Overflow Flag

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Outline

- Based on
- The Overflow flag
 - TOC: The Overflow flag
 - The overflow flag in unsigned and signed computations
 - Rules for the overflow flag
 - Method 1 for computing the overflow flag
 - Method 2 for computing the overflow flag
 - More examples of the overflow flag

Based on

 The CARRY flag and OVERFLOW flag in binary arithmetic lan! D. Allen - idallen@idallen.ca - www.idallen.com https://teaching.idallen.com/dat2343/10f/notes/ 040 overflow.ttx

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Compling 32-bit program on 64-bit gcc

- gcc -v
- gcc -m32 t.c
- sudo apt-get install gcc-multilib
- sudo apt-get install g++-multilib
- gcc-multilib
- g++-multilib
- gcc -m32
- objdump -m i386

TOC: Overflow flag

- Overflow flag in unsigned and signed computations
- Rules for the overflow flag
- Method 1 for computing the overflow flag
- Method 2 for computing the overflow flag
- More examples of the overflow flag

TOC Overflow flag in unsigned and signed computations

Overflow flag

Overflow flag (1)

- overflow flag is based on signed arithmetic
- to decide if the overflow flag is turned on or off, only need to look at the sign bits (leftmost) of the three numbers

```
augend + addend = sum
minuend - subrahend = difference
```

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Overflow flag (2)

- in signed arithmetic,
 - watch the overflow flag to detect errors
 - overflow flag on means the result is wrong
 - errors can be detected by examining the <u>sign</u> of the result, in the 2's complement arithmetic (P + P → N or N + N → P)
- in unsigned arithmetic,
 - the overflow flag tells you nothing interesting

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Overflow flag (3)

- when two positive numbers are added
 - if the result is a positive, $(P + P \rightarrow P)$, then no overflow
 - if the result is a negative, $(P + P \rightarrow N)$, then overflow
- when two negative numbers are added
 - if the result is a <u>negative</u>, $(N + N \rightarrow N)$, then <u>no overflow</u>
 - if the result is a positive, $(N + N \rightarrow P)$, then overflow

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Overflow flag (4)

- adding negative (N) and positive (P) numbers cannot be wrong, because the sum is between the addends ([N, P]) .
 - if opposite signed numbers are added, then no overflow
 - both of the addends lies in the allowable range of numbers. their sum is between the opposite signed addends, therefore the sum lies also in the allowable range
 - $(P + N \rightarrow P \text{ or } N)$ no overflow always
 - $(N + P \rightarrow P \text{ or } N)$ no overflow always

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TOC Rules for the overflow flag

- the 1st rule for setting OF
- the 2nd rule for setting OF
- ullet cases for clearing OF $(1\sim6)$

Overflow flag setting and clearing conditions

		Method 1		Method 2
		ADD conditions	SUB conditions	
1	0F=1	$P + P \rightarrow N$	$P - N \rightarrow N$	$c_n \bigoplus c_{n-1} = 1$
2	0F=1	$N + N \rightarrow P$	N-P o P	$c_n \bigoplus c_{n-1} = 1$
3	0F=0	$P + P \rightarrow P$	$P - N \rightarrow P$	$c_n \bigoplus c_{n-1} = 0$
4	0F=0	$N + N \rightarrow N$	N-P o N	$c_n \bigoplus c_{n-1} = 0$
5	0F=0	$P + N \rightarrow P$	$P - N \rightarrow P$	$c_n \bigoplus c_{n-1} = 0$
6	0F=0	$P + N \rightarrow N$	P-P o N	$c_n \bigoplus c_{n-1} = 0$
7	0F=0	$N + P \rightarrow P$	$N - N \rightarrow P$	$c_n \bigoplus c_{n-1} = 0$
8	0F=0	$N + P \rightarrow N$	$N - P \rightarrow N$	$c_n \bigoplus c_{n-1} = 0$

$$+P = -(-P) = -N$$

 $+N = -(-N) = -P$

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The 1st case - setting the overflow flag

If the sum of two signed numbers with the sign bits off (0, 0) yields a result number with the sign bit on (1),

the overflow flag is turned on

signed	addition	signed	subtraction	unsigned	l addition
0100	carries				
0100	(+4)	0100	(+4)	0100	(4)
+0100	+(+4)	-1100	-(-4)	+0100 +	·(4)
01000	(-8)	01000	(-8)	01000	(8)

```
Method 1 OF = 1 when \overline{a_{n-1}} \cdot \overline{b_{n-1}} \cdot s_{n-1} for addition and \overline{a_{n-1}} \cdot b_{n-1} \cdot s_{n-1} for subtraction Method 2 OF = c_n \bigoplus c_{n-1} c_4 \bigoplus c_3 = 0 \bigoplus 1 = 1
```

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The 2nd case - setting the overflow flag

If the sum of two numbers with the sign bits on (1, 1) yields a result number with the sign bit off (0)

the overflow flag is turned on.

signed	addition	signed s	subtraction	unsigne	ed addition
1001	carries				
1001	(-7)	1001	(-7)	1001	(9)
+1001	+(-7)	-0111 -	-(+7)	+1001	+(9)
10010	(2)	10010	(2)	10010	(18)

```
\begin{array}{lll} \text{Method 1} & \text{OF = 1} & \text{when } a_{n-1} \cdot \underline{b_{n-1}} \cdot \overline{s_{n-1}} \text{ for addition} \\ & \text{and } a_{n-1} \cdot \overline{b_{n-1}} \cdot \overline{s_{n-1}} \text{ for subtraction} \\ \text{Method 2} & \text{OF = } c_n \bigoplus c_{n-1} & c_4 \bigoplus c_3 = 1 \bigoplus 0 = 1 \end{array}
```

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The 3rd case - clearing the overflow flag

If the sum of two signed numbers with the sign bits off (0, 0) yields a result number with the sign bit off (0),

the overflow flag is turned off

signed addition signed subtraction unsigned addition 0011 carries 0011 (+3) (+3)0011 (3) 0011 +0011 + (+3)-1101 - (-3)+0011 + (3)00110 (+6) 00110 (+6)00110 (6)

```
Method 1 OF = 0
                                                 when \overline{a_{n-1}} \cdot \overline{b_{n-1}} \cdot \overline{s_{n-1}} for addition
                                                  and \overline{a_{n-1}} \cdot b_{n-1} \cdot \overline{s_{n-1}} for subtraction
Method 2 OF = c_n \oplus c_{n-1} c_4 \oplus c_3 = 0 \oplus 0 = 0
```

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The 4th case - clearing the overflow flag

If the sum of two signed numbers with the sign bits on (1, 1) yields a result number with the sign bit on (1),

the overflow flag is turned off

addition	signed	subtraction	unsign	ed addition
carries				
(-3)	1101	(-3)	1101	(13)
+(-3)	-0011	-(+3)	+1101	+(13)
(-6)	11010	(-6)	11010	(26)
	carries (-3) +(-3)	carries (-3) 1101 +(-3) -0011	carries (-3) 1101 (-3) +(-3) -0011 -(+3)	carries (-3) 1101 (-3) 1101 +(-3) -0011 -(+3) +1101

```
Method 1 OF = 0 when a_{n-1} \cdot b_{n-1} \cdot s_{n-1} for addition and a_{n-1} \cdot \frac{b_{n-1}}{b_{n-1}} \cdot s_{n-1} for subtraction Method 2 OF = c_n \oplus c_{n-1} or c_4 \oplus c_3 = 1 \oplus 1 = 0
```

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The 5th case - clearing the overflow flag

If the sum of two signed numbers with the sign bits off and on (0, 1) yields a result number with the sign bit off (0),

the overflow flag is turned off

signed	l addition	signed s	ubtraction	unsigne	ed addition
1100	carries				
	(+4)	0100	` '	0100	. ,
+1101	+(-3)	-0011 -	(+3)	+1101	+(13)
10001	(+1)	10001	(+1)	10001	(17)

$$\begin{array}{ll} \text{Method 1} & \text{OF = 0} & \text{when } \overline{a_{n-1}} \cdot \underline{b_{n-1}} \cdot \overline{s_{n-1}} \text{ for addition} \\ & \text{and } \overline{a_{n-1}} \cdot \overline{b_{n-1}} \cdot \overline{s_{n-1}} \text{ for subtraction} \\ \text{Method 2} & \text{OF = } c_n \bigoplus c_{n-1} & c_4 \bigoplus c_3 = 1 \bigoplus 1 = 0 \end{array}$$

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The 6th case - clearing the overflow flag

If the sum of two signed numbers with the sign bits off and on (0, 1) yields a result number with the sign bit on (1),

the overflow flag is turned off

signed	addition	signed	subtraction	unsign	ed addition
0000	carries				
0011	(+3)	0011	(+3)	0011	(3)
+1100	+(-4)	-0100	-(+4)	+1100	+(12)
01111	(-1)	01111	(-1)	01111	(15)

Method 1 OF = 0 when
$$\overline{a_{n-1}} \cdot \underline{b_{n-1}} \cdot s_{n-1}$$
 for addition and $\overline{a_{n-1}} \cdot \overline{b_{n-1}} \cdot s_{n-1}$ for subtraction Method 2 OF = $c_n \oplus c_{n-1}$ or $c_4 \oplus c_3 = 0 \oplus 0 = 0$

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The 7th case - clearing the overflow flag

If the sum of two signed numbers with the sign bits on and off (1, 0) yields a result number with the sign bit off (0),

the overflow flag is turned off

signed	addition	signed sub	otraction	unsigne	ed addition
1100	carries				
1101	(-3)	0011 (-	-3)	1101	(13)
+0100	(+4)	-1100 -(-	-4)	+0100	+(4)
10001	(+1)	10001 (+	-1)	10001	(17)

$$\begin{array}{ll} \text{Method 1} & \text{OF = 0} & \text{when } a_{n-1} \cdot \overline{b_{n-1}} \cdot \overline{s_{n-1}} \text{ for addition} \\ & \text{and } a_{n-1} \cdot b_{n-1} \cdot \overline{s_{n-1}} \text{ for subtraction} \\ \text{Method 2} & \text{OF = } c_n \bigoplus c_{n-1} & c_4 \bigoplus c_3 = 1 \bigoplus 1 = 0 \end{array}$$

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The 8th case - clearing the overflow flag

If the sum of two signed numbers with the sign bits on and off (1, 0) yields a result number with the sign bit on (1),

the overflow flag is turned off

signed add:	ition signe	d subtraction	unsigned addition
0000 car	ries		
1100 (-4	4) 0100	(-4)	1100 (12)
+0011 +(+;	-1101	-(-3)	+0011 +(3)
01111 (-:	1) 01111	(-1)	01111 (15)

```
Method 1 OF = 0 when a_{n-1} \cdot \overline{b_{n-1}} \cdot s_{n-1} and a_{n-1} \cdot b_{n-1} \cdot s_{n-1}
c_4 \bigoplus c_3 = 0 \bigoplus 0 = 0
```

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TOC Method 1 for computing the overflow flag

- Adding two numbers with the same sign
- Overflow conditions for additions and subtractions
- Overflow condition for an addition
- Overflow conditions for a subtraction
- Overflow in signed computations

Adding two numbers with the same sign

- overflow can only happen when adding two numbers of the same sign results in a different sign $(P + P \rightarrow N, N + N \rightarrow P)$
- *n*-bit signed binary arithmetic A + B = S

$$A = (a_{n-1}, \dots, a_1, a_0)$$

$$B = (b_{n-1}, \dots, b_1, b_0)$$

$$S = (s_{n-1}, \dots, s_1, s_0)$$

- to detect overflow
 - only the sign bits are considered
 - msb (most significant bit) $a_{n-1}, b_{n-1}, s_{n-1}$
 - the other bits are ignored

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Overflow conditions for additions and subtractions

- with two operands (A and B) and one result (S), three sign bits $(a_{n-1}, b_{n-1}, s_{n-1})$ are considered $\rightarrow 2^3 = 8$ possible combinations
- only two cases result in overflow for an addition
 - 0 0 1 $(p+p \rightarrow n)$
 - 1 1 0 $(n+n\to p)$
- only two cases are considered as overflow for an subtraction
 - 0 1 1 $(p-n \to n)$
 - 100 $(n-p \to p)$

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Overflow condition for an addition

• Overflow in an addition (A + B)

	a_{n-1}	b_{n-1}	c_{n-1}	
	0	0	0	$\overline{p+p o p}$
OVER	0	0	1	p+p o n
	0	1	0	$p + n \rightarrow p$
	0	1	1	$p + n \rightarrow n$
	1	0	0	n+p o p
	1	0	1	$n+p \rightarrow n$
OVER	1	1	0	$n + n \rightarrow p$
	1	1	1	$n + n \rightarrow n$

- adding two positives should be positive
- adding two negatives should be negative

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Overflow conditions for a subtraction

• Overflow in a subtraction (A - B)

	a_{n-1}	b_{n-1}	c_{n-1}	
	0	0	0	p-p o p
	0	0	1	p-p ightarrow n
	0	1	0	p-n o p
OVER	0	1	1	p-n ightarrow n
OVER	1	0	0	n-p o p
	1	0	1	n-p ightarrow n
	1	1	0	n-n o p
	1	1	1	$n-n \rightarrow n$
-				

- subtracting a negative is the same as adding a positive
- subtracting a positive is the same as adding a negative

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Overflow in signed computations

- ALU might contain a small logic that <u>sets</u> the <u>overflow</u> flag to "1" if and only if any one of the above four <u>OV</u> conditions is met.
- in signed computations, <u>adding</u> two numbers of the <u>same sign</u> must produce a <u>result</u> of the <u>same sign</u>, otherwise overflow happened.

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TOC Method 2 for computing the overflow flag

- Carry into and carry out of the sign bit
- Overflow in 2's complement arithmetic
- Overflow flag = $c_n \bigoplus c_{n-1}$
- Examples of 4-bit signed additions
- c_n and c_{n-1} in a n-bit addition
- Overflow flag computation
- Examples of computing overflow flag
- Hexadecimal carry, octal carry, decimal carry
- No carry into the sign bit

Carry into and carry out of the sign bit

- When <u>adding</u> two *n*-bit binary values, consider
 - the carry coming into the most significant bit (msb)
 c_{n-1}: carry into the sign bit
 - the carry going out of the most significant bit (msb)
 c_n: carry out of the sign bit

this is the carry flag (CF) in the processor

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Overflow in 2's complement arithmetic

- overflow in 2's complement happens (OF=1) when
 - there is a carry into the sign bit $(c_{n-1} = 1)$ but no carry out of the sign bit $(c_n = 0)$
 - there is no carry into the sign bit $(c_{n-1} = 0)$ but a carry out of the sign bit $(c_n = 1)$

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Overflow flag = $c_n \bigoplus c_{n-1}$

- the overflow flag is the XOR $(c_n \bigoplus c_{n-1})$ of
 - of the carry coming into the sign bit (c_{n-1})
 - with the carry going out of the sign bit (c_n)
- overflow happens when the carry in (c_{n-1}) does <u>not</u> equal to the carry out (c_n)

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Examples of 4-bit signed additions (1)

4-bit 2's complement addition examples

```
0000
                                0100
0100 (+4) (pos sign 0)
                                0100 (+4) (pos sign 0)
+1000 (-8) (neg sign 1)
                                +0100 (+4) (pos sign 0)
01100 (-4) (neg sign 1)
                                01000 (-8) (neg sign 1)
C4 carry out 0 (1+0+0)
                               C4 carry out 0 (0+0+1)
C3 carry in 0 (0+1+0)
                                C3 carry in 1 (1+1+0)
O XOR O = NO OVERFI.OW
                                O XOR 1 = OVERFI.OW!
1100
                                1000
1100 (-4) (neg sign 1)
                               1100 (-4) (neg sign 1)
+0100 (+4) (pos sign 0)
                                +1000 (-8) (neg sign 1)
10000 (0) (pos sign 0)
                                10100 (+4) (pos sign 0)
C4 carry out 1 (1+0+1)
                                C4 carry out 1 (1+1+0)
C3 carry in 1 (1+1+0)
                                C3 carry in 0 (1+0+0)
1 XOR 1 = NO OVERFI.OW
                                1 XOR O = OVERFI.OW!
```

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Examples of 4-bit signed additions (2)

ullet same sign addition o possible overflow

+ +, -	, +	+ +, +	, -
+5	-5	+5	-5
+5	-5	+1	-1
-6(0	F) +6(0	F) +6	-6
0101	1011	0001	1111
0101	1011	0101	1011
0101	1011	0001	1111
01010	10110	00110	11010
C4 = 0	C4 = 1	C4 = 0	C4 = 1
C3 = 1	C3 = 0	C3 = 0	C3 = 1
0F = 1	OF = 1	OF = O	OF = O

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Examples of 4-bit signed additions (3)

ullet mixed sign addition o no overflow

+ -, +	+ -, -	- +, +	- +, -
+5	+5	-5	-5
-1	-6	+6	+1
+4	-1	+1	-4
1111	0000	1110	0011
0101	0101	1011	1011
1111	1010	0110	0001
10100	01111	10001	01100
C4 = 1	C4 = 0	C4 = 1	C4 = 0
C3 = 1	C3 = 0	C3 = 1	C3 = 0
OF = O	OF = O	OF = O	OF = O

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c_n and c_{n-1} in a *n*-bit addition

$(n-1)^{th}$ bit – MSB

- adding operations at the (n-1) bit position
- $\{c_n, s_{n-1}\} = a_{n-1} + b_{n-1} + c_{n-1}$

$$\begin{array}{c}
\text{msb} \\
a_{n-1} \\
b_{n-1} \\
c_{n-1}
\end{array}$$

● C_n: carry coming out of the msb

$(n-2)^{th}$ bit

- adding operations at the (n-2) bit position
- $\{c_{n-1}, s_{n-2}\} =$ $a_{n-2} + b_{n-2} + c_{n-2}$

msb

$$a_{n-2} \\ b_{n-2} \\ c_{n-2} \\ c_{n-1} \\ s_{n-2}$$

 \circ c_{n-1} : carry coming into the msb

Overflow flag computation

ADD (addition)	SUB (subtraction)
$0F = c_n \bigoplus c_{n-1}$	OF = $c_n \bigoplus c_{n-1}$
a 2's complement addition $A + B = A + B + 0$ ($c_0 = 0$)	the transformed addition $A-B=A+\overline{B}+1$ $(c_0=1)$
${c_n, s_{n-1}}$ = $a_{n-1} + b_{n-1} + c_{n-1}$	$\{c_n, s_{n-1}\}\$ = $a_{n-1} + \overline{b_{n-1}} + c_{n-1}$
${c_{n-1}, s_{n-2}}$ = $a_{n-2} + b_{n-2} + c_{n-2}$	$ \{c_{n-1}, s_{n-2}\} $ $= a_{n-2} + \overline{b_{n-2}} + c_{n-2} $

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Hexadecimal carry, octal carry, decimal carry

- Note that this XOR method only works with the binary carry that goes into the sign bit.
- not works with hexadecimal carry decimal carry, octal carry
 - the carry doesn't go into the sign bit
 - can't XOR that non-binary carry with the outgoing carry.

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No carry into the sign bit

Hexadecimal addition example
 (showing that XOR doesn't work for hex carry):
 ^{8Ah}
 +8Ah
 ====
 114h

- The hexadecimal carry of 1 resulting from A+A does not affect the sign bit.
- If you do the math in binary, you'll see
 that there is no carry into the sign bit;
 but, there is carry out of the sign bit.
 Therefore, the above example sets OVERFLOW on.
 (The example adds two negative numbers and gets a positive number.)

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Summary I

unsigned add	l/sub	signed	additio	n	signed	subtrac	tion	CF	OF
1101 (13) +1110 +(14)	ADD		(-3) +(-2)			(-3) -(+2)			
11011 (11)	(+16)	1 11011	(-5)		11011	(-5)		1	0
0011 (3) -1110 -(14)	SUB	0011	(+3) +(+2)		0011 -1110	(+3) -(-2)	SUB		
10101 (5)	(-16)	00101	(+5)		00101	(+5)		 1 	0
0011 (3) +0010 +(2)	ADD	0011	(+3) +(+2)	ADD	0011 -1110	(+3) -(-2)			
00101 (5)	(+ 0)	 00101	(+5)		00101	(+5)		 0 	0
1101 (13) -0010 -(2)	SUB	1101 +1110	(-3) +(-2)		1101 -0010	(-3) -(+2)	SUB		
11011 (11)	(-16)	11011	(-5)		11011	(-5)		0	0

Summary II

		/								
unsign	iea aaa,	sub	signed	additio	n 	signea	subtrac	tion	CF	10
1011	(11)		1011	(-5)		1011	(-5)		l	
+1100	+(12)	ADD	+1100	+(-4)	ADD	-0100	-(+4)			
10111	(7)	(116)		(+7)		10111	(+7)		 1	1
10111	(1)	(+16)	10111	(+1)		10111	(+7)		 T	1
0101	(5)		0101	(+5)		0101	(+5)		i	
-1100	-(12)	SUB	+0100	+(+4)		-1100	-(-4)	SUB	I	
									<u> </u>	
11001	(9)	(-16)	01001	(-7)		01001	(-7)		1 1	1
0101	(5)		0101	(+5)		0101	(+5)		İ	
+0100	+(4)	ADD	+0100	+(+4)	ADD	-1100	-(-4)		I	
			l						l	
01001	(9)	(+ 0)	01001 	(-7)		01001	(-7)		0 	1
1011	(11)		1011	(-5)		1011	(-5)		İ	
-0100	-(4)	SUB	+1100	+(-4)		-0100	-(+4)	SUB		
		(0)		(.7)			(.7)			
00111	(7)	(0)	10111	(+7)		10111	(+7)		0	1

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Cases for setting the overflow flag (1) CF=1, OF=1

singed integer overflow (OF=1 means incorrect S)

* unsigned addition	* signed addition	signed subtraction
1011 (11) +1100 +(12) ADD	1000 1011 (-5) +1100 +(-4) ADD	1011 (-5) -0100 -(+4)
10111 (7) (+16)	10111 (+7)	10111 (+7)
0F=1	n + n -> p (OF=1)	n - p -> p (OF=1)
OF meaningless	-> incorrect S	-> incorrect S
S = 0111	S = 0111	S = 0111
* think hand addition	* 0F <- C4 XOR C3 = 1 XO	

 $\boldsymbol{*}$ CF=1, S=0111, OF=1 for all three interpretations

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Cases for setting the overflow flag (2) CF=1, OF=1

singed integer overflow (OF=1 means incorrect S)

* unsigned subtraction	Ι	signed addition	* signed subtraction
0101 (5) -1100 -(12) SUB 11001 (9) (-16)	 	0100 0101 (+5) +0100 +(+4) 01001 (-7)	0101 (+5) -1100 -(-4) SUB 01001 (-7)
0F=1	 	p + p -> n (OF=1)	p - n -> n (OF=1)
OF meaningless S = 1001	 	-> incorrect S S = 1001	-> incorrect S S = 1001
* think hand subtraction	* *	OF <- C4 XOR C3 = 0 XOR of signed addition	

^{*} CF=1, S=1001, OF=1 for all three interpretations

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Cases for setting the overflow flag (3) CF=0, OF=1

singed integer overflow (OF=1 means incorrect S)

* unsigned addition	* signed addition	signed subtraction
0101 (5) +0100 +(4) ADD	0100 0101 (+5) +0100 +(+4) ADD	0101 (+5) -1100 -(-4)
01001 (9) (+ 0)	01001 (-7)	01001 (-7)
0F=1	p + p -> n (0F=1)	p - n -> n (0F=1)
OF meaningless S = 1001	-> incorrect S S = 1001	-> incorrect S S = 1001
* think hand addition	* OF <- C4 XOR C3 = O XOR of signed addition	= =

^{*} CF=0, S=1001, OF=1 for all three interpretations

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Cases for setting the overflow flag (4) CF=0, OF=1

• singed integer overflow (OF=1 means incorrect S)

unsigned subtraction	1	signed addition	* signed subtraction
1011 (11) -0100 -(4) SUB	 	1000 1011 (-5) +1100 +(-4)	1011 (-5) -0100 -(+4) SUB
00111 (7) (0)	 	10111 (+7)	10111 (+7)
0F=1	i	$n + n \rightarrow p (0F=1)$	n - p -> p (0F=1)
OF meaningless S = 0111	i I	-> incorrect S S = 0111	-> incorrect S S = 0111
think hand subtraction	; 	* OF <- C4 XOR C3 = 1 X	

 \ast CF=0, S=0111, OF=1 for all three interpretations

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Cases for clearing the overflow flag (1) CF=1, OF=0

• no signed integer overflow (CF=0 means correct S)

unsigned addition	* signed addition	signed subtraction
1101 (13) +1110 +(14) ADD	1100 1101 (-3) +1110 +(-2) ADD	1101 (-3) -0010 -(+2)
11011 (11) (+16)	 11011 (-5)	11011 (-5)
0F=0	n + n -> n (0F=0)	n - p -> n (OF=0)
OF meaningless	-> correct S	-> correct S
S = 0000	S = 0000	S = 0000
think hand addition	* OF <- C4 XOR C3 = 1 XC of signed additi	

 $\boldsymbol{*}$ CF=1, S=1011, OF=0 for all three interpretations

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Cases for clearing the overflow flag (2) CF=1, OF=0

• no signed integer overflow (CF=0 means correct S)

unsigned subtraction				signed		
		0010				
0011 (3)	1	0011 (+3)		0011	(+3)	
-1110 -(14) SUB		+0010 +(+2)		-1110	-(-2)	SUE
	1					
10101 (5) (-16)	 	00101 (+5)		00101	(+5)	
CF=1	i i	p + p -> p (OF=0)		p - n	-> p (0)F=0)
OF meaningless	İ	-> correct S		-> cor:	rect S	
S = 0101	I	S = 0101		S =	0101	
think hand	*	OF <- C4 XOR C3 =	O XOR	0 = 0		
subtraction	1	of signed ad	dition			

* CF=1, S=0101, OF=0 for all three interpretations

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Cases for clearing the overflow flag (3) CF=0, OF=0

• no signed integer overflow (CF=0 means correct S)

unsigned addition	* signed addition	signed subtractio
	0010	
0011 (3)	0011 (+3)	0011 (+3)
+0010 +(2) ADD	+0010 +(+2) ADD	-1110 -(-2)
00101 (5) (+ 0)	00101 (+5)	00101 (+5)
0F=0	p + p -> p (OF=0)	p - n -> p (OF=0)
OF meaningless	-> correct S	-> correct S
S = 0101	S = 0101	S = 0101
think hand	* OF <- C4 XOR C3 = 0 X	DR 0 = 0
addition	of signed addit:	ion

 \ast CF=0, S=0101, OF=0 for all three interpretations

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Cases for clearing the overflow flag (4) CF=0, OF=0

• no signed integer overflow (CF=0 means correct S)

unsigned addition	<pre>* signed addition</pre>	signed subtraction
1101 (13) -0010 -(2) SUB	1100 1101 (-3) +1110 +(-2)	1101 (-3) -0010 -(+2) SUB
11011 (11) (-16)	 11011 (-5)	11011 (-5)
0F=0	n + n -> n (OF=0)	n - p -> n (0F=0)
OF meaningless S = 1011	-> correct S S = 1011	-> correct S S = 1011
think hand subtraction	* OF <- C4 XOR C3 = 1 X	

 $\boldsymbol{*}$ CF=0, S=1011, OF=0 for all three interpretations

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