## Applicatives Methods (3B)

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#### The definition of Applicative

```
class (Functor f) => Applicative f where

pure :: a -> f a

(<*>) :: f (a -> b) -> f a -> f b
```

The class has a two methods:

pure brings arbitrary values into the functor

(<\*>) takes a <u>function</u> wrapped in a functor f and a <u>value</u> wrapped in a functor f and returns the <u>result</u> of the <u>application</u> which is also wrapped in a functor f

#### The Maybe instance of Applicative

```
instance Applicative Maybe where

pure = Just

(Just f) <*> (Just x) = Just (f x)

_ <*> _ = Nothing
```

```
pure wraps the value with Just;
```

```
(<*>) applies
the <u>function</u> wrapped in <u>Just</u>
to the <u>value</u> wrapped in <u>Just</u> if both exist,
and results in <u>Nothing</u> otherwise.
```

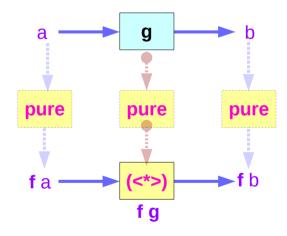
#### An Instance of the Applicative Typeclass

f : Functor, Applicative

(Functor f) => Applicative f

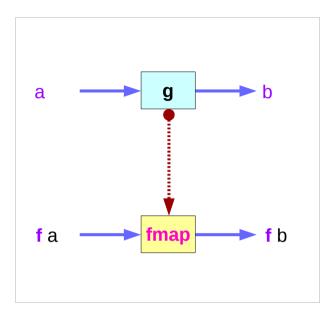
# instance Applicative Maybe where pure = Just Nothing <\*> \_ = Nothing (Just <u>f</u>) <\*> something = fmap <u>f</u> something

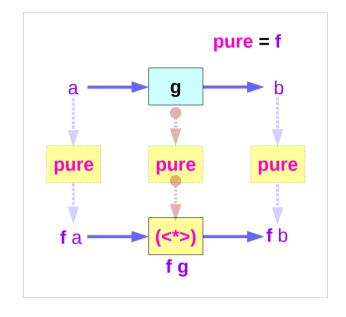
 $\underline{\mathbf{f}}$ : function in a context

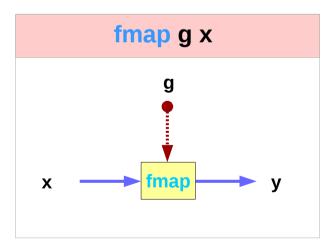


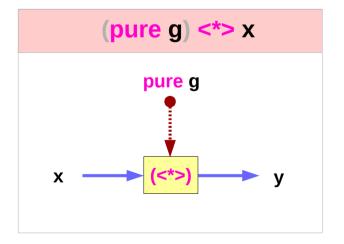
(Functor f) => Applicative f

## fmap g x = (pure g) <\*> x







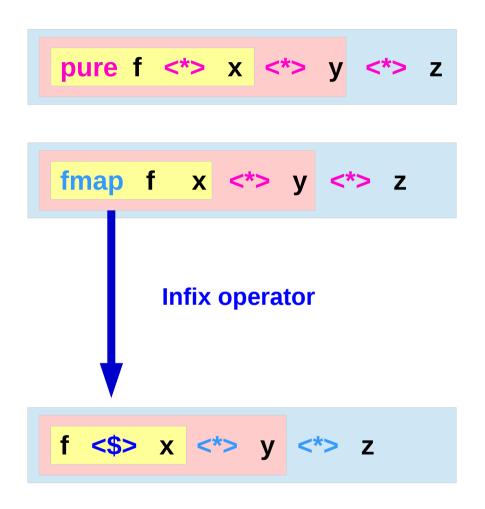


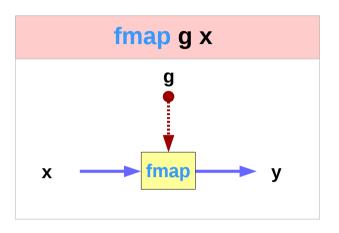
http://learnyouahaskell.com/functors-applicative-functors-and-monoids

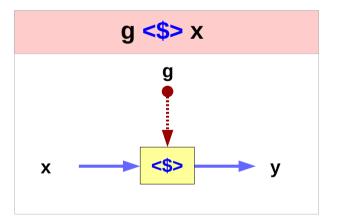
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#### Left Associative <\*>

## Infix Operator <\$>



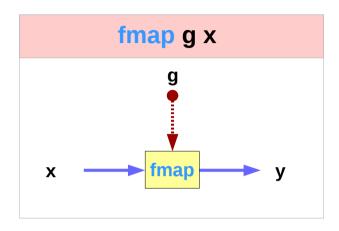


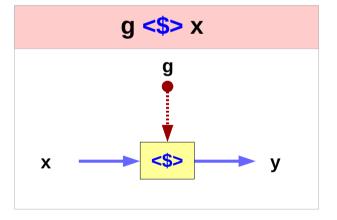


#### Infix Operator <\$>: not a class method

```
class (Functor f) => Applicative f where
  pure :: a -> f a
  (<*>) :: f (a -> b) -> f a -> f b
```

#### not a class method





#### The Applicative Typeclass

Applicative is a <u>superclass</u> of Monad. every Monad is also a Functor and an Applicative fmap, pure, (<\*>) can all be used with monads.

a **Monad** instance also requires **Functor** and **Applicative** instances.

the types and roles of return and (>>)

## (\*> v.s. >>) and (pure v.s. return)

```
(*>) :: Applicative f \Rightarrow fa \Rightarrow fb \Rightarrow fb
```

$$(>>) :: Monad m => m a -> m b -> m b$$

the constraint changes from **Applicative** to **Monad**.

## The Applicative Laws

The identity law: pure id <\*>v = v

Homomorphism: pure f <\*> pure x = pure (f x)

Interchange:  $u \le pure y = pure (\$ y) \le u$ 

Composition: u <\*> (v <\*> w) = pure (.) <\*> u <\*> v <\*> w

#### The Identity Law

#### The identity law

pure id <\*> v = v

pure to inject <u>values</u> into the functorin a default, featureless way,so that the result is as close as possible to the <u>plain</u> value.

applying the **pure id** morphism does nothing, exactly like with the plain **id** function.

#### The Homomorphism Law

#### The homomorphism law

pure f <\*> pure x = pure (f x)

applying a "pure" <u>function</u> to a "pure" <u>value</u> is the same as applying the function to the <u>value</u> in the normal way and then using **pure** on the result.

means **pure** preserves function application.

#### applying a non-effectful function f

to a <u>non-effectful</u> argument x in an <u>effectful</u> context **pure** is the same as just **applying** the function f to the argument x and then injecting the result (f x) into the <u>context</u> with **pure**.

#### The Interchange Law

#### The interchange law

```
applying a morphism u to a "pure" value pure y is the same as applying pure ($ y) to the morphism u
```

(\$ y) is the function that supplies y as <u>argument</u> to another function

the higher order functions

when evaluating the application of an <u>effectful function</u> **u** to a <u>pure argument pure</u> **y**, the <u>order</u> in which we evaluate the <u>function</u> **u** and its <u>argument pure</u> **y** <u>doesn't matter</u>.

#### The Composition Law

The composition law

pure (.) composes morphisms similarly

(f.g) x = f(gx)

to how (.) composes functions:

= pure f <\*> (pure g <\*> pure x)

u = pure f

v = pure g

w = pure x

applying the composed morphism pure (.) <\*> u <\*> v to w

gives the same result as applying **u** 

to the result of applying **v** to **w** 

(v <\*> w)

it is expressing a sort of associativity property of (<\*>).

## <\$> related operators

Functor map <\$>

(<\$>) :: Functor f => (a -> b) -> f a -> f b

(<\$) :: Functor f => a -> f b -> f a

(\$>) :: Functor f => f a -> b -> f b

The **<\$>** operator is just a synonym for the **fmap** function from the Functor typeclass.

This function generalizes the **map** function for lists to many other data types, such as **Maybe**, **IO**, and **Map**.

#### <\$> exammples

```
#!/usr/bin/env stack
-- stack --resolver ghc-7.10.3 runghc
import Data.Monoid ((<>))

main :: IO ()
main = do
    putStrLn "Enter your year of birth"
    year <- read <$> getLine
    let age :: Int
        age = 2020 - year
    putStrLn $ "Age in 2020: " <> show age
```

## <\$, \$> operators

In addition, there are two additional operators provided which replace a value inside a Functor instead of applying a function.

This can be both more convenient in some cases, as well as for some Functors be more efficient.

$$x < y = y > x$$

$$x > y = y < x$$

#### <\*> related operators

Applicative function application <\*>

(<\*>) :: Applicative f => f (a -> b) -> f a -> f b

(\*>) :: Applicative f => f a -> f b -> f b

(<\*) :: Applicative f => f a -> f b -> f a

Commonly seen with <\$>, <\*> is an operator that applies a wrapped function to a wrapped value. It is part of the Applicative typeclass, and is very often seen in code like the following:

foo <\$> bar <\*> baz

#### <\*> examples

For cases when you're dealing with a Monad, this is equivalent to:

```
do x <- bar
y <- baz
return (foo x y)</pre>
```

Other common examples including parsers and serialization libraries.

Here's an example you might see using the aeson package:

```
data Person = Person { name :: Text, age :: Int } deriving Show
```

-- We expect a JSON object, so we fail at any non-Object value.

instance FromJSON Person where

```
parseJSON (Object v) = Person <$> v .: "name" <*> v .: "age"
parseJSON _ = empty
```

#### \*> operator

To go along with this, we have two helper operators that are less frequently used:

\*> ignores the value from the first argument. It can be defined as:

Or in do-notation:

For Monads, this is completely equivalent to >>.

## <\* operator

<\* is the same thing in reverse: perform the first action then the second,

but only take the value from the first action.

Again, definitions in terms of <\*> and do-notation:

res <- a1

\_ <- a2

return res

liftA2 :: (a -> b -> c) -> f a -> f b -> f c

Lift a binary function to actions.

Some functors support an implementation of liftA2 that is more efficient than the default one.

In particular, if fmap is an expensive operation, it is likely better to use liftA2 than to fmap over the structure and then use <\*>.

http://hackage.haskell.org/package/base-4.10.1.0/docs/Control-Applicative.html#v:liftA2

If you have the variables

```
f :: a -> b -> c
a :: f a
```

b :: f b

you can combine them in the following ways with the same result of type f  $\ensuremath{\text{c}}$  .

But how to cope with let and sharing in the presence of effects?

https://wiki.haskell.org/Applicative\_functor

Consider the non-functorial expression:

```
x :: x
g :: x -> y
h :: y -> y -> z
let y = g x
in h y y
```

Very simple. Now we like to generalize this to

```
fx :: f x
fg :: f (x -> y)
fh :: f (y -> y -> z)
```

https://wiki.haskell.org/Applicative\_functor

However, we note that

```
let fy = fg <*> fx
in fh <*> fy <*> fy
```

runs the effect of fy

twice. E.g. if fy

writes something to the terminal then fh <\*> fy <\*> fy

writes twice. This could be intended, but how can we achieve, that the effect is run only once and the result is used twice? Actually, using the liftA

commands we can pull results of applicative functors into a scope where we can talk exclusively about functor results and not about effects. Note that functor results can also be functions. This scope is simply a function, which contains the code that we used in the non-functorial setting.

#### liftA3

```
(x g h \rightarrow let y = g x in h y y)
fx fg fh
```

The order of effects is entirely determined by the order of arguments to liftA3

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https://wiki.haskell.org/Applicative\_functor

#### References

- [1] ftp://ftp.geoinfo.tuwien.ac.at/navratil/HaskellTutorial.pdf
- [2] https://www.umiacs.umd.edu/~hal/docs/daume02yaht.pdf