Applicative Sequencing (3C)

Young Won Lim 5/26/18 Copyright (c) 2016 - 2018 Young W. Lim.

Permission is granted to copy, distribute and/or modify this document under the terms of the GNU Free Documentation License, Version 1.2 or any later version published by the Free Software Foundation; with no Invariant Sections, no Front-Cover Texts, and no Back-Cover Texts. A copy of the license is included in the section entitled "GNU Free Documentation License".

Please send corrections (or suggestions) to youngwlim@hotmail.com.

This document was produced by using LibreOffice.

Young Won Lim 5/26/18 http://learnyouahaskell.com/making-our-own-types-and-typeclasses#the-functor-typeclass

http://learnyouahaskell.com/functors-applicative-functors-and-monoids

Haskell in 5 steps https://wiki.haskell.org/Haskell_in_5_steps

<\$> related operators

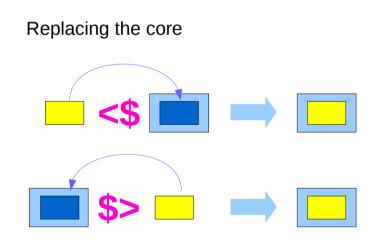
Functor map <\$>

- (<\$>) :: Functor f => (a -> b) -> f a -> f b
- (<\$) :: Functor f => a -> f b -> f a
- (\$>) :: Functor f => f a -> b -> f b

replace b in f b with a ... f a replace a in f a with b ... f b

The <\$> operator is just a synonym for the fmap function in the Functor typeclass.

fmap generalizes **map** for **lists** to other data types : **Maybe**, **IO**, **Map**.



4

https://haskell-lang.org/tutorial/operators

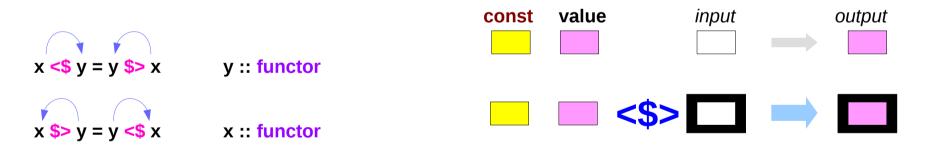
Young Won Lim 5/26/18

<**\$ / <\$> / \$>** operators

there are two additional operators provided which <u>replace</u> a **value** inside a Functor instead of applying a function. This can be both more convenient in some cases,

as well as for some Functors be more efficient.

value <\$ functor = const value <\$> functor
functor \$> value = const value <\$> functor



App	licatives	
Sequ	uencing	(3C)

<\$ / <\$> / \$> operators examples

import Data.Functor	import Data.Functor
Prelude> Just 1 \$> 2	Prelude> (+1) <\$> Just 2
Just 2	Just 3
Prelude> Just 2 \$> 1	Prelude> (+1) <\$> Just 3
Just 1	Just 4
Prelude> 1 <\$ Just 3	Prelude> (+1) <\$> Nothing
Just 1	Nothing
Prelude> 3 <\$ Just 1	
Just 3	
Prelude> 1 <\$ Just 3	Prelude> const 2 <\$> Just 111
Just 1	Just 2
Prelude> 3 <\$ Just 1	
Just 3	

https://www.schoolofhaskell.com/school/to-infinity-and-beyond/pick-of-the-week/Simple%20examples

Applicatives	
Sequencing	(3C)

6

<\$> examples

```
#!/usr/bin/env stack-- stack --resolver ghc-7.10.3 runghcimport Data.Monoid ((<>))
```

main :: IO ()

main = do

putStrLn "Enter your year of birth"

year <- read <\$> getLine

let age :: Int

age = 2020 - year

putStrLn \$ "Age in 2020: " <> show age

getLine :: IO String

Input: read "12"::Double Output: 12.0

-- this infix synonym for mappend is found in Data.Monoid x <> y = mappend x y infixr 6 <>

https://haskell-lang.org/tutorial/operators

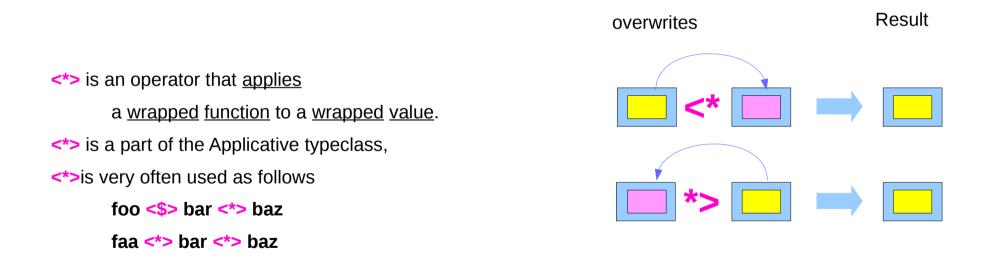
Applicatives Sequencing (3C)

Young Won Lim 5/26/18

<*> related operators

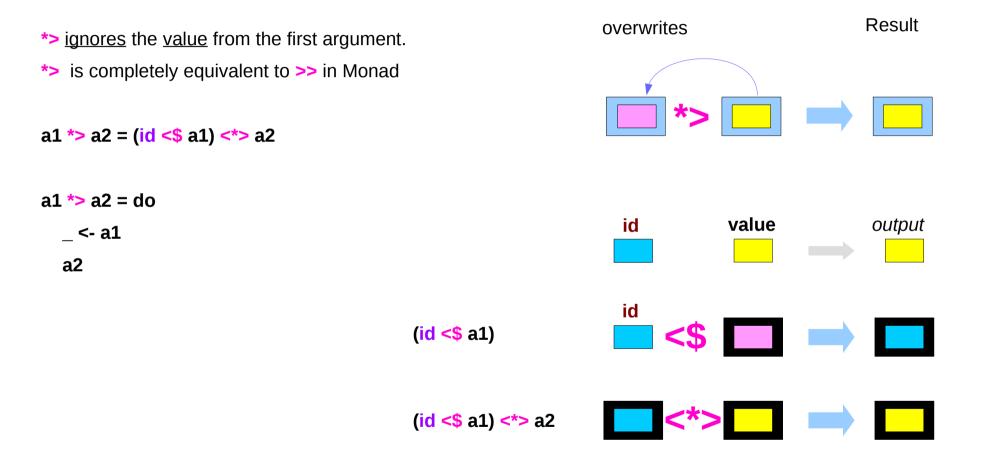
Applicative function application <*>

- (<*>) :: Applicative f => f (a -> b) -> f a -> f b
- (*>) :: Applicative f => f a -> f b -> f b
- (<*) :: Applicative f => f a -> f b -> f a



***>** operator

two helper operators

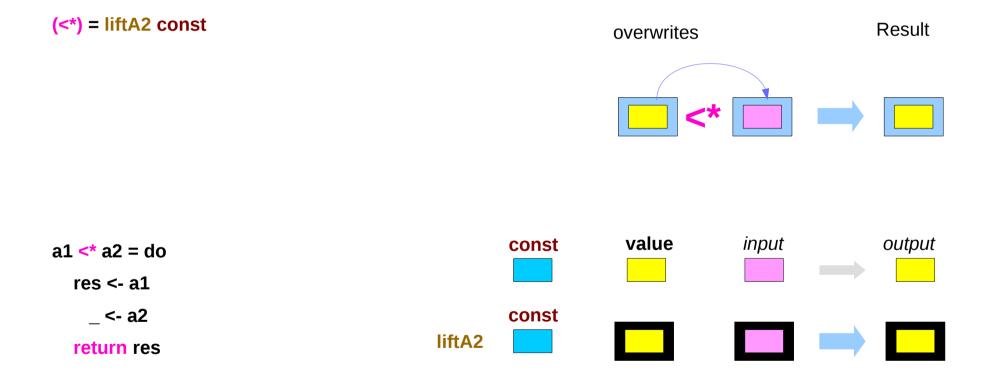


https://haskell-lang.org/tutorial/operators

Applicatives Sequencing (3C)

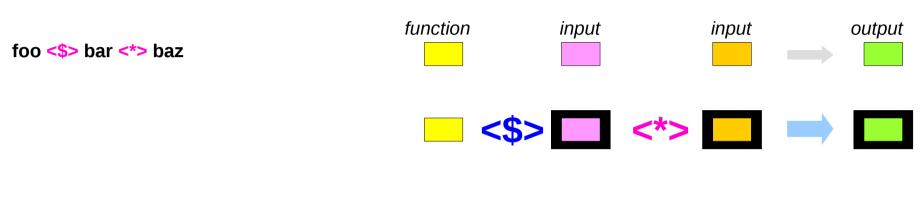
<* operator

<* is the same thing in reverse: <u>perform</u> the <u>first action</u> then the <u>second</u>, but only <u>take</u> the <u>value</u> from the <u>first</u> action.

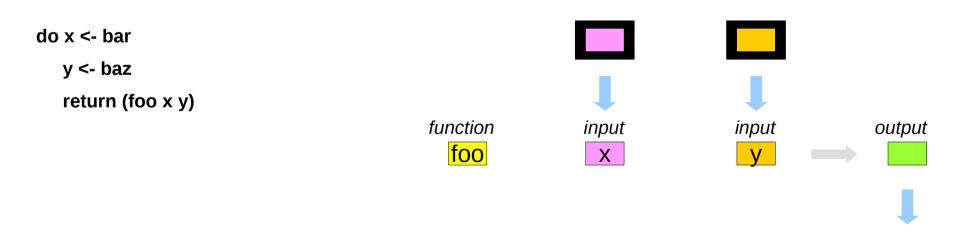


Applicatives	
Sequencing	(3C)





With a Monad, this is equivalent to:



Applicatives	
Sequencing	(3C)

<*> examples

```
examples including parsers and serialization libraries.
using the aeson package: (handling JSON data)
```

data Person = Person { name :: Text, age :: Int } deriving Show

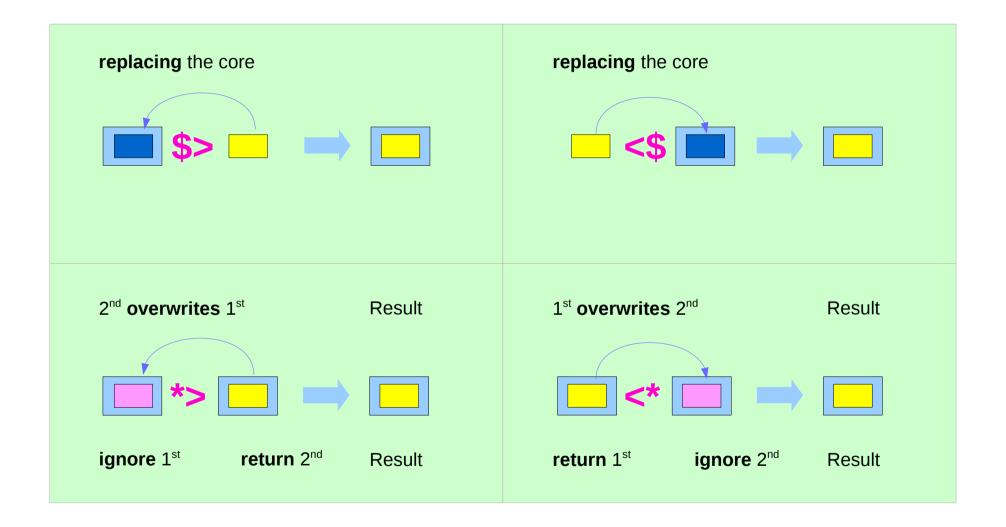
-- We expect a JSON object, so we fail at any non-Object value.

instance FromJSON Person where

parseJSON (Object v) = Person <\$> v .: "name" <*> v .: "age"
parseJSON _ = empty

- : append-head operator (cons)
- . function composition operators
- . name qualifier

(\$> v.s. <\$) and (*> v.s. <*)

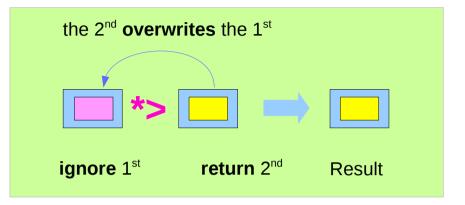


https://en.wikibooks.org/wiki/Haskell/Applicative_functors

Applicatives Sequencing (3C)

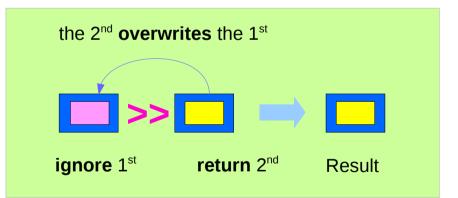
(*> v.s. >>) and (pure v.s. return)

(*>) ::	Applicative f	=>	f a ->	f b ->	f b
(>>) ::	Monad m	=>	<mark>m</mark> a ->	m b ->	m b
pure ::	Applicative f	=>	a ->	fa	
return ::	Monad m	=>	a ->	m a	



the constraint changes from Applicative to Monad.

(*>) in <u>Applicative</u> (>>) in <u>Monad</u>
pure in <u>Applicative</u> (>>) in <u>Monad</u>



https://en.wikibooks.org/wiki/Haskell/Applicative_functors

Applicatives Sequencing (3C)

Sequencing of Effects

commutative monads in Haskell,

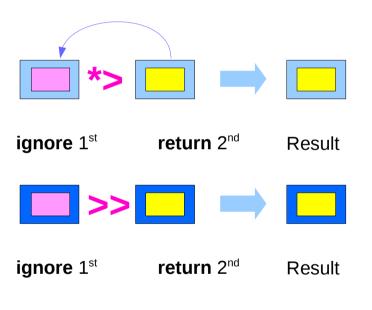
the concept involved is the same, only specialised to Monad.

Commutativity (or the lack thereof) affects
other functions which are derived from (<*>) as well.
(*>) is a clear example:

(*>) :: Applicative f => f a -> f b -> f b

(*>) <u>combines</u> effects while <u>preserving</u> only the values of its <u>second</u> argument.

For **monads**, it is equivalent to (>>). Here is a demonstration of it using Maybe, which is commutative:



Sequencing of Effects

Prelude> [(2*),(3*)] <*> [4,5]	
1) [8,10,12,15] correct	answer
2) [8,12,10,15]	
The difference is that for the first (and	l correct) answer
the result is obtained	
by taking the skeleton of the first list	
and replacing each element	(2*), (3*)
by all possible combinations	[(2*) <*> 4, (2*) <*> 5, (3*) <*> 4, (3*) <*> 5]
with elements of the second list,	
while for the other possibility	
the starting point is the second list.	

[(2*) <*> 4, (3*) <*> 4, (2*) <*> 5, (3*) <*> 5]

sequencing effects

Non-commutative Functors

by **effects** we mean the functorial **context**, as opposed to the **values** within the functor

some examples:

the <u>skeleton</u> of a list, <u>actions</u> performed in the real world in IO, the <u>existence</u> of a value in Maybe

The existence of two legal implementations of (<*>) for lists only differ in the sequencing of effects
[] is a non-commutative applicative functor.

Prelude> [(2*),(3*)] <*> [4,5]

- 1) [8,10,12,15]
- 2) [8,12,10,15]

Commutative Functors

a commutative applicative functor is one for which the following holds:

liftA2 f u v = liftA2 (flip f) v u -- Commutativity

Or, equivalently,

f <\$> u <*> v = flip f <\$> v <*> u

https://en.wikibooks.org/wiki/Haskell/Applicative_functors

Applicatives Sequencing (3C)

Commutative Monads (1)

do	do
a <- actA	b <- actB
b <- actB	a <- actA
return (a + b)	return (a + b)

commutative if the order of side effects is not important.

there are many monads that commute (e.g. Maybe, Random). If the monad is commutative, then the operations captured within it can be computed in parallel.

No good syntax for monads that commute still an open research problem

https://stackoverflow.com/questions/5897845/relax-ordering-constraints-in-monadic-computation

Commutative Monads (2)

Commutative monads are monads for which the order of actions makes no difference (they commute), that is when following code:

do	do	
a <- actA	b <- actB	
b <- actB	a <- actA	
m a b	m a b	

commutative if the order of side effects is not important.

Examples of commutative include:

Reader monad Maybe monad

https://wiki.haskell.org/Monad#Commutative_monads https://stackoverflow.com/questions/6089997/how-do-i-find-out-whether-a-monad-is-commutative

Applicatives Sequencing (3C)



Left-to-right sequencing

The convention in Haskell is to always implement (<*>) and other applicative operators using **left-to-right sequencing**.

Even though this convention helps reducing confusion, it also means appearances sometimes are misleading.

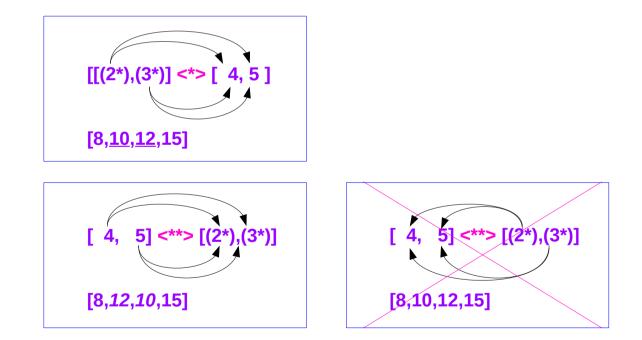
For instance, the (<*) function is <u>not flip</u> (*>), as it sequences effects <u>from left to right</u> just like (*>):



(<**>) :: Applicative f => f a-> f (a -> b)-> f b(<*>) :: Applicative f => f(a -> b)-> f a-> f b

from **Control.Applicative** not flip (<*>)

a way of inverting the sequencing



https://en.wikibooks.org/wiki/Haskell/Applicative_functors

Applicatives Sequencing (3C)

Sequencing examples (1)

```
Prelude> [(2*),(3*)] <*> [4,5]
[8,10,12,15]
Prelude> [4,5] <**> [(2*),(3*)]
[8,12,10,15]
Prelude> Just 2 *> Just 3
Just 3
Prelude> Just 3 *> Just 2
Just 2
Prelude> Just 2 *> Nothing
Nothing
Prelude> Nothing *> Just 2
Nothing
```

[(2*)] <*> [4,5], [(3*)] <*> [4,5]

[4] <**> [(2*),(3*)], [5] <**> [(2*),(3*)]

Sequencing examples (2)

```
Prelude> (print "foo" *> pure 2) *> (print "bar" *> pure 3)

"foo"

"bar"

3

Prelude> (print "bar" *> pure 3) *> (print "foo" *> pure 2)

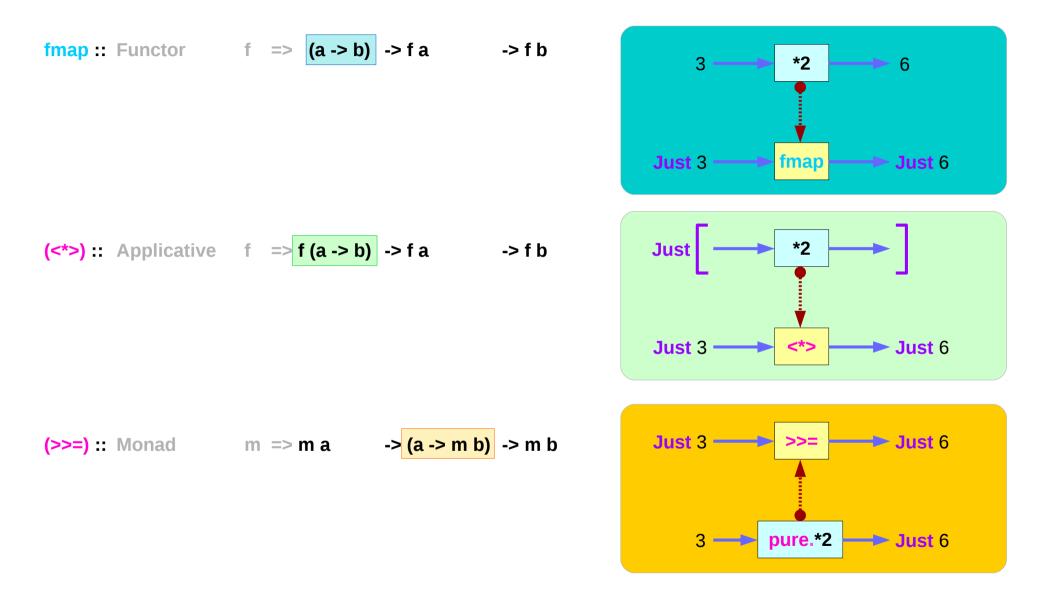
"bar"

"foo"

2
```

```
Prelude> <u>(print "foo" *> pure 2)</u> <* (print "bar" *> pure 3)
"foo"
"bar"
2
```

Functors, Applicative, and Monad



https://en.wikibooks.org/wiki/Haskell/Applicative_functors

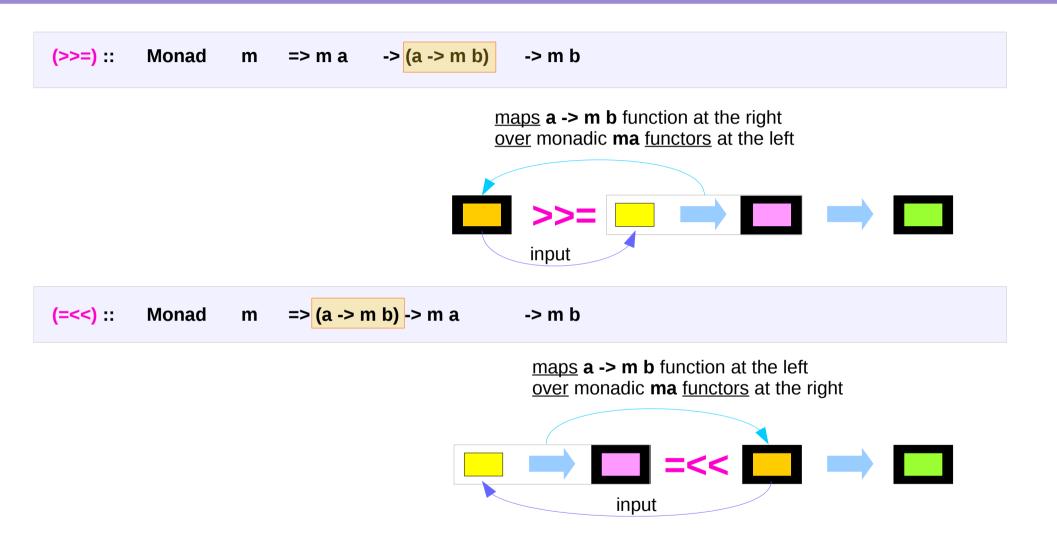
Applicatives Sequencing (3C)

25

Functors, Applicative, and Monad Examples

fmap :: Functor	f => (a -> b)	-> f a	-> f b	Prelude> <mark>fmap (*2) (Just 3)</mark> Just 6
(<*>) :: Applicative	f => <mark>f(a->b)</mark>	-> f a	-> f b	Prelude> (Just (*2)) <*> (Just 3) Just 6
(>>=) :: Monad	m => m a	-> <mark>(a -> m b)</mark>	-> m b	Prelude> (Just 3) >>= (pure . (*2)) Just 6 Prelude> (Just 3) >>= (return . (*2)) Just 6

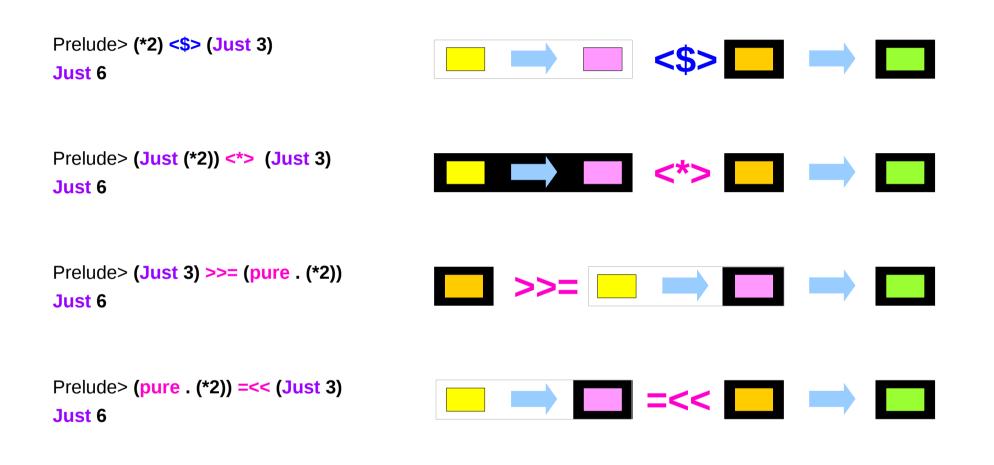
(=<<) : the flipped version of (>>=)



https://en.wikibooks.org/wiki/Haskell/Applicative_functors

Applicatives Sequencing (3C)

<\$>, <*>, >>=, and =<< examples</p>



Comparing the three characteristic methods

replace fmap by its infix synonym, (<\$>) replace (>>=) by its flipped version, (=<<)

fmap :: Functor	f	=>	(a -> b)	->	fa	-> f b
<pre>(<*>) :: Applicative</pre>	f	=>	f (a -> b)	->	fa	-> f b
(>>=) :: Monad	m	=>	m a	->	(a -> m b)	-> m b
(<\$>) :: Functor	t	=>	(a -> b))	-> <mark>(t a -></mark>	<mark>> t b)</mark>
<pre>(<*>) :: Applicative</pre>	t	=>	t (a -> b))	-> <mark>(t a -></mark>	<mark>> t b)</mark>
(=<<) :: Monad	t	=>	(a -> t b))	-> <mark>(t a -></mark>	<mark>> t b)</mark>

All mapping functions over Functors

fmap, (<*>) and (=<<) are all mapping functions over Functors.

The differences between them are in what is being mapped (functions) over in each case:

(<\$>) :: Functor t	=>	(a -> b)	->	(t a -> t b)
(<*>) :: Applicative t	=>	t (a -> b)	->	(t a -> t b)
(=<<) :: Monad t	=>	(a -> t b)	->	(t a -> t b)

fmap maps (a -> b) arbitrary functions <u>over functors</u>.
(<*>) maps t (a -> b) morphisms <u>over</u> (applicative) <u>functors</u>.
(=<<) maps a -> t b functions <u>over</u> (monadic) <u>functors</u>.

From fmap to (<*>) and then to (>>=)

The differences of **Functor**, **Applicative** and **Monad** follow from what these three mapping functions allow you to do.

As you move from fmap to (<*>) and then to (>>=), you gain in power, versatility and control, at the cost of guarantees about the results.

We will now slide along this scale. While doing so, we will use the contrasting terms

> **values** to refer to <u>plain values</u> within a functor **contexts** to <u>whatever surrounds</u> them, respectively.

No changing the context

The type of **fmap** ensures that it is impossible to use it to **change the context**, no matter which function it is given.

In (a -> b) -> t a -> t b, the (a -> b) function has <u>nothing to do with</u> the t <u>context</u> of the t a functorial <u>value</u>, and so applying it <u>cannot affect</u> the <u>context</u>. For that reason, if you do fmap f xs on some list xs the number of elements of the list will never change.

Changing the context

fmap cannot change the context

the **(a -> b)** function has no relation with the **t context** the application of this function does <u>not affect</u> the **context t** for example, the <u>number</u> of <u>elements</u> of the list will never change

Prelude> fmap (2*) [2,5,6] [4,10,12] a list with 3 elements a list with 3 elements

That could be a safety guarantee or an unfortunate restriction depending on your purpose

(<*>) is clearly able to <u>change</u> the **context**:

Prelude> [(2*),(3*)] <*> [2,5,6] [4,10,12,6,15,18]

two lists each with 3 elements a list with 6 elements

Carrying a context

The t (a -> b) morphism <u>carries</u> a context of its own,

which is combined (applied) with the **context** of the **t** a functorial **value**.

(<*>), however, is subject to a more subtle <u>restriction</u>

while **t** (**a** -> **b**) morphisms <u>carry</u> **context**, within them there are <u>plain</u> (**a** -> **b**), which are still <u>unable</u> to <u>modify</u> the **context**.

this means the changes to the context (<*>) performsare fully determined by the context of its arguments,t (a->b) or t band the values have no influence over the resulting context.(a->b) or a

Prelude> [(2*),(3*)] <*> [2,5,6] [4,10,12,6,15,18] two lists each with 3 elements a list with 6 elements

Carrying a context examples

```
      Prelude> (print "foo" *> pure (2*)) <*> (print "bar" *> pure 3)
      (pure (2*)) <*> (pure 3)

      "foo"

      "bar"

      6

      Prelude> (print "foo" *> pure 2) *> (print "bar" *> pure 3)
      (pure 2) *> (pure 3)

      "foo"

      "bar"

      3

      Prelude> (print "foo" *> pure undefined) *> (print "bar" *> pure 3)

      "foo"

      "foo"

      "bar"

      3

      Prelude> (print "foo" *> pure undefined) *> (print "bar" *> pure 3)

      "foo"

      "foo"

      "bar"

      3

      "foo"

      "bar"

      3

      "foo"

      "bar"

      3

  </tbr>
```

Creating a context

Prelude> [(2*),(3*)] <*> [2,5,6] [4,10,12,6,15,18] two lists each with 3 elements a list with 6 elements

with **list (<*>)** you know that the <u>length</u> of the resulting list will be the <u>product</u> of the <u>lengths</u> of the original lists,

with **IO** (<*>) you know that all real world effect will happen as long as the <u>evaluation</u> <u>terminates</u>, and so forth.

with Monad, however, it is very different
(>>=) takes a (a -> t b) function, and so it is able
to create context from values
which means a lot of <u>flexibility</u>:

a -> t b creaing context t

Creating a context examples

```
Prelude> [1,2,5] >>= \x -> replicate x x
[1,2,2,5,5,5,5,5]
```

Prelude> [1,2,5] >>= \x -> replicate x x [replicate 1 1, replicate 2 2, replicate 5,5]

```
Prelude> [0,0,0] >>= \x -> replicate x x [ replicate 0 0, replicate 0 0, replicate 0,0 ]
[]
```

Prelude> return 3 >>= \x -> print \$ if x < 10 then "Too small" else "OK"

"Too small"

Prelude> return 42 >>= \x -> print \$ if x < 10 then "Too small" else "OK" "OK"

Flexibility

the extra flexibility

the less guarantees about

- whether your functions are able to unexpectedly erase parts of a data structure for pathological inputs
- whether the control flow in your application remains intelligible

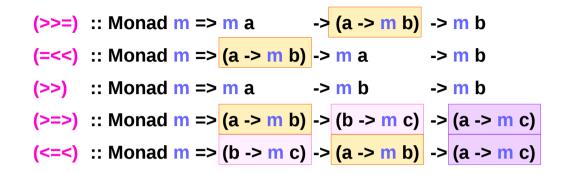
performance implications

• the <u>complex</u> data dependencies of monadic codes might <u>prevent</u> refactoring and optimizations.

it is a good idea to only use <u>as much power as</u> needed for the task at hand.

If you do need the extra capabilities, use **Monad** however, it is often worth it to check whether **Applicative** or **Functor** are sufficient.

Monadic binding / composition operators



Monadic binding operators (1)

(>>=)	::	Monad	m =>	m a		-> (a	a -> m b)	-> m b
(=<<)	::	Monad	m =>	(a ->	• m b)	-> n	1 a	-> m b
(>>)	::	Monad	m =>	m a		-> n	ו b	-> m b

monadic binding operators

The two most basic are >>= and >>

>>=, >>, =<< can be expressed in **do-notation**

>> is just a <u>synonym</u> for *> from the **Applicative** class

=<< is just >>= with the <u>arguments</u> <u>reversed</u>

Monadic binding operators (2)

(>>=) :: Monad m	=> m a	-> m b
(=<<) :: Monad m	=> (a -> m b) -> m a	-> m b
(>>) :: Monad m	=> m a -> m b	-> m b
m1 >>= func = do		m1 :: m a
x <- m1	extract the value	x :: a
func x		func :: a -> m b
m1 >> m2 = do		m1 :: m a
_<- m1	side effect only, igno	
_<- m1 m2	side effect only, igno	ore the value m2 :: m b
_<- m1 m2 func =<< m1 = do	side effect only, igno	

Monadic composition operators (1)

(>=>) :: Monad m => (a -> m b) -> (b -> m c) -> (a -> m c)
(<=<) :: Monad m => (b -> m c) -> (a -> m b) -> (a -> m c)

composition operators for two monadic functions

>=>=, <=< can be expressed in do-notation
>=> pipes the result from the left side to the right side
<=< pipes the result from the right side to the left side</pre>

Monadic composition operators (2)

f >=> g = \x -> do y <- f x g y	f::a->mb, x::a, fx::mb g::b->mc, y::b, gy::mc	
g <=< f = \x -> do y <- f x g y	f::a->mb, x::a, fx::mb g::b->mc, y::b, gy::mc	
f >=> g = g <=< f g >=> f = f <=< g	First f Then g	

References

- [1] ftp://ftp.geoinfo.tuwien.ac.at/navratil/HaskellTutorial.pdf
- [2] https://www.umiacs.umd.edu/~hal/docs/daume02yaht.pdf