# Applicative Methods (3B)

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#### The definition of Applicative

```
class (Functor f) => Applicative f where

pure :: a -> f a

(<*>) :: f (a -> b) -> f a -> f b
```

f (a -> b) :: a function wrapped in f
f a :: a value wrapped in f

The class has a two methods:

pure brings arbitrary values into the functor

(<\*>) takes a function wrapped in a functor f and a value wrapped in a functor f and returns the result of the application which is also wrapped in a functor f

#### The Maybe instance of Applicative

```
instance Applicative Maybe where

pure = Just

(Just f) <*> (Just x) = Just (f x)

_ <*> _ = Nothing
```

```
pure wraps the value with Just;
```

```
(<*>) applies
the <u>function</u> wrapped in <u>Just</u>
to the <u>value</u> wrapped in <u>Just</u> if both exist,
and results in <u>Nothing</u> otherwise.
```

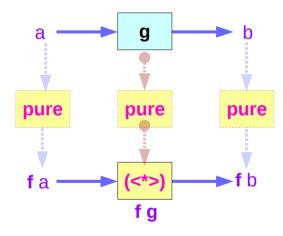
#### An Instance of the Applicative Typeclass

f : Functor, Applicative

(Functor f) => Applicative f

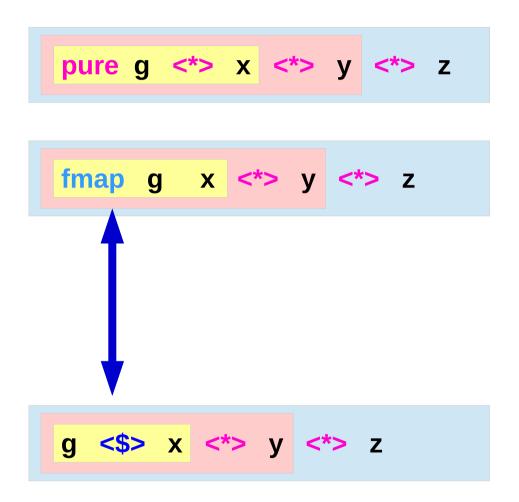
# instance Applicative Maybe where pure = Just Nothing <\*> \_ = Nothing (Just <u>f</u>) <\*> something = fmap <u>f</u> something

 $\underline{\mathbf{f}}$ : function in a context



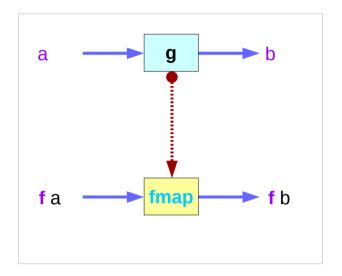
(Functor f) => Applicative f

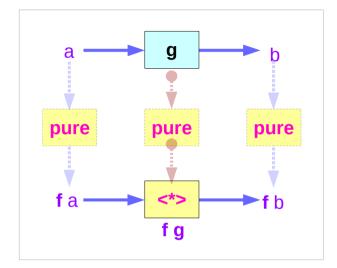
#### Left associative <\*>, fmap, and <\$>



infix operator <\$>

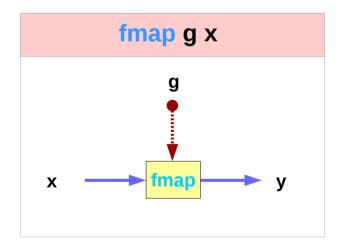
#### fmap g x = (pure g) <\*> x



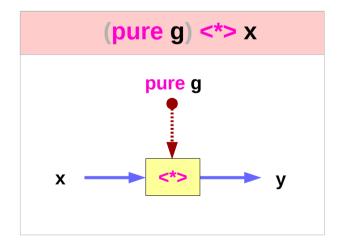


pure = f

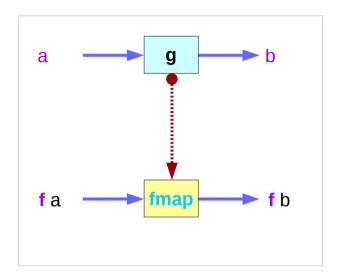
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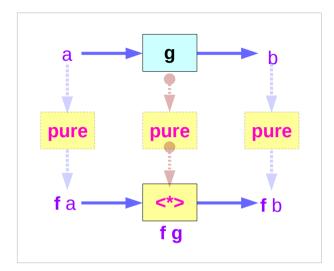


x :: f a y :: f b

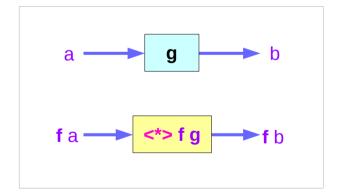


## fa <\*> fg





pure = f



#### Left associative <\*> examples

## Infix Operators <\*> vs <\$> - a type view

## h <\*> x <\*> y

h :: f (a -> b -> c)

x :: f a

y :: **f** b

x :: f a

h <\*> x :: f (b -> c)

x :: f a

h <\*> x :: f (b -> c)

y :: **f** b

h <\*> x <\*> y :: f c

x :: f a

y :: **f** b

x :: f a

g :: (a -> b -> c)

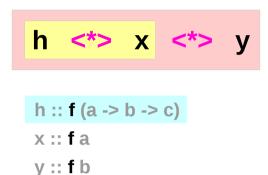
x :: f a

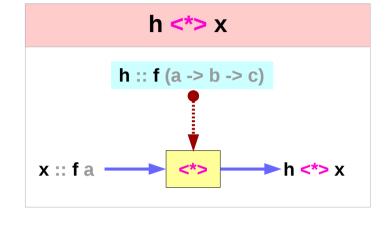
g <\$> x :: **f** (b -> c)

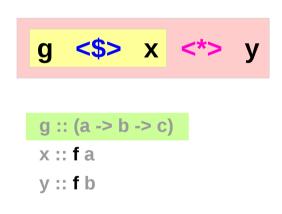
y :: **f** b

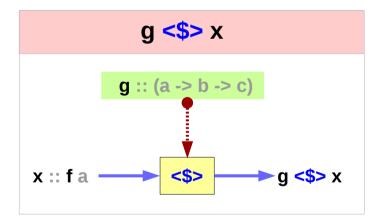
g <\$> x <\*> y :: f c

#### Infix Operators <\*> vs <\$> - a curried function view





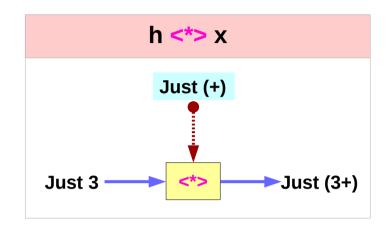




#### Infix Operators <\*> vs <\$> examples

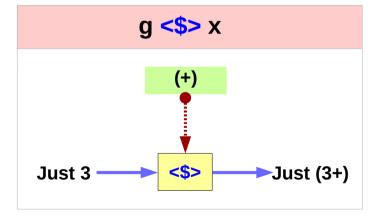


Just (+) <\*> Just 3 <\*> Just 2 Just (+3) <\*> Just 2 Just 5





(+) <\$> Just 3 <\*> Just 2 Just (+3) <\*> Just 2 Just 5



#### the minimal complete definition

```
class (Functor f) => Applicative f where
  pure :: a -> f a
  (<*>) :: f (a -> b) -> f a -> f b
```

the minimal complete definition

```
(<$>) :: (Functor f) => (a -> b) -> f a -> f b
g <$> x = fmap g x
```

Not in the minimal complete definition

```
instance Applicative Maybe where
pure = Just
Nothing <*> _ = Nothing
(Just g) <*> something = fmap g something
```

#### The Applicative Typeclass

```
Applicative is a <u>superclass</u> of Monad. every Monad is also a Functor and an Applicative
```

fmap, pure, (<\*>) can all be used with monads.

#### a Monad instance

requires **Functor** and **Applicative** instances. defines the types and roles of **return** and (>>)

fmap: defined in Functors

pure, (<\*>): defined in Applicatives

return, (>>): defined in Monads



#### (<\$>) infix operator

```
(<$>) :: (Functor f) => (a -> b) -> f a -> f b
g <$> x = fmap g x
```

#### The \$ operator is for avoiding parentheses

```
putStrLn (show (1 + 1))
putStrLn $ show (1 + 1)
putStrLn $ show $ 1 + 1 — right associative
```

(\$) calls the <u>function</u> which is its left-hand argument of \$ on the <u>value</u> which is its right-hand argument of \$

#### The Applicative Laws

The identity law: **pure** id <\*> v = v

id :: a -> a v :: f a

**Homomorphism:** 

pure g < > pure x = pure (g x)

q :: a -> b

x :: a

Interchange:

u <\*> pure y = pure (\$ y) <\*> u

u :: f (a -> b)

y :: a

**Composition:** 

Left associative

 $u <^*> v <^*> w = (u <^*> v) <^*> w$ 

u :: f (c -> b -> a)

v :: f c

u < > v :: f (b -> a)

w :: f b

u < > v < > w = fa

#### The Identity Law

The identity law

**pure** id <\*> v = v

pure to inject values into the functor in a default, featureless way, so that the result is as close as possible to the <u>plain</u> value.

applying the pure id morphism does nothing, exactly like with the plain **id** function.

#### The Homomorphism Law

The homomorphism law

**pure g <\*> pure x = pure (g x)** 

g :: a -> b

x :: a

applying a "pure" <u>function</u> to a "pure" <u>value</u> is the same as applying the <u>function</u> to the <u>value</u> in the <u>ordinary way</u> and then using **pure** on the result.

means **pure** <u>preserves</u> function application.

applying a non-effectful function g

to a <u>non-effectful</u> argument x in an <u>effectful</u> <u>context</u> <u>pure</u> is the same as just **applying** the function g to the argument x and then injecting the result (f x) into the <u>effectual context</u> with <u>pure</u>.

#### The Interchange Law

The interchange law

(\$ y) is the function that supplies yas argument to another functiona higher order function

**Function \$ Argument \$ y**( **y** ) as a single argument

applying a <u>morphism</u> **u** to a <u>"pure" value</u> <u>pure</u> **y** is the same as applying <u>pure</u> (\$ **y**) to the <u>morphism</u> **u** 

Just (+3) <\*> Just 2 Just (\$ 2) <\*> Just (+3)

when evaluating the application of an <u>effectful function</u> (**u**) to a <u>pure argument</u> (<u>pure y</u>), the <u>order doesn't matter</u> – commutative.

#### The Composition Law

The composition law pure (.) <\*> u <\*> v <\*> w = u <\*> (v <\*> w) <math>w :: fa v :: f(a -> b) u :: f(b -> c)

pure (.) <u>composes</u> morphisms similarly to how (.) <u>composes</u> functions:

applying the composed mourphism

it is expressing a sort of <u>associativity</u> property of (<\*>).

```
w :: f a -- value
v :: f (a -> b) -- func1
u : f (b -> c) -- func2
```

```
v <*> w :: f b
u <*> (v <*> w) :: f c
```

```
pure (.) <*> u <*> v :: f (a -> c )
pure (.) <*> u <*> v <*> w :: f c
```

#### The Composition Law and Left Associativity

```
The composition law
                             pure (.) <*> u <*> v <*> w = u <*> (v <*> w)
                                                                              w :: fa v :: f(a -> b) u :: f(b -> c)
                                                                             f(b \rightarrow c)
                                                                                              f (a -> b)
                                                                                                 pure h
                                                                               pure g
pure (.) <*> pure g <*> pure h <*> pure x
                                                     (g.h) x
                                                                                       pure (.)
((pure (.) <*> pure q) <*> pure h) <*> pure x
                                                                                      f(a \rightarrow c)
= pure q <*> (pure h <*> pure x)
                                                     g (h x)
                                                                              u = pure g :: f (b -> c) g :: (b -> c)
                                                                              v = pure h :: f (a -> b) h :: (a -> b)
                                                                              w = pure x :: f a
                                                                                                          x :: a
Left associative
                       u <*> v <*> w = (u <*> v) <*> w
                                                                             u :: f(c -> b -> a) v :: fc w :: fb
                                                                             u :: f (c -> b -> a)
                                                                             v :: f c
                                                                             u < > v :: f (b -> a)
                                                                             w :: f b
                                                                             u < > v < > w = f a
  https://en.wikibooks.org/wiki/Haskell/Applicative functors
```

#### liftA2

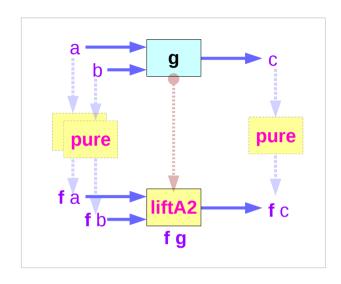
liftA2 :: (a -> b -> c) -> f a -> f b -> f c

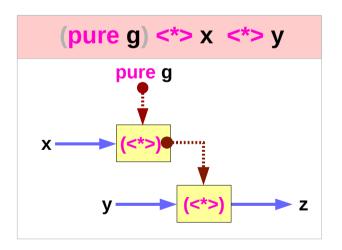
<u>lift</u> a <u>binary</u> <u>function</u> (**a->b->c**) to actions.

Some functors support an implementation of **liftA2** that is more efficient than the default one.

**liftA2** may have an <u>efficient</u> implementation whereas **fmap** is an <u>expensive</u> operation,

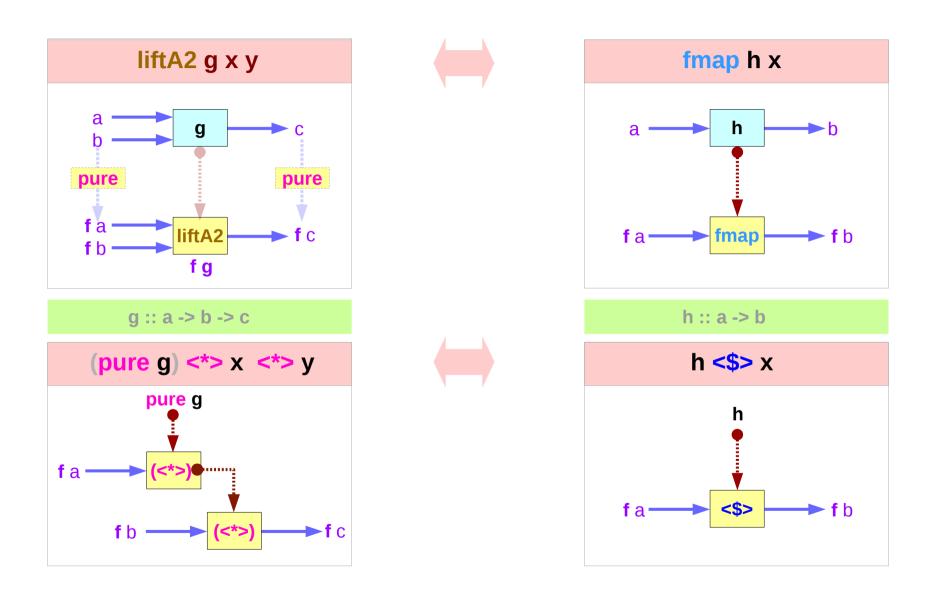
sometimes better to use **liftA2** than to use **fmap** over the structure and then use <\*>.



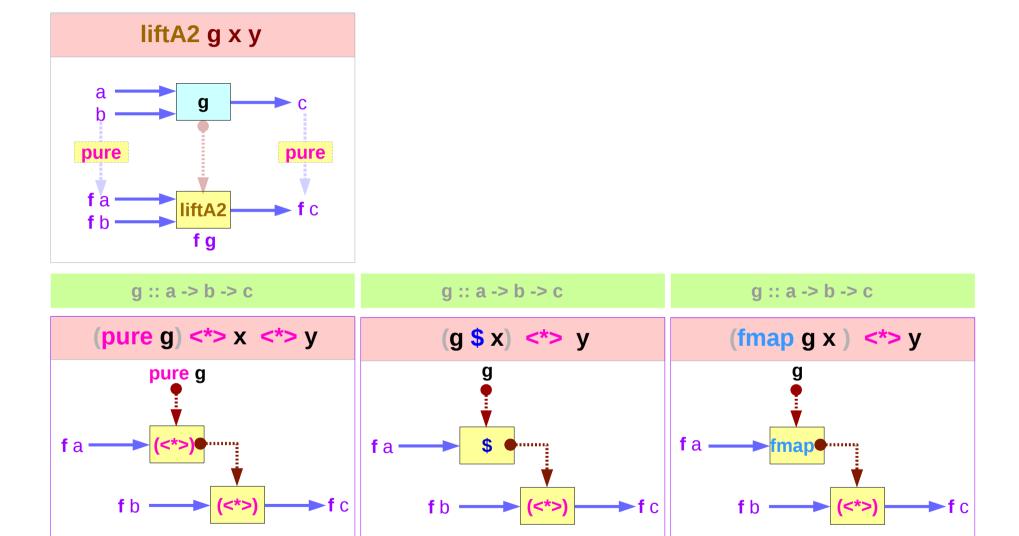


http://hackage.haskell.org/package/base-4.10.1.0/docs/Control-Applicative.html#v:liftA2

## liftA2, <\*>, fmap, <\$>



## pure g <\*> x <\*> y equivalent



#### liftA2

```
liftA2 g x y
liftA2 :: (a -> b -> c) -> f a -> f b -> f c
g:: a -> b -> c
x :: f a
y :: f b
liftA2 g x y :: f c
pure g <*> x <*> y
g:: a -> b -> c
x :: f a
y :: f b
z :: f c
pure g <*> x <*> y :: f c
```

liftA2 g x y g pure pure liftA2 f g g :: a -> b -> c (pure g) <\*> x <\*> y pure g

https://wiki.haskell.org/Applicative functor

#### **Limitations of Functors**

```
(a \rightarrow b \rightarrow c) \rightarrow (f a \rightarrow f b \rightarrow f c) – let fmap2 Functor as an extension of fmap
fmap :: (a -> b) -> (f a -> f b)
fmap2 :: Functor f => (a -> b -> c) -> (f a -> f b -> f c)
fmap2 h fa fb = undefined
h :: a -> b -> c
fa :: f a
fb :: f b
h
        :: a -> (b -> c)
fmap h :: fa \rightarrow f(b \rightarrow c)
fmap h fa :: f(b \rightarrow c) — now f(b \rightarrow c) must be applied to f(b)
fmap gives us a way to <u>apply</u> functions (\mathbf{a} \rightarrow \mathbf{b}) to <u>values</u> (\mathbf{f} \mathbf{a}) inside a <u>Functor</u> context,
but fmap cannot be used to apply a functions f (b -> c)
```

which are themselves in a Functor context to values **f b** in a Functor context.

http://www.openhaskell.com/lectures/applicative.html

#### pure, fmap, and liftA2

```
class Functor f => Applicative f where
 pure :: a -> f a
 (<*>) :: f (a -> b) -> f a -> f b
pure
      :: a -> f a
                                                  - fmap0 → pure
fmap :: (a -> b) -> f a -> f b
                                         – fmap1 → fmap
fmap2 :: (a -> b -> c) -> fa -> fb -> fc
                                          – fmap2 → liftA2
liftA2 :: Applicative f => (a -> b -> c) -> f a -> f b -> f c
liftA2 h fa fb = (h `fmap` fa) <*> fb
liftA2 h fa fb = h < $> fa <*> fb
(<$>) :: Functor f => (a -> b) -> f a -> f b
(<$>) = fmap
liftA2 :: Applicative f => (a -> b -> c -> d) -> f a -> f b -> f c -> f d
liftA3 h fa fb fc = ((h <$> fa) <*> fb) <*> fc
```

http://www.openhaskell.com/lectures/applicative.html

#### liftA2 examples

```
liftA2 :: Applicative f => (a -> b -> c) -> f a -> f b -> f c
```

#### <\*> or liftA2 implementations

liftA2 :: (a -> b -> c) -> f a -> f b -> f c

A minimal complete definition:

either one of the two

- 1) pure and <\*>
- 2) pure and liftA2

If it defines <u>both</u>, then they must behave the same as their default definitions:

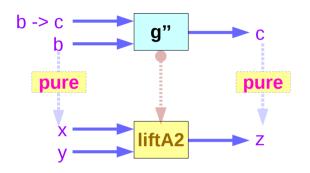
2) liftA2 g x y

#### (<\*>) = liftA2 id

**liftA2 id** 
$$x y = id < x < y = x < y$$

liftA2 id 
$$x y = x <*> y$$



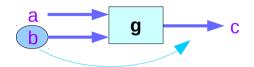


<b>liftA2</b> g x y = g < $x < y$	g :: a -> b -> c	x :: f a y :: f b
liftA2 g" x y = g" <\$> x <*> y	g" :: (b -> c) -> b -> c	x :: f (b -> c) y :: f b
liftA2 id x y = id $<$ x $<$ y = x $<$ y	id :: (b -> c) -> (b -> c)	x :: f (b -> c) y :: f b

(<\*>) = liftA2 id

http://hackage.haskell.org/package/base-4.10.1.0/docs/Control-Applicative.html#v:liftA2

## g":: (b->c) -> b->c



view the function as having one input only

consider the case when  $\mathbf{a}$  is  $(\mathbf{b} \rightarrow \mathbf{c})$ 

Then **g**" is the same as **id** 

$$g'' :: (b -> c) -> (b -> c)$$

#### Results and effects in a scope

Actually, using the **liftA** commands

we can pull <u>results</u> of applicative functors

into a scope where we can talk

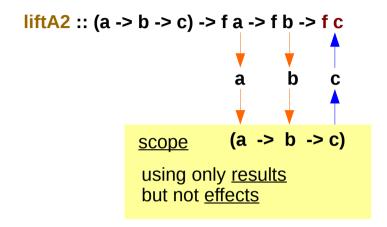
exclusively about <u>functor</u> <u>results</u>

and not about <u>effects</u>. **f c** 

Note that <u>functor results</u> can also be <u>functions</u>. **c** 

This scope is simply a function,

which contains the code that we used in the non-functorial setting.



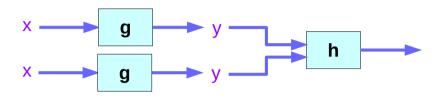
http://hackage.haskell.org/package/base-4.10.1.0/docs/Control-Applicative.html#v:liftA2

C

#### liftA3 – a non-functorial expression

#### Consider the non-functorial expression:

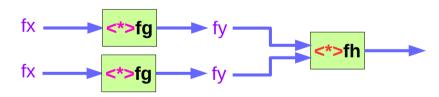
x :: x g :: x -> y h :: y -> y -> z let y = g x



generalization

in hyy

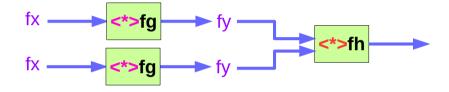
fx :: f x fg :: f (x -> y) fh :: f (y -> y -> z)



https://wiki.haskell.org/Applicative\_functor

## liftA3 – using <\*> only

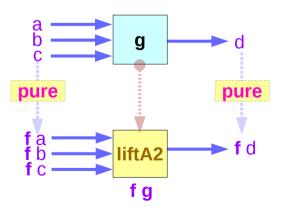
if **fy** writes something to the terminal then **fh** <\*> **fy** <\*> **fy** writes twice.



this runs the effect of fy twice.

How the <u>effect</u> is run only <u>once</u> and the <u>result</u> is used <u>twice</u>?

→ utilize liftA3



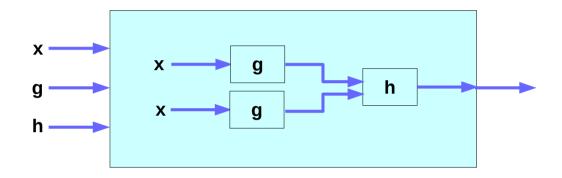
https://wiki.haskell.org/Applicative\_functor

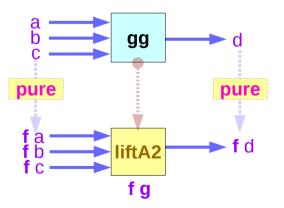
#### **liftA3** – using three input function

liftA3 :: Applicative f => (a -> b -> c -> d) -> f a -> f b -> f c -> f dliftA3 h fa fb fc = ((h <\$> fa) <\*> fb) <\*> fc

#### liftA3

$$(x g h \rightarrow let y = g x in h y y)$$
  
fx fg fh





http://hackage.haskell.org/package/base-4.10.1.0/docs/Control-Applicative.html#v:liftA2

#### liftA3 – effects, results and scopes

Actually, using the **liftA** commands

we can pull <u>results</u> of applicative functors **y** from **fy** 

into a scope where we can talk  $y \rightarrow y \rightarrow z$ 

exclusively about <u>functor results</u> **y** 

and not about <u>effects</u>. **fy** 

Note that <u>functor results</u> can also be <u>functions</u>. **y** 

This scope is simply a function,  $y \rightarrow y \rightarrow z$ 

which contains the code that we used in the non-functorial setting.

#### liftA3

```
(x g h \rightarrow let y = g x in h y y)
fx fg fh
```

The order of effects is entirely determined by the order of arguments to liftA3

https://wiki.haskell.org/Applicative\_functor

# liftA2 (<\*>)

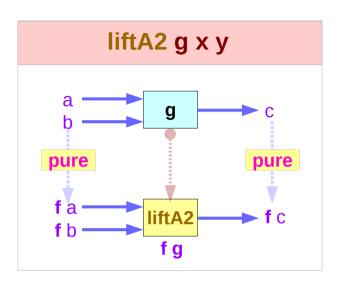
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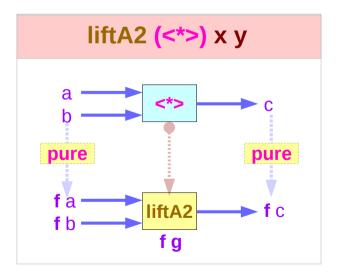
The wiki article says that **liftA2** (<\*>) can be used to <u>compose applicative functors</u>. It's easy to see how to use it from its type:

```
o :: (Applicative f, Applicative f1) =>
f (f1 (a -> b)) -> f (f1 a) -> f (f1 b)
o = liftA2 (<*>)
```

### liftA2 (<\*>) for composite applicative functors

liftA2 (<\*>) can be used to compose applicative functors.





## **liftA2 (<\*>)** Examples (1)

```
if f is Maybe and f1 is [] we get:
```

Just [(+1),(+6)] `o` Just [1, 6] Just [2,7,7,12]

(+1) [1, 6]

(+6) [1, 6]

### **liftA2 (<\*>)** Examples (2)

```
[Just (+1),Just (+6)] `o` [Just 1, Just 6]

[Just (+1),Just (+6)] `o` [Just 1, Just 6]

[Just 2, Just 7, Just 7, Just 12]

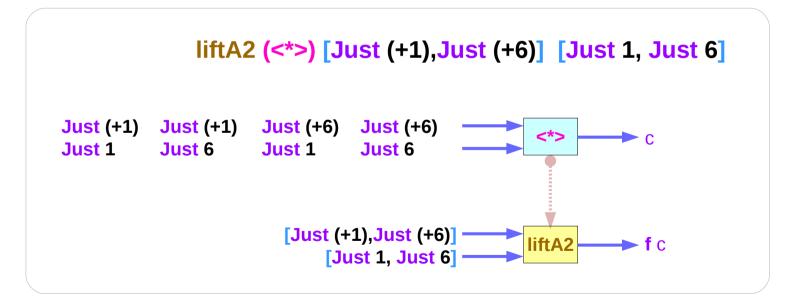
[Just (+1) [Just 1, Just 6]

[Just (+6) [Just 1, Just 6]

[Just (+1),Just (+6)] [Just 1, Just 6]

[Just (+1) <*> Just 1, Just (+6) <*> Just 1, Just (+6) <*> Just 6]

[Just 2, Just 7, Just 7, Just 12]
```



## liftA2 (:)

```
liftA2 (:) "abc" ["pqr", "xyz"]
["apqr", "axyz", "bpqr", "bxyz", "cpqr", "cxyz"]
```

```
liftA2 (:) "abc" ["pqr", "xyz"]

"a" "a" "b" "b" "c" "c" "c" "yz" "pqr" "xyz" "pqr" "xyz" | iftA2 | f c
```

### liftA2 (:)

```
(liftA2 . liftA2) (:) (Just "abc") (Just ["pqr", "xyz"])

Just ["apqr", "axyz", "bpqr", "bxyz", "cpqr", "cxyz"]
```

```
(liftA2 . liftA2) (:) "abc" ["pqr", "xyz"]
"xyz" "pqr" "xyz" "pqr"
            ["pqr", "xyz"]
        Just ["pqr", "xyz"]
```

### related operators

#### Functor map <\$>

```
(<) :: Functor f => (a -> b) -> f a -> f b
```

(
$$$>$$
) :: Functor  $f => fa$  ->  $b$  ->  $fb$ 

functor \$> value = const value <\$> functor

The <\$> operator is just a synonym for the fmap function from the Functor typeclass.

This function generalizes the map function for lists to many other data types, such as Maybe, IO, and Map.

#### **<\$, <\$>, \$>** operators

there are two additional operators provided which <u>replace</u> a **value** inside a Functor instead of applying a function.

This can be both more convenient in some cases, as well as for some Functors be more efficient.

$$x < y = y > x$$
  $y :: functor$ 

$$x > y = y < x$$
 x :: functor

### <\$, <\$>, \$> operators examples

#### import Data.Functor

Prelude> Just 1 \$> 2

Just 2

Prelude> Just 2 \$> 1

Just 1

Prelude> 1 <\$ Just 3

Just 1

Prelude> 3 <\$ Just 1

Just 3

Prelude> 1 <\$ Just 3

Just 1

Prelude> 3 <\$ Just 1

Just 3

#### import Data.Functor

Prelude> (+1) <\$> Just 2

Just 3

Prelude> (+1) <\$> Just 3

Just 4

Prelude> (+1) <\$> Nothing

Nothing

Prelude> const 2 <\$> Just 111

Just 2

https://www.schoolofhaskell.com/school/to-infinity-and-beyond/pick-of-the-week/Simple%20examples

#### <\$> examples

```
#!/usr/bin/env stack
-- stack --resolver ghc-7.10.3 runghc
import Data.Monoid ((<>))

main :: IO ()
main = do
    putStrLn "Enter your year of birth"
    year <- read <$> getLine
    let age :: Int
        age = 2020 - year
    putStrLn $ "Age in 2020: " <> show age
```

#### <\*> related operators

Applicative function application <\*>

```
(<*>) :: Applicative f => f(a -> b) -> fa -> fb
```

(<\*) :: Applicative 
$$f => f a$$
 ->  $f b -> f a$ 

Commonly seen with <\$>, <\*> is an operator that applies a wrapped function to a wrapped value. It is part of the Applicative typeclass, and is very often seen in code like the following:

## \*> operator

To go along with this, we have two helper operators that are less frequently used:

\*> ignores the value from the first argument. It can be defined as:

Or in do-notation:

For Monads, this is completely equivalent to >>.



<\* is the same thing in reverse: perform the first action then the second,

but only take the value from the first action.

Again, definitions in terms of <\*> and do-notation:

res <- a1

\_ <- a2

return res



For cases when you're dealing with a Monad, this is equivalent to:

```
do x <- bar
y <- baz
return (foo x y)</pre>
```

Other common examples including parsers and serialization libraries.

Here's an example you might see using the aeson package:

```
data Person = Person { name :: Text, age :: Int } deriving Show
```

-- We expect a JSON object, so we fail at any non-Object value.

instance FromJSON Person where

```
parseJSON (Object v) = Person <$> v .: "name" <*> v .: "age"
parseJSON _ = empty
```

### (\*> v.s. >>) and (pure v.s. return)

```
(*>) :: Applicative f \Rightarrow fa \Rightarrow fb \Rightarrow fb
```

$$(>>) :: Monad m => m a -> m b -> m b$$

the constraint changes from **Applicative** to **Monad**.

https://en.wikibooks.org/wiki/Haskell/Applicative\_functors

#### References

- [1] ftp://ftp.geoinfo.tuwien.ac.at/navratil/HaskellTutorial.pdf
- [2] https://www.umiacs.umd.edu/~hal/docs/daume02yaht.pdf