

GAS Tutorial - 4. Sections & Relocation

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1 Sections and Relocation

“Using as”, Dean Elsner, Jay Fenlason & friends

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- a section is a range of addresses, with no gaps
- the linker ld
 - reads many object files (initial address 0)
 - combines them to an executable
 - moves blocks of bytes (i.e, sections) of your program to their run-time addresses
- relocation: assigning run-time addresses to sections

- an object file written by `as` has at least three sections
 - text section
 - data section
 - bss section
- these sections can be empty
- in an object file: the text section starts at address 0, the data section follows, and

finally the bss section.

- COFF or ELF output
 - named section (`.section directives`)

- when the sections are relocated, the ld should know
 - which data changes
 - how to change that data
- whenever an address in the object file is referenced,
 - the beginning of this reference to an address?
 - the length (in bytes) of this reference?
 - which section does the address refer to?
 - (address) (start-address of section)?
 - “Program-Counter relative”?

every address is expressed as

- (section) + (offset into section)
- every expression has this section-relative nature
- {secname N } notation: “offset N into section sec-name.”

Absolute Section

- addresses in the absolute section remain unchanged
- address {absolute 0} is “relocated” to run-time address 0 by ld
- generally, linker never use overlapping addresses
- address in absolute sections must overlap

{undefined U }

- any address whose section is unknown at assembly time
- U is to be filled
- to generate an undefined address
 - using an undefined symbol.
 - using a named common block

ld puts

- all partial programs' text sections
- in contiguous addresses in the linked program.
- all partial programs' data sections
- in contiguous addresses in the linked program.
- all partial programs' bss sections
- in contiguous addresses in the linked program.

Linker's view of section types

- 1 text section, data section, named section
- 2 bss section
- 3 absolute section
- 4 undefind section

Linker Section - text/data/named sections

- these sections hold your program
- as and ld treat them as separate but equal sections
- these sections are differentiated when the program is running
 - the text section : unalterable
 - often shared among processes
 - contains instructions, constants
 - the data section : alterable:
 - C variables

- contains zeroed bytes when your program begins running
- used to hold uninitialized variables or common storage
- the length of each partial program's bss section is important
- there is no need to store explicit zero bytes in the object file
- because the program starts out containing zeroed bytes
- bss section was invented to eliminate those explicit zeros from object files.

Linker Section - absolute sections

- Address 0 of this section is always “relocated” to runtime address 0
- useful when referring to an address that ld must not change
- being “unrelocatable”: addresses do not change during relocation.

- This “section” is a catch-all for address references to objects not in the preceding sections.

used to locate separate groups of data in named sections close to each other in the object file

- a section can be divided into numbered subsections
- subsection number ranging from 0 to 8192
- default subsection number 0
- bytes with the same subsection number are assembled together

Subsection Usage

- `.text` expression
- `.data` expression
- `.section name , expression [COFF]`
- `.subsection expression [ELF]`

- expression represents subsection number
- expression should use absolute address

- `.text` (`.text 0`, equivalently)
- `.data` (`.data 0`, equivalently)

Subsection Examples (1)

For example, to store constants in the text section not interspersed

- `'.text 0'` before each section of code being output
- `'.text 1'` before each group of constants being output.

Subsection Examples (2)

- .text 0
- .ascii "1st text subsection - (1)"
- .text 1
- .ascii "2nd text subsection."
- .data 0
- .ascii "1st data subsection,"
- .text 0
- .ascii "1st text subsection - (2)"

Location Counter

- each section has a location counter
- incremented by one for every byte assembled into that section
- subsection do not have its own location counter
- by using `.align`, a location counter is manipulated indirectly
- by using `label`, the value of a location counter can be captured
- the location counter of a section which are being assembled is said to be active

- used for local common variable storage
- allocate address space without data loading
- when the programming starts running,

all the contents of the bss section are zeroed

- `.lcomm` defines a local common symbol in the bss section
- `.comm` can be used to declare a common symbol
- `.section name, "b" [COFF]`
- `.section name, "a" [ELF]`
- `.skip size, 0`

