Carry and Overflow

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Outline



Overview

Overview

3 Carry flag

- TOC: Carry flag
- Examples of signed and unsigned integer arithmetic
- Carry flag in unsigned and signed computations
- Rules for the carry flag
- Method for computing the carry flag

Overflow flag

- TOC: Overflow flag
- Overflow flag in unsigned and signed computations
- Rules for the overflow flag
- Method 1 for computing the overflow flag
- Method 2 for computing the overflow flag

 "Self-service Linux: Mastering the Art of Problem Determination", Mark Wilding

Computer Architecture: A Programmer's Perspective", Bryant & O'Hallaron

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- gcc -v
- gcc -m32 t.c
- sudo apt-get install gcc-multilib
- sudo apt-get install g++-multilib
- gcc-multilib
- g++-multilib
- gcc -m32
- objdump -m i386

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- Carry flag and overflow flag
- Signed and unsigned computations
- Flags for an <u>unsigned</u> number
- Flags for a signed number
- Detecting errors in usigned and signed arithmetic
- The verb to overflow v.s. the overflow flag

- considering carry and overflow flags in x86
- do not confuse the carry flag with the overflow flag in integer arithmetic.
- the ALU always sets these flags appropriately when doing any integer math.
- these flags can occur on its *own*, or *both* together.

- the CPU's ALU <u>doesn't</u> care or know whether signed or <u>unsigned</u> computations are performed;
- the <u>ALU</u> just performs integer arithmetic and sets the flags appropriately.
- It's up to the <u>programmer</u> to know which flag to check after the arithmetic is done.

- if a word is treated as an unsigned number,
 - the carry flag must be used to check if the result is fit into *n*-bit or (*n*+1)-bit number
 - the overflow flag is *irrelevant* to an <u>unsigned</u> number arithmetic

- if a word is treated as an signed number,
 - the carry flag is *irrelevant* to an signed number arithmetic
 - the overflow flag must be used to check if the result is wrong or not

	unsigned integer arithmetic	<mark>signed</mark> integer arithmetic
CF Carry Flag	detects <i>overflows</i>	
	extends an <i>n-bit</i> result	
	into an (<i>n+1</i>)-bit result	
OF Overflow Flag		detects <i>overflows</i>
		errors
		the result cannot be used

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- unsigned integer arithmetic overflow is indicated by the carry flag
 - P + P CF=1 \rightarrow carry out the result is too large for an *n*-bit integer
 - P P CF=1 \rightarrow borrow in the result is too small for an *n*-bit integer
- signed integer arithmetic overflow is indicated by the overflow flag
 - $P + P \rightarrow N$ OF=1 \rightarrow overflow the result is not correct
 - $N + N \rightarrow P$ OF=1 \rightarrow overflow the result is not correct
- P (positive), N (negative)

Detecting errors in usigned and signed arithmetic (3)

- unsigned integer arithmetic *overflow* is indicated by the carry flag
 - the overflowed n-bit result can be extended into (n+1)-bit result by using the carry flag
- signed integer arithmetic *overflow* is indicated by the overflow flag
 - the overflowed n-bit result cannot be used

- Do not confuse the English verb to overflow with the overflow flag in the ALU.
- The verb to overflow is used casually to indicate that some math result doesn't fit in the number of bits available;
- it could be integer math, or floating-point math, or whatever.
- The overflow flag is set specifically by the ALU it isn't the same as the casual English verb "to overflow"

The verb to overflow v.s. the overflow flag (2)

- In English, we may say
 "the binary/integer math overflowed
 the number of bits available for the result,
 causing the carry flag to come on".
- Note how this English usage of the verb "to overflow" is not the same as saying the overflow flag is on".
- A math result can <u>overflow</u> (the <u>verb</u>) the number of bits available without turning on the ALU <u>overflow</u> flag

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt

CF (carry flag) and OF (overflow flag) computation

ADD (addition)	SUB (subtraction)
$CF = C_n$	$CF = \overline{C_n}$
$OF = C_n \bigoplus C_{n-1}$	$OF = C_n \bigoplus C_{n-1}$
a 2's complement addition A + B = A + B + 0	a transformed addition $A - B = A + \overline{B} + 1$
$\{C_n, S_{n-1}\} = a_{n-1} + b_{n-1} + c_{n-1}$	$\{C_n, S_{n-1}\} = a_{n-1} + \overline{b_{n-1}} + c_{n-1}$
$\{C_{n-1}, S_{n-2}\} = a_{n-2} + b_{n-2} + c_{n-2}$	$\{C_{n-1}, S_{n-2}\} = a_{n-2} + \overline{b_{n-2}} + c_{n-2}$

https://www.csie.ntu.edu.tw/~cyy/courses/assembly/12fall/lectures/handouts/lec14_:

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- Carry flag in unsigned and signed computations
- Rules for the carry flag
- Method for computing the carry flag

TOC: Examples of signed and unsigned integer arithmetic

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Examples of signed and unsigned integer arithmetic (1)

• 0xFFFFBDC3 as a signed (negative) number -0x0000423D (-16957₁₀)

	F	F	F	F	В	D	С	3		
-0xFFFFBDC3	0x1111	_1111	_1111	_1111	_1011	_1101_	_1100_	0011		
0x0000423D	-0x0000	_0000	_0000	_0000_	_0100	_0010_	_0011_	1100	(1's	complement)
0x0000423D	-0x0000	_0000	_0000	_0000_	_0100	_0010	_0011_	1101	(1's	complement)
	0	0	0	0	4	2	3	D		
	0	0	0	0	4	2	3	D		
-0x0000423D	-0x0000	_0000	_0000	_0000_	_0100	_0010	_0011_	1101		
0x0000BDC2	0x1111	_1111	_1111	_1111	_1011	_1101	_1100_	0010	(1's	complement)
0xFFFFBDC3	0x1111.	_1111	_1111	_1111_	_1011	_1101	_1100_	_0011	(2's	complement)
	F	F	F	F	В	D	С	3		

• 0xFFFFBDC3 as an unsigned (positive) number (+4294950339₁₀)

https://stackoverflow.com/questions/47333458/assembly-x86-64-setting-carry-flag-f

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Examples of signed and unsigned integer arithmetic (2)

• 0x0000195D - 0x0000618D : unsigned subtraction subtraction by hand

0 5 0 0 0 1 9 D 0x0000 0000 0000 0000 0001 1001 0101 1101 0x0000195D - 0x0000618D 0x0000 0000 0000 0000 0110 0001 1000 1101 0 0 0 0 6 1 8 0xFFFFB7D0 1 0x1111_1111_1111_1111_011_0111_1101_0000 (hand subtraction) F F F F 7 D 1 В 0 V borrow (CF=1) : unsigned integer overflow

- A borrow is indicated by the carry flag (CF=1)
 - whenever an unsigned integer overflow happened
 - A B, when A < B, for non-negative integers A, B

https://stackoverflow.com/questions/47333458/assembly-x86-64-setting-carry-flag-f

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Examples of signed and unsigned integer arithmetic (3)

 0x0000195D + (-0x0000618D) : signed subtraction the transformed addition using the 2's complement of subtrahend

		0	0	0	0	6	1	8	D	
-0x0000618D		-0x0000_	0000	_0000_	_0000	_0110_	0001	_1000	_1101	
0xFFFF9E73		0x1111_	1111	1111	1111	_1001_	1110	_0111	_0011	(2's complement)
		F	F	F	F	8	E	7	3	
0x0000195D		0x0000_	0000	0000	_0000	_0001_	1001	_0101	_1101	(+0x0000195D)
+ 0xFFFF9E73		0x1111_	1111	1111	1111	_1001_	1110	_0111_	_0011	(-0x0000618D)
										-
0xFFFFB7D0	0	0x1111_	1111	1111	1111	_1011_	0111	_1101	_0000	(hand addition)
	0	F	F	F	F	В	7	D	0	
-0x00004830		0x0000_	0000	0000	_0000	_0100_	1000	_0011	_0000	(2's complement)
		0	0	0	0	4	8	3	0	
	V	no carr	y in	the 1	transi	formed	l add:	ition	(Cn=	0)> (CF=1)

• signed integer overflow is indicated not by the carry flag (CF), but by the overflow flag (OF)

• the carry flag is set by the inverted carry of a transformed addition

https://stackoverflow.com/questions/47333458/assembly-x86-64-setting-carry-flag-f

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• 0x0000195D - 0x0000618D

- 0x0000195D 0x0000618D : hand subtraction unsigned integer subtraction
- 0x0000195D + (-0x0000618D) : the transformed addition using the 2's complement of the <u>subtrahend</u> signed integer subtraction
- the result is 0xFFFFB7D0 (the two methods have the same bit pattern)
 - interpreting as a unsigned integer 4294948816₁₀ 0xFFFB7D0 with a borrow (CF=1)
 - interpreting as a signed integer -18480₁₀
 -0x00004830 (meaningless CF=1)

0xFFFFB7D0the result of unsigned subtraction429494881610with CF=1with unsigned integer overflow

-0x00004830 the result of signed subtraction -18480₁₀

https://stackoverflow.com/questions/47333458/assembly-x86-64-setting-carry-flag-fo

Examples of signed and unsigned integer arithmetic (6)

- 0x0000195D 0x0000618D : unsigned subtraction
 - there is an unsigned integer overflow so the carry flag will be set (CF=1) to indicate a borrow (A - B, when A < B, for non-negative integers A, B) (unsigned integers can't be negative),
- 0x0000195D + (-0x0000618D) : signed subtraction
 - there is no signed integer overflow the overflow flag won't be set (OF=0)
 - signed overflw occurrs , in the transformed addition,
 - two *positive* numbers are added and the result is a *negative*, $(P + P \rightarrow N)$, or
 - two negative numbers are added and the result is a positive, (N + N → P)

- Using the Carry Flag as a borrow
- Examples of signed and unsigned integer arithmetic

0111	(+7)	1000	(-8)
0110	(+6)	1001	(-7)
0101	(+5)	1010	(-6)
0100	(+4)	1011	(-5)
0011	(+3)	1100	(-4)
0010	(+2)	1101	(-3)
0001	(+1)	1110	(-2)
0000	(0)	1111	(-1)

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n	bits	addened	Α	$\{a_{n-1}, a_{n-2}, \cdots, a_1, a_0\}$
n	bits	augend	В	$\{b_{n-1}, b_{n-2}, \cdots, b_1, b_0\}$
(n+1)	bits	carry bits	С	$\{C_n, C_{n-1}, C_{n-2}, \cdots, C_1, C_0\}$
n	bits	sum bits	S	$\{S_{n-1}, S_{n-2}, \cdots, S_1, S_0\}$

external carry bits : C_n carry out, C_0 carry in

	a_{n-1}	a_{n-2}	• • • • • • • • • •	a_1	a ₀
	b_{n-1}	b_{n-2}		b_1	b_0
					C ₀
Cn	S_{n-1}	S_{n-2}		S_1	S_0

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full adder operation in the i^{th} bit position

$$\{C_{i+1},S_i\}=a_i+b_i+C_i$$

$$\begin{array}{c} a_i\\ b_i\\ C_i\\ \hline C_{i+1} \quad S_i \end{array}$$

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt

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Internal and external carry bits

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Addition and Subtraction

addition

$$\{C_n, S\} = A + B = A + B + 0$$

$$a_{n-1} \quad a_{n-2} \quad \dots \quad a_1 \quad a_0$$

$$b_{n-1} \quad b_{n-2} \quad \dots \quad b_1 \quad b_0$$

$$C_{n-1} \quad C_{n-2} \quad \dots \quad \dots \quad C_1 \quad 0$$

	C_{n-1}	C_{n-2}	• • • • • • • • • •	C_1	0
Cn	S_{n-1}	S_{n-2}	• • • • • • • • • •	S_1	S_0

• subtraction - transformed addition

$$\{C_n,S\}=A-B=A+\overline{B}+1$$

	a_{n-1}	a_{n-2}	• • • • • • • • • •	a_1	a ₀
	$\overline{b_{n-1}}$	$\overline{b_{n-2}}$		$\overline{b_1}$	$\overline{b_0}$
	C_{n-1}	C_{n-2}		C_1	1
Cn	S_{n-1}	S_{n-2}		S_1	S_0

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- a borrow (CF=1) occurs in the subtraction A – B when b is larger than a (A < B) as unsigned numbers
- Computer hardware can detect

 a borrow (CF=1) in subtraction
 by looking at whether a carry out (Cn) occurred
 in the transformed addition

a borrow (CF=1) occurs
 in the subtraction A - B (A < B)
 as unsigned numbers

• a carry out (Cn) in the transformed addition

- If there is <u>no</u> carry (Cn=0) then there is a borrow (CF=1)
- If there is a carry (Cn=1) then there is no borrow (CF=0)

• CF = !Cn

- the same *addition* and *subtraction* instructions are used for both <u>unsigned</u> and <u>signed</u> integer arithmetic.
 - no special *addition* and *subtraction* instructions for <u>unsigned</u> and <u>signed</u> integer arithmetic
- the only difference is
 - which flags you test afterwards and
 - how you *interpret* the result

- 2's complement numbers : 4-bit
- The 1st rule for setting the carry flag
- The 2nd rule for setting the carry flag
- Cases for clearing the carry flag

0111	(+7)	1000	(-8)
0110	(+6)	1001	(-7)
0101	(+5)	1010	(-6)
0100	(+4)	1011	(-5)
0011	(+3)	1100	(-4)
0010	(+2)	1101	(-3)
0001	(+1)	1110	(-2)
0000	(0)	1111	(-1)

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The <u>1st</u> rule for <u>setting</u> the carry flag

The carry flag is set (CF = 1 : carry in addition) if the addition of two unsigned numbers causes a carry out of the most significant (leftmost) bits added. (hand addition rule)

signed addition	signed subtraction	unsigned addition
1111 (-1) +0001 +(+1)	1111 (-1) -1111 -(-1)	1111 (15) +0001 +(1)
10000 (0)	10000 (0)	10000 (16)
Cn=1 -> CF=1	Cn=1 -> CF=1	CF=1
CF is not 16 S = 0000	CF is not 16 S = 0000	$\begin{array}{r} \text{CF means } 16\\ \text{S} = 0000 \end{array}$

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The 2nd rule for setting the carry flag

The carry flag is also set (CF = 1 : borrow in subtraction) if the subtraction of two unsigned numbers requires a borrow into the most significant (leftmost) bits subtracted. (hand subtraction rule)

signed addition	signed subtraction	unsigned subtraction
0000 (0) +1111 +(-1)	0000 (0) -0001 -(+1)	0000 (0) -0001 -(1)
01111 (-1)	01111 (-1)	01111 (15) (-16)
Cn=0 -> CF=1	Cn=0 -> CF=1	CF=1
CF is not -16 S = 1111	CF is not -16 S = 1111	CF means -16 S = 1111

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• Otherwise, the carry flag is turned off (zero).

signed addition signed subtraction		unsigned addition	
0111 (+7) +0001 +(+1)	0111 (+7) -1111 -(-1)	0111 (7) +0001 +(1)	
01000 (-8)	01000 (-8)	01000 (8)	
Cn=0 -> CF=0	Cn=0 -> CF=0	CF=0	

• Otherwise, the carry flag is turned off (zero).

signed addition	signed subtraction	unsigned subtraction
1000 (-8) +1111 +(-1)	1000 (-8) -0001 -(+1)	1000 (8) -0001 -(1)
10111 (+7)	10111 (+7)	00111 (7)
Cn=1 -> CF=0	Cn=1 -> CF=0	CF=0

• Carry flag computation

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ADD (addition)	SUB (subtraction)
$CF = C_n$	$CF = \overline{C_n}$
normal carry of a 2's complement addition	inverted carry of a transformed addition

A + B = A + B + 0

A - B = A + B + 1

$$\{C_n, S_{n-1}\} \qquad \{C_n, S_{n-1}\} \\ = a_{n-1} + b_{n-1} + c_{n-1} \qquad = a_{n-1} + \overline{b_{n-1}} + c_{n-1}$$

https://www.csie.ntu.edu.tw/~cyy/courses/assembly/12fall/lectures/handouts/lec14_:

Carry flag computation (2)

• In unsigned arithmetic,

- the carry flag is used to detect overflow
- the carry flag is used to extend *n*-bit result into (*n*+1)-bit result
- for addition, the carry flag is a carry out
- for subtraction, the carry flag is a borrow in

In signed arithmetic,

- the carry flag is useless
- the carry flag neither detects overflow nor extends n-bit result

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In unsigned arithmetic,

Addition	CF = 1 means carry out	when $Cn = 1$
Subtraction	CF = 1 means borrow in	when $Cn = 0$

- CF Carry Flag in x86
- Cn the normal carry out
 - the carry out of a 2's complement addition for ADD
 - the carry out of a transformed addition for SUB
- In signed arithmetic,
 - the carry flag is useless

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- Overflow flag in unsigned and signed computations
- Rules for the overflow flag
- Method 1 for computing the overflow flag
- Method 2 for computing the overflow flag

TOC Overflow flag in unsigned and signed computations

• Overflow flag

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Carry and Overflow

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- only need to look at the sign bits (leftmost) of the three numbers to decide if the overflow flag is turned on or off.
- overflow flag is based on signed arithmetic
- in signed arithmetic, watch the overflow flag to detect errors.
- in unsigned arithmetic, the overflow flag tells you nothing interesting

- for signed (two's complement) arithmetic, overflow flag on means the answer is wrong
 - two positive numbers are added and the result is a negative, (P + P \rightarrow N), or
 - two negative numbers are added and the result is a positive, $(N + N \rightarrow P)$
 - opposite signed numbers are added, then no overflow

•
$$(P + N \rightarrow P \text{ or } N)$$

- $(N + P \rightarrow P \text{ or } N)$
- for unsigned arithmetic,

the overflow flag means nothing and should be ignored

- the rules for two's complement detect errors by examining the sign of the result.
- a <u>negative</u> and <u>positive</u> added together <u>cannot</u> be <u>wrong</u>, because the sum is between the addends.
- mixed-sign addition never turns on the overflow flag.
- since both of the <u>addends</u> fit within the <u>allowable range</u> of numbers, and their sum is between them, it must fit as well.

- The <u>1st</u> rule for setting the overflow flag
- The <u>2nd</u> rule for setting the overflow flag
- Cases for clearing the overflow flag

 If the sum of two signed numbers with the sign bits off (0, 0) yields a result number with the sign bit on (1) the overflow flag is turned on (OF =1 : +, + → -)

signed	addition	signed	l subtraction	n unsign	ed addition
0100 +0100	(+4) +(+4)		(+4) -(-4)	0100 +0100	.,
01000	(-8)	01000	(-8)	01000	(8)

If the sum of two numbers with the sign bits on (1, 1) yields a result number with the sign bit off (0) the overflow flag is turned on. (OF =1 : -, - → +) signed addition (-7) + (-7) = (2) with a borrow (-16) 1001 + 1001 = 1 0010 (2's complement addition) (-, - → +) signed subtraction (-7) - (7) = (2) with a borrow (-16) 1001 - 0111 = 1 0010 (transformed subtraction)

unsigned addition (9) + (9) = (18) 1001 + 1001 = 1 0010

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• overflow flag is turned off. (OF = $0: +, + \rightarrow +$)

signed addition	signed subtraction	unsigned addition	
0011 (+3)	0011 (+3)	0011 (3)	
+0011 +(+3)	-1101 -(-3)	+0011 +(3)	
00110 (+6)	00110 (+6)	00110 (6)	

• overflow flag is turned off. (OF = 0 : $-, - \rightarrow -$)

signed addition	signed subtraction	unsigned addition	
1101 (-3)	1101 (-3)	1101 (13)	
+1101 +(-3)	-0011 - (+3)	+1101 +(13)	
11010 (-6)	11010 (-6)	11010 (26)	

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Image: Image:

• overflow flag is turned off. (DF = 0 : $+, - \rightarrow +$)

signed addition	signed subtraction	unsigned addition	
0100 (+4) +1101 +(-3)	0100 (+4) -0011 -(+3)	0100 (4) +1101 +(13)	
10001 (+1)	10001 (+1)	10001 (17)	

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• overflow flag is turned off. (OF = 0 : $+, - \rightarrow -$)

signed addition	signed subtraction	unsigned addition
0011 (+3) +1100 +(-4)	0011 (+3) -0100 -(+4)	0011 (3) +1100 +(12)
01111 (-1)	01111 (-1)	01111 (15)

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Image: Image:

• overflow flag is turned off. (OF = 0 : $-, + \rightarrow +$)

signed addition	signed subtraction	unsigned addition
1101 +(-3) +0100 (+4)	0011 - (+3) -0100 (+4)	1101 +(13) +0100 (4)
10001 (+1)	10001 (+1)	10001 (17)

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• overflow flag is turned off. (OF = 0 : $-, + \rightarrow -$)

signed addition	signed subtraction	unsigned addition
1100 +(-4)	0100 -(+4)	1100 +(12)
+0011 (+3)	-0011 (+3)	+0011 (3)
01111 (-1)	01111 (-1)	01111 (15)

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- Adding two numbers with the same sign
- Overflow conditions for additions and subtractions
- Overflow condition for an addition
- Overflow conditions for a subtraction
- Overflow in signed computations

- overflow can only happen when <u>adding</u> two numbers of the <u>same</u> <u>sign</u> results in a different sign.
- signed binary arithmetic
- to detect overflow
 - only the sign bits are considered
 - the other bits are ignored

- with two <u>operands</u> and one <u>result</u>, three sign bits are considered $2^3 = 8$ possible combinations
- only two cases result in overflow for an addition

• 0 0 1
$$(p+p \rightarrow n)$$

- 1 1 0 $(n+n \rightarrow p)$
- only two cases are considered as overflow for an subtraction

• 0 1 1
$$(p - n \to n)$$

• 100 $(n - p \to p)$

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Overflow in an addition (num1 + num2)

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• Overflow in a subtraction (num1 - num2)

```
num1 num2 sub (num1 - num2)
  sign sign sign
                _____
     0 0 0
     0 0 1
     0 1 0
*OVER* 0 1 1 (subtracting a negative is the same as adding a positive)
*OVER* 1 0 0 (subtracting a positive is the same as adding a negative)
     101
      1 1 0
     1 1 1
```

- subtracting a *positive* number is the same as adding a *negative*
- subtracting a *negative* number is the same as adding a *positive*

- A computer might contain a small logic gate array that <u>sets</u> the overflow flag to "1" iff any one of the above four OV conditions is met.
- in signed computations, <u>adding</u> two numbers of the same sign must produce a result of the same sign, otherwise overflow happened.

- Carry into and carry out of the sign bit
- Overflow in 2's complement arithmetic
- Overflow flag = $C_n \bigoplus C_{n-1}$
- Examples
- C_n and C_{n-1} in a *n*-bit addition
- Overflow flag computation
- Examples of computing overflow flag
- Hexadecimal carry, octal carry, decimal carry
- No carry into the sign bit

• When adding two binary values, consider

- the carry *coming into* the <u>leftmost</u> place (carry into the sign bit)
- the carry going out of that leftmost place. (carry out of the sign bit) this is the carry flag in the ALU

- overflow in 2's complement happens when
 - there is a carry *into* the sign bit but <u>no</u> carry *out of* the sign bit.
 - there is <u>no carry</u> into the sign bit but a carry out of the sign bit.

- the overflow flag is the XOR
 - of the carry coming into the sign bit
 - with the carry going out of the sign bit
- overflow happens when the carry in does <u>not</u> equal to the carry out

Examples

4-bit signed 2's complement addition examples

```
1100 (-4) (neg sign 1)
                            1100 (-4) (neg sign 1)
+0100 (+4) (pos sign 0)
                              +1000 (-8) (neg sign 1)
_____
                              _____
10000 ( 0) (pos sign 0)
                              10100 (+4) (pos sign 0)
carry in 1 (1+1+0)
                            carrv in 0 (1+0+0)
carry out 1 (1+0+1)
                           carry out 1 (1+1+0)
1 \text{ XOR } 1 = \text{NO OVERFLOW}
                              0 \text{ XOR } 1 = 0 \text{VERFLOW}!
0100 (+4) (pos sign 0)
                           1000 (-8) (neg sign 1)
+0100 (+4) (pos sign 0)
                              +0100 (+4) (pos sign 0)
  _____
                                 _____
01000 (-8) (neg sign 1)
                              01100 (-4) (neg sign 1)
carry in 1 (1+1+0)
                              carry in 0 (0+1+0)
carry out 0 (0+0+1)
                              carry out 0 (1+0+0)
1 \text{ XOR } 0 = 0 \text{VERFLOW}!
                              O XOR O = NO OVERFLOW
```

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$(n-1)^{th}$ bit – MSB

 adding operations at the (n - 1) bit position

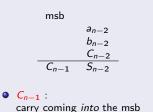
•
$$\{C_n, S_{n-1}\} = a_{n-1} + b_{n-1} + c_{n-1}$$

$$\begin{array}{c}
\text{msb}\\
a_{n-1}\\
b_{n-1}\\
C_{n-1}\\
\hline
C_n & S_{n-1}\\
\end{array}$$

• C_n: carry coming <u>out of</u> the msb

$(n-2)^{th}$ bit

- adding operations at the (n-2) bit position
- $\{C_{n-1}, S_{n-2}\} = a_{n-2} + b_{n-2} + c_{n-2}$



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ADD (addition) SUB (subtraction) OF = $C_n \bigoplus C_{n-1}$ OF = $C_n \bigoplus C_{n-1}$ a 2's complement addition the transformed addition A + B = A + B + 0 $A - B = A + \overline{B} + 1$ $\{C_n, S_{n-1}\}$ $\{C_n, S_{n-1}\}$ $= a_{n-1} + b_{n-1} + c_{n-1}$ $= a_{n-1} + \overline{b_{n-1}} + c_{n-1}$ $\{C_{n-1}, S_{n-2}\}$ $\{C_{n-1}, S_{n-2}\}$ $=a_n + \overline{b_n} + c_n$ $= a_{n-2} + b_{n-2} + c_{n-2}$

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• 4-bit signed addition examples

+4	-7	+4	+6	-8	-4
+4	-7	+1	-7	+1	-4
-8	+2	+5	-1	-7	-8
0100	1001	0100	0110	1000	1100
0100	1001	0001	1001	0001	1100
01000	10010	00000	00000	00000	11000
1000	0010	0101	1111	1001	1000
C4 = 0	C4 = 1	C4 = 0	C4 = 0	C4 = 0	C4 = 1
C3 = 1	C3 = 0	C3 = 0	C3 = 0	C3 = 0	C3 = 1
+ +, -	, +	+ +, +	+ -, -	- +, -	, -
OF = 1	OF = 1	OF = O	OF = 0	OF = O	OF = O

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- Note that this XOR method only works with the binary carry that goes into the sign bit.
- not works with hexadecimal carry decimal carry, octal carry
 - the carry doesn't go into the sign bit
 - can't XOR that non-binary carry with the outgoing carry.

No carry into the sign bit

- Hexadecimal addition example (showing that XOR doesn't work for hex carry): 8Ah +8Ah ==== 114h
- The hexadecimal carry of 1 resulting from A+A does not affect the sign bit.
- If you do the math in binary, you'll see that there is no carry into the sign bit; but, there is carry out of the sign bit. Therefore, the above example sets OVERFLOW on. (The example adds two negative numbers and gets a positive number.)