

Pointers (1A)

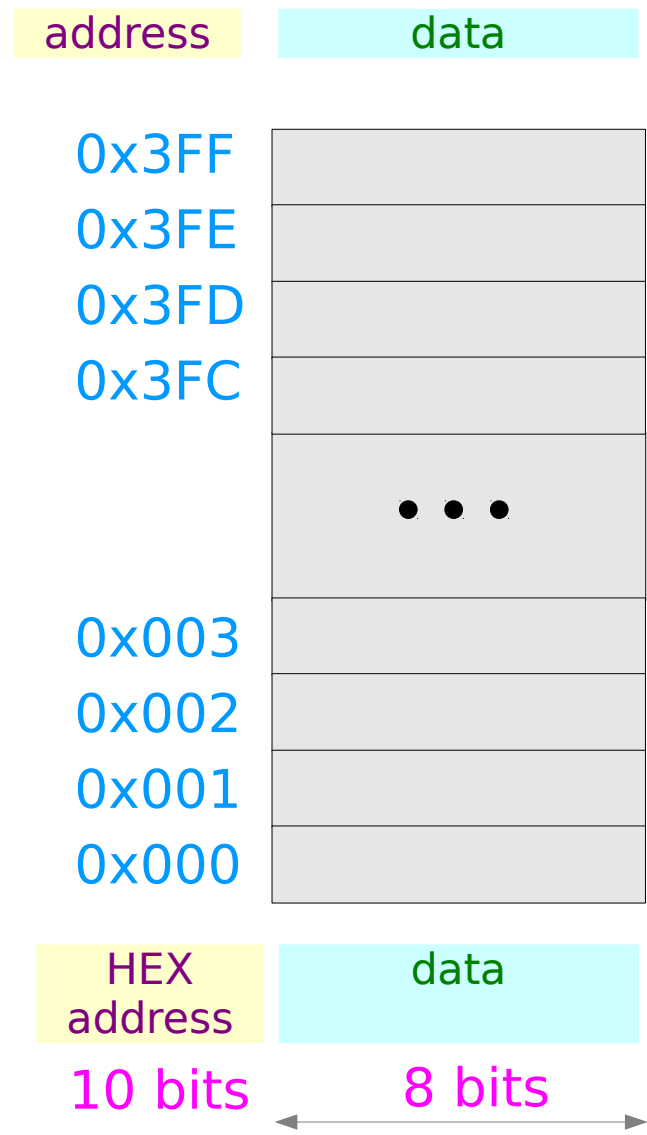
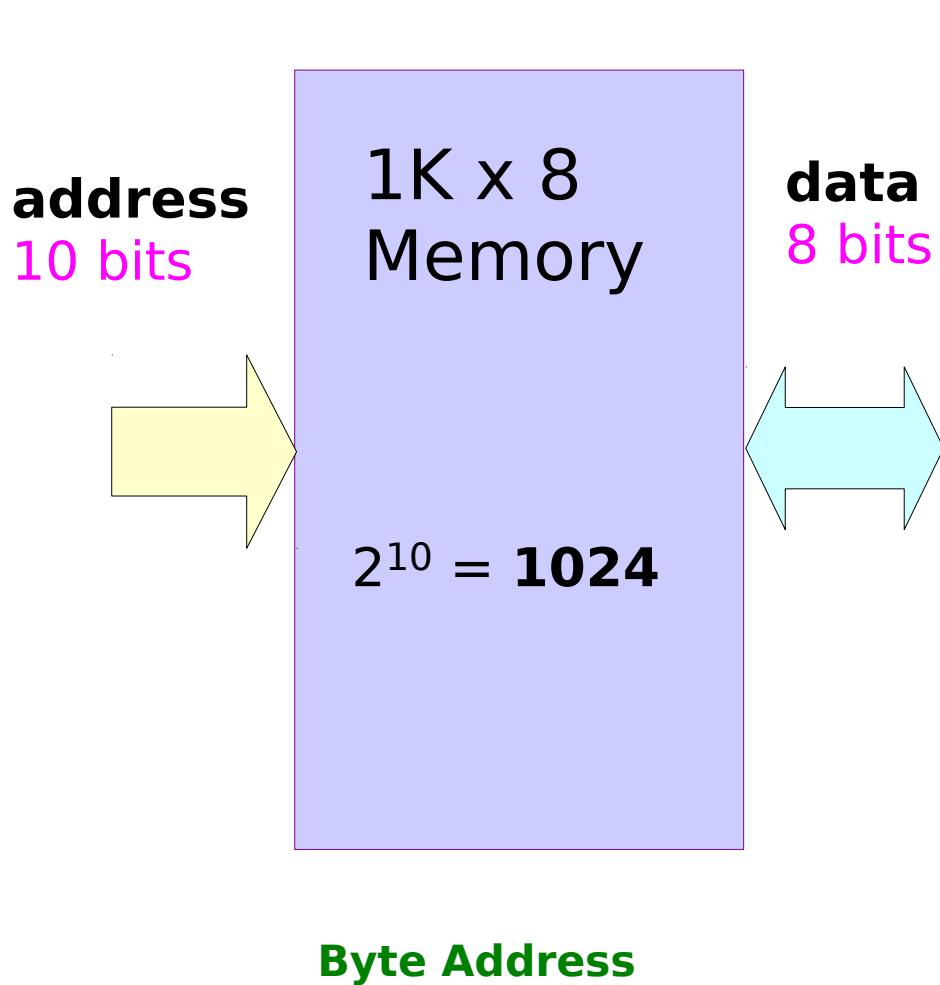
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Byte Address and Data in a Memory



Variables

```
int a;
```

a can hold an *integer* value

address

data

&a

a

```
a = 100;
```

a holds the *integer* 100

address

data

&a

a ← 100

Pointer Variables

```
int * p;
```

`p` holds an address

`p` can hold the address of an `int` data

`*p` can hold an integer value

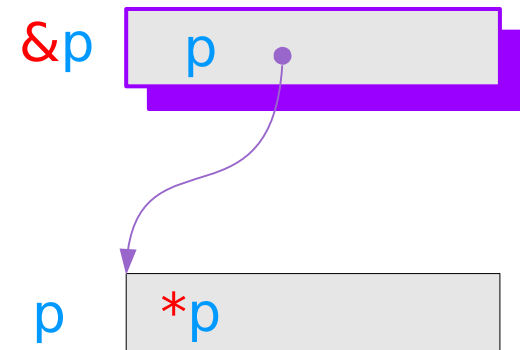
type **variable**

```
int * p;
```

pointer to int

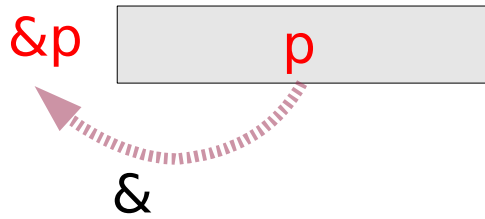
```
int * p;
```

int

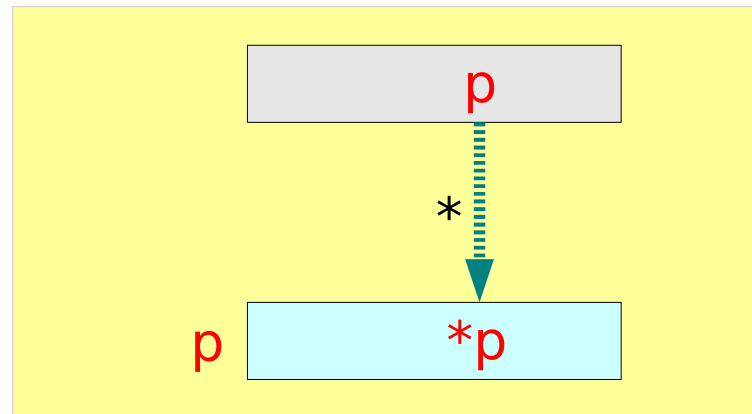
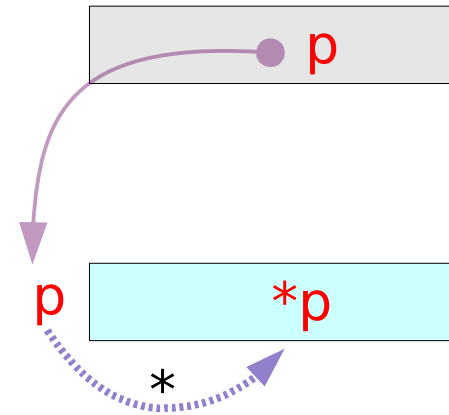


Dereferencing

*The address of a variable :
Address of operator &*



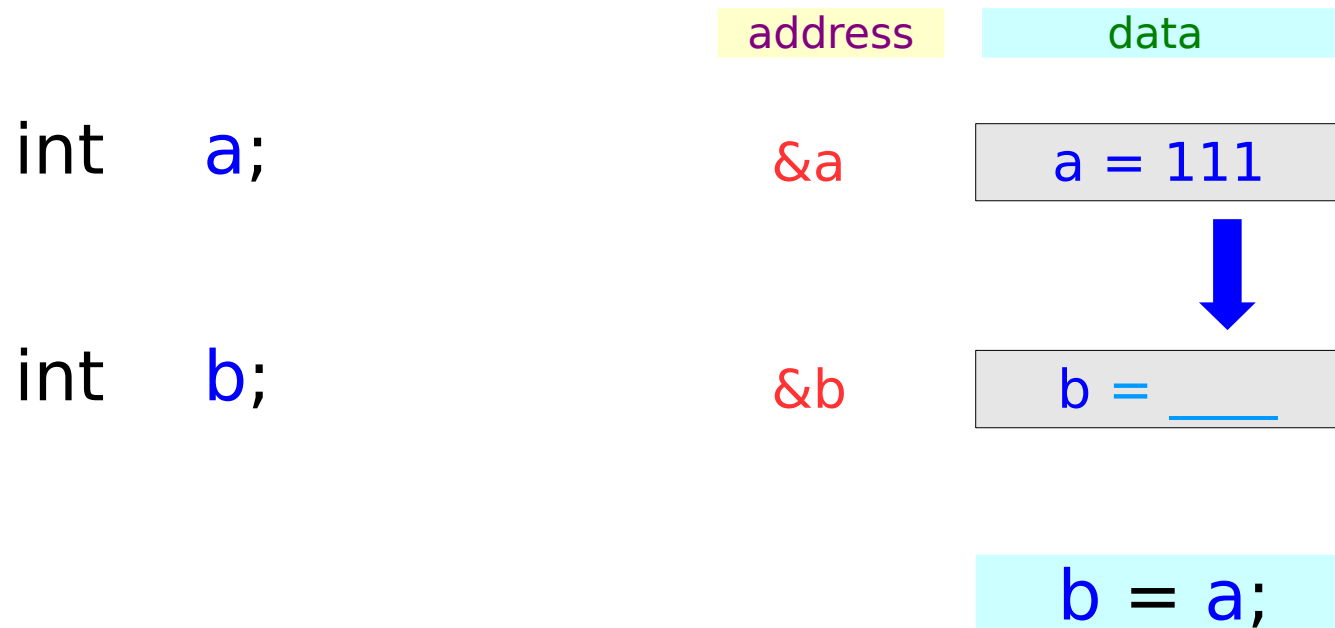
*The content of a pointed location :
Dereferencing operator **



Variables and their addresses

	address	data
<code>int a;</code>	<code>&a</code>	<code>a</code>
<code>int *p;</code>	<code>&p</code>	<code>p</code>

Assignment of a value



Assignment of an address

```
int a;
```

```
int *p;
```

address

data

&a

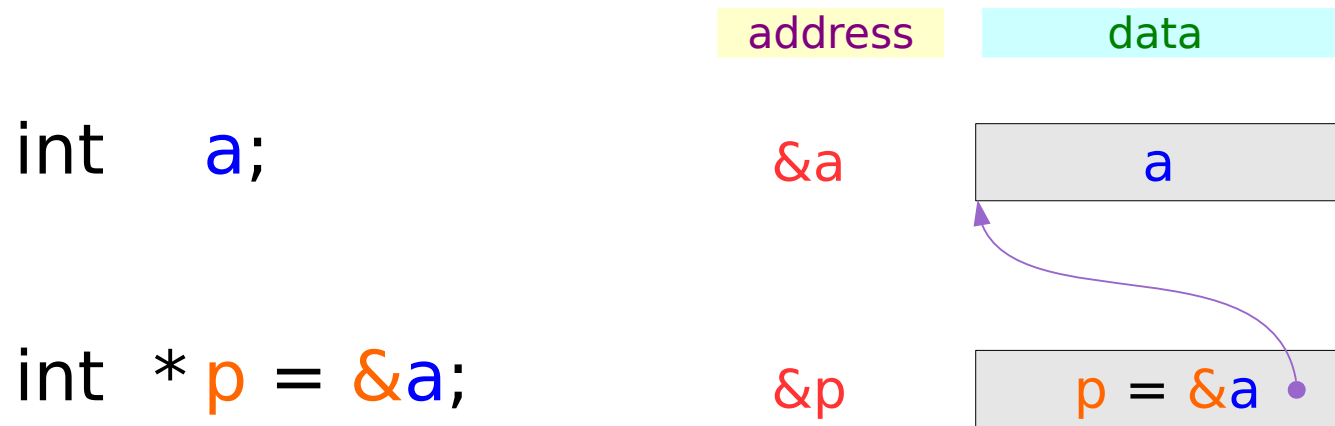
a = 111

&p

p = _____

p = &a;

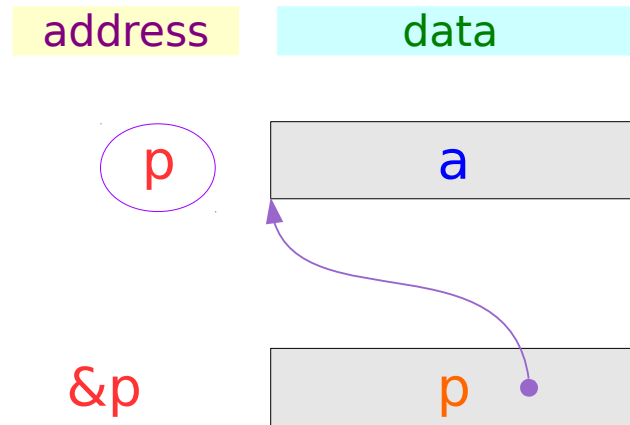
Variables with initializations



Pointed addresses : p

```
int a;
```

```
int *p = &a;
```

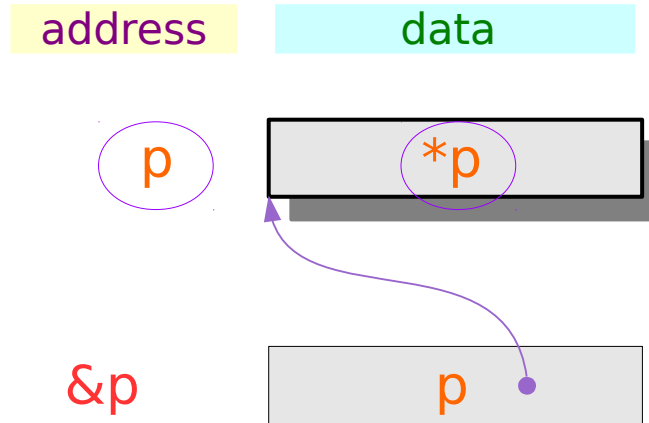


$p \equiv \&a$

Dereferenced Variable : *p

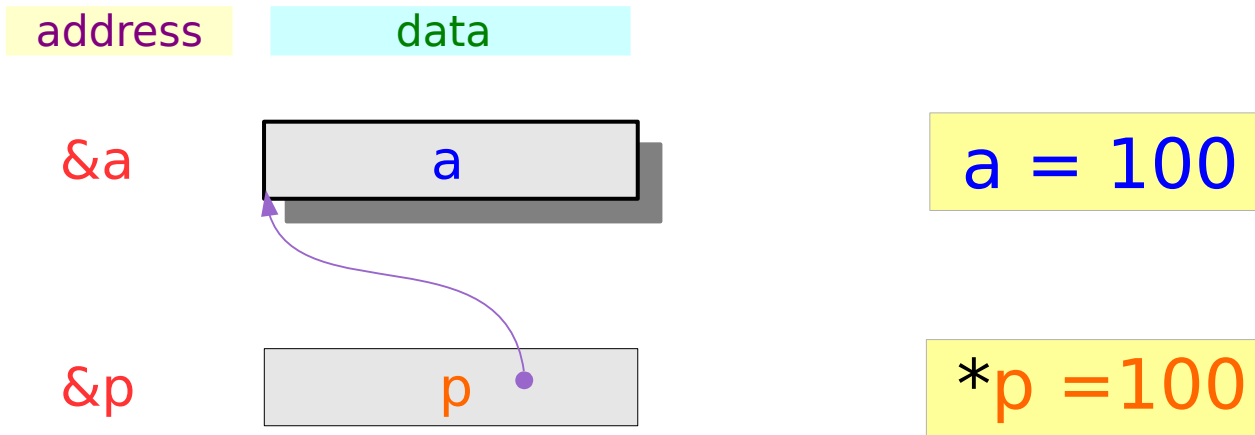
```
int a;
```

```
int *p = &a;
```



$p \equiv \&a$
 $*p \equiv a$

Two way to access: `a` and `*p`



- 1) Read/Write `a`
- 2) Read/Write `*p`

-
1. Pass by Reference
 2. Arrays

Pass by Reference

Variable Scopes

```
int func1 (int a, int b)  
{  
    int i, int j;  
    ...  
    ...  
    ...  
    ...  
}
```

i and **j**'s
variable scope



cannot access
each other

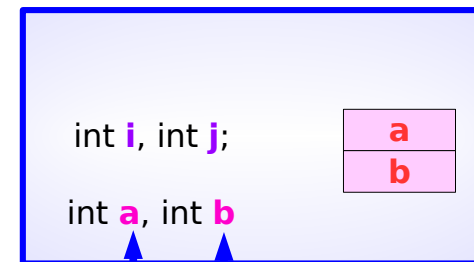
```
int main ()  
{  
    int x, int y;  
    ...  
    ...  
    func1 ( 10, 20 );  
    ...  
    ...  
}
```

x and **y**'s
variable scope

Only **top** stack frame is active
and its variable can be accessed

Communications are performed
only through the **parameter** variables

func1's
Stack
Frame



(10, 20)

main's
Stack
Frame

int **x**, int **y**;

Pass by Reference

```
int func1 (int* a, int* b)  
{  
    int i, int j;  
    ...  
    ...  
    ...  
    ...  
}
```

x and **y** are made known to **func1**
func1 can read / write **x** and **y**
through their addresses

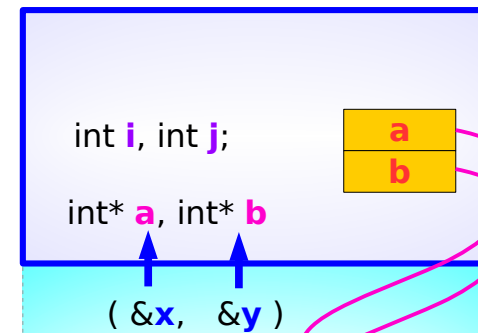
a=&**x**
b=&**y**

x and **y**'s
variable scope

```
int main ()  
{  
    int x, int y;  
    ...  
    ...  
    func1 ( &x, &y );  
    ...  
    ...  
}
```

***a**
***b**

func1's
Stack
Frame

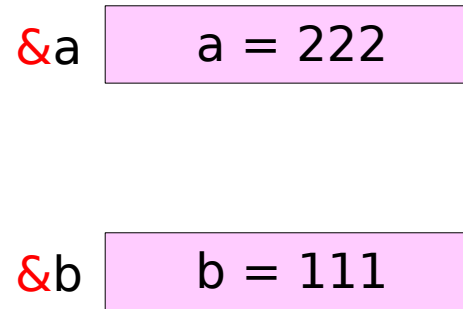
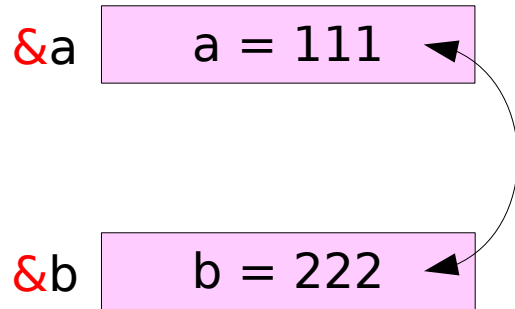


main's
Stack
Frame

int **x**, int **y**;

***a**
***b**

Swapping integers



```
int a, b;
```

```
swap( &a, &b );
```

```
swap( int *, int * );
```

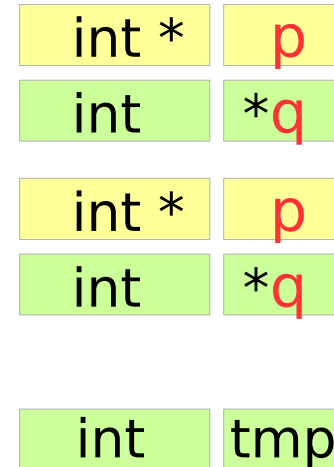
function call

function prototype

Pass by integer reference

```
void swap(int *p, int *q) {  
    int tmp;  
  
    tmp = *p;  
    *p = *q;  
    *q = tmp;  
}
```

```
int a, b;  
...  
swap( &a, &b );
```



Integer and Integer Pointer Types

```
int *m  
int *n
```

integer pointer declarations



a way of thinking

```
int * m  
int * n  
int *m  
int *n
```



```
m  
n  
*m  
*n
```

integer pointer variables

treated as integer variables

variables

```
int *
```

```
int
```

types

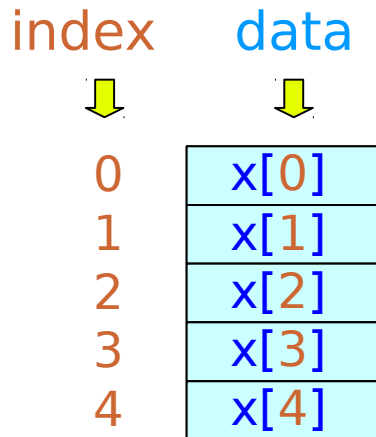
Arrays

Accessing array elements - using an address

```
int    x[5];
```

x holds the *starting address* of **5** consecutive **int** variables

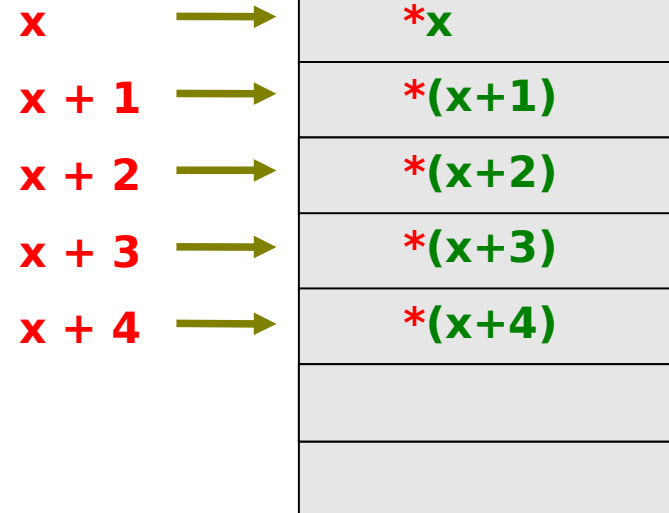
5 int variables



cannot change
address x
(constant)

address

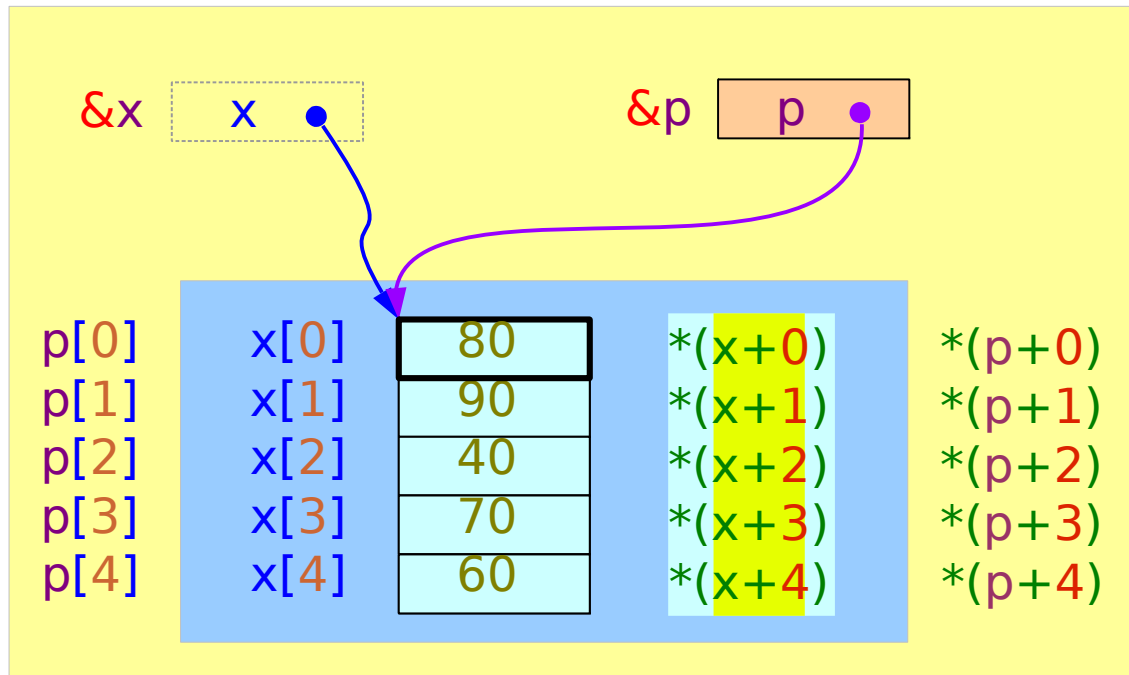
data



Accessing an Array with a Pointer Variable

```
int x [5] = { 1, 2, 3, 4, 5 };
```

```
int *p = x;
```



`x` is a constant symbol
cannot be changed

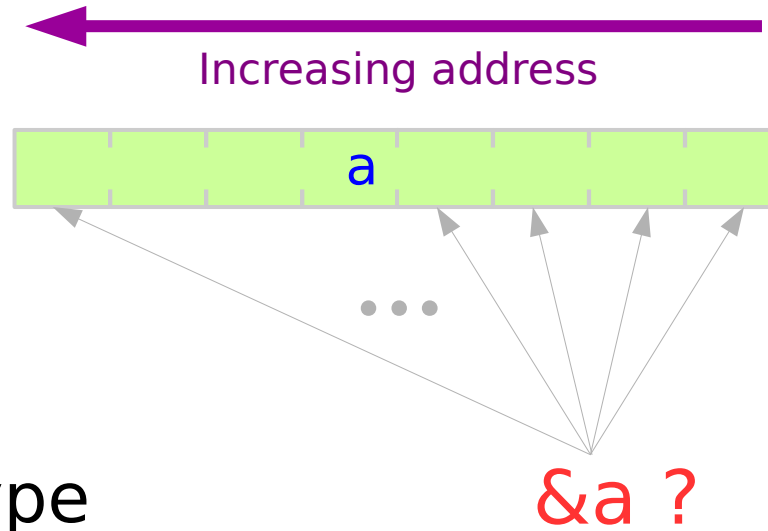
`p` is a variable
can point to other addresses

Byte Address
Little Endian
Big Endian

Byte Address

long a;

8-byte size data type



Numbers in Positional Notation

```
long a = 0x1020304050607080;
```

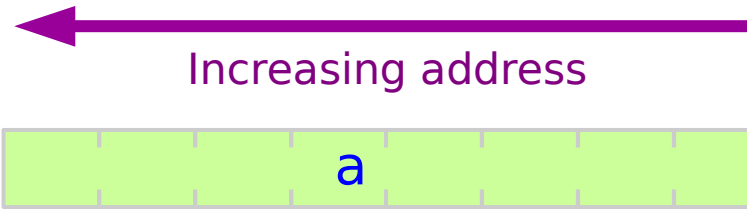
8 (bytes)

a_7 a_6 a_5 a_4 a_3 a_2 a_1 a_0

Most Significant Byte	$a_7 = 0x10 \dots 16^7$	the highest weight
	$a_6 = 0x20 \dots 16^6$	
	$a_5 = 0x30 \dots 16^5$	
	$a_4 = 0x40 \dots 16^4$	
	$a_3 = 0x50 \dots 16^3$	
	$a_2 = 0x60 \dots 16^2$	
Least Significant Byte	$a_1 = 0x70 \dots 16^1$	the lowest weight
	$a_0 = 0x80 \dots 16^0$	

Little / Big Endian Ordering of Bytes

long a;



MSByte Little Endian LSByte



$a_7 a_6 a_5 a_4 a_3 a_2 a_1 a_0$

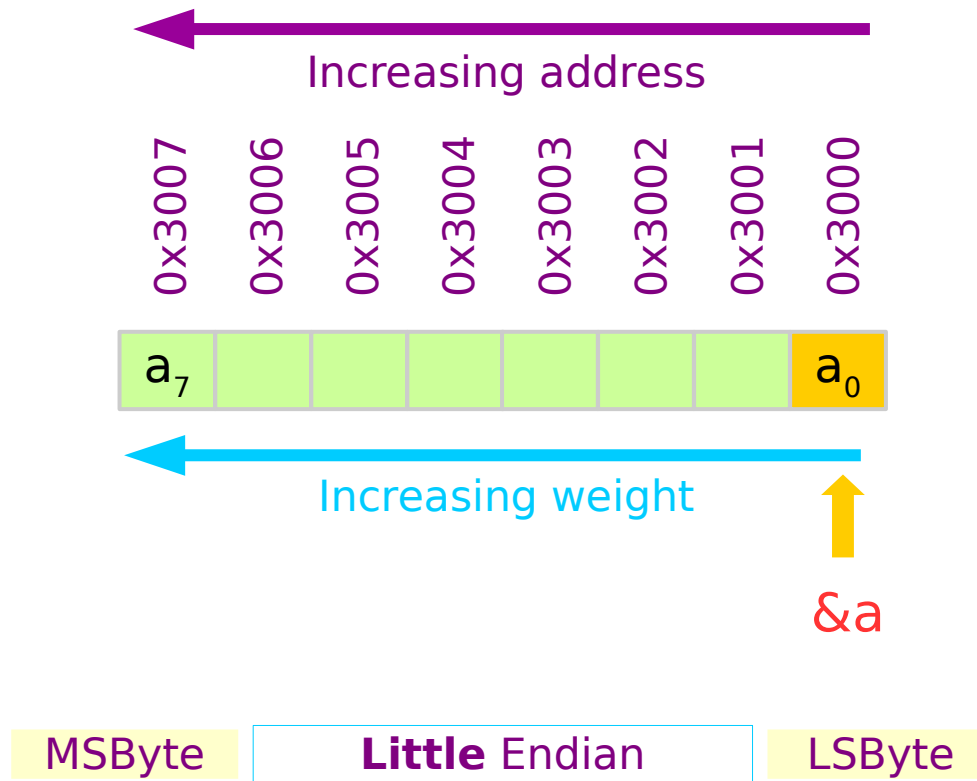
LSByte Big Endian MSByte



$a_0 a_1 a_2 a_3 a_4 a_5 a_6 a_7$

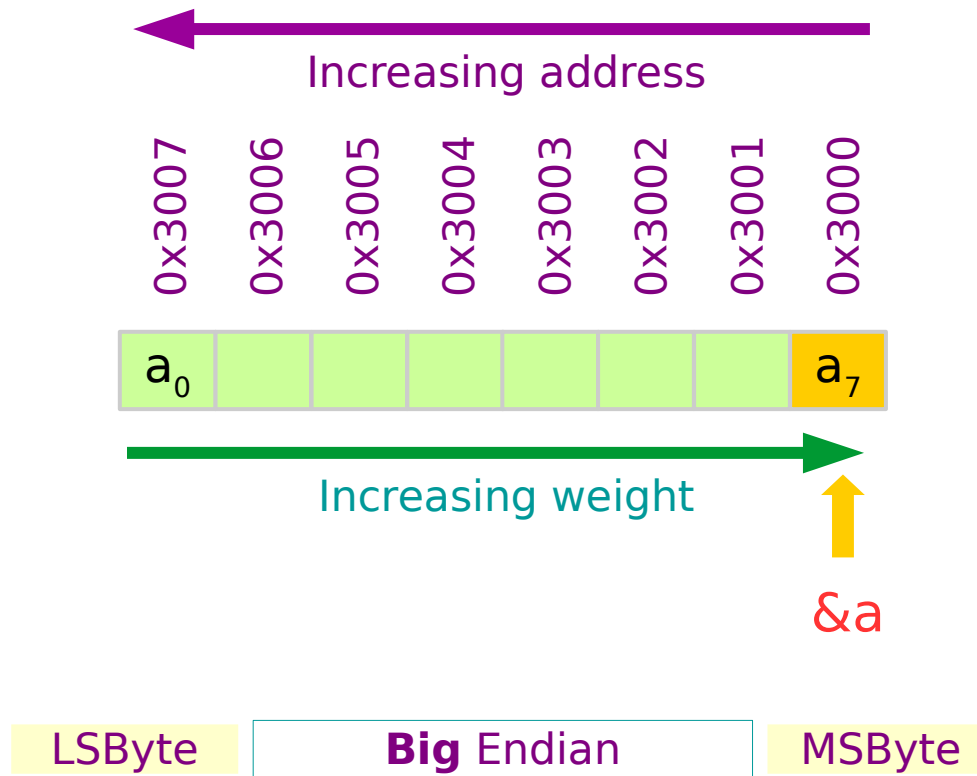
Little Endian Byte Address Example

long a;



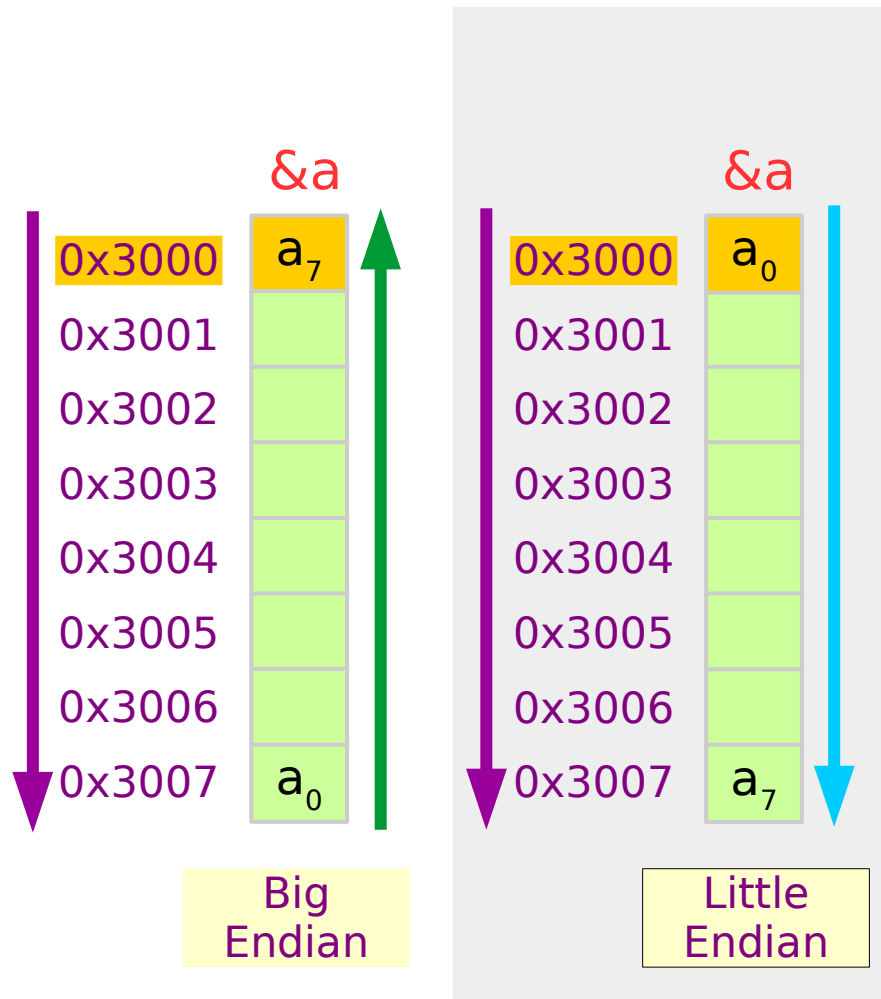
Big Endian Byte Address Example

long a;

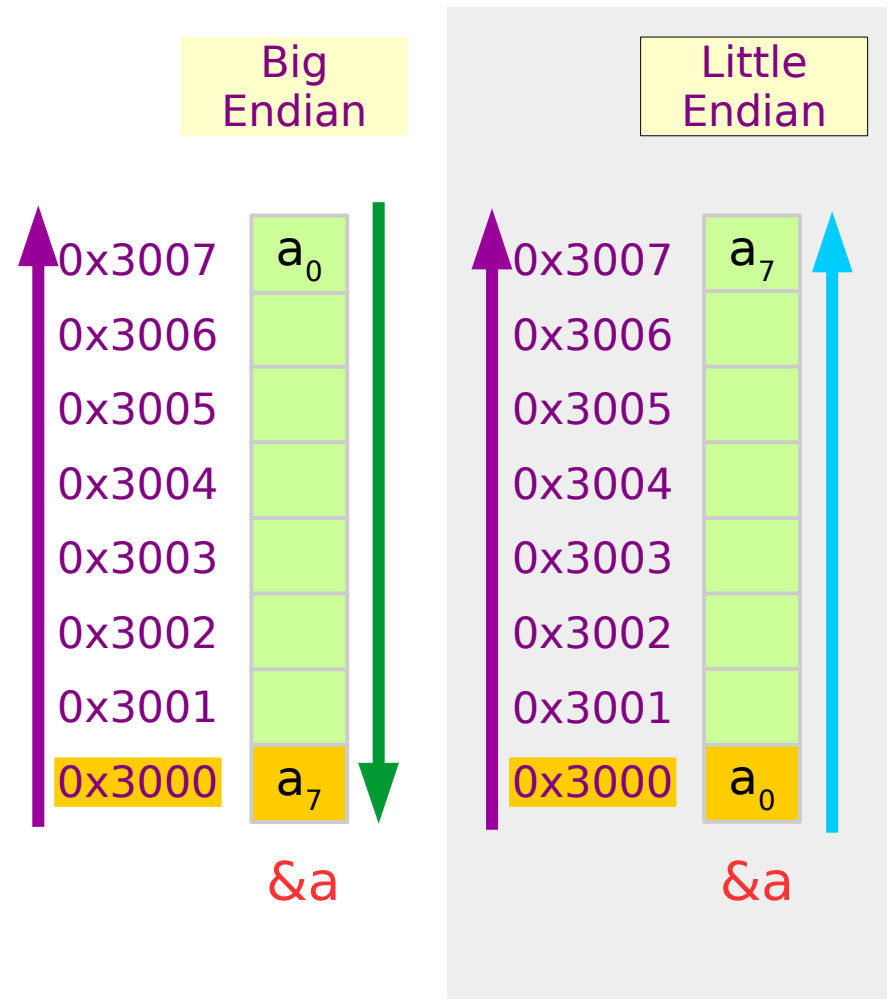


Representations of Endianness

downward, increasing address



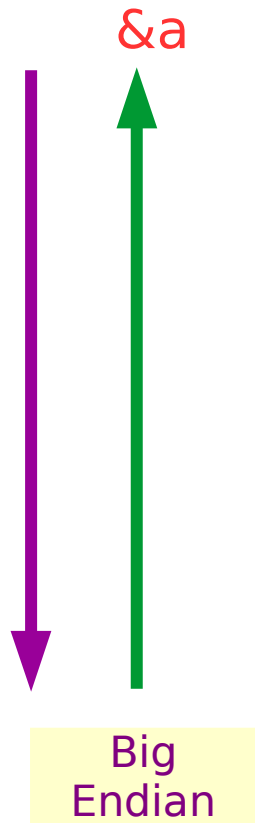
upward, increasing address



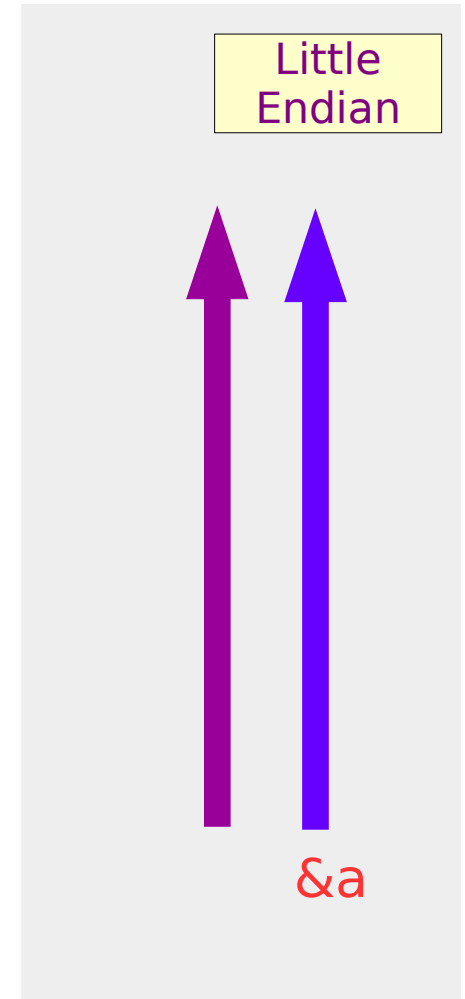
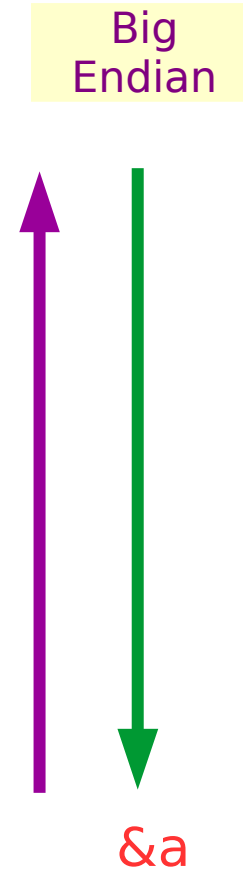
<https://stackoverflow.com/questions/15620673/which-bit-is-the-address-of-an-integer>

Increasing address, Increasing byte weight

downward, increasing address



upward, increasing address



<https://stackoverflow.com/questions/15620673/which-bit-is-the-address-of-an-integer>

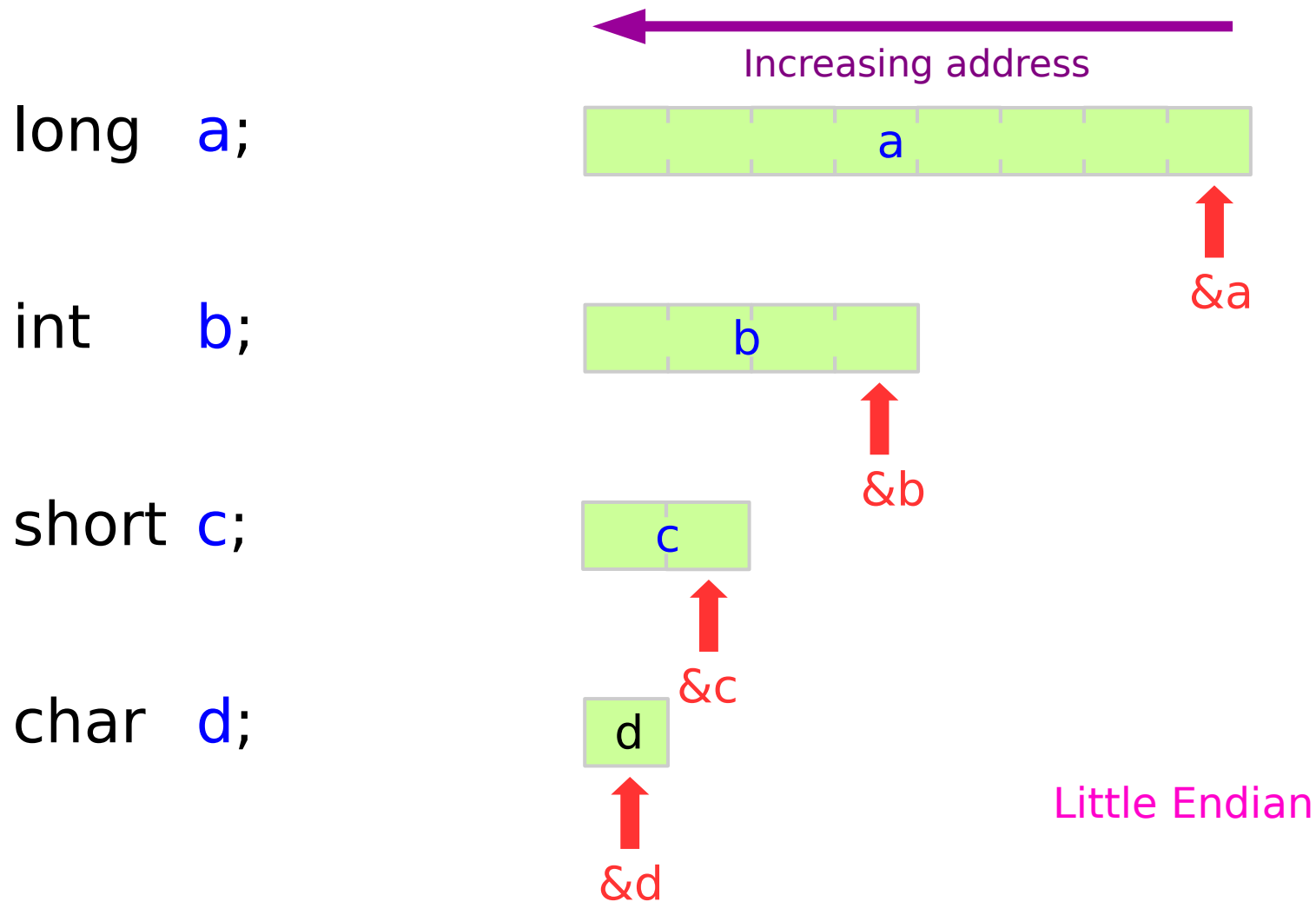
Little / Big Endian Processors

Processor	Endianness
Motorola 68000	Big Endian
PowerPC (PPC)	Big Endian
Sun Sparc	Big Endian
IBM S/390	Big Endian
Intel x86 (32 bit)	Little Endian
Intel x86_64 (64 bit)	Little Endian
Dec VAX	Little Endian
Alpha	(Big/Little) Endian
ARM	(Big/Little) Endian
IA-64 (64 bit)	(Big/Little) Endian
MIPS	(Big/Little) Endian

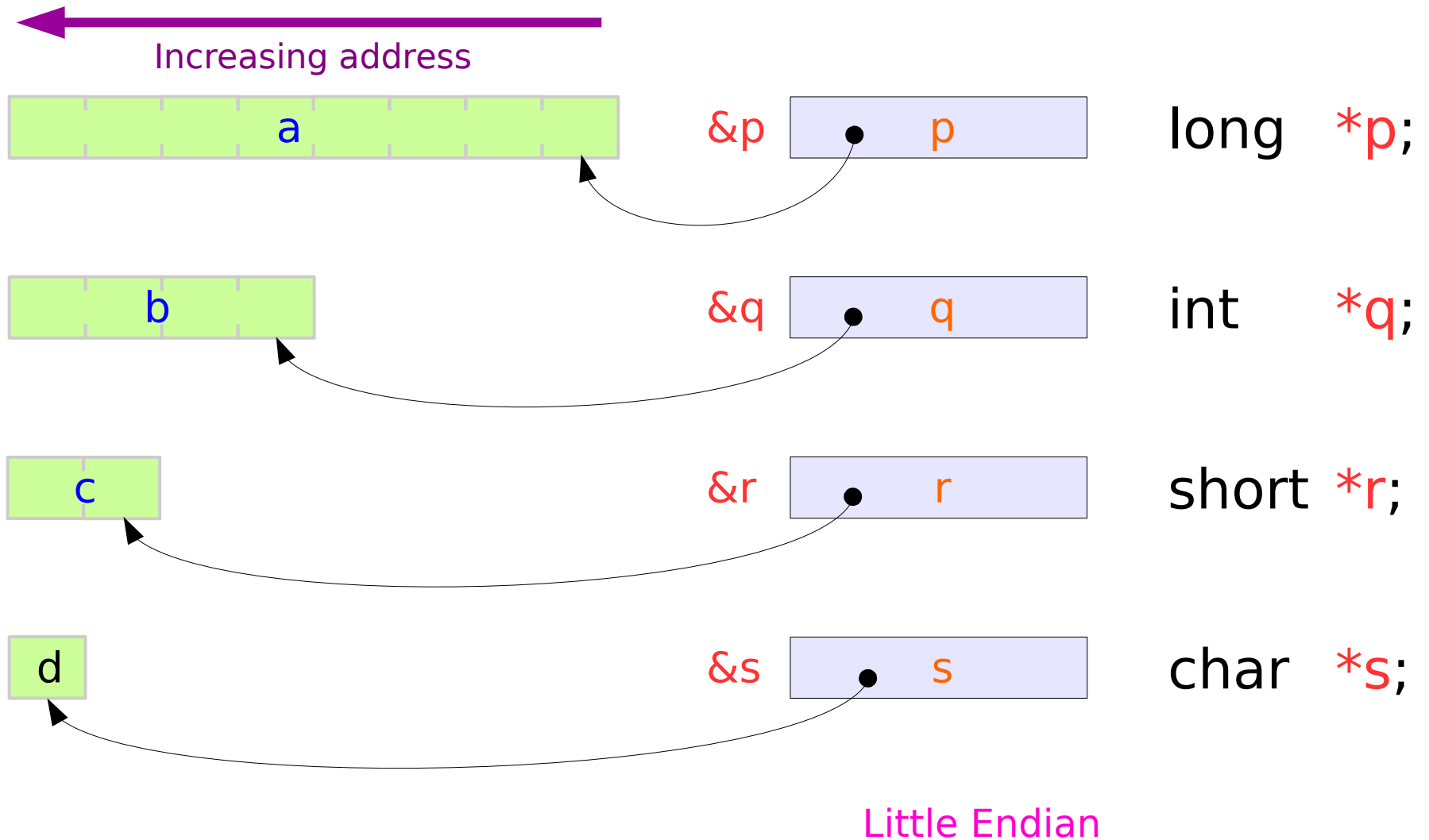
<http://www.yolinux.com/TUTORIALS/Endian-Byte-Order.html>

Pointer Types

Integer Type Variables and Their Addresses

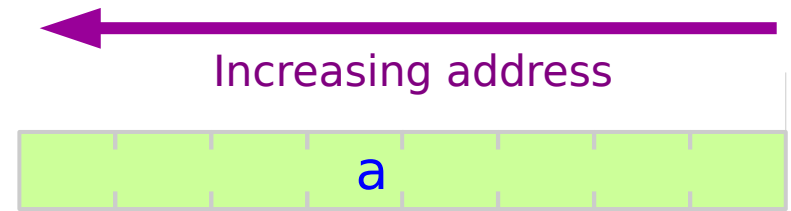


Points to the LSByte



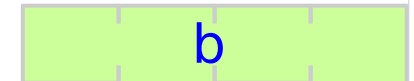
Aligning variables of different sizes

`long a;` `sizeof(long)` → 8 (bytes)



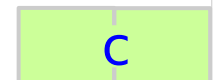
`&a` ↑

`int b;` `sizeof(int)` → 4 (bytes)



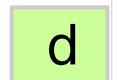
`&b` ↑

`short c;` `sizeof(short)` → 2 (bytes)



`&c` ↑

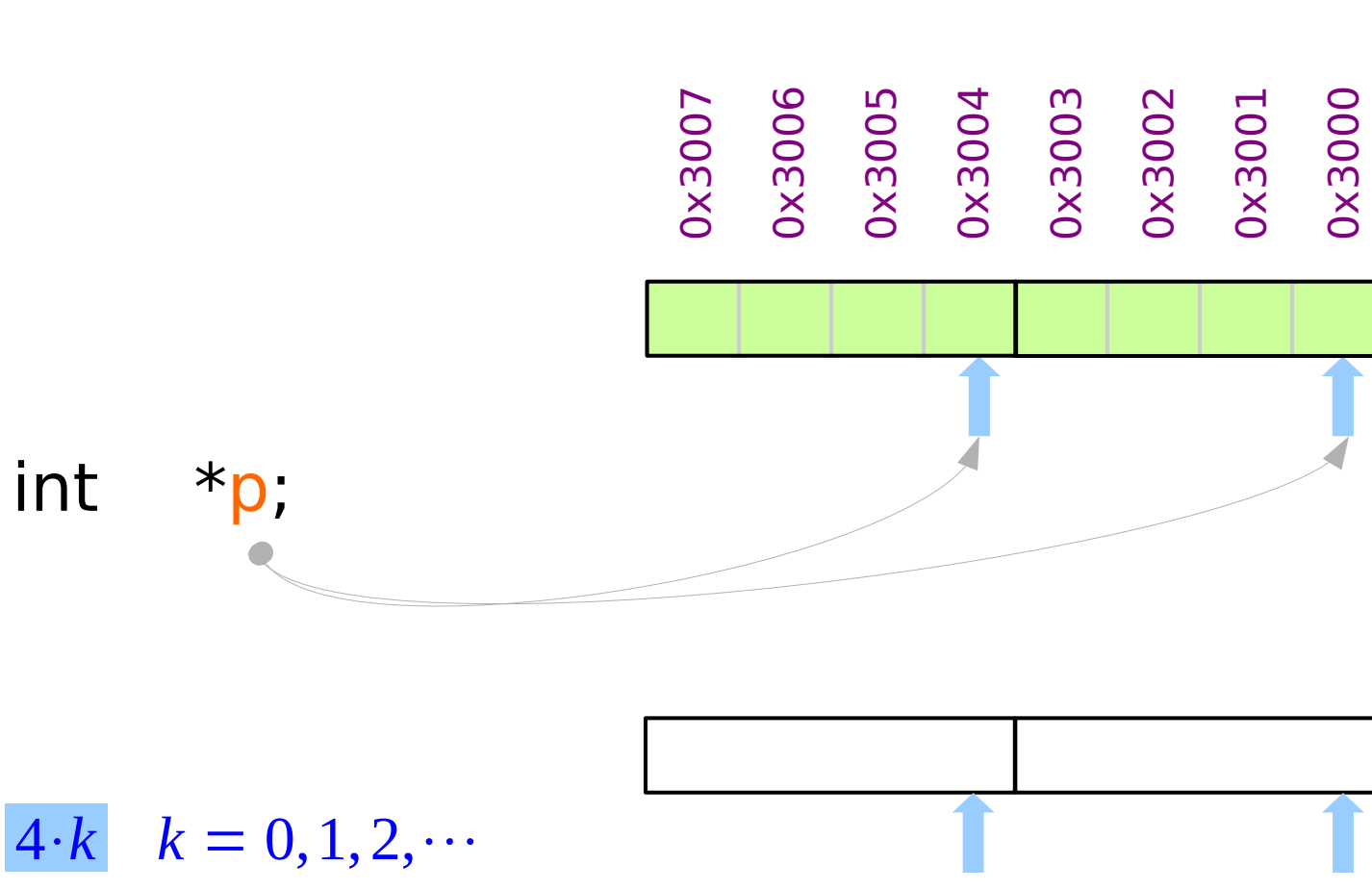
`char d;` `sizeof(char)` → 1 (bytes)



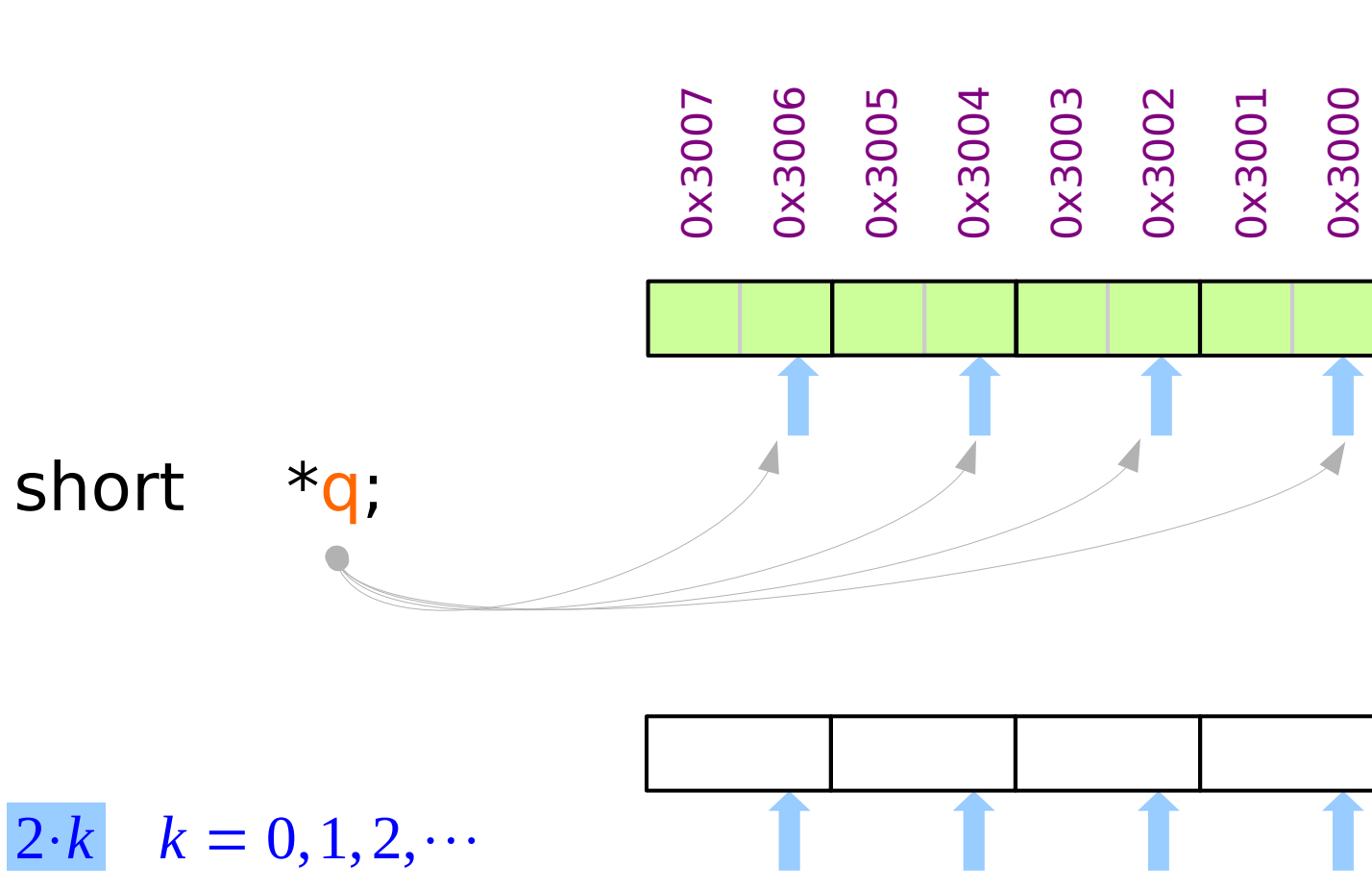
`&d` ↑

Memory Alignment
in the Little Endian

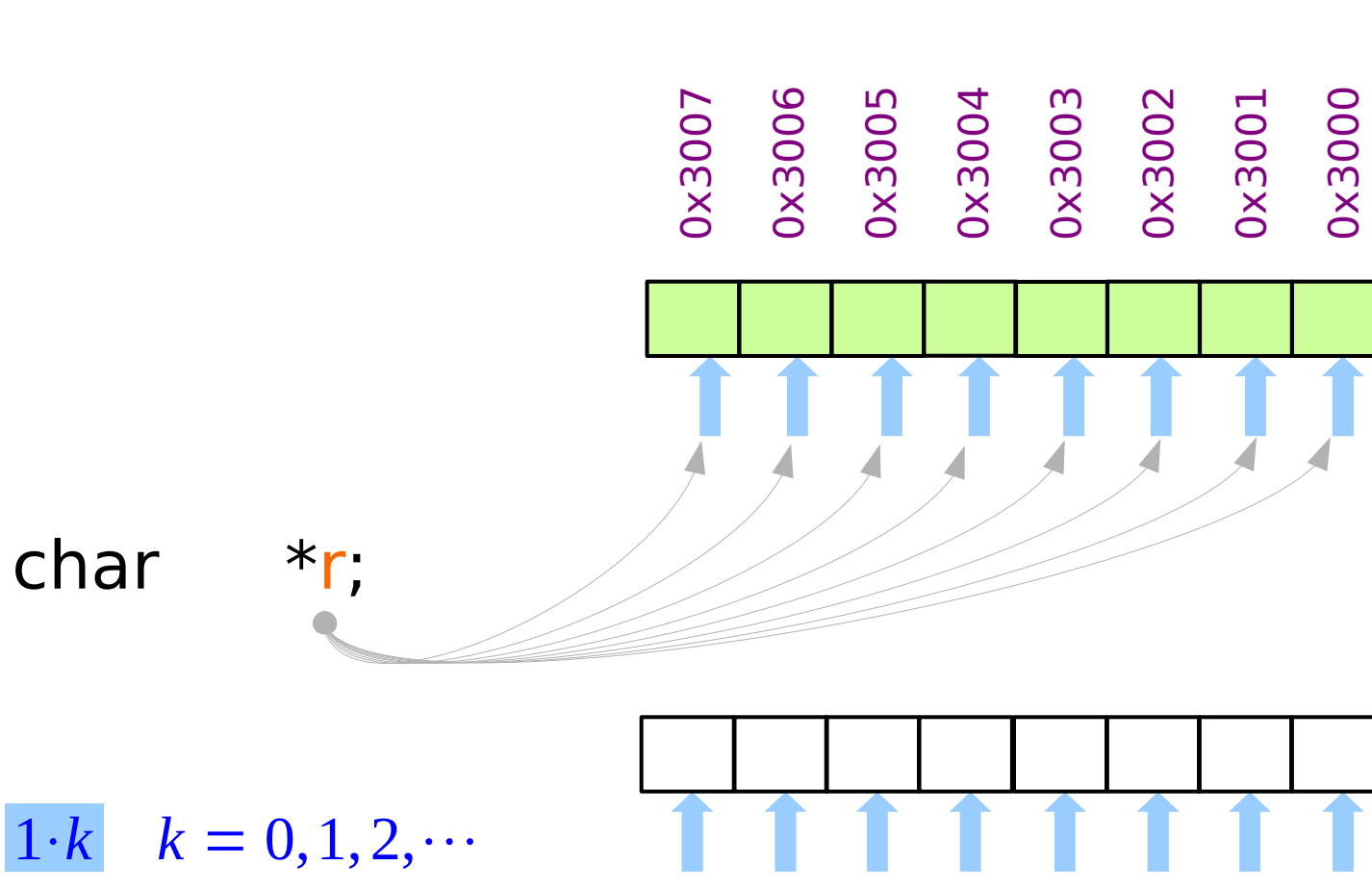
Possible addresses for **int** values



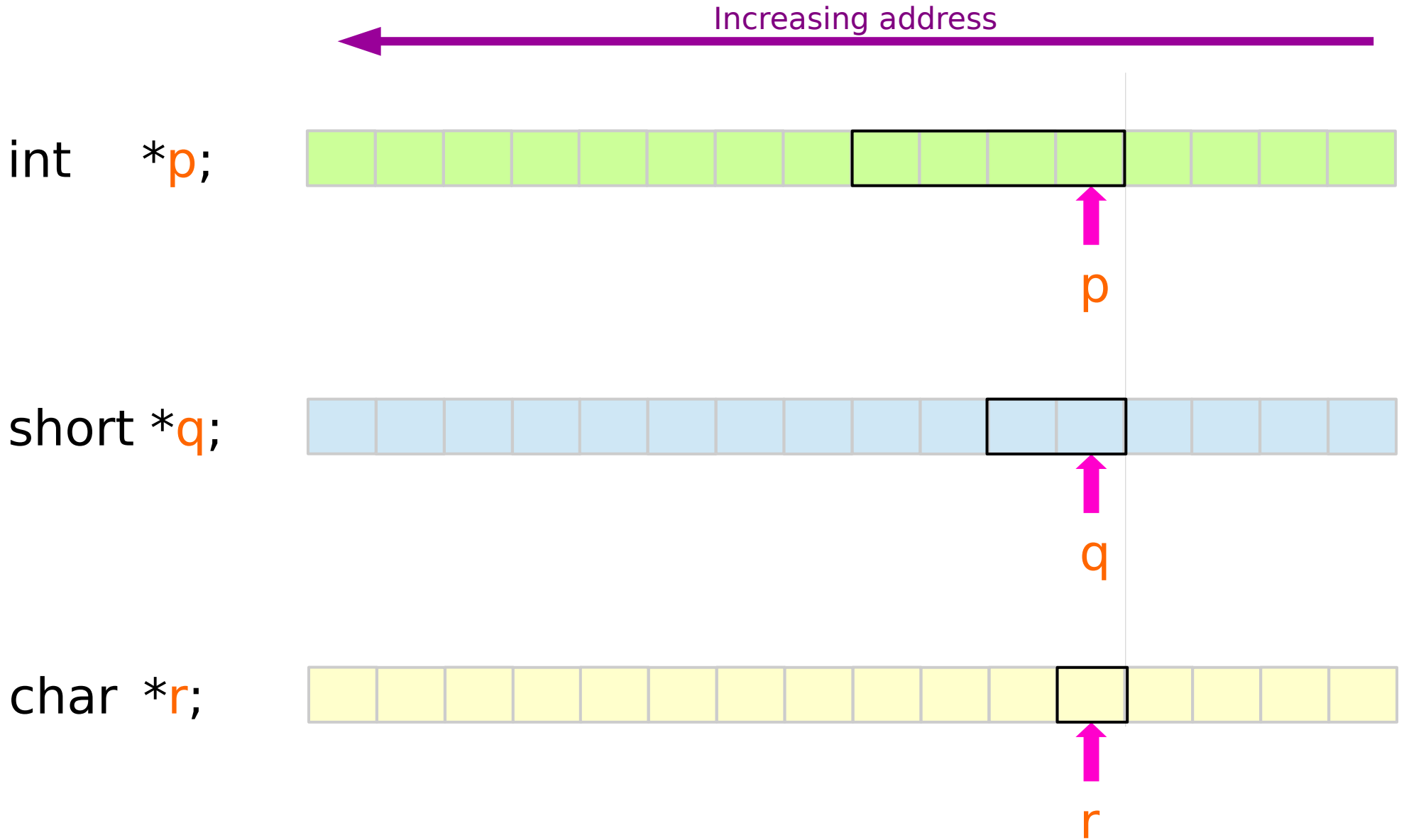
Possible addresses for **short** values



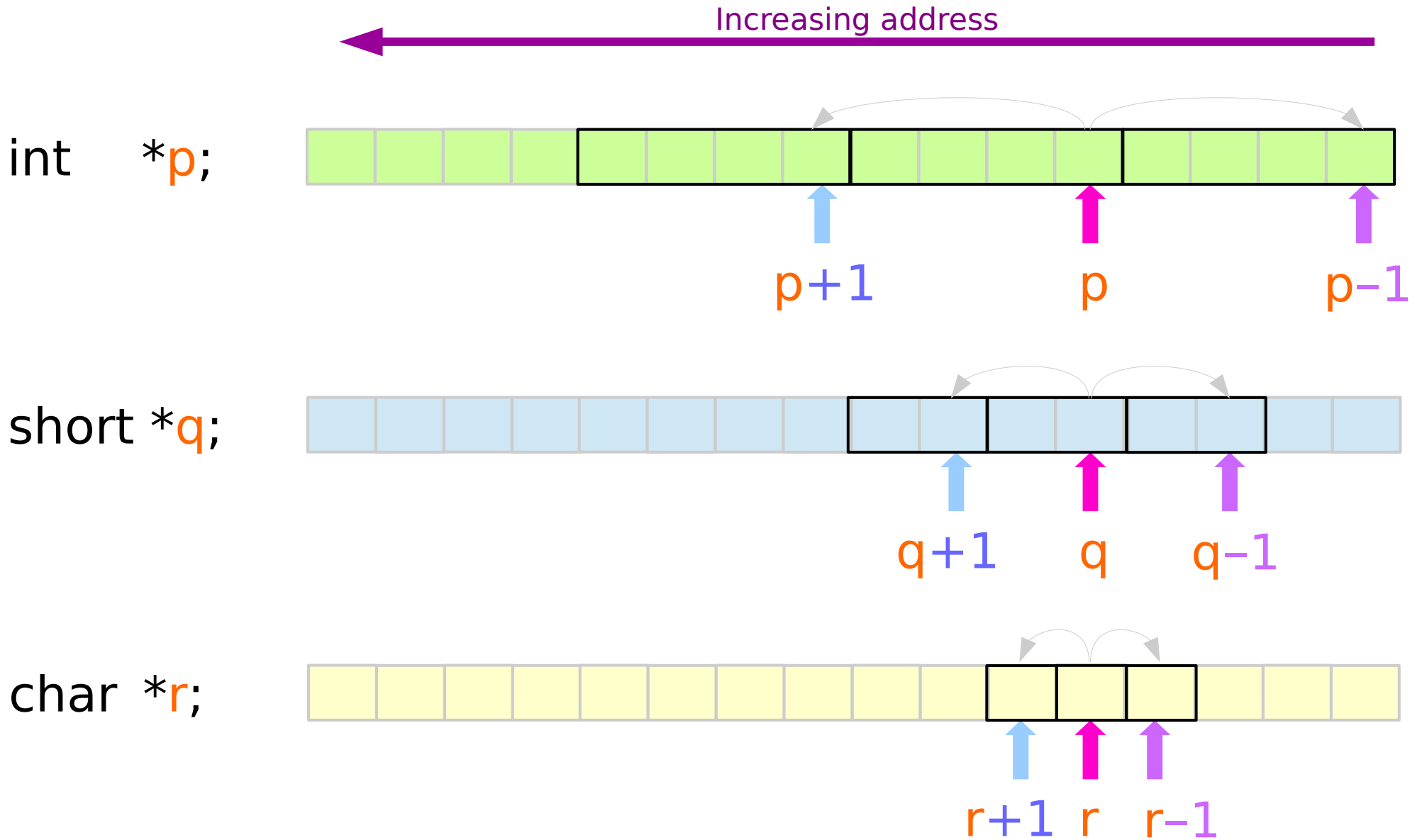
Possible addresses for **char** values



Data size at the pointed address



Incrementing / decrementing pointers



Memory Alignment (1) - allocation of variables

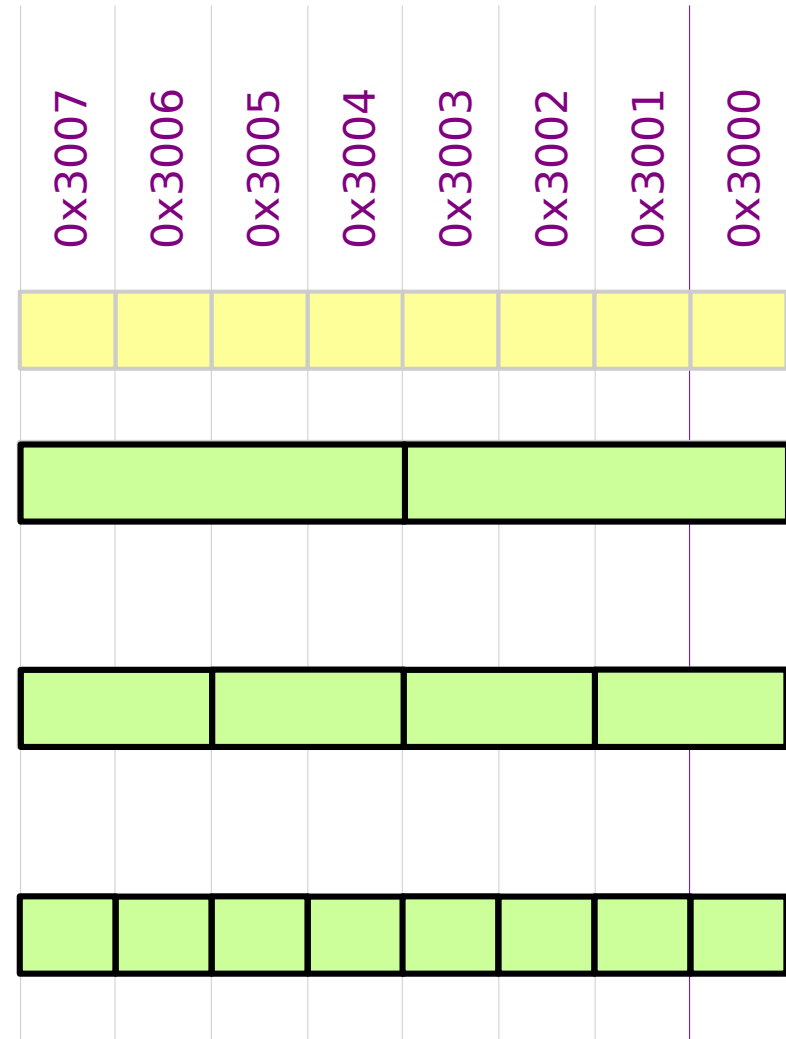
enforced by compilers

efficient memory access

int **a**;

short **b**;

char **c**;



Memory Alignment (2) - integer multiple addresses

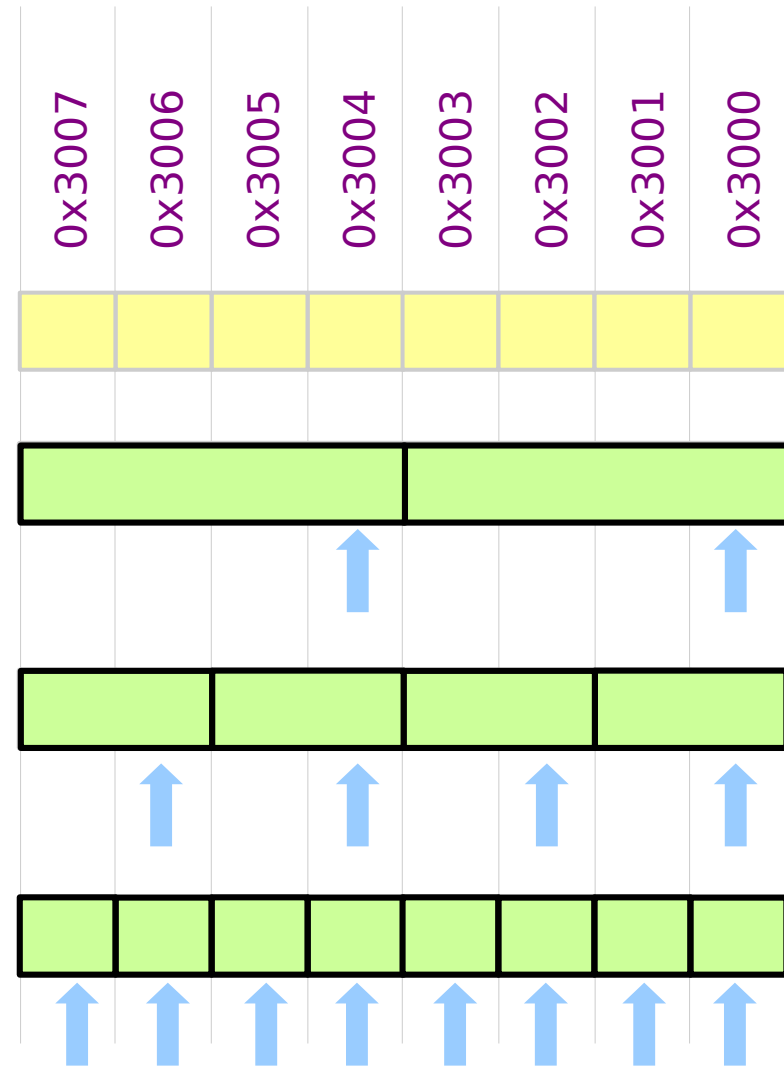
Memory Alignment:
the data **address** is a
multiple of the data **size**.

$$k = 0, 1, 2, \dots$$

$$\text{integer addresses} = 4 \cdot k$$

$$\text{short addresses} = 2 \cdot k$$

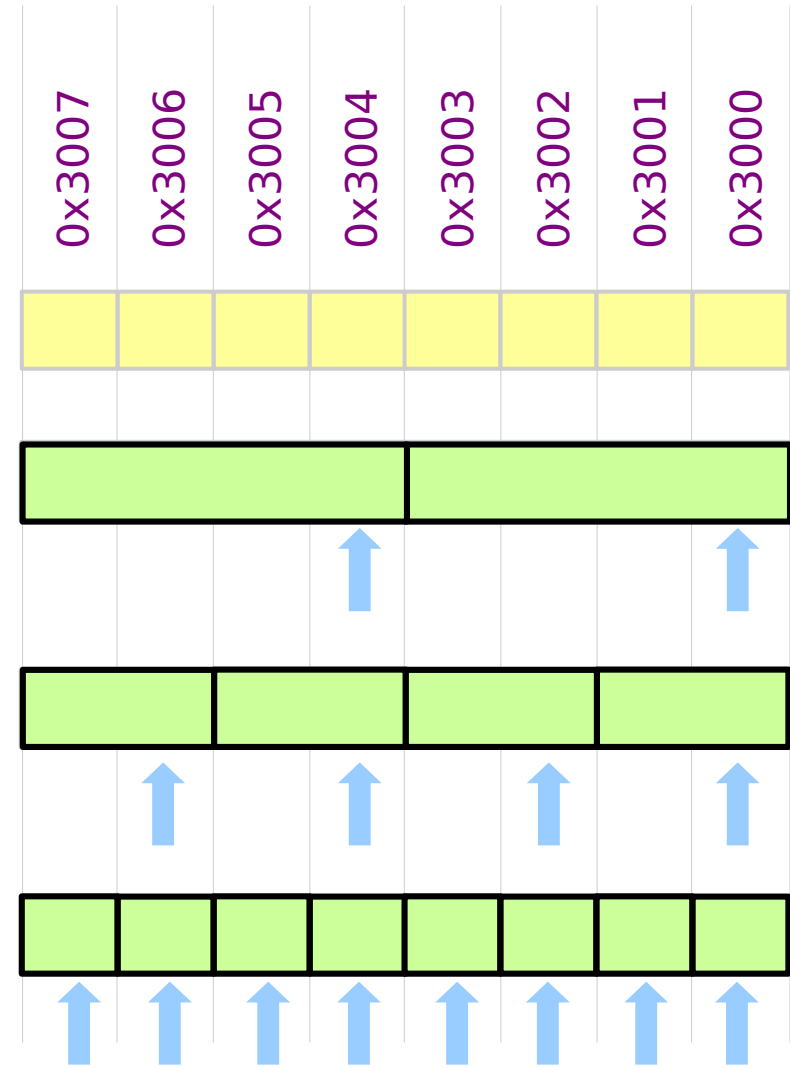
$$\text{character addresses} = 1 \cdot k$$



Memory Alignment (3) - pointed addresses

int *p;
short *q;
char *r;

$4 \cdot k$
 $2 \cdot k$
 $1 \cdot k$



Memory Alignment (4) - non-pointed addresses

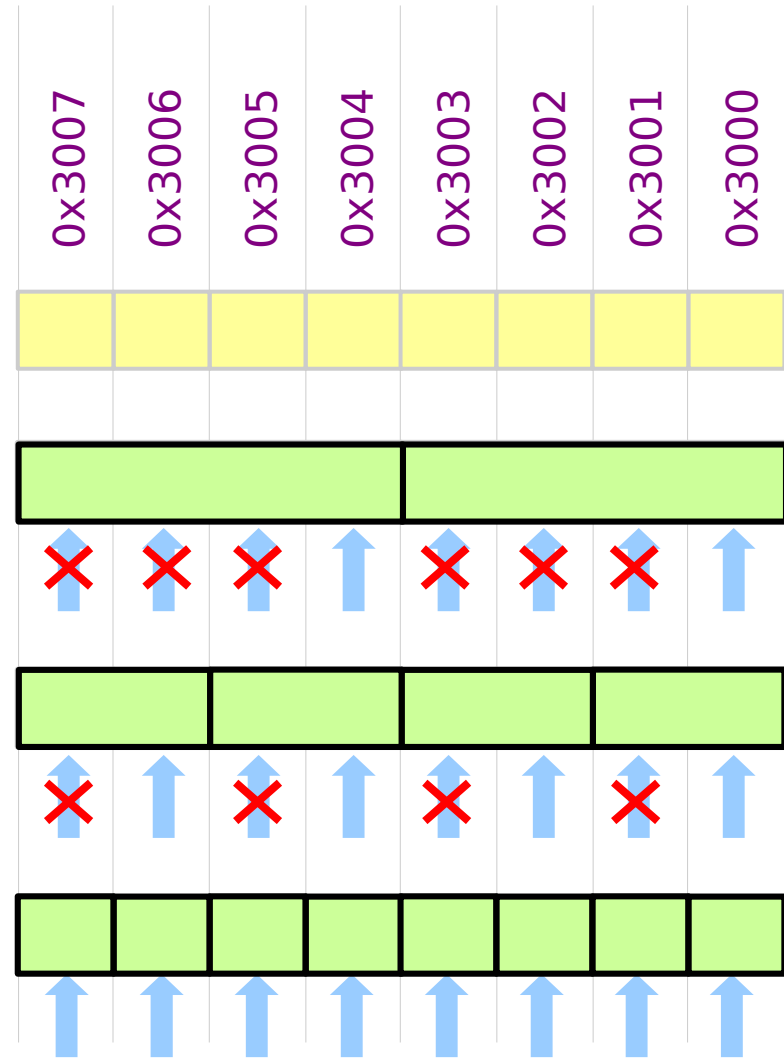
int *p;

$$4 \cdot k + 1, 2, 3$$

short *q;

$$2 \cdot k + 1$$

char *r;



Memory Alignment (5) - broken alignment

Memory access is still possible
but it takes **longer** time to access

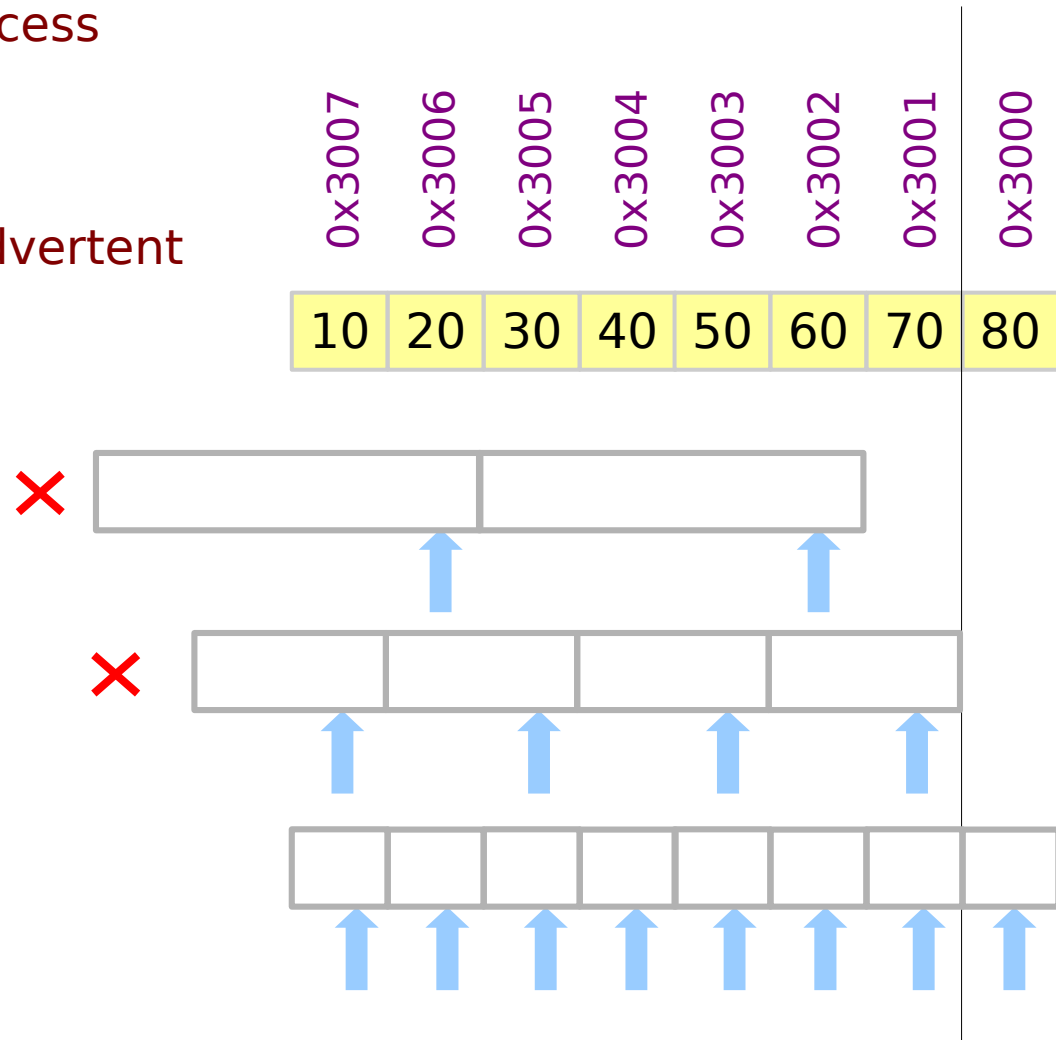
(Low Efficiency)

This can happen by using inadvertent
pointer type casting

`int *p;`

`short *q;`

`char *r;`



Pointer Type Cast

Changing the associated data type of an address

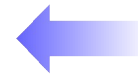
`long a;`

address of a `long` value



`int * p;`

address of an `int` value



`short * q;`

address of a `short` value



`char * r;`

address of a `char` value



Pointer Type Casting

long a;

address of a long value



&a



int * p;

address of an int value



p = (int *) &a

short * q;

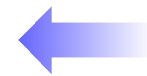
address of a short value



q = (short *) &a

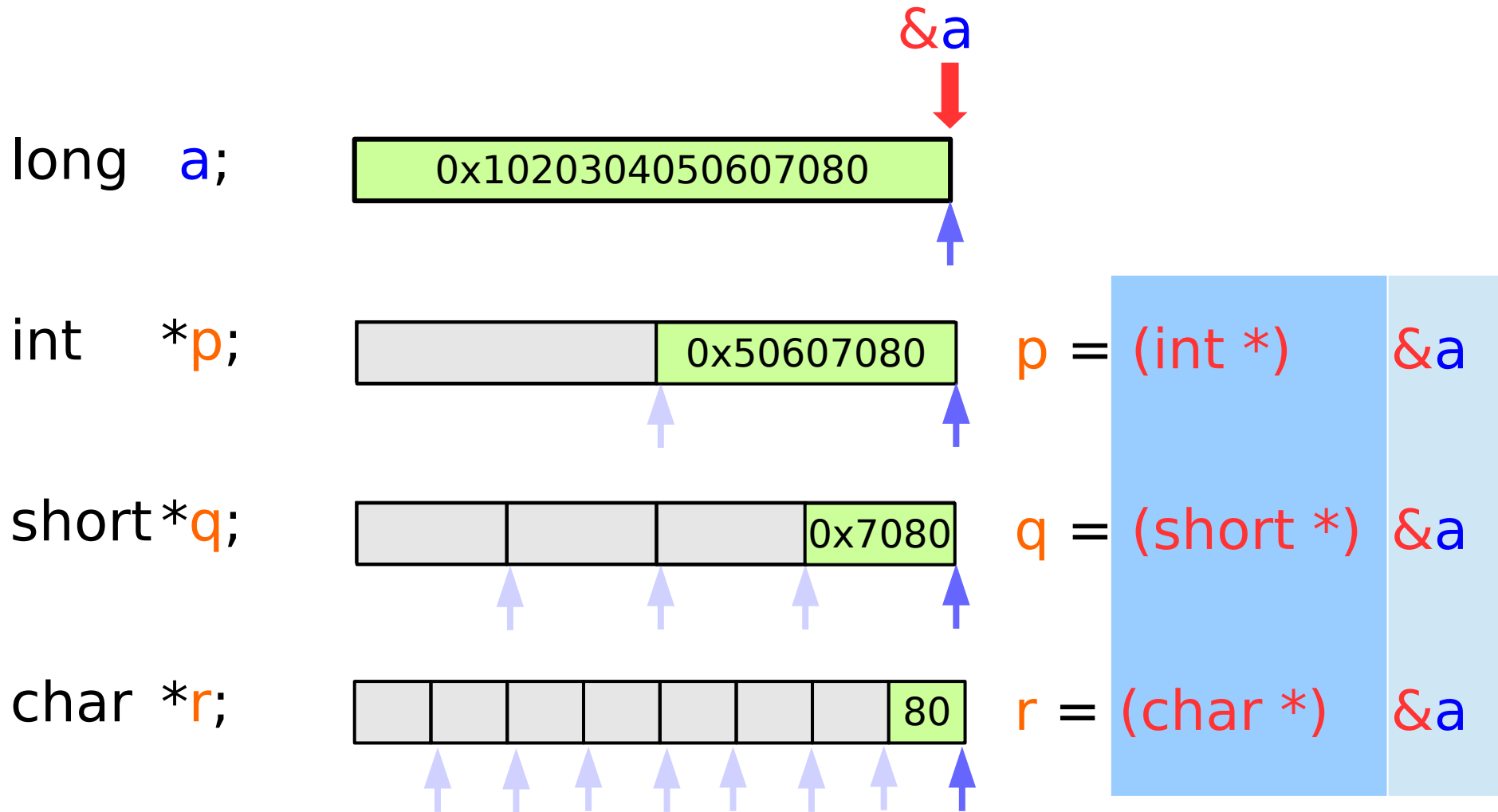
char * r;

address of a char value



r = (char *) &a

Re-interpretation of memory data - case I



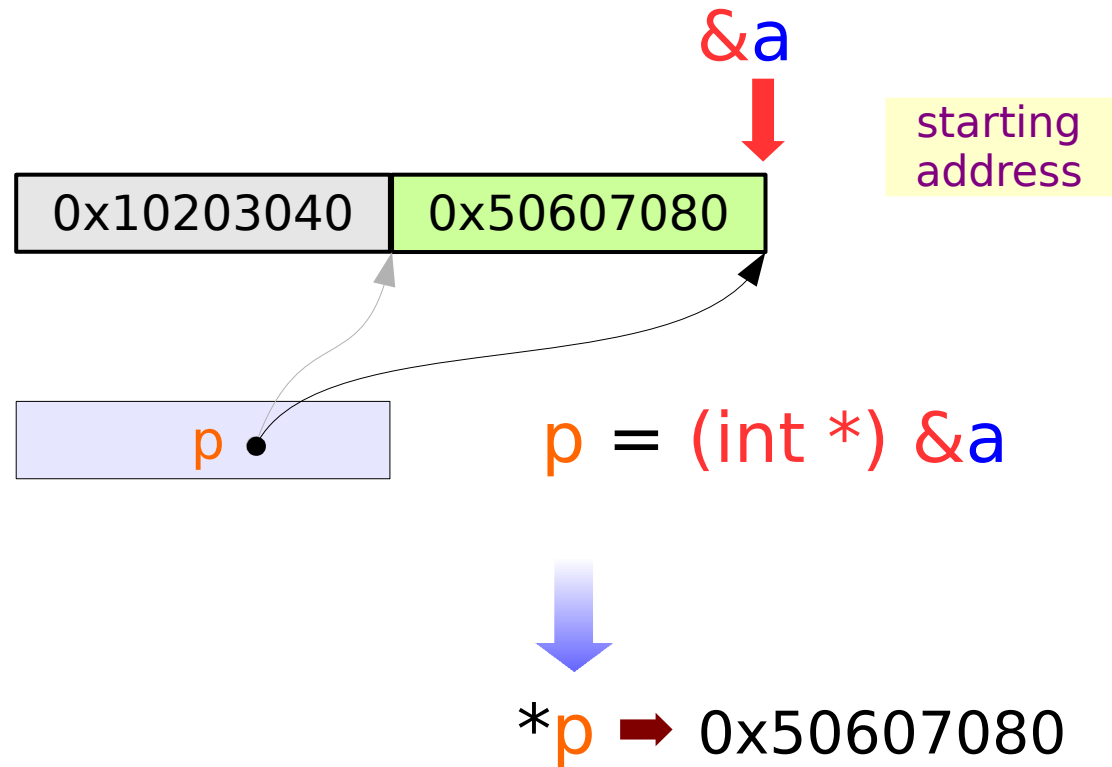
Pointer Type Cast

```
long a;
```

```
int *p;
```

```
short *q;
```

```
char *r;
```



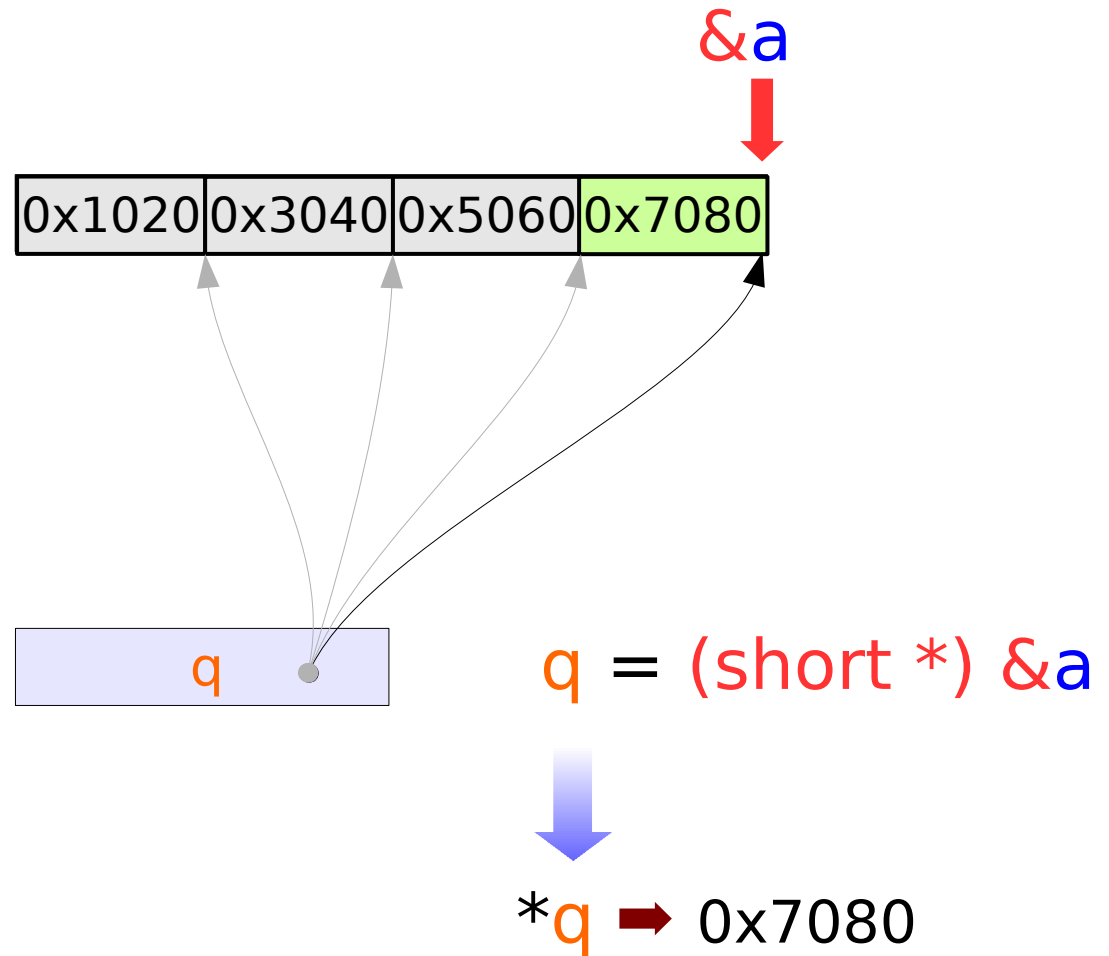
Integer Pointer Types

long a;

int *p;

short *q;

char *r;



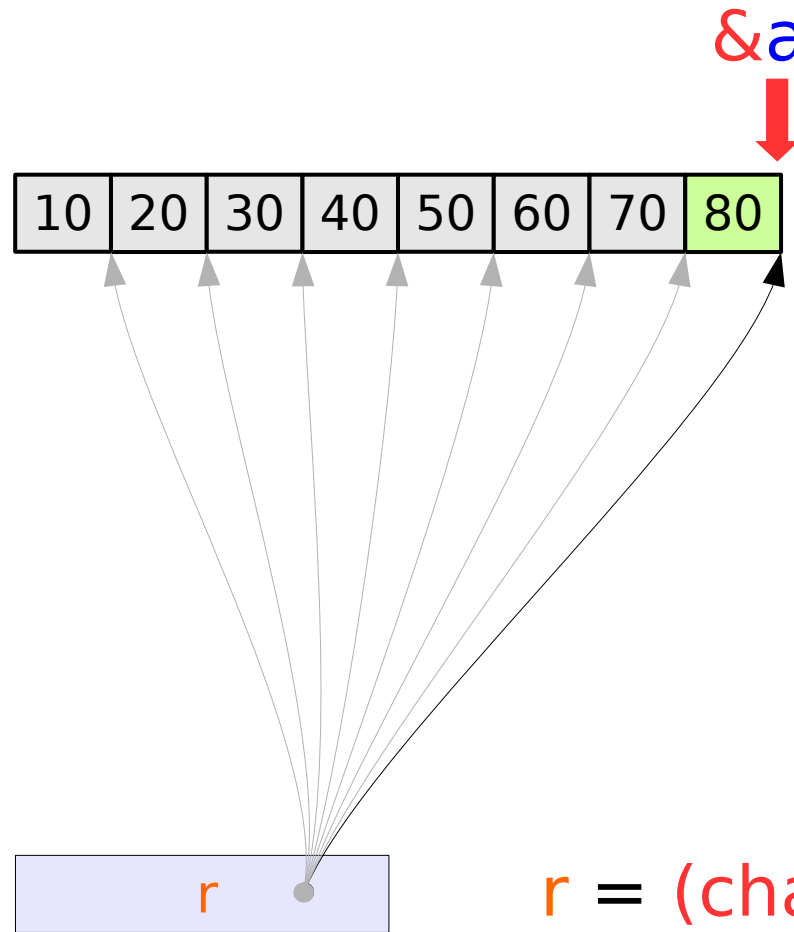
Integer Pointer Types

long **a**;

int ***p**;

short ***q**;

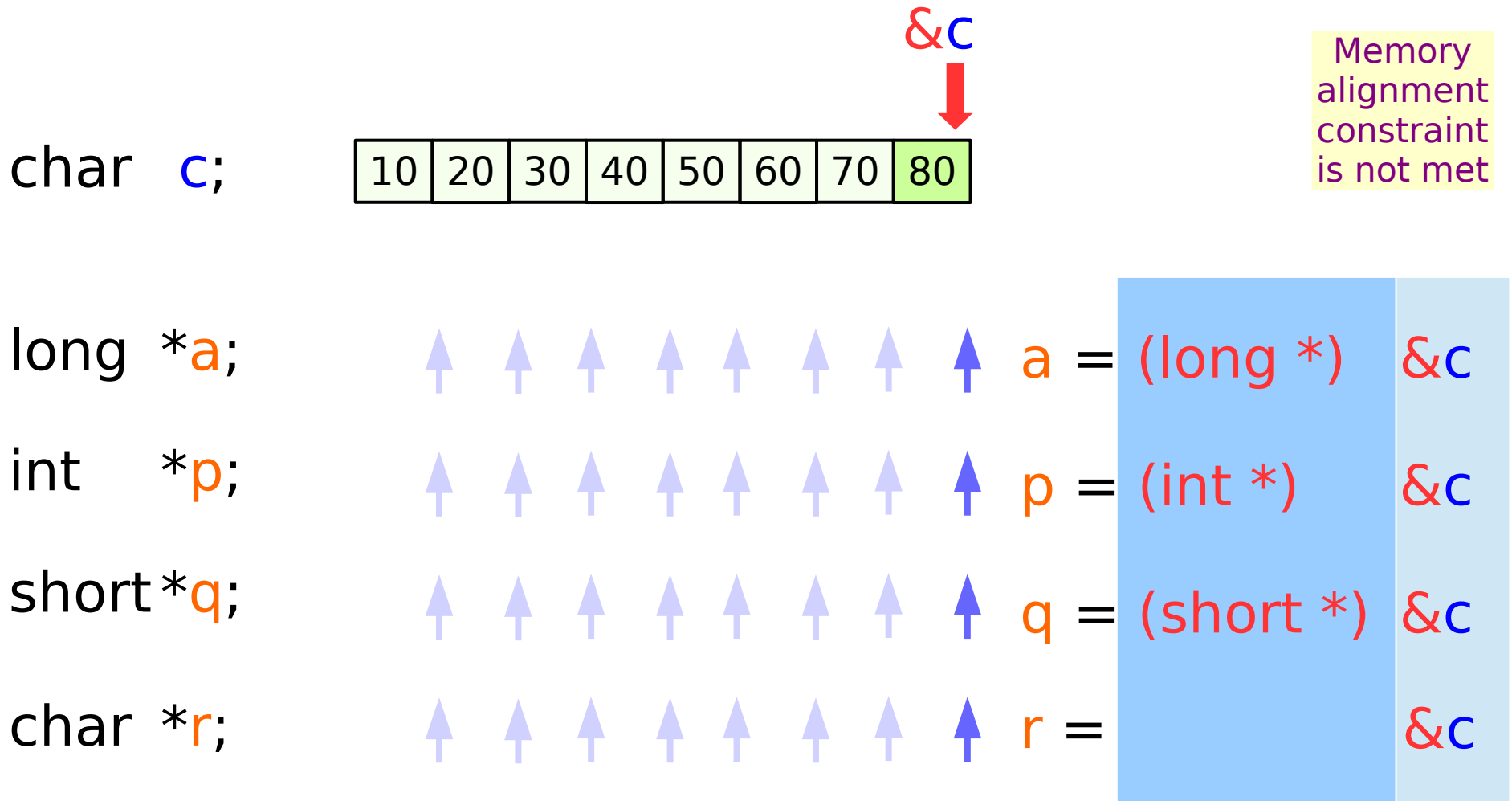
char ***r**;



r = (char *) &a

***r** → 0x80

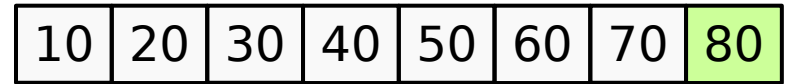
Re-interpretation of memory data – case II



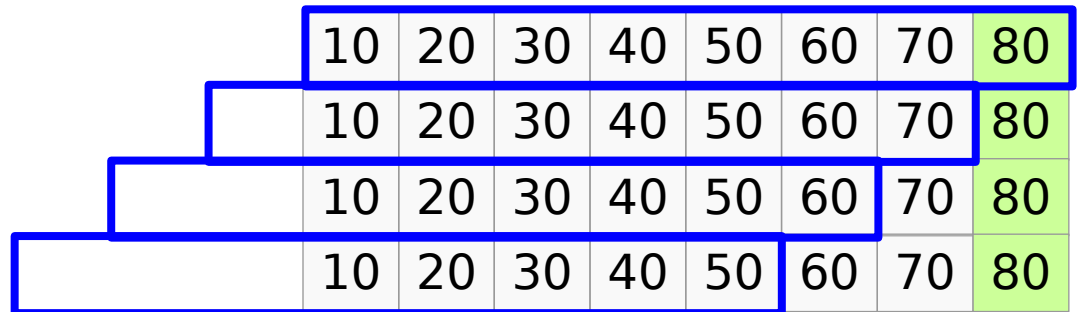
In this case, the memory alignment constraint can be broken

Pointer Type Cast

char c;

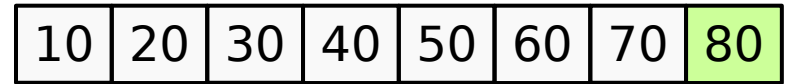


long *a;

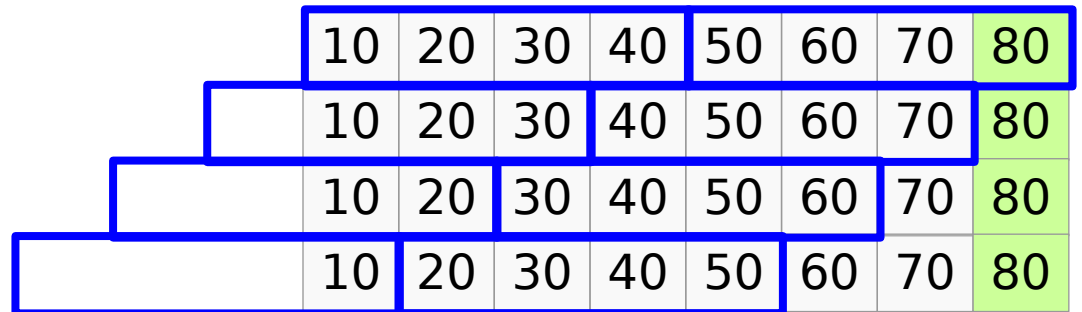


Pointer Type Cast

char c;

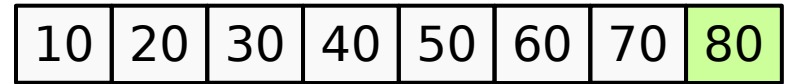


int *p;

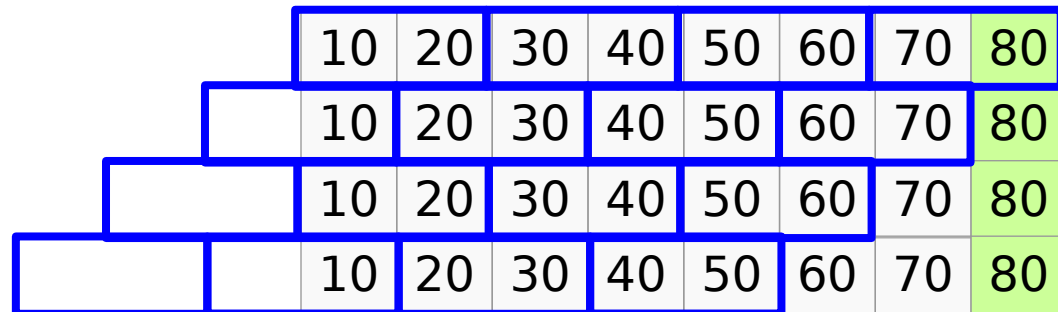


Pointer Type Cast

char c;



short*q;



const pointers

const type, const pointer type (1)

```
const int * p;
```

read only integer value

```
int * const q ;
```

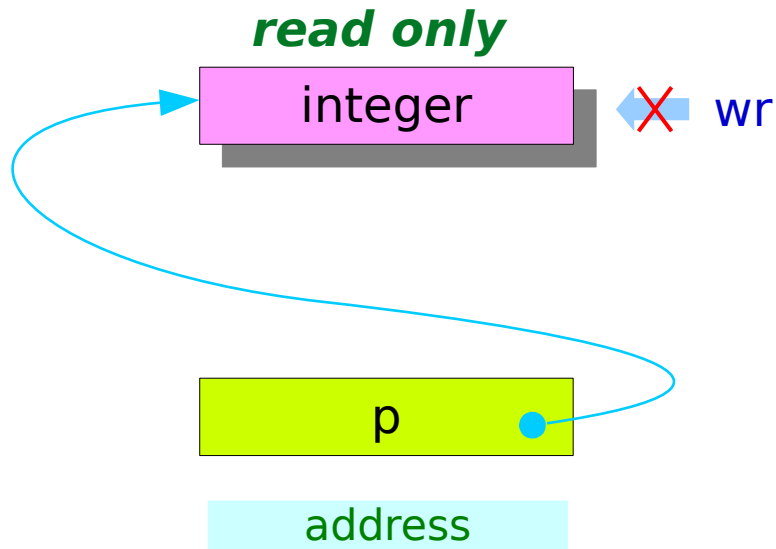
read only integer pointer

```
const int * const r ;
```

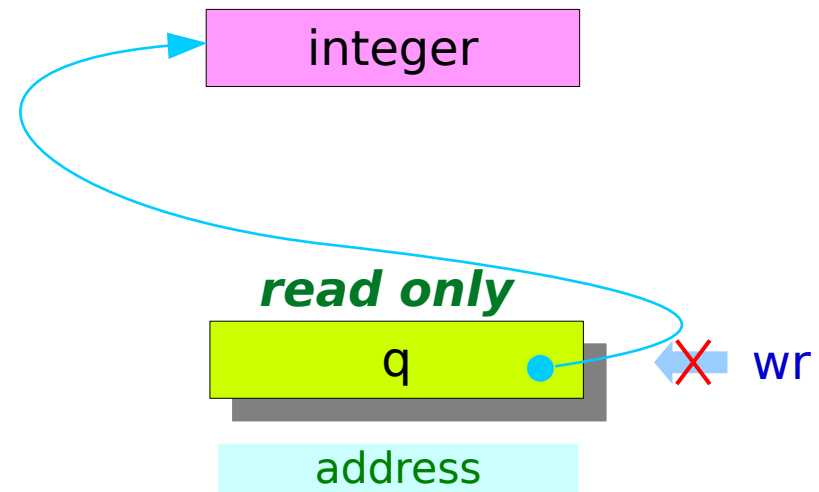
*read only integer value
read only integer pointer*

const type, const pointer type (2)

`const int * p;`

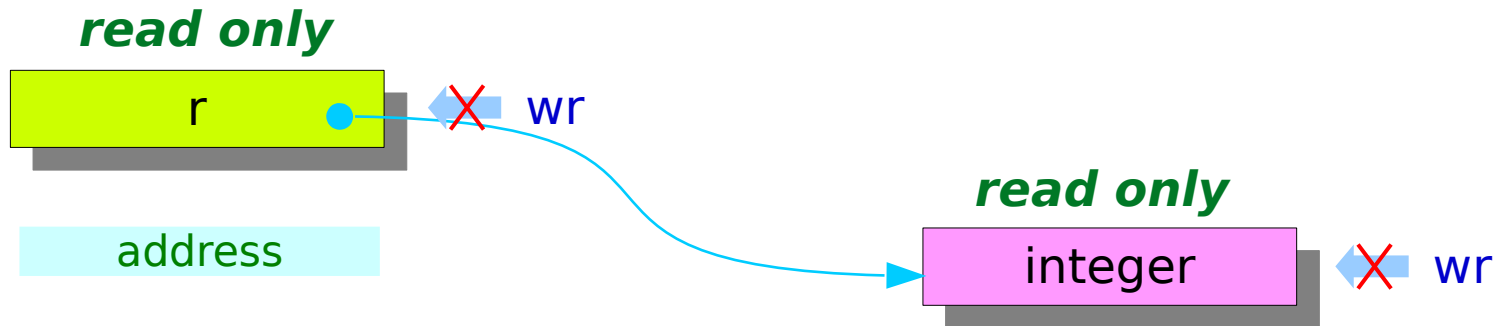


`int * const q;`

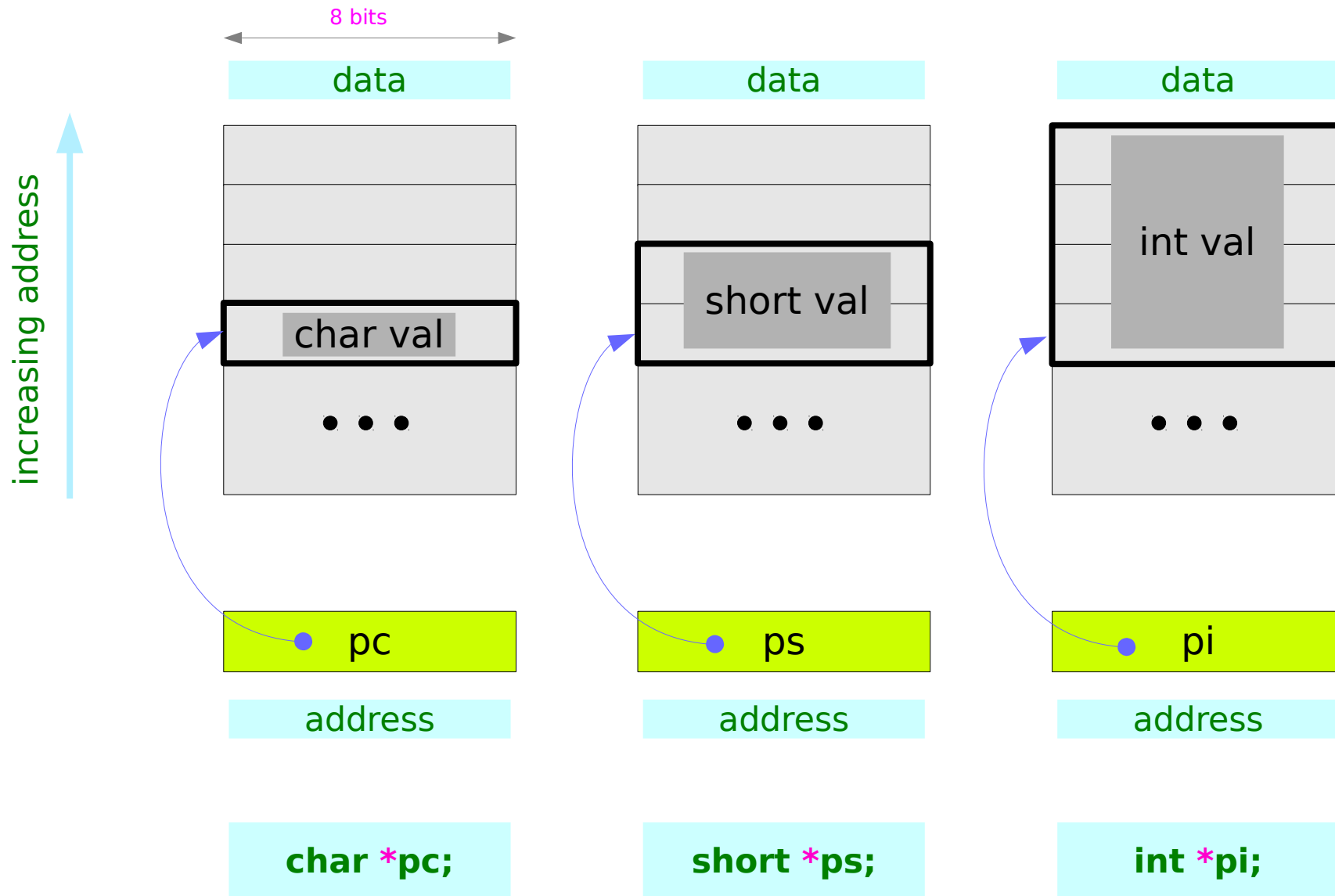


const type, const pointer type (3)

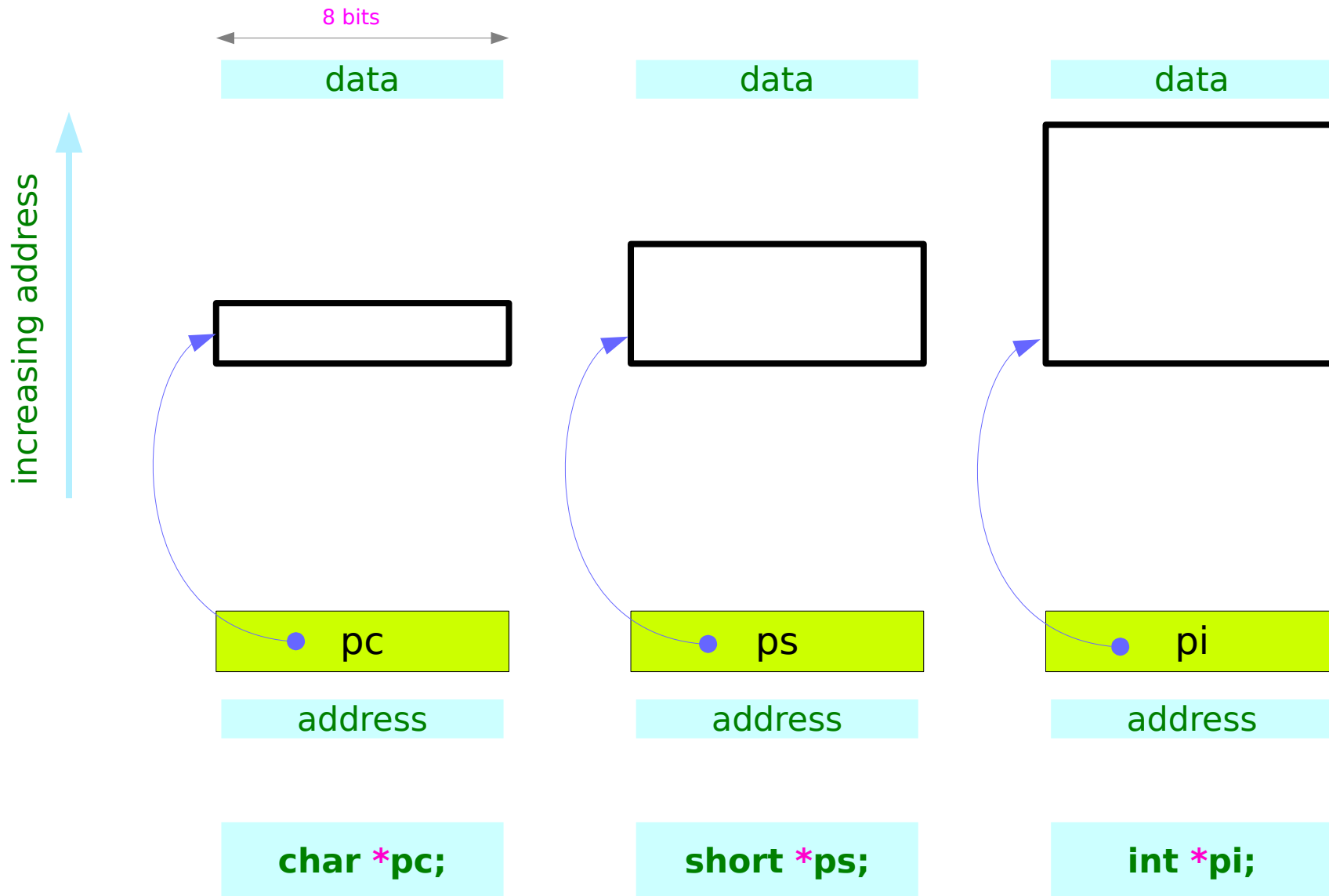
```
const int * const r ;
```



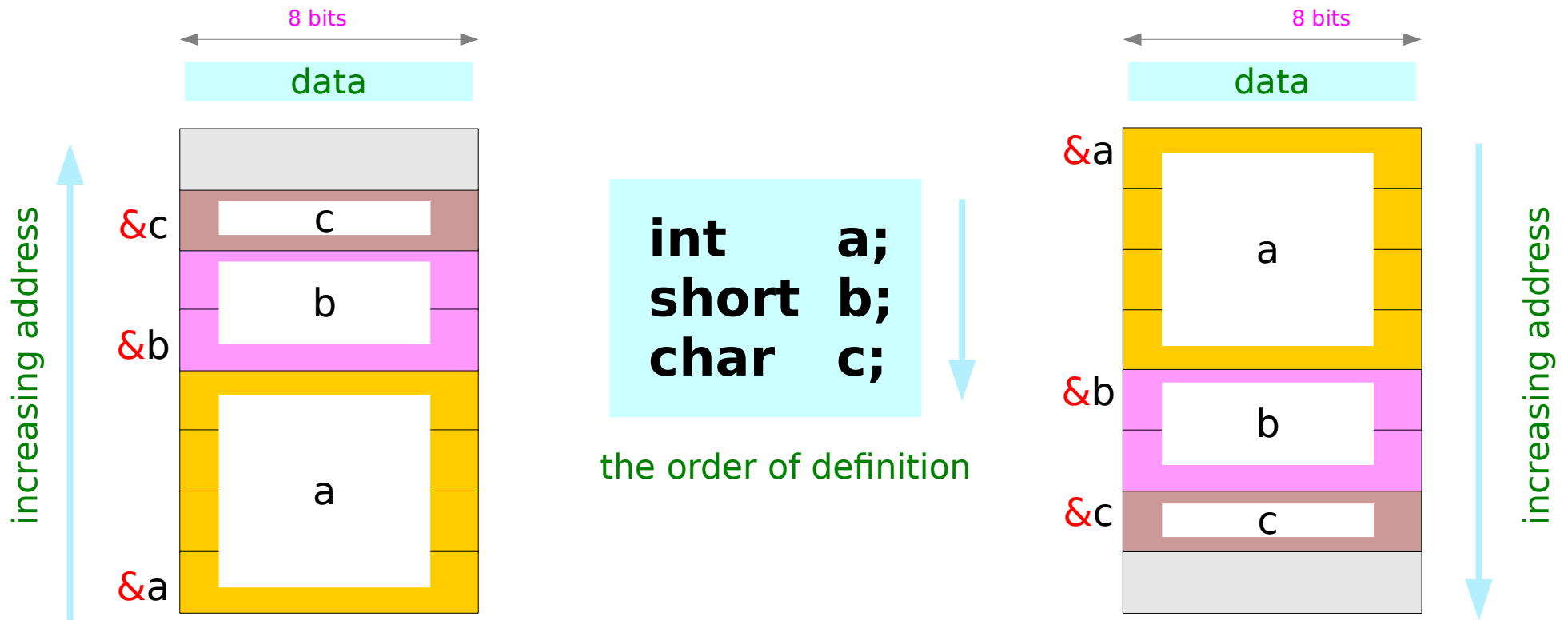
Pointer Types and Associated Data



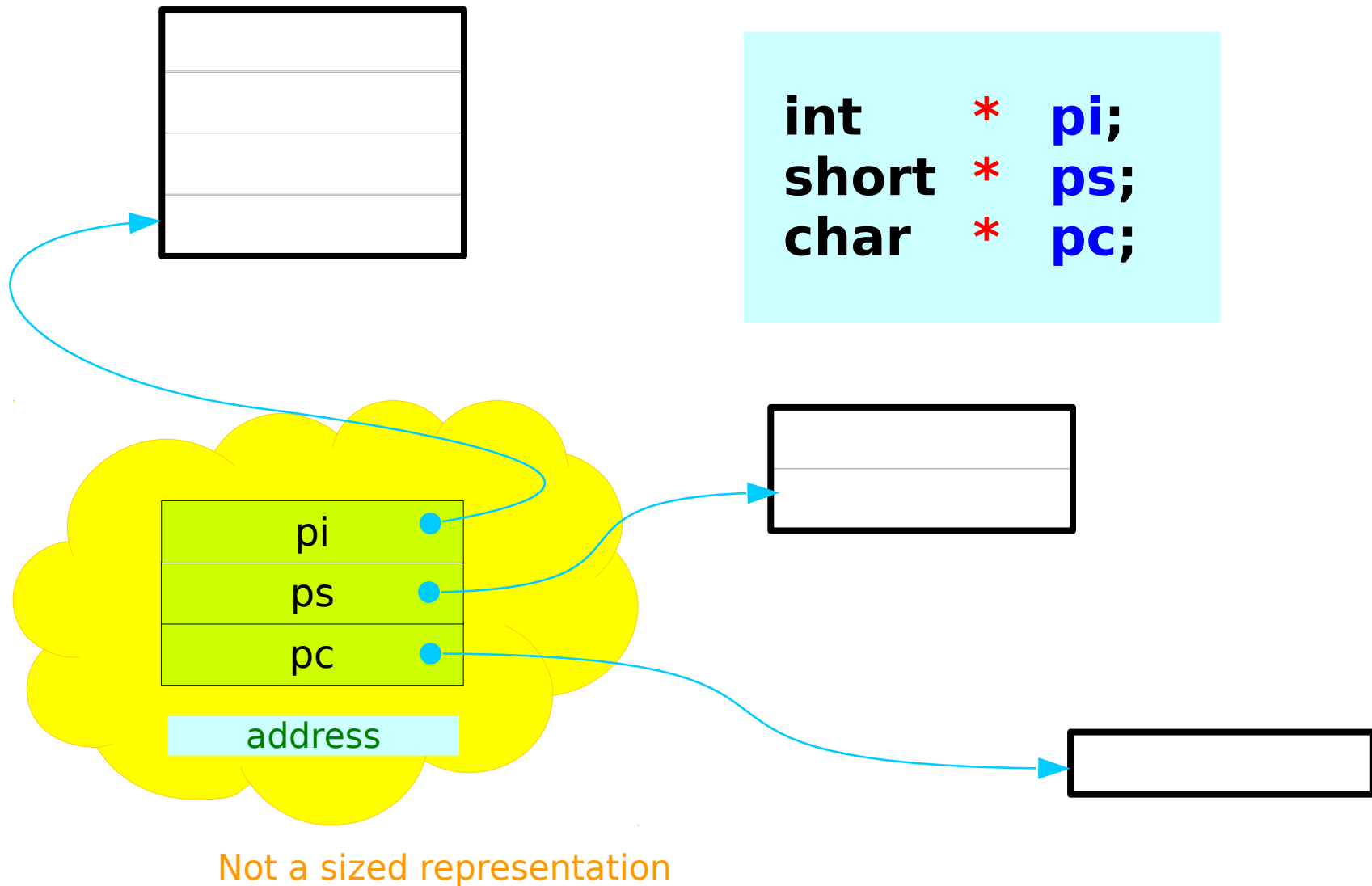
Pointer Types



Little Endian Example

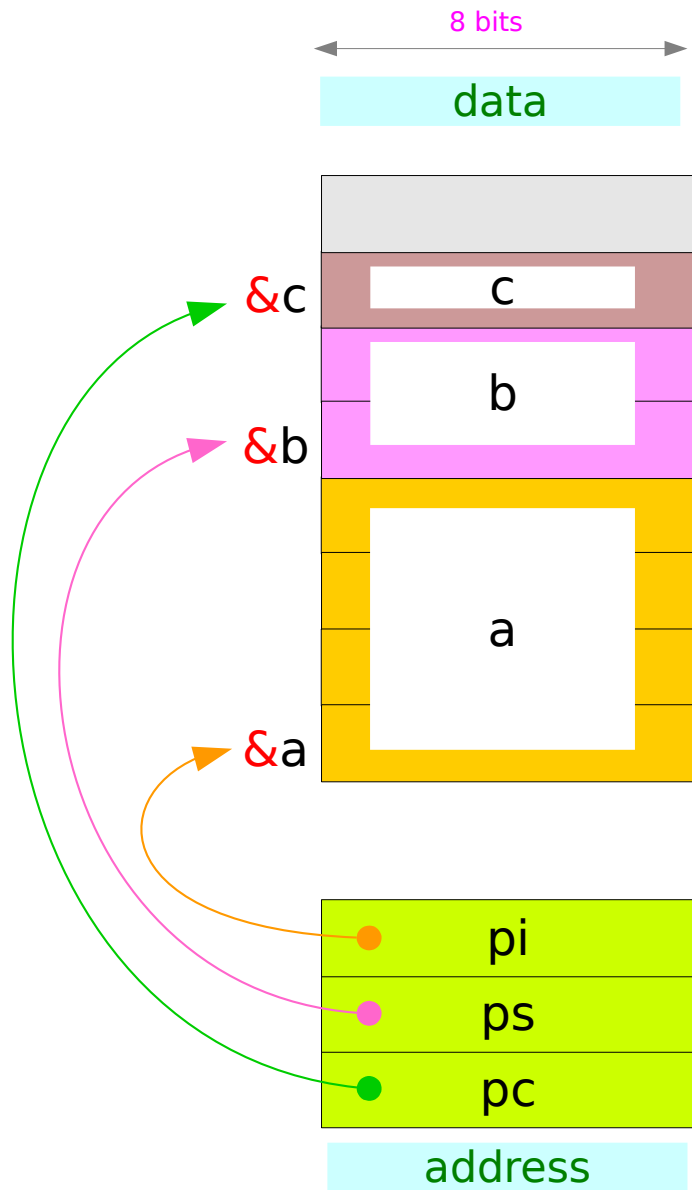


int *, short *, char * type variables



Not a sized representation

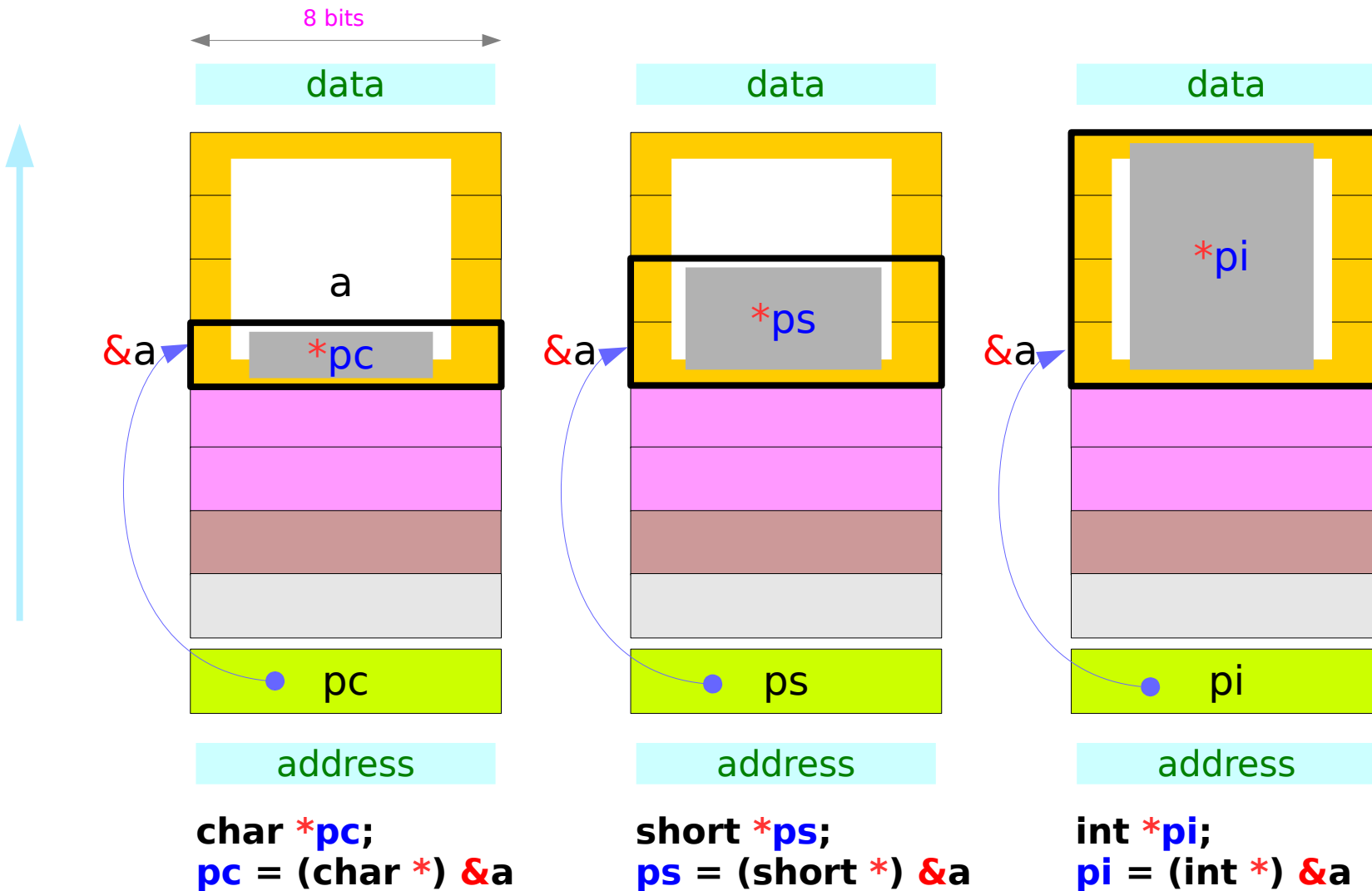
Pointer Variable Assignment



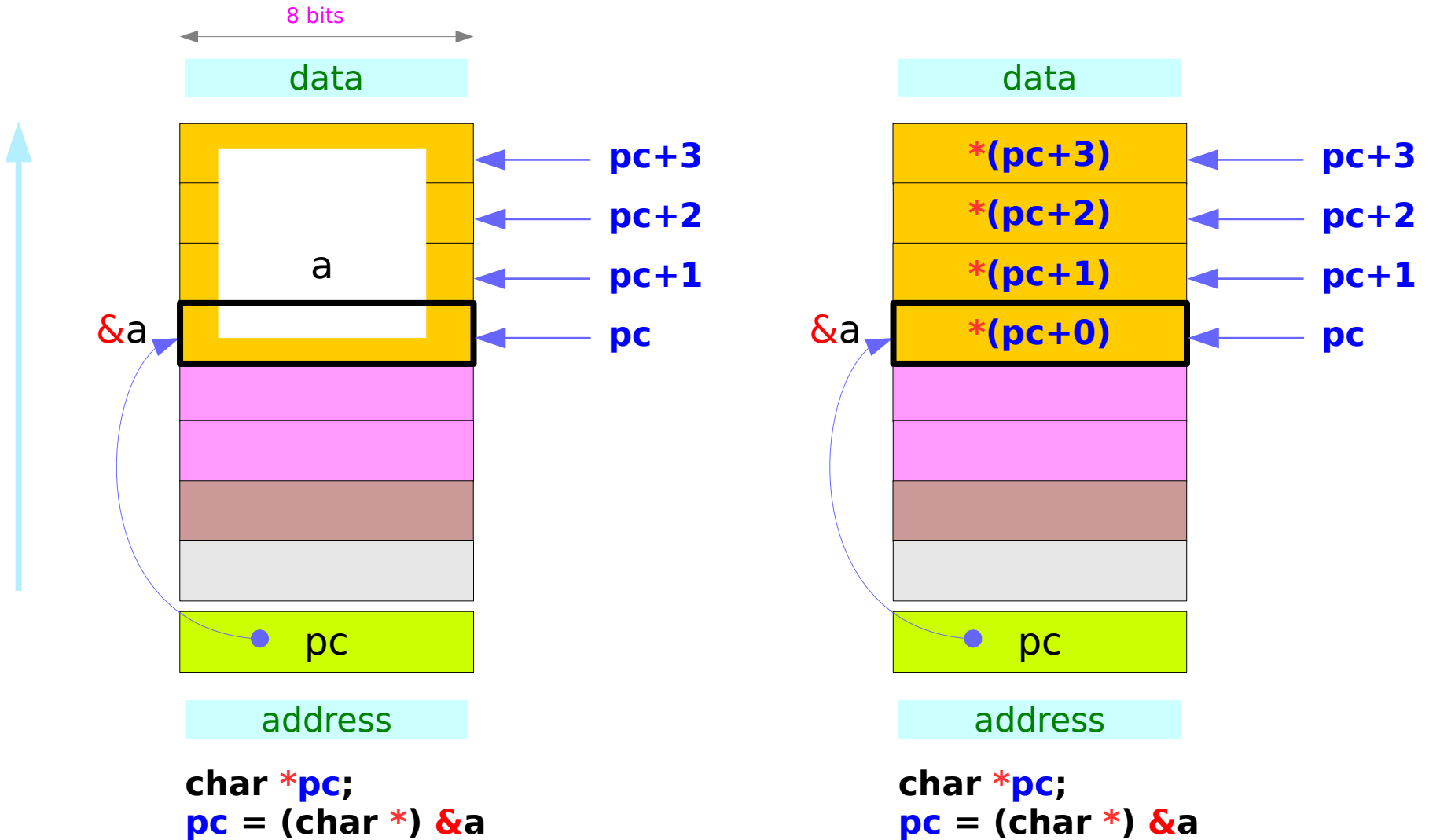
```
char * pc;  
short * ps;  
int * pi;  
  
int a;  
short b;  
char c;
```

```
pi = &a;  
ps = &b;  
pc = &c;
```

Pointer Type Casting



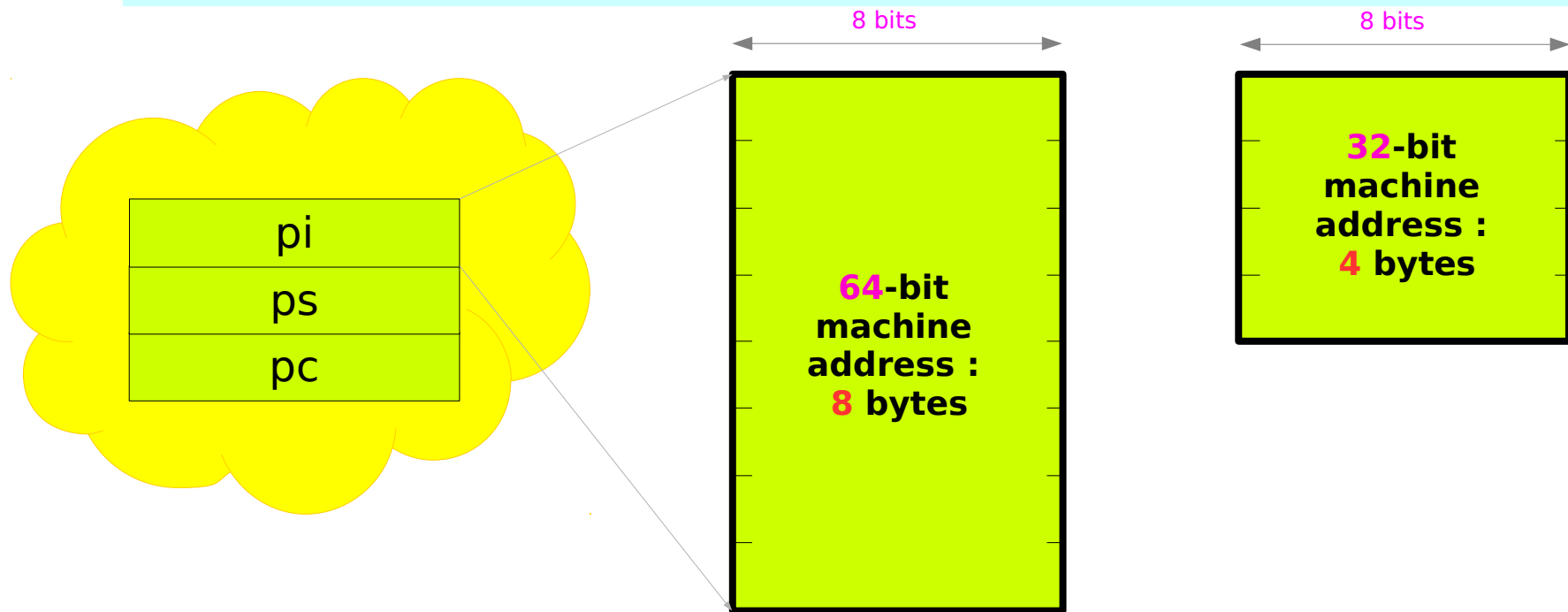
Accessing bytes of a variable



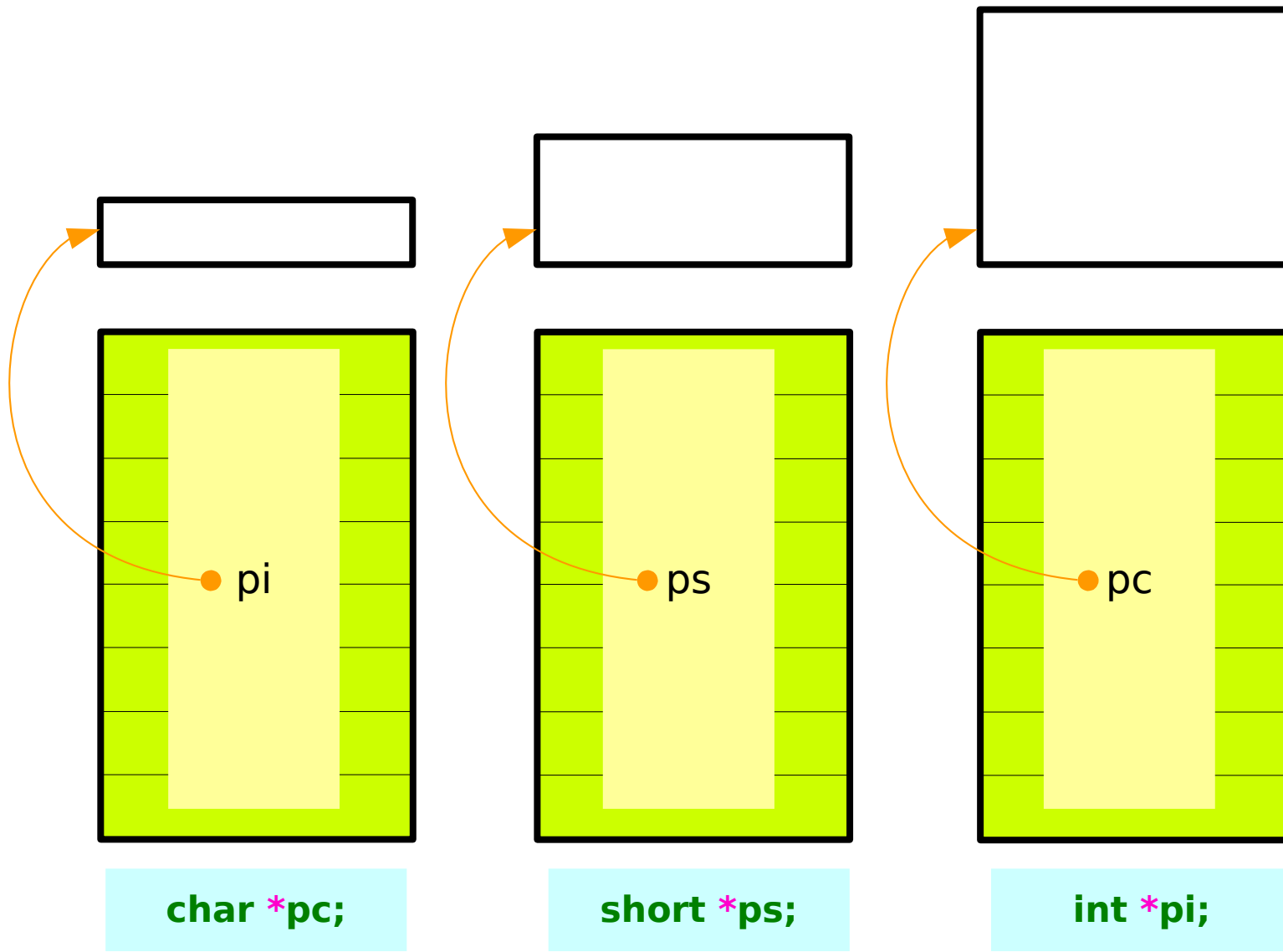
32-bit and 64-bit Address

32-bit machine : address : 4 bytes

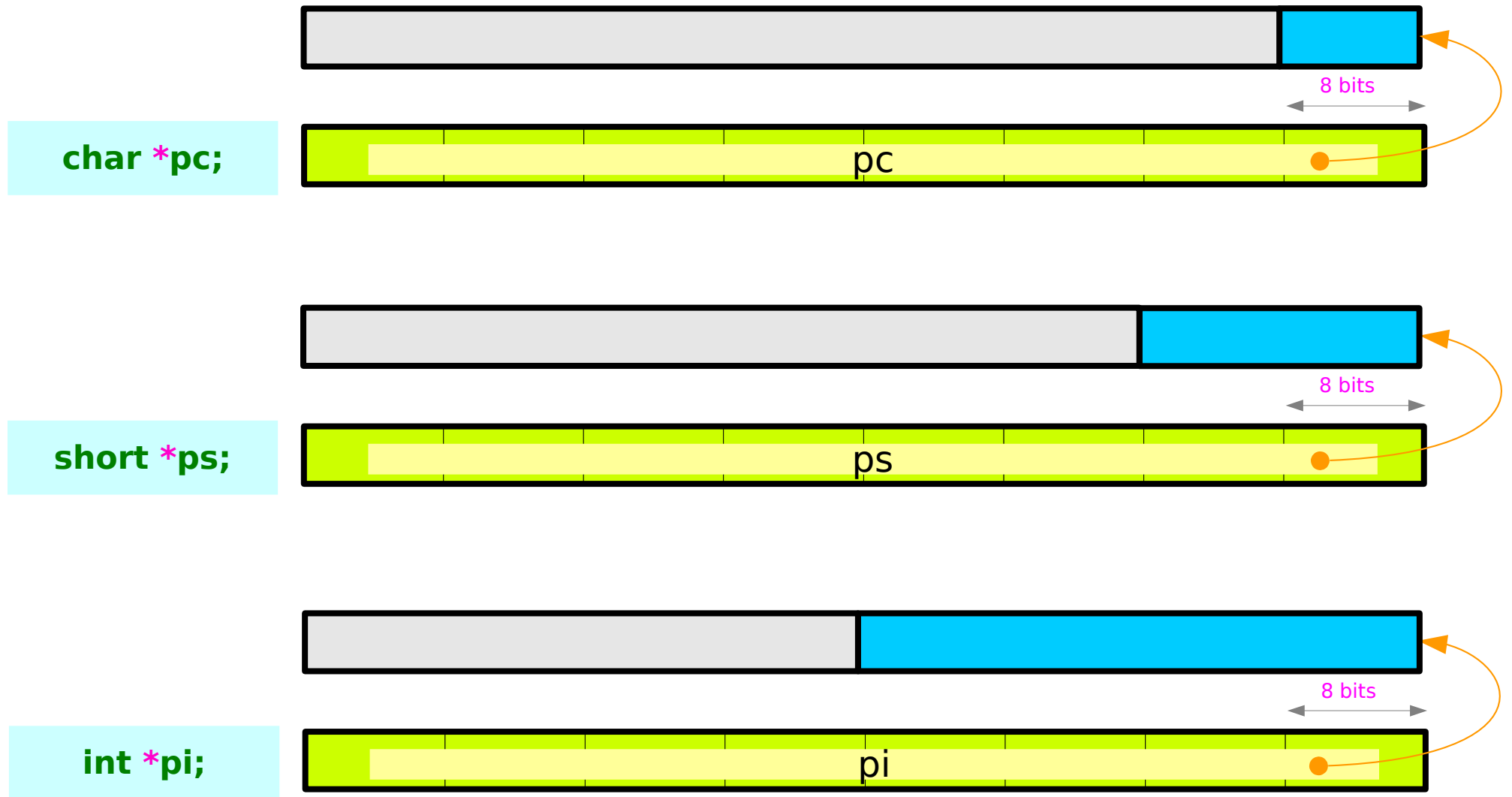
64-bit machine : address : 8 bytes



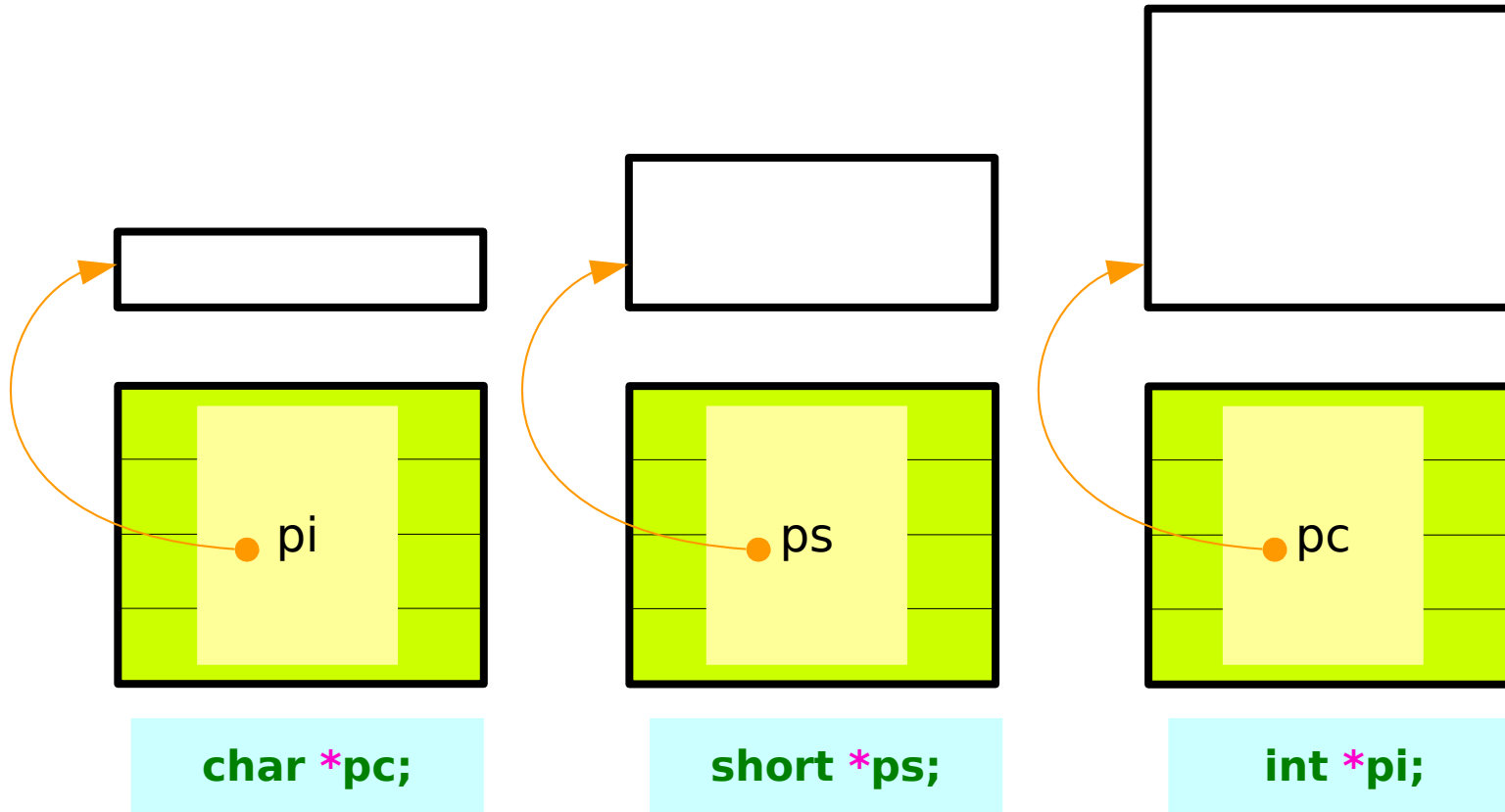
64-bit machine : 8-byte address



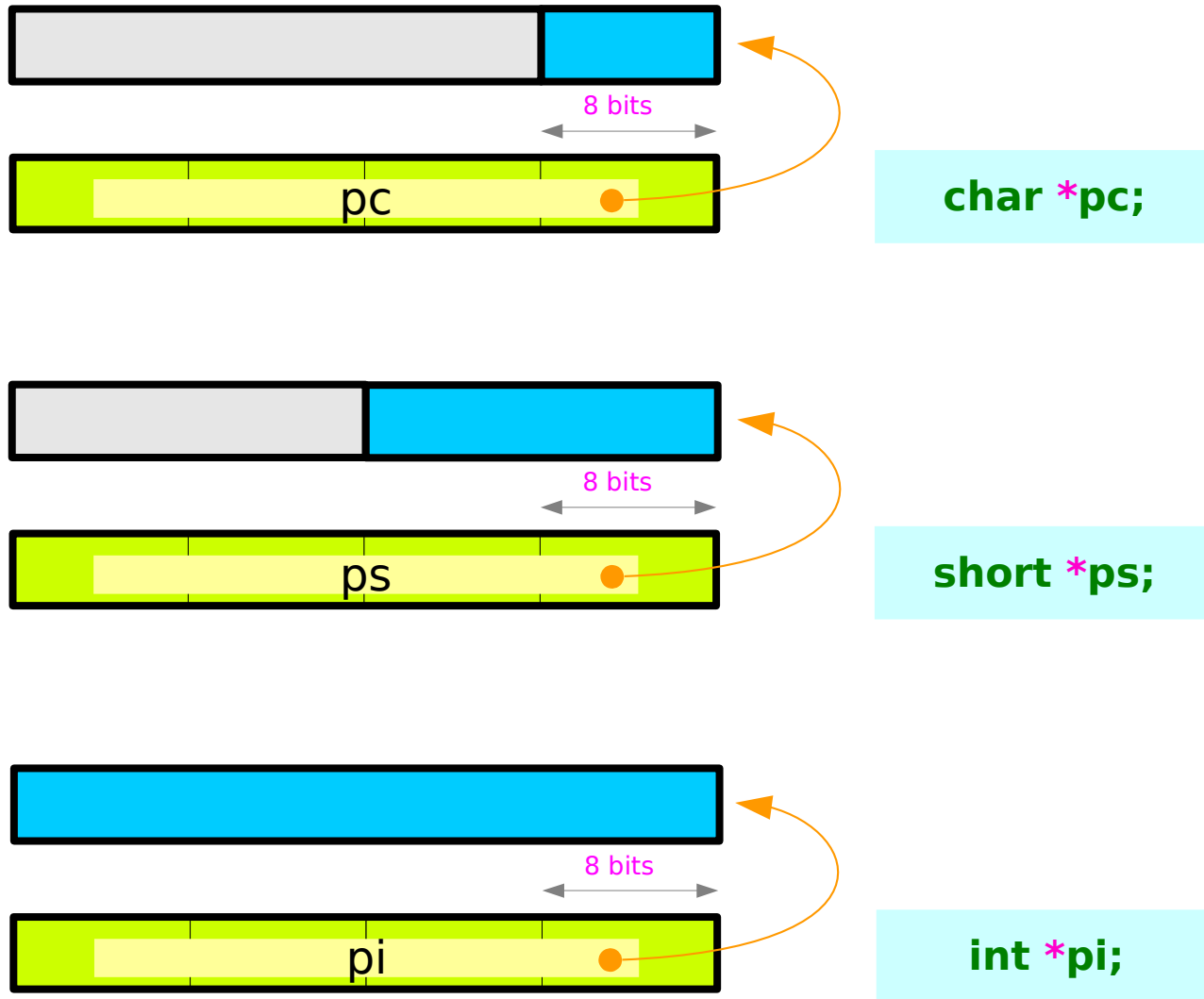
64-bit machine : 8-byte address & data buses



32-bit machine : 4-byte address



64-bit machine : 8-byte address and data buses



References

- [1] Essential C, Nick Parlante
- [2] Efficient C Programming, Mark A. Weiss
- [3] C A Reference Manual, Samuel P. Harbison & Guy L. Steele Jr.
- [4] C Language Express, I. K. Chun