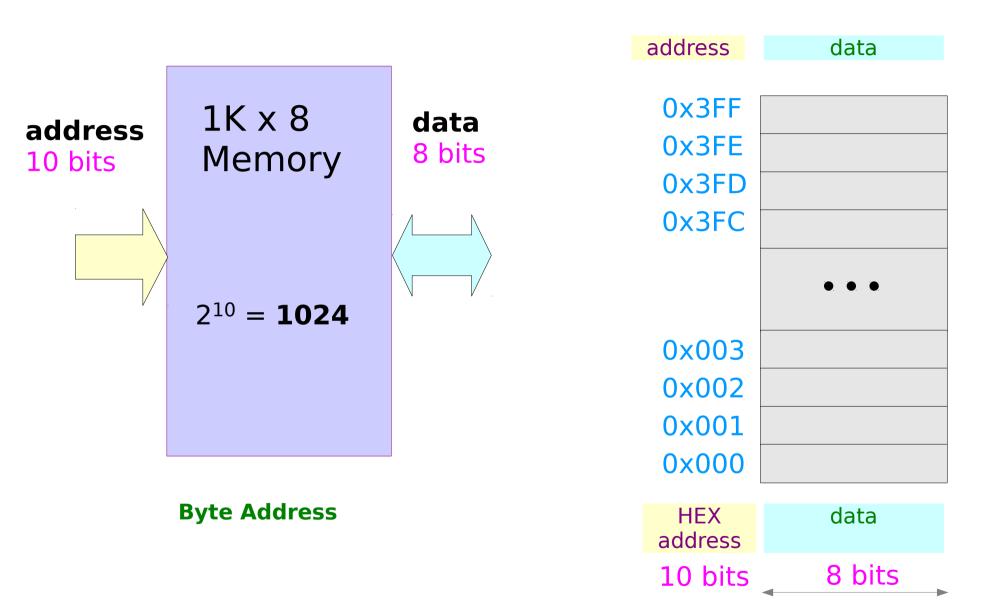
Pointers (1A)

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Byte Address and Data in a Memory



Variables

int a; a can hold an <u>integer</u>

address data

a = 100;

a holds the <u>integer</u> 100

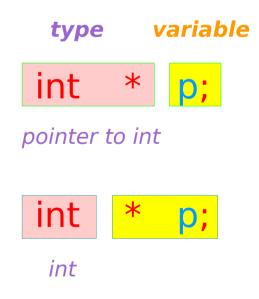
address data
&a = 100

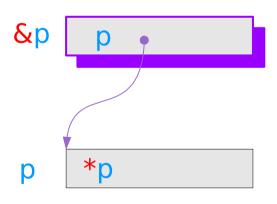
Pointer Variables

```
int * p;
p holds an <u>address</u>
```

p can hold the <u>address</u> of an int data

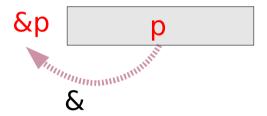
*p can hold an <u>integer</u>



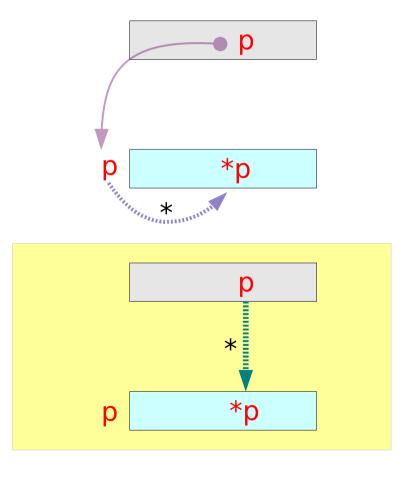


Dereferencing

The address of a variable : Address of operator &



The content of a pointed location : Dereferencing operator *



Variables and their addresses

	address	data
int a;	&a	a
int *p;	&p	p

Assignment of a value

int a;

int b;

address data

&a = 111

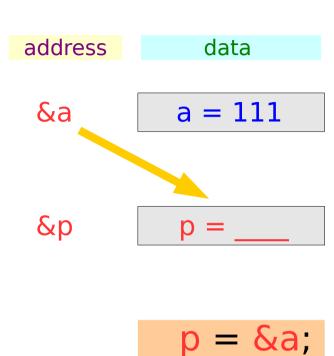
&b =

b = a;

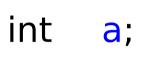
Assignment of an address

int a;

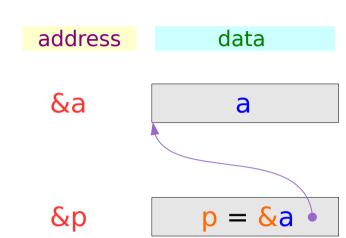
int *p;



Variables with initializations



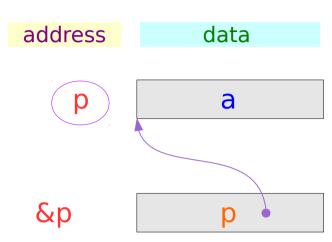
int
$$*p = \&a$$



Pointed addresses: p

int a;

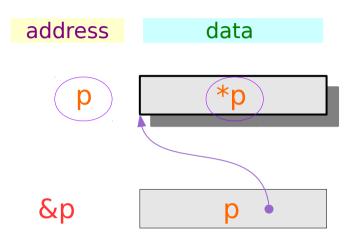
int *p = &a;



$$p \equiv \&a$$

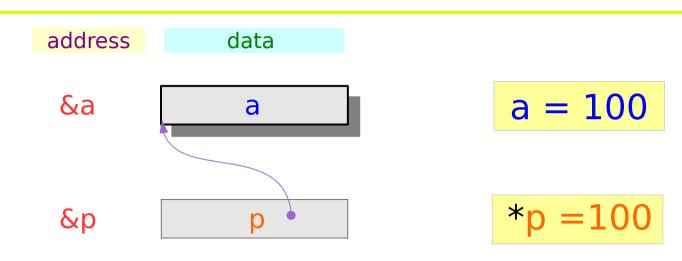
Dereferenced Variable: *p

int
$$*p = \&a$$



$$p \equiv &a$$
* $p \equiv a$

Two way to access: a and *p



- 1) Read/Write a
- 2) Read/Write *p

- 1. Pass by Reference
- 2. Arrays

Pass by Reference

Variable Scopes

```
int func1 (int a, int b)
{
   int i, int j;
   ...
   ...
}
```

```
int main ()
{
    int x, int y;
    ...
    ...
    func1 ( 10, 20 );
    ...
    ...
}
```

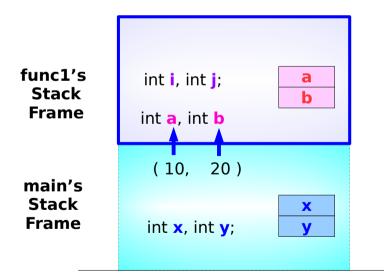
```
i and j's
variable scope

cannot access
each other

x and y's
variable scope
```

Only **top** stack frame is <u>active</u> and its variable can be accessed

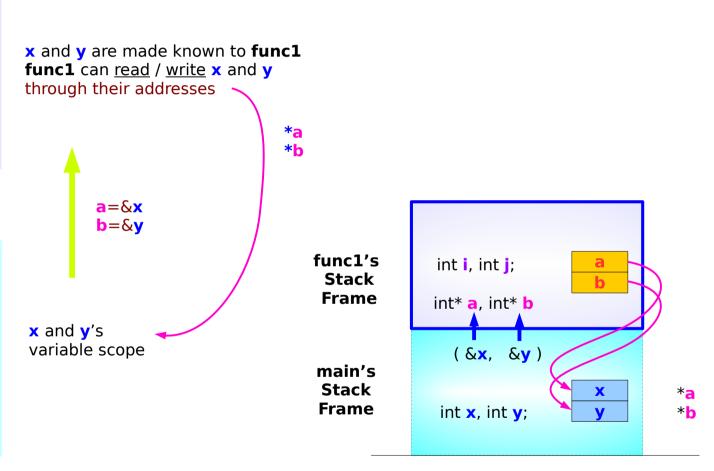
Communications are performed only through the **parameter** variables



Pass by Reference

```
int func1 (int* a, int* b)
{
    int i, int j;
    ...
    ...
    ...
}
```

```
int main ()
{
    int x, int y;
    ...
    ...
    func1 ( &x, &y );
    ...
    ...
}
```



Swapping integers

function call

function prototype

Pass by integer reference

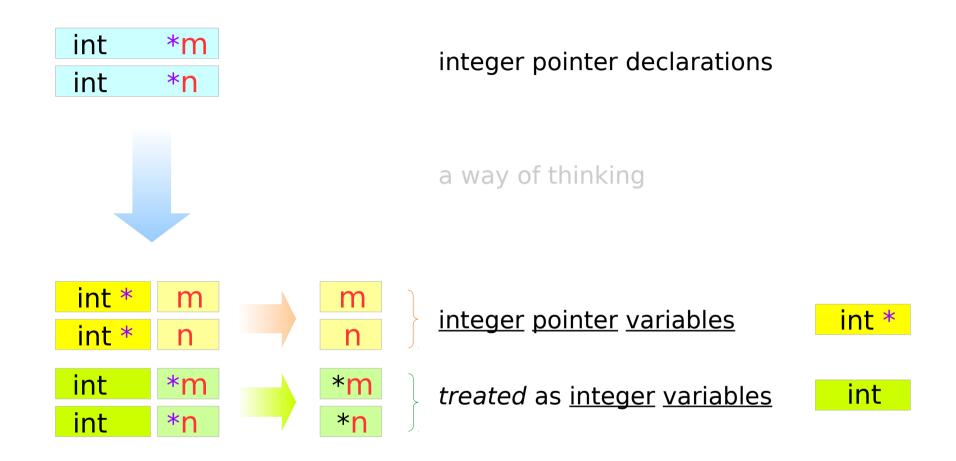
```
void swap(int *p, int *q) {
   int tmp;

tmp = *p;
   *p = *q;
   *q = tmp;
}
```

```
int * p
int *q
int * p
int * q
int tmp
```

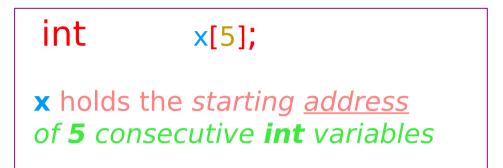
```
int a, b;
...
swap( &a, &b );
```

Integer and Integer Pointer Types



Arrays

Accessing array elements – using an address



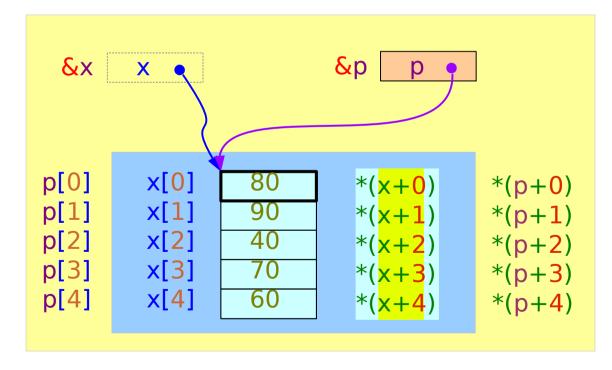
5 int variables

index data

0 x[0]
1 x[1]
2 x[2]
3 x[3]
4 x[4]

address data Ţ cannot change *x address x (constant) *(x+1) x + 1*(x+2)x + 2*(x+3)x + 3*(x+4)x + 4

Accessing an Array with a Pointer Variable



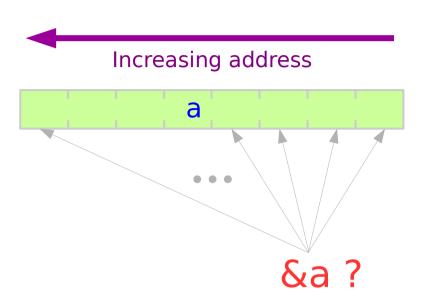
x is a constant symbol cannot be changed

p is a variable can point to other addresses

Byte Address Little Endian Big Endian

Byte Address

long a;



Numbers in Positional Notation

long
$$a = 0x1020304050607080$$
;

8 (bytes)

$$a_7$$
 a_6 a_5 a_4 a_3 a_2 a_1 a_0

Most Significant Byte
$$a_7 = 0 \times 10 + 16^7$$
 the highest weight

$$a_6 = 0 \times 20 \quad \cdots \quad 16^6$$

$$a_5 = 0 \times 30 \quad \cdots \quad 16^5$$

$$a_4 = 0 \times 40 \quad \cdots \quad 16^4$$

$$a_3 = 0 \times 50 \quad \cdots \quad 16^3$$

$$a_2 = 0 \times 60 \quad \cdots \quad 16^2$$

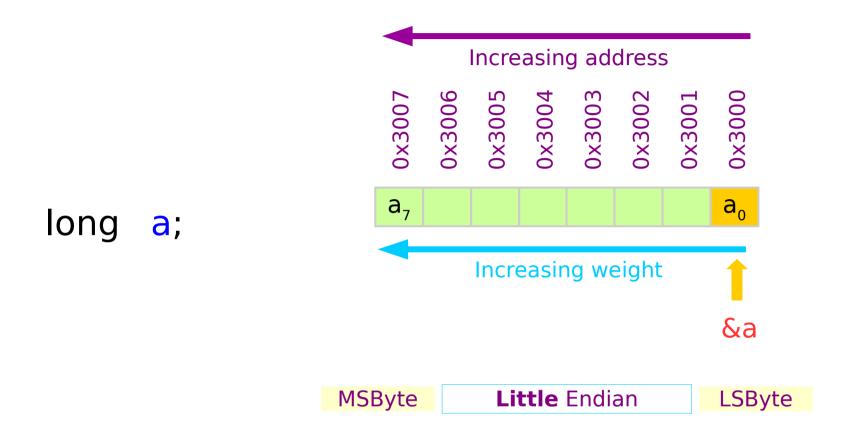
$$a_1 = 0 \times 70 \quad \dots \quad 16^1$$

Least Significant Byte
$$a_0 = 0 \times 80 \dots 16^0$$
 the lowest weight

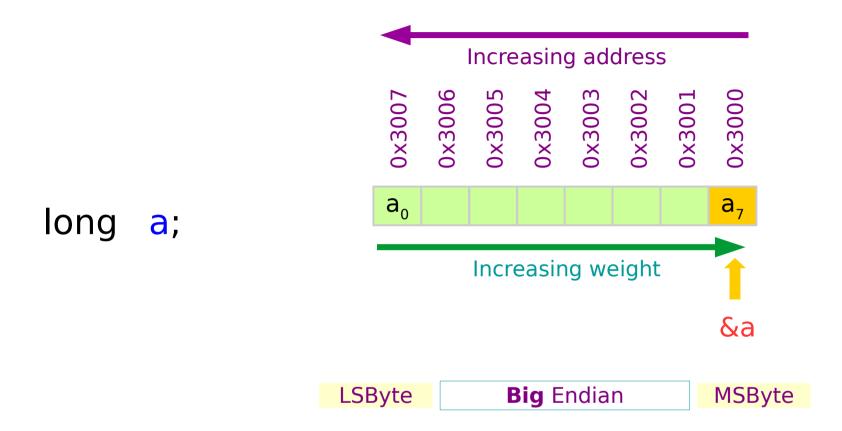
Little / Big Endian

Increasing address long **a**; a **MSByte** LSByte Little Endian $a_7 \ a_6 \ a_5 \ a_4 \ a_3 \ a_2 \ a_1 \ a_0$ Big Endian LSByte **MSByte** $a_0 \ a_1 \ a_2 \ a_3 \ a_4 \ a_5 \ a_6 \ a_7$

Little Endian Byte Address Example



Big Endian Byte Address Example



Representations of Endianness

downward, increasing address

&a &a 0x3000 a₇ a_o 0x3000 0x3001 0x3001 0x3002 0x3002 0x3003 0x3003 0x3004 0x3004 0x3005 0x3005

0x3006

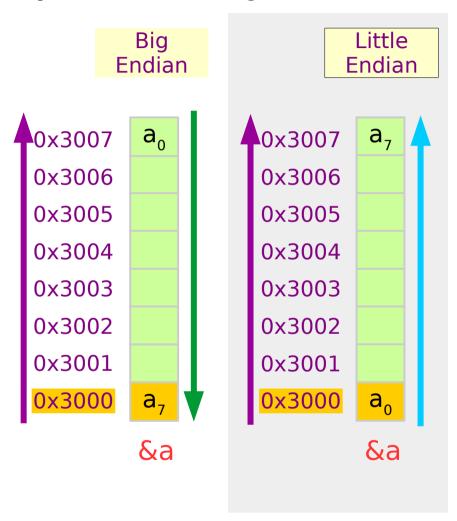
0x3007

 a_7

Little

Endian

upward, increasing address



https://stackoverflow.com/questions/15620673/which-bit-is-the-address-of-an-integer

0x3006

0x3007

 a_0

Big

Endian

Little / Big Endian Processors

Processor	Endianness
Motorola 68000	Big Endian
PowerPC (PPC)	Big Endian
Sun Sparc	Big Endian
IBM S/390	Big Endian
<pre>Intel x86 (32 bit)</pre>	Little Endian
<pre>Intel x86_64 (64 bit)</pre>	Little Endian
Dec VAX	Little Endian
Alpha	(Big/Little) Endian
ARM	(Big/Little) Endian
IA-64 (64 bit)	(Big/Little) Endian
MIPS	(Big/Little) Endian

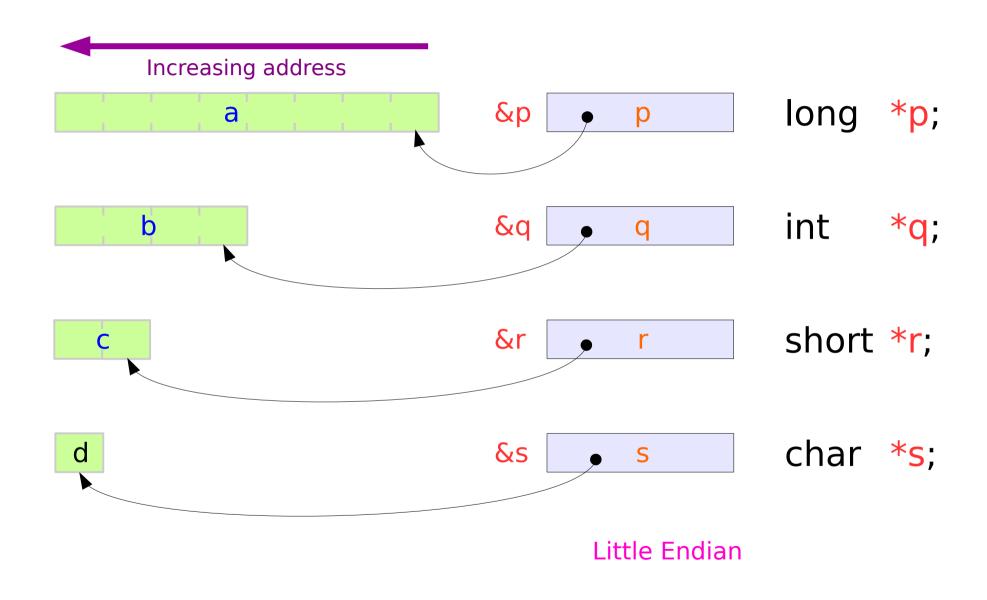
http://www.yolinux.com/TUTORIALS/Endian-Byte-Order.html

Pointer Types

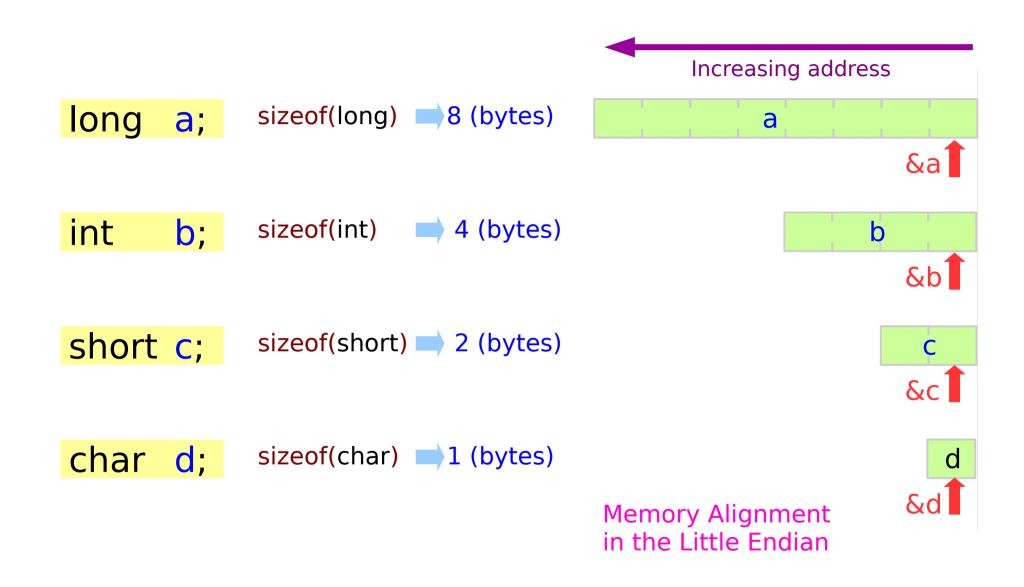
Integer Type Variables and Their Addresses

Increasing address long **a**; a &a int b; b short c; &c char d; d Little Endian &d

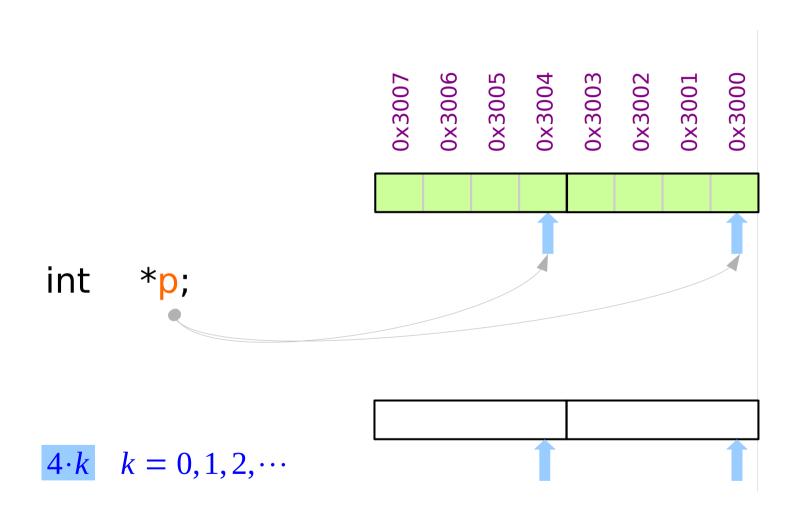
Points to the LSByte



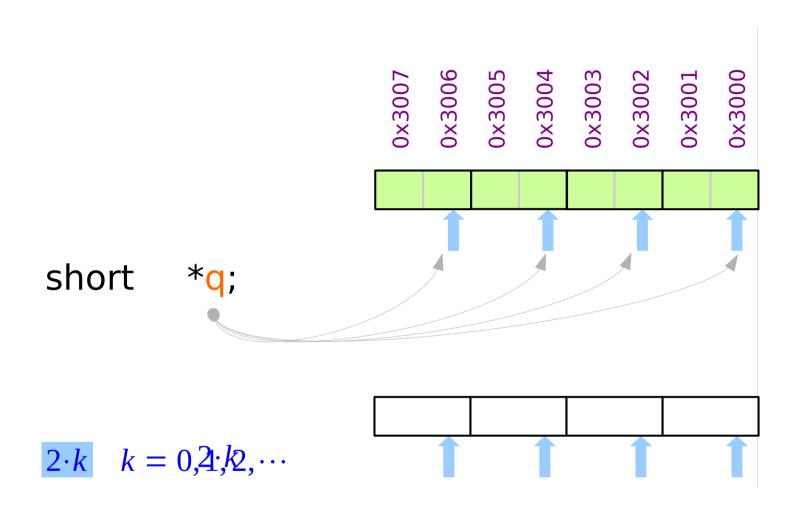
Sizes of Integer Types



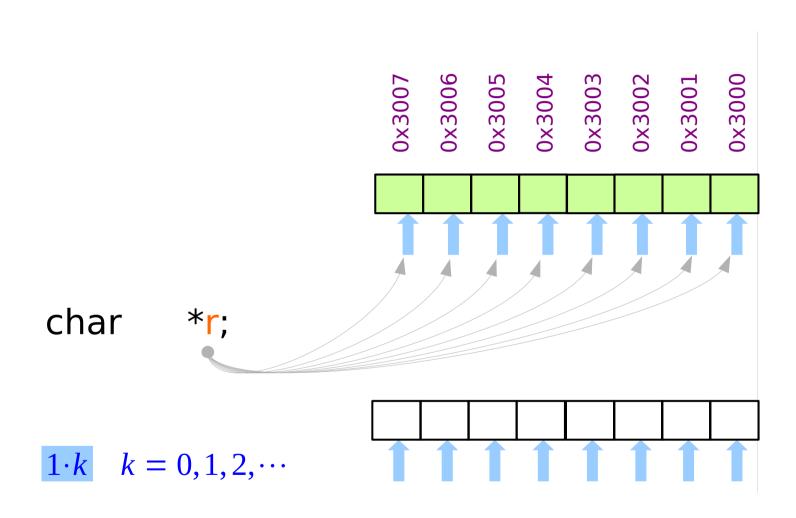
Possible addresses for integer values



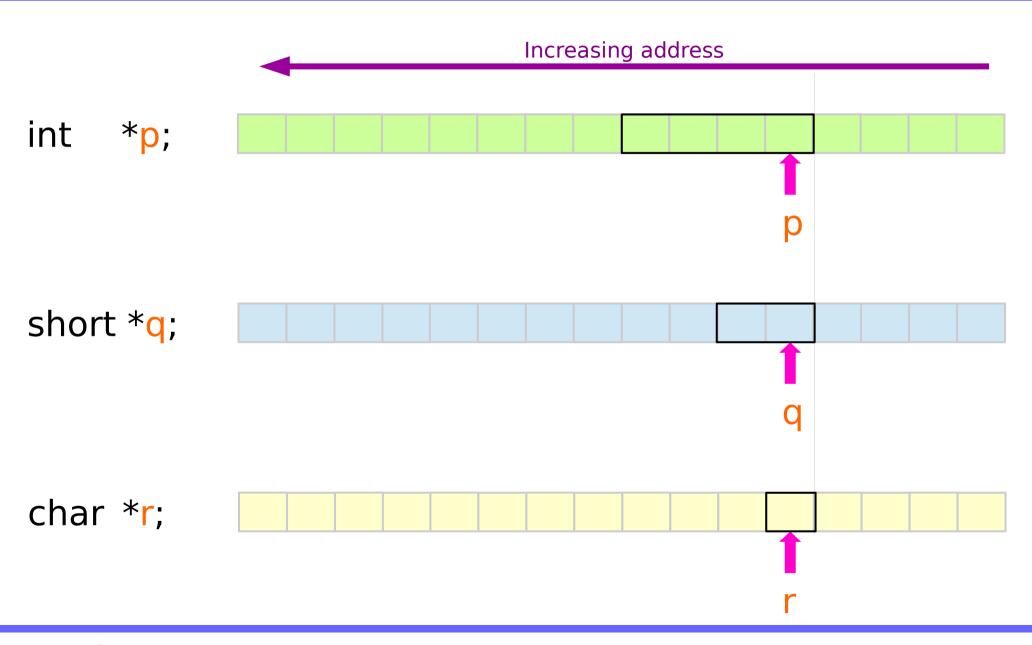
Possible addresses for short values



Possible address for char values



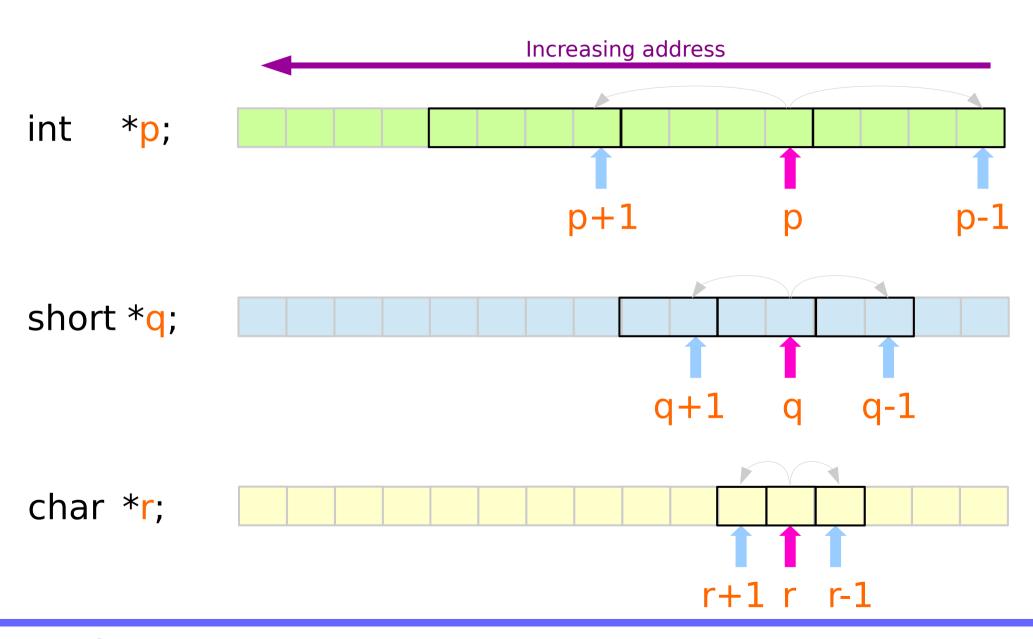
Pointed Addresses



Series: 2. Pointers

Young Won Lim 2/10/18

Incrementing / decrementing pointers



Memory Alignment (1) - allocation of variables

Memory Alignment: the data address is a <u>multiple</u> of the data size.

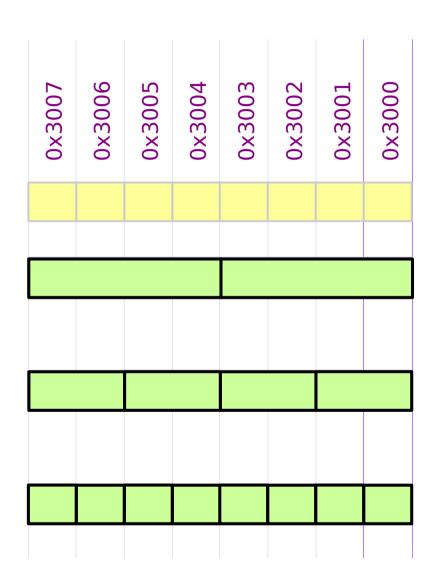
enforced by compilers

efficient memory access

int a;

short b;

char c;



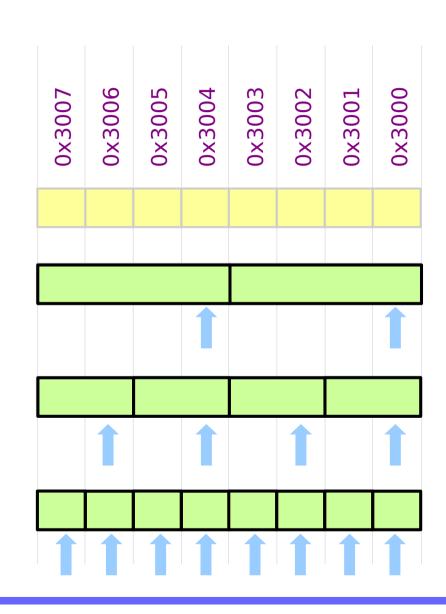
Memory Alignment (2) – integer multiple addresses

$$k = 0, 1, 2, \cdots$$

integer addresses = $4 \cdot k$

short addresses = $2 \cdot k$

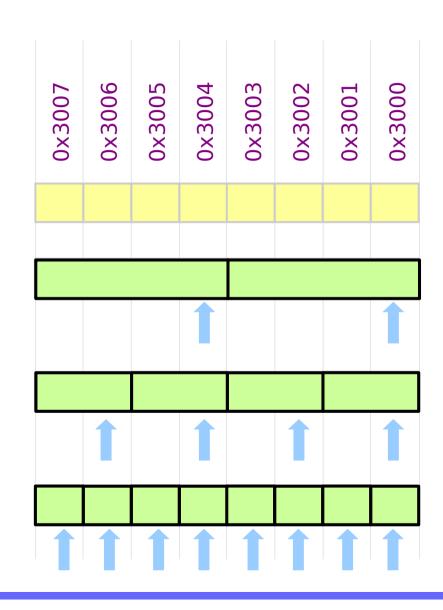
character addresses = $1 \cdot k$



Memory Alignment (3) – pointable addresses

int *p;

short *q;



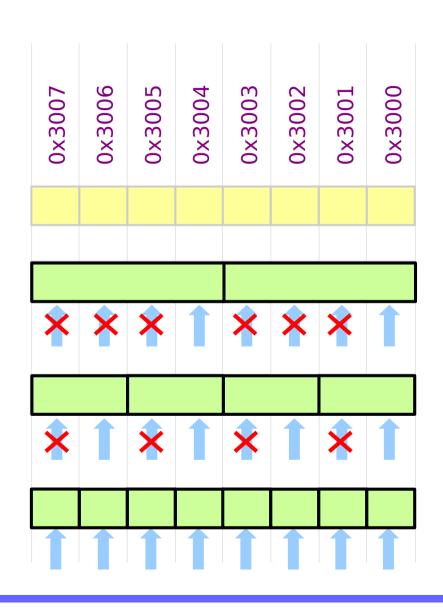
Memory Alignment (4) – non-pointed addresses



 $4 \cdot k + 1, 2, 3$

short *q;

 $2 \cdot k + 1$



Memory Alignment (5) - breaking alignment

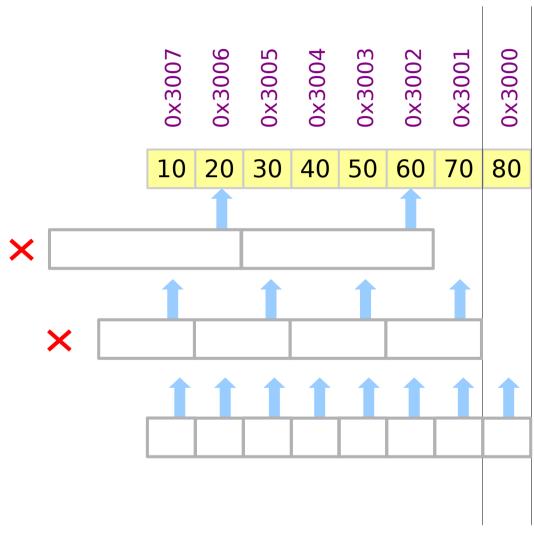
Memory access is still possible but it takes **longer** to access

(Low Efficiency)

This can happen by using inadvertent pointer type casting

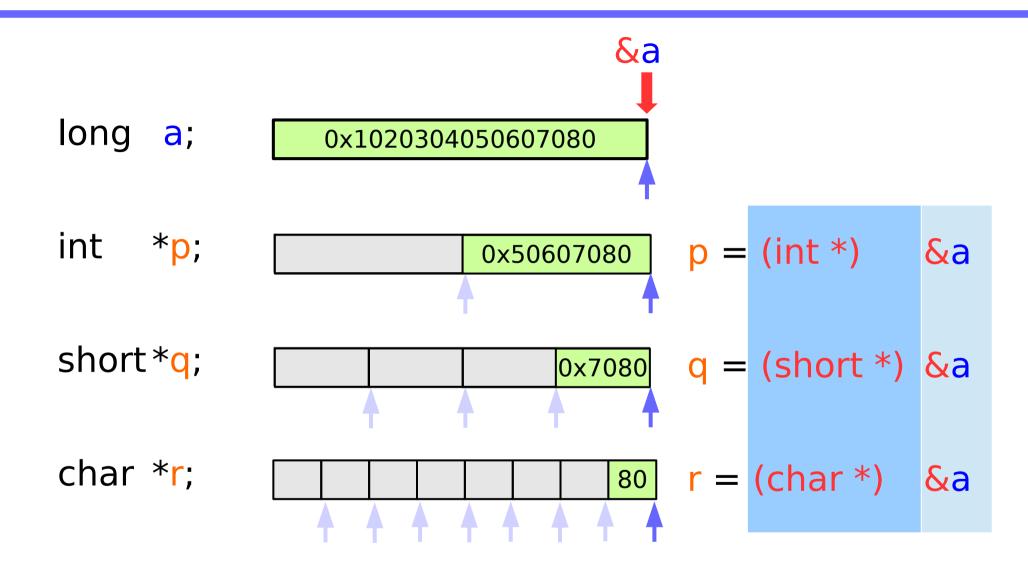
int *p;

short*q;



Pointer Type Cast

Re-interpretation of memory data - case I



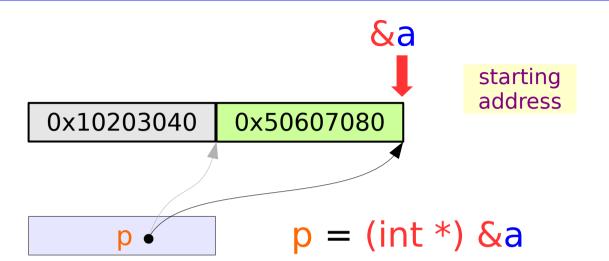
47

Pointer Type Cast

long a;

int *p;

short*q;

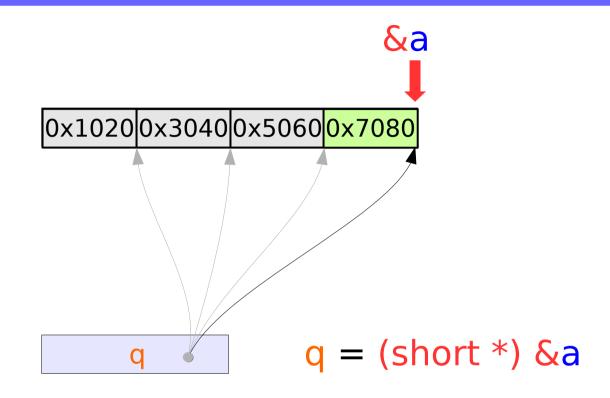


Integer Pointer Types

long a;

int *p;

short*q;

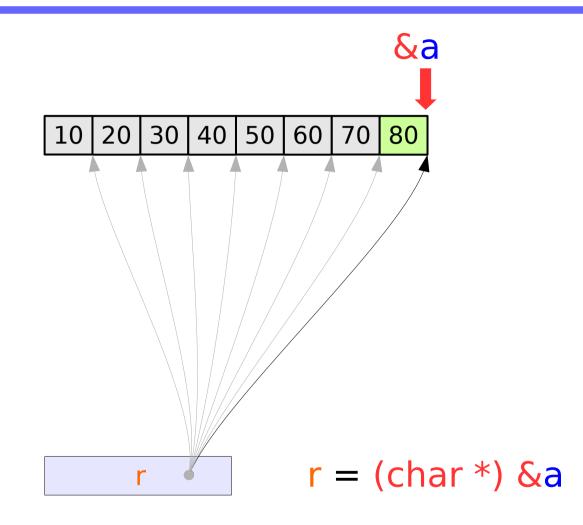


Integer Pointer Types

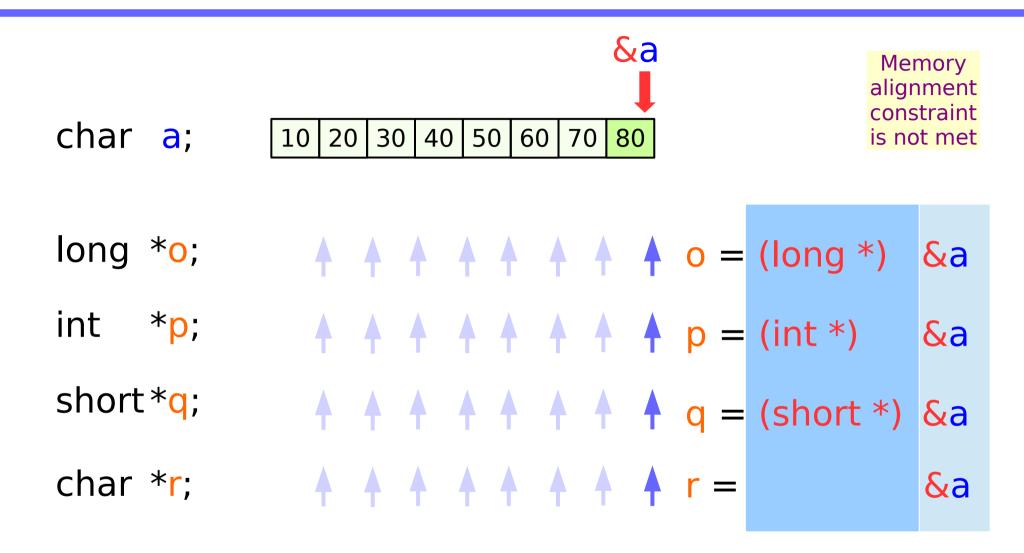
```
long a;
```

int *p;

short*q;



Re-interpretation of memory data - case II



In this case, the memory alignment constraint can be broken

const pointers

const type, const pointer type (1)

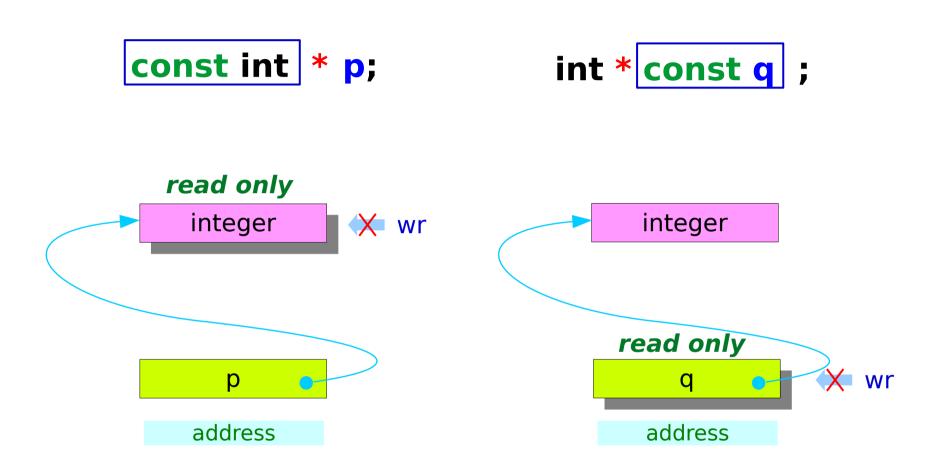
```
const int * p;
int * const q ;
const int * const r ;
```

```
read only integer <u>value</u>
```

read only integer pointer

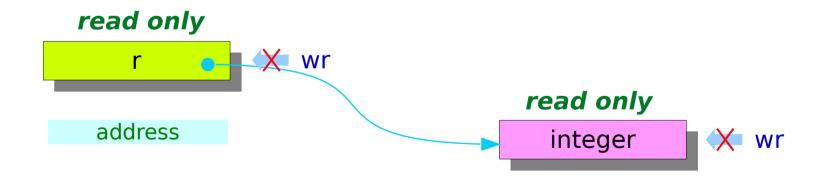
read only integer <u>value</u> read only integer <u>pointer</u>

const type, const pointer type (2)

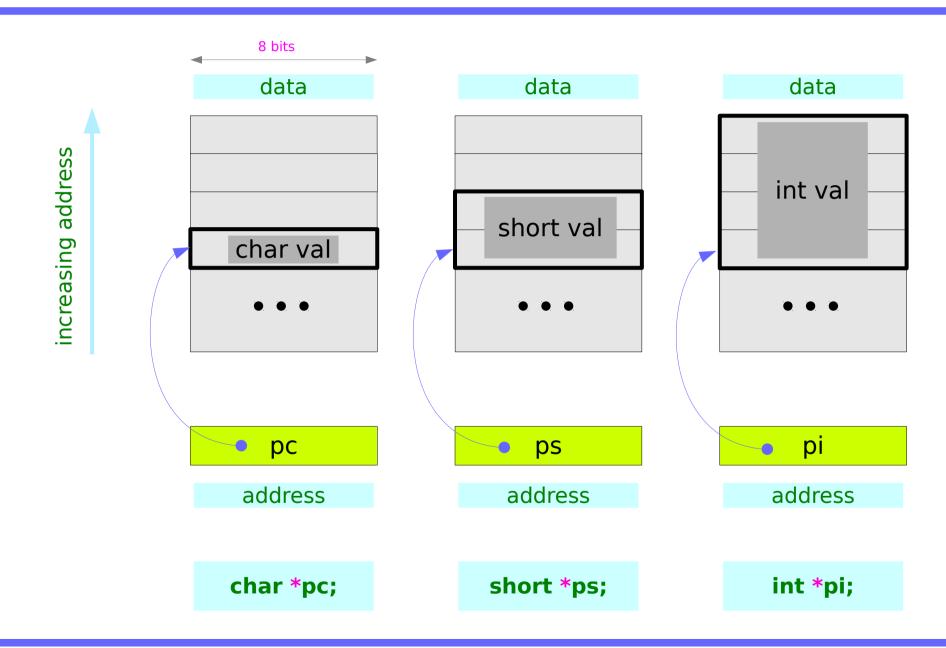


const type, const pointer type (3)

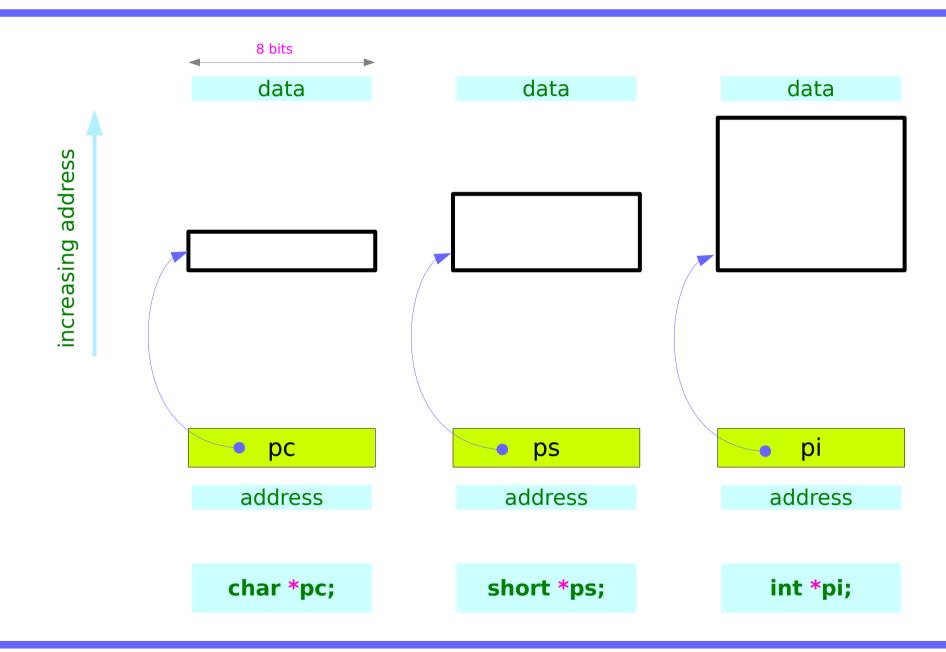
```
const int * const r ;
```



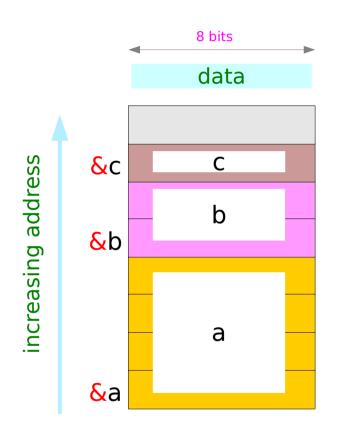
Pointer Types and Associated Data

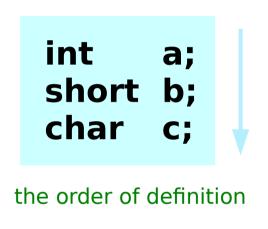


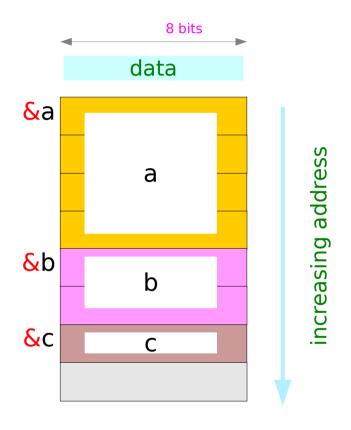
Pointer Types



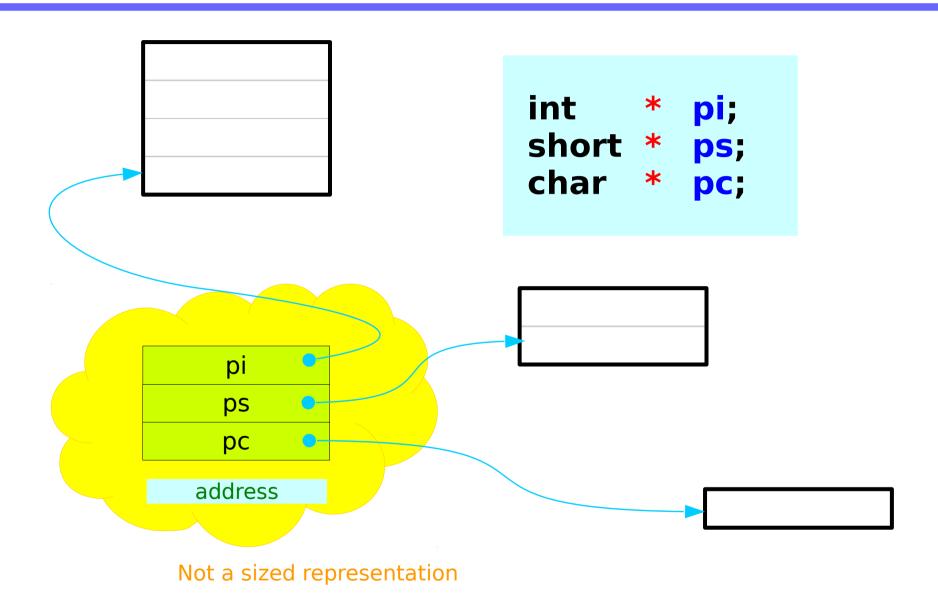
Little Endian Example



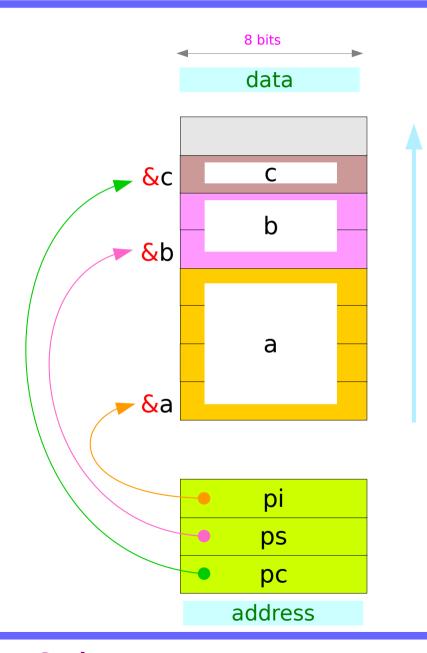




int *, short *, char * type variables



Pointer Variable Assignment

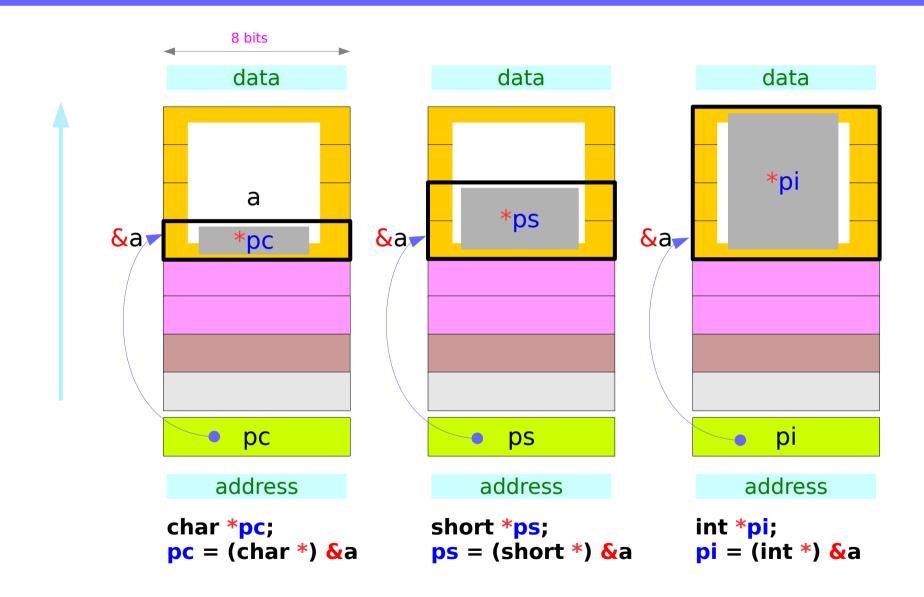


```
char * pc;
short * ps;
int * pi;

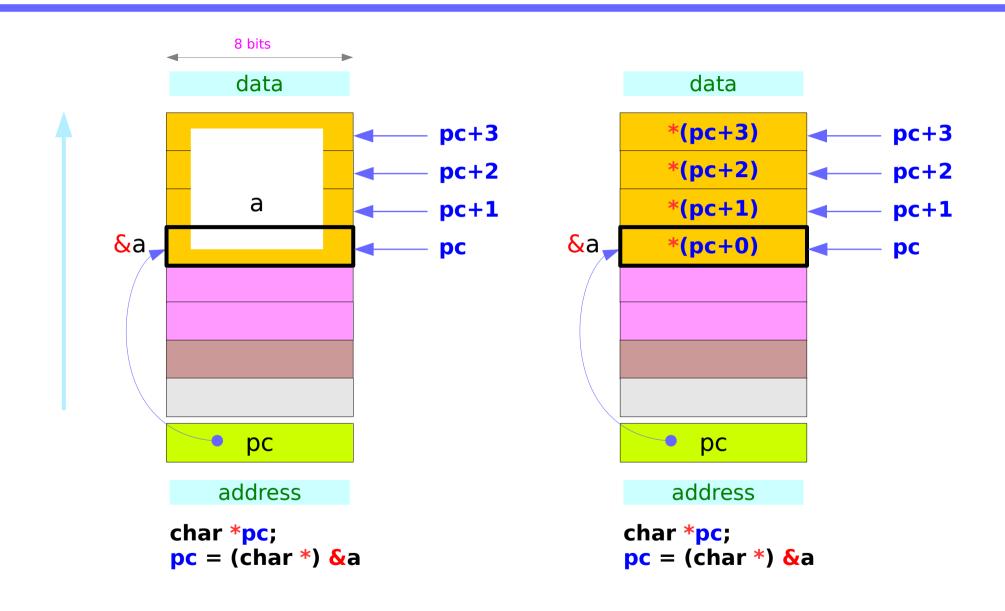
int a;
short b;
char c;
```

```
pi = &a;
ps = &b;
pc = &c;
```

Pointer Type Casting



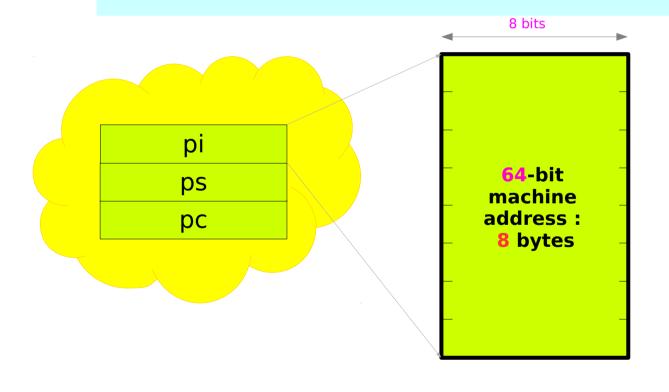
Accessing bytes of a variable

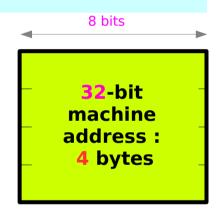


32-bit and 64-bit Address

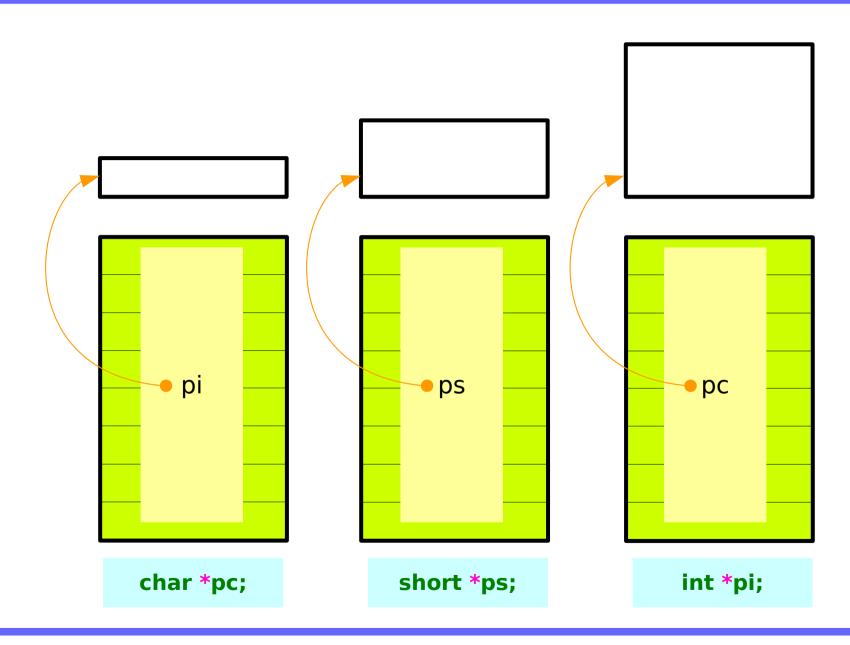


64-bit machine : address : 8 bytes

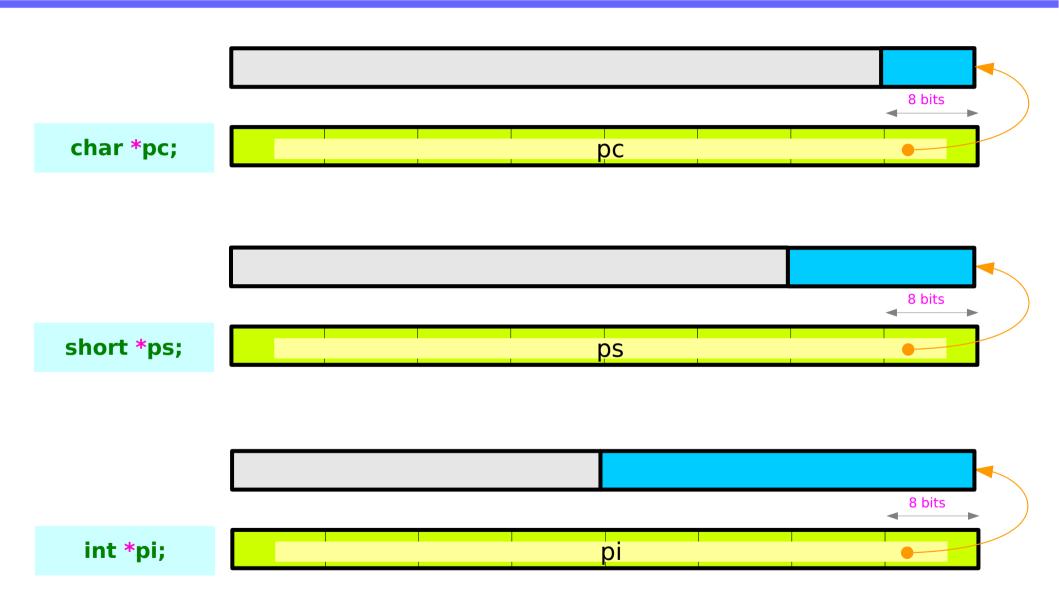




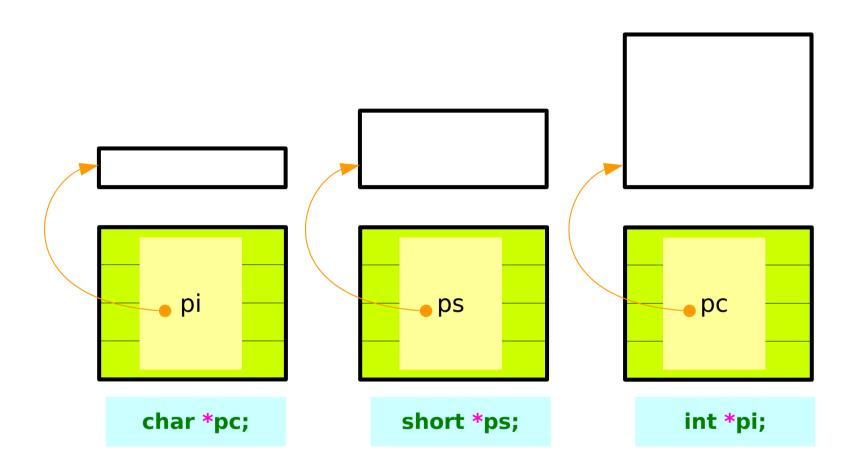
64-bit machine: 8-byte address



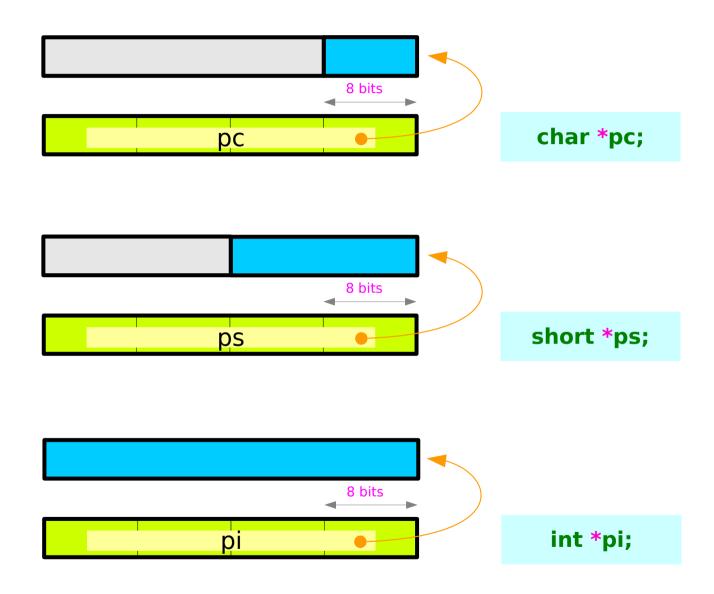
64-bit machine: 8-byte address & data buses



32-bit machine : 4-byte address



64-bit machine: 8-byte address and data buses



References

- [1] Essential C, Nick Parlante
- [2] Efficient C Programming, Mark A. Weiss
- [3] C A Reference Manual, Samuel P. Harbison & Guy L. Steele Jr.
- [4] C Language Express, I. K. Chun