Reduction Clause (5A)

- Loop
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Young Won Lim 9/29/21 reduction (operator: list)

Performs a **reduction** on all **scalar variables** in list using the <u>specified</u> **operator**. Reduction variables in list are separated by commas.

A private copy of each variable in list is created for each thread.

At the <u>end</u> of the statement block,

the **final** values of all **private copies** of the reduction variable are <u>combined</u> in a manner appropriate to the **operator**, and the **result** is placed back in the **original value** of the **shared reduction variable**.

https://www.ibm.com/docs/en/xl-c-aix/13.1.2?topic=processing-pragma-omp-section-pragma-omp-sections

Reduction (1)

#include <stdio.h>
#include <omp.h>

```
int main(void)
{
    int sum = 100;
    #pragma omp parallel for
    for (int i = 1; i <= 4; i++)
    {
        sum += i;
    }
    printf("sum is %d\n", sum);
    return 0;
}</pre>
```

Reduction (2)

Because sum is a shared variable in threads, so we need to use synchronization to protect accessing it:

```
#pragma omp parallel for
for (int i = 1; i <= 4; i++)
{
    #pragma omp critical
    sum += i;
}</pre>
```

But this will cause losing the advantage of using parallelism. The other method is using reduction clause:

reduction(reduction-identifier : list)

Reduction (3)

#include <stdio.h>
#include <omp.h>

```
int main(void)
```

```
{
```

}

```
int sum = 100;
```

printf("Before parallelism, sum's address is %p\n", &sum);

```
#pragma omp parallel for reduction(+ : sum)
for (int i = 1; i <= 4; i++)
{
    printf("sum's address in thread %d is %p, value is %d\n", omp_get_thread_num(), &sum, sum);
    sum += i;
}
printf("After parallelism, sum's address is %p, and value is %d\n", &sum, sum);
return 0;</pre>
```

Reduction (4)

gcc -fopenmp parallel.c
./a.out
Before parallelism, sum's address is 0x7ffcc880baf0
sum's address in thread 3 is 0x7f6baea5ee20, value is 0
sum's address in thread 2 is 0x7f6baf25fe20, value is 0
sum's address in thread 0 is 0x7ffcc880ba90, value is 0
sum's address in thread 1 is 0x7f6bafa60e20, value is 0
After parallelism, sum's address is 0x7ffcc880baf0, and value is 110

reduction (operator: list)

For example, when the **max operator** is specified, the original reduction variable value combines with the <u>final</u> **values** of the **private copies** by using the following expression:

original_reduction_variable =
 original_reduction_variable < private_copy
 ? private_copy
 : original_reduction_variable;</pre>

https://www.ibm.com/docs/en/xl-c-aix/13.1.2?topic=processing-pragma-omp-section-pragma-omp-sections

References

- [1] en.wikipedia.org
- [2] M Harris, http://beowulf.lcs.mit.edu/18.337-2008/lectslides/scan.pdf