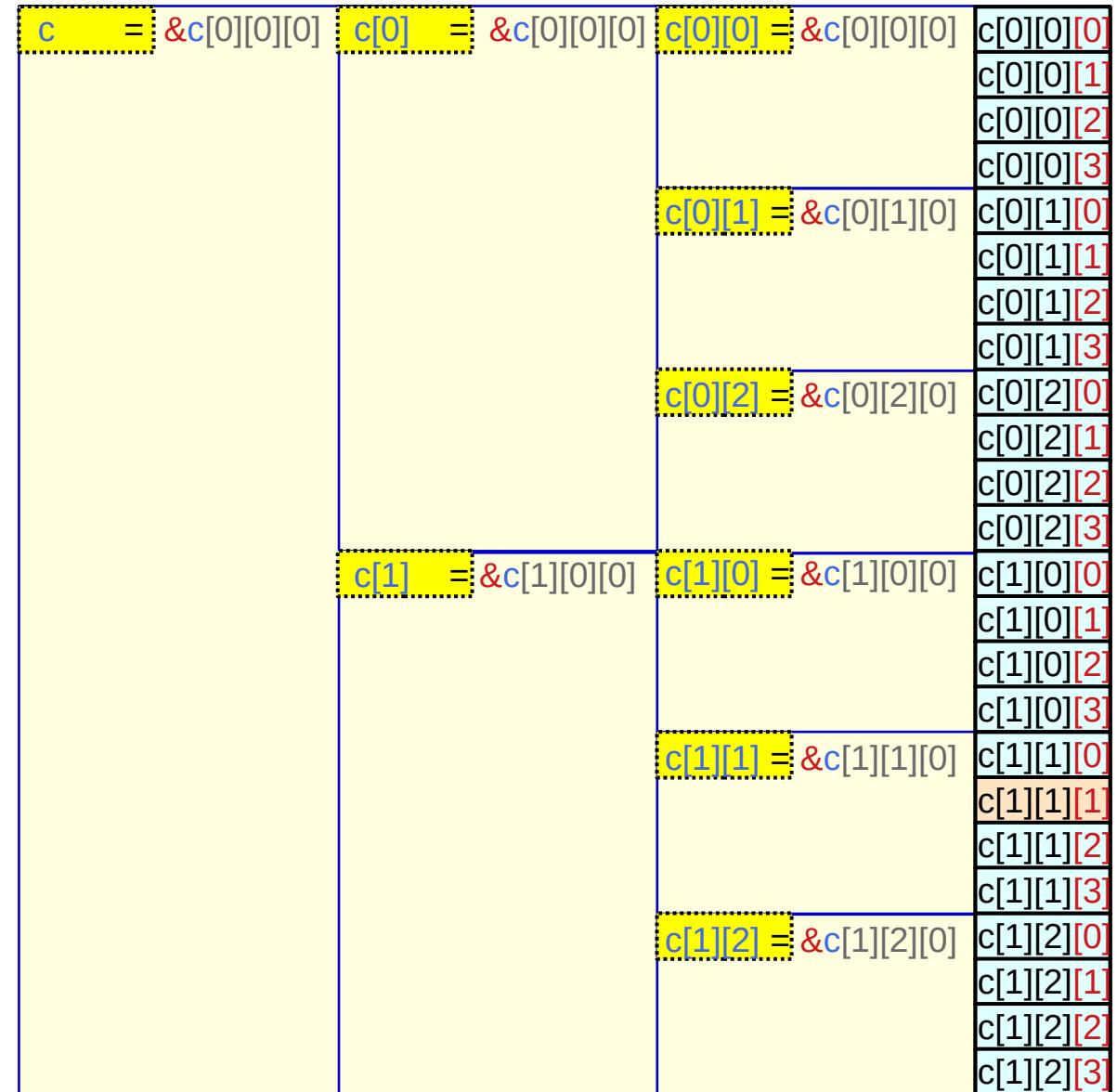


```
c [i][j][k];
```

virtual assignments

c	←	&c[0][0][0]
c[i]	←	&c[i][0][0]
c[i][j]	←	&c[i][j][0]

row major ordering
contiguous linear layout



Virtual assignments and type casts

```
int c [2][3][4];
```



```
c [i][j][k];
```

c \leftarrow (int (*)[3][4])
c[i] \leftarrow (int (*)[4])
c[i][j] \leftarrow (int (*))

type casts

&c[0][0][0]
&c[i][0][0]
&c[i][j][0]

address values

virtual assignments

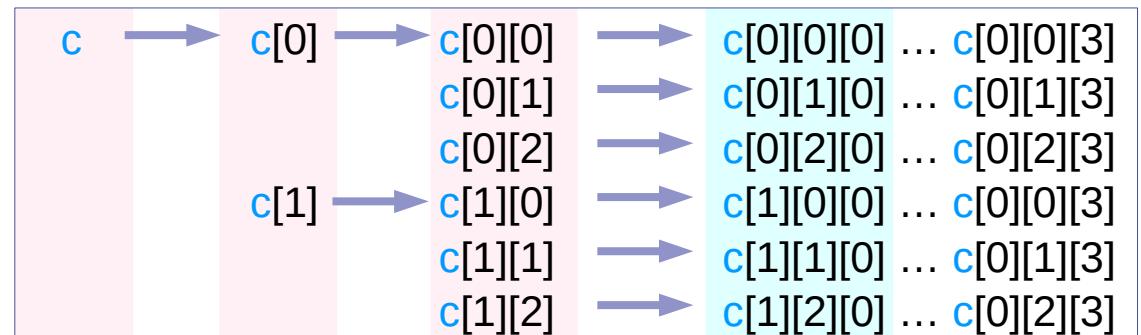
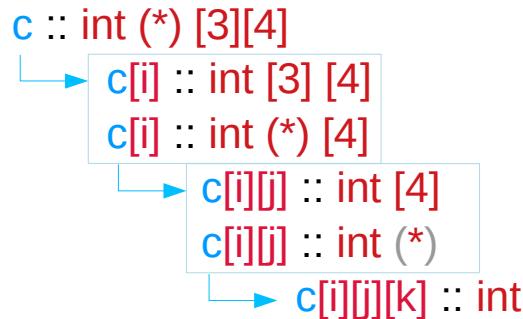
c \leftarrow &c[0][0][0]
c[i] \leftarrow &c[i][0][0]
c[i][j] \leftarrow &c[i][j][0]

row major ordering
contiguous linear layout

if c, c[i], c[i][j] were real pointer variables,
type casts would be needed

Dual types of **c**, **c[i]**, **c[i][j]**

int **c [2][3][4];**



int [2] [3][4] int [3] [4]
int (*) [3][4] int (*) [4]

pointers to
a 2-d array

int [4] int (*)
int (*)

pointers to
a 1-d array 1-d array
names

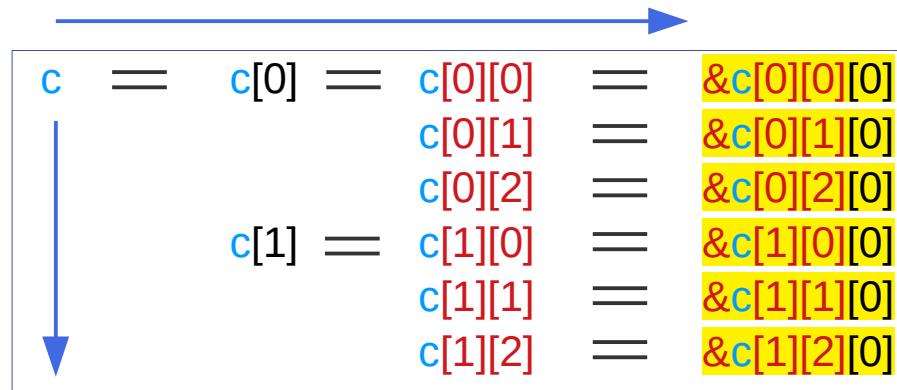
int ... int
int ... int

leading element
of 4-integer array

Values of c , $c[i]$, $c[i][j]$

```
int c [2][3][4];
```

virtual array pointers have address values
in each row in the following figure
have the same address value



Horizontal displacements are not counted
only vertical displacements are considered
for address values

virtual assignments

```
c[i][j] = &c[i][j][0]
c[i]   = &c[i][0][0]
c     = &c[0][0][0]
```

Finding address values of **c**, **c[i]**, **c[i][j]**

```
int c [2][3][4];
```

append [0] to the right

c	\equiv	c[0]	\equiv	c[0][0]	\equiv	[0]	&c[0][0][0]
				c[0][1]	\equiv	[0]	&c[0][1][0]
				c[0][2]	\equiv	[0]	&c[0][2][0]
		c[1]	\equiv	c[1][0]	\equiv	[0]	&c[1][0][0]
				c[1][1]	\equiv	[0]	&c[1][1][0]
				c[1][2]	\equiv	[0]	&c[1][2][0]

int (*) [3][4]

int (*) [4]

int [4]

int

c[0][0][0] :
leading
elements
of c

&c[0][0][0]

c[i][0][0] :
leading
elements
of c[i]

&c[0][0][0]

&c[1][0][0]

c[i][j][0] :
leading
elements
of c[i][j]

&c[0][0][0]

&c[0][1][0]

&c[0][2][0]

&c[1][0][0]

&c[1][1][0]

&c[1][2][0]

Finding sub-array names with the address $\&c[i][j][0]$

```
int c [2][3][4];
```

delete [0] from the right

$\&c[0][0][0]$	$\stackrel{-[0]}{=}$	$c[0][0]$	$\stackrel{-[0]}{=}$	$c[0]$	$\stackrel{-[0]}{=}$	c
$\&c[0][1][0]$	$\stackrel{-[0]}{=}$	$c[0][1]$				
$\&c[0][2][0]$	$\stackrel{-[0]}{=}$	$c[0][2]$				
$\&c[1][0][0]$	$\stackrel{-[0]}{=}$	$c[1][0]$	$\stackrel{-[0]}{=}$	$c[1]$		
$\&c[1][1][0]$	$\stackrel{-[0]}{=}$	$c[1][1]$				
$\&c[1][2][0]$	$\stackrel{-[0]}{=}$	$c[1][2]$				

int

int [4]

int (*) [4]

int (*) [3][4]

$c[0][0][0]$ is the leading element of $c[0][0]$, $c[0]$, c

$c[0][1][0]$ is the leading element of $c[0][1]$

$c[0][2][0]$ is the leading element of $c[0][2]$

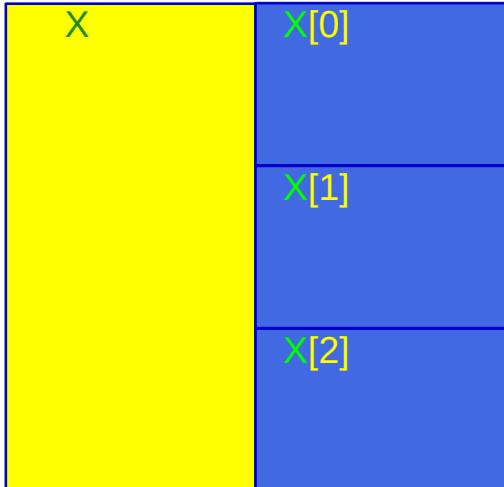
$c[1][0][0]$ is the leading element of $c[1][0]$, $c[1]$

$c[1][1][0]$ is the leading element of $c[1][1]$

$c[1][2][0]$ is the leading element of $c[1][2]$

Dual type – virtual array pointer and abstract array

Abstract data (array) $p[i]$

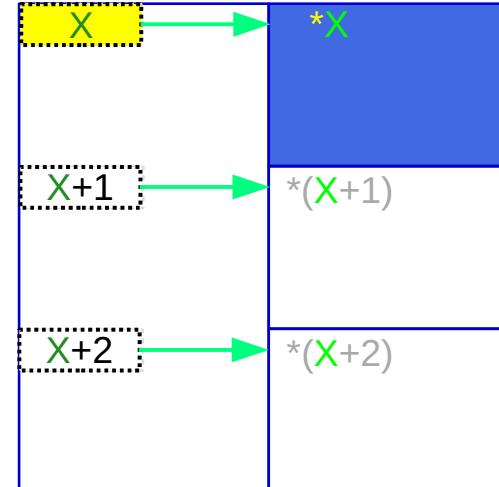


array (abstract data)

$c[i][j]$ starts from $\&c[i][j][0]$
 $c[i]$ starts from $\&c[i][0]$
 c starts from $\&c[0]$

$$\begin{aligned} \&c[i][j] &= & \&c[i][j][0] \\ \&c[i] &= & \&c[i][0] \\ \&c &= & \&c[0] \end{aligned}$$

Virtual array pointer $p[i]$



array pointer

$c[i][j]$ points to $c[i][j][0]$
 $c[i]$ points to $c[i][0]$
 c points to $c[0]$

address value

$$\begin{aligned} c[i][j] &= \&c[i][j][0] \\ c[i] &= \&c[i][0] \\ c &= \&c[0] \end{aligned}$$

Virtual array pointer values and addresses

$\&c[i][j][0]$	$= c[i][j]$
$\&c[i][0]$	$= c[i]$
$\&c[0]$	$= c$



$\&c[i][j][0]$	$= \&c[i][j]$
$\&c[i][0]$	$= \&c[i]$
$\&c[0]$	$= \&c$



$c[i][j]$	$= \&c[i][j]$
$c[i]$	$= \&c[i]$
c	$= \&c$

Virtual array pointer

array pointer	array (abstract data)
$c[i][j]$	points to $c[i][j][0]$
$c[i]$	points to $c[i][0]$
c	points to $c[0]$

address value

Abstract data (array)

array (abstract data)	array (abstract data)
$c[i][j]$	starts from $\&c[i][j][0]$
$c[i]$	starts from $\&c[i][0]$
c	starts from $\&c[0]$

array (abstract data)	Address of an array pointer
$c[i][j]$	pointer value = pointer address $\&c[i][j]$
$c[i]$	pointer value = pointer address $\&c[i]$
c	pointer value = pointer address $\&c$

$c[0] = c[0][0]$ relation

```
int    c [2][3][4];
```

$c = c[0] = c[0][0] = \&c[0][0][0]$

$\text{value}(c[0]) = \&c[0][0][0]$

$\text{value}(c[0][0]) = \&c[0][0][0]$

$\text{type}(c[0]) = \text{int } (*)[4]$

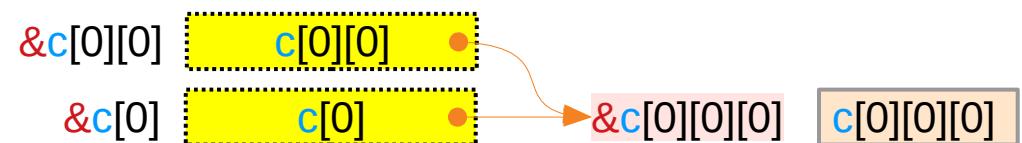
$\text{type}(c[0][0]) = \text{int } [4]$

$c[0] = c[0][0]$ means

$\text{value}(c[0]) = \text{value}(c[0][0])$

$c[0] = c[0][0]$ does not mean

$\text{type}(c[0]) = \text{type}(c[0][0])$



Addresses and Values of $c[0]$ and $c[0][0]$

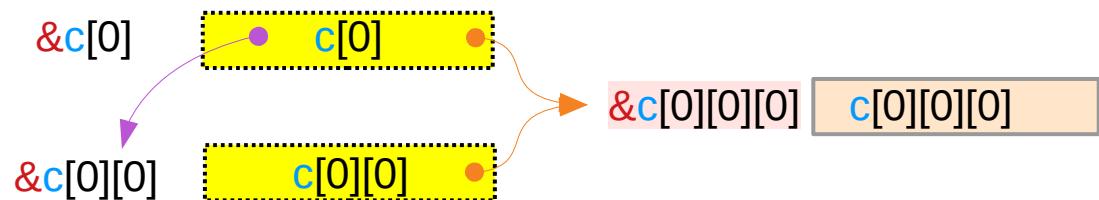
```
int c [2][3][4];
```

$c[0] \rightarrow c[0][0]$

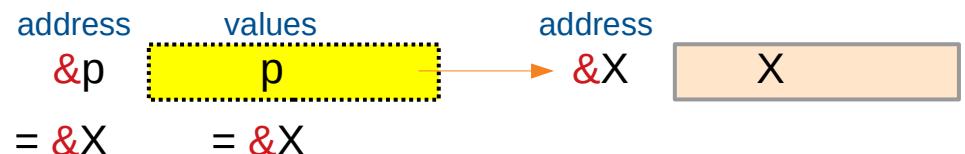


$\&c[0] \rightarrow \&c[0][0]$

$c[0] == c[0][0] == \&c[0][0][0]$



A virtual pointer's address and value are the same



c[i] and **c[i][0]** point to the same **c[i][0][0]**

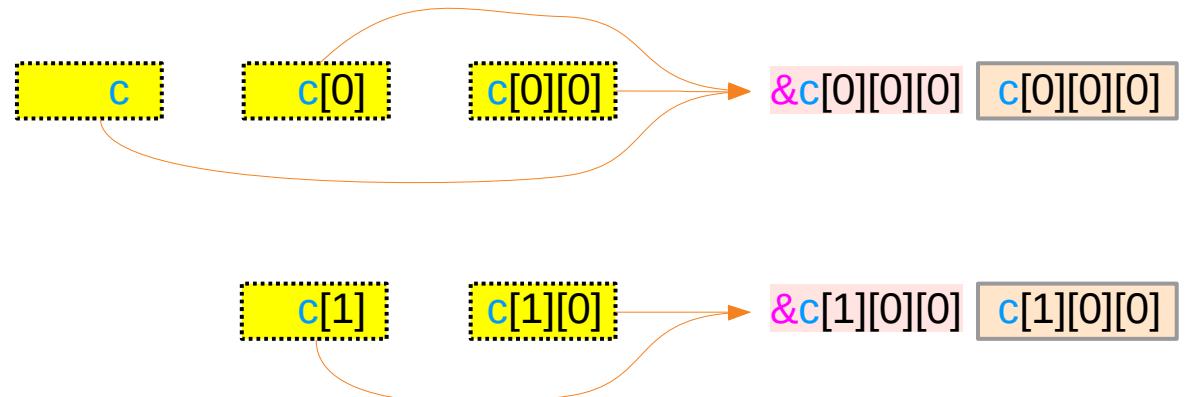
```
int c [2][3][4];
```

c = c[0] = c[0][0] = &c[0][0][0]

int(*)[3][4] int(*)[4] int(*) int ← value
← type

c[1] = c[1][0] = &c[1][0][0]

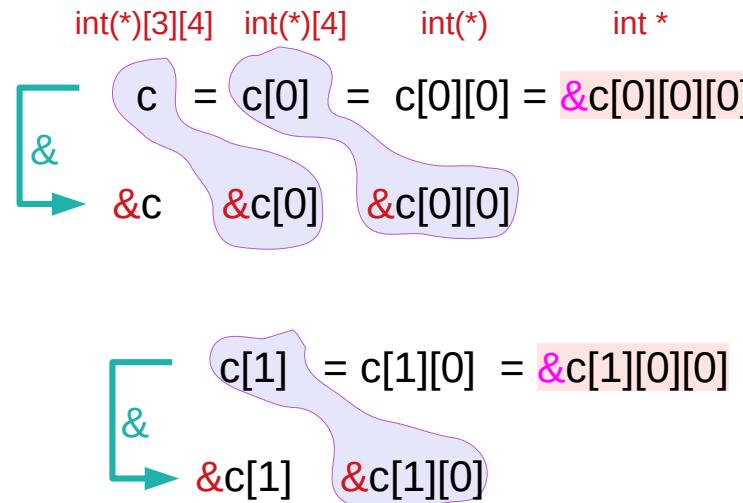
int(*)[4] int(*) int ← value
← type



These virtual pointers have different types
but the same value (address)

`&c[i][0]` and `&c[i][0][0]` – equivalence relations

```
int      c [2][3][4];
```



equivalences

$$\begin{aligned}c &\equiv \&c[0], \\c[0] &\equiv \&c[0][0] \\c[0][0] &\equiv \&c[0][0][0]\end{aligned}$$

equivalences

$$\begin{aligned}c[1] &\equiv \&c[1][0] \\c[1][0] &\equiv \&c[1][0][0]\end{aligned}$$

Horizontal displacements are not counted
only vertical displacements are considered
for address values

equivalences

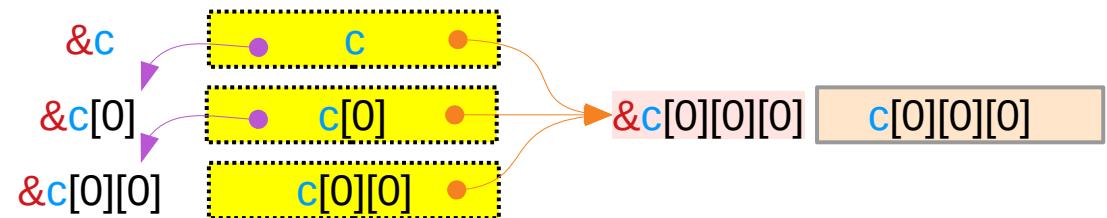
$$\begin{aligned}c &\equiv \&c[0], \\c[i] &\equiv \&c[i][0] \\c[i][0] &\equiv \&c[i][0][0]\end{aligned}$$

$c[i] = &c[i]$ and $c[i][0] = &c[i][0]$

```
int c [2][3][4];
```

$$\begin{aligned} c &= \boxed{c} & c[0] &= \boxed{c[0]} & c[0][0] &= \boxed{c[0][0]} \\ \parallel & & \parallel & & \parallel & \&c[0][0][0] \\ \&c &= \boxed{\&c} &= \boxed{\&c[0]} &= \boxed{\&c[0][0]} \end{aligned}$$

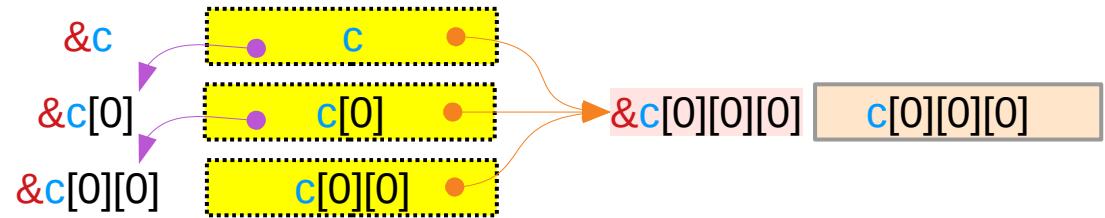
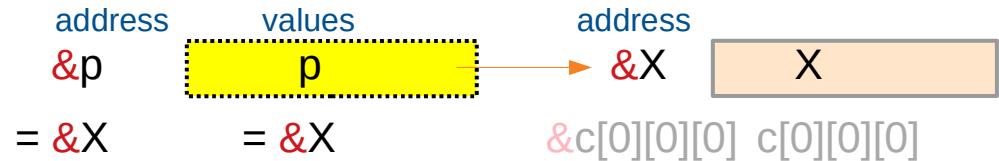
$$\begin{aligned} c[1] &= \boxed{c[1]} & c[1][0] &= \boxed{c[1][0]} & c[1][0][0] &= \boxed{c[1][0][0]} \\ \parallel & & \parallel & & \parallel & \&c[1][0][0] \\ \&c[1] &= \boxed{\&c[1]} &= \boxed{\&c[1][0]} \end{aligned}$$



$c[i] = &c[i]$ and $c[i][0] = &c[i][0]$

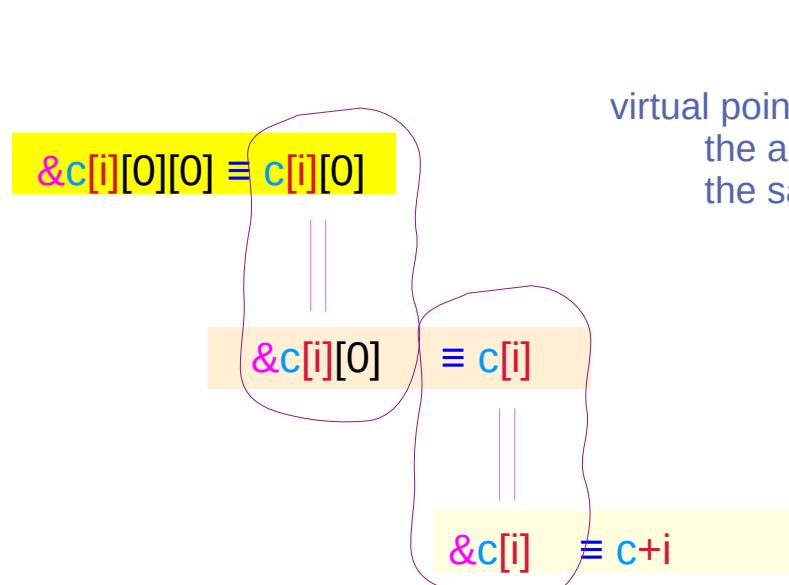
```
int c [2][3][4];
```

A virtual pointer's address and value are the same



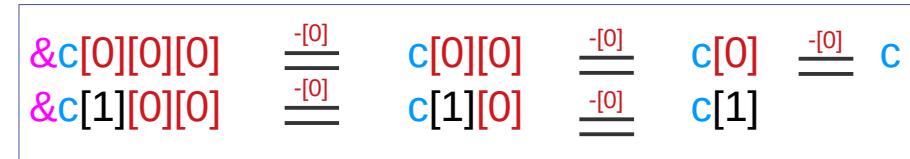
c[i]

Array Pointers to $c[i][0][0]$



$$\begin{aligned} &= c + i * \text{sizeof}(*c) \\ &= \&c[0][0][0] + i * 3 * 4 \end{aligned}$$

delete [0] from the right



Array Pointers to $c[i][j][0]$

$$\&c[i][j][0] = c[i][j]$$

$$\&c[i][j]$$

$$\equiv c[i] + j$$

$$= c[i] + j * \text{sizeof}(*c[i])$$

$$= c + i * \text{sizeof}(c) + j * 4$$

$$= \&c[0][0][0] + i * 3 * 4 + j * 4$$

delete [0] from the right

$\&c[0][0][0]$	$\equiv^{[0]}$	$c[0][0]$	$\equiv^{[0]}$	$c[0]$	$\equiv^{[0]}$	c
$\&c[0][1][0]$	$\equiv^{[0]}$	$c[0][1]$				
$\&c[0][2][0]$	$\equiv^{[0]}$	$c[0][2]$				
$\&c[1][0][0]$	$\equiv^{[0]}$	$c[1][0]$	$\equiv^{[0]}$	$c[1]$		
$\&c[1][1][0]$	$\equiv^{[0]}$	$c[1][1]$				
$\&c[1][2][0]$	$\equiv^{[0]}$	$c[1][2]$				

Sub-array properties in multi-dimensional arrays

int c [2][3][4];  3-d access c [i][j][k]

2-d array pointer c int (*) [3][4]

1-d array pointers c[i] int (*) [4]

0-d array pointers c[i][j] int (*)

Hierarchical sub-arrays in a 3-d array

```
int    c [L][M][N];
```

c [i][j][k]

left-to-right associativity

Array Names and Types

Pointers to hierarchical sub-arrays

c [i] [j][k]
c[i] [j] [k]
c[i][j] [k]

c 3-d array names
c[i] 2-d array names
c[i][j] 1-d array names

int (*) [M][N]	2-d array pointer
int (*) [N]	1-d array pointer
int (*)	0-d array pointer

Associativity and Equivalence Relations

left-to-right associativity

$$((c[i])[j])[k]$$

\equiv

left-to-right associativity

$$*(*(*(c+i)+j)+k)$$

$$X[n]$$

\equiv

$$*(X+n)$$

given $c[i][j]$

$$c[i][j][k]$$

\equiv

$$*(c[i][j]+k)$$

for all k

given $c[i]$

$$c[i][j]$$

\equiv

$$*(c[i]+j)$$

for all j

given c

$$c[i]$$

\equiv

$$*(c+i)$$

for all i

Requirements for the expression $c[i][j][k]$

3 contiguity requirements

for a given $c[i][j]$, for all k

for a given $c[i]$, for all j

for a given c, for all i

$$\begin{aligned}c[i][j][k] &= *(c[i][j]+k) \\c[i][j] &= *(c[i]+j) \\c[i] &= *(c+i)\end{aligned}$$

for a given $c[i][j]$, contiguous $c[i][j][k]$'s

for a given $c[i]$, contiguous $c[i][j]$'s

for a given c, contiguous $c[i]$'s

for a given
subarray pointer contiguous
subarrays

Equivalent requirements for the expression $c[i][j][k]$

for all k

$$c[i][j][k] = *(c[i][j]+k)$$

for all j

$$c[i][j] = *(c[i]+j)$$

for all i

$$c[i] = *(c+i)$$



$$\&c[i][j][k] = c[i][j]+k$$

$$\&c[i][j] = c[i]+j$$

$$\&c[i] = c+i$$

for all k

for all j

for all i



$$c[i][j][0] = *(c[i][j])$$

$$c[i][0] = *(c[i])$$

$$c[0] = *(c)$$

with contiguous subarrays



$$\&c[i][j][0] = c[i][j]$$

$$\&c[i][0] = c[i]$$

$$\&c[0] = c$$

with contiguous subarrays

3-d access pattern $c[i][j][k]$

General requirements

$c[i][j][k]$



$$\begin{aligned}\&c[i][j][k] &= c[i][j]+k \\ \&c[i][j] &= c[i]+j \\ \&c[i] &= c+i\end{aligned}$$

for all k
for all j
for all i

Pointer array approach

```
int** c[2];
int* b[2*3];
int c[2*3*4];
```

$c[i][j][k]$:: int
$c[i][j]$:: int *
$c[i]$:: int **
c	:: int ***

```
c[i] ← &b[i*3]
b[j] ← &a[j*4]
```

with contiguous a, b, c

Explicit
Arrays of pointers with
Multiple Indirection

N-dim Array approach

```
int c[2][3][4];
```

$c[i][j][k]$:: int
$c[i][j]$:: int [4]
$c[i]$:: int [3][4]
c	:: int [2][3][4]

```
c[i][j] ← &c[i][j][0]
c[i] ← &c[i][0][0]
c ← &c[0][0][0]
```

with contiguous c

Implicit
Nested
Virtual Array Pointers

3-d access pattern $c[i][j][k]$ – array pointer approach

General requirements

$c[i][j][k]$



$$\begin{aligned}\&c[i][j][k] &= c[i][j]+k \\ \&c[i][j] &= c[i]+j \\ \&c[i] &= c+i\end{aligned}$$

for all k
for all j
for all i

N-dim array approach

`int c[2][3][4];`

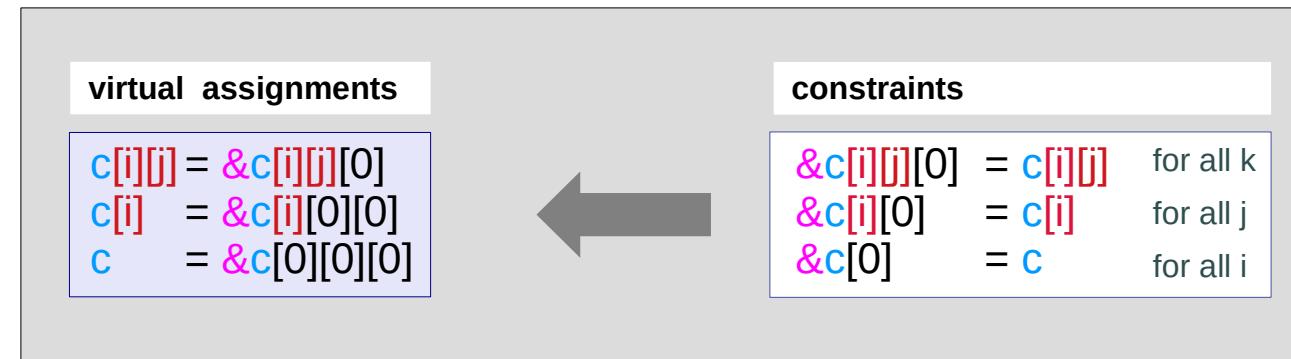
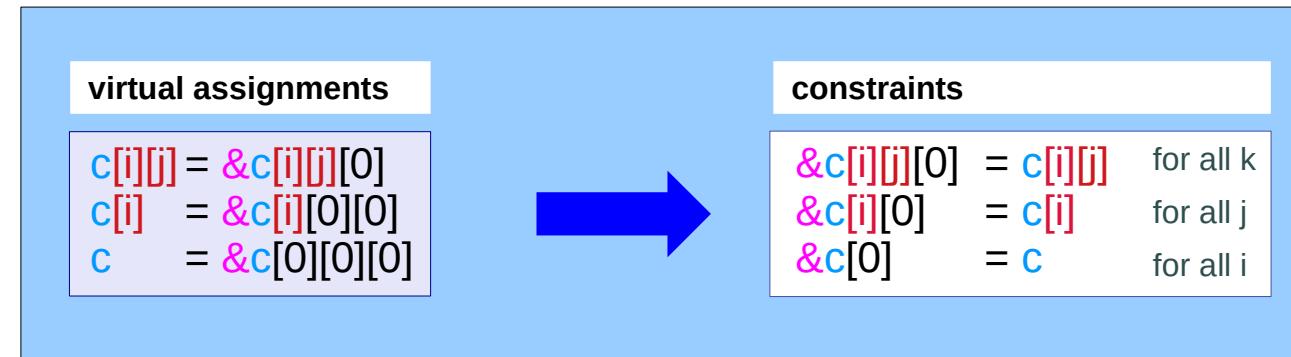
$c[i][j][k] :: \text{int}$
 $c[i][j] :: \text{int}[4]$
 $c[i] :: \text{int}[3][4]$
 $c :: \text{int}[2][3][4]$

$c[i][j] \leftarrow \&c[i][j][0]$
 $c[i] \leftarrow \&c[i][0][0]$
 $c \leftarrow \&c[0][0][0]$

with contiguous c

Implicit
Nested
Virtual Array Pointers

multi-dimensional arrays



General requirements for $c[i][j][k]$

```
int    c [2][3][4];
```

$$\begin{aligned}\&c[i][j][0] &= c[i][j] \\ \&c[i][0] &= c[i] \\ \&c[0] &= c\end{aligned}$$

with contiguous subarrays

$$\begin{aligned}\&c[i][j][k] &= c[i][j]+k \\ \&c[i][j] &= c[i]+j \\ \&c[i] &= c+i\end{aligned}$$

for all k
for all j
for all i



virtual assignments

$$\begin{aligned}c[i][j] &= \&c[i][j][0] \\ c[i] &= \&c[i][0][0] \\ c &= \&c[0][0][0]\end{aligned}$$

Virtual assignments

$$\begin{array}{lll}\text{int (*)} & c[i][j] & = (\text{int (*)}) & \&c[i][j][0] \\ \text{int (*) [4]} & c[i] & = (\text{int (*) [4]}) & \&c[i][0][0] \\ \text{int (*) [3][4]} & c & = (\text{int (*) [3][4]}) & \&c[0][0][0]\end{array}$$

Pointer
Types

Sizes of abstract data types

$$\begin{array}{lll}\text{int [4]} & c[i][j] & \text{size } = 4^*4 \\ \text{int [3][4]} & c[i] & \text{size } = 3^*4^*4 \\ \text{int [2][3][4]} & c & \text{size } = 2^*3^*4^*4\end{array}$$

Abstract Data
Types

Strides of array elements

$$\begin{array}{lll}\text{c[i][j][0]} & \text{stride } = 4^*4 \\ \text{c[i][0][0]} & \text{stride } = 3^*4^*4 \\ \text{c[0][0][0]} & \text{stride } = 2^*3^*4^*4\end{array}$$

$$\begin{array}{lll}k=[0:3] & \text{c[i][j][k]} & 4 \text{ integers} \\ j=[0:2], k=[0:3] & \text{c[i][j][k]} & 3*4 \text{ integers} \\ l=[0:1], j=[0:2], k=[0:3] & \text{c[i][j][k]} & 2*3*4 \text{ integers}\end{array}$$

General requirements for $c[i][j][k]$

```
int    c [2][3][4];
```

$$\begin{aligned}\&c[i][j][0] &= c[i][j] \\ \&c[i][0] &= c[i] \\ \&c[0] &= c\end{aligned}$$

with contiguous subarrays

$$\begin{aligned}\&c[i][j][k] &= c[i][j]+k \\ \&c[i][j] &= c[i]+j \\ \&c[i] &= c+i\end{aligned}$$



virtual assignments

$$\begin{aligned}c[i][j] &= \&c[i][j][0] \\ c[i] &= \&c[i][0][0] \\ c &= \&c[0][0][0]\end{aligned}$$

c[i][j]
c[i]
c

Pointer
Types

c[i][j]+1
c[i]+1
c+1

Pointer
Types

c[i][j]+k
c[i]+j
c+i

Pointer
Types

has an address of
has an address of
has an address of

c[i][j][0]
c[i][0][0]
c[0][0][0]

1 integer away
4 integers away
3*4 integers away

c[i][j][1]
c[i][1][0]
c[1][0][0]

1*k integers away
4*j integers away
3*4*i integers away

has an address of
has an address of
has an address of

General requirements for $c[i][j][k]$

```
int    c [2][3][4];
```

$$\begin{aligned}\&c[i][j][0] &= c[i][j] \\ \&c[i][0] &= c[i] \\ \&c[0] &= c\end{aligned}$$

with contiguous subarrays

$$\begin{aligned}\&c[i][j][k] &= c[i][j]+k \\ \&c[i][j] &= c[i]+j \\ \&c[i] &= c+i\end{aligned}$$



virtual assignments

$$\begin{aligned}c[i][j] &= \&c[i][j][0] \\ c[i] &= \&c[i][0][0] \\ c &= \&c[0][0][0]\end{aligned}$$

c[i][j]
c[i]
c

Abstract
Data

c[i][j]
c[i]
c

Abstract
Data

starts from
starts from
starts from

\&c[i][j][0]
\&c[i][0][0]
\&c[0][0][0]

c[i][j][k]
c[i][j][k]
c[i][j][k]

has a size of
has a size of
has a size of

4
3*4
2*3*4

integers
integers
integers

General requirements for $c[i][j][k]$

```
int    c [2][3][4];
```

$$\begin{aligned}\&c[i][j][0] &= c[i][j] \\ \&c[i][0] &= c[i] \\ \&c[0] &= c\end{aligned}$$

with contiguous subarrays

$$\begin{aligned}\&c[i][j][k] &= c[i][j]+k \\ \&c[i][j] &= c[i]+j \\ \&c[i] &= c+i\end{aligned}$$

for all k
for all j
for all i



virtual assignments

$$\begin{aligned}c[i][j] &= \&c[i][j][0] \\ c[i] &= \&c[i][0][0] \\ c &= \&c[0][0][0]\end{aligned}$$

$$\begin{aligned}c[i][j] &= \&c[i][j][0] \\ c[i] &= \&c[i][0][0] \\ c &= \&c[0][0][0]\end{aligned}$$

c[i]
c[i]
c

Abstract
Data

c[i]
c

Abstract
Data

c[i]
c

Abstract
Data

\&c[i][j][0]
\&c[i][0][0]
\&c[0][0][0]

start from
start from

\&c[i][0][0]
\&c[0][0][0]

starts from
starts from

\&c[i][0]
\&c[0]

Abstract
Data

General requirements for $c[i][j][k]$

```
int c [2][3][4];
```

$$\begin{aligned}\&c[i][j][0] &= c[i][j] \\ \&c[i][0] &= c[i] \\ \&c[0] &= c\end{aligned}$$

with contiguous subarrays

$$\begin{aligned}\&c[i][j][k] &= c[i][j]+k \\ \&c[i][j] &= c[i]+j \\ \&c[i] &= c+i\end{aligned}$$

for all k
for all j
for all i

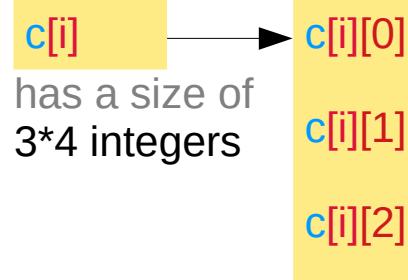


virtual assignments

$$\begin{aligned}c[i][j] &= \&c[i][j][0] \\ c[i] &= \&c[i][0][0] \\ c &= \&c[0][0][0]\end{aligned}$$

$c[i]$ and c start from $c[i][0]$ and $c[0]$ start from $\&c[i][0][0]$ and $\&c[0][0][0]$

Abstract Data



Abstract Data

has a size of 4 integers
has a size of 4 integers
has a size of 4 integers

$c[i]$ is the name of an array which has 3 elements
 $c[i][0]$ $c[i][1]$ $c[i][2]$

points to $c[i][0]$

$\&c[i][0]$ = $c[i]$

General requirements for $c[i][j][k]$

```
int    c [2][3][4];
```

$$\begin{aligned}\&c[i][j][0] &= c[i][j] \\ \&c[i][0] &= c[i] \\ \&c[0] &= c\end{aligned}$$

with contiguous subarrays

$$\begin{aligned}\&c[i][j][k] &= c[i][j]+k \\ \&c[i][j] &= c[i]+j \\ \&c[i] &= c+i\end{aligned}$$



virtual assignments

$$\begin{aligned}c[i][j] &= \&c[i][j][0] \\ c[i] &= \&c[i][0][0] \\ c &= \&c[0][0][0]\end{aligned}$$

Each
Each
Each

$c[i][j]$
 $c[i]$
 c

Abstract Data
Types

Each
Each
Each

$c[i][j]+1$
 $c[i]+1$
 $c+1$

Pointer
Types

Each
Each
Each

$c[i][j]+k$
 $c[i]+j$
 $c+i$

Pointer
Types

sub-array contains
sub-array contains
sub-array contains

4 integers
3*4 integers
2*3*4 integers

starts at
starts at
starts at

$c[i][j][1]$
 $c[i][1][0]$
 $c[1][0][0]$

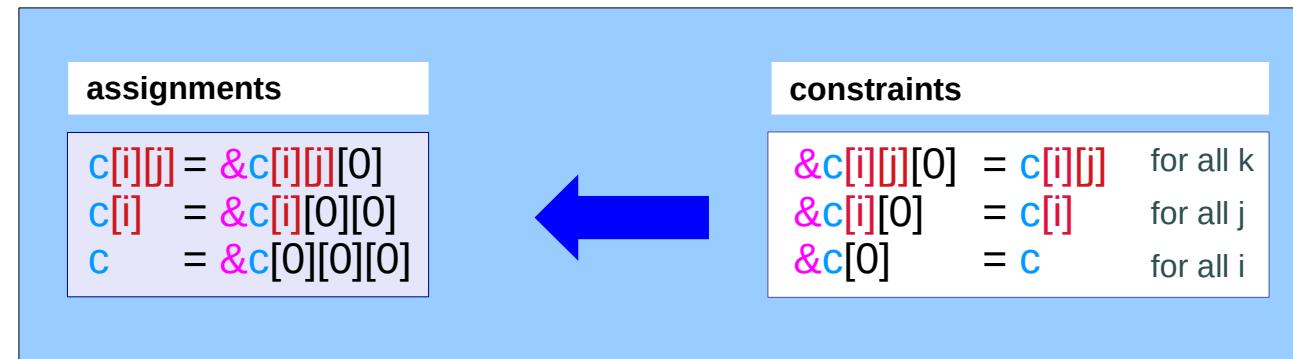
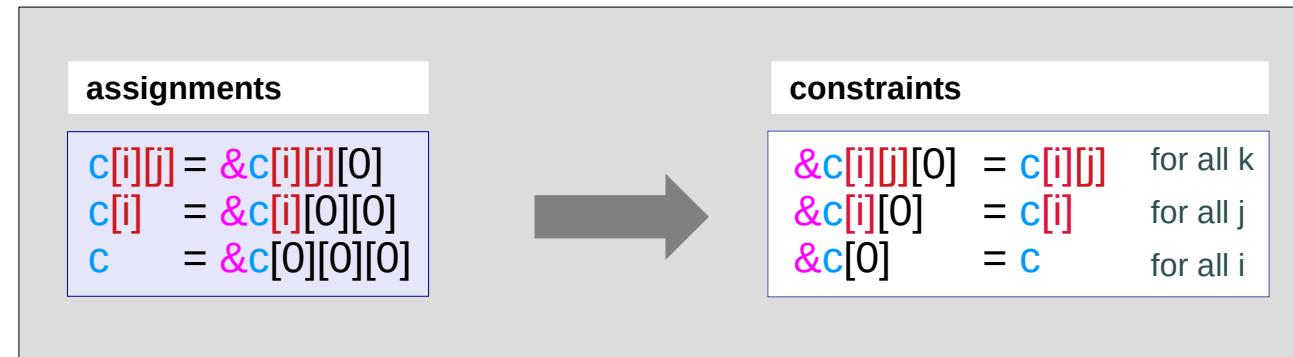
1 integer away
4 integers away
3*4 integers away

starts at
starts at
starts at

$c[i][j][k]$
 $c[i][j][0]$
 $c[i][0][0]$

$1*k$ integers away
 $4*j$ integers away
 $3*4*i$ integers away

multi-dimensional arrays



Subarray starting addresses

```
int    c [2][3][4];
```

Virtual array pointer

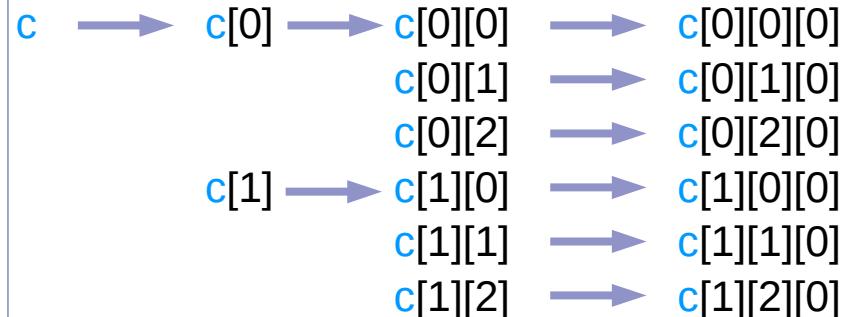
$$\begin{aligned}\&c[i][j][0] &= c[i][j] \\ \&c[i][0] &= c[i] \\ \&c[0] &= c\end{aligned}$$

$c[i][j]$ points to $c[i][j][0]$
 $c[i]$ points to $c[i][0]$
 c points to $c[0]$

Abstract data (array)

$$\begin{aligned}\&c[i][j][0] &= \&c[i][j] \\ \&c[i][0] &= \&c[i] \\ \&c[0] &= \&c\end{aligned}$$

$c[i][j]$ starts from $\&c[i][j][0]$
 $c[i]$ starts from $\&c[i][0]$
 c starts from $\&c[0]$



$$c[i][j] = \&c[i][j][0]$$

$$\begin{aligned}c[0][0] &= \&c[0][0][0] \\ c[0][1] &= \&c[0][1][0] \\ c[0][2] &= \&c[0][2][0] \\ c[1][0] &= \&c[1][0][0] \\ c[1][1] &= \&c[1][1][0] \\ c[1][2] &= \&c[1][2][0]\end{aligned}$$

$$c[i] = \&c[i][0]$$

$$\begin{aligned}c[0] &= \&c[0][0] = \&c[0][0][0] \\ c[1] &= \&c[1][0] = \&c[1][0][0]\end{aligned}$$

$$c = \&c[0]$$

$$c = \&c[0] = \&c[0][0] = \&c[0][0][0]$$

General requirements for $c[i][j][k]$

```
int    c [2][3][4];
```

$$\begin{aligned}\&c[i][j][0] &= c[i][j] \\ \&c[i][0] &= c[i] \\ \&c[0] &= c\end{aligned}$$

with contiguous subarrays

$$\begin{aligned}\&c[i][j][k] &= c[i][j]+k \\ \&c[i][j] &= c[i]+j \\ \&c[i] &= c+i\end{aligned}$$



virtual assignments

$$\begin{aligned}c[i][j] &= \&c[i][j][0] \\ c[i] &= \&c[i][0][0] \\ c &= \&c[0][0][0]\end{aligned}$$

Contiguity Constraints

c [i][j][k];

Virtual Array Pointers and Contiguity

Using array pointers

`int (*) [N]` `int (*) [M][N]` `int (*) [L][M][N]`



`int [N]`

`int (*)`

`int [M][N]`

`int (*) [N]`

`int (*) [L][M][N]`



`int [L][M][N]`

`int (*)[M][N]`

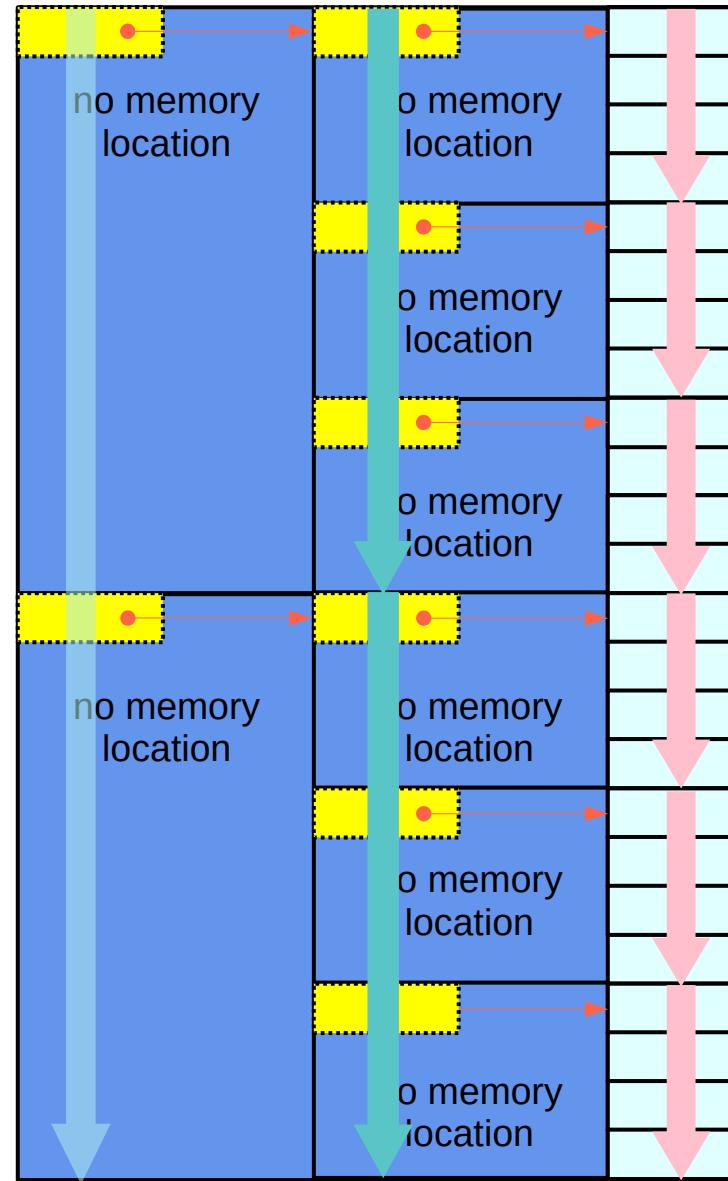
Array Pointer

Array Name

Array Pointer

Array pointer approach – contiguity constraints

Array Pointer Approach
(pointer to arrays)



Equivalence and contiguity (1)

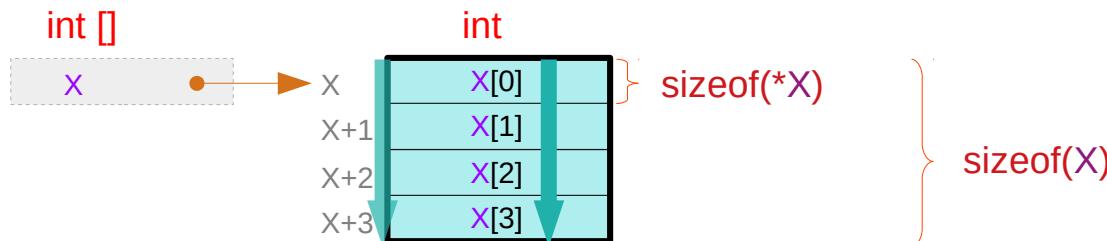
consecutive address

$*(\text{X} + n)$

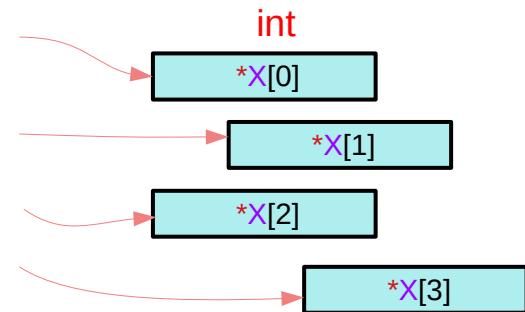
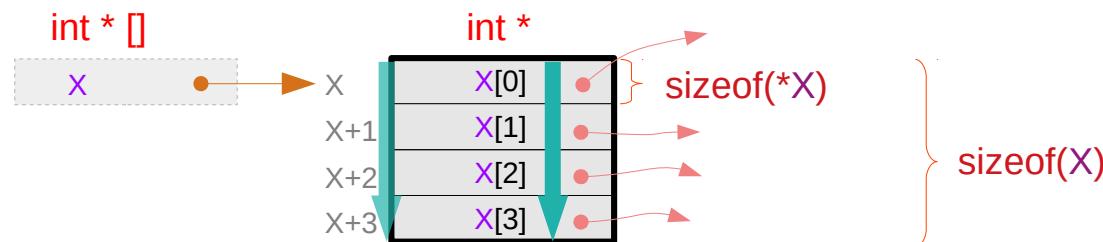
consecutive data

$\equiv \text{X}[n]$

contiguous index : n



int X[4] ; contiguous X[i] for a given X : primitive types



$\text{int } * \text{X[4]}$; contiguous X[i] for a given X : pointer types

Equivalence and contiguity (2)

consecutive address

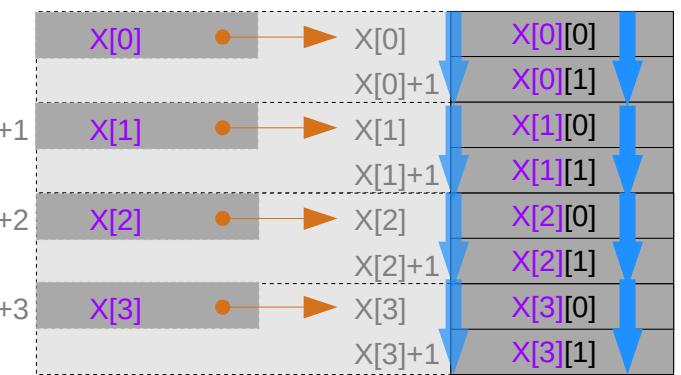
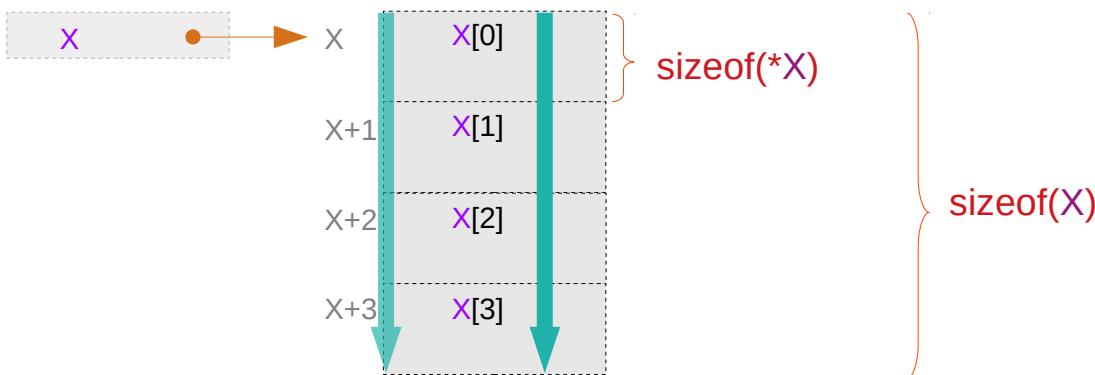
$*(\text{X}+n)$

consecutive data

$\equiv \text{X}[n]$

contiguous index : n

can be recursively applied



atype * $\text{X}[4]$; contiguous $\text{X}[i]$ for a given X : abstract data types

Recursive applications of equivalences

By definition, contiguous memory locations are assumed

consecutive address

consecutive data

$$*(\textcolor{red}{X} + \textcolor{blue}{n}) \equiv \textcolor{red}{X}[n]$$

contiguous index : n

$$*(\textcolor{red}{p[m]} + \textcolor{blue}{n}) \leftrightarrow \textcolor{red}{p[m]}[\textcolor{blue}{n}]$$

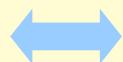
$\textcolor{red}{X} = \textcolor{red}{p[m]}$ contiguous index : n

$$(*(\textcolor{red}{p} + \textcolor{blue}{m}))[n]; \leftrightarrow \textcolor{red}{p}[m][n];$$

$\textcolor{red}{X} = \textcolor{red}{p}$ contiguous index : m

Equivalence for a given $p[m]$ (1)

$*(p[m]+n)$



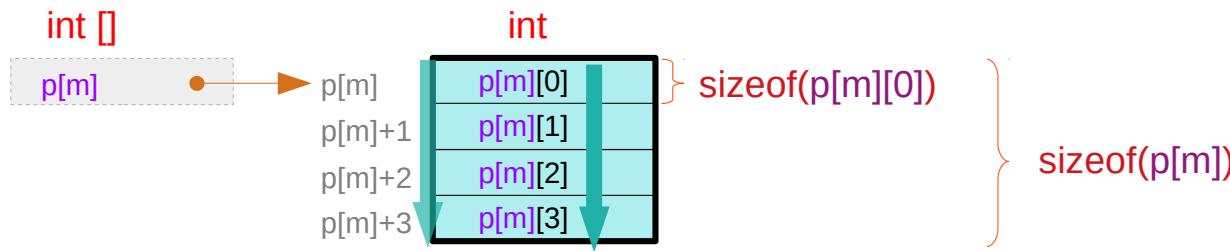
$p[m][n]$

for a given $p[m]$

contiguous index : n

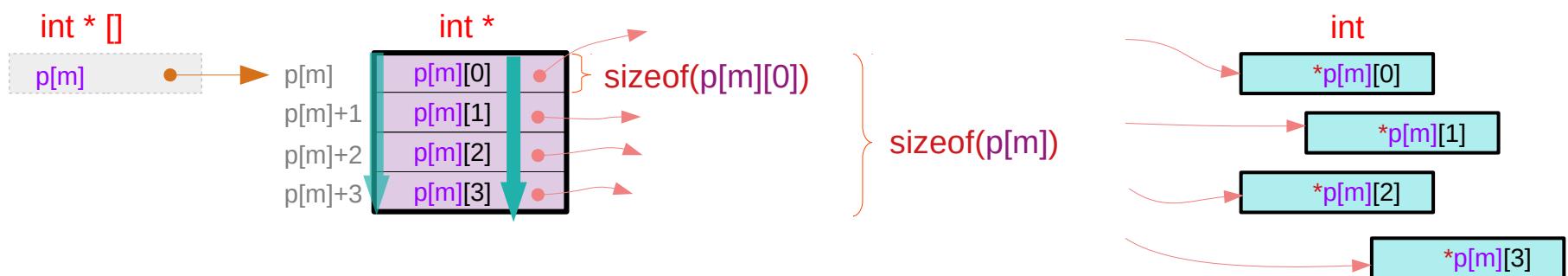
int p[M][4]; contiguous $p[m][n]$ for a given $p[m]$: primitive types

$m = 0, 1, \dots, M-1$



int * p[M][4]; contiguous $p[m][n]$ for a given $p[m]$: pointer types

$m = 0, 1, \dots, M-1$

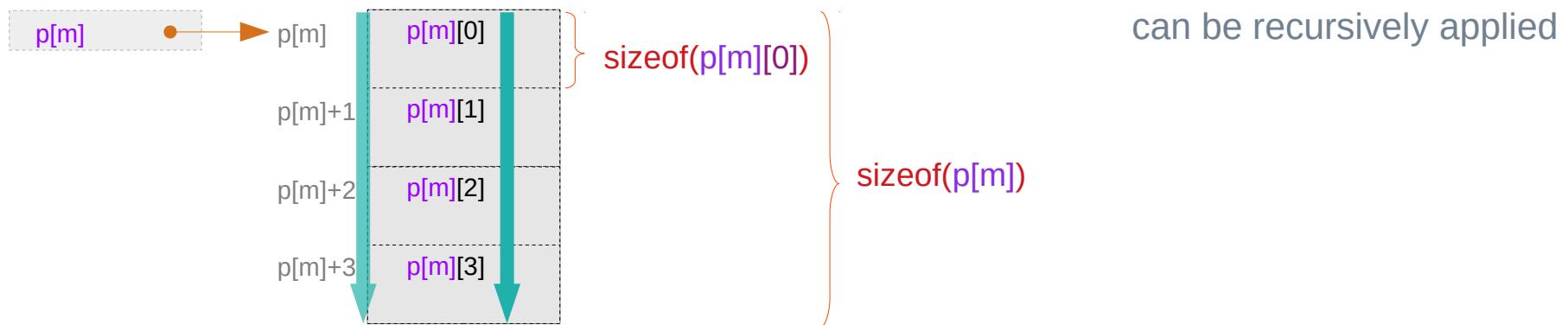


Equivalence for a given $p[m]$ (2)

$$*(p[m]+n) \leftrightarrow p[m][n]$$

for a given $p[m]$ contiguous index : n

atype * p[M][4]; contiguous $p[m][n]$ for a given $p[m]$: abstract data types $m = 0, 1, \dots, M-1$

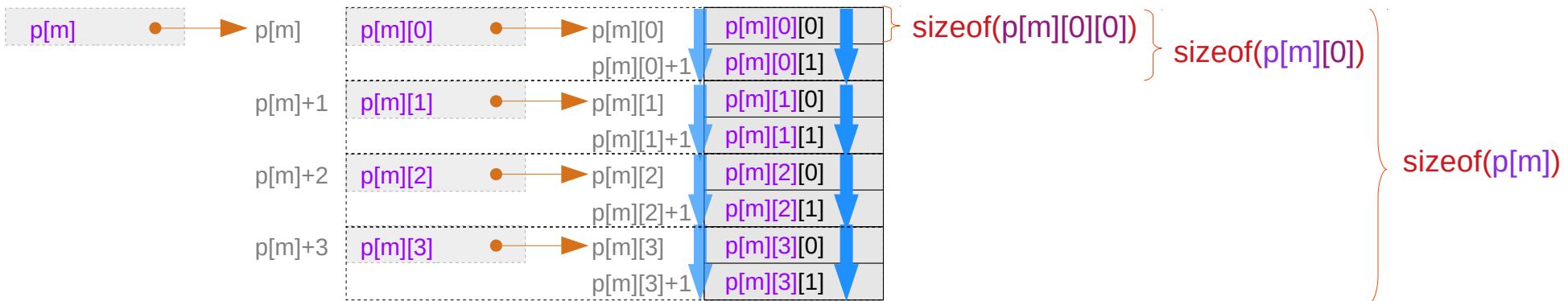


Equivalence for a given $p[m][n]$

$$*(p[m][n]+k) \leftrightarrow p[m][n][k]$$

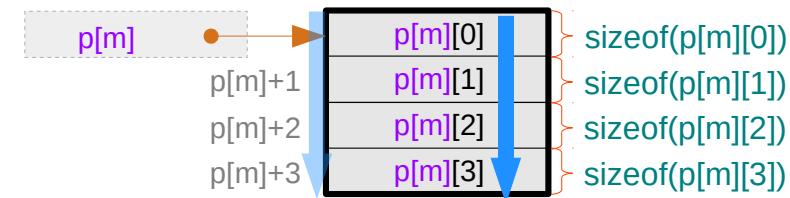
for a given $p[m][n]$ contiguous index : k

atype * p[M][4][2]; contiguous $p[m][n][k]$ for a given $p[m][n]$: abstract data types $m = 0, 1, \dots, M-1$



Contiguity constraints in multi-dimensional arrays

$$*(p[m] + n) \leftrightarrow p[m][n]$$



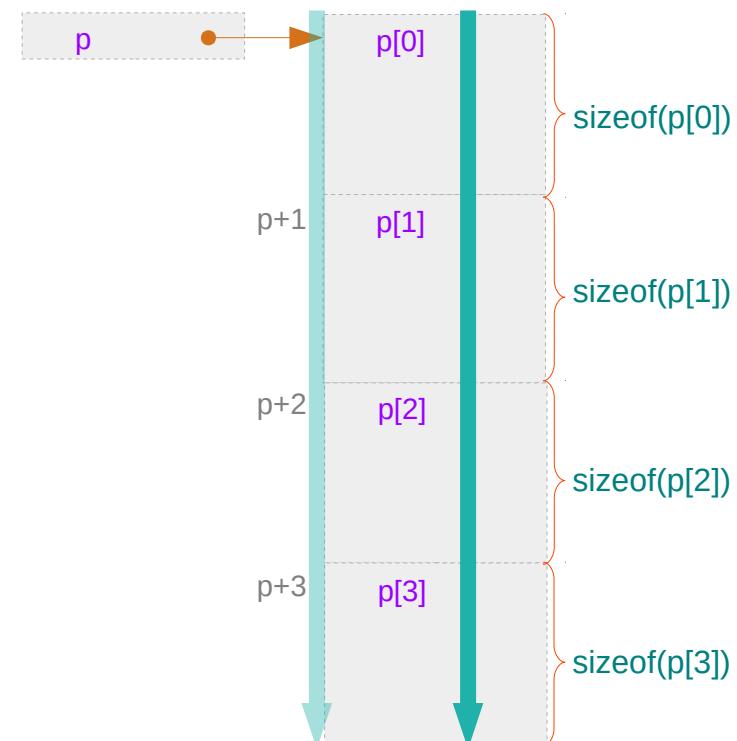
for a given $p[m]$, thus for a given p and m ,
 $p[m][n]$'s must be contiguous for all n .
 $p[m][0], p[m][1], \dots, p[m][N-1]$

contiguous index : n

$$*(p+m) \leftrightarrow p[m]$$

for a given p ,
 $p[m]$'s must be contiguous for all m .
 $p[0], p[1], \dots, p[M-1]$

contiguous index : m



Contiguity constraints for p

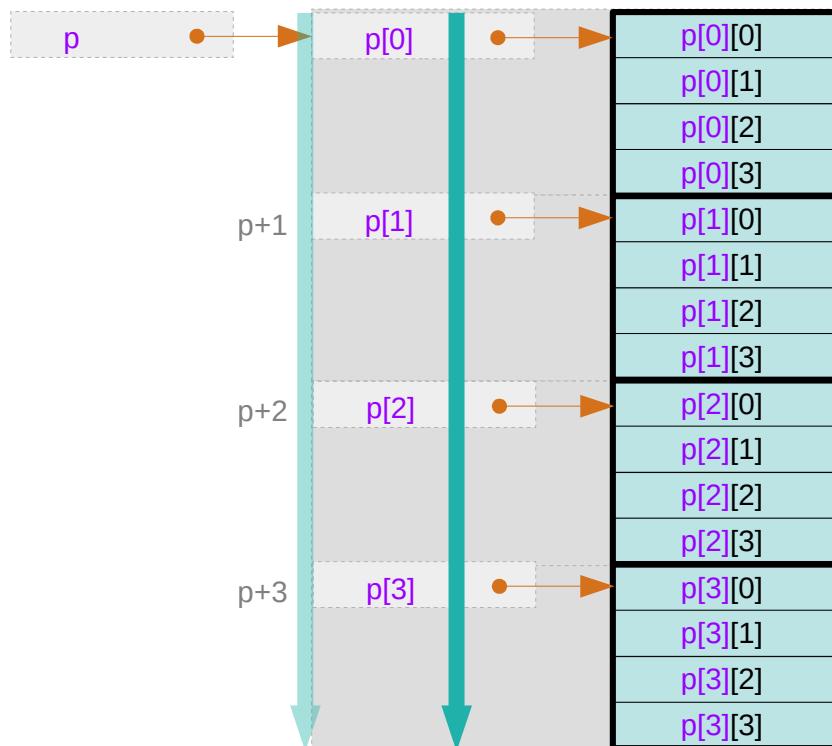
$$*(p+m) \leftrightarrow p[m]$$

for a given p

contiguous index : m

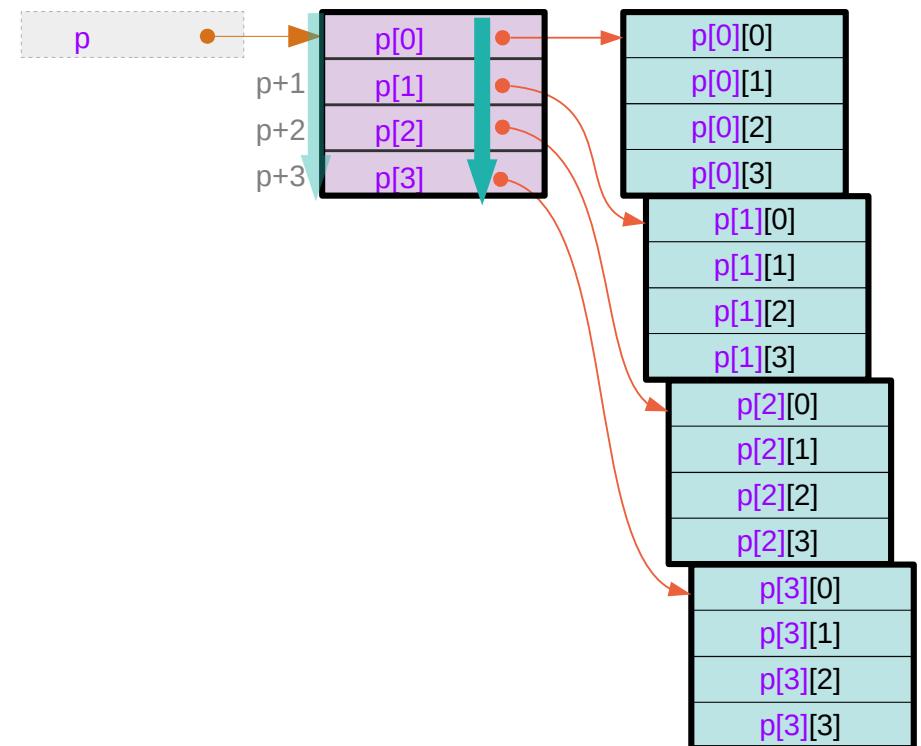
2-d array name

1-d array names



contiguous $p[m]$ → contiguous $p[m][n]$

1-d array of pointers



contiguous $p[m]$ → contiguous $p[m][n]$

Not necessarily

Contiguity constraints for $p[m]$ – using array pointers

$*(p[m]+n)$

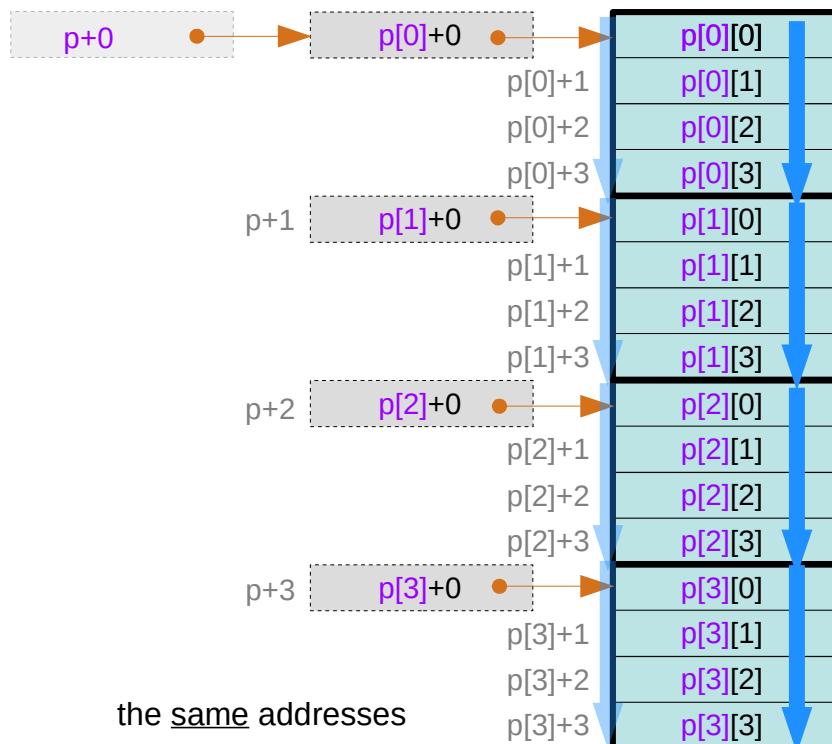


$p[m][n]$

for a given $p[m]$ contiguous index : n

2-d array name

1-d array names



contiguous $p[m]$  contiguous $p[m][n]$

$p[0][0] = *(p[0]+0)$

addr

$\&p[0][0] = p[0]$

addr

$p+0$



$p[1][0] = *(p[1]+0)$

addr

$\&p[1][0] = p[1]$

addr

$p+1$



$p[2][0] = *(p[2]+0)$

addr

$\&p[2][0] = p[2]$

addr

$p+2$



$p[3][0] = *(p[3]+0)$

addr

$\&p[3][0] = p[3]$

addr

$p+3$

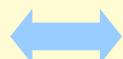


the same addresses

virtual array pointer  no real memory locations

Contiguity constraints for $p[m]$ – using pointer arrays

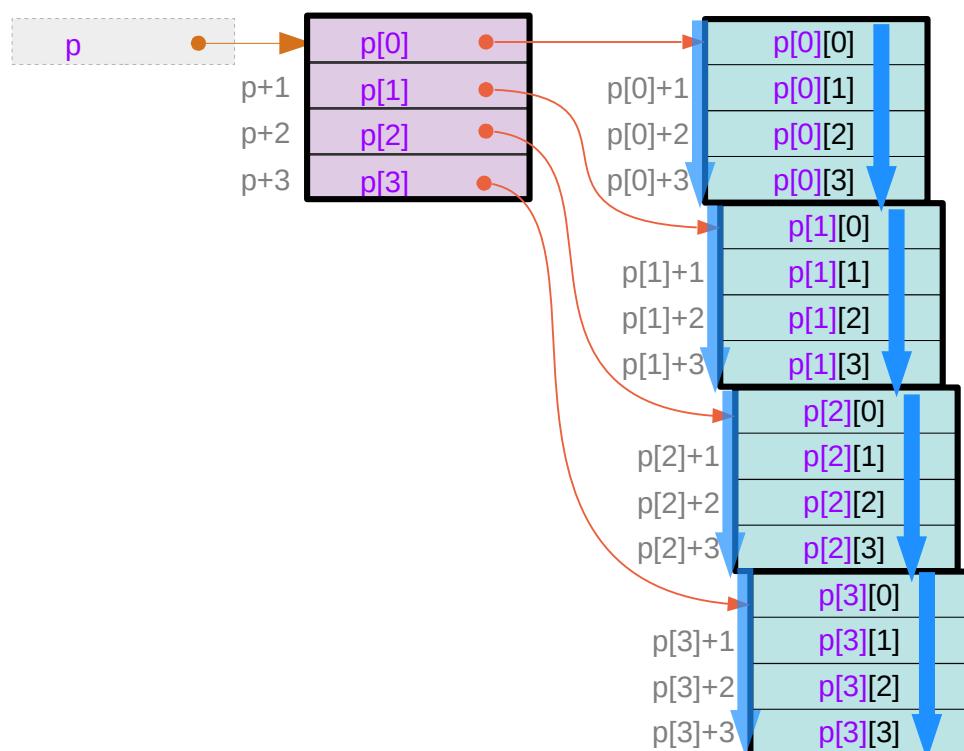
$*(p[m]+n)$



$p[m][n]$

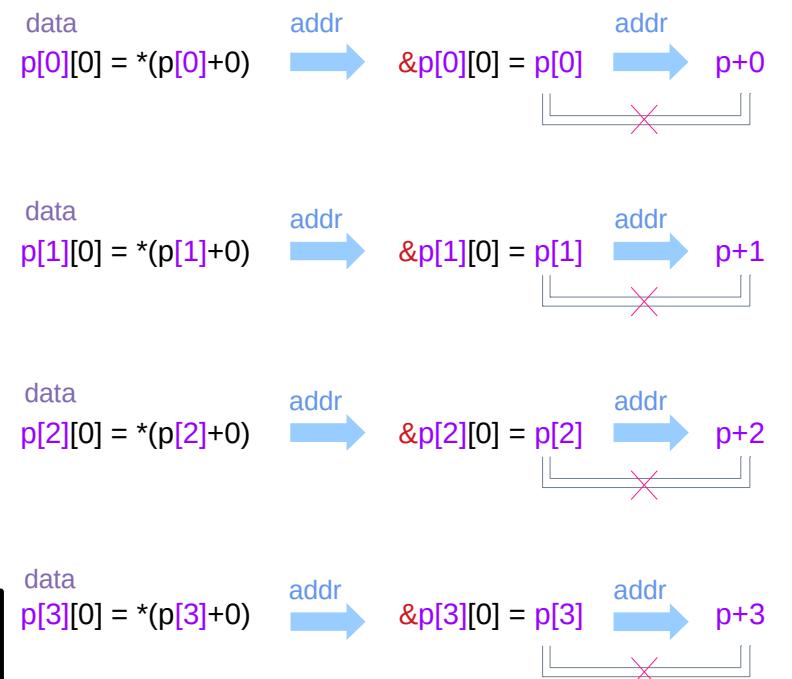
for a given $p[m]$ contiguous index : n

1-d array of pointers



contiguous $p[m] \rightarrow$ contiguous $p[m][n]$

Not necessarily



the different addresses

Contiguity constraints for 2-d arrays

`int a[M][N] ;`

$$*(a+m) \leftrightarrow a[m]$$

$a[0], a[1], \dots, a[M-1]$
are contiguous

$$*(a[m]+n) \leftrightarrow a[m][n]$$

$a[m][0], a[m][1], \dots, a[m][N-1]$
are contiguous

`int (*b)[N] ;`

$$*(b+m) \leftrightarrow b[m]$$

$b[0], b[1], \dots, b[M-1]$
are contiguous

$$*(b[m]+n) \leftrightarrow b[m][n]$$

$b[m][0], b[m][1], \dots, b[m][N-1]$
are contiguous

`int * c[M] ;`

$$*(c+m) \leftrightarrow c[m]$$

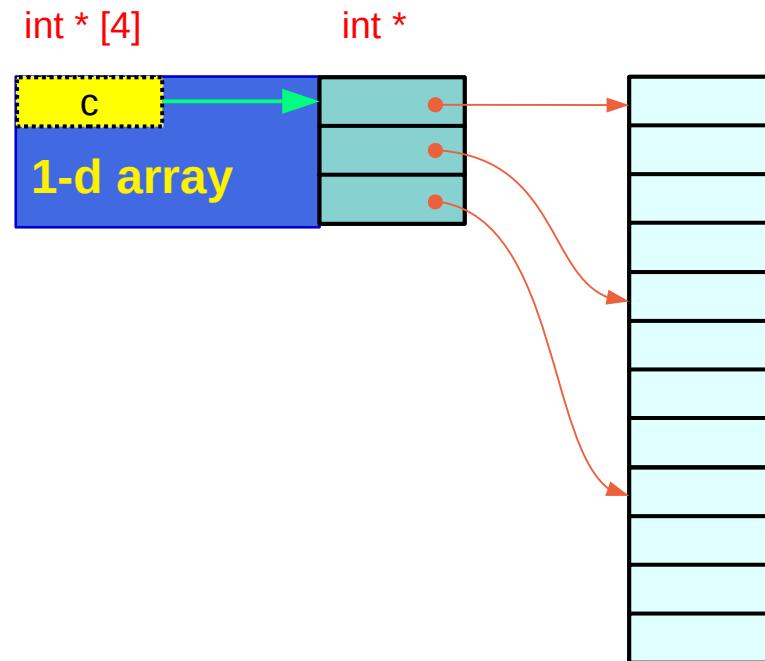
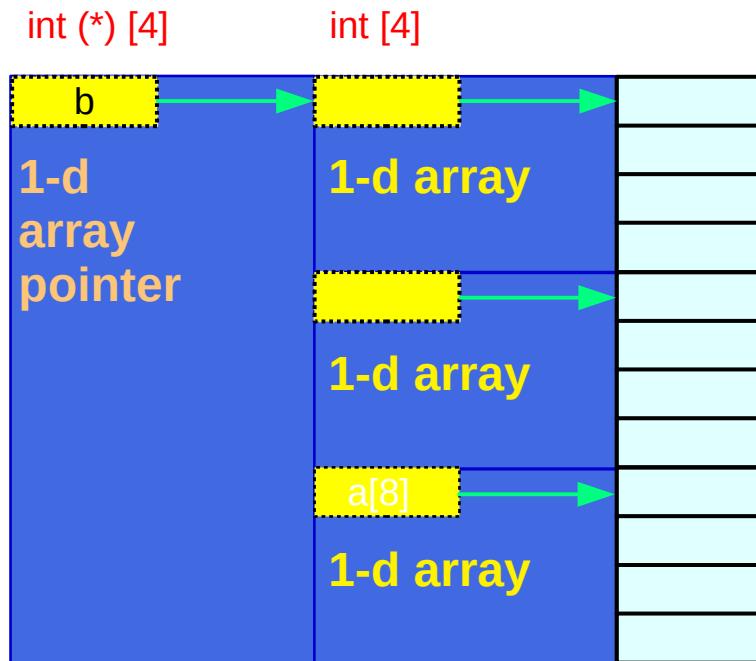
$c[0], c[1], \dots, c[M-1]$
are contiguous

$$*(c[m]+n) \leftrightarrow c[m][n]$$

$c[m][0], c[m][1], \dots, c[m][N-1]$
are contiguous

a set of assignments of pointers
are necessary for this contiguity

Pointer Arrays vs Array Pointers



`int (*b)[N] ;`

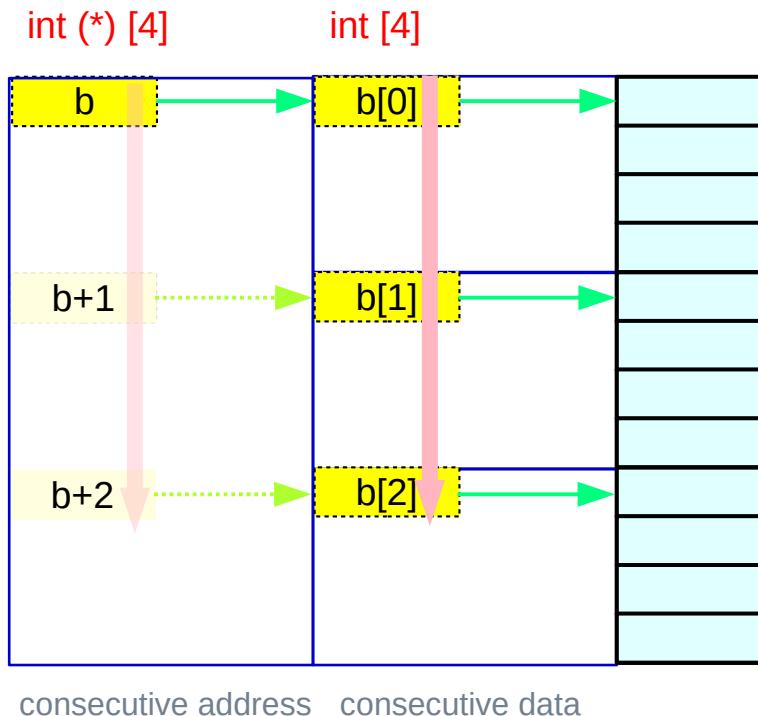
$*(b+m)$ \leftrightarrow $b[m]$
 $*(b[m]+n)$ \leftrightarrow $b[m][n]$

`int * c[M] ;`

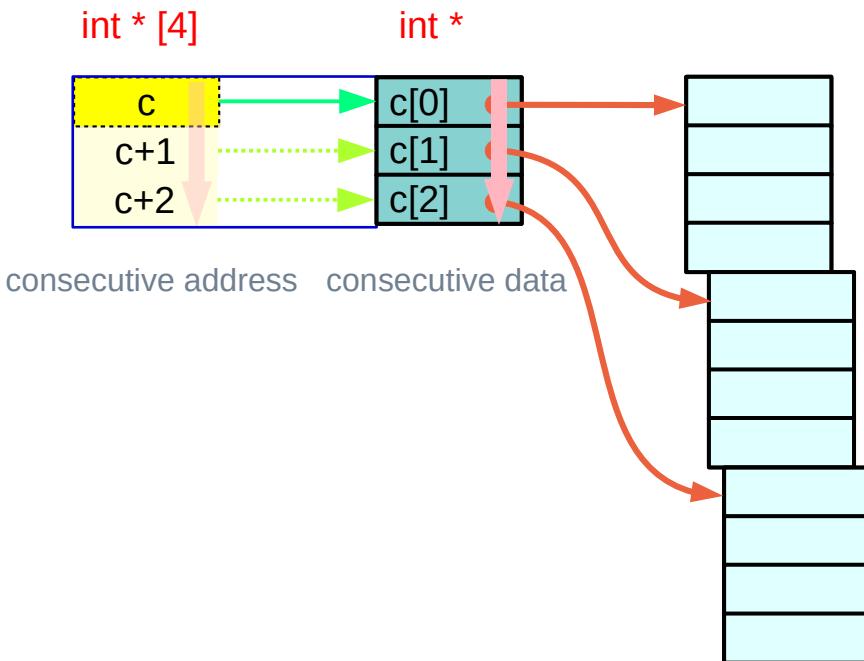
with proper assignments

$*(c+m)$ \leftrightarrow $c[m]$ or
 $*(c[m]+n)$ \leftrightarrow $c[m][n]$

Pointer Arrays vs Array Pointers



consecutive address consecutive data



consecutive address consecutive data

`int (*b)[N] ;`

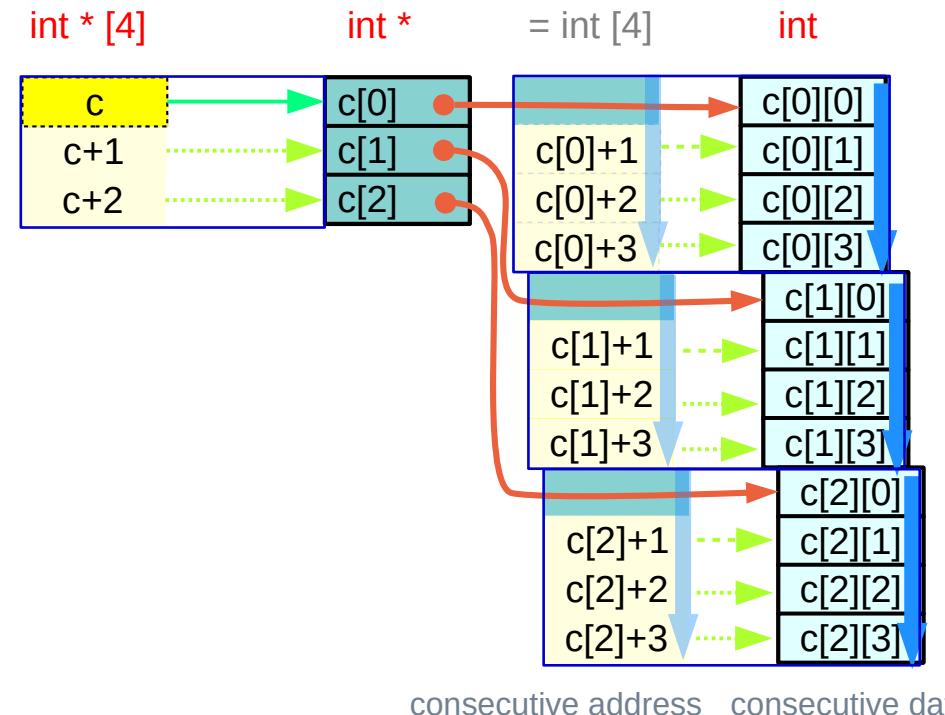
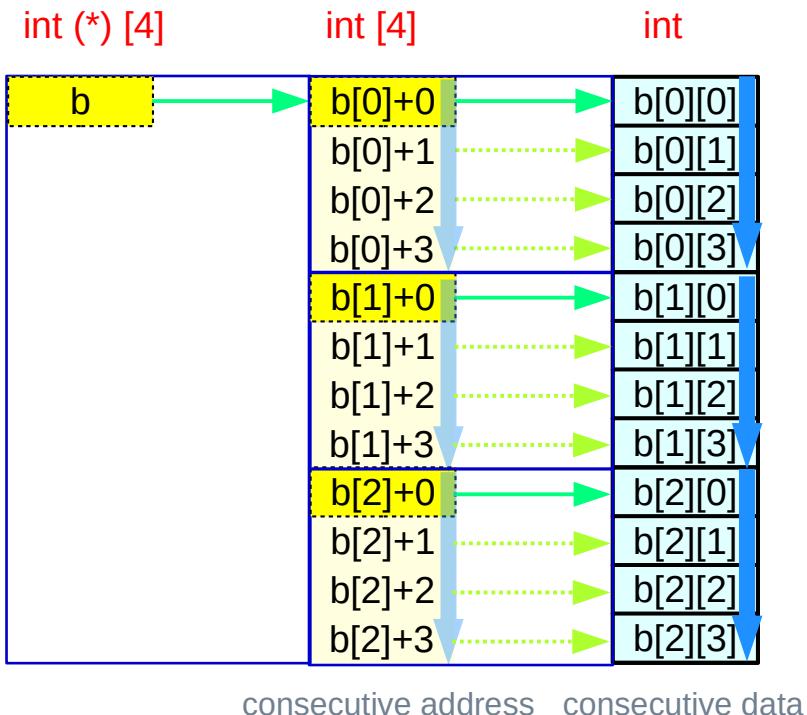
$\ast(b+m)$ \longleftrightarrow $b[m]$
 $\ast(b[m]+n)$ \longleftrightarrow $b[m][n]$

`int * c[M] ;`

with proper assignments

$\ast(c+m)$ \longleftrightarrow $c[m]$ or
 $\ast(c[m]+n)$ \longleftrightarrow $c[m][n]$

Pointer Arrays vs Array Pointers



`int (*b)[N] ;`

$$\begin{array}{ccc} *(b+m) & \leftrightarrow & b[m] \\ *(b[m]+n) & \leftrightarrow & b[m][n] \end{array}$$

`int * c[M] ;`

$$\begin{array}{ccc} *(c+m) & \leftrightarrow & c[m] \\ *(c[m]+n) & \leftrightarrow & c[m][n] \end{array}$$

Three contiguity constraints for 3-d arrays

Pointer Array Approach (array of pointers)

$c[i][j][k]$	\rightarrow	$*(c[i][j] + k)$
$*(c[i][j] + k)$	\rightarrow	$*(*(c[i] + j) + k)$
$*(*(c[i] + j) + k)$	\rightarrow	$*(*(*(c + i) + j) + k)$

contiguous int	int
contiguous pointers to int	int *
contiguous double pointers to int	int **

the contiguity constraints are satisfied by allocating arrays of pointers

Array Pointer Approach (pointer to arrays)

$c[i][j][k]$	\rightarrow	$*(c[i][j] + k)$
$*(c[i][j] + k)$	\rightarrow	$*(*(c[i] + j) + k)$
$*(*(c[i] + j) + k)$	\rightarrow	$*(*(*(c + i) + j) + k)$

contiguous 0-d arrays	int	int
contiguous 1-d arrays	int [4]	int *
contiguous 2-d arrays	int [3][4]	int (*) [4]

The contiguity constraints are satisfied by row major ordered linear data layout

Contiguous array pointers $c[i][j][k] \equiv *(*c[i][j] + k)$

$c[0][0][0] = *(c[0][0] + 0)$
$c[0][0][1] = *(c[0][0] + 1)$
$c[0][0][2] = *(c[0][0] + 2)$
$c[0][0][3] = *(c[0][0] + 3)$
$c[0][1][0] = *(c[0][1] + 0)$
$c[0][1][1] = *(c[0][1] + 1)$
$c[0][1][2] = *(c[0][1] + 2)$
$c[0][1][3] = *(c[0][1] + 3)$
⋮
⋮

contiguous 1-d array elements

$c[i][j]$

$\text{int } [4]$

$\text{int } *$

int

$\text{sizeof}(c[i][j])$
 $[k]$
 $\text{sizeof}(c[i][j][k]) * 4$
 $\text{sizeof}(\text{int}) * 4$

int $c[2][3][4];$

Address Value

$c[i][j] + k$

$\&c[i][j][0] + k * \text{sizeof}(*c[i][j])$

$\&c[i][j][0] + k * \text{sizeof}(c[i][j][0])$

$\&c[i][j][0] + k * 4$

pointer $\text{int } (*)$

$c[0][0]$

$c[0][0][0]$

$c[0][0][1]$

$c[0][0][2]$

$c[0][0][3]$

$c[0][2]$

$c[0][2][0]$

$c[0][2][1]$

$c[0][2][2]$

$c[0][2][3]$

$c[1][1]$

$c[1][1][0]$

$c[1][1][1]$

$c[1][1][2]$

$c[1][1][3]$

$c[1][2]$

$c[1][2][0]$

$c[1][2][1]$

$c[1][2][2]$

$c[1][2][3]$

int primitive data

$c[0][0]$

$c[0][0][0]$

$c[0][0][1]$

$c[0][0][2]$

$c[0][0][3]$

$c[0][1]$

$c[0][1][0]$

$c[0][1][1]$

$c[0][1][2]$

$c[0][1][3]$

$c[1][0]$

$c[1][0][0]$

$c[1][0][1]$

$c[1][0][2]$

$c[1][0][3]$

$c[1][1]$

$c[1][1][0]$

$c[1][1][1]$

$c[1][1][2]$

$c[1][1][3]$

$c[1][2]$

$c[1][2][0]$

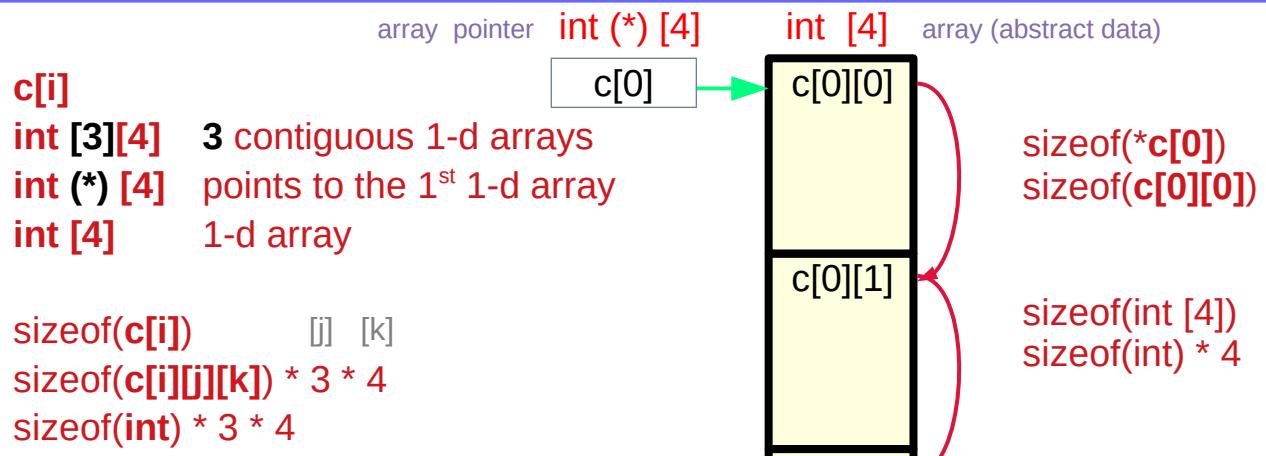
$c[1][2][1]$

$c[1][2][2]$

$c[1][2][3]$

Contiguous array pointers $c[i][j] \equiv *(*c[i] + j)$

```
c[0][0] = *(c[0] + 0)
c[0][1] = *(c[0] + 1)
c[0][2] = *(c[0] + 2)
c[1][0] = *(c[1] + 0)
c[1][1] = *(c[2] + 1)
c[1][2] = *(c[3] + 2)
```



```
int c[2][3][4];
```

Address Value

$c[i] + j$

$\&c[i][0][0] + j * \text{sizeof}(*c[i])$

$\&c[i][0][0] + j * \text{sizeof}(c[i][0])$

$\&c[i][0][0] + j * 4 * 4$

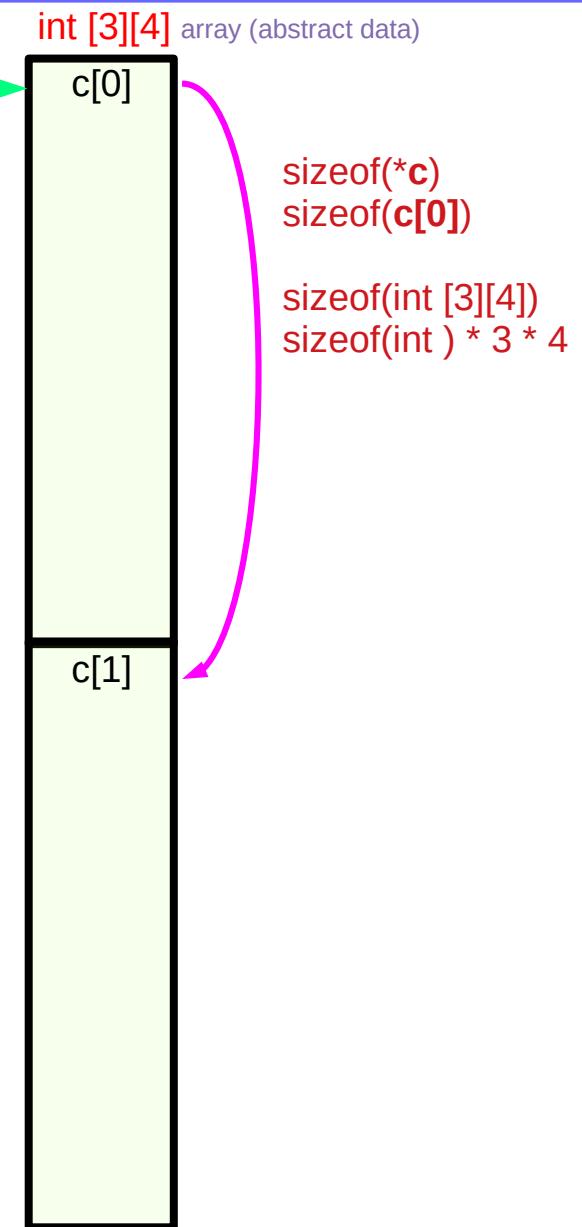
Contiguous array pointers $c[i] \equiv *(c + i)$

$c[0] = *(c + 0)$
$c[1] = *(c + 1)$

array pointer $\text{int } (*) [3][4]$
c
 $\text{int } [2][3][4]$ 2 contiguous 2-d arrays
 $\text{Int } (*) [3][4]$ points to the 1st 2-d array
 $\text{int } [3][4]$ 2-d array

$\text{sizeof}(c)$ $[i] \quad [j] \quad [k]$
 $\text{sizeof}(c[i][j][k]) * 2 * 3 * 4$
 $\text{sizeof}(\text{int}) * 2 * 3 * 4$

int $c[2][3][4];$



Address Value

$c + i$
 $\&c[0][0][0] + i * \text{sizeof}(*c)$
 $\&c[0][0][0] + i * \text{sizeof}(c[0])$
 $\&c[0][0][0] + i * 4 * 3 * 4$

Contiguous linear layout

```
int c [L][M][N];
```

```
c [i][j][k];
```

L	M	N
i	j	k
i^*M^*N	j^*N	k

Base Index = 0

Offset Index 1 (i=1)

i^*M^*N

Offset Index 2 (j=1)

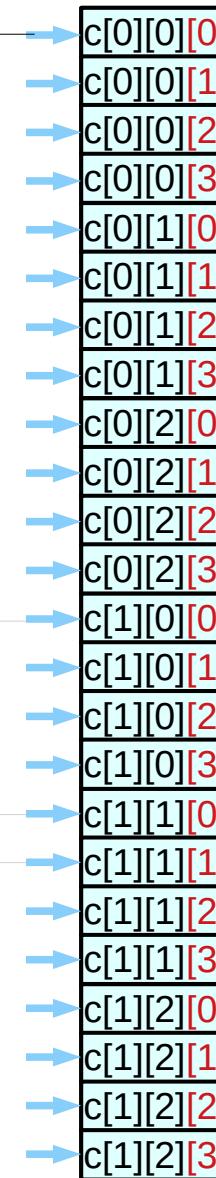
j^*N

Offset Index 3 (k=1)

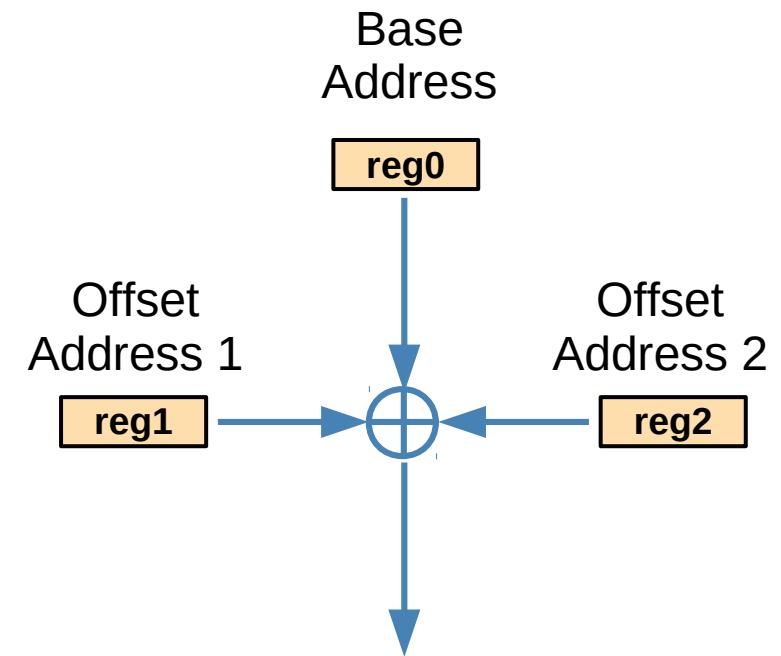
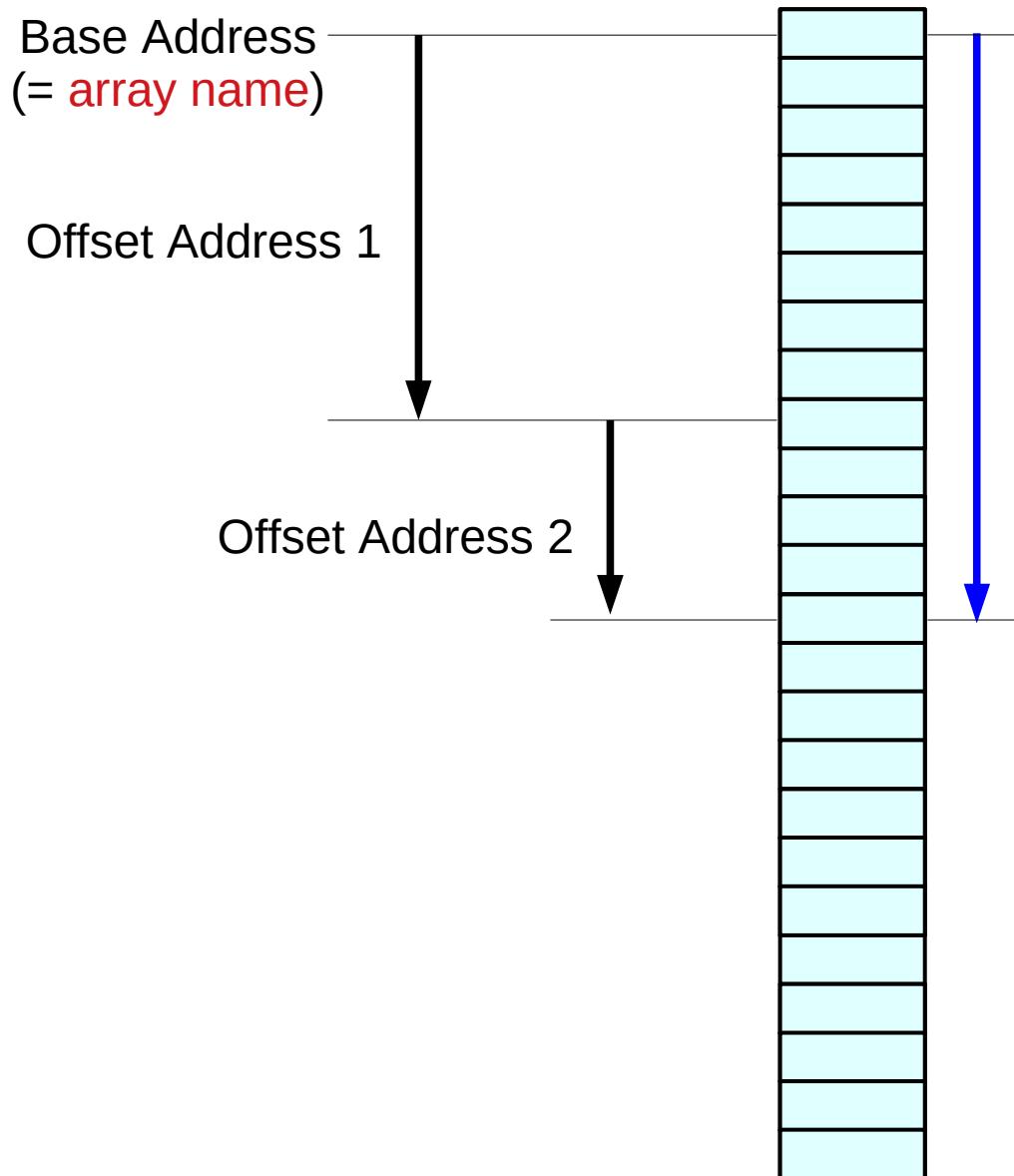
k

$$(i^*M^*N + j^*N + k)$$

$$((i^*M + j)^*N + k)$$

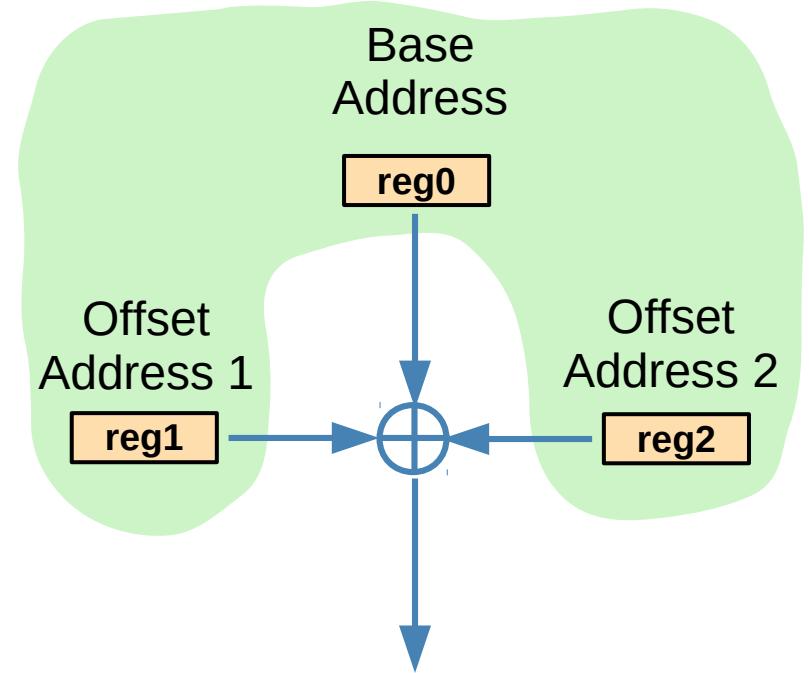
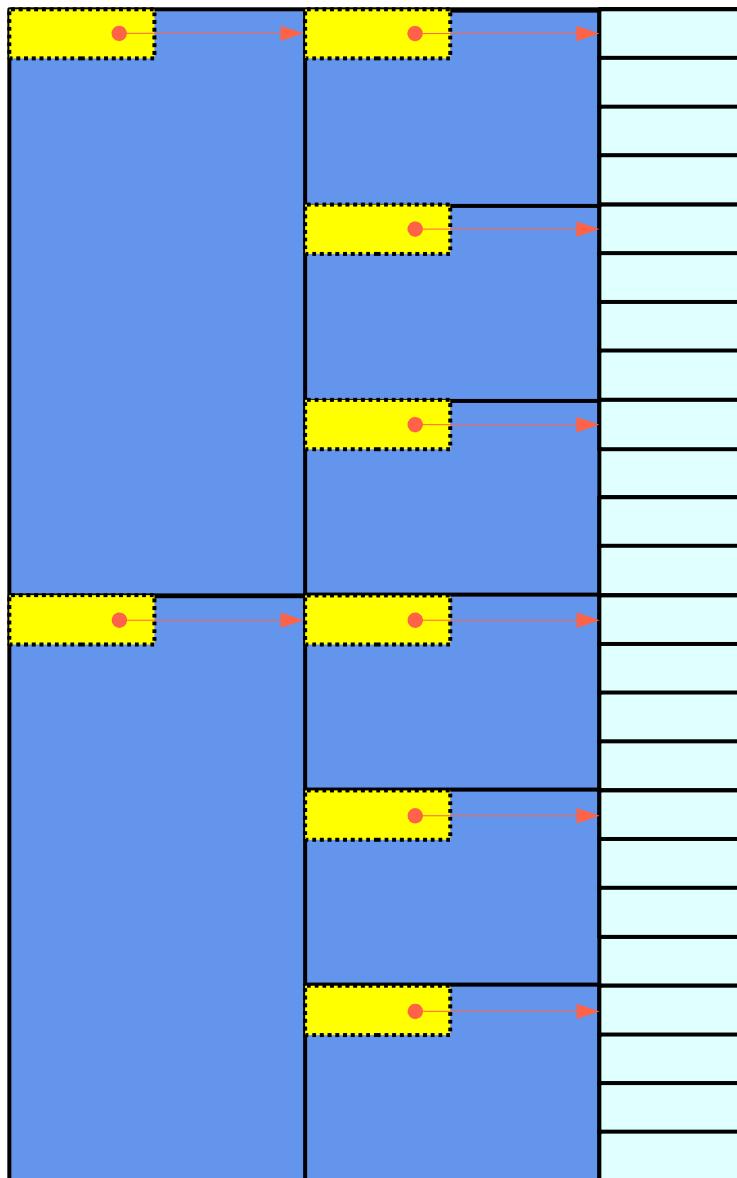


Base and Offset Addressing



compiler
assembly instruction
registers in the CPU

Array Pointer Approach



Array Pointer Approach
(pointer to arrays)

References

- [1] Essential C, Nick Parlante
- [2] Efficient C Programming, Mark A. Weiss
- [3] C A Reference Manual, Samuel P. Harbison & Guy L. Steele Jr.
- [4] C Language Express, I. K. Chun