OpenMP Synchronization (5A)

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Based on

https://www.openmp.org/wp-content/uploads/OpenMP-4.0-C.pdf

Synchronization (1)

Synchronization I

• Threads communicate through shared variables.

Uncoordinated access of these variables can lead to undesired effects.

– E.g. two threads update (write) a shared variable in the same step of execution, the result is dependent on the way this variable is accessed. This is called a race condition.

https://www3.nd.edu/~zxu2/acms60212-40212-S12/Lec-11-02.pdf

Synchronization (2)

- To prevent race condition, the access to shared variables must be synchronized.
- Synchronization can be time consuming.
- The barrier directive is set to synchronize all threads. All threads wait at the barrier until all of them have arrived.

https://www3.nd.edu/~zxu2/acms60212-40212-S12/Lec-11-02.pdf

Synchronization (3)

Synchronization II

• Synchronization imposes order constraints and is used to protect access to shared data

• High level synchronization:

– critical

– atomic

– barrier

– ordered

• Low level synchronization

– flush

– locks (both simple and nested)

https://www3.nd.edu/~zxu2/acms60212-40212-S12/Lec-11-02.pdf

Critical (1)

Synchronization: critical

• Mutual exclusion: only one thread at a time can enter a critical region.

```
{
double res;
#pragma omp parallel
{
double B;
int i, id, nthrds;
id = omp_get_thread_num();
nthrds = omp_get_num_threads();
for(i=id; i<niters; i+=nthrds){
B = some_work(i);
#pragma omp critical
consume(B,res);
}
}
}
```

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Critical (2)

Threads wait here: only one thread at a time calls consume(). So this is a piece of sequential code inside the for loop.

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Critical (3)

```
Sum = 0;
#pragma omp parallel shared(n,a,sum) private(TID,sumLocal)
{
        TID = omp_get_thread_num();
        sumLocal = 0;
        #pragma omp for
                for (i=0; I<n; i++)
                        sumLocal += a[i];
        #pragma omp critical (update_sum)
        {
                sum += sumLocal;
                printf("TID=%d: sumLocal=%d sum=%d\n", TID, sumLocal, sum)
        }
} /* --- End of parallel region --- */
```

https://www3.nd.edu/~zxu2/acms60212-40212-S12/Lec-11-02.pdf

Critical (4)

```
{
…
#pragma omp parallel
{
#pragma omp for nowait shared(best_cost)
for(i=0; i<N; i++){
int my_cost;
my_cost = estimate(i);
#pragma omp critical
{
if(best_cost < my_cost)
best_cost = my_cost;
}
}
}
}
```

Only one thread at a time executes if() statement. This ensures mutual exclusion when accessing shared data. Without critical, this will set up a race condition, in which the computation exhibits nondeterministic behavior when performed by multiple threads accessing a shared variable

https://www3.nd.edu/~zxu2/acms60212-40212-S12/Lec-11-02.pdf

References

[1] ftp://ftp.geoinfo.tuwien.ac.at/navratil/HaskellTutorial.pdf

[2] https://www.umiacs.umd.edu/~hal/docs/daume02yaht.pdf