Redundant CORDIC Timmermann (C)

20170104

Termination Algorithms
Modified CORDIC
CSD (Canonic Sign Digit) Encoding
Booth Encoding

Low Latency Time CORDIC Algorithms - Timmermann (1992)
Redundant and on-line CORDIC - Ercegovac & Lang (1990)

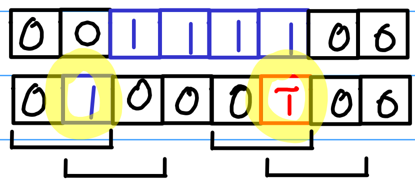
CSD (Canonic Signed Digit)

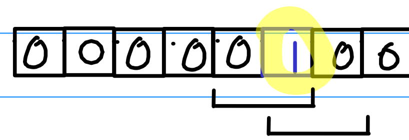
like Booth encoding (not modified Booth)

all non-zero digits are separated by zeros

$\Rightarrow$ $\boxed{\sigma_i \sigma_{i+1} = 0}$

·1-run





$0 \cdot 1 = 0$ $\quad$ $0 \cdot \bar{1} = 0$

$1 \cdot 0 = 0$ $\quad$ $\bar{1} \cdot 0 = 0$

$$
\boxed{
\begin{aligned}
x_{i+1} &= x_i - m\,\sigma_i\,2^{-S(m,i)}\,y_i \\
y_{i+1} &= y_i + \sigma_i\,2^{-S(m,i)}\,x_i \\
z_{i+1} &= z_i - \sigma_i\,\alpha_{m,i}
\end{aligned}
}
$$

$i+1$

$$
\begin{bmatrix} x_{i+1} \\ y_{i+1} \end{bmatrix} = \begin{bmatrix} 1 & -m\,\sigma_i 2^{-i} \\ \sigma_i 2^{-i} & 1 \end{bmatrix} \begin{bmatrix} x_i \\ y_i \end{bmatrix}
$$

$i+2$

$$
\begin{bmatrix} x_{i+2} \\ y_{i+2} \end{bmatrix} = \begin{bmatrix} 1 & -m\,\sigma_{i+1} 2^{-i-1} \\ \sigma_{i+1} 2^{-i-1} & 1 \end{bmatrix} \begin{bmatrix} x_{i+1} \\ y_{i+1} \end{bmatrix}
$$

$$
= \begin{bmatrix} 1 & -m\,\sigma_{i+1} 2^{-i-1} \\ \sigma_{i+1} 2^{-i-1} & 1 \end{bmatrix} \begin{bmatrix} 1 & -m\,\sigma_i 2^{-i} \\ \sigma_i 2^{-i} & 1 \end{bmatrix} \begin{bmatrix} x_i \\ y_i \end{bmatrix}
$$

$$
= \begin{bmatrix} 1 - m\,\sigma_{i+1}\sigma_i\,2^{-2i-1} & -m\,\sigma_i 2^{-i} - m\,\sigma_{i+1} 2^{-i-1} \\ \sigma_{i+1} 2^{-i-1} + \sigma_i 2^{-i} & -m\,\sigma_{i+1}\sigma_i\,2^{-2i-1} + 1 \end{bmatrix} \begin{bmatrix} x_i \\ y_i \end{bmatrix}
$$

$$
= \begin{bmatrix} 1 - m\,\sigma_i\sigma_{i+1}\,2^{-2i-1} & -m\left(\sigma_i 2^{-i} + \sigma_{i+1} 2^{-i-1}\right) \\ \left(\sigma_i 2^{-i} + \sigma_{i+1} 2^{-i-1}\right) & 1 - m\,\sigma_i\sigma_{i+1}\,2^{-2i-1} \end{bmatrix} \begin{bmatrix} x_i \\ y_i \end{bmatrix}
$$

$$
\begin{bmatrix} 1 - m\,\boxed{\sigma_i\sigma_{i+1}}\,2^{-2i-1} & -m\left(\sigma_i 2^{-i} + \sigma_{i+1} 2^{-i-1}\right) \\ \left(\sigma_i 2^{-i} + \sigma_{i+1} 2^{-i-1}\right) & 1 - m\,\boxed{\sigma_i\sigma_{i+1}}\,2^{-2i-1} \end{bmatrix}
$$

$\boxed{\sigma_i \neq 0} \longrightarrow \boxed{\sigma_i \sigma_{i+1} = 0}$ property of Booth encoding

$$\begin{bmatrix} x_{i+1} \\ y_{i+1} \end{bmatrix} = \begin{bmatrix} 1 & -m\,\sigma_i 2^{-i} \\ \sigma_i 2^{-i} & 1 \end{bmatrix} \begin{bmatrix} x_i \\ y_i \end{bmatrix}$$

$$\begin{bmatrix} x_{i+2} \\ y_{i+2} \end{bmatrix} = \begin{bmatrix} 1 - m \boxed{\sigma_i \sigma_{i+1}} 2^{-2i-1} & -m(\sigma_i 2^{-i} + \sigma_{i+1} 2^{-i-1}) \\ (\sigma_i 2^{-i} + \sigma_{i+1} 2^{-i-1}) & 1 - m \boxed{\sigma_i \sigma_{i+1}} 2^{-2i-1} \end{bmatrix} \begin{bmatrix} x_i \\ y_i \end{bmatrix}$$

$$\Downarrow$$

$$\begin{bmatrix} x_{i+2} \\ y_{i+2} \end{bmatrix} = \begin{bmatrix} 1 & -m(\sigma_i 2^{-i} + \sigma_{i+1} 2^{-i-1}) \\ (\sigma_i 2^{-i} + \sigma_{i+1} 2^{-i-1}) & 1 \end{bmatrix} \begin{bmatrix} x_i \\ y_i \end{bmatrix}$$

$$x_{i+2} = x_i - m(\sigma_i 2^{-i} + \sigma_{i+1} 2^{-i-1})\, y_i$$

$$y_{i+2} = (\sigma_i 2^{-i} + \sigma_{i+1} 2^{-i-1})\, x_i + y_i$$

$$\boxed{\begin{aligned} x_{i+2} &= x_i - m\,\sigma_i 2^{-i} y_i - m\,\sigma_{i+1} 2^{-i-1} y_i \\ y_{i+2} &= y_i + \sigma_i 2^{-i} x_i + \sigma_{i+1} 2^{-i-1} x_i \\ z_{i+2} &= z_i - \sigma_i \alpha_{m,i} - \sigma_{i+1} \alpha_{m,i+1} \end{aligned}}$$

$\boxed{\sigma_i = 0} \rightarrow$ inc/dec no rotation,

but compensate the scale factor.

$$\boxed{\begin{aligned} x_{i+1} &= x_i + m \cdot 2^{-2i-1} x_i \\ y_{i+1} &= y_i + m \cdot 2^{-2i-1} y_i \\ z_{i+1} &= z_i \end{aligned}}$$

inc/dec

inc/dec

$m=+1$ / $m=-1$

$$\begin{aligned} x_{i+1} &= (1 + m \cdot 2^{-2i-1}) x_i \\ y_{i+1} &= (1 + m \cdot 2^{-2i-1}) y_i \\ z_{i+1} &= z_i \end{aligned}$$

$$\begin{aligned} x_{i+2} &= (1 + m \cdot 2^{-2i-3}) x_{i+1} = (1 + m \cdot 2^{-2i-3})(1 + m \cdot 2^{-2i-1}) x_i \\ y_{i+2} &= (1 + m \cdot 2^{-2i-3}) y_{i+1} = (1 + m \cdot 2^{-2i-3})(1 + m \cdot 2^{-2i-1}) y_i \\ z_{i+1} &= z_i \end{aligned}$$

$$2^{-2i-3} \cdot 2^{-2i-1} = 2^{-4i-4} \ll 1$$

$$\boxed{\begin{aligned} x_{i+2} &= (1 + m \cdot 2^{-2i-3} + m \cdot 2^{-2i-1}) x_i \\ y_{i+2} &= (1 + m \cdot 2^{-2i-3} + m \cdot 2^{-2i-1}) y_i \\ z_{i+1} &= z_i \end{aligned}}$$

$$m=1, \quad s(m,i) = i$$

Cond (I) $\quad 0 \leq i \leq \frac{1}{4}(n-3)$

$$\begin{aligned} x_{i+1} &= x_i - \sigma_i 2^{-i} y_i \\ y_{i+1} &= y_i + \sigma_i 2^{-i} x_i \\ z_{i+1} &= z_i - \sigma_i \tan^{-1}(2^{-i}) \end{aligned}$$

$$\begin{aligned} x_{i+1} &= x_i - \sigma_i 2^{-i} y_i \\ y_{i+1} &= y_i + \sigma_i 2^{-i} x_i \\ z_{i+1} &= z_i - \sigma_i \tan^{-1}(2^{-i}) \end{aligned}$$

Cond (II) $\quad \frac{1}{4}(n-3) < i \leq \frac{1}{2}(n+1)$

$\sigma_i \neq 0$

$$\begin{aligned} x_{i+1} &= x_i - \sigma_i 2^{-i} y_i \\ y_{i+1} &= y_i + \sigma_i 2^{-i} x_i \\ z_{i+1} &= z_i - \sigma_i \tan^{-1}(2^{-i}) \end{aligned}$$

$$\begin{aligned} x_{i+2} &= x_i - m\,\sigma_i 2^{-i} y_i - m\,\sigma_{i+1} 2^{-i-1} y_i \\ y_{i+2} &= y_i + \sigma_i 2^{-i} x_i + \sigma_{i+1} 2^{-i-1} x_i \\ z_{i+2} &= z_i - \sigma_i \alpha_{m,i} - \sigma_{i+1} \alpha_{m,i+1} \end{aligned}$$

$\sigma_i = 0$

$$\begin{aligned} x_{i+1} &= x_i + m \cdot 2^{-2i-1} x_i \\ y_{i+1} &= y_i + m \cdot 2^{-2i-1} y_i \\ z_{i+1} &= z_i \end{aligned}$$

$$\begin{aligned} x_{i+2} &= (1 + m \cdot 2^{-2i-3} + m \cdot 2^{-2i-1}) x_i \\ y_{i+2} &= (1 + m \cdot 2^{-2i-3} + m \cdot 2^{-2i-1}) y_i \\ z_{i+1} &= z_i \end{aligned}$$

Cond (III) $\quad \frac{1}{2}(n+1) < i$

$\sigma_i \neq 0$

$$\begin{aligned} x_{i+1} &= x_i - \sigma_i 2^{-i} y_i \\ y_{i+1} &= y_i + \sigma_i 2^{-i} x_i \\ z_{i+1} &= z_i - \sigma_i \tan^{-1}(2^{-i}) \end{aligned}$$

$$\begin{aligned} x_{i+2} &= x_i - m\,\sigma_i 2^{-i} y_i - m\,\sigma_{i+1} 2^{-i-1} y_i \\ y_{i+2} &= y_i + \sigma_i 2^{-i} x_i + \sigma_{i+1} 2^{-i-1} x_i \\ z_{i+2} &= z_i - \sigma_i \alpha_{m,i} - \sigma_{i+1} \alpha_{m,i+1} \end{aligned}$$

$\sigma_i = 0$

$$\begin{aligned} x_{i+1} &= x_i \\ y_{i+1} &= y_i \\ z_{i+1} &= z_i \end{aligned}$$

– *$0 \le i \le (n-3)/4$* :

use the prediction algorithm, generate σ_i from z_i using Table 1 in the same manner as in Fig. 1, $\sigma_i \in \{-1,1\}$, execute iterations according to Eqs. 1–3

– *$(n-3)/4 < i \le (n+1)/2$* :

use the prediction algorithm, generate σ_i from z_i by special recoding (explained later), $\sigma_i \in \{0,1,-1\}$, after each iteration increment *i* by 2

$\sigma_i <> 0$: $x_{i+2} = x_i - m\,\sigma_i\, 2^{-i}\, y_i - m\,\sigma_{i+1}\, 2^{-i-1}\, y_i$

$y_{i+2} = y_i + \sigma_i 2^{-i}\, x_i + \sigma_{i+1} 2^{-i-1} x_i$

$z_{i+2} = z_i - \sigma_i \alpha_{m,i} - \sigma_{i+1} \alpha_{m,i+1}$

$\sigma_i = 0$: $x_{i+2} = x_i + m\, 2^{-2i-1} x_i + m\, 2^{-2i-2} x_i \quad (11)$

$y_{i+2} = y_i + m\, 2^{-2i-1} y_i + m\, 2^{-2i-2} y_i \quad (12)$

$z_{i+2} = z_i \quad (13)$

$$\begin{bmatrix} x_{i+1} \\ y_{i+1} \end{bmatrix} = \begin{bmatrix} 1 & -m\,\sigma_i 2^{-i} \\ \sigma_i 2^{-i} & 1 \end{bmatrix} \begin{bmatrix} x_i \\ y_i \end{bmatrix}$$

$$\begin{bmatrix} x_{i+2} \\ y_{i+2} \end{bmatrix} = \begin{bmatrix} 1-m\boxed{\sigma_i \sigma_{i+1}}2^{-2i-1} & -m(\sigma_i 2^{-i} + \sigma_{i+1} 2^{-i-1}) \\ (\sigma_i 2^{-i} + \sigma_{i+1} 2^{-i-1}) & 1-m\boxed{\sigma_i \sigma_{i+1}}2^{-2i-1} \end{bmatrix} \begin{bmatrix} x_i \\ y_i \end{bmatrix}$$

$\boxed{\sigma_i = 0} \rightarrow \boxed{\sigma_i \sigma_{i+1} = 0}$

$$\begin{bmatrix} x_{i+2} \\ y_{i+2} \end{bmatrix} = \begin{bmatrix} 1 & -m(\cancelto{0}{\sigma_i 2^{-i}} + \sigma_{i+1} 2^{-i-1}) \\ (\cancelto{0}{\sigma_i 2^{-i}} + \sigma_{i+1} 2^{-i-1}) & 1 \end{bmatrix} \begin{bmatrix} x_i \\ y_i \end{bmatrix}$$

$$\begin{aligned} x_{i+2} &= x_i - m(\sigma_i 2^{-i} + \sigma_{i+1} 2^{-i-1})\, y_i \\ y_{i+2} &= (\sigma_i 2^{-i} + \sigma_{i+1} 2^{-i-1})\, x_i + y_i \end{aligned}$$

$$\begin{aligned} x_{i+2} &= x_i - m\,\sigma_i 2^{-i} y_i - m\,\sigma_{i+1} 2^{-i-1} y_i \\ y_{i+2} &= y_i + \sigma_i 2^{-i} x_i + \sigma_{i+1} 2^{-i-1} x_i \end{aligned}$$

- *$0 \le i \le (n-3)/4$:*
  use the prediction algorithm, generate σ_i from z_i using Table 1 in the same manner as in Fig. 1, $\sigma_i \in \{-1,1\}$, execute iterations according to Eqs. 1–3
- *$(n-3)/4 < i \le (n+1)/2$:*
  use the prediction algorithm, generate σ_i from z_i by special recoding (explained later), $\sigma_i \in \{0,1,-1\}$, after each iteration increment i by 2

$\sigma_i <> 0$:
$$\begin{aligned} x_{i+2} &= x_i - m\,\sigma_i\, 2^{-i} y_i - m\,\sigma_{i+1}\, 2^{-i-1} y_i \\ y_{i+2} &= y_i + \sigma_i 2^{-i} x_i + \sigma_{i+1} 2^{-i-1} x_i \\ z_{i+2} &= z_i - \sigma_i \alpha_{m,i} - \sigma_{i+1} \alpha_{m,i+1} \end{aligned}$$

$\sigma_i = 0$:
$$x_{i+2} = x_i + m\, 2^{-2i-1} x_i + m\, 2^{-2i-2} x_i \quad (11)$$
$$y_{i+2} = y_i + m\, 2^{-2i-1} y_i + m\, 2^{-2i-2} y_i \quad (12)$$
$$z_{i+2} = z_i \quad (13)$$

$$\lambda(-1) = 0$$

$$\lambda(1) = 0$$

$$\lambda(0) = 1$$

$$\lambda(\sigma_i) = 0 \quad \sigma_i \in \{1, -1\}$$

$$= 1 \quad \sigma_i = 0$$

Modified CORDIC

Timmermann 1989 Electronics Letters

$$x_n = k_m \{ x_0 \cos[\sqrt{(m)}\,\alpha] - \sqrt{(m)}\, y_0 \sin[\sqrt{(m)}\,\alpha] \}$$

$$y_n = k_m \{ 1/\sqrt{(m)}\, x_0 \sin[\sqrt{(m)}\,\alpha] + y_0 \cos[\sqrt{(m)}\,\alpha] \}$$

$$z_n = z_0 + \alpha$$

k_m : the scaling factor

m : the coordinate system $(0, -1, +1)$

α : the rotation angle

$\begin{pmatrix} x_0 \\ y_0 \\ z_0 \end{pmatrix}$: the initial values depends on the iteration goal

Data dependency across iteration

$\rightarrow$ CSA no benefit

1st half iterations : the most significant contribution

the rotation angle $\alpha_i = \frac{1}{\sqrt{(m)}} \tan^{-1}[\sqrt{(m)}\, 2^{-S(m,i)}]$

$S(m, i)$ the iteration shift values

α_i decreases with the increasing iteration index i

2nd half iterations : can improve the accuracy only by one bit each rotation $z_n \to 0$

$$x_n = k_m \{ x_0 \cos[\sqrt{(m)}\,\alpha] - \sqrt{(m)}\, y_0 \sin[\sqrt{(m)}\,\alpha] \}$$
$$y_n = k_m \{ 1/\sqrt{(m)}\, x_0 \sin[\sqrt{(m)}\,\alpha] + y_0 \cos[\sqrt{(m)}\,\alpha] \}$$
$$z_n = z_0 + \alpha$$

Vectoring $y_n \to 0$

$$x_n = k_m \sqrt{x_0^2 + m y_0^2}$$
$$z_n = z_0 + 1/\sqrt{(m)} \tan^{-1}[\sqrt{(m)}\, y_0 / x_0]$$

2nd half iterations : can improve the accuracy
only by one bit each

replace these iterations by a single rotation
after the remaining rotation angle
has been reduced using a fixed number of
pure CORDIC iterations

this truncation reduces the latency time and saves area
although the truncation requires extra hardware

the necessary minimum number of iterations

Rotation mode $(z \to 0)$

after j CORDIC rotations have been performed
the z path contains the remaining rotation angle z_j

$$\begin{bmatrix} x_n \\ y_n \end{bmatrix} = \begin{bmatrix} \cos[\sqrt{(m)}\, z_j] & -\sqrt{(m)}\, \sin[\sqrt{(m)}\, z_j] \\ 1/\sqrt{(m)}\, \sin[\sqrt{(m)}\, z_j] & \cos[\sqrt{(m)}\, z_j] \end{bmatrix} \begin{bmatrix} x_j \\ y_j \end{bmatrix}$$

$$x_n = k_m \{ x_0 \cos[\sqrt{(m)}\, \alpha] - \sqrt{(m)}\, y_0 \sin[\sqrt{(m)}\, \alpha] \}$$
$$y_n = k_m \{ 1/\sqrt{(m)}\, x_0 \sin[\sqrt{(m)}\, \alpha] + y_0 \cos[\sqrt{(m)}\, \alpha] \}$$
$$z_n = z_0 + \alpha$$

assume $k_m = 1$

$$\begin{bmatrix} x_n \\ y_n \end{bmatrix} = \begin{bmatrix} \cos[\sqrt{(m)}\, z_j] & -\sqrt{(m)}\ \sin[\sqrt{(m)}\, z_j] \\ 1/\sqrt{(m)}\ \sin[\sqrt{(m)}\, z_j] & \cos[\sqrt{(m)}\, z_j] \end{bmatrix} \begin{bmatrix} x_j \\ y_j \end{bmatrix}$$

Taylor Series expansions to $\sin\theta$, $\cos\theta$
take only the first terms

$$\sin\theta = x - \frac{1}{3!}x^3 + \frac{1}{5!}x^5 - \frac{1}{7!}x^7 + \cdots$$

$$\cos\theta = 1 - \frac{1}{2!}x^2 + \frac{1}{4!}x^4 - \frac{1}{6!}x^6 + \cdots$$

$$\begin{bmatrix} x_n \\ y_n \end{bmatrix} = \begin{bmatrix} 1 & -\sqrt{(m)}\cdot\sqrt{(m)}\, z_j \\ 1/\sqrt{(m)}\cdot\sqrt{(m)}\, z_j & 1 \end{bmatrix} \begin{bmatrix} x_j \\ y_j \end{bmatrix}$$

$$\begin{bmatrix} x_n \\ y_n \end{bmatrix} = \begin{bmatrix} 1 & -m\, z_j \\ z_j & 1 \end{bmatrix} \begin{bmatrix} x_j \\ y_j \end{bmatrix}$$

for a sufficiently small z_j

the required precision of n-bit
the upper limit on the maximal remainder

$$z_j \leqslant \frac{1}{\sqrt{(m)}} \tan^{-1}[\sqrt{(m)}\, 2^{-j+1}]$$

$$\begin{bmatrix} x_n \\ y_n \end{bmatrix} = \begin{bmatrix} 1 & -m z_j \\ z_j & 1 \end{bmatrix} \begin{bmatrix} x_j \\ y_j \end{bmatrix}$$

for a sufficiently small z_j

the required precision of n-bit
the upper limit on the maximal remainder

$$\frac{1}{2} z_j^2 \leqslant 2^{-n} \qquad z_j^{\,2} \leqslant 2^{-n+1} \qquad z_j \leqslant 2^{\frac{-n+1}{2}}$$

$$z_j \leqslant \frac{1}{\sqrt{(m)}} \tan^{-1}\left[\sqrt{(m)}\, 2^{-j+1}\right]$$

$$j > \frac{n+1}{2}$$

Rotation mode

$$x_n = x_j - m z_j y_j \qquad (j > (n+1)/2)$$
$$y_n = z_j x_j + y_j \qquad (j > (n+1)/2)$$

Vectoring mode

$$x_n = z_j \qquad j > (n+1)/2$$
$$z_n = z_j + y_j/x_j \qquad j > (n/3) + 0.472$$

the Truncated CORDIC Algorithm

- reduces the number of CORDIC iterations

- multiplication / division hardware

  Booth Technique halves the amount of partial products
  Carry Save Architecture

$k_m \neq 1 \Rightarrow$ multiplication $\Rightarrow$ multiplier anyway