

Pointers (1A)

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Pointers

Variables

```
int a ;
```

a can hold an integer data

type : int
size : 4 bytes
value : integer

```
a = 100 ;
```

a holds the integer 100

address

data

&a

a ← 100

Pointer Variables

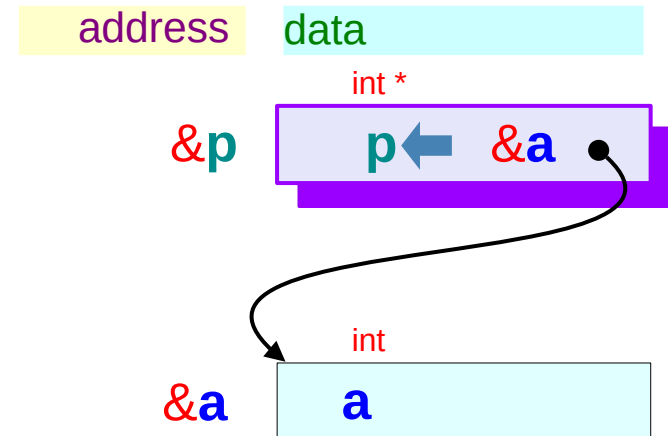
```
int * p ;
```

p can hold an address of an integer data

type : int *
size : 4 bytes (32-bit system)
 : 8 bytes (64-bit system)
value : address

```
p = &a ;
```

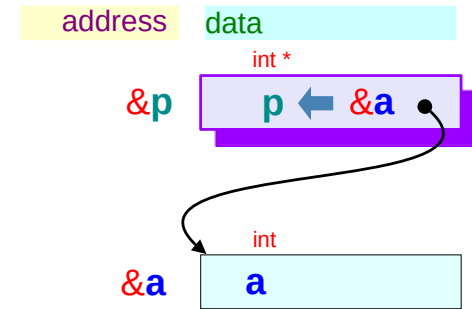
p can hold an address (**&a**) of an integer data (**a**)



Dereferenced Pointer variables

```
int * p ;
```

p can hold an address of an integer data



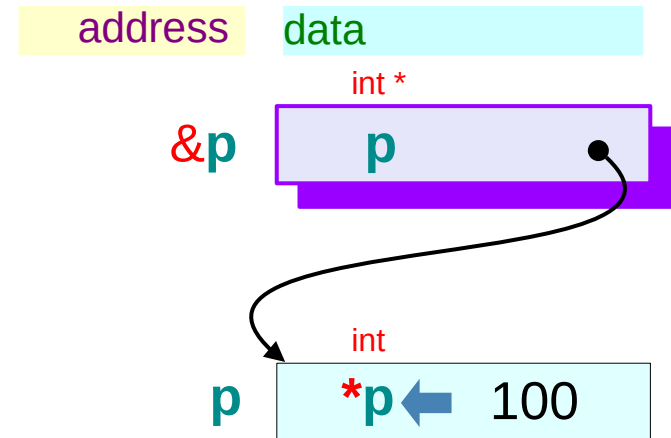
type *variable*

```
int * p ;    p holds the address
```

pointer to int

```
int * p ;    *p holds an integer data
```

int



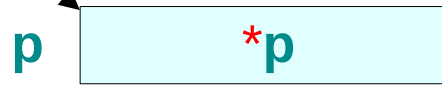
Pointer variable **p** and dereferenced variable ***p**

*&p : the address of
pointer variable **p***



*pointer
variable **p***

*pointer **p** points here,
thus the address is **p***

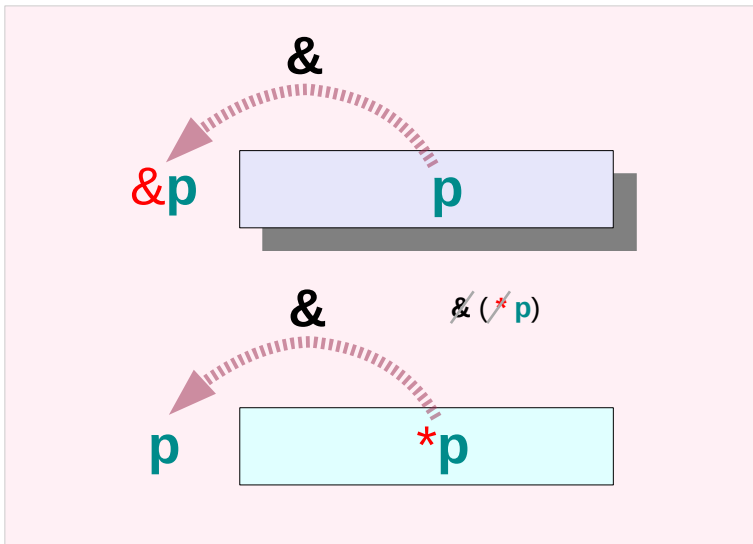


*dereferenced
variable ***p***

****p** and **p** are variables
&**p** is an address value*

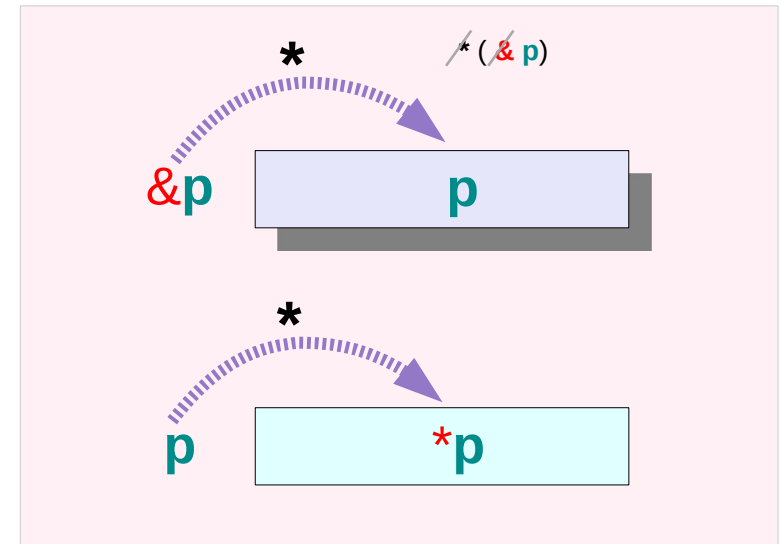
Address-of operator and dereferencing operator (1)

the address of a variable :
address-of operator $\&$



$\& (p) \equiv \text{value}(\&p)$... value
 $\& (*p) \equiv \text{value}(p)$... value

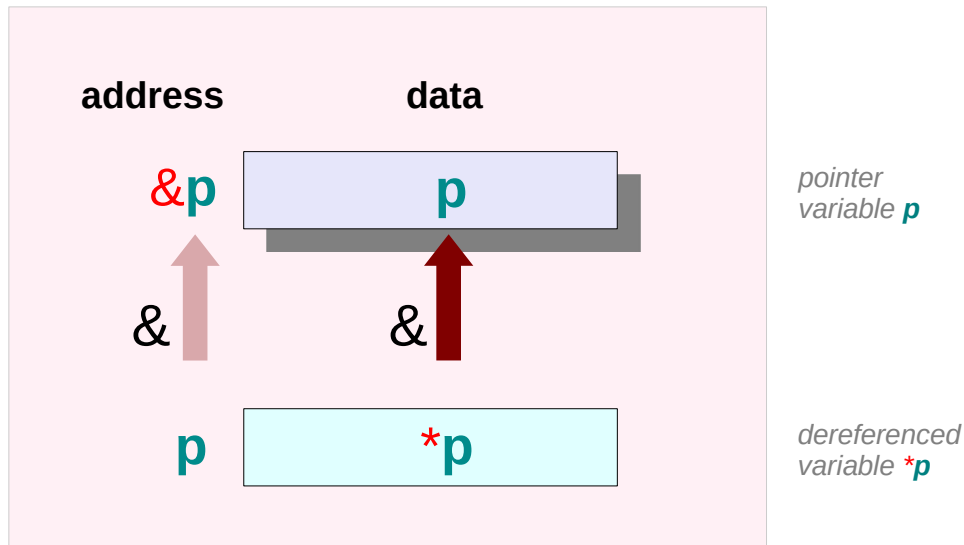
the content at an address :
dereferencing operator $*$



$* (\&p) \equiv p$... variable
 $* (p) \equiv *p$... variable

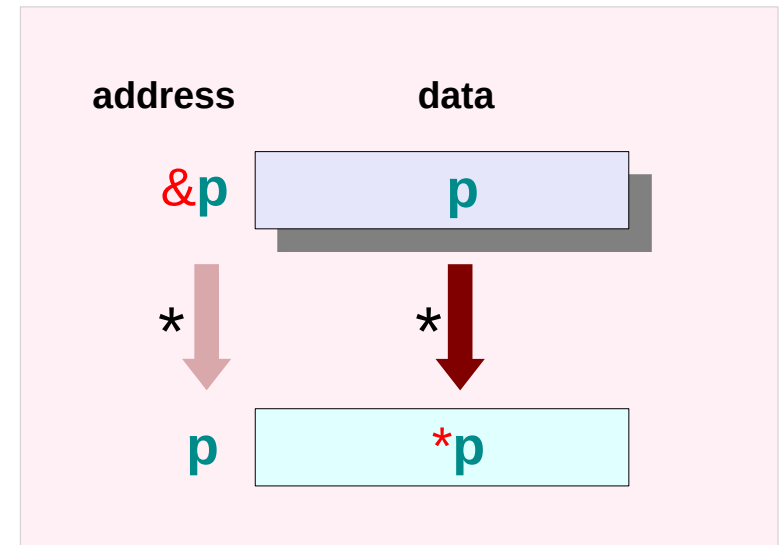
Address-of operator and dereferencing operator (2)

the address of a variable :
address-of operator $\&$



$\& (p) \equiv \text{value}(\&p)$... value
 $\& (*p) \equiv \text{value}(p)$... value

the content at an address :
dereferencing operator $*$

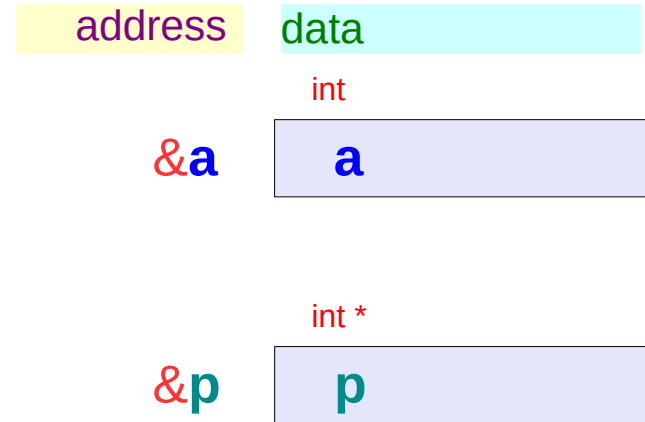


$* (\&p) \equiv p$... variable
 $* (p) \equiv *p$... variable

Variables and their addresses

```
int a ;
```

```
int * p ;
```

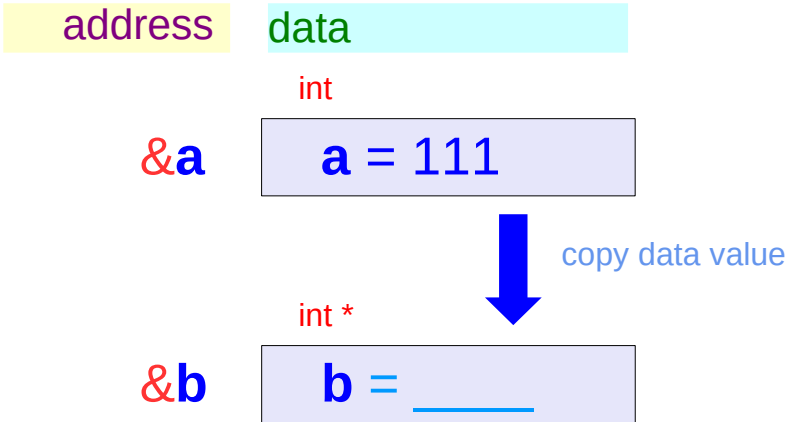


Assignment of a value

```
int a ;
```

```
int b ;
```

```
b = a ;
```

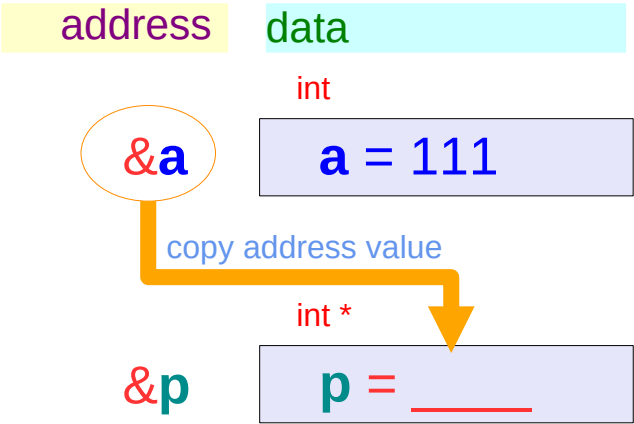


Assignment of an address value

```
int a ;
```

```
int * p ;
```

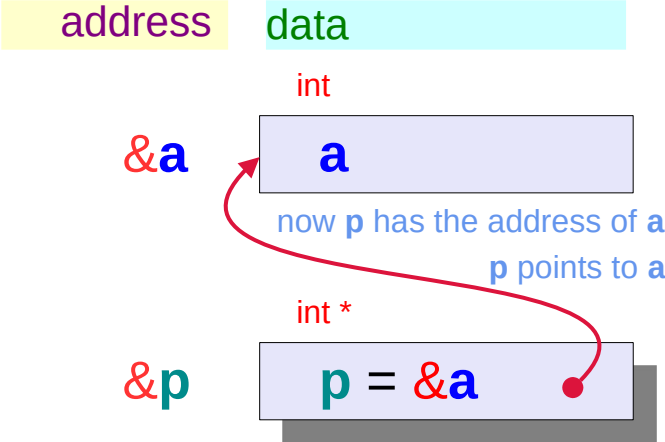
```
p = &a;
```



Arrow notations

```
int a ;
```

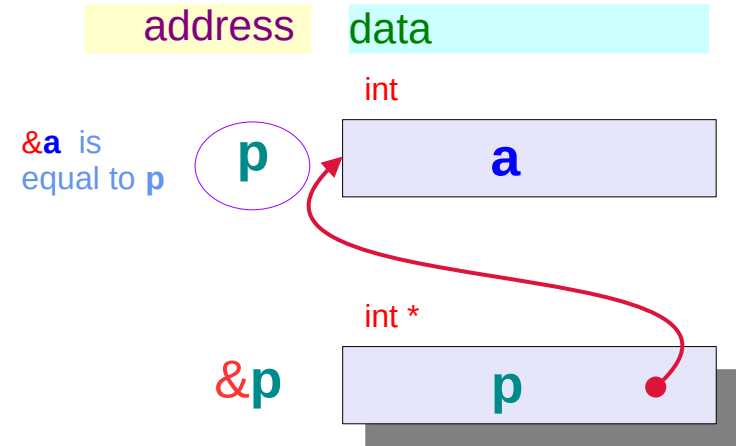
```
int * p = &a;      p = &a;
```



Pointed address : p

```
int a;
```

```
int * p = &a;    p = &a;
```

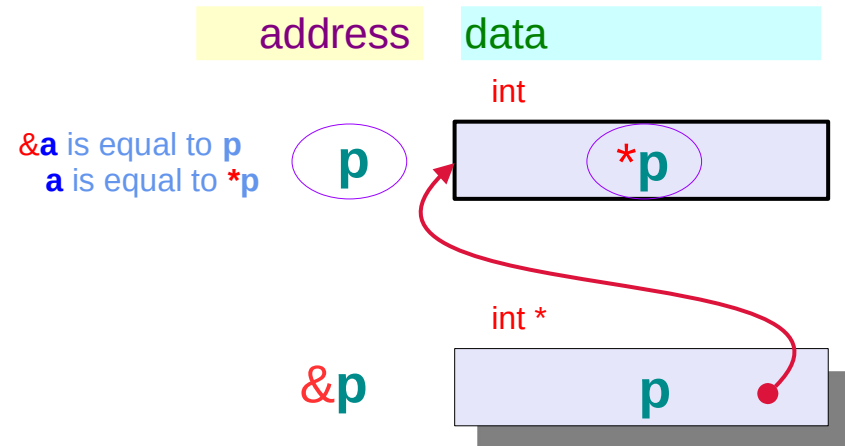


$p \equiv \&a$

Dereferenced variable : *p

```
int a;
```

```
int * p = &a;    p = &a;
```



equivalence

$p \equiv \&a$

$*(p) \equiv *(\&a)$

$*p \equiv a$

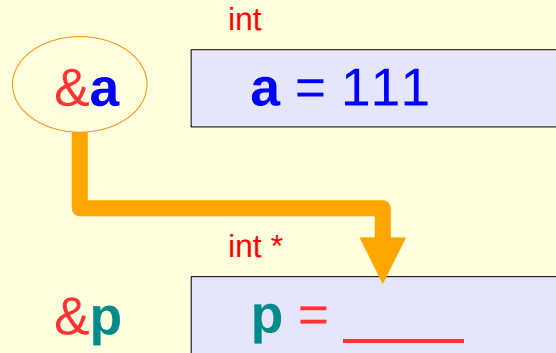
Variable `p`, `*p`

```
int a;
```

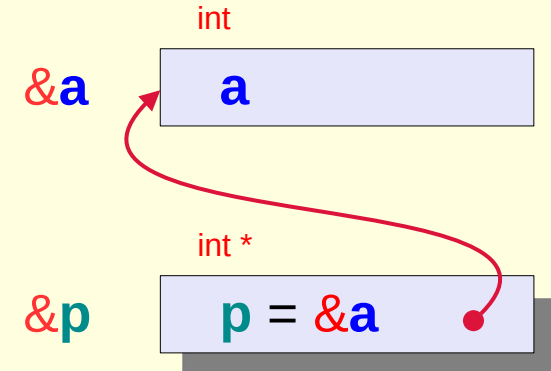
```
int *p;
```

```
p = &a;
```

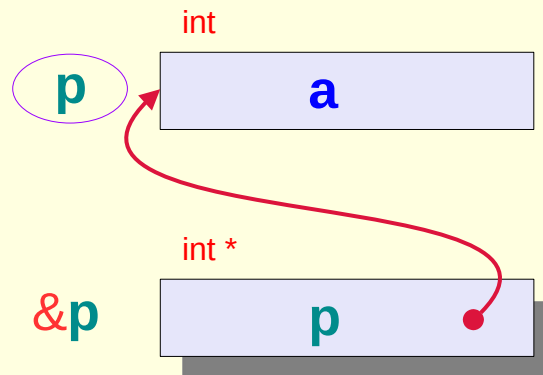
(1) copy address value



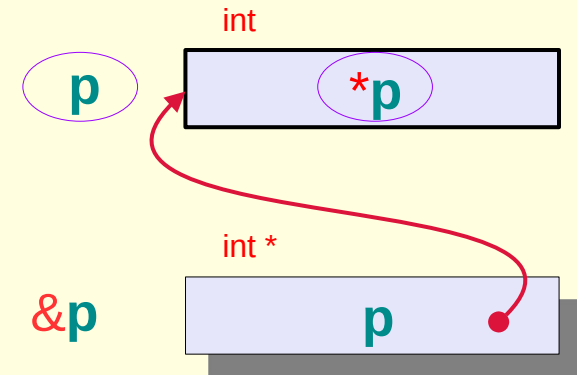
(2) `p` points to `a`



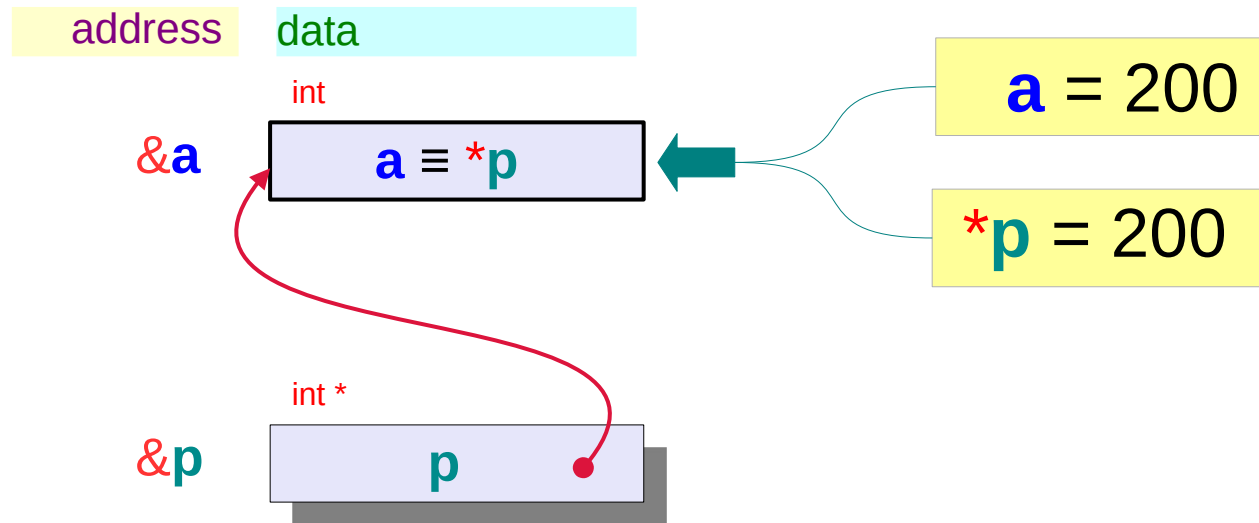
(3) `&a` is equal to `p`



(4) `a` is equal to `*p`



Two way to access: **a** and ***p**



- 1) Read / Write **a**
- 2) Read / Write ***p**

Double Pointers

Double Pointer Variable Definition

```
int ** q ;
```

q holds an address

```
int **
```

```
q ;
```

a pointer to
an integer pointer

q holds an address of
an integer pointer data

```
int *
```

```
*q ;
```

a pointer to
an integer

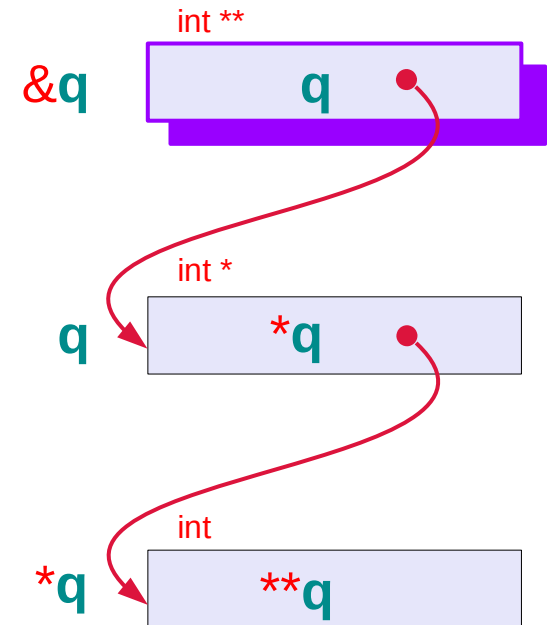
***q** holds an address of
an integer data

```
int
```

```
**q ;
```

an integer

****q** holds an integer data



Variables and their addresses

```
int a ;
```

```
int * p ;
```

```
int ** q ;
```

address	data
<code>&a</code>	<code>int</code> <code>a</code>
<code>&p</code>	<code>int *</code> <code>p</code>
<code>&q</code>	<code>int **</code> <code>q</code>

Assignments of address values

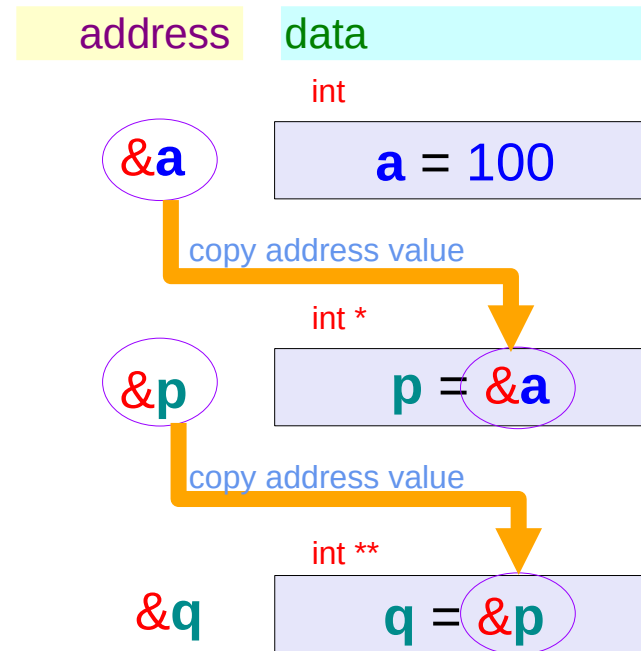
```
int a = 100 ;
```

```
int * p = &a ;
```

```
int ** q = &p ;
```

```
p = &a;
```

```
q = &p;
```



Arrow notations

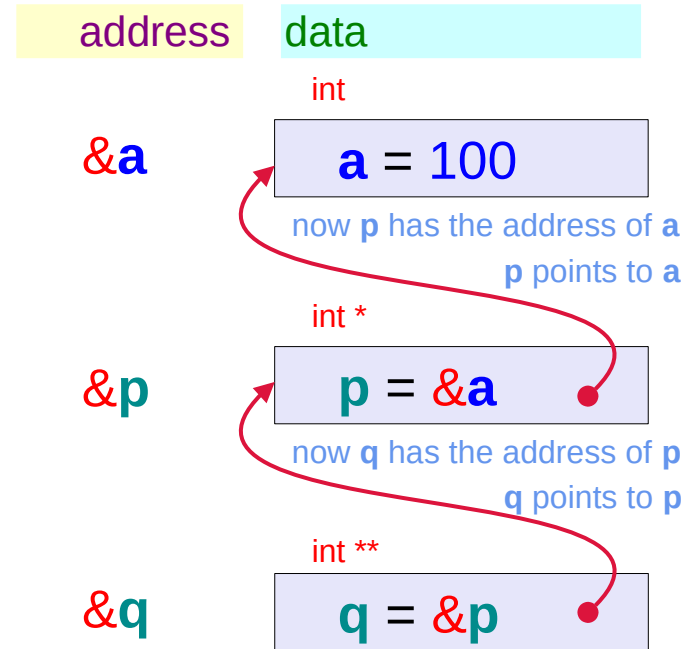
```
int a = 100 ;
```

```
int * p = &a ;
```

```
int ** q = &p ;
```

```
p = &a;
```

```
q = &p;
```



Pointed addresses : p, q

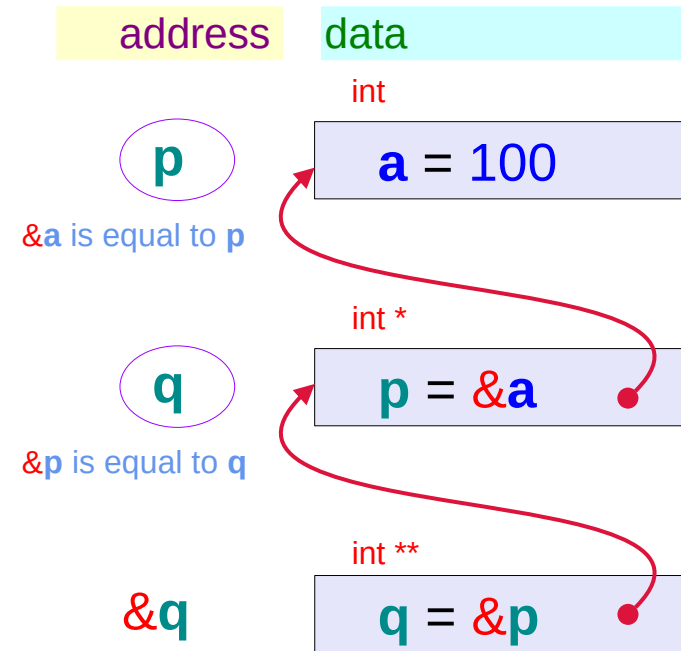
```
int a = 100 ;
```

```
int * p = &a ;
```

```
int ** q = &p ;
```

```
p = &a;
```

```
q = &p;
```



Dereferenced variables : *q, **q

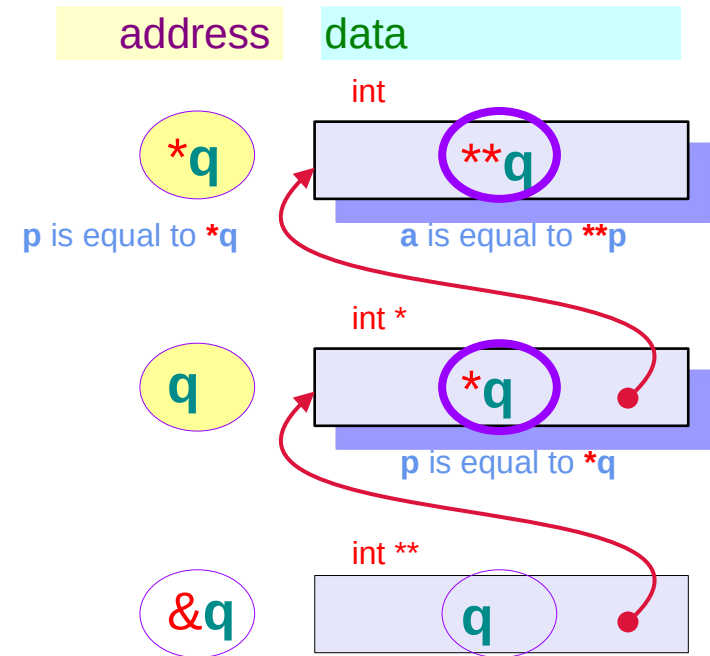
```
int a = 100 ;
```

****q ≡ a**

```
int * p = &a ;
```

***q ≡ p**

```
int ** q = &p ;
```

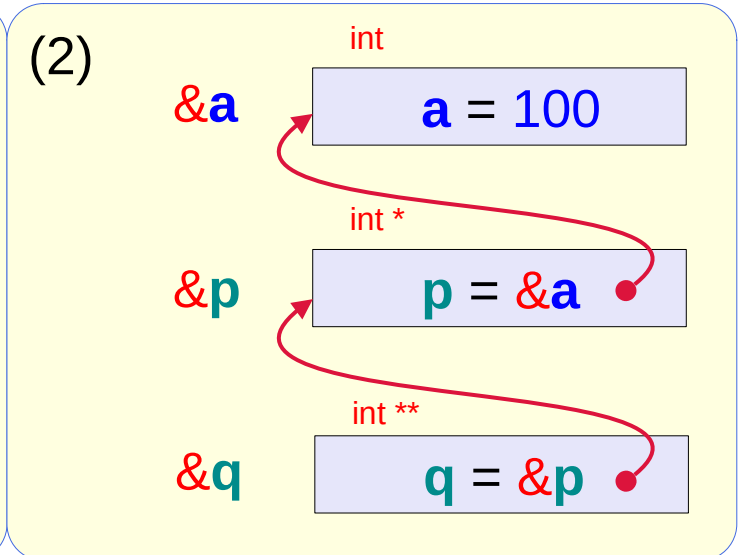
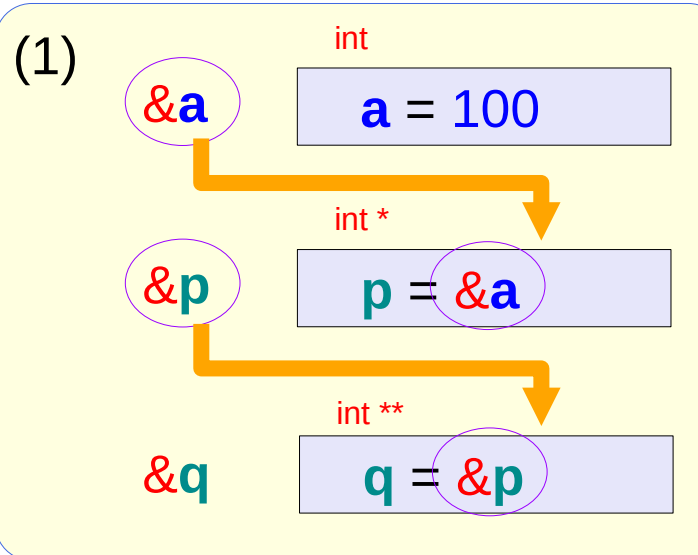


Variable `p`, `*p`, and `q`, `*q`, `**q`

```
int a ;
```

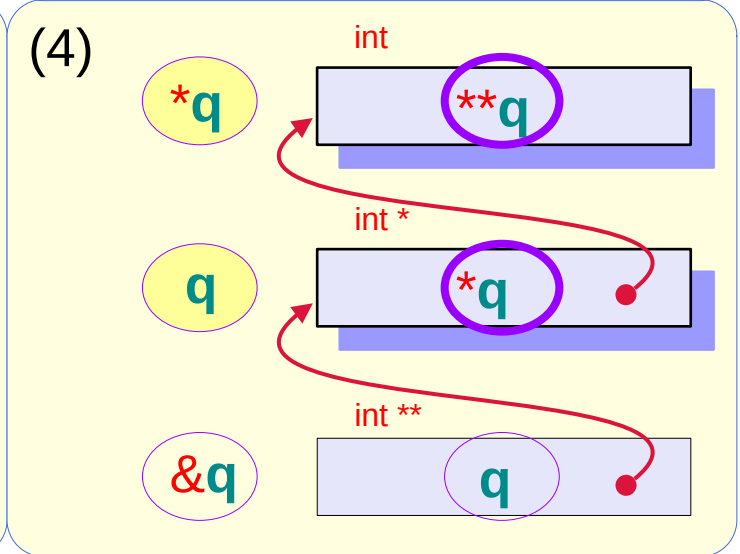
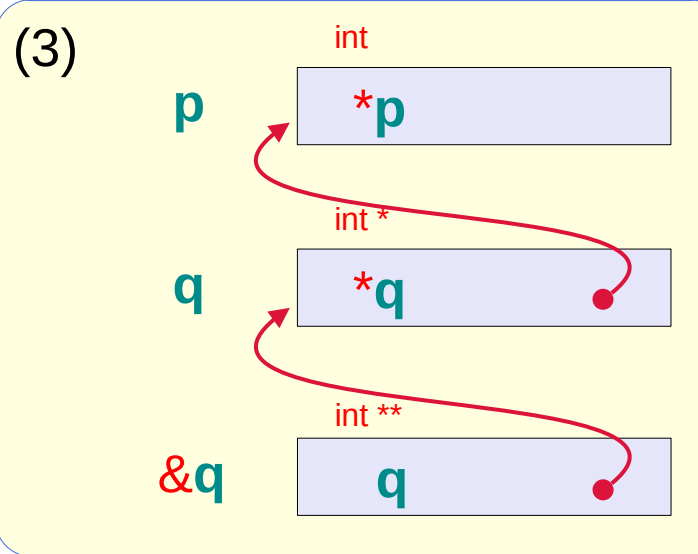
```
int * p ;
```

```
int ** q ;
```



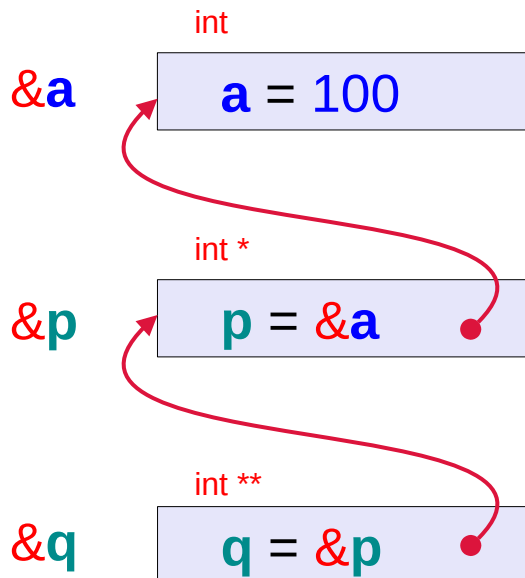
```
p = &a;
```

```
q = &p;
```



Aliased variables : *p,*q, **q

```
int    a = 100 ;  
int    *p = &a ;  
int    **q = &p ;
```



Address assignment

```
p = &a ;
```

Variable aliasing

$$*p \equiv a$$

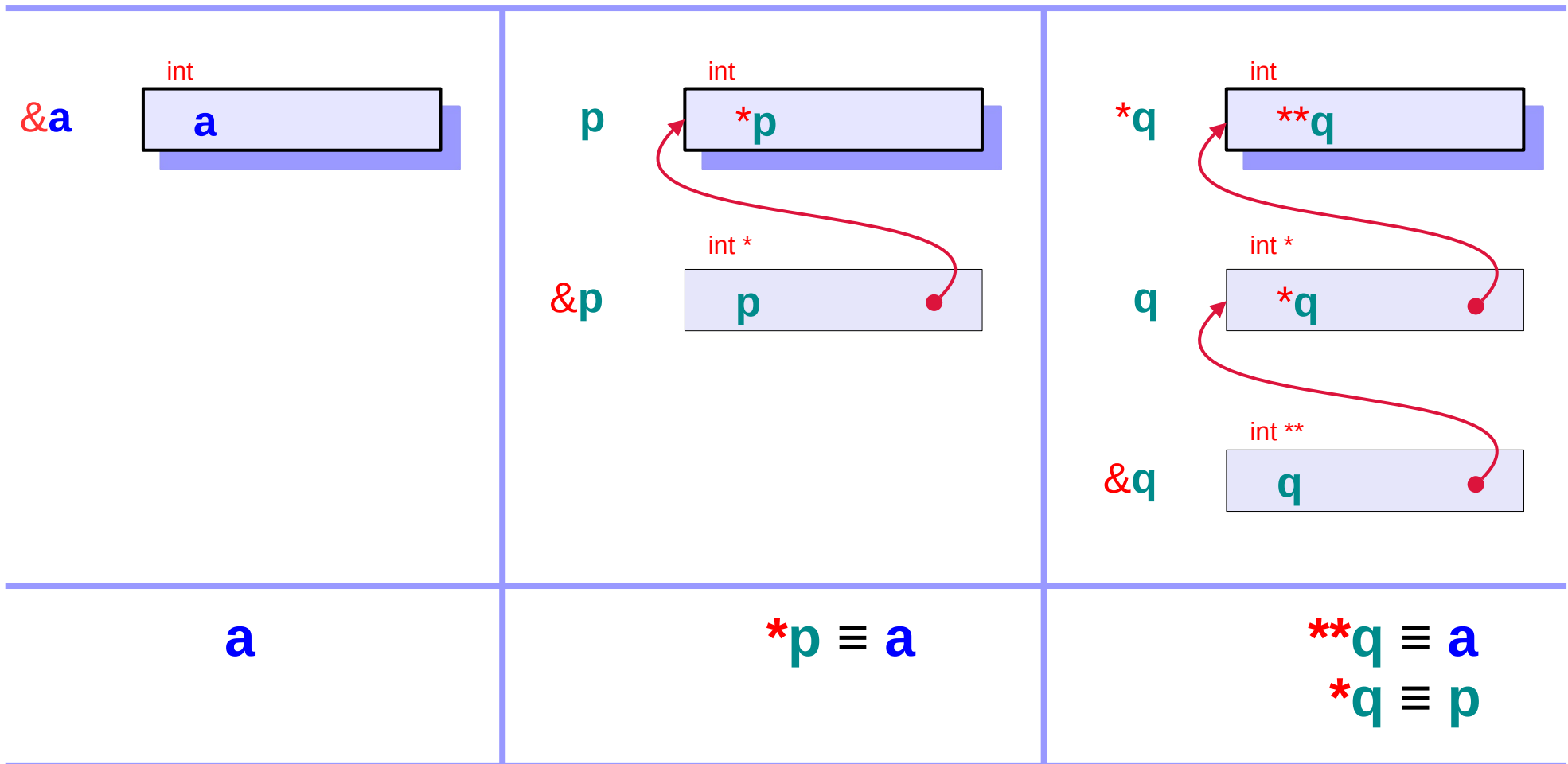
```
p ≡ &a  
*(p) ≡ *(&a)  
*p ≡ a
```

```
q = &p ;  
*q ≡ p  
**q ≡ a
```

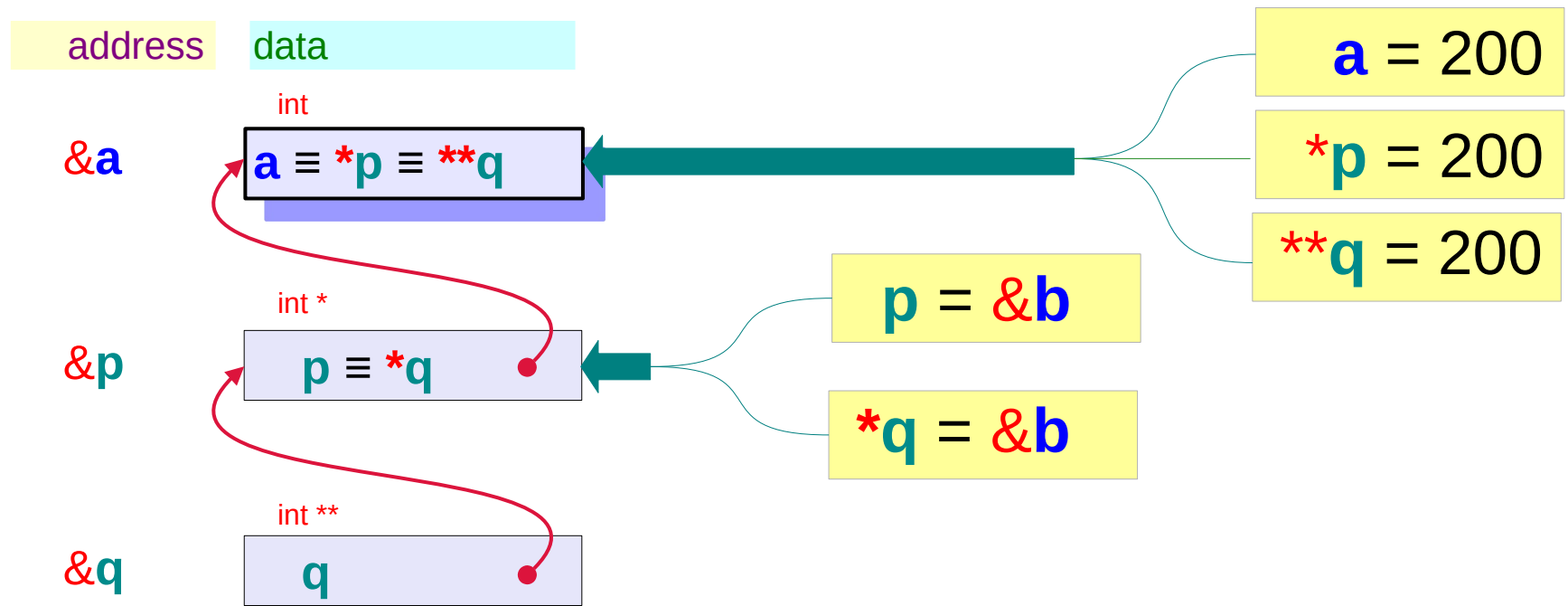
```
q ≡ &p  
*(q) ≡ *(&p)  
*q ≡ p  
**q ≡ *p  
**q ≡ a
```

equivalent relations after address assignment

Two aliased variables of $a : *p, **q$



Two more ways to access **a** : ***p**, ****q**



- 1) Read / Write **p**
- 2) Read / Write ***q**

- 1) Read / Write **a**
- 2) Read / Write ***p**
- 3) Read / Write ****q**

Single and Double Pointers

Pointed Addresses and Data

`int a ;` `&a` ^{int} `a = 100`

The integer variable `a` holds an integer data

`int * p ;` `&p` ^{int *} `p` → `p` ^{int} `*p = 200`

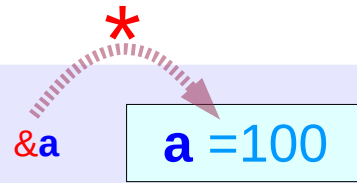
The **pointer** variable `p` holds an address, at this address, an integer data is stored

`int ** q ;` `&q` ^{int **} `q` → `q` ^{int *} `*q` → `*q` ^{int} `**q = 30`

The **pointer** variable `q` holds an address, at the address `q`, another address `*q` is stored, at the address `*q`, an integer data `**q` is stored

Dereferencing Operator *

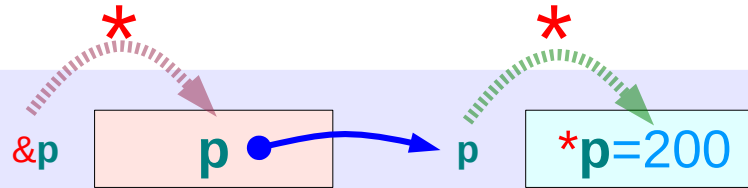
`int a ;`



the expression `a` is a *variable* that can be assigned

$$*(\&a) \equiv a$$

`int * p ;`

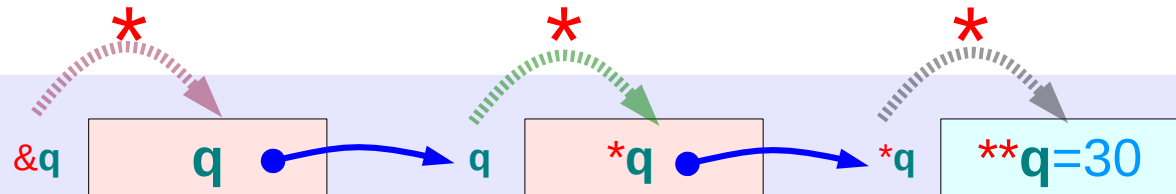


the expression `*p` is a *variable* that can be assigned

$$*(\&p) \equiv p$$

$$*(p) \equiv *p$$

`int ** q ;`



the expression `*q`, `**q` are *variables* that can be assigned

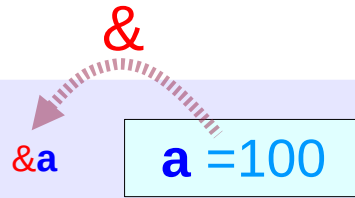
$$*(\&q) \equiv q$$

$$*(q) \equiv *q$$

$$*(*q) \equiv **q$$

Address-of Operator &

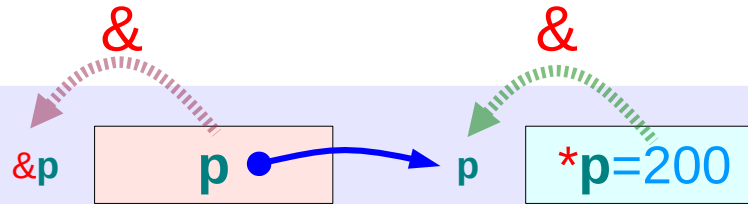
`int a ;`



the expression `&a` returns
an *address value*
that cannot be assigned

`&a` \equiv `value(&a)`

`int * p ;`

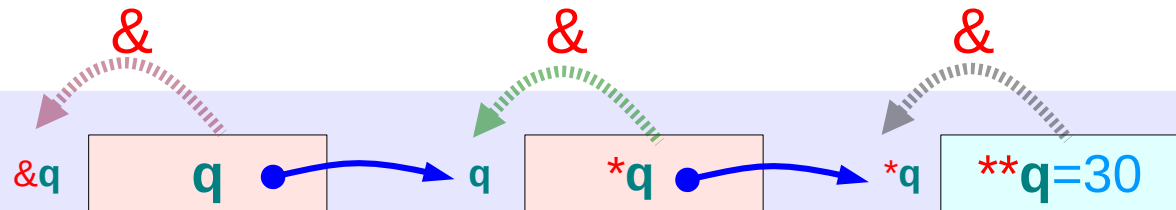


the expression `&p` returns
an *address value*
that cannot be assigned

`&p` \equiv `value(&p)`

`&>(*p)` \equiv `value(p)`

`int ** q ;`



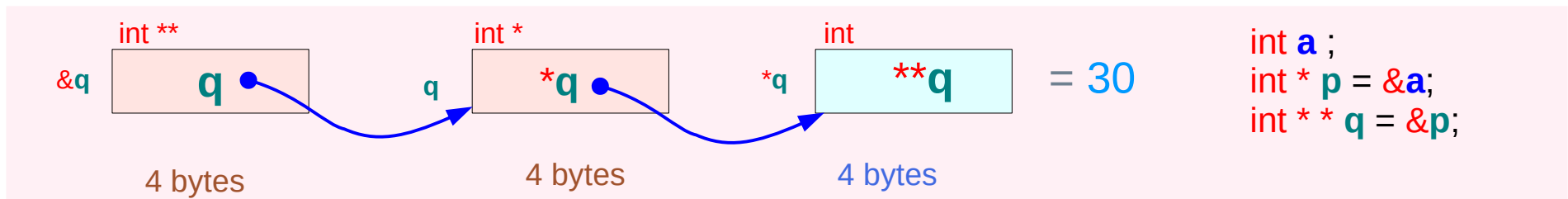
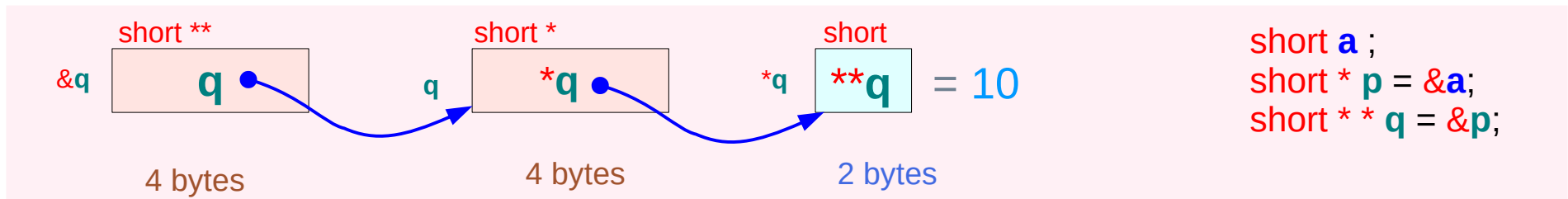
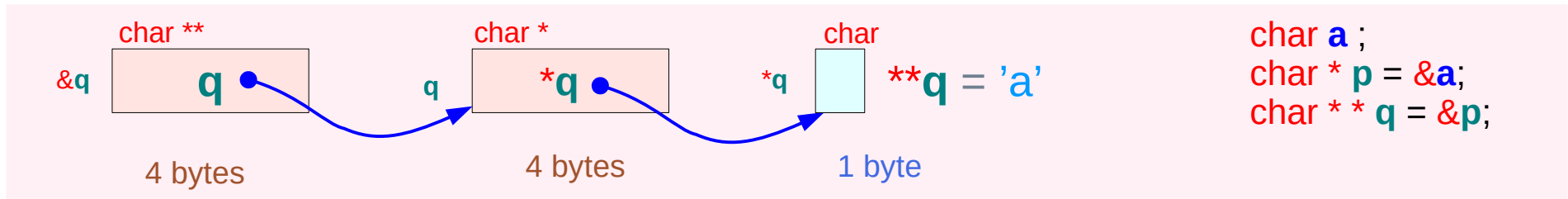
the expression `&q` returns
an *address value*
that cannot be assigned

`&q` \equiv `value(&q)`

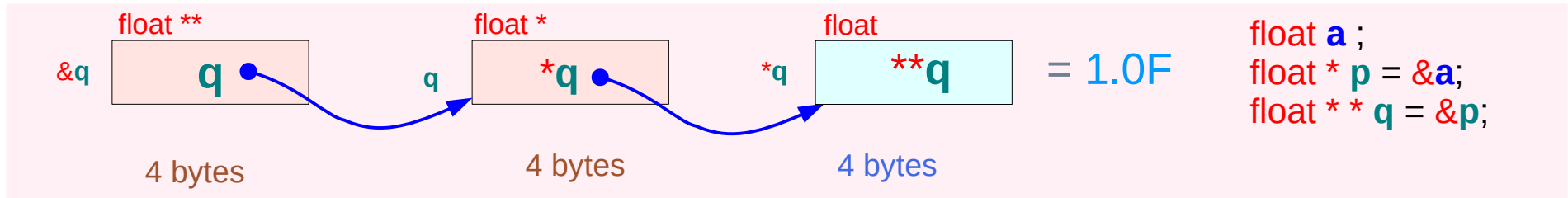
`&>(*q)` \equiv `value(q)`

`&**q` \equiv `value(*q)`

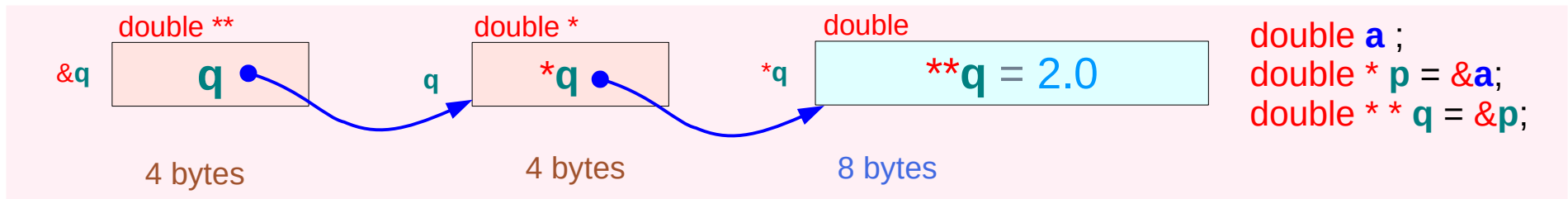
Pointers to various types on 32-bit system (1)



Pointers to various types on 32-bit system (2)



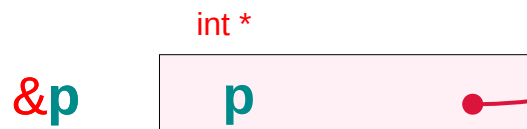
on a 32-bit system



on a 32-bit system

Neighborhood of *p

```
int a = 100 ;  
int * p = &a ;
```



abstract
address

int *

p-3

p-2

p-1

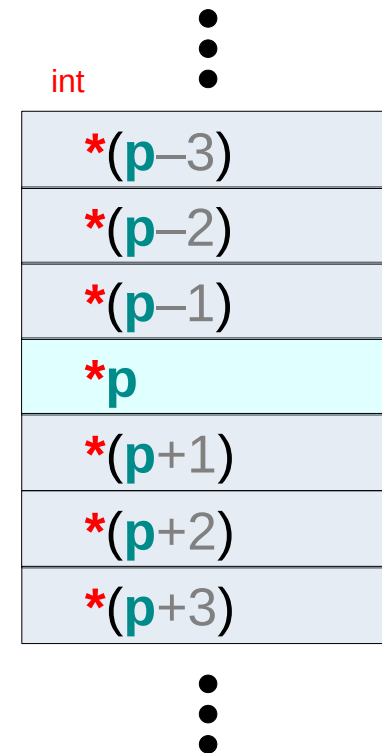
p

p+1

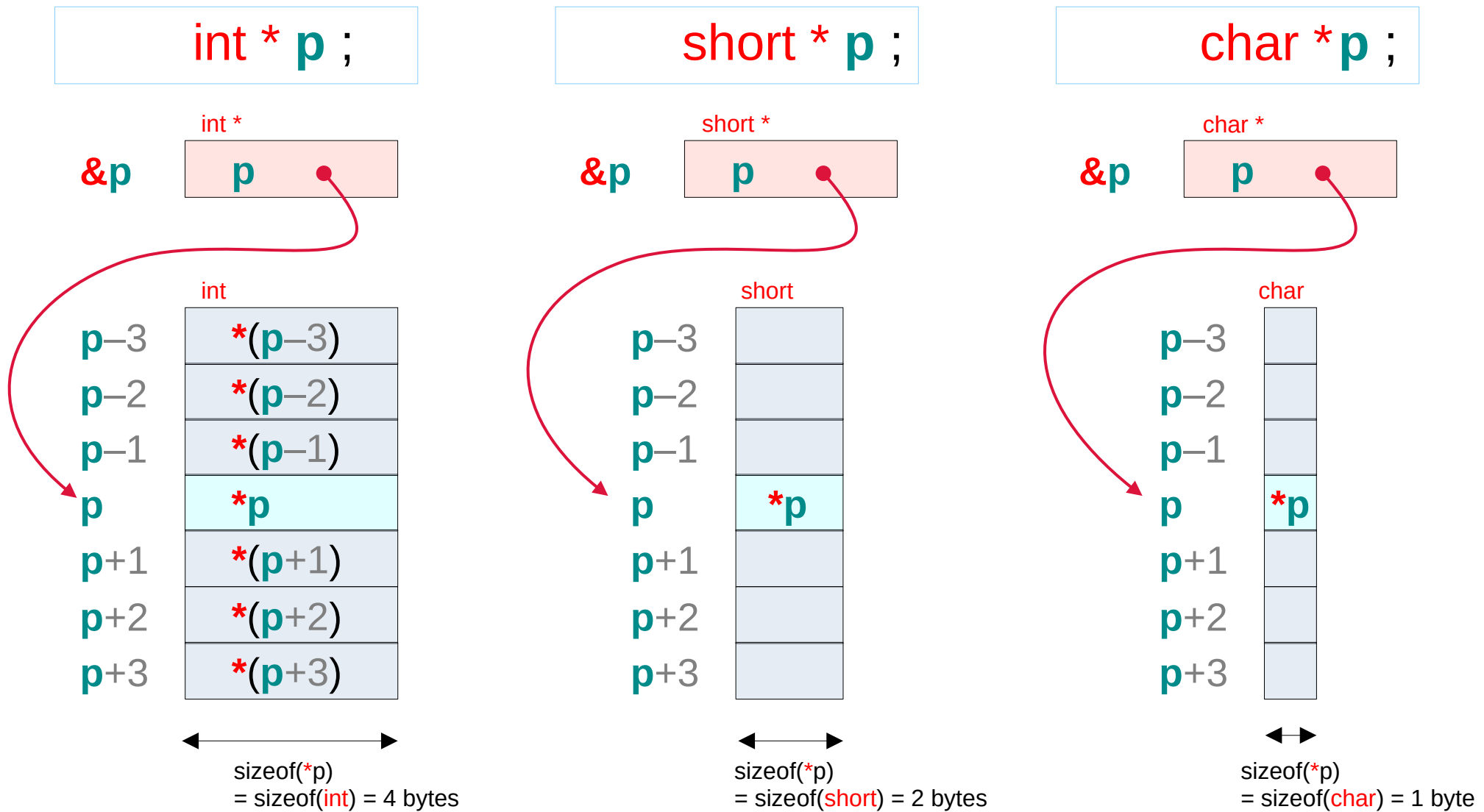
p+2

p+3

referenced
data



Neighborhood of *p – various sizes of referenced data

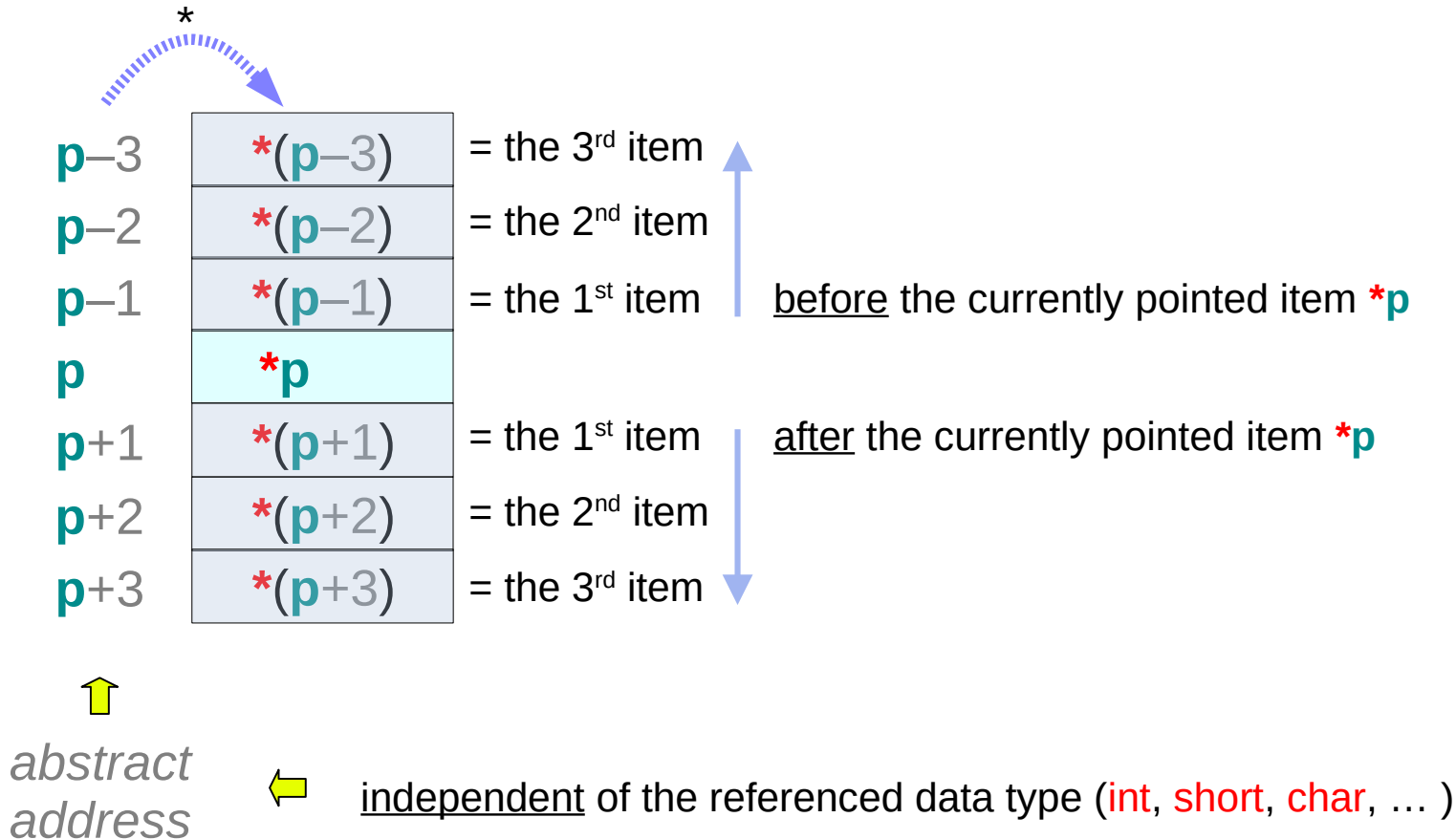


Neighborhood of $*p$ – the same abstract address

`int * p ;`

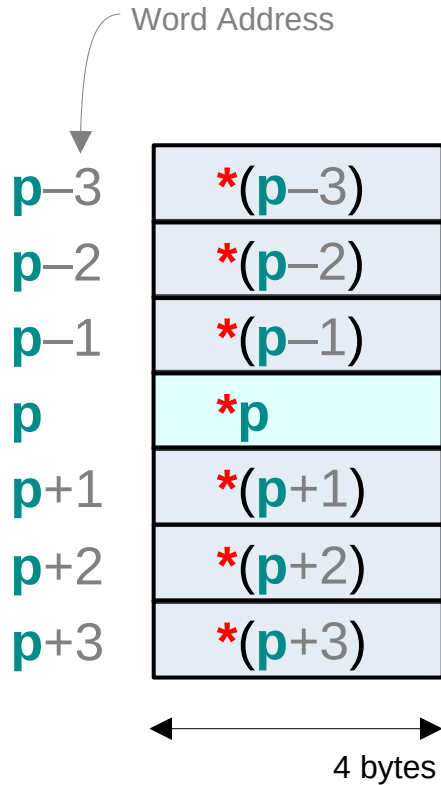
`short * p ;`

`char * p ;`

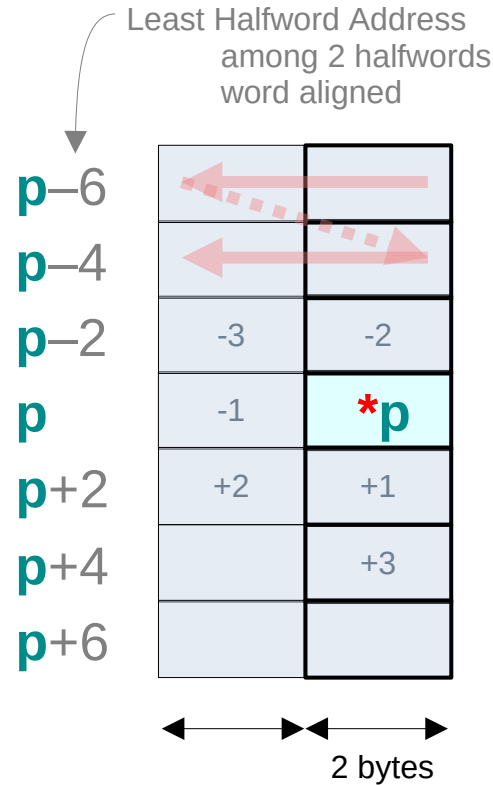


Neighborhood of `*p` – word aligned view

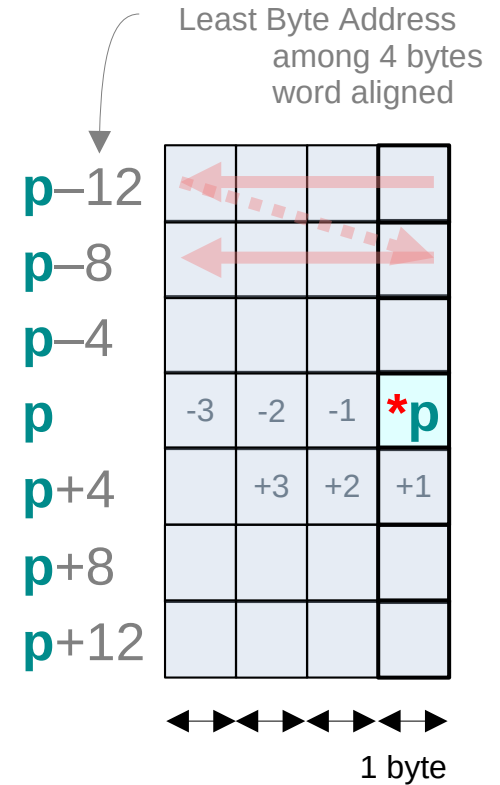
`int * p ;`



`short * p ;`



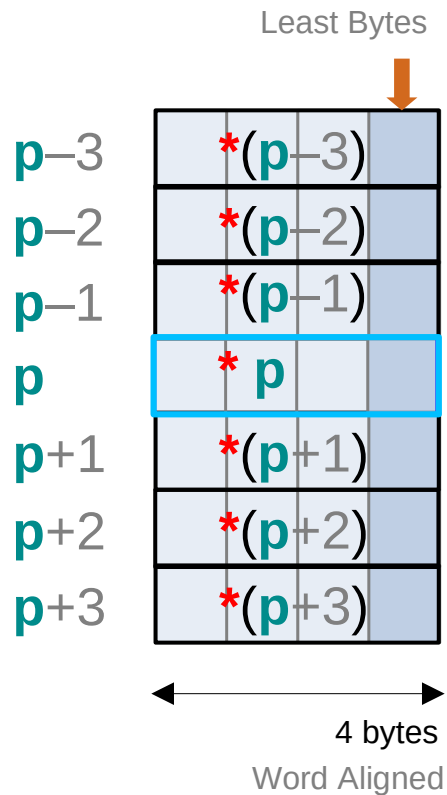
`char * p ;`



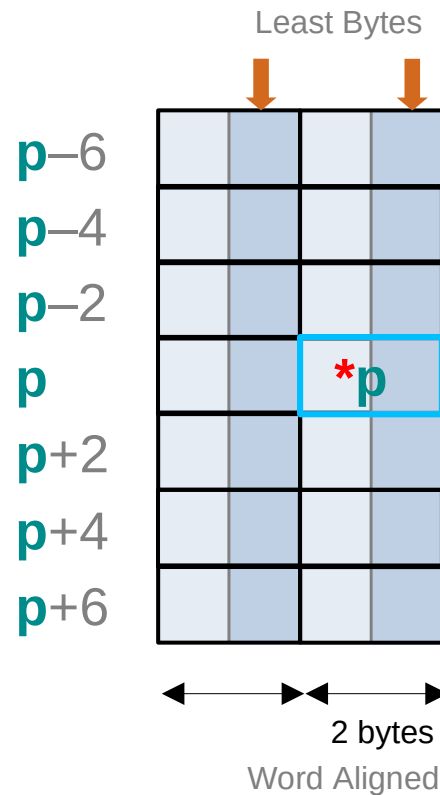
word view

Neighborhood of `*p` – least bytes

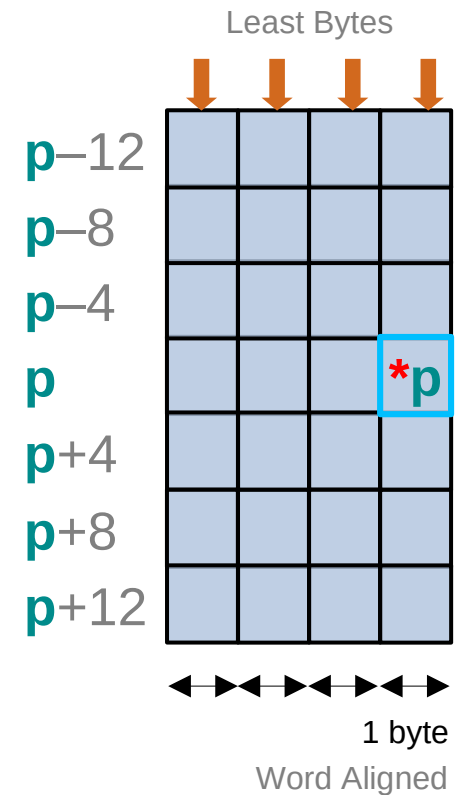
```
int * p ;
```



```
short * p ;
```

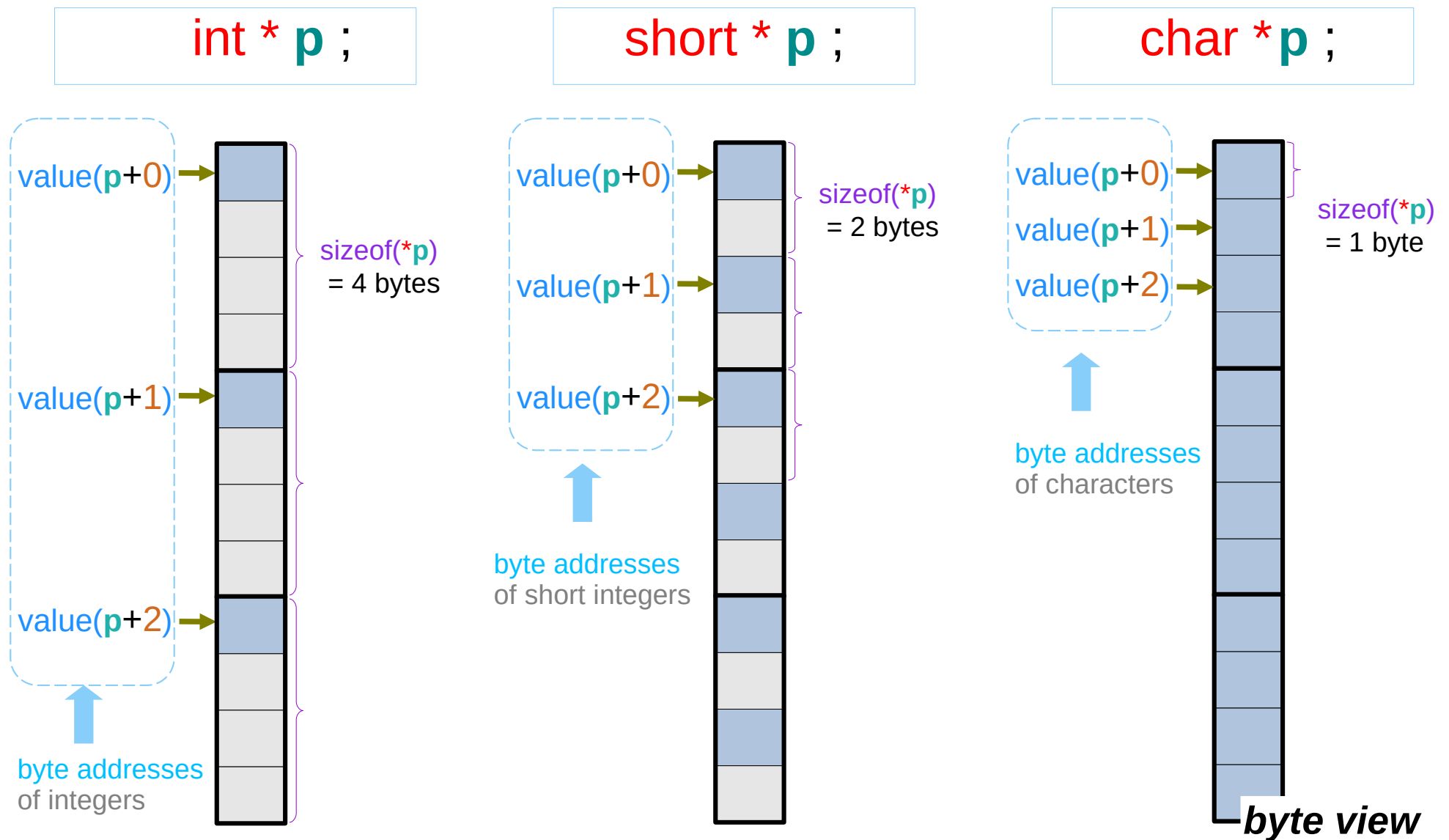


```
char * p ;
```



word view

Neighborhood of `*p` – least byte addresses



Neighborhood of $*p$ – incremented pointer values

```
int * p ;
```

abstract address

expression $p+1$

least byte address
of int type $*(p+1)$

$$\text{value}(p+1) = \text{value}(p) + \text{sizeof}(*p)$$
$$\text{value}(p+i) = \text{value}(p) + i*4$$

```
short * p ;
```

abstract address

expression $p+1$

least byte address
of short type $*(p+1)$

$$\text{value}(p+1) = \text{value}(p) + \text{sizeof}(*p)$$
$$\text{value}(p+i) = \text{value}(p) + i*2$$

```
char * p ;
```

abstract address

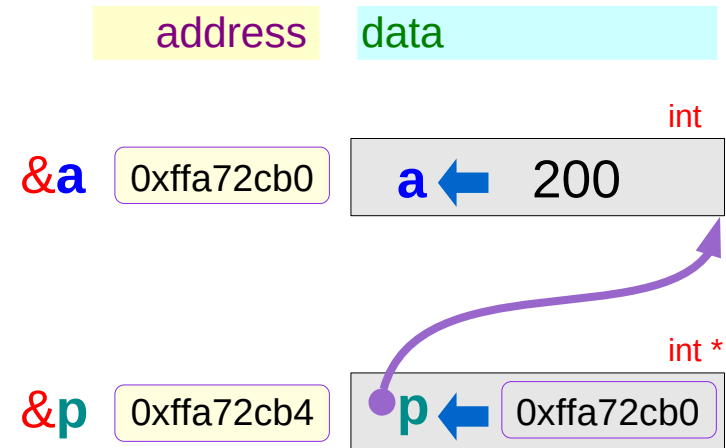
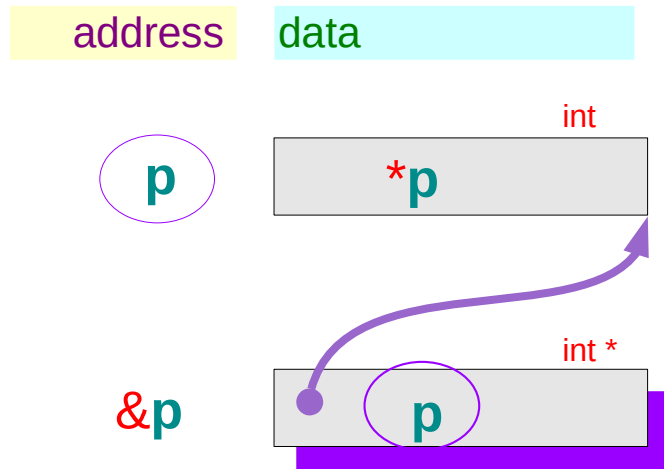
expression $p+1$

least byte address
of char type $*(p+1)$

$$\text{value}(p+1) = \text{value}(p) + \text{sizeof}(*p)$$
$$\text{value}(p+i) = \text{value}(p) + i*1$$

Pointer Variable Example

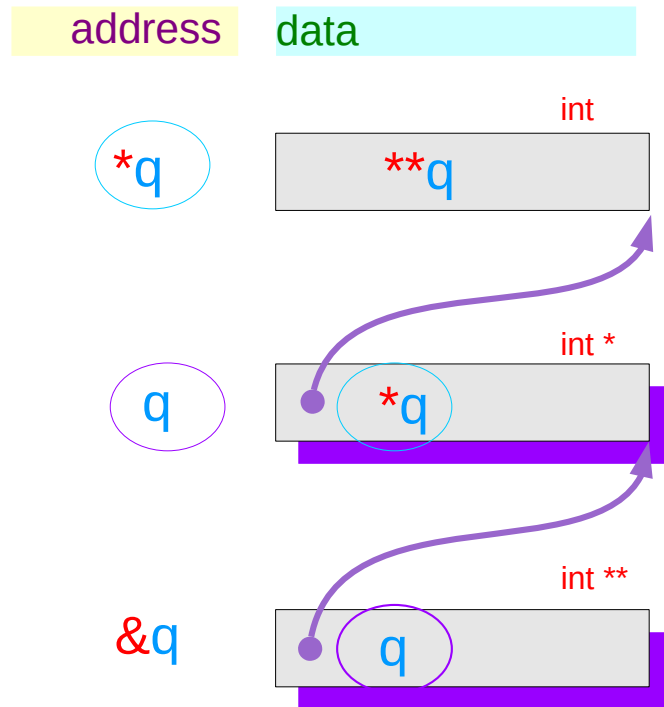
Pointer variable **p** of the type **int ***



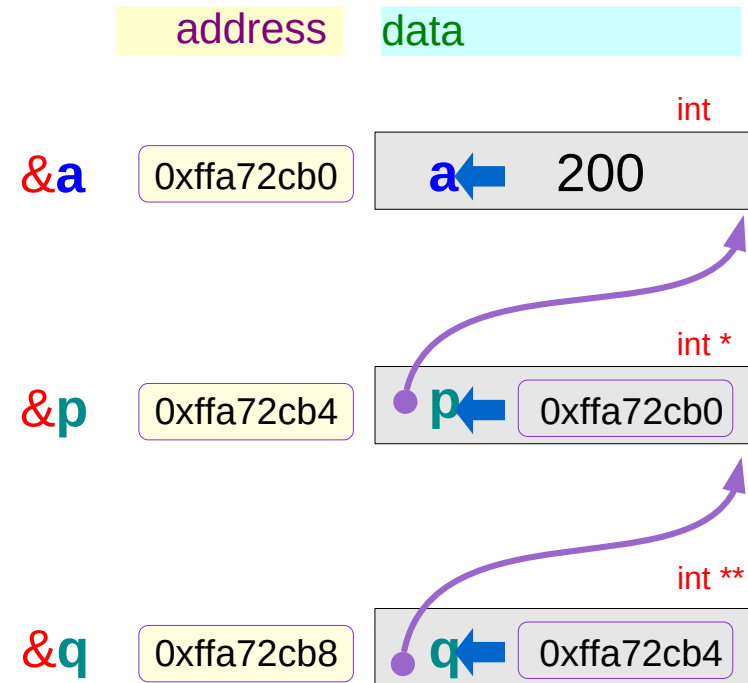
```
int    a = 200;  
int *  p = &a;
```

```
&p → 0xffa72cb4  
p  → 0xffa72cb0  
*p → 200
```

Pointer variable **q** of the type **int ****



```
int a = 200;  
int * p = &a;  
int ** q = &p;
```




```
&q ← 0xffa72cb8  
q ← 0xffa72cb4  
*q ← 0xffa72cb0  
**q ← 200
```

Pointer variable example codes

```
// t.c
#include <stdio.h>

int main(void) {
    int    a = 200;
    int    *p = &a;
    int    **q = &p;

    printf("&a=%p a=%d \n", &a, a);
    printf("&p=%p p=%p \n", &p, p);
    printf("&q=%p q=%p \n", &q, q);

    
}
```

```
gcc -Wall -m32 t.c
./a.out
```

```
&a= 0xffa72cb0 a=      200
&p= 0xffa72cb4 p= 0xffa72cb0
&q= 0xffa72cb8 q= 0xffa72cb4
```

```
printf("  &p=%12p \n", &p);
printf("  p=%12p \n",  p);
printf("  *p=%12d \n",  *p);
printf("\n");

printf("  &q=%12p \n", &q);
printf("  q=%12p \n",  q);
printf("  *q=%12p \n",  *q);
printf("  **q=%12d \n", **q);
```

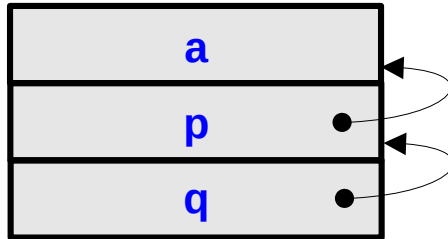
```
&p= 0xffa72cb4
p= 0xffa72cb0
*p=      200

&q= 0xffa72cb8
q= 0xffa72cb4
*q= 0xffa72cb0
**q=      200
```

Byte addresses of variables

abstract
address

&a



&p

&q

sizeof(**a**) = sizeof(int) =
sizeof(**p**) = sizeof(int *) =
sizeof(**q**) = sizeof(int **) =
4 bytes

```
int a = 200 ;  
int *p = &a ;  
int **q = &p ;
```

```
&a= 0xffa72cb0 ➡ value (&a)  
&p= 0xffa72cb4 ➡ value (&p)  
&q= 0xffa72cb8 ➡ value (&q)
```

here, the values of **&a**, **&p**, **&q**
are displayed

byte
address

MSByte

LSByte

0xffa72cb0

00	00	00	c8
----	----	----	----

0xffa72cb4

ff	a7	2c	b0
----	----	----	----

0xffa72cb8

ff	a7	2c	b4
----	----	----	----

byte
address

0xffa72cb0

c8

0xffa72cb1

00

0xffa72cb2

00

0xffa72cb3

00

0xffa72cb4

b0

0xffa72cb5

2c

0xffa72cb6

a7

0xffa72cb7

ff

0xffa72cb8

b4

0xffa72cb9

2c

0xffa72cba

a7

0xffa72cbb

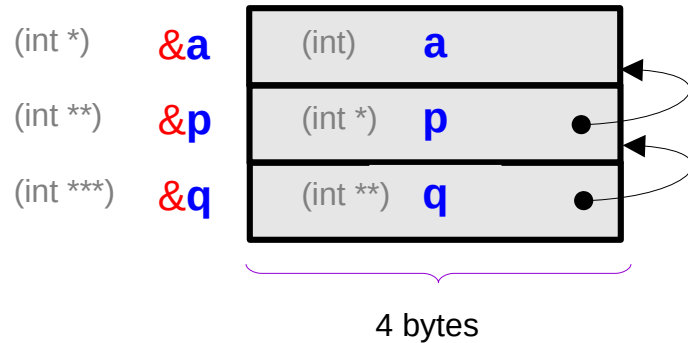
ff

address value

byte address of
the least significant byte
in a little endian system

Abstract vs. byte addresses

*abstract
address*



*abstract address contains
information of **type**, **size**, **value***

*byte
address*

	MSByte			LSByte
0xffa72cb0	00	00	00	c8
0xffa72cb4	ff	a7	2c	b0
0xffa72cb8	ff	a7	2c	b4

*byte address contains
value information only*

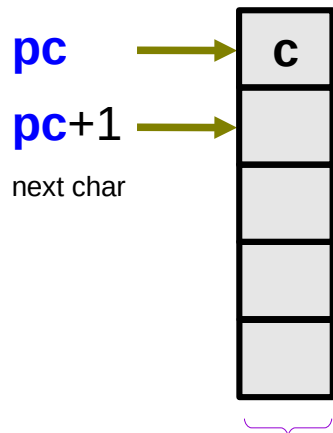
Abstract Address	Type	Size	Value (address value)
&a	int *	4 bytes	0xffa72cb0
&p	int **	4 bytes	0xffa72cb4
&q	int ***	4 bytes	0xffa72cb8

Byte Address	Value
value(&a)	0xffa72cb0
value(&p)	0xffa72cb4
value(&q)	0xffa72cb8

Endianness and memory alignment

Pointers and abstract addresses

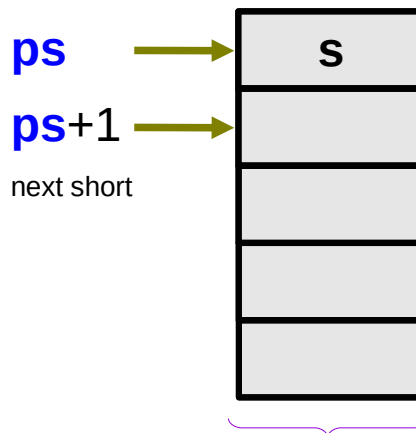
```
char c ;  
char *pc = &c ;
```



next char

$\text{sizeof}(*pc) =$
 $\text{sizeof}(c) =$
 $\text{sizeof}(int) = 1$ byte

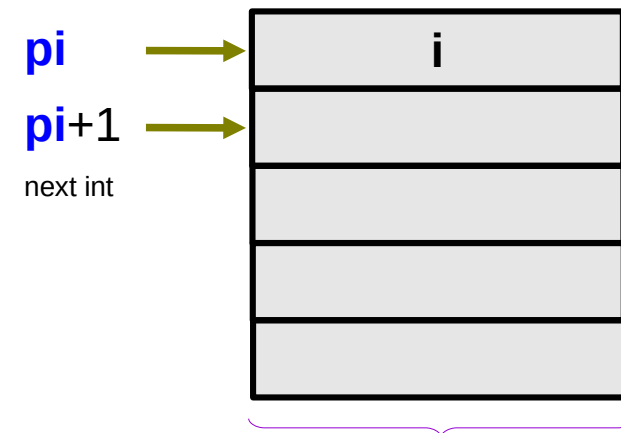
```
short s ;  
short *ps = &s ;
```



next short

$\text{sizeof}(*ps) =$
 $\text{sizeof}(s) =$
 $\text{sizeof}(short) = 2$ bytes

```
int i ;  
int *pi = &i ;
```



next int

$\text{sizeof}(*pi) =$
 $\text{sizeof}(i) =$
 $\text{sizeof}(int) = 4$ bytes

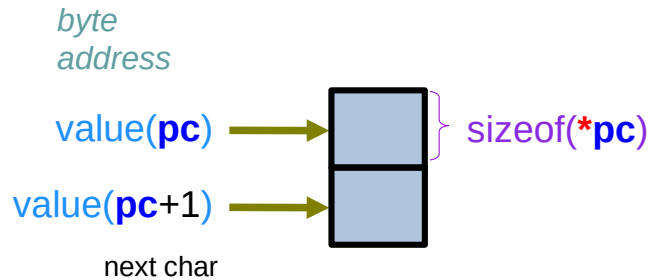
```
sizeof(pc) = 4 or 8 bytes  
type(pc) = char *  
value(pc) = byte address
```

```
sizeof(ps) = 4 or 8 bytes  
type(ps) = short *  
value(ps) = byte address
```

```
sizeof(pi) = 4 or 8 bytes  
type(pi) = int *  
value(pi) = byte address
```

Pointer values and byte addresses

char *pc ;

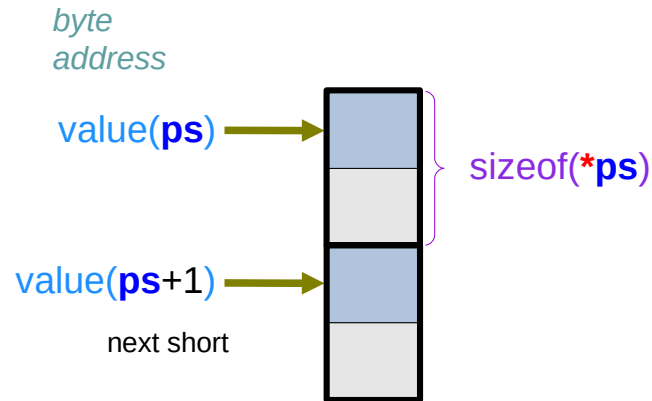


The **value** of a *pointer* (address value) is the **byte address** of the least significant byte in a little endian system

pc points to the least significant byte, i.e., the **1** byte of **char**

$\text{value}(\text{pc}+1) = \text{value}(\text{pc}) + \text{sizeof}(*\text{pc})$

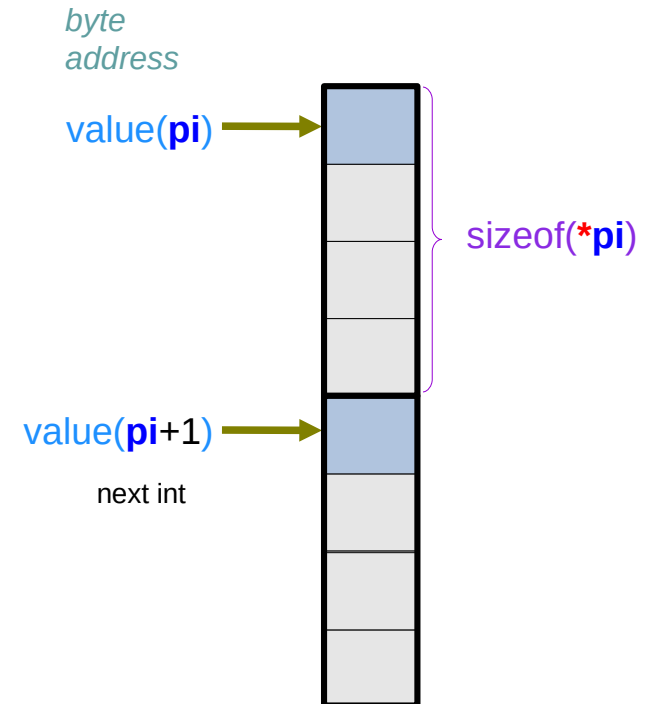
short *ps ;



ps points to the least significant byte among the **2** bytes of **short**

$\text{value}(\text{ps}+1) = \text{value}(\text{ps}) + \text{sizeof}(*\text{ps})$

int *pi ;

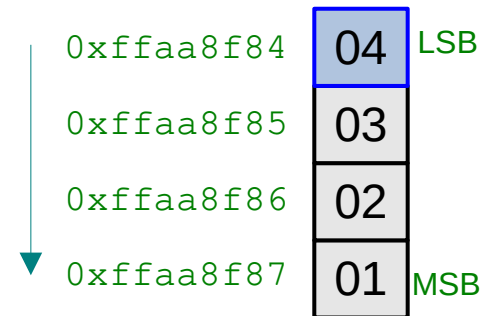
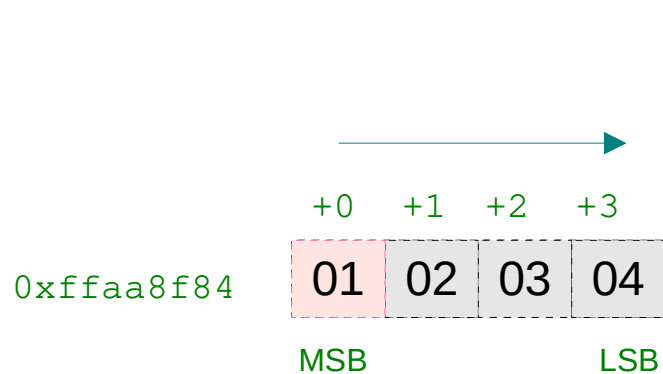


pi points to the least significant byte among the **4** bytes of an **int**

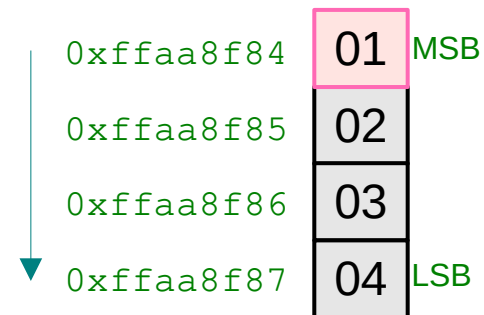
$\text{value}(\text{pi}+1) = \text{value}(\text{pi}) + \text{sizeof}(*\text{pi})$

Little Endian and Big Endian

```
int i = 0x01020304;
```



Little Endian
stores LSByte first

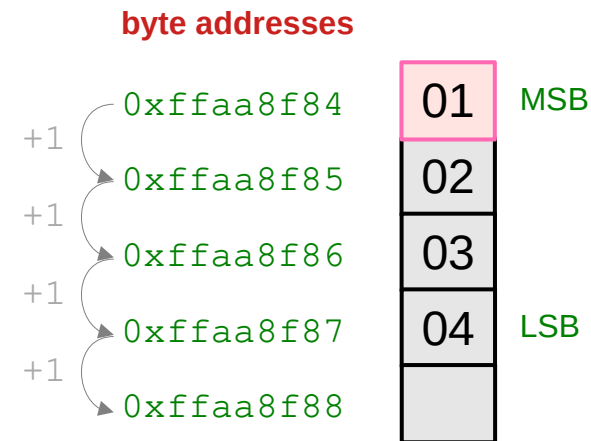
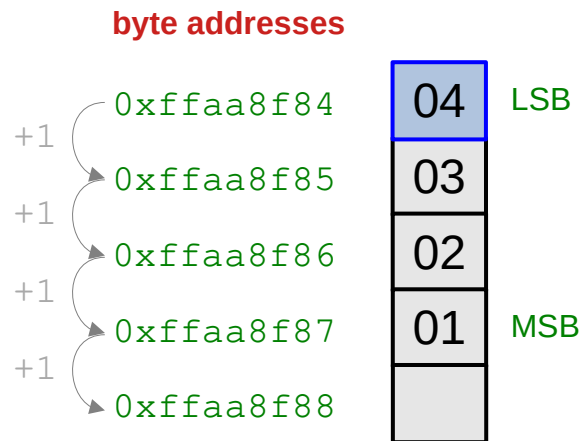
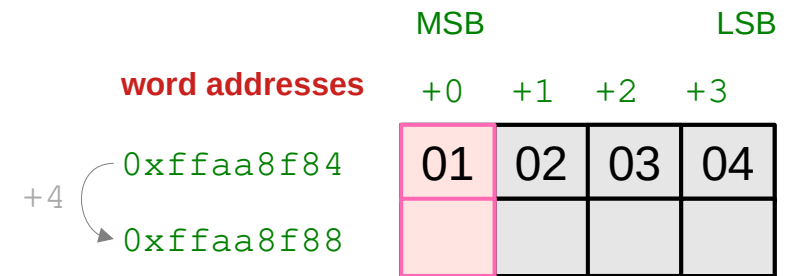
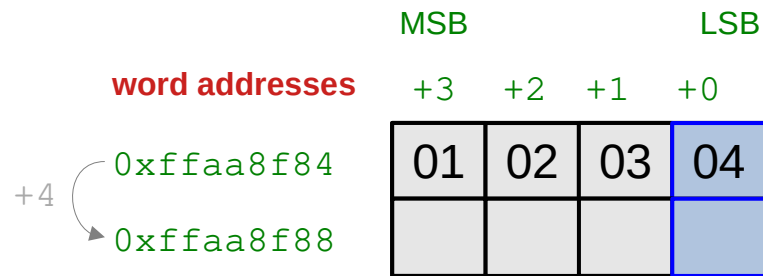


Big Endian
stores MSByte first

Word addresses and byte addresses

Little Endian store LSByte first

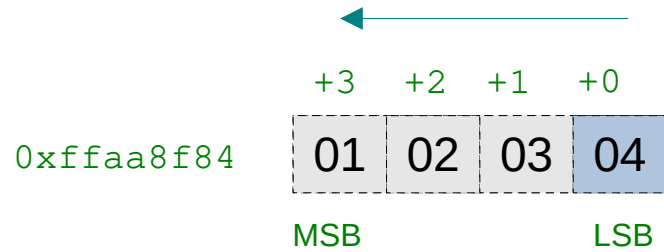
Big Endian store MSByte first



Little Endian Examples

```
int i = 0x01020304;
```

MSB LSB



Little Endian

store LSB first

Least Significant Byte (LSB)

```
int i = 0x01020304;
char *p = (char *)&i;

printf("%p \n", &i);

printf("%p : %x \n", p+0, p[0]);
printf("%p : %x \n", p+1, p[1]);
printf("%p : %x \n", p+2, p[2]);
printf("%p : %x \n", p+3, p[3]);
```

MSB LSB

	0xffaa8f84	
LSB address	0xffaa8f84 : 4	LSB data
	0xffaa8f85 : 3	
	0xffaa8f86 : 2	
MSB address	0xffaa8f87 : 1	MSB data

assume 32-bit CPU
32-bit address
32-bit data (1 word, 4 bytes)

Storing values in a little endian system

long `a = 0x1020304050607080 ;`

MSB

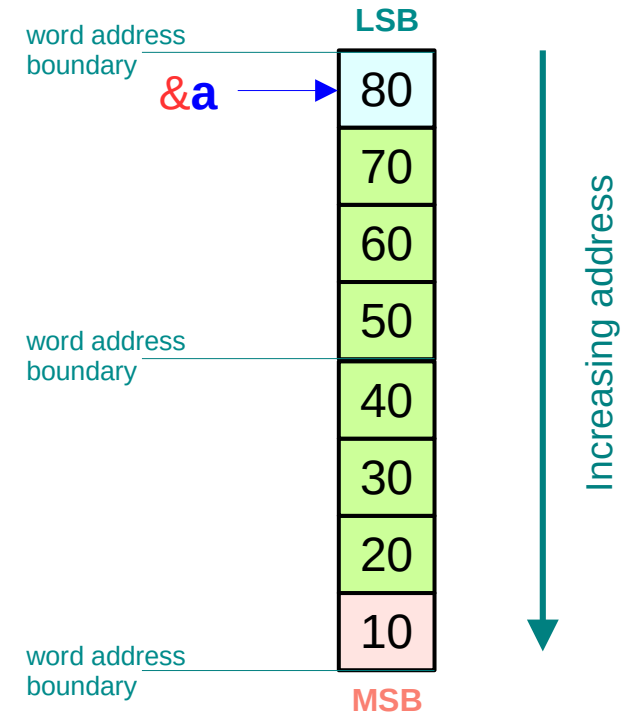
LSB

MSB (Most Significant Byte)

LSB (Least Significant Byte)

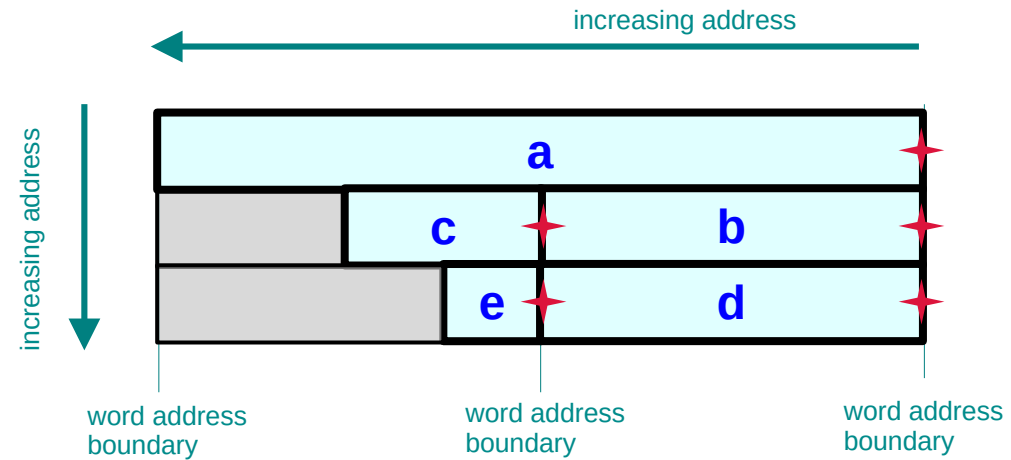
Little Endian System

store the LSB first

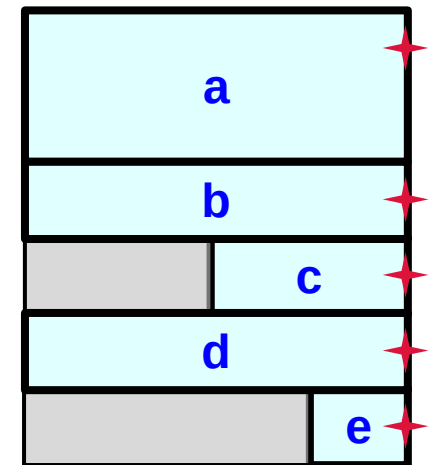


Memory alignment in a little endian system

10	20	30	40	50	60	70	80
	12	34	11	22	33	44	
		99	01	02	03	04	

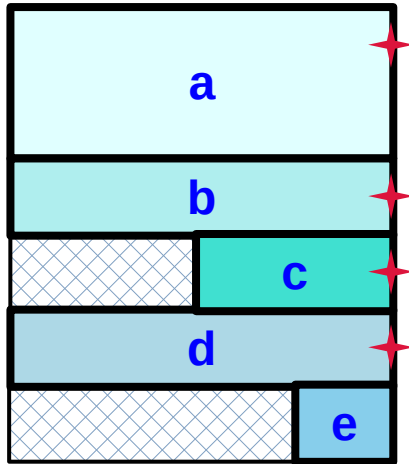


long **a** = 0x^{MSB}10203040506070^{LSB}80 ;
int **b** = 0x112233^{MSB}44^{LSB} ;
short **c** = 0x12^{MSB}34^{LSB} ;
int **d** = 0x010203^{MSB}04^{LSB} ;
char **e** = 0x99^{LSB} ;

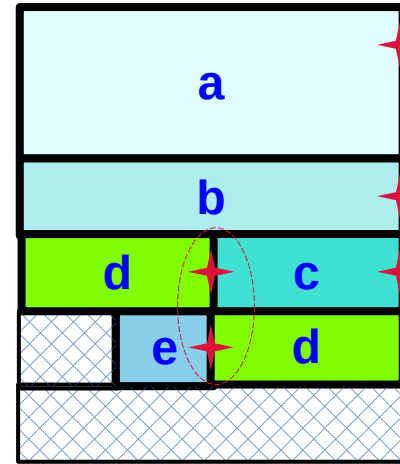


Memory alignment – performance

Memory alignment



No memory alignment



saves memory

loses performance :

accessing **d** may require two memory cycles instead of one

memory hardware constraints

```
long   a = 0x1020304050607080 ;
int    b =      0x11223344 ;
short  c =      0x1234 ;
int    d =      0x01020304 ;
char   e =      0x99 ;
```

MSB (pink) and LSB (blue) labels are placed above and below the hex values. Red stars are placed above and below the semicolons.

Pointer Type Cast

Assigning pointer variables

`long a;` `&a` address of a long value

`int * p;` address of an int value `p`  address of a long value `&a`

`short * q;` address of a short value `q`  address of a long value `&a`

`char * r;` address of a char value `r`  address of a long value `&a`

Type cast pointers

`long a;`

address of a long value `&a`

`int * p ;`

address of an int value

`p`

int address conversion

`(int *) &a`



`short * q ;`

address of a short value

`q`

short address conversion

`(short *) &a`



`char * r ;`

address of a char value

`r`

char address conversion

`(char *) &a`

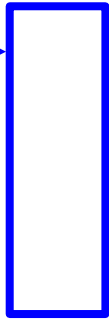


View windows – type cast pointer

integer
view window

int *

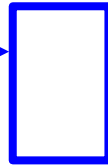
p



short
view window

short *

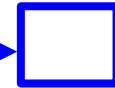
q



character
view window

char *

r



Little Endian
Memory System

&a

LSB

80

70

60

50

40

30

20

10

MSB

Increasing address

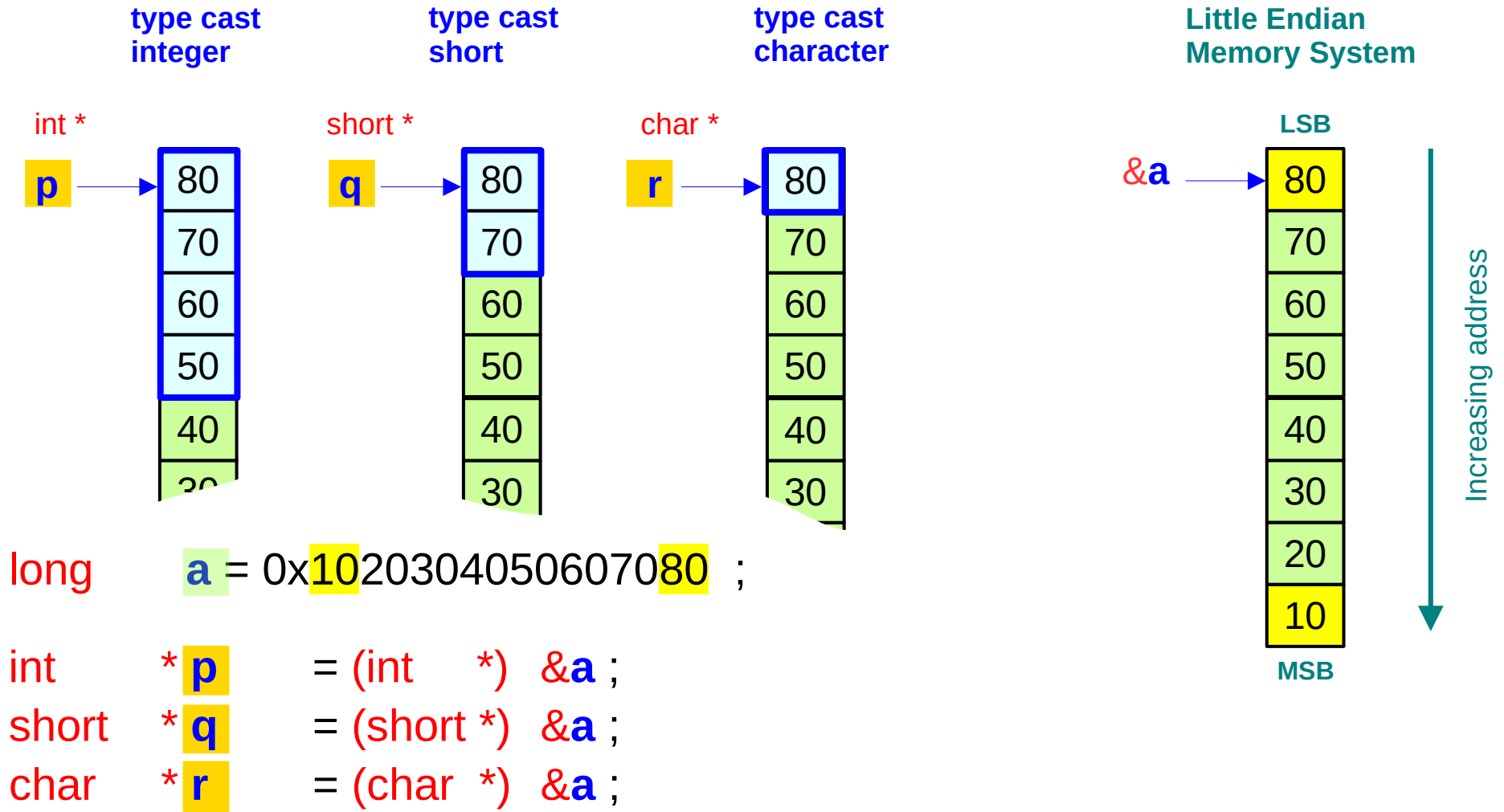
```
long    a = 0x1020304050607080 ;
```

```
int     * p ;
```

```
short   * q ;
```

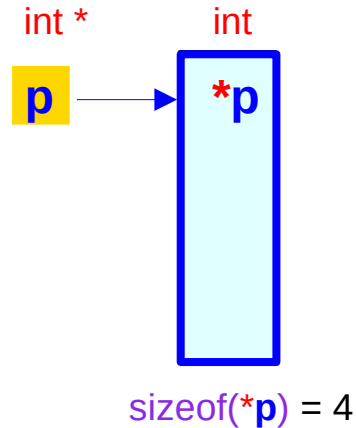
```
char    * r ;
```

Applying type cast pointers to a memory location

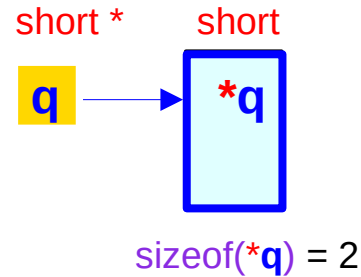


De-referencing type cast pointers

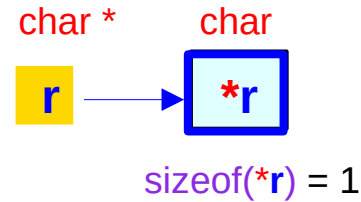
type cast
integer pointer



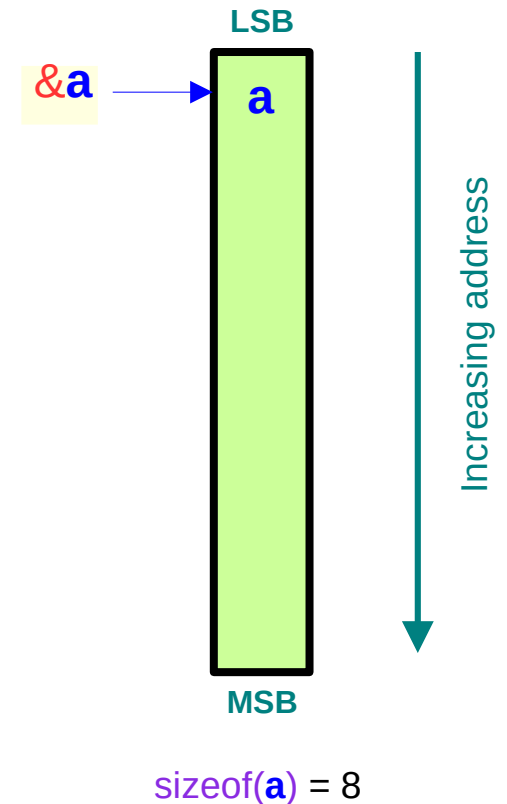
type cast
short pointer



type cast
character pointer



Little Endian
Memory System

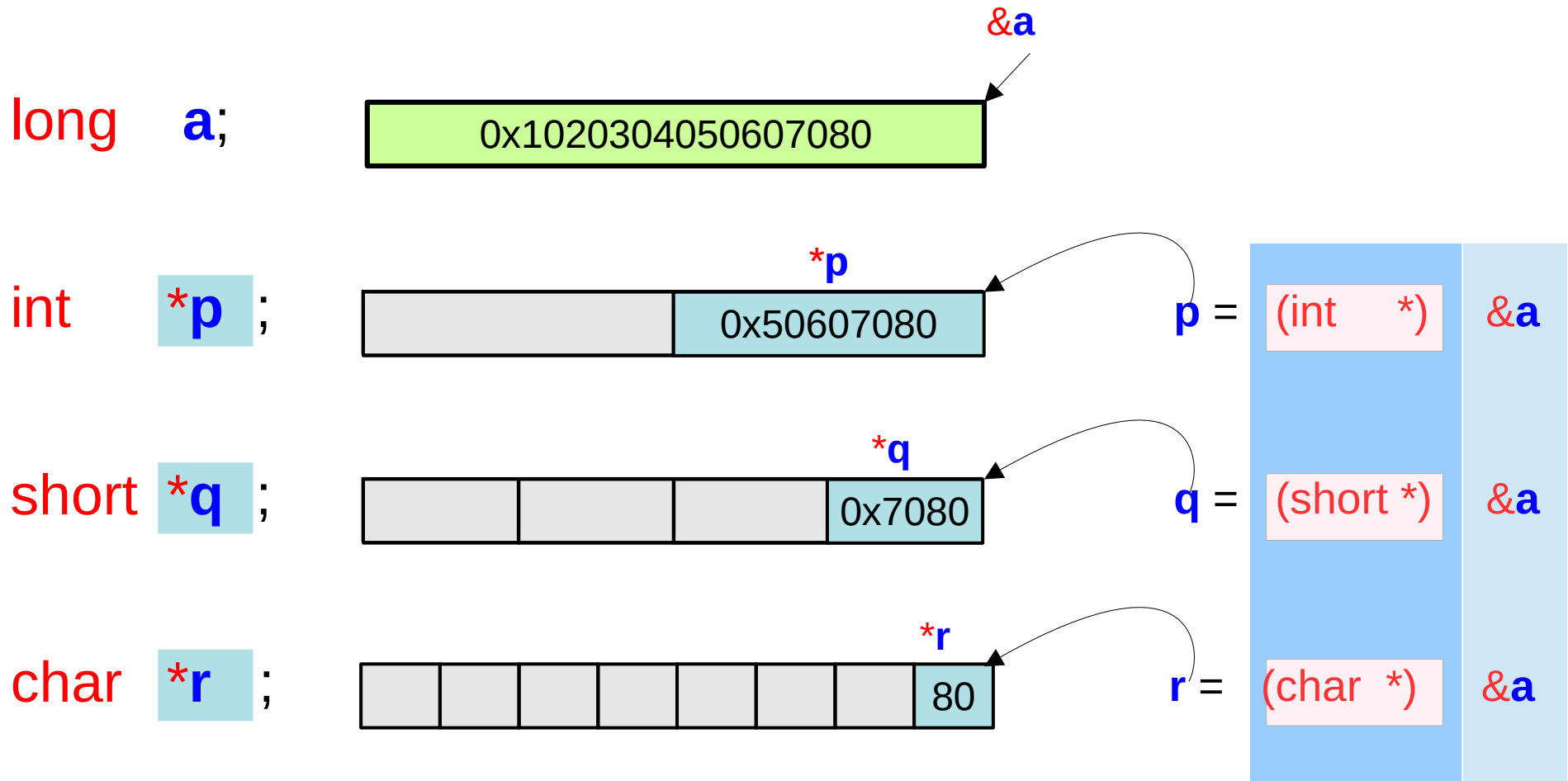


```
long a = 0x1020304050607080 ;
```

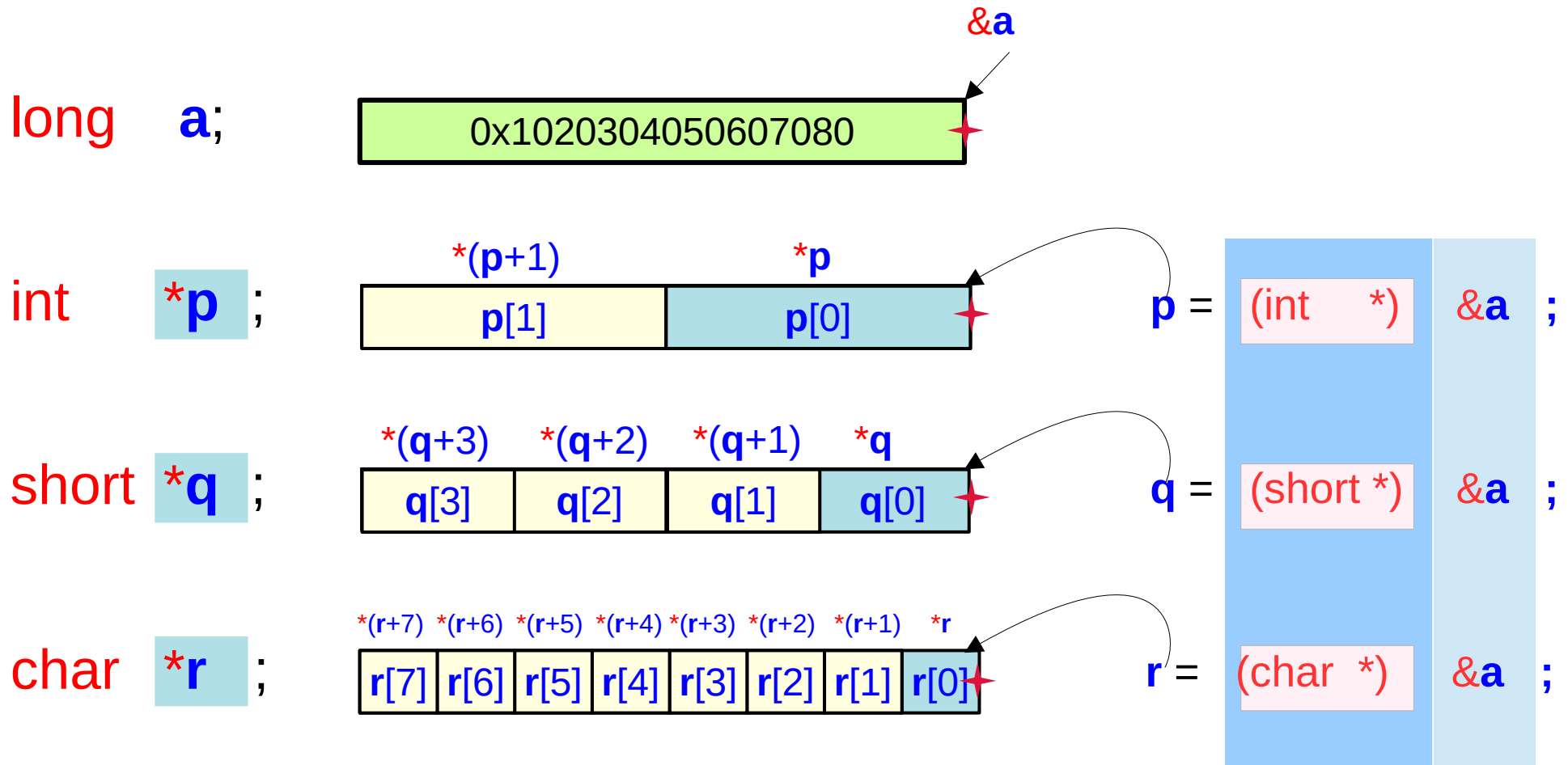
```
int *p = (int *) &a ;  
short *q = (short *) &a ;  
char *r = (char *) &a ;
```

```
*p ≡ 0x50607080  
*q ≡ 0x7080  
*r ≡ 0x80
```

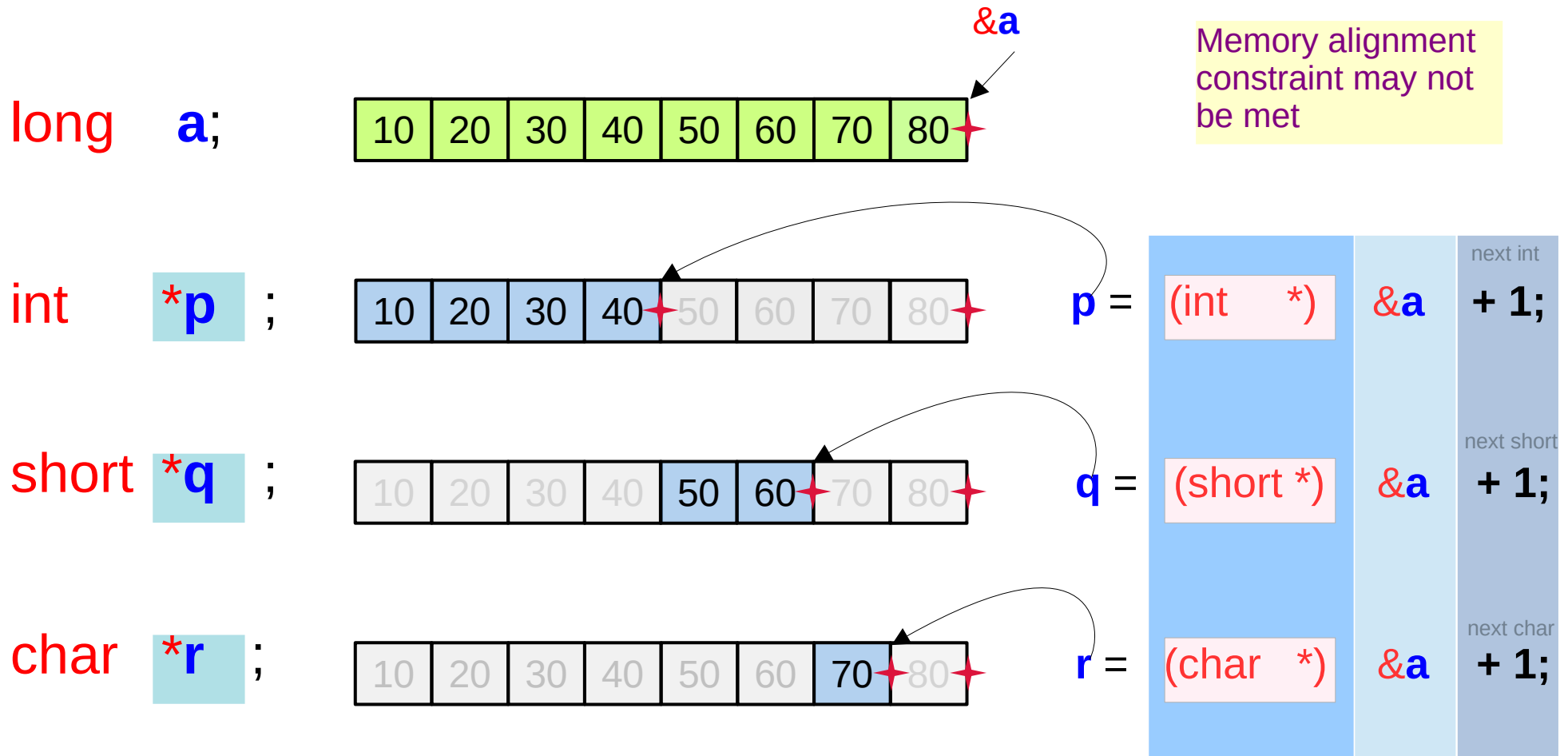
Re-interpretation of memory data – case I



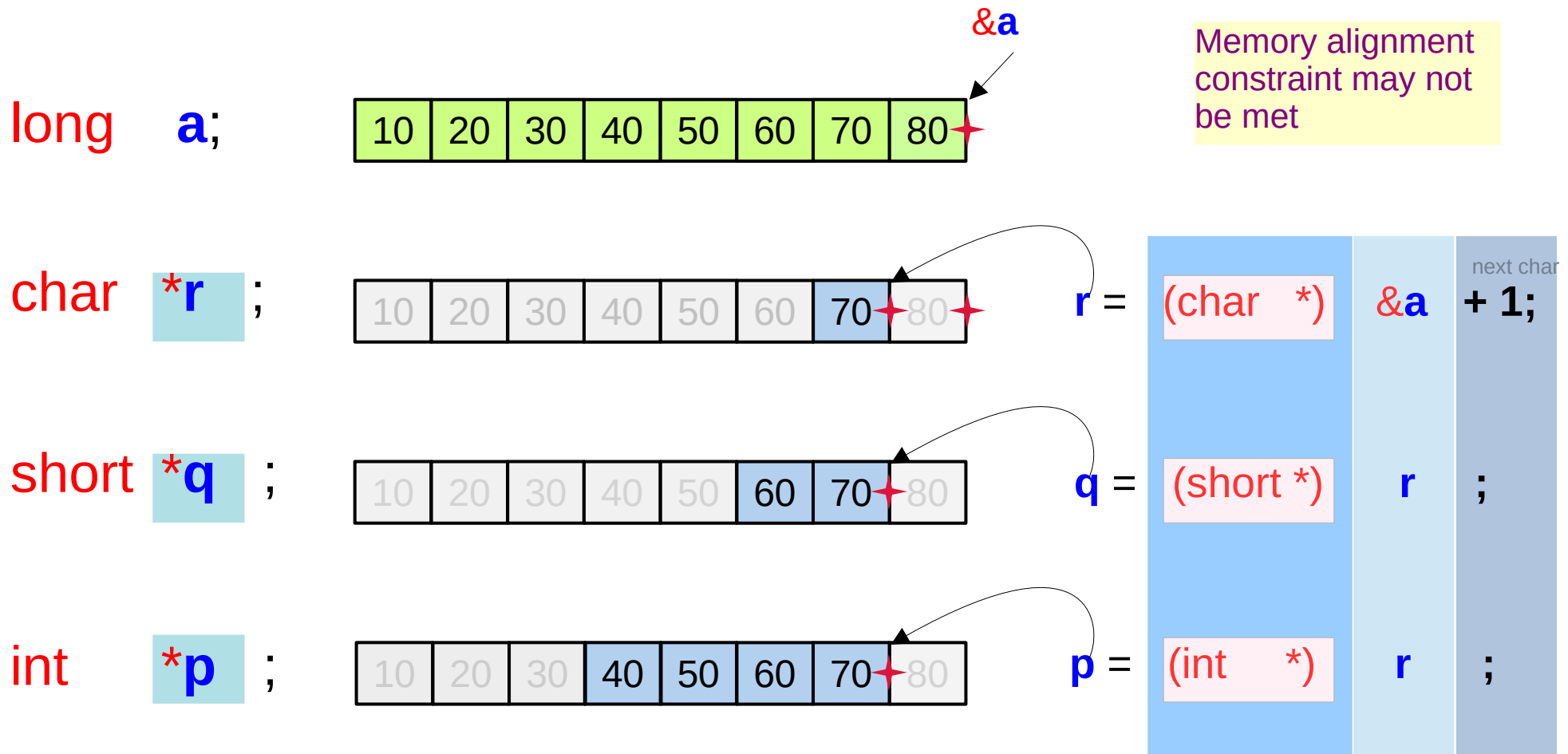
Re-interpretation of memory data – case I



Re-interpretation of memory data – case II



Re-interpretation of memory data – case III



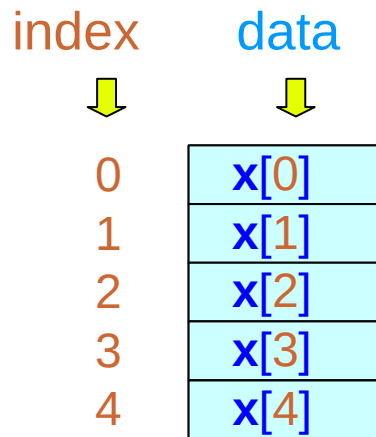
Arrays

Accessing array elements – using abstract addresses

```
int      x[5] ;
```

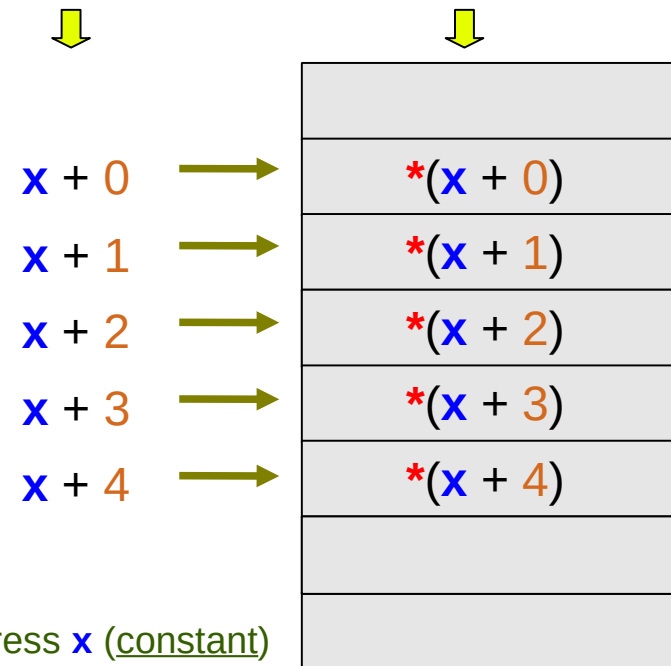
x holds the *starting address*
of 5 consecutive *int* variables

5 int variables



abstract
address

data



cannot change address x (constant)

Accessing array elements – using byte addresses

```
int    x[5];
```

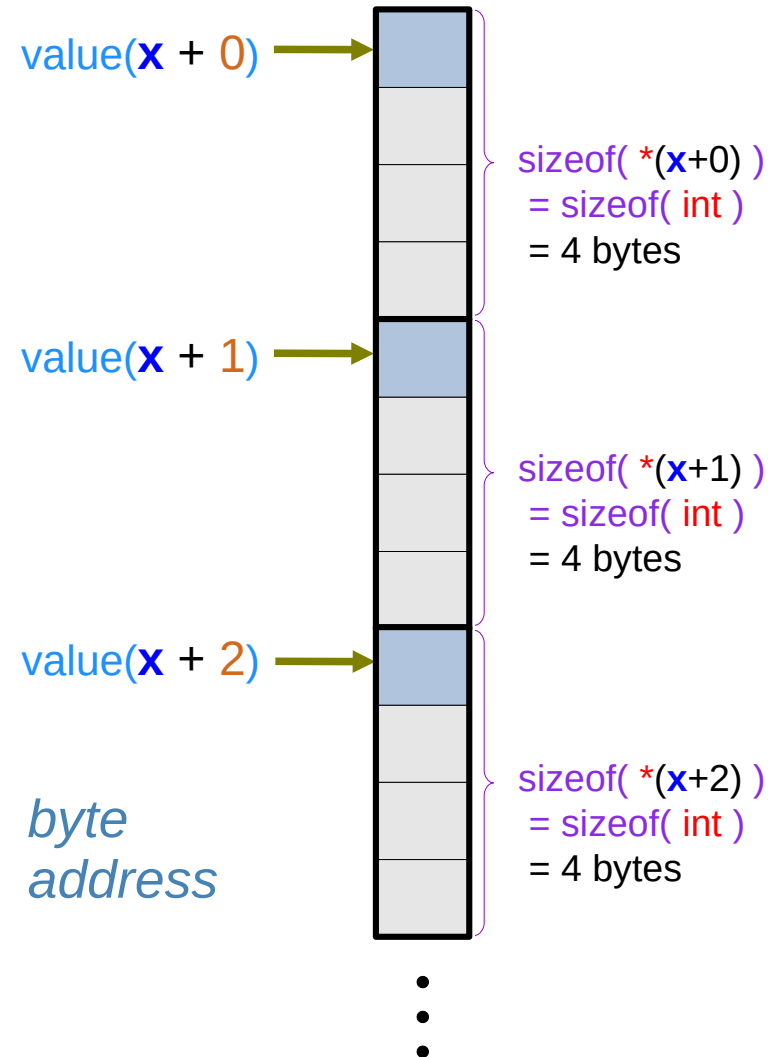
x holds the *starting address* of 5 consecutive *int* variables

5 int variables

$value(x + 0)$	=	$value(x)$	+	$0 * sizeof(*x)$
$value(x + 1)$	=	$value(x)$	+	$1 * sizeof(*x)$
$value(x + 2)$	=	$value(x)$	+	$2 * sizeof(*x)$
⋮		⋮		⋮

↓ ↓ ↓

byte address byte address $sizeof(int) = 4$ bytes



Two aspects of **x**

```
int      x[5];
```

x is an array of 5 integer elements

abstract data **x size**

pointer **x starting address**

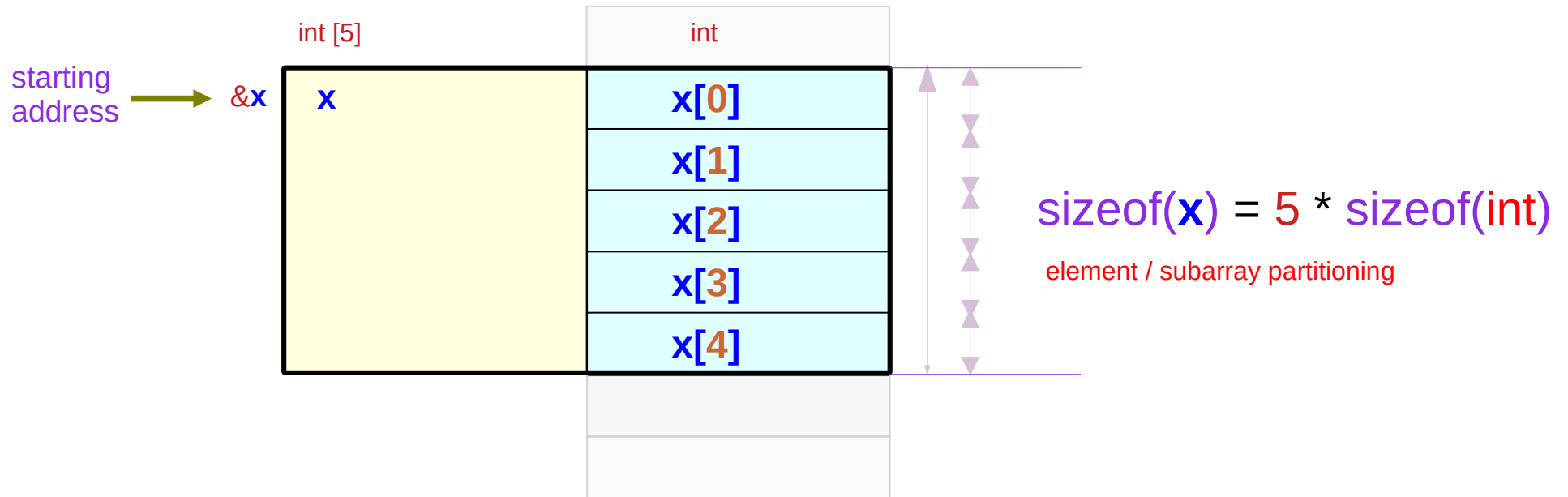
Abstract data **x**

```
int      x[5];
```

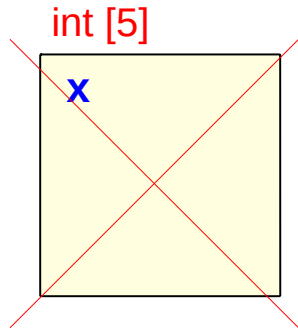
the size of **x** is the total size of 5 consecutive `int` variables

abstract data **x**

size



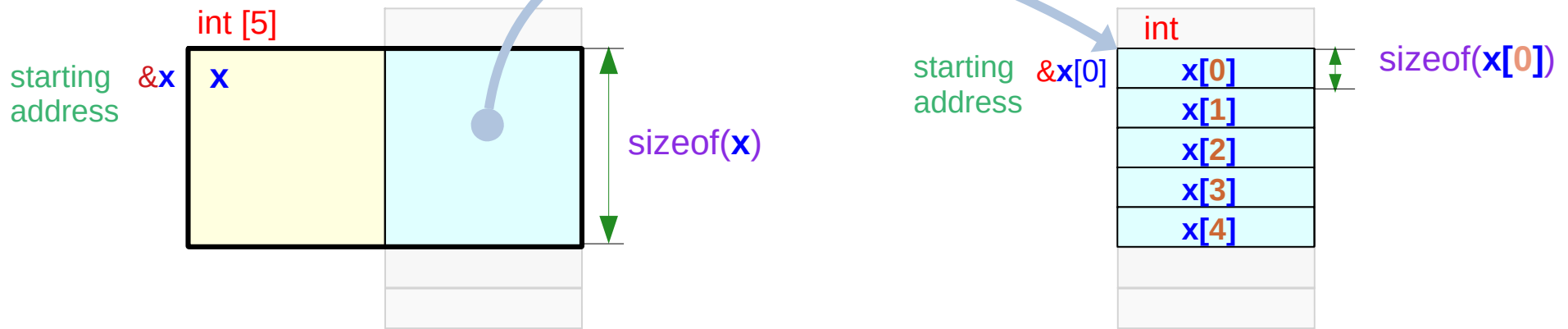
Recursive structure



`int [5]` type data must include 5 `int` elements (primitive data)

nested (recursive) structure

in a multi-dimensional array, nested sub-arrays instead of primitive data



abstract data `x`

primitive data `x[i]`

Pointer x

```
int x[5];
```

x holds the starting address of 5 consecutive int variables

$$x[i] \equiv *(x+i)$$

only for $i = 0, \dots, 4$

$$x[0] \equiv *x$$
$$\&x[0] \equiv x$$

equivalence relations

starting address

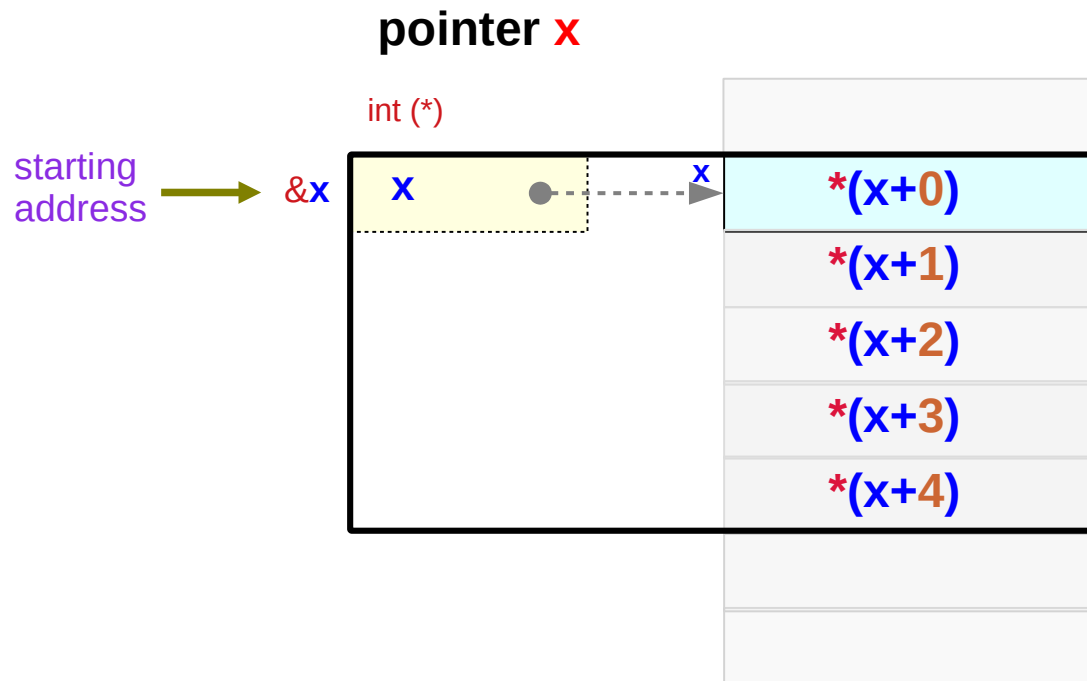
the starting address of 5 consecutive int variables

$$\text{value}(x) = \text{value}(\&x[0])$$

the starting address ($\&x$) of the array x

$$\text{value}(\&x) = \text{value}(x)$$

address replication



Aggregate Data Type – 1-d Array

```
int    x[5];
```

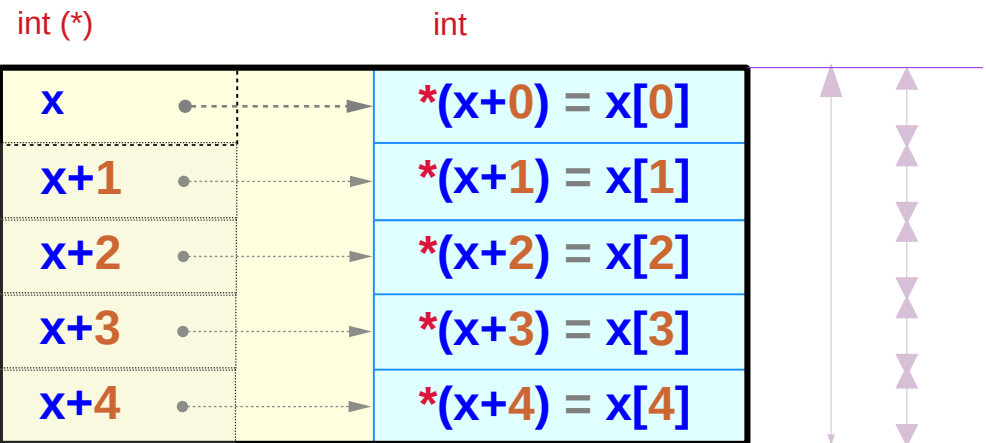
x holds the starting address of **5** consecutive **int** variables

Aggregate data with 5 elements

the start address of each element :

$\text{value}(\mathbf{x+0})$	=	$\text{value}(\mathbf{x})$	+	0	*	$\text{sizeof}(*\mathbf{x})$
$\text{value}(\mathbf{x+1})$	=	$\text{value}(\mathbf{x})$	+	1	*	$\text{sizeof}(*\mathbf{x})$
$\text{value}(\mathbf{x+2})$	=	$\text{value}(\mathbf{x})$	+	2	*	$\text{sizeof}(*\mathbf{x})$
$\text{value}(\mathbf{x+3})$	=	$\text{value}(\mathbf{x})$	+	3	*	$\text{sizeof}(*\mathbf{x})$
$\text{value}(\mathbf{x+4})$	=	$\text{value}(\mathbf{x})$	+	4	*	$\text{sizeof}(*\mathbf{x})$

$\text{value}(\mathbf{x})$ = *the start address of the array*

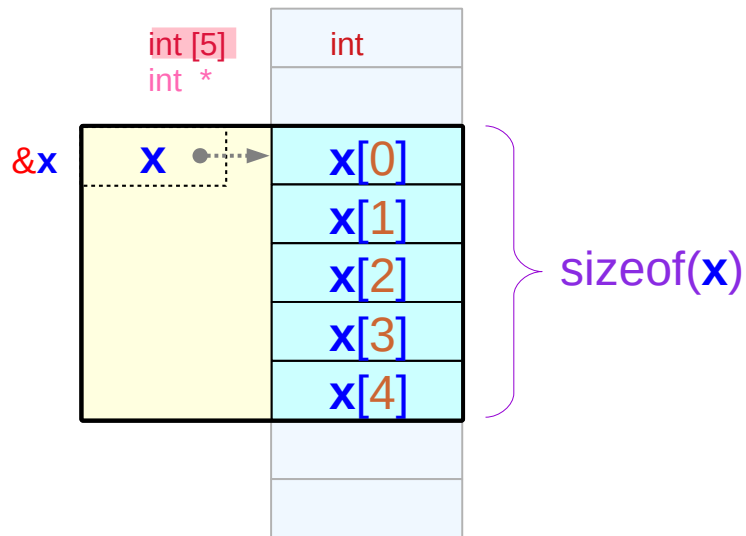


$\text{sizeof}(\mathbf{x}) = \text{sizeof}(*\mathbf{x}) * 5$

-
- **An array x and a pointer p**

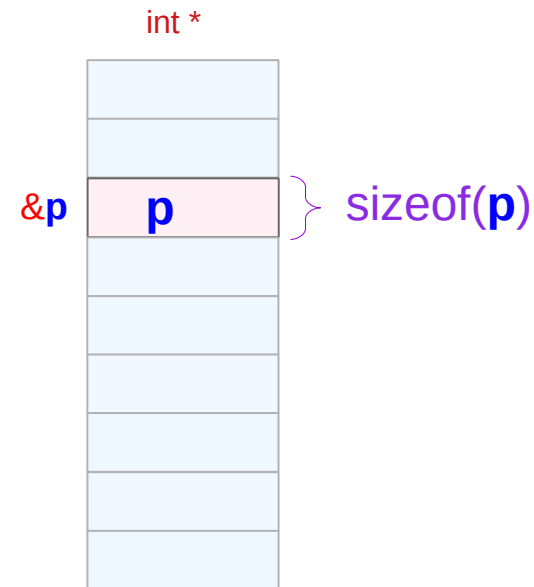
Sizes of an array **x** and a pointer **p**

```
int x [5] ;
```



$$\text{sizeof}(x) = 5 * \text{sizeof}(\text{int})$$

```
int * p ;
```



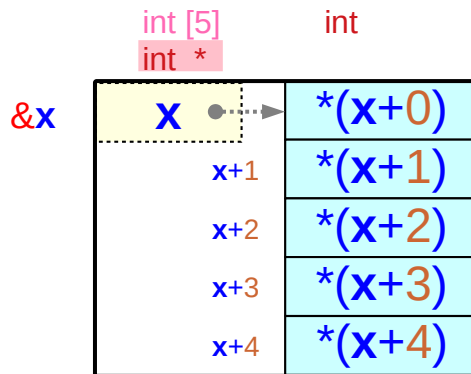
$$\begin{aligned} \text{sizeof}(p) &= \text{size of a pointer} \\ &= 4 \text{ bytes (32-bit system)} \\ &= 8 \text{ bytes (64-bit system)} \end{aligned}$$

Address values of an array **x** and a pointer **p**

```
int x [5] ;
```

x : an array variable name (constant)

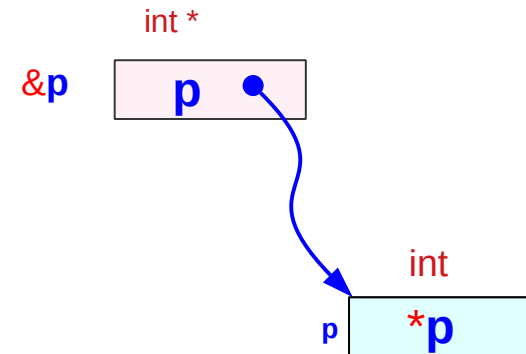
value(x) : the starting address of
5 consecutive **int** variables



```
int * p ;
```

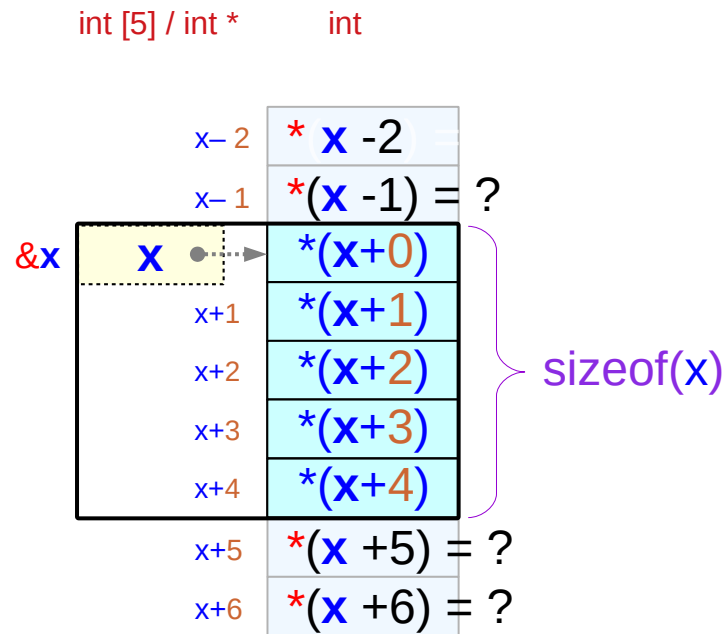
p : an variable name

value(p) : the address
of an **int** variable

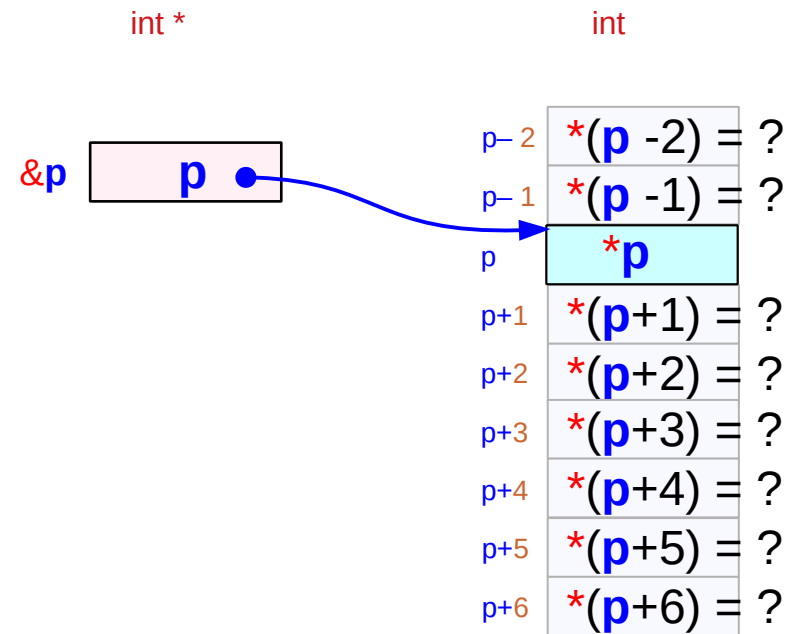


Out of range index

```
int x [5] ;
```



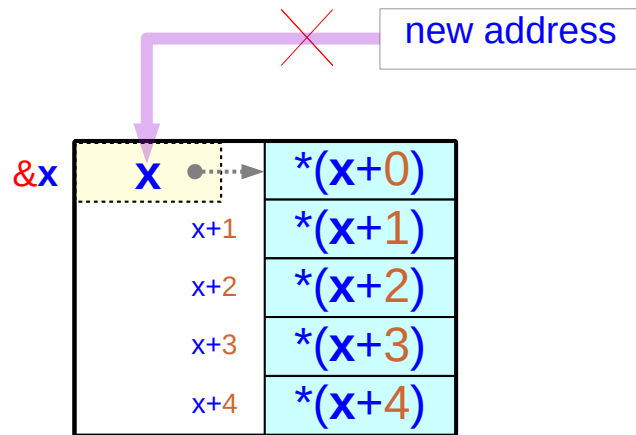
```
int * p ;
```



A programmer's responsibility

Assignment of a new address

```
int x [5] ;
```

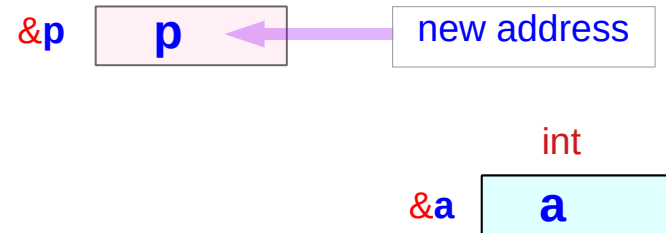


x is a constant variable
(not an *lvalue*)

x and &x give the same
value of &x[0] (address replication)

this address value(x) is assigned by
the compiler and cannot be changed

```
int * p ;
```



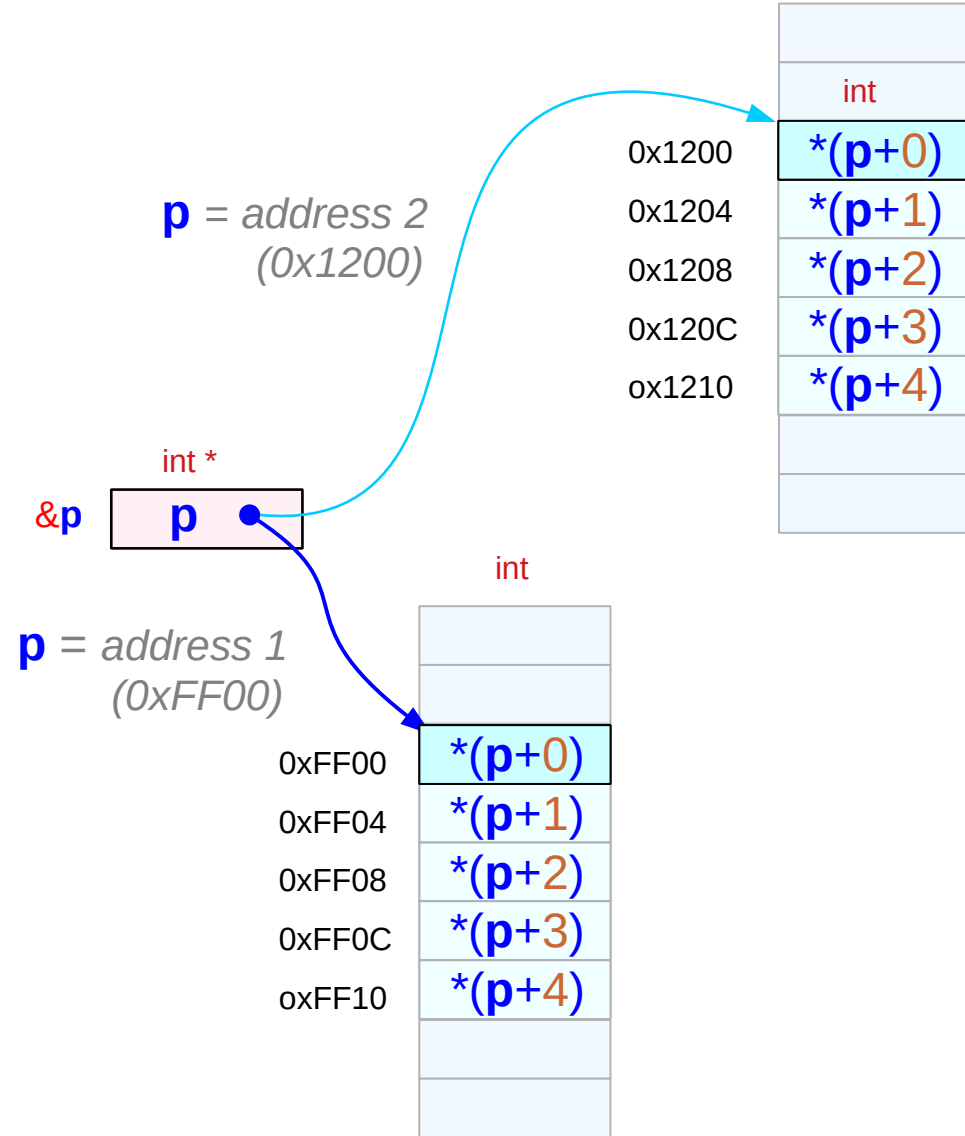
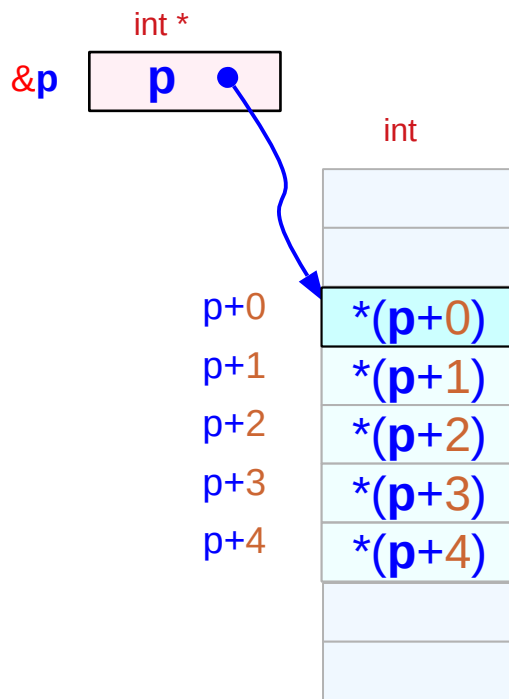
p is a pointer variable
(a *lvalue*)

a memory location is allocated
for a variable p

value(p) can be changed by an
assignment statement, i.e.,
p = &a

Pointer variable can point different locations

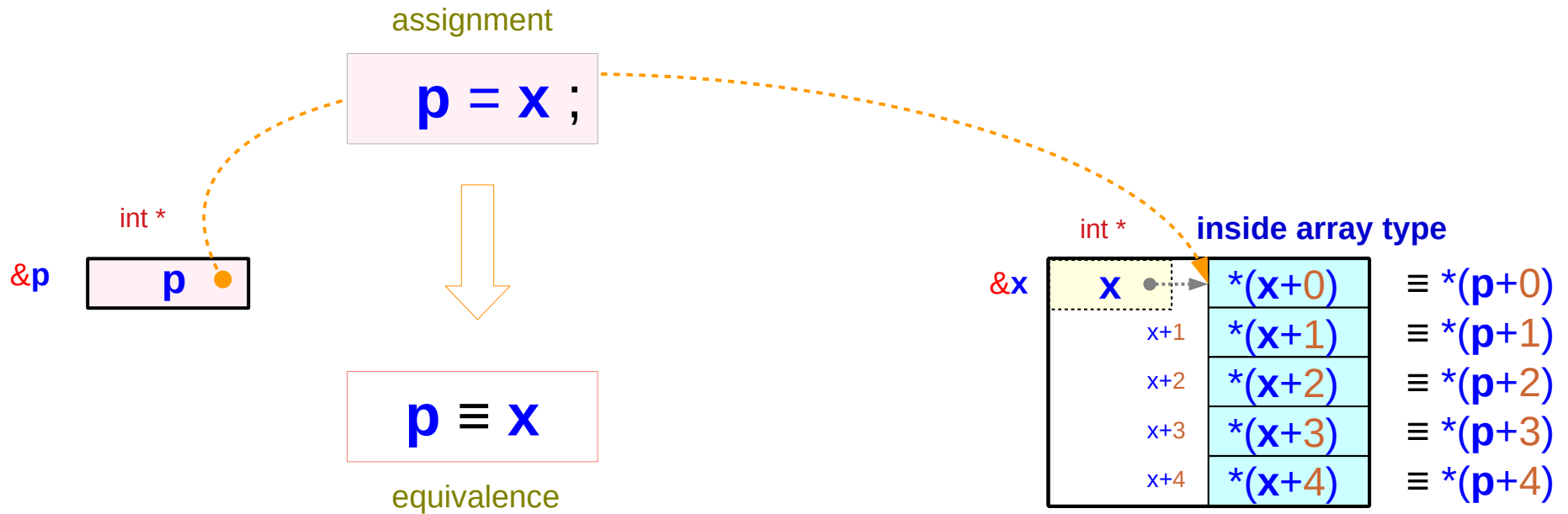
```
int * p ;
```



Pointer to an array element

```
int (*p) ;
```

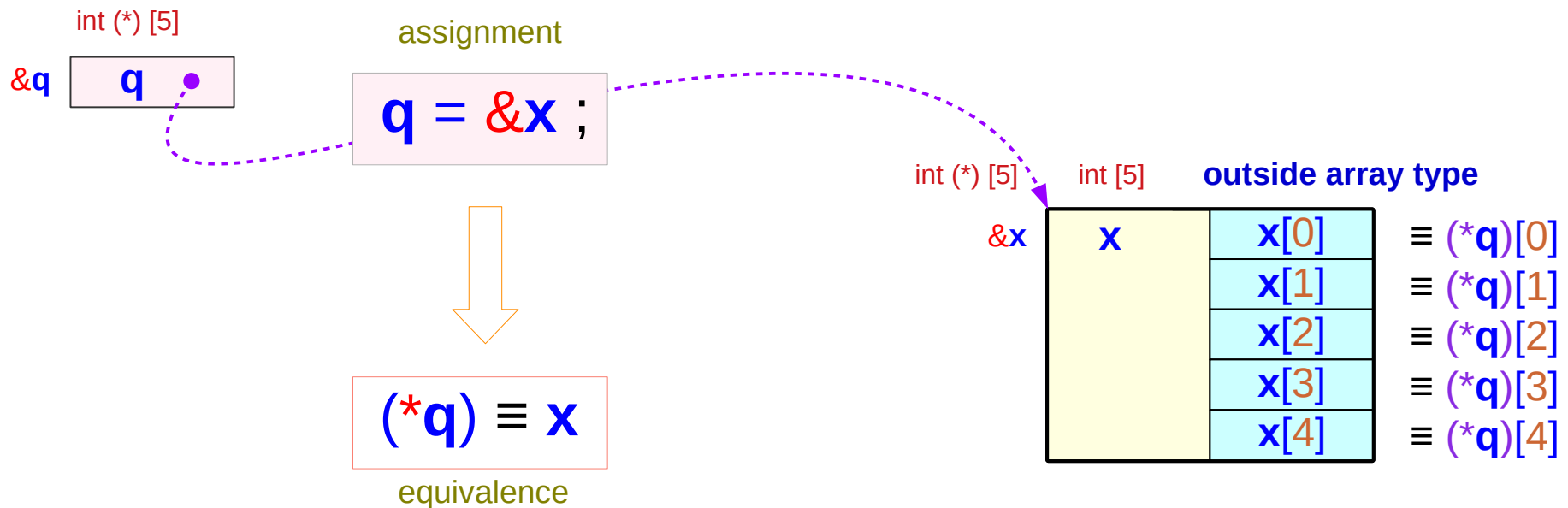
```
int x [5] ;
```



Pointer to an array

```
int (*q) [5] ;
```

```
int x [5] ;
```



`q` is a pointer variable that points to an array with 5 integer variables

Pass by Reference

Variable Scopes

```
int func1 (int a, int b)
{
    int i, int j;
    ...
    ...
    ...
}
```

i and j's
variable scope



cannot access
each other

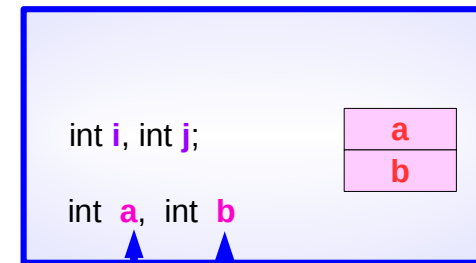
```
int main ()
{
    int x, int y;
    ...
    ...
    func1 ( 10, 20 );
    ...
    ...
}
```

x and y's
variable scope

Only **top** stack frame is active
and its variable can be accessed

Communications are performed
only through the **parameter** variables

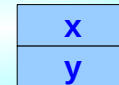
func1's
Stack
Frame



(10, 20)

main's
Stack
Frame

int x, int y;



Pass by Reference

```
int func1 (int* a, int* b)
{
  int i, int j;
  ...
  ...
  ...
  ...
}
```

x and **y** are made known to **func1**
func1 can read / write **x** and **y**
through their addresses



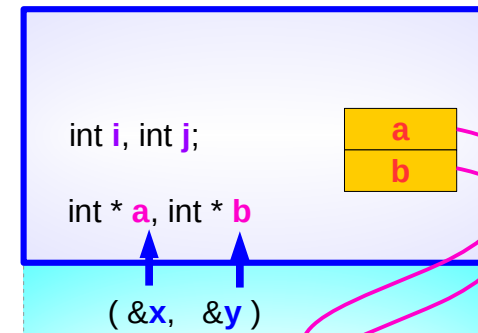
a = &**x**
b = &**y**

```
int main ()
{
  int x, int y;
  ...
  ...
  func1 ( &x, &y );
  ...
  ...
}
```

x and **y**'s
variable scope

***a**
***b**

func1's
Stack
Frame



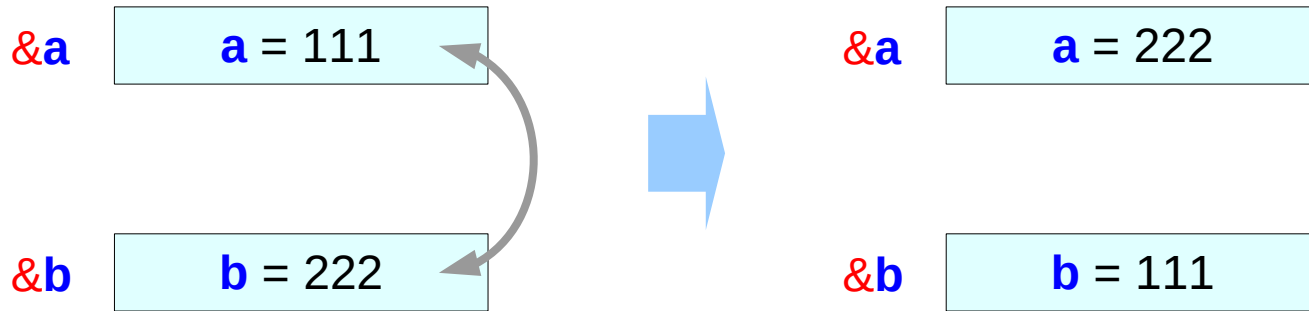
main's
Stack
Frame



***a**
***b**

Example :
swapping integers
swapping pointers

Swapping integers



```
int a, b;
```

```
swap( &a, &b );
```

function call

```
swap( int *, int * );
```

function prototype

Pass by integer reference

```
void swap(int *p, int *q) {  
    int tmp;  
  
    tmp = *p;  
    *p = *q;  
    *q = tmp;  
}
```

```
int a = 111, b = 222;
```

...

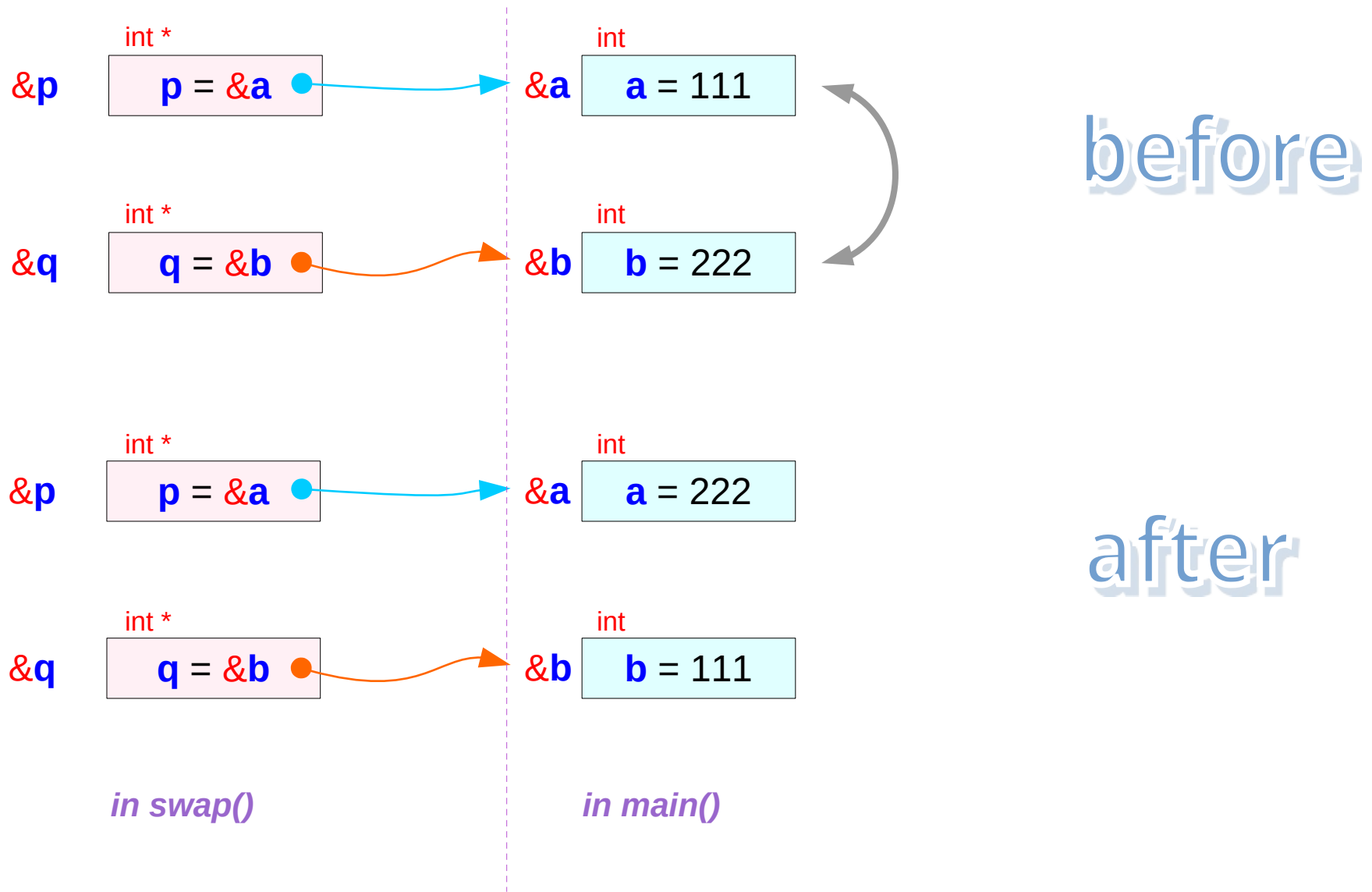
```
swap( &a, &b );
```

int *	p
int	*p

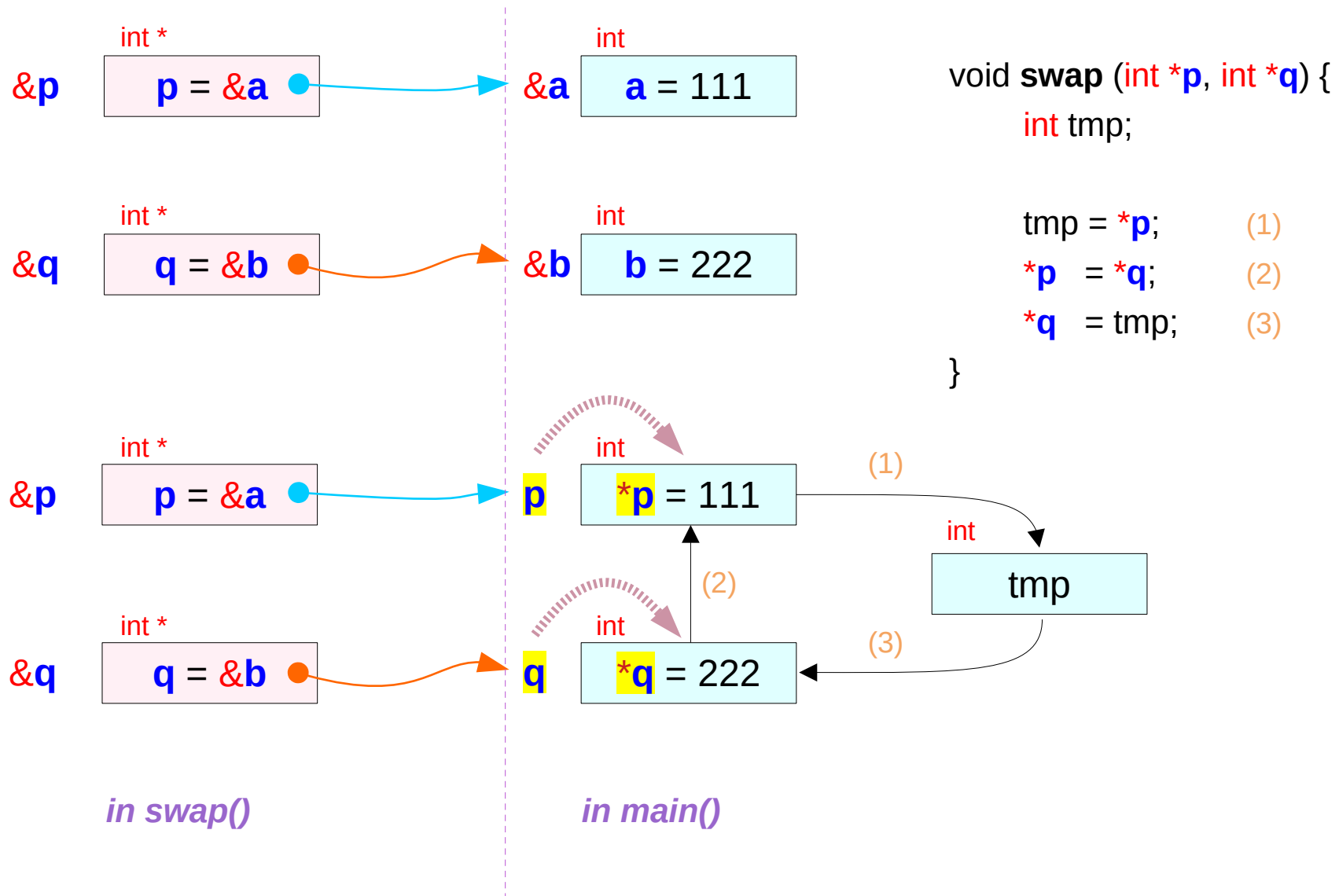
int *	q
int	*q

int	tmp
-----	-----

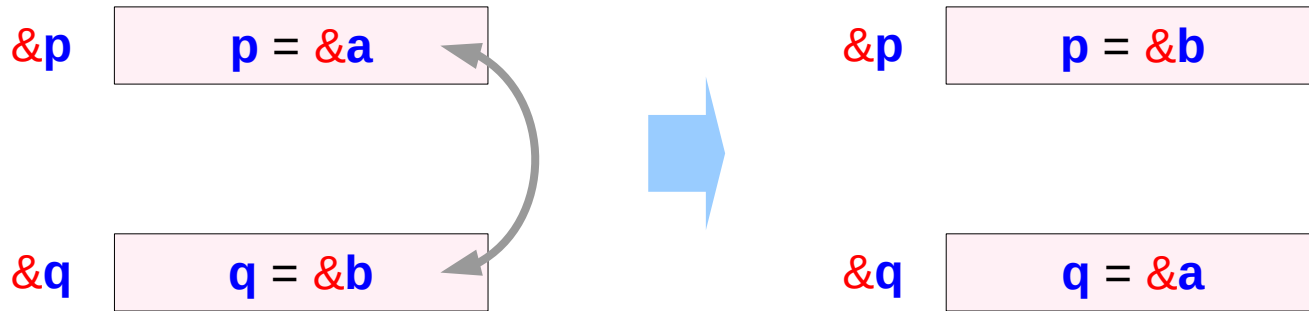
Swapping integers



Swapping integers via integer pointers



Swapping integer pointers



```
int *p, *q ;  
pswap ( &p, &q );  
void pswap( int **, int ** );
```

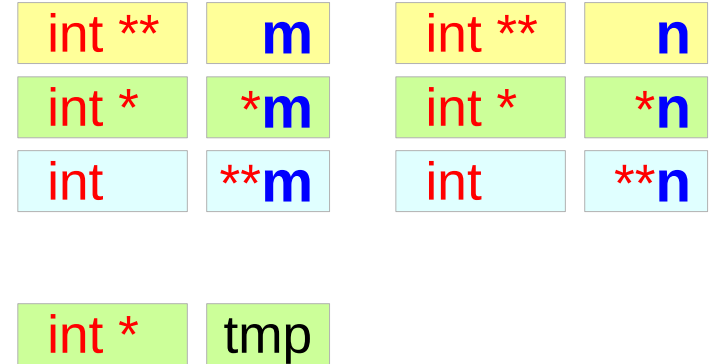
function call

function prototype

Pass by integer pointer reference

```
void pswap (int **m, int **n)
{
    int * tmp;

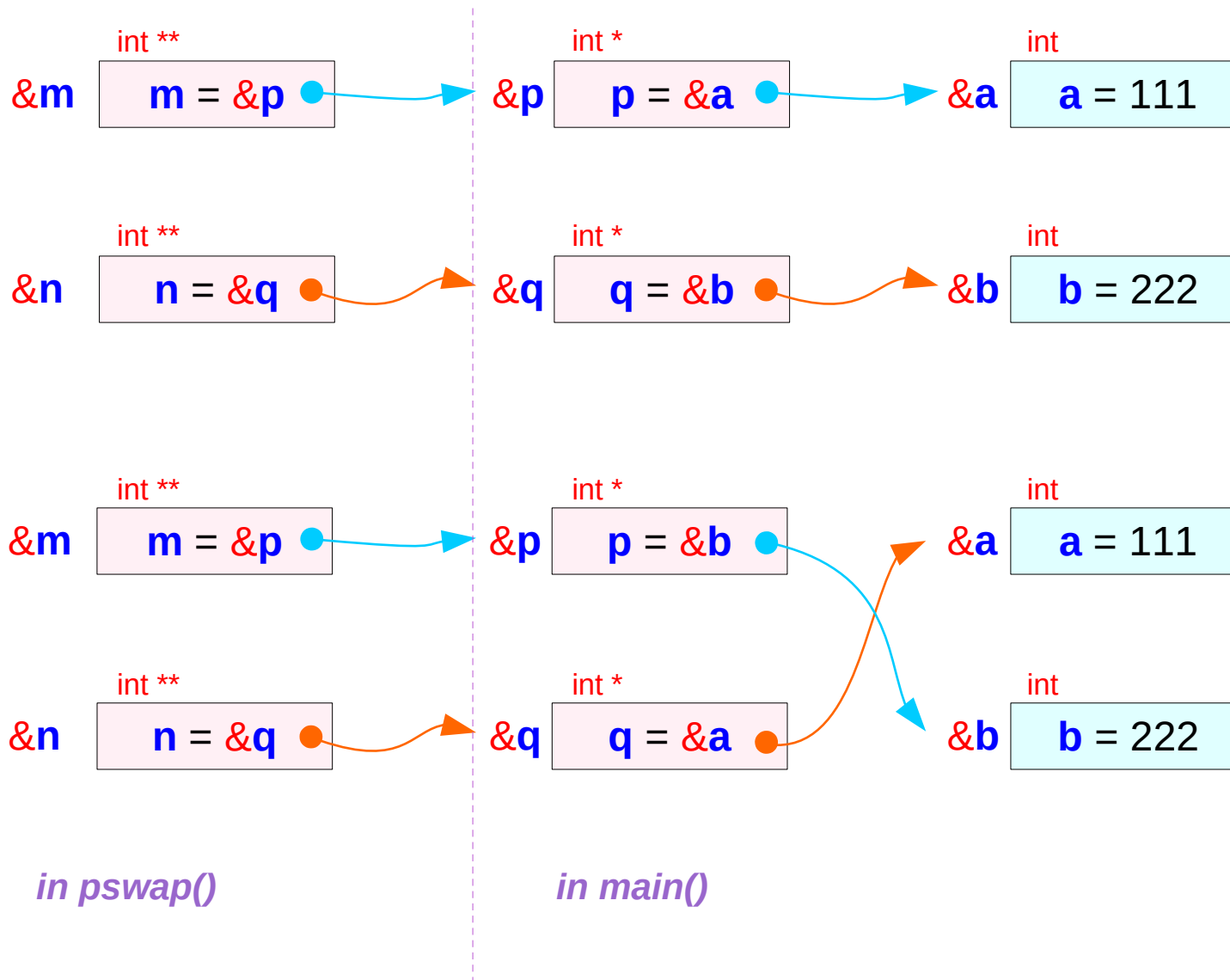
    tmp = *m;
    *m = *n;
    *n = tmp;
}
```



```
int  a = 111, b = 222;
int *p = &a, *q = &b;
...
pswap ( &p, &q );
```

```
int **  m
int *   *m
int     **m
```

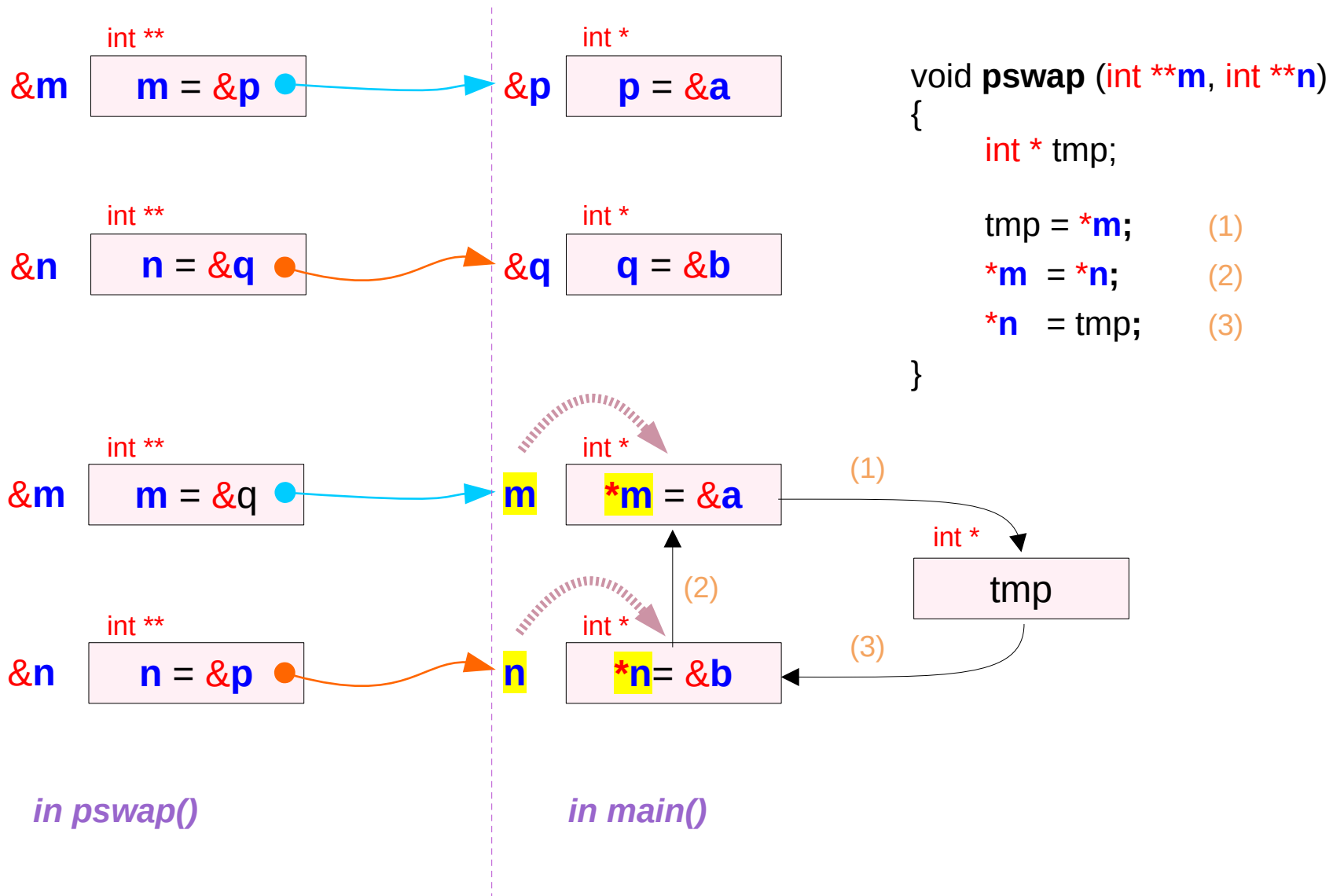
Swapping integer pointers



before

after

Swapping integer pointers via double pointers



References

- [1] Essential C, Nick Parlante
- [2] Efficient C Programming, Mark A. Weiss
- [3] C A Reference Manual, Samuel P. Harbison & Guy L. Steele Jr.
- [4] C Language Express, I. K. Chun