

Applications of Array Pointers (1A)

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Assumption

assume that

value(c) returns the hexadecimal number that is obtained by `printf("%p", c)`, when the variable `c` contains an address as its value

type(c) can be determined by the warning message of `printf("%d", c)`, when the variable `c` contains an address as its value

```
#include <stdio.h>
int main(void) {
    int c[3];
    printf ("c= %p \n", &c);
}
```

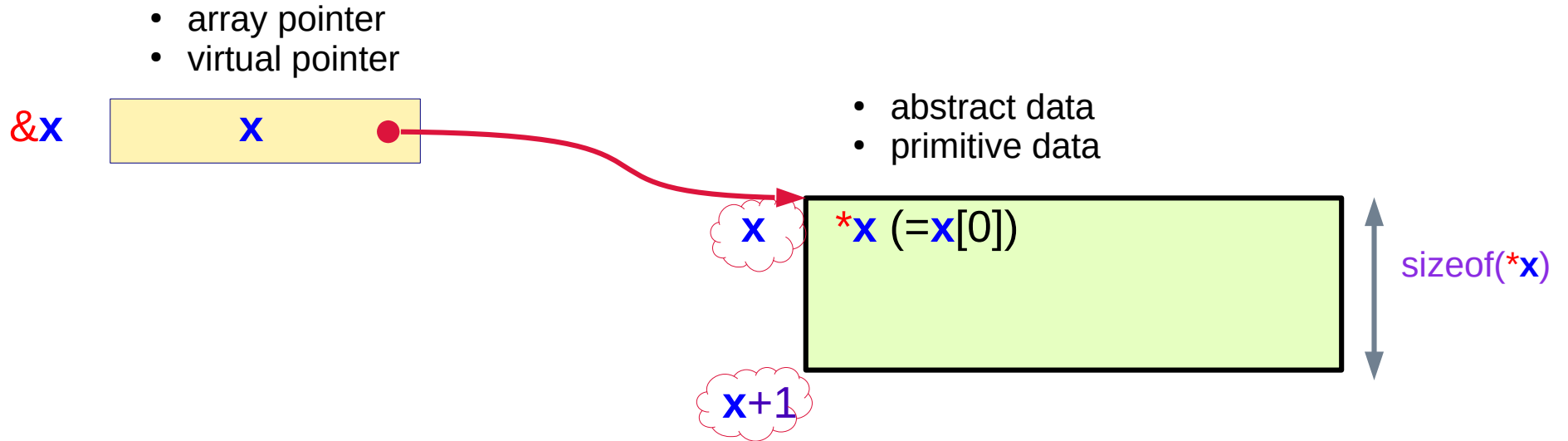
`c= 0x7fffd923487c`

```
#include <stdio.h>
int main(void) {
    int c[3];
    printf ("c= %d \n", &c);
}
```

t.c: In function 'main':
t.c:5:16: warning: format '%d' expects argument of type 'int',
but argument 2 has type 'int (*)[3]' [-Wformat=]
printf ("c= %d \n", &c);

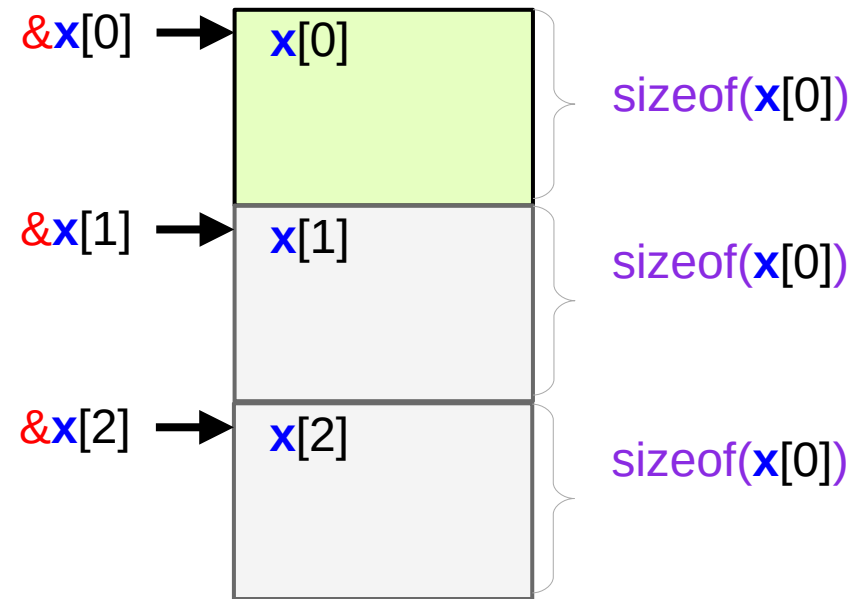
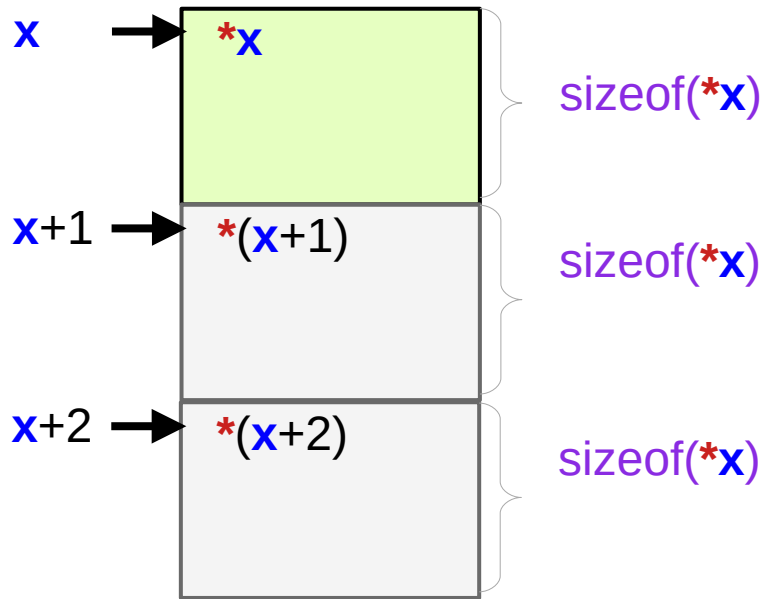
Byte Address

Pointer x and $x+1$ relationship



$$\text{value}(x+1) = \text{value}(x) + \text{sizeof}(*x)$$

Byte addresses in an array



$$\begin{aligned} \text{value}(x+1) &= \text{value}(x) + 1 * \text{sizeof}(*x) \\ \text{value}(x+2) &= \text{value}(x) + 2 * \text{sizeof}(*x) \end{aligned}$$

byte address byte address byte size

$$\begin{aligned} \text{value}(\&x[1]) &= \text{value}(x) + 1 * \text{sizeof}(x[0]) \\ \text{value}(\&x[2]) &= \text{value}(x) + 2 * \text{sizeof}(x[0]) \end{aligned}$$

byte address byte address byte size

Four cases of array pointers

Virtual pointers vs. real pointers

Case 1

virtual pointer **c**

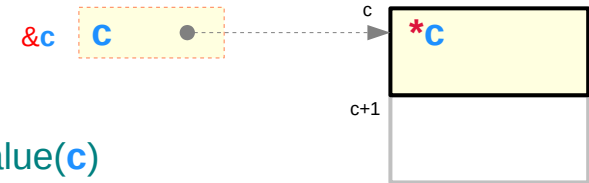
primitive data ***c**



Case 2

virtual pointer **c**

abstract data ***c**



subarray partitioning

$$\star \text{sizeof}(\mathbf{c}) = \text{sizeof}(\star \mathbf{c}) * \mathbf{N}$$

$$\text{value}(\mathbf{c}+1) = \text{value}(\mathbf{c}) + \text{sizeof}(\star \mathbf{c})$$

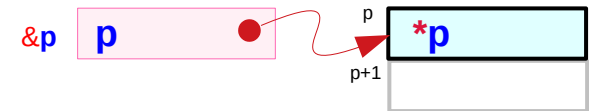
address replication

$$\star \text{value}(\&\mathbf{c}) = \text{value}(\mathbf{c})$$

Case 3

real pointer **p**

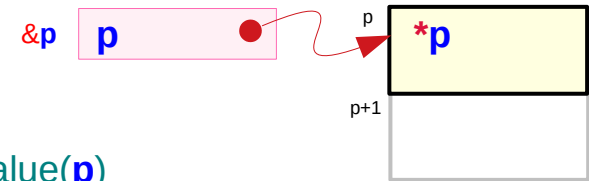
primitive data ***p**



Case 4

real pointer **p**

abstract data ***p**



$$\text{sizeof}(\mathbf{p}) = \text{pointer size (4/8 bytes)}$$

$$\text{value}(\&\mathbf{p}) \neq \text{value}(\mathbf{p})$$

$$\text{value}(\mathbf{p}+1) = \text{value}(\mathbf{p}) + \text{sizeof}(\star \mathbf{p})$$

Primitive data vs. abstract data

Case 1

virtual pointer **c**

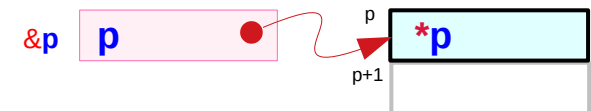
primitive data ***c**



Case 3

real pointer **p**

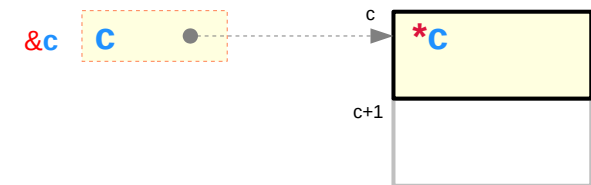
primitive data ***p**



Case 2

virtual pointer **c**

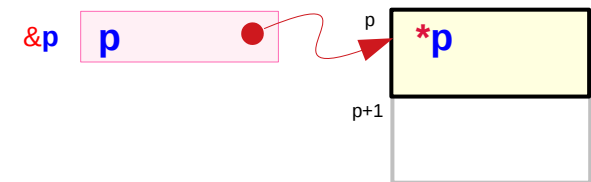
abstract data ***c**



Case 4

real pointer **p**

abstract data ***p**



subarray partitioning

★ $\text{sizeof}(\mathbf{c}) = \text{sizeof}(\mathbf{*c}) * N$
 $\text{sizeof}(\mathbf{p}) = \text{pointer size (4/8 bytes)}$

$\text{value}(\mathbf{c}+1) = \text{value}(\mathbf{c}) + \text{sizeof}(\mathbf{*c})$
 $\text{value}(\mathbf{p}+1) = \text{value}(\mathbf{p}) + \text{sizeof}(\mathbf{*p})$

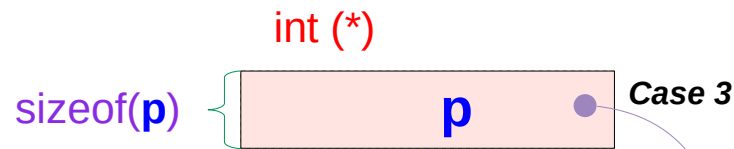
address replication

★ $\text{value}(\mathbf{\&c}) = \text{value}(\mathbf{c})$
 $\text{value}(\mathbf{\&p}) \neq \text{value}(\mathbf{p})$

Sizes of integer pointers

a pointer to an `int`

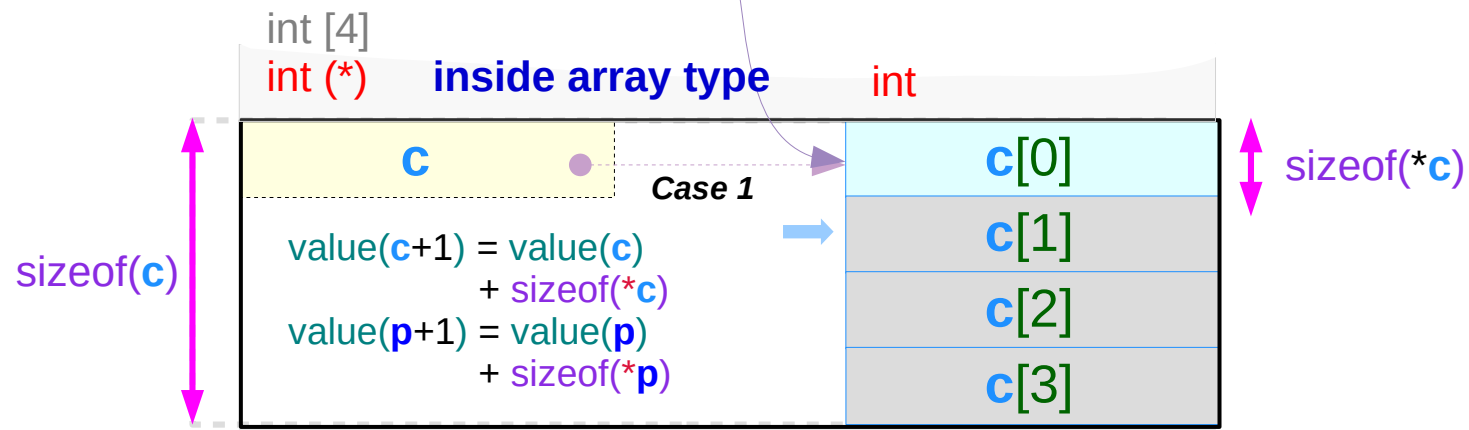
`sizeof(p)` = pointer size
= 8 bytes (64-bit CPU)
= 4 bytes (32-bit CPU)



`int (*p) ;`
`int c[4] ;`

an `int` array

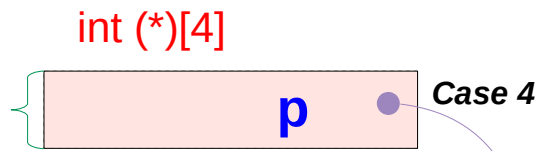
`sizeof(c)`
= `sizeof(*c) * 4`
= `sizeof(int) * 4`
= $4 * 4 = 16$ bytes



Case 1: virtual pointer `c` to primitive data `*c`
Case 3: real pointer `p` to primitive data `*p`

$$\text{value}(p+1) = \text{value}(p) + \text{sizeof}(*p)$$

Sizes of integer pointers



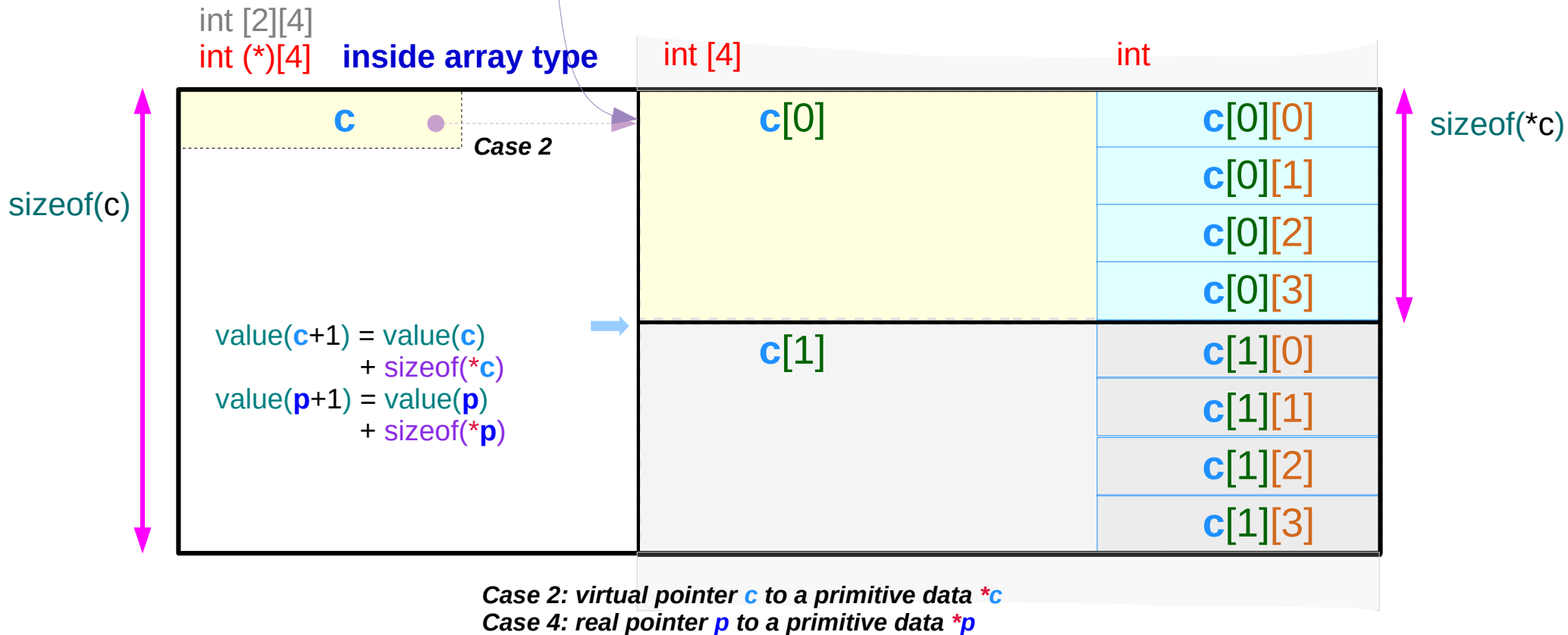
a pointer to an int

$\text{sizeof}(p) = \text{pointer size}$
 = 8 bytes (64-bit CPU)
 = 4 bytes (32-bit CPU)

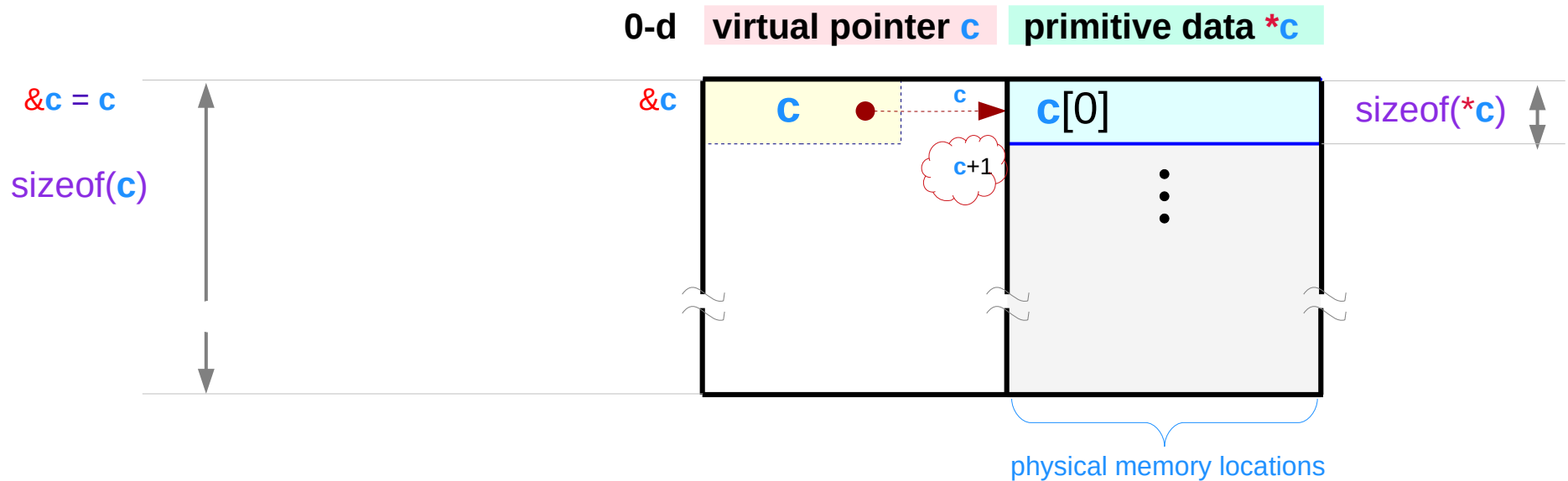
an int array

$\text{sizeof}(c)$
 = $\text{sizeof}(*c) * 2$
 = $\text{sizeof}(c[0]) * 2$
 = $\text{sizeof}(*c[0]) * 4 * 2$
 = $\text{sizeof}(\text{int}) * 4 * 2$
 = $4 * 4 * 2 = 32$ bytes

`int (*p) [4] ;`
`int c[2][4] ;`



Case 1: virtual pointer **c** to a primitive data ***c**



Abstract data **c** subarray partitioning

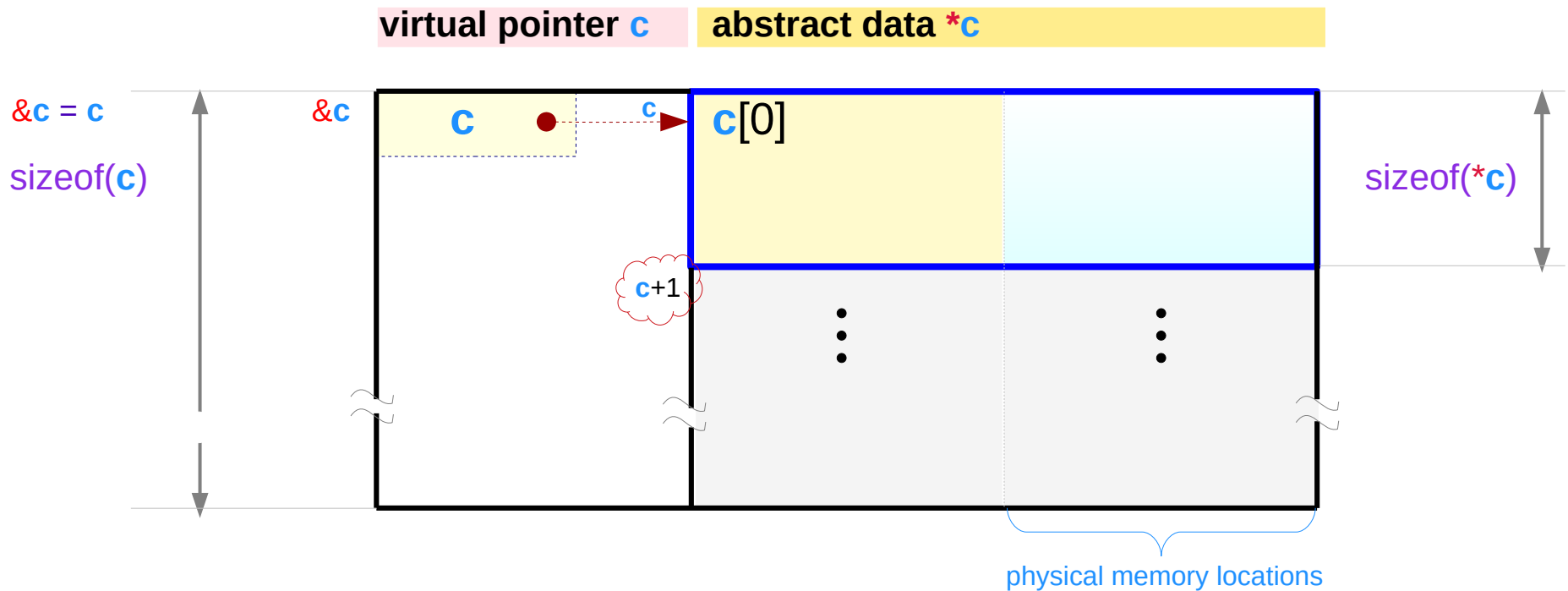
$$\star \text{sizeof}(c) = \text{sizeof}(*c) * N$$

$$\text{value}(c+1) = \text{value}(c) + \text{sizeof}(*c)$$

Virtual pointer **c** address replication

$$\star \text{value}(\&c) = \text{value}(c)$$

Case 2: virtual pointer **c** to an abstract data ***c**



Abstract data **c**, ***c** subarray partitioning

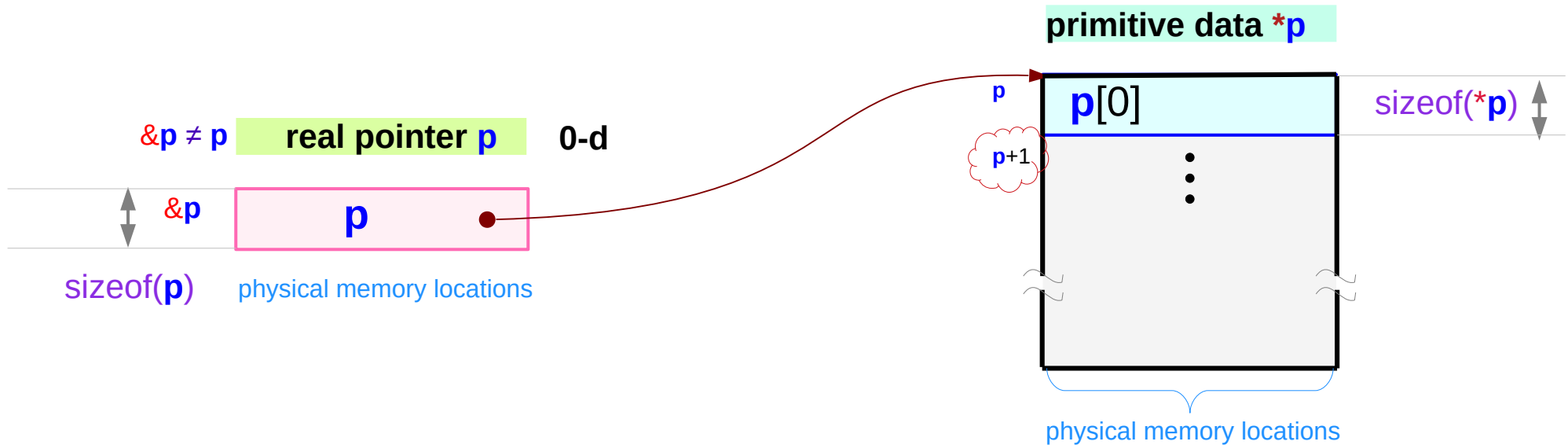
$$\star \text{sizeof}(c) = \text{sizeof}(*c) * N$$

$$\text{value}(c+1) = \text{value}(c) + \text{sizeof}(*c)$$

Virtual pointer **c** address replication

$$\star \text{value}(\&c) = \text{value}(c)$$

Case 3: real pointer **p** to a primitive data ***p**



real pointer size

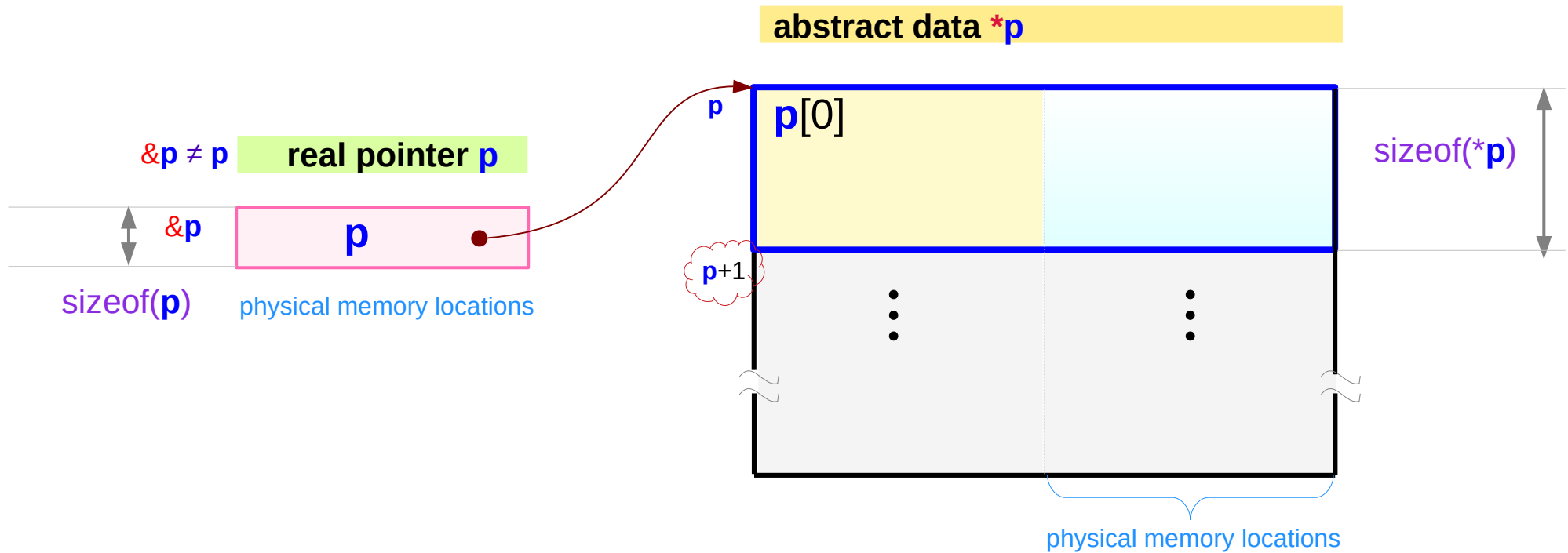
$\text{sizeof}(\mathbf{p}) = \text{pointer size (4/8 bytes)}$

$\text{value}(\mathbf{p}+1) = \text{value}(\mathbf{p}) + \text{sizeof}(*\mathbf{p})$

unique pointer value and address

$\text{value}(\&\mathbf{p}) \neq \text{value}(\mathbf{p})$

Case 4: real pointer p to an abstract data $*p$



real pointer size

$\text{sizeof}(p)$ = pointer size (4/8 bytes)

$\text{value}(p+1) = \text{value}(p) + \text{sizeof}(*p)$

unique pointer value and address

$\text{value}(\&p) \neq \text{value}(p)$

Properties of array pointers

virtual pointer **c** in an array **c**

Abstract data **c[N][] ... []** subarray partitioning

$$\star \text{sizeof}(\mathbf{c}) = \text{sizeof}(\star\mathbf{c}) * \mathbf{N}$$

$$\text{value}(\mathbf{c}+1) = \text{value}(\mathbf{c}) + \text{sizeof}(\star\mathbf{c})$$

implicit array pointer

Virtual pointer **(*c)[] ... []** address replication

$$\star \text{value}(\&\mathbf{c}) = \text{value}(\mathbf{c})$$

array pointer **p**

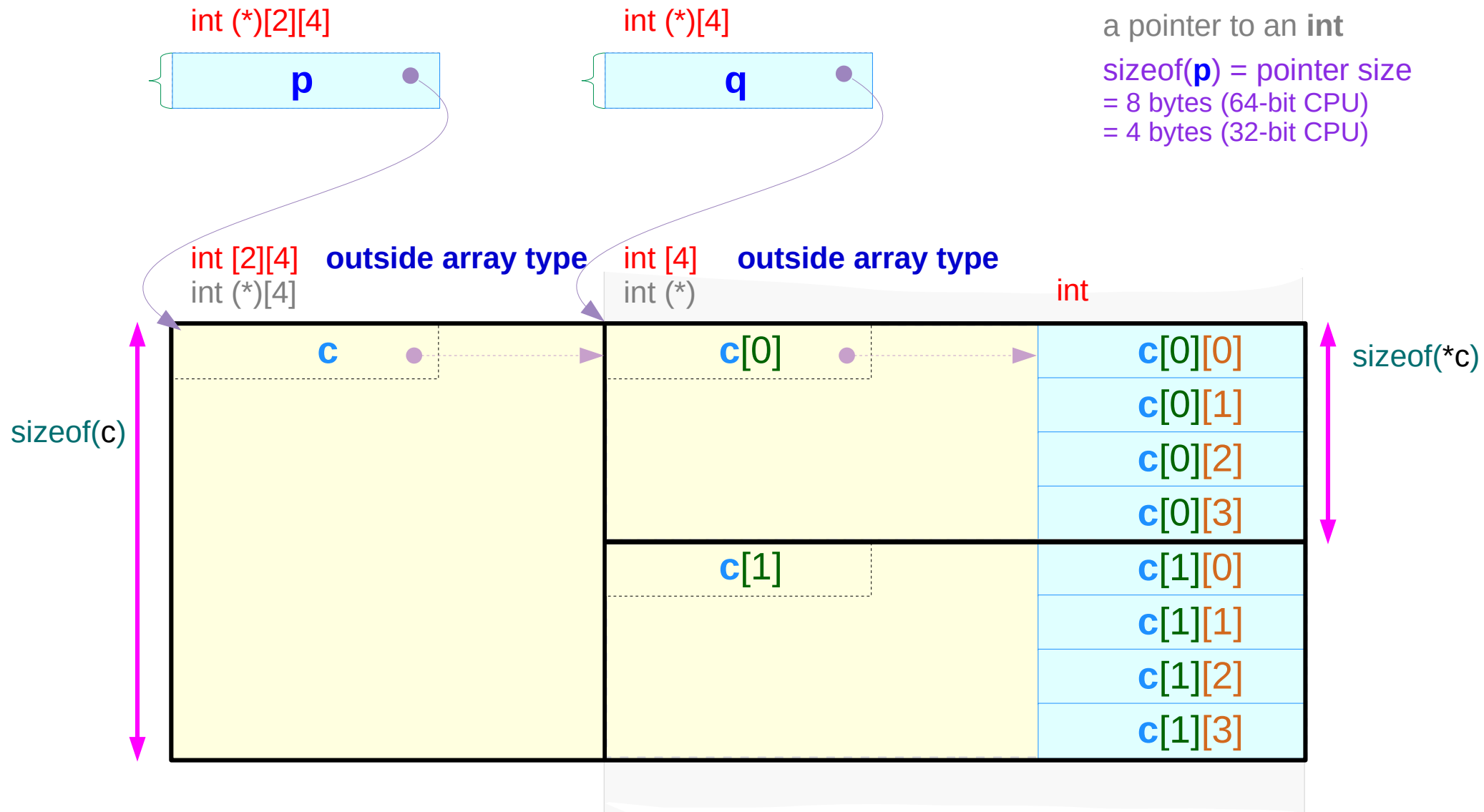
$$\text{sizeof}(\mathbf{p}) = \text{pointer size (4/8 bytes)}$$

$$\text{value}(\mathbf{p}+1) = \text{value}(\mathbf{p}) + \text{sizeof}(\star\mathbf{p})$$

explicit array pointer

$$\text{value}(\&\mathbf{p}) \neq \text{value}(\mathbf{p})$$

Sizes of integer pointers



Array element notation $p[i]$

Dereference notation $*(p+i)$

`int a[4];` `int (*p) [4] = &a;`

1-d array, 1-d array pointer

`int c[3][4];` `int (*p) [3][4] = &c;`

2-d array, 2-d array pointer

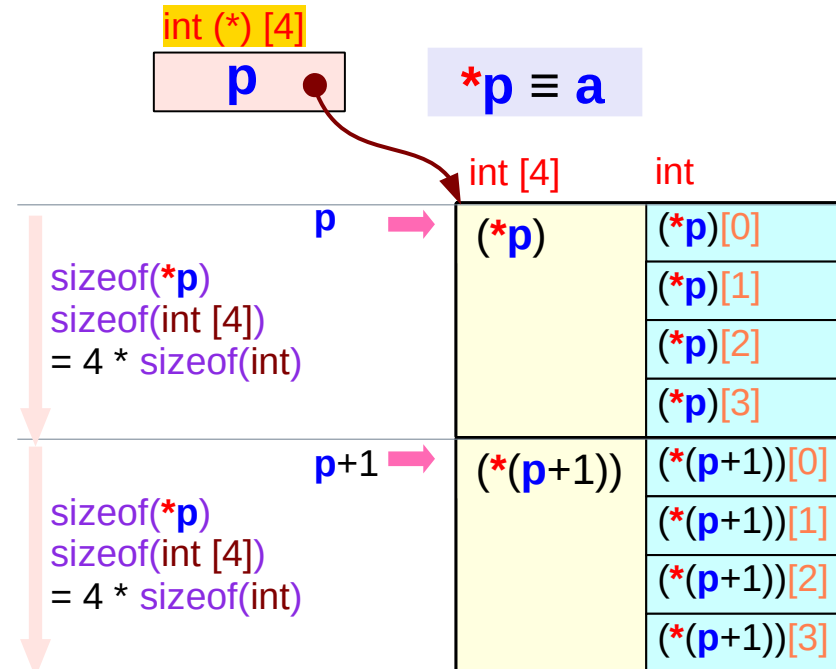
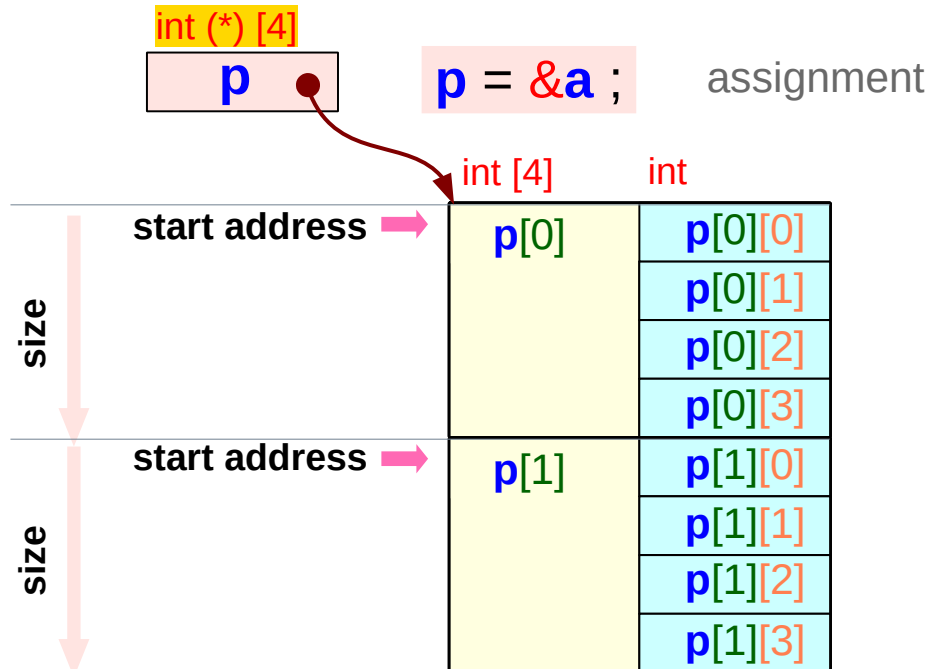
`int c[3][4];` `int (* q)[4] = c;`

2-d array, 1-d array pointer

1-d array, 1-d array pointer **p** – size and start address

```
int (*p) [4] = &a;
```

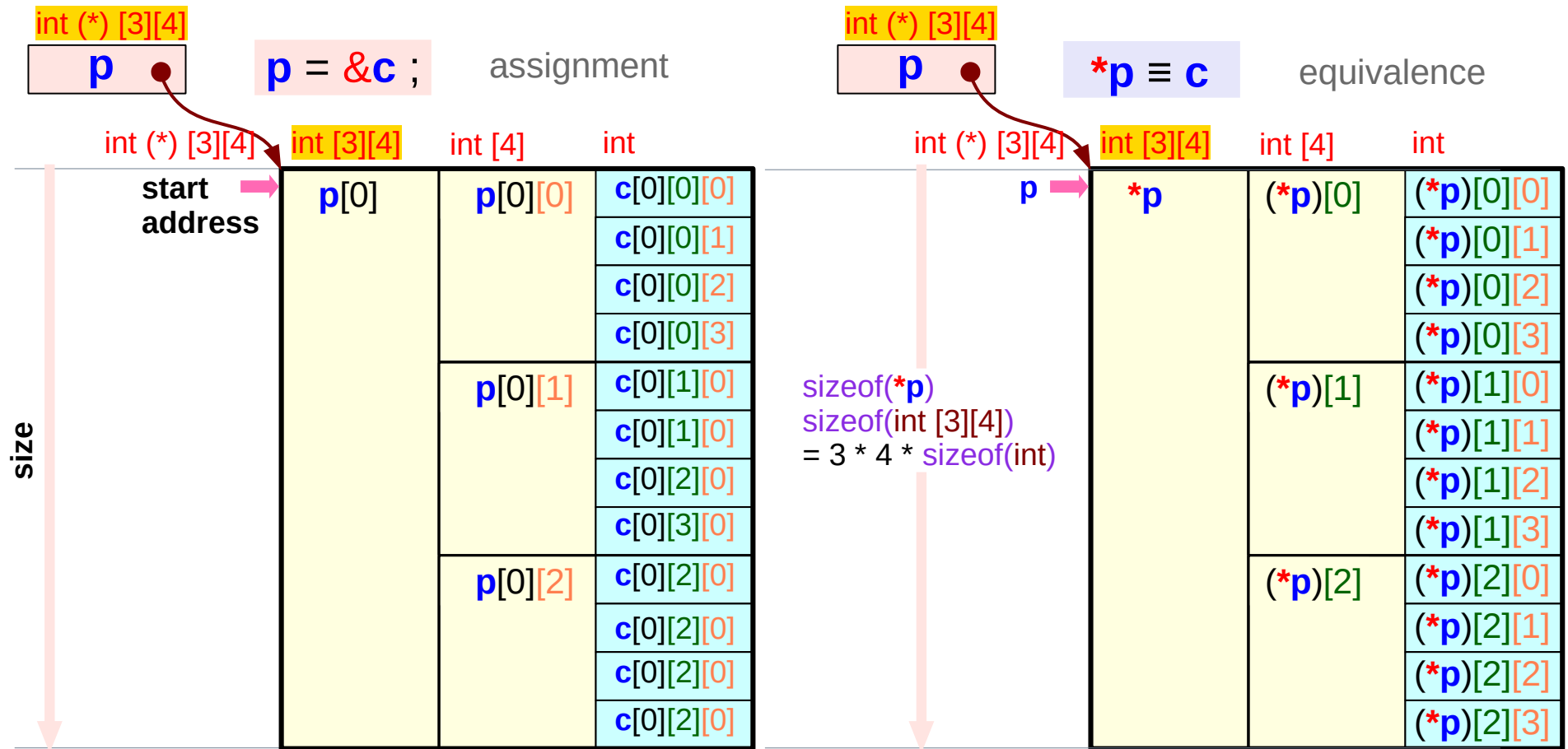
$$\begin{aligned} \text{value}(p+1) - \text{value}(p) &= \text{sizeof}(*p) = \text{sizeof}(\text{int } [4]) \\ &= (\text{long})(p+1) - (\text{long})(p) = 4 * \text{sizeof}(\text{int}) \end{aligned}$$



2-d array, 2-d array pointer **p** – size and start address

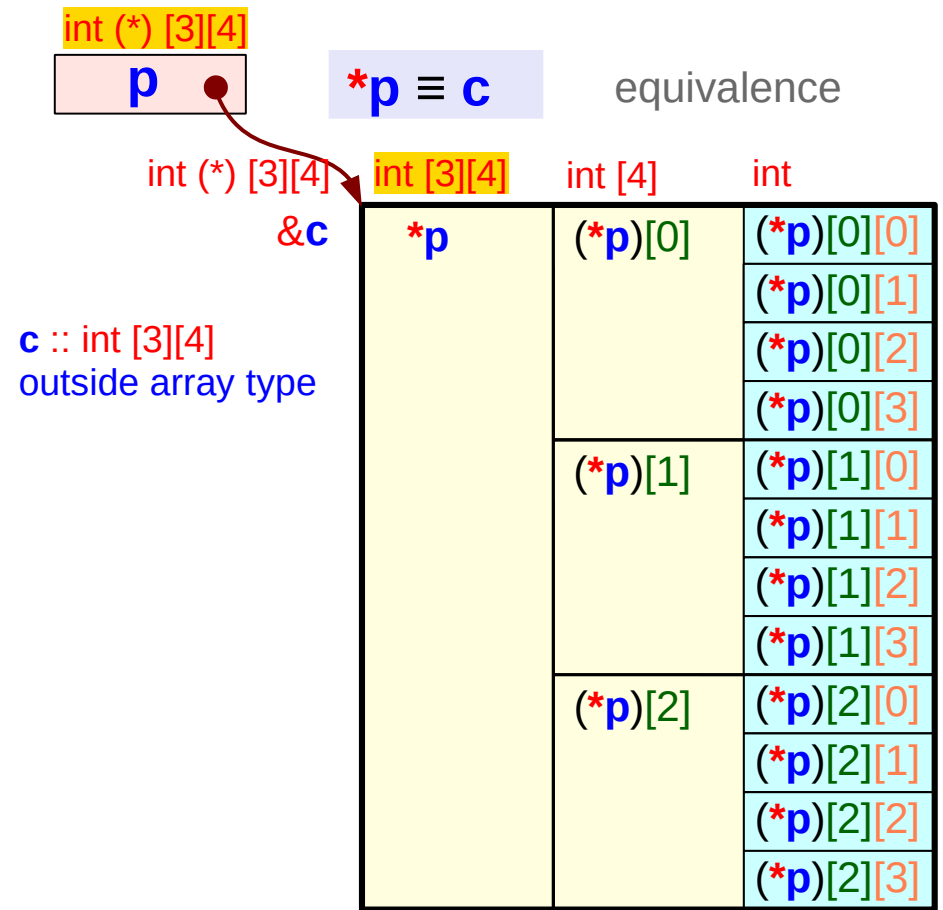
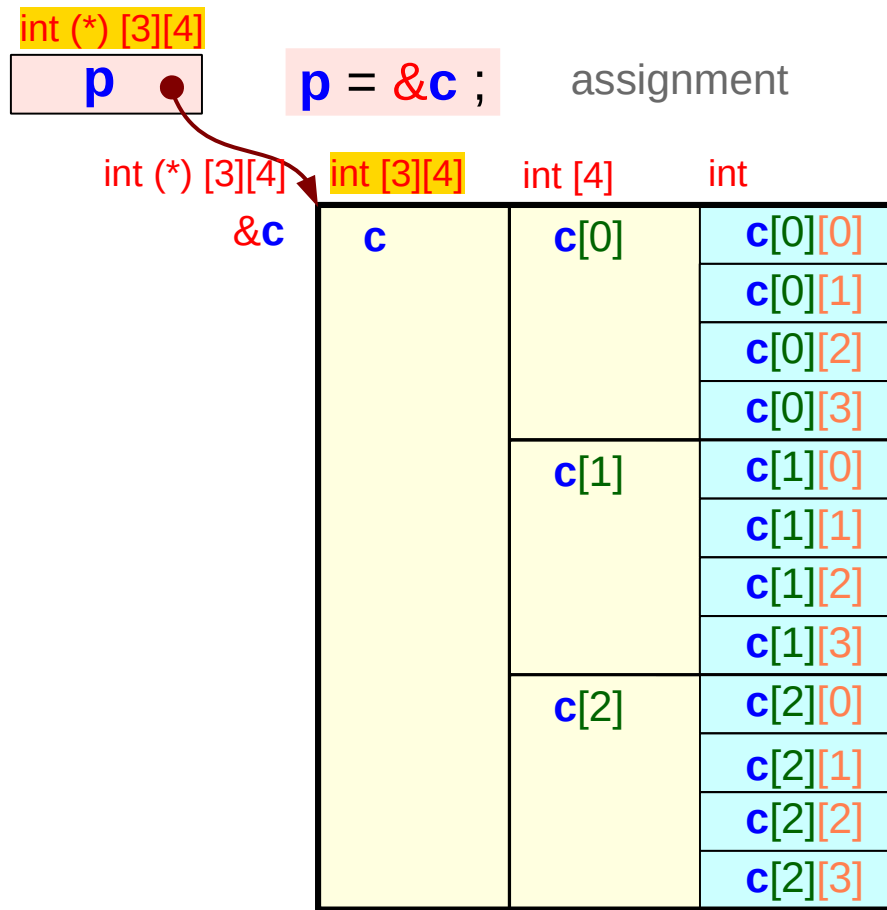
```
int (*p) [3][4] = &c;
```

$$\begin{aligned} \text{value}(p+1) - \text{value}(p) &= \text{sizeof}(*p) = \text{sizeof}(\text{int } [3][4]) \\ &= (\text{long}) (p+1) - (\text{long}) (p) = 3 * 4 * \text{sizeof}(\text{int}) \end{aligned}$$



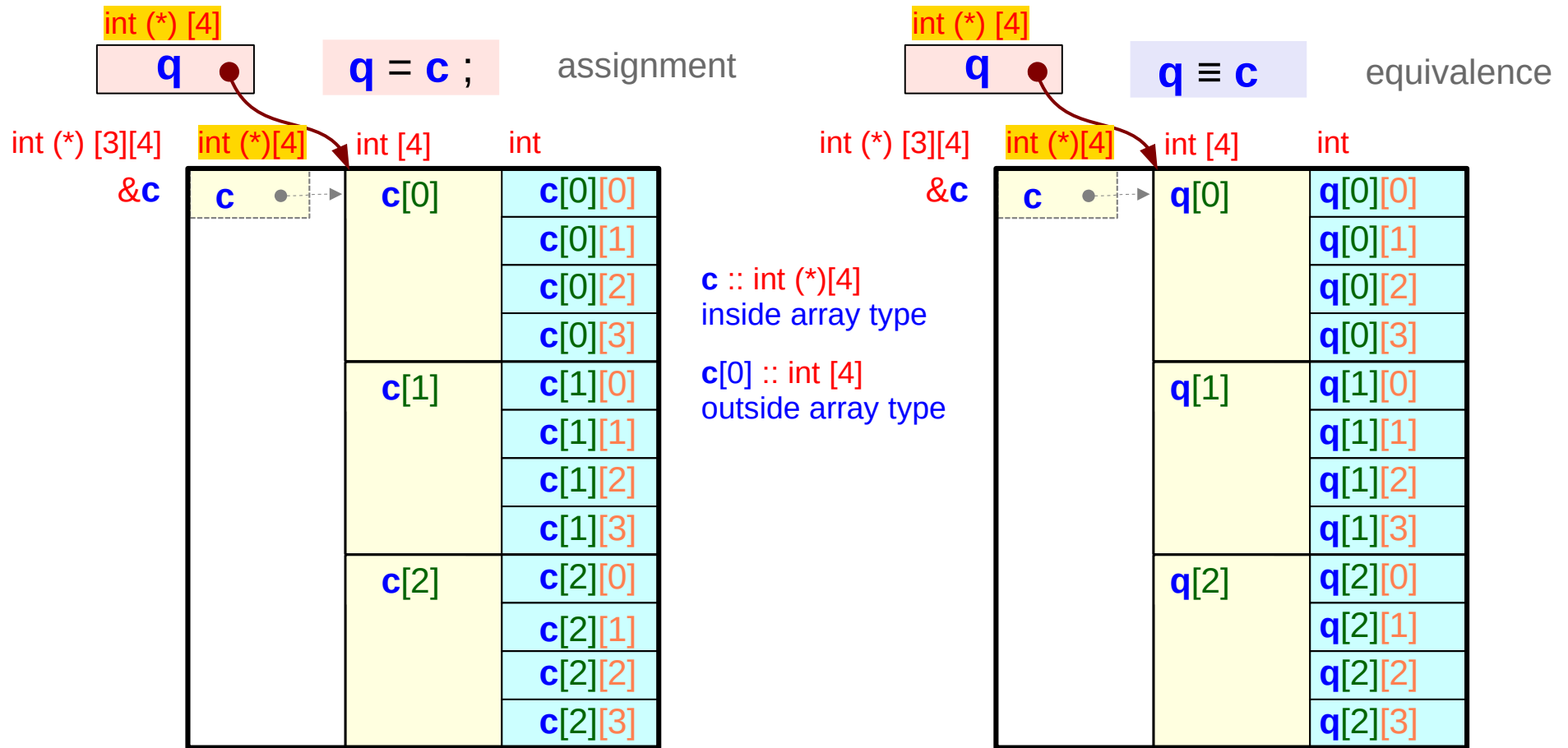
2-d array, 2-d array pointer **p** – outside array type

<code>int c [3][4];</code>	assignment	dereference	equivalence	equivalence
<code>int (*p) [3][4];</code>	<code>p = &c</code>	<code>*p ≡ c</code>	<code>c ≡ &c[0]</code>	<code>c[0] ≡ &c[0][0]</code>



2-d array, 1-d array pointer **q**– inside array type

<code>int c [3][4];</code>	assignment	dereference	equivalence	equivalence
<code>int (*q) [4];</code>	<code>q = c</code>	<code>q ≡ c</code>	<code>c ≡ &c[0]</code>	<code>c[0] ≡ &c[0][0]</code>



Incrementing a 2-d array pointer

Incrementing a 1-d array pointer

Subarray sizes referenced by array pointers **p** and **q**

```
int c [3][4] ;
```

2-d array **c**

`sizeof(c) = array size`
3·4·4 bytes

```
int c [3][4]
```

`c : sizeof(int [3][4])`

```
int c [3][4]            i = 0,1,2
```

`c[i] : sizeof(int [4])`

```
int c [3][4]            i = 0,1,2  
                         j = 0,1,2,3
```

`c[i][j] : sizeof(int)`

```
int (*p) [3][4] ;
```

2-d array pointer **p**

`sizeof(p) = pointer size`
4 or 8 bytes

```
int (*p) [3][4]
```

`(*p) or p[0] : sizeof(int [3][4])`

```
int (*p) [3][4]            i = 0,1,2
```

`(*p)[i] or p[0][i] : sizeof(int [4])`

```
int (*p) [3][4]            i = 0,1,2  
                                 j = 0,1,2,3
```

`(*p)[i][j] or p[0][i][j] : sizeof(int)`

```
int (*q) [4] ;
```

1-d array pointer **q**

`sizeof(q) = pointer size`
4 or 8 bytes

```
int (*q) [4]
```

`(*q) or q[0] : sizeof(int [4])`

```
int (*q) [4]            i = 0,1,2
```

`(*q+i) or q[i] : sizeof(int [4])`

```
int (*q) [4]            i = 0,1,2  
                                 j = 0,1,2,3
```

`(*q+i)+j) or q[i][j] : sizeof(int)`

2-d array **c**, 2-d array pointer **p**, 1-d array pointer **q**

```
int c [3][4] ;
```

```
int (*p) [3][4] = &c;
```

```
int (*q) [4] = c ;
```

2-d

(*p)[i][j]

i = 0,1,2
j = 0,1,2,3

3-d

p[0][i][j]

i = 0,1,2
j = 0,1,2,3



***(p+m)[i][j]**

i = 0,1,2
j = 0,1,2,3

2-d



p[m][i][j]

i = 0,1,2
j = 0,1,2,3

3-d

Basic Reference

Extended Reference

1-d

(*q)[j]

j = 0,1,2,3

2-d

q[0][j]

j = 0,1,2,3



***(q+i)[j]**

i = 0,1,2
j = 0,1,2,3

1-d



q[i][j]

i = 0,1,2
j = 0,1,2,3

2-d

Basic Reference

Extended Reference

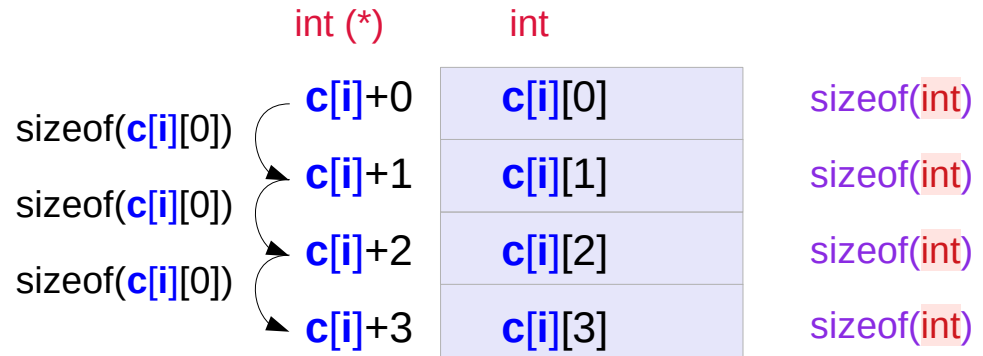
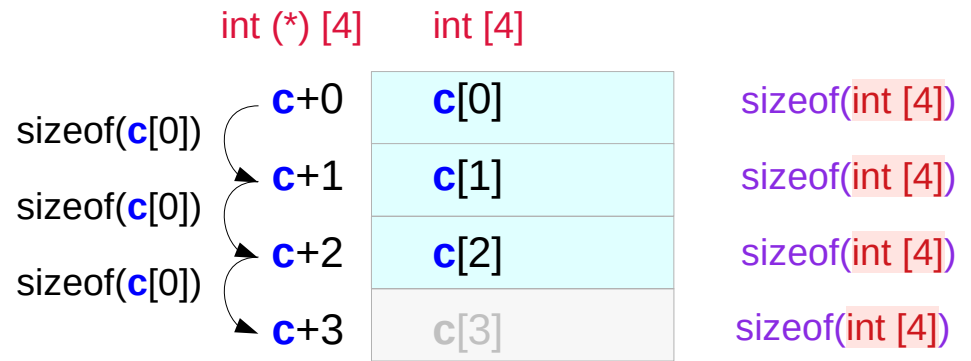
2-d array c

```
int c [3][4] ;
```

(int [4]) **c[i]** i = 0,1,2

```
int c [3][4] ;
```

(int) **c[i][j]** i = 0,1,2
j = 0,1,2,3



$$\text{value}(\mathbf{c} + \mathbf{i}) = \text{value}(\mathbf{c}) + \mathbf{i} * \text{sizeof}(*\mathbf{c})$$

$$\text{value}(\mathbf{c}[\mathbf{i}] + \mathbf{j}) = \text{value}(\mathbf{c}[\mathbf{i}]) + \mathbf{j} * \text{sizeof}(*\mathbf{c}[\mathbf{0}])$$

$$\text{value}(\mathbf{c} + \mathbf{i}) = \text{value}(\mathbf{c}[\mathbf{i}]) \quad (\text{address replication})$$

$$\&\mathbf{c}[\mathbf{i}][\mathbf{j}] = \text{value}(\mathbf{c}) + \mathbf{i} * \text{sizeof}(*\mathbf{c}) + \mathbf{j} * \text{sizeof}(*\mathbf{c}[\mathbf{0}])$$

Incrementing a 2-d array pointer p

```
int (*p) [3][4] = &c ;
```

```
int c [3][4] ;
```

p+0		sizeof(*p)	sizeof(int [3][4])	sizeof(2-d array)
p+1		sizeof(*p)	sizeof(int [3][4])	sizeof(2-d array)
p+2		sizeof(*p)	sizeof(int [3][4])	sizeof(2-d array)
p+3		sizeof(*p)	sizeof(int [3][4])	sizeof(2-d array)

(*p)[i][j] $\begin{matrix} i = 0,1,2 \\ j = 0,1,2,3 \end{matrix}$
p[0][i][j] $\begin{matrix} i = 0,1,2 \\ j = 0,1,2,3 \end{matrix}$

*(p+0)	↔	p[0]
*(p+1)	↔	p[1]
*(p+2)	↔	p[2]
*(p+3)	↔	p[3]

(*p)[i][j]	↔	p[0][i][j]	used for c[3][4]
*(p+1)[i][j]	↔	p[1][i][j]	
*(p+2)[i][j]	↔	p[2][i][j]	
*(p+3)[i][j]	↔	p[3][i][j]	

Incrementing a 1-d array pointer **q**

```
int (*q) [4] = c ;
```

```
int c [3][4] ;
```

q+0	}	sizeof(*q)	sizeof(int [4])	sizeof(1-d array)
q+1		sizeof(*q)	sizeof(int [4])	sizeof(1-d array)
q+2		sizeof(*q)	sizeof(int [4])	sizeof(1-d array)
q+3		sizeof(*q)	sizeof(int [4])	sizeof(1-d array)

$(*(q+i))[j]$ $i = 0,1,2$
 $j = 0,1,2,3$

$q[i][j]$ $i = 0,1,2$
 $j = 0,1,2,3$

$*(q+0)$	↔	$q[0]$
$*(q+1)$	↔	$q[1]$
$*(q+2)$	↔	$q[2]$
$*(q+3)$	↔	$q[3]$

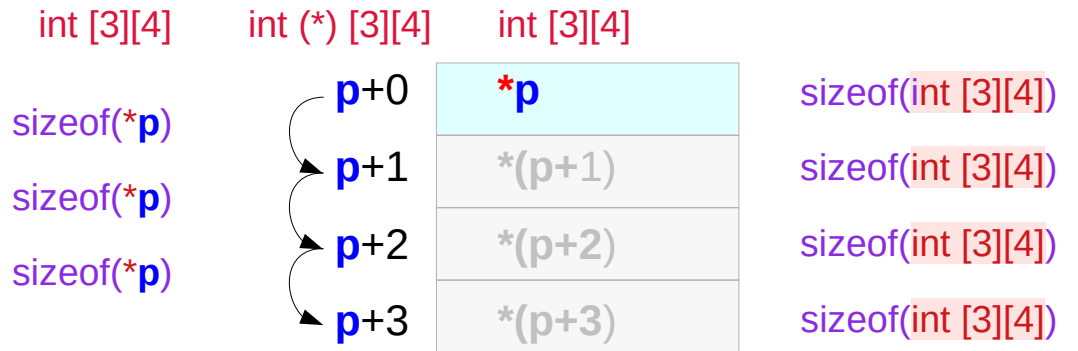
$(*(q+0))[j]$	↔	$q[0][j]$
$(*(q+1))[j]$	↔	$q[1][j]$
$(*(q+2))[j]$	↔	$q[2][j]$
$(*(q+3))[j]$	↔	$q[3][j]$

used for $c[3][4]$

2-d array pointer p

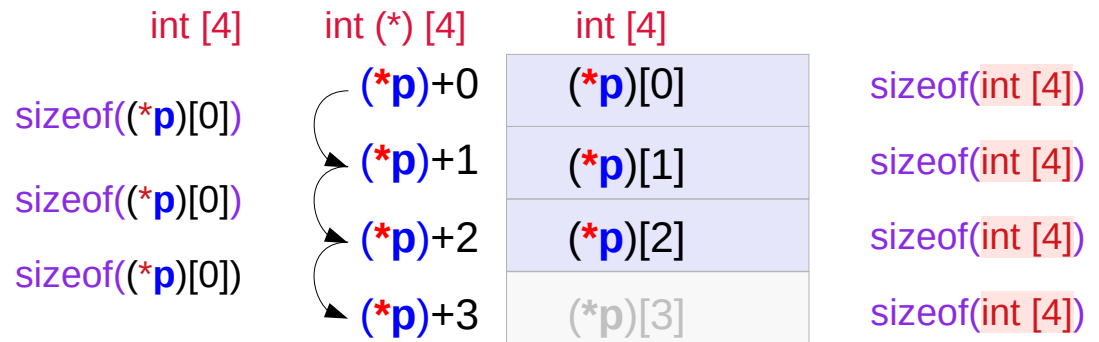
```
int (*p) [3][4] ;
```

(int [3][4]) (*p)



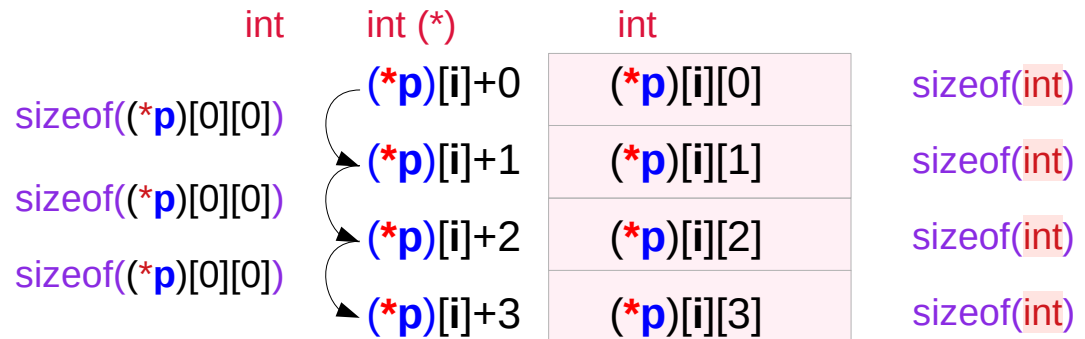
```
int (*p) [3][4] ;
```

(int [4]) (*p)[i] i = 0,1,2



```
int (*p) [3][4] ;
```

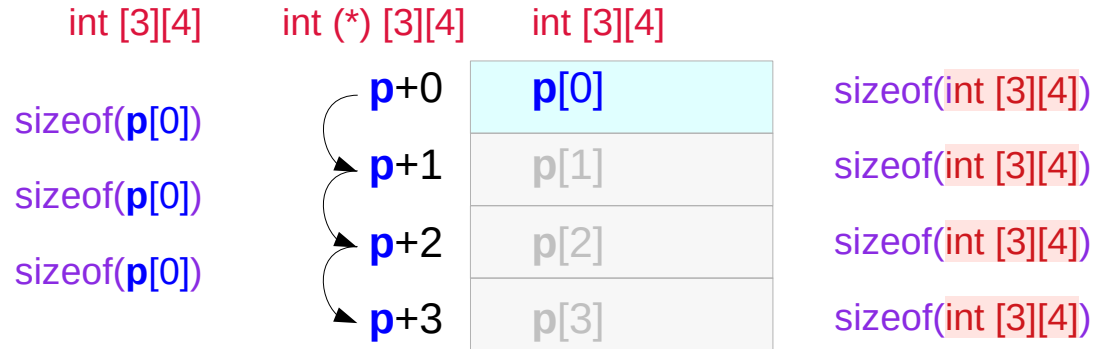
(int) (*p)[i][j] i = 0,1,2
j = 0,1,2,3



2-d array pointer p

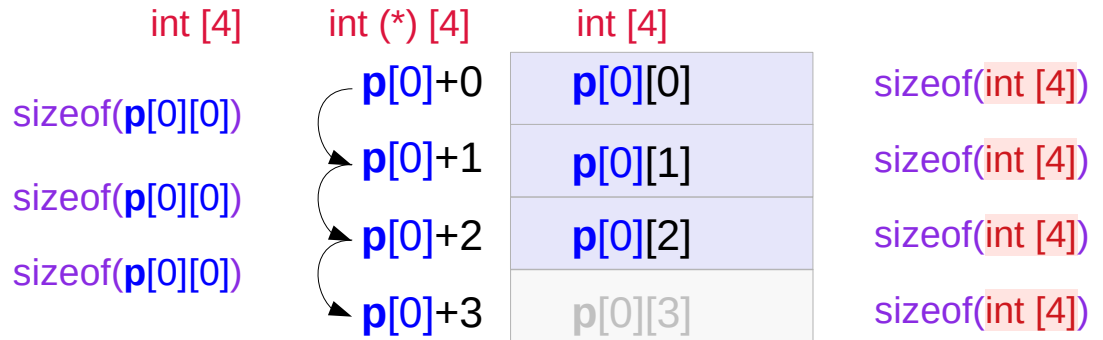
```
int (*p) [3][4] ;
```

(int [3][4]) p[0]



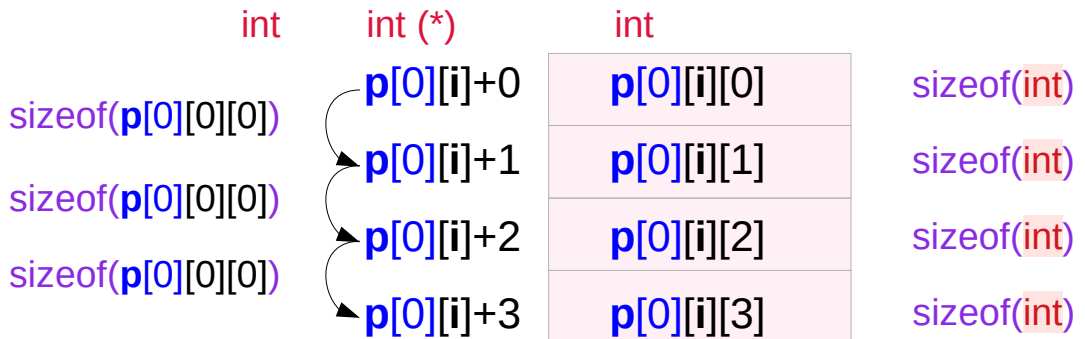
```
int (*p) [3][4] ;
```

(int [4]) p[0][i] i = 0,1,2



```
int (*p) [3][4] ;
```

(int) p[0][i][j] i = 0,1,2 j = 0,1,2,3



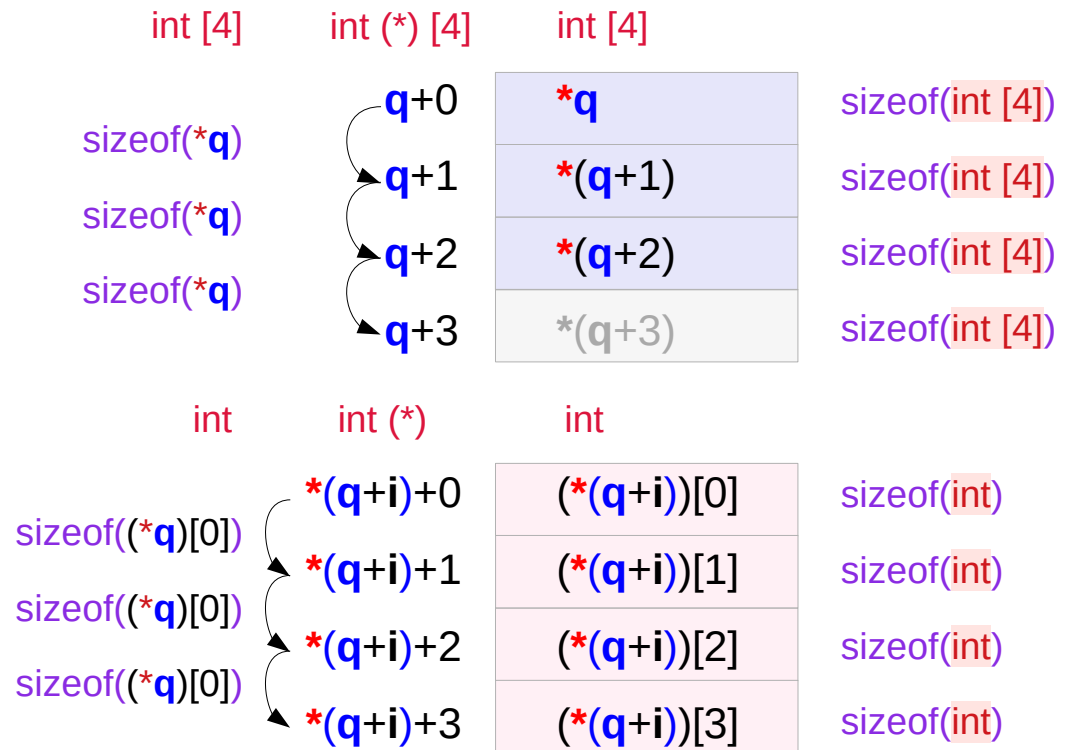
1-d array pointer q

```
int (*q) [4] ;
```

```
(int [4]) *(q+i)  i = 0,1,2
```

```
int (*q) [4] ;
```

```
(int) (*(q+i))[j]  i = 0,1,2  
j = 0,1,2,3
```



1-d array pointer q

```
int (*q) [4] ;
```

(int [4]) **q**[i] i = 0,1,2

```
int (*q) [4] ;
```

(int) **q**[i][j] i = 0,1,2
 j = 0,1,2,3

int [4]

sizeof(q[0])

sizeof(q[0])

sizeof(q[0])

int (*) [4]

q+0

q+1

q+2

q+3

int [4]

q[0]

q[1]

q[2]

q[3]

sizeof(int [4])

sizeof(int [4])

sizeof(int [4])

sizeof(int [4])

int

sizeof(q[0][0])

sizeof(q[0][0])

sizeof(q[0][0])

int (*)

q[i]+0

q[i]+1

q[i]+2

q[i]+3

int

q[i][0]

q[i][1]

q[i][2]

q[i][3]

sizeof(int)

sizeof(int)

sizeof(int)

sizeof(int)

Array element address

```
int c [3][4] ;
```

$$\text{value}(\mathbf{c} + \mathbf{i}) = \text{value}(\mathbf{c}) + \mathbf{i} * \text{sizeof}(*\mathbf{c})$$
$$\text{value}(\mathbf{c}[\mathbf{i}] + \mathbf{j}) = \text{value}(\mathbf{c}[\mathbf{i}]) + \mathbf{j} * \text{sizeof}(*\mathbf{c}[\mathbf{i}])$$
$$\text{value}(\mathbf{c} + \mathbf{i}) = \text{value}(\mathbf{c}[\mathbf{i}])$$

address replication

$$\begin{aligned} \&\mathbf{c}[\mathbf{i}][\mathbf{j}] &= \text{value}(\mathbf{c}[\mathbf{i}] + \mathbf{j}) \\ &= \text{value}(\mathbf{c}[\mathbf{i}]) + \mathbf{j} * \text{sizeof}(*\mathbf{c}[\mathbf{i}]) \\ &= \text{value}(\mathbf{c} + \mathbf{i}) + \mathbf{j} * \text{sizeof}(*\mathbf{c}[\mathbf{i}]) \\ &= \text{value}(\mathbf{c}) + \mathbf{i} * \text{sizeof}(*\mathbf{c}) \\ &\quad + \mathbf{j} * \text{sizeof}(*\mathbf{c}[\mathbf{i}]) \\ &= \text{value}(\mathbf{c}) + \mathbf{i} * 4 * 4 + \mathbf{j} * 4 \end{aligned}$$

```
int (*p) [3][4] = &c;
```

$$\text{value}(*\mathbf{p} + \mathbf{i}) = \text{value}(*\mathbf{p}) + \mathbf{i} * \text{sizeof}(**\mathbf{p})$$
$$\text{value}((*\mathbf{p})[\mathbf{i}] + \mathbf{j}) = \text{value}((*\mathbf{p})[\mathbf{i}]) + \mathbf{j} * \text{sizeof}>(*(*\mathbf{p})[\mathbf{i}])$$
$$\text{value}(*\mathbf{p} + \mathbf{i}) = \text{value}((*\mathbf{p})[\mathbf{i}])$$

address replication

$$\begin{aligned} \&(*\mathbf{p})[\mathbf{i}][\mathbf{j}] &= \text{value}((*\mathbf{p})[\mathbf{i}] + \mathbf{j}) \\ &= \text{value}((*\mathbf{p})[\mathbf{i}]) + \mathbf{j} * \text{sizeof}>(*(*\mathbf{p})[\mathbf{i}]) \\ &= \text{value}(*\mathbf{p} + \mathbf{i}) + \mathbf{j} * \text{sizeof}(*(*\mathbf{p})[\mathbf{i}]) \\ &= \text{value}(*\mathbf{p}) + \mathbf{i} * \text{sizeof}(**\mathbf{p}) \\ &\quad + \mathbf{j} * \text{sizeof}(*(*\mathbf{p})[\mathbf{i}]) \\ &= \text{value}(*\mathbf{p}) + \mathbf{i} * 4 * 4 + \mathbf{j} * 4 \end{aligned}$$

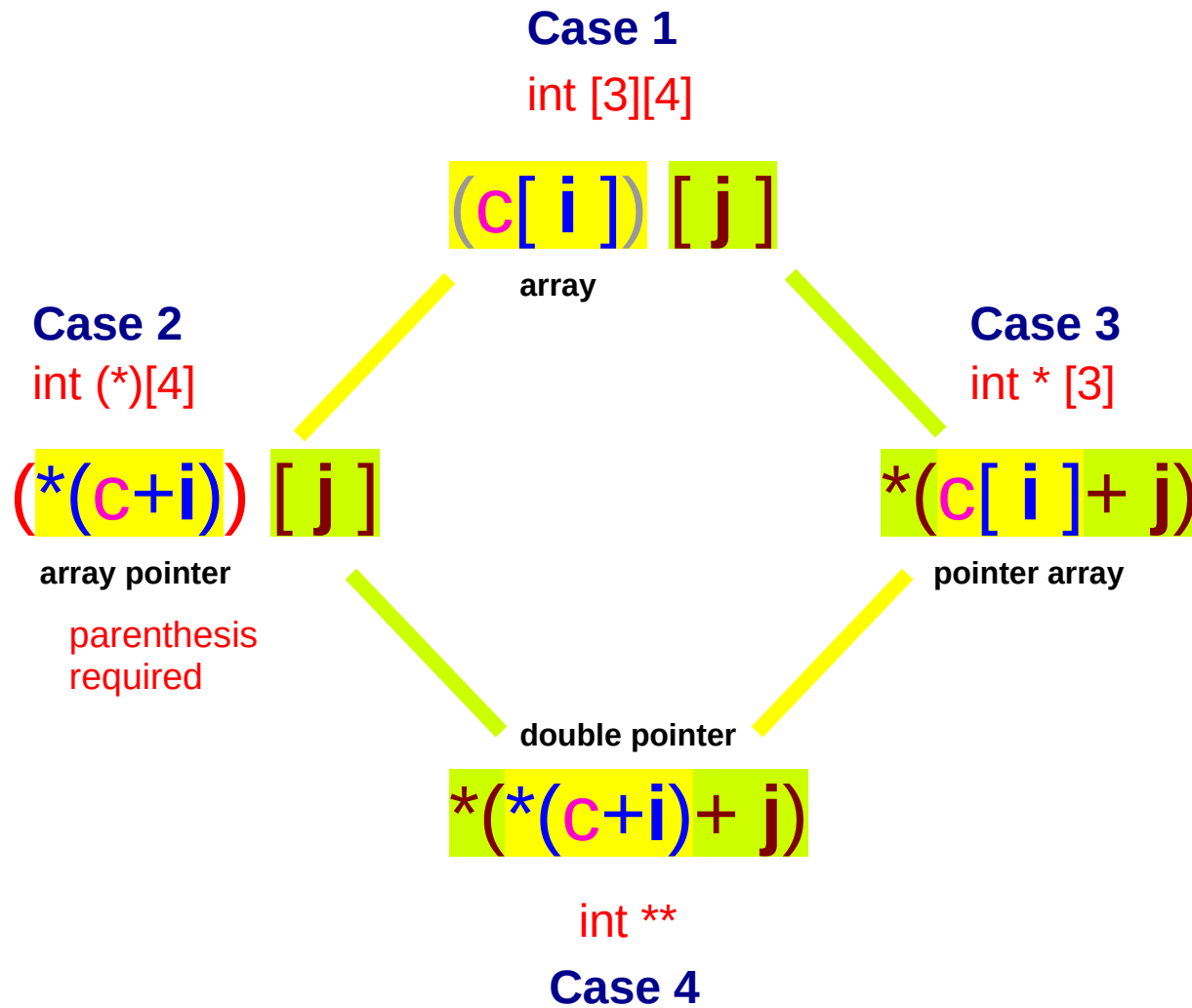
```
int (*q) [4] = c ;
```

$$\text{value}(\mathbf{q} + \mathbf{i}) = \text{value}(\mathbf{q}) + \mathbf{i} * \text{sizeof}(*\mathbf{q})$$
$$\text{value}(\mathbf{q}[\mathbf{i}] + \mathbf{j}) = \text{value}(\mathbf{q}[\mathbf{i}]) + \mathbf{j} * \text{sizeof}(*\mathbf{q}[\mathbf{i}])$$
$$\text{value}(\mathbf{q} + \mathbf{i}) = \text{value}(\mathbf{q}[\mathbf{i}])$$

address replication

$$\begin{aligned} \&\mathbf{q}[\mathbf{i}][\mathbf{j}] &= \text{value}(\mathbf{q}[\mathbf{i}] + \mathbf{j}) \\ &= \text{value}(\mathbf{q}[\mathbf{i}]) + \mathbf{j} * \text{sizeof}(*\mathbf{q}[\mathbf{i}]) \\ &= \text{value}(\mathbf{q} + \mathbf{i}) + \mathbf{j} * \text{sizeof}(*\mathbf{q}[\mathbf{i}]) \\ &= \text{value}(\mathbf{q}) + \mathbf{i} * \text{sizeof}(*\mathbf{q}) \\ &\quad + \mathbf{j} * \text{sizeof}(*\mathbf{q}[\mathbf{i}]) \\ &= \text{value}(\mathbf{q}) + \mathbf{i} * 4 * 4 + \mathbf{j} * 4 \end{aligned}$$

Array-pointer conversions from a 2-d array type



2-d array-pointer conversion types

Case 1 `int [3][4]`

`c[i][j]`

2-d array

Case 2 `int (*) [4]`

`(*(c+i))[j]`

1-d array pointer

Case 3 `int * [3]`

`*(c[i]+j)`

1-d array of pointers

Case 4 `int **`

`*(*(c+i)+j)`

double pointer

Types of **c**

Case 1	Case 2	Case 3	Case 4
<code>int [3][4]</code>	<code>int (*)[4]</code> array pointer <code>c</code> points to an array of 4 integers	<code>int * [3]</code> pointer array <code>c</code> is an array of 3 integer pointers	<code>int **</code> double pointer <code>c</code> points to an integer pointer
<code>c[i][j]</code>	$\equiv (*(\mathbf{c+i}))[\mathbf{j}]$	$\equiv *(\mathbf{c[i]+j})$	$\equiv *(*(\mathbf{c+i})+\mathbf{j})$
<code>&c[i][j]</code>	$\equiv *(\mathbf{c+i})+\mathbf{j}$	$\equiv \mathbf{c[i]+j}$	$\equiv *(\mathbf{c+i})+\mathbf{j}$

the address of `c[i][j]` is `*(c+i)+j` or `c[i]+j`

The row address is `*(c+i)` or `c[i]`

-
- **Pointer conversions in array types**
 - Simulating array accesses by real pointers
 - Dynamic memory allocation

Pointer conversions in 2-d and 1-d array types

Case 1 `int [3][4]`

`c[i][j]`

2-d array `c`

relaxing the 1st dimension



Case 2 `int (*) [4]`

`(*(c+i))[j]`

1-d array pointer

Case 3 `int * [3]`

`*(c[i]+j)`

1-d array `c`

relaxing the 1st dimension



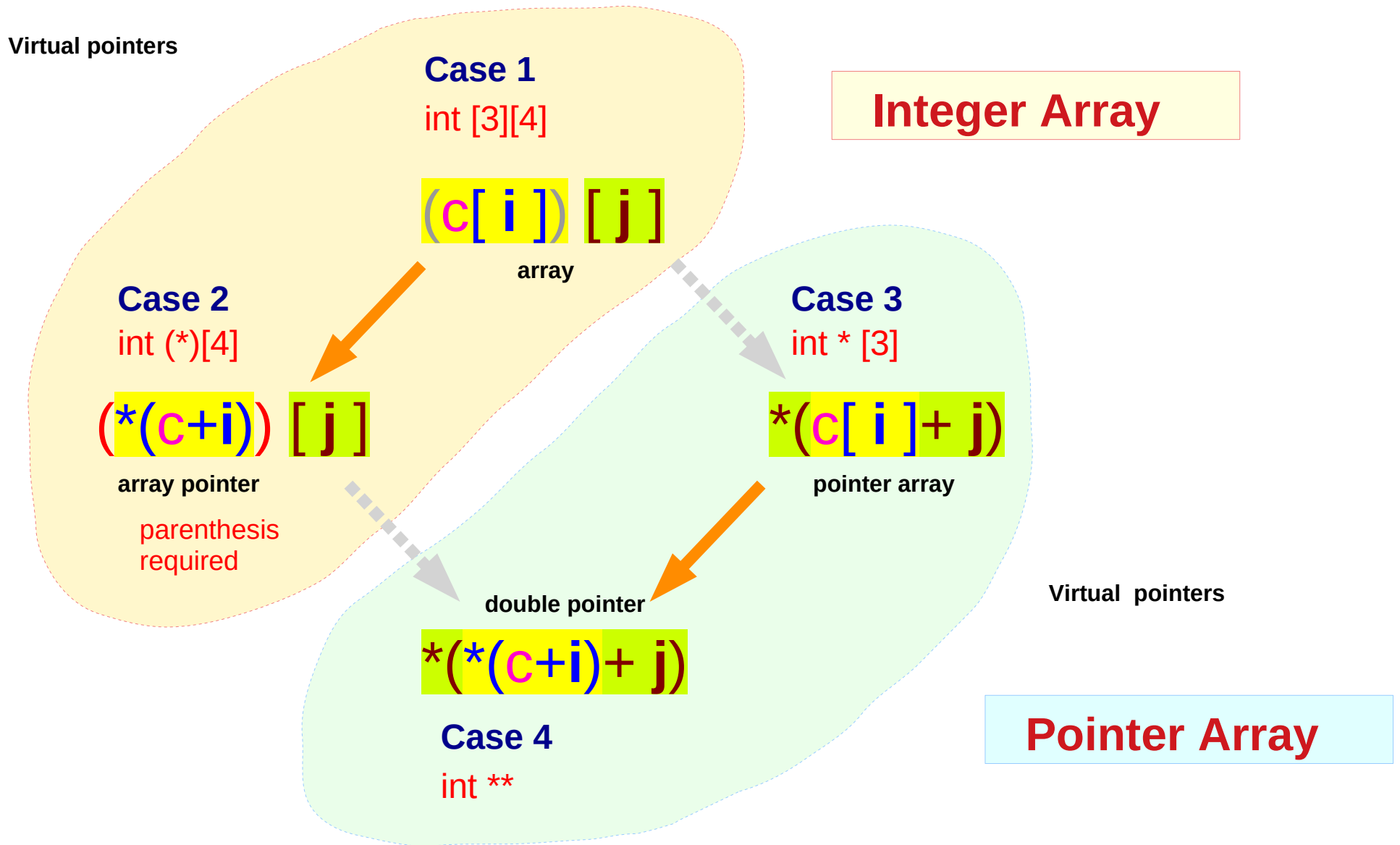
Case 4 `int **`

`**(*(c+i)+j)`

double pointer

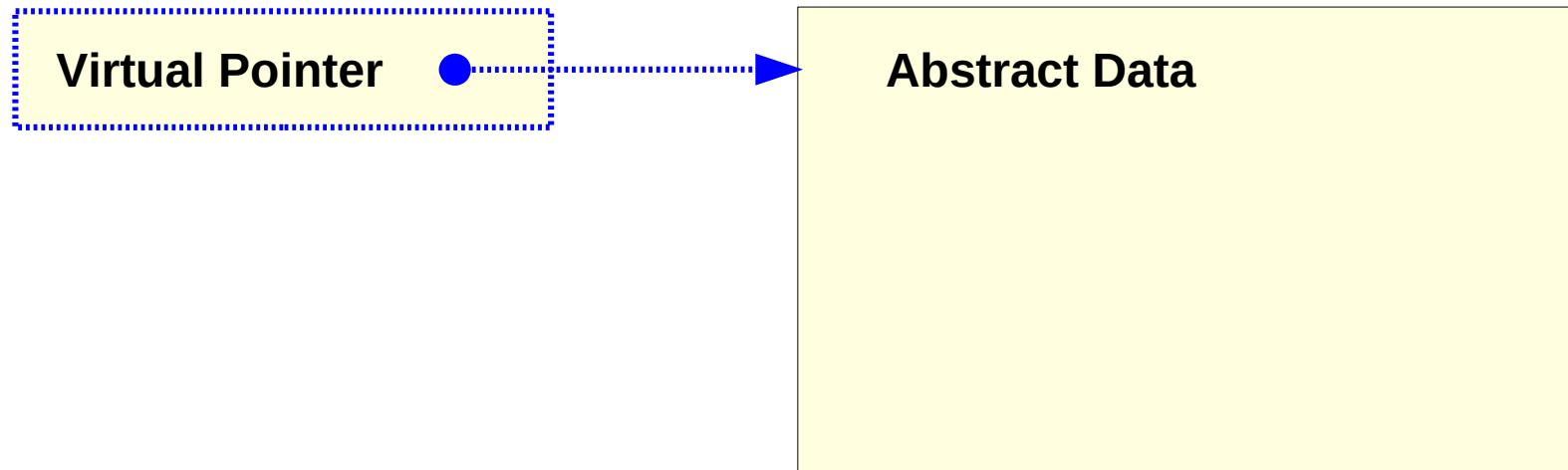
relaxing the 1st dimension

Relaxing the 1st dimension of an array



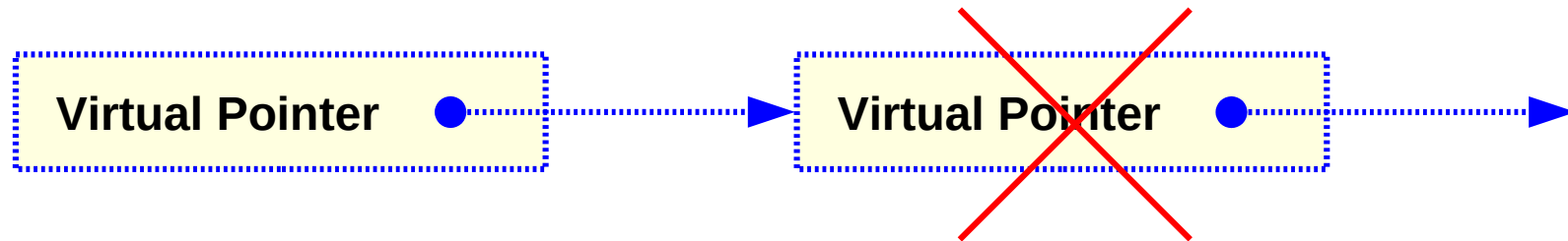
Pointer conversions in **2-d** and **1-d** array types

Only the 1st dimension can be relaxed



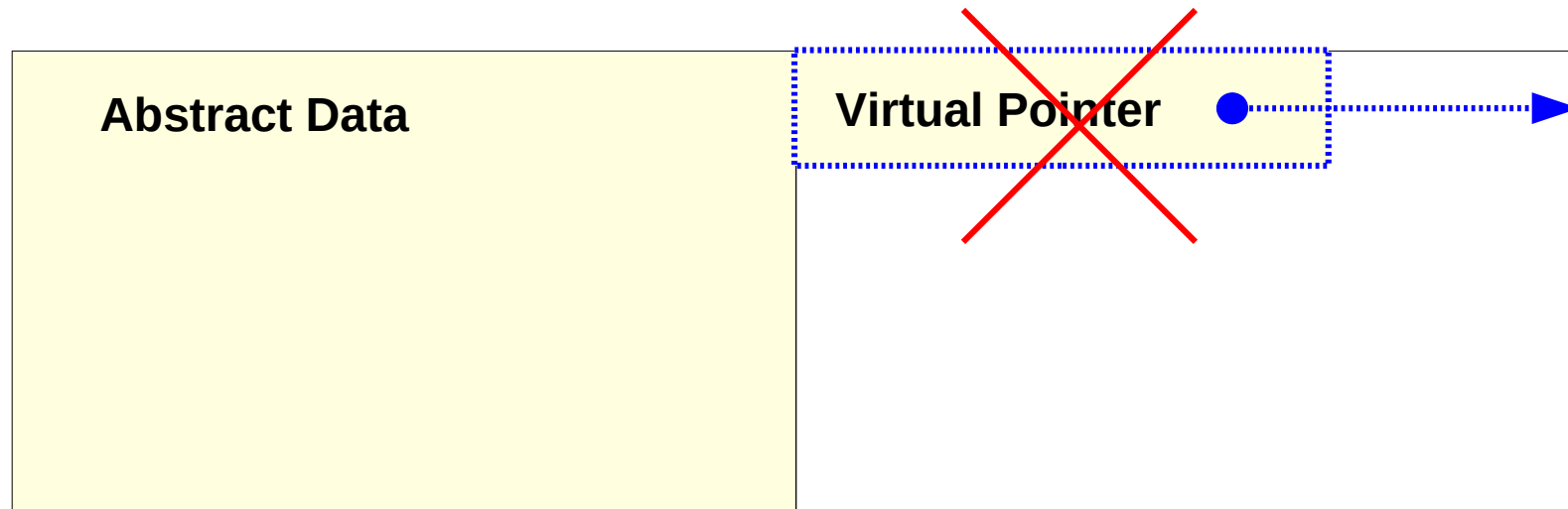
Pointer conversions in **2-d** and **1-d** array types

Only the 1st dimension can be relaxed

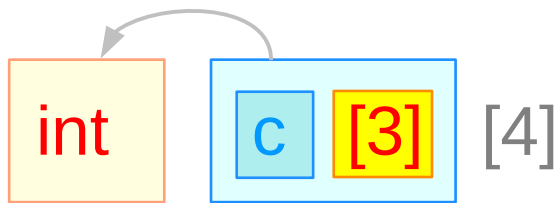


Pointer conversions in **2-d** and **1-d** array types

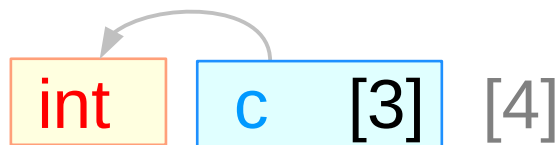
Only the 1st dimension can be relaxed



Case 3) 1-d array **c**, pointer **c[i]**



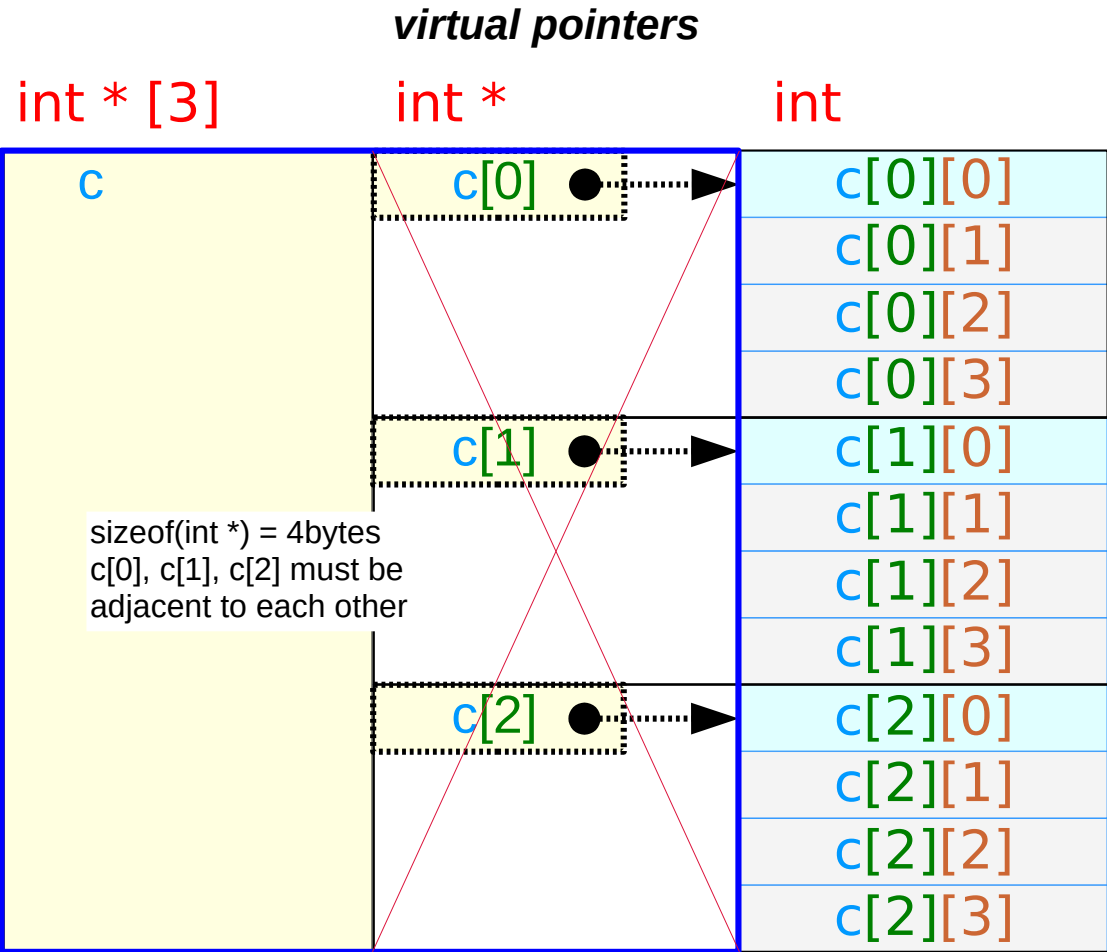
c 1-d array
type : `int * [3]`



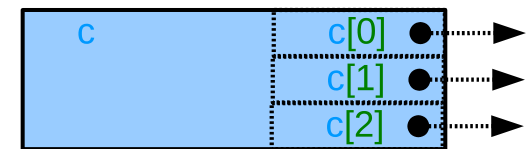
c[i] pointer
type : `int *`

Int pointer

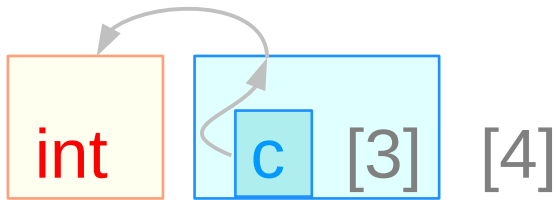
`*(c[i] + j)`



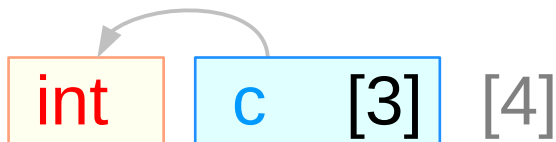
c is an array of 3 integer pointers



Case 4) double pointer **c**, pointer **c[i]**



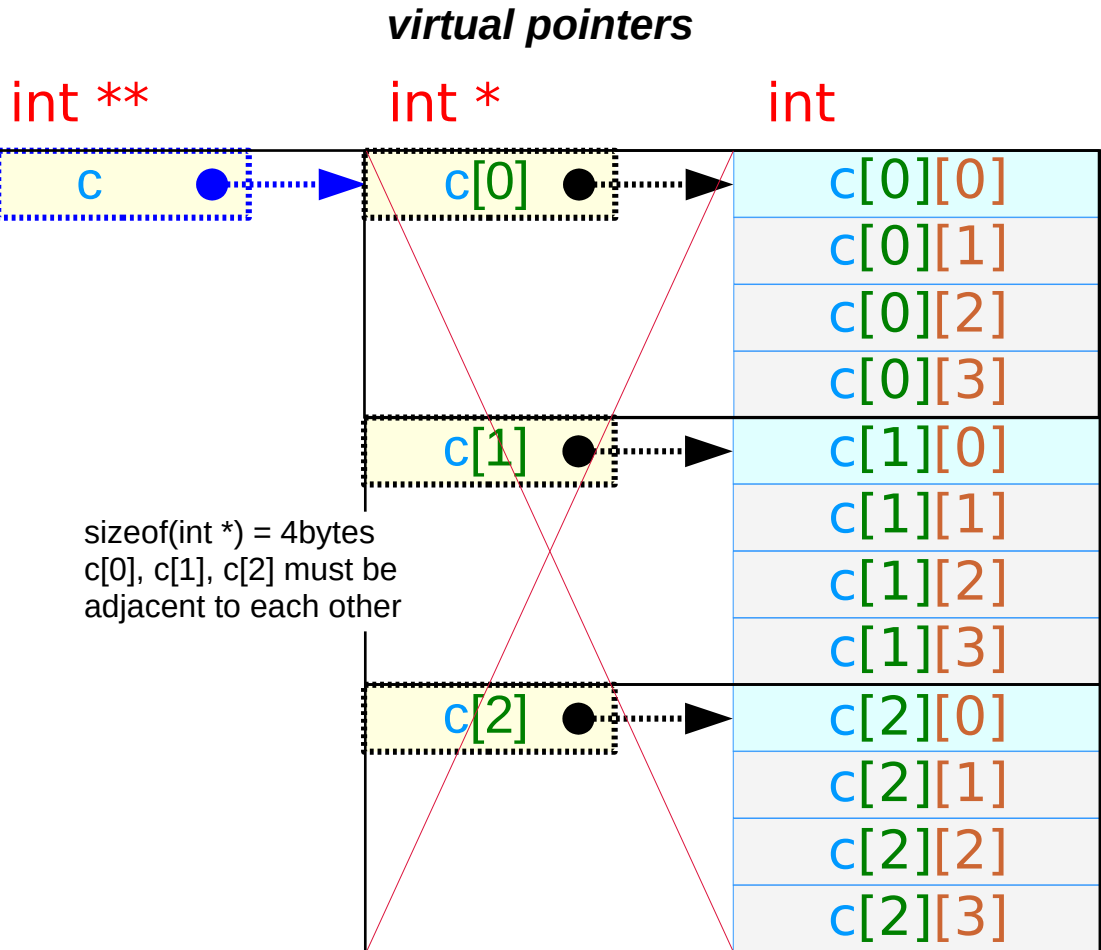
c double pointer
type : **int ****



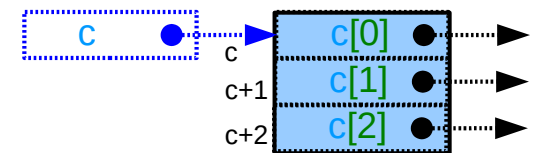
c[i] pointer
type : **int ***

Double pointer

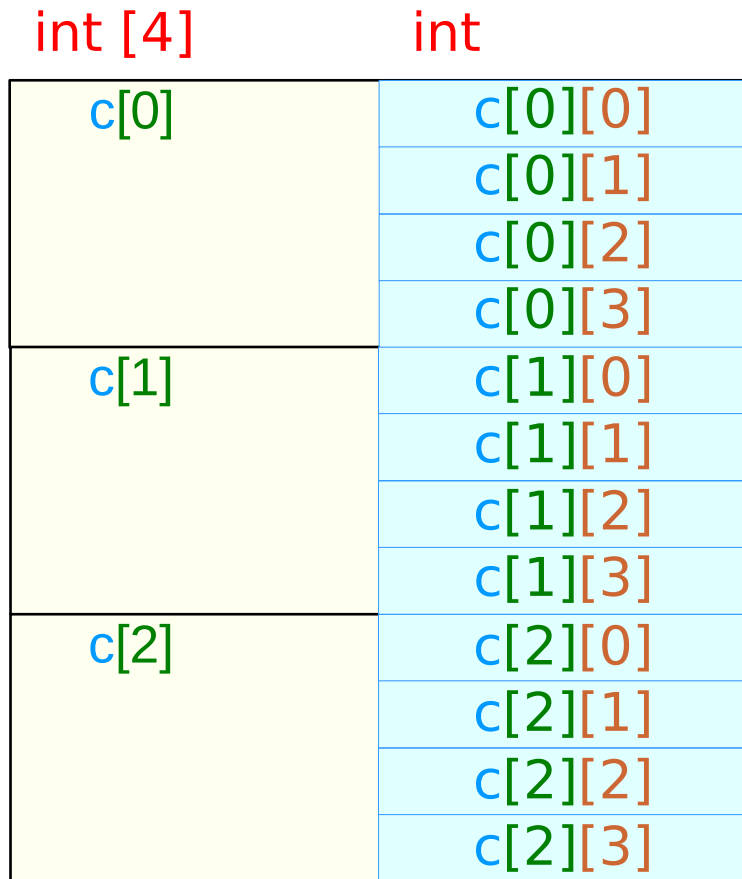
$*(*(c+i)+j)$



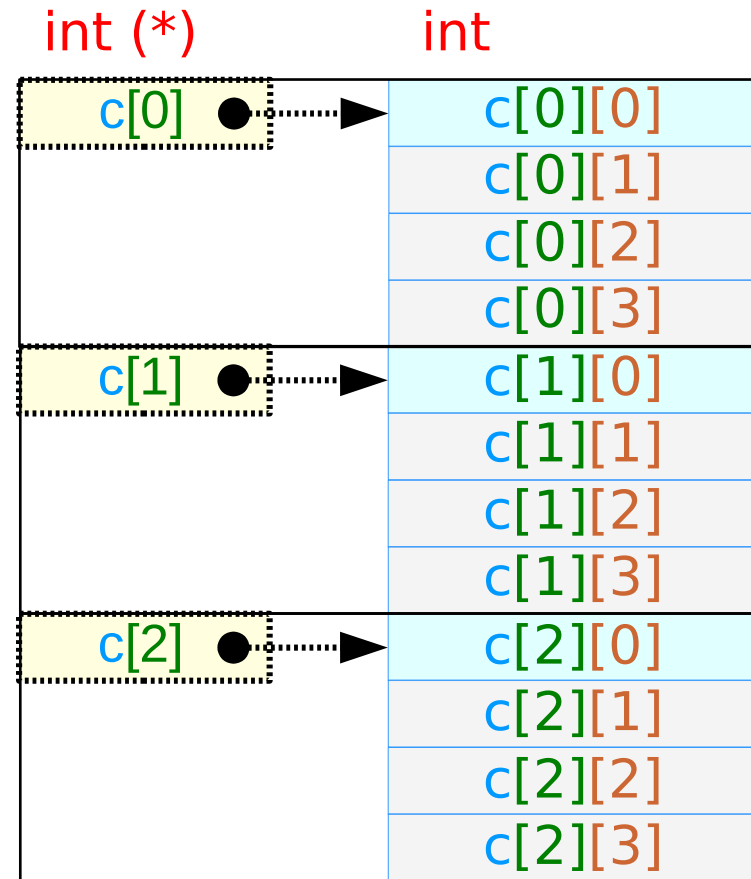
c points to an integer pointer



Case



virtual pointers



$*(*(\mathbf{c}+\mathbf{i})+\mathbf{j})$

c points to an integer pointer

Types in a 2-d array

int c [3] [4]

c 2-d array

type : int [3][4]

size : 3 * 4 * 4

value : &c[0][0]

relaxing the 1st dimension

int c [3] [4]

c 1-d array pointer (virtual)

type : int (*) [4]

size : 3 * 4 * 4

value : &c[0][0]

int c [3] [4]

c[i] 1-d array

type : int [4]

size : 4 * 4

value : &c[i][0]

relaxing the 1st dimension

int c [3] [4]

c[i] 0-d array pointer (virtual)

type : int (*)

size : 4 * 4

value : &c[i][0]

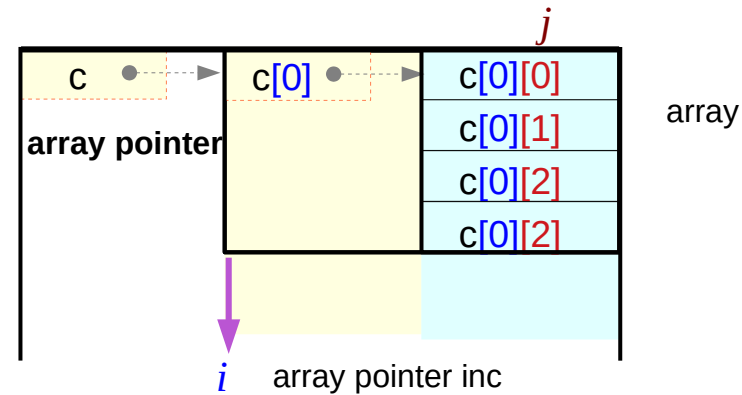
Abstract data and virtual pointer types

Case 2

$(*(c+i))[j]$

$int (*)[4]$

c points to a **1-d** array with 4 elements

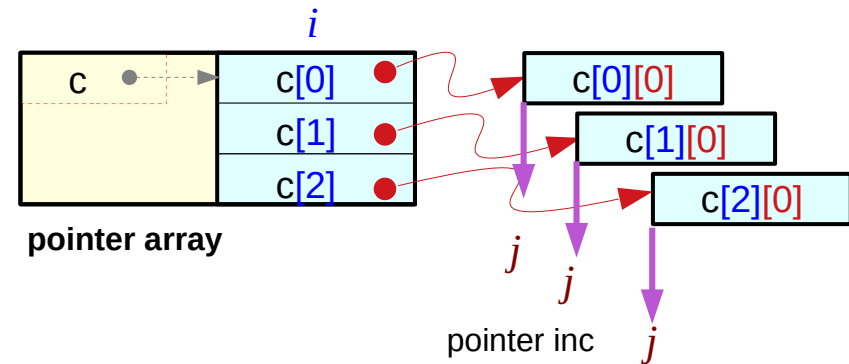


Case 3

$*(c[i]+j)$

$int * [3]$

c is a **1-d** array of integer pointers

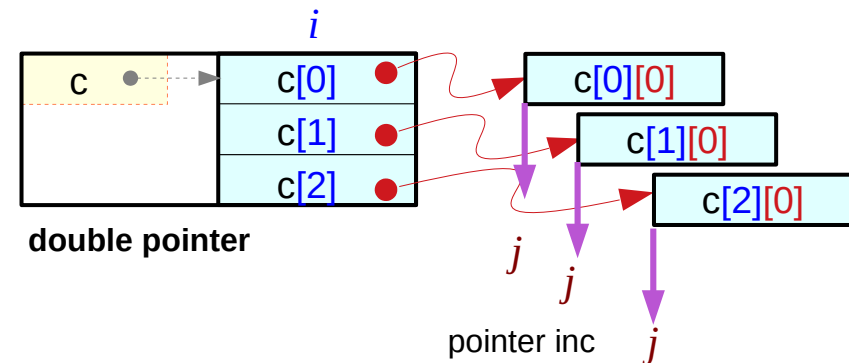


Case 4

$*(*(c+i)+j)$

$int **$

c points to an integer pointer



Case 1) 2-d array **c**, 1-d array **c[i]**

int **c** [3] [4]

c 2-d array

type : int [3][4]

int **c** [3] [4]

c[i] 1-d array

type : int [4]

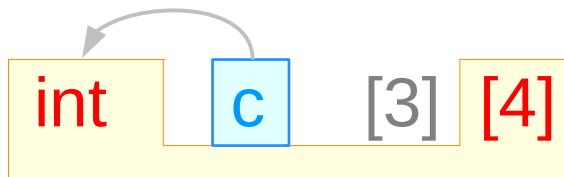
Abstract Data

(**c**[**i**]) [**j**]

int [3][4]	int [4]	int
c	c[0]	c[0][0]
		c[0][1]
		c[0][2]
		c[0][3]
	c[1]	c[1][0]
		c[1][1]
		c[1][2]
		c[1][3]
	c[2]	c[2][0]
		c[2][1]
		c[2][2]
		c[2][3]

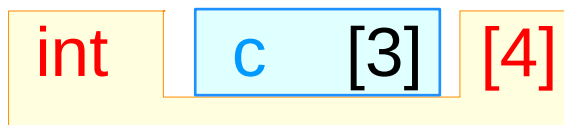
2-d array c

Case 2) 1-d array pointer **c**, 1-d array **c[i]**



c 1-d array pointer

type : `int (*) [4]`

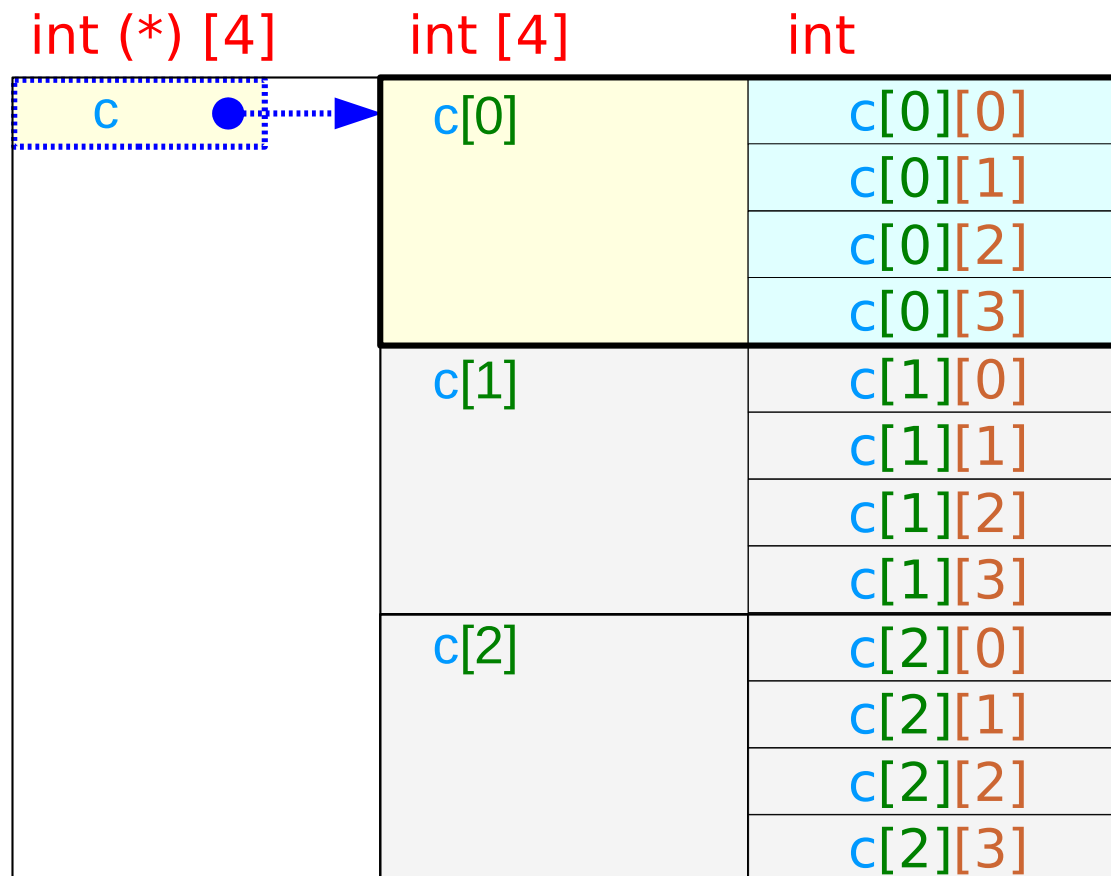


c[i] 1-d array

type : `int [4]`

Abstract Data

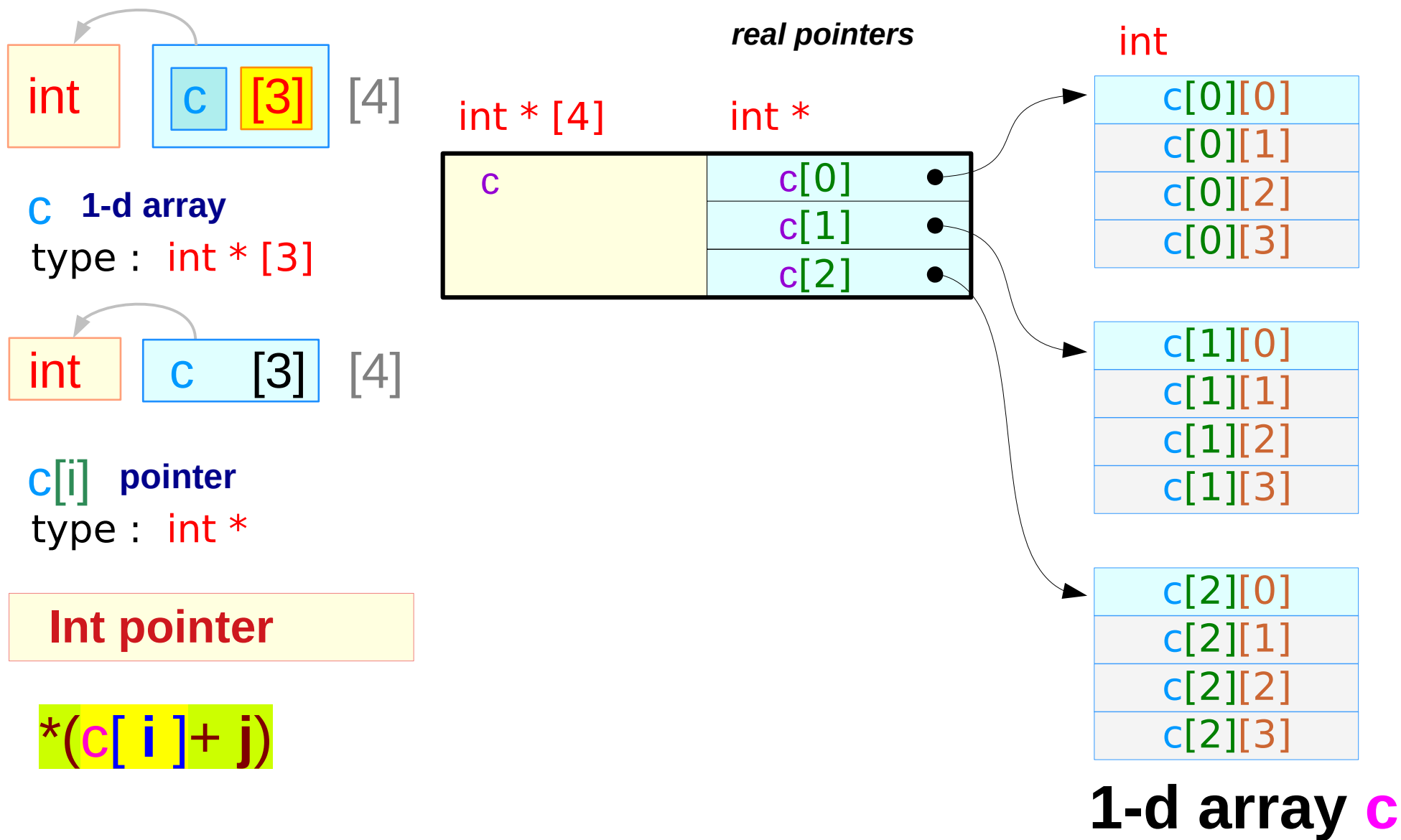
`(*(c+i)) [j]`



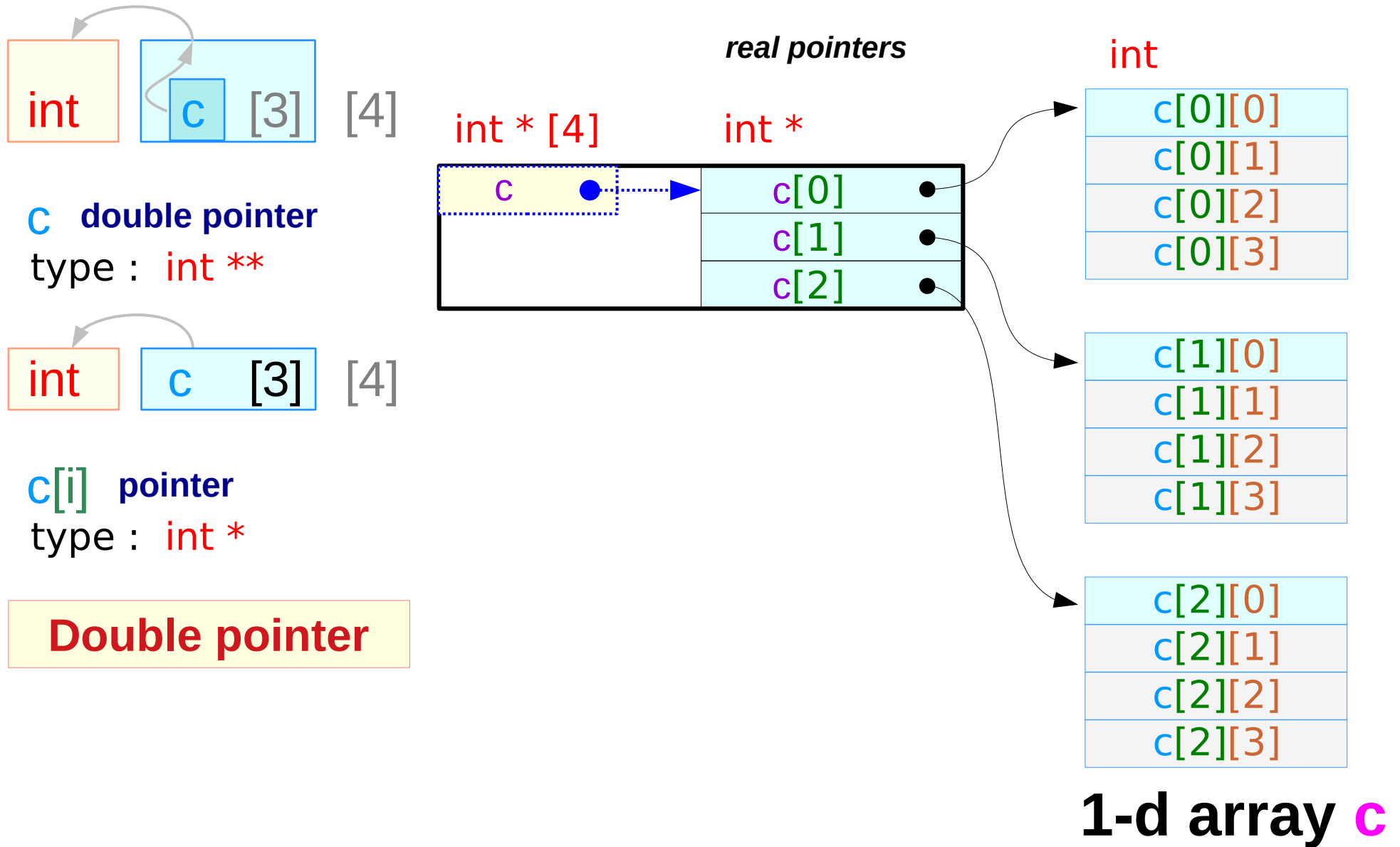
c points to an array
of 4 integers

2-d array **c**

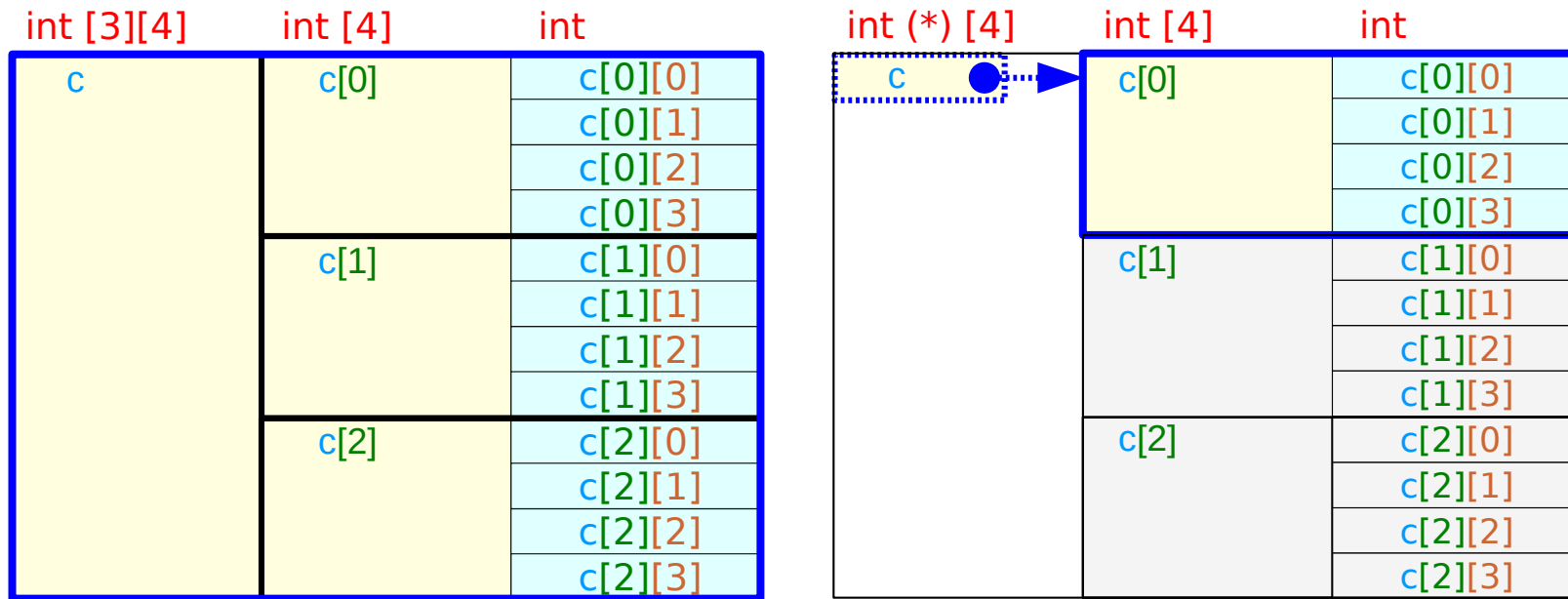
Case 3) 1-d array **c**, pointer **c[i]**



Case 4) double pointer **c**, pointer **c[i]**



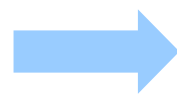
Nested Structure – abstract data, virtual pointer



Case 1

Case 2

`(c[i])[j]`

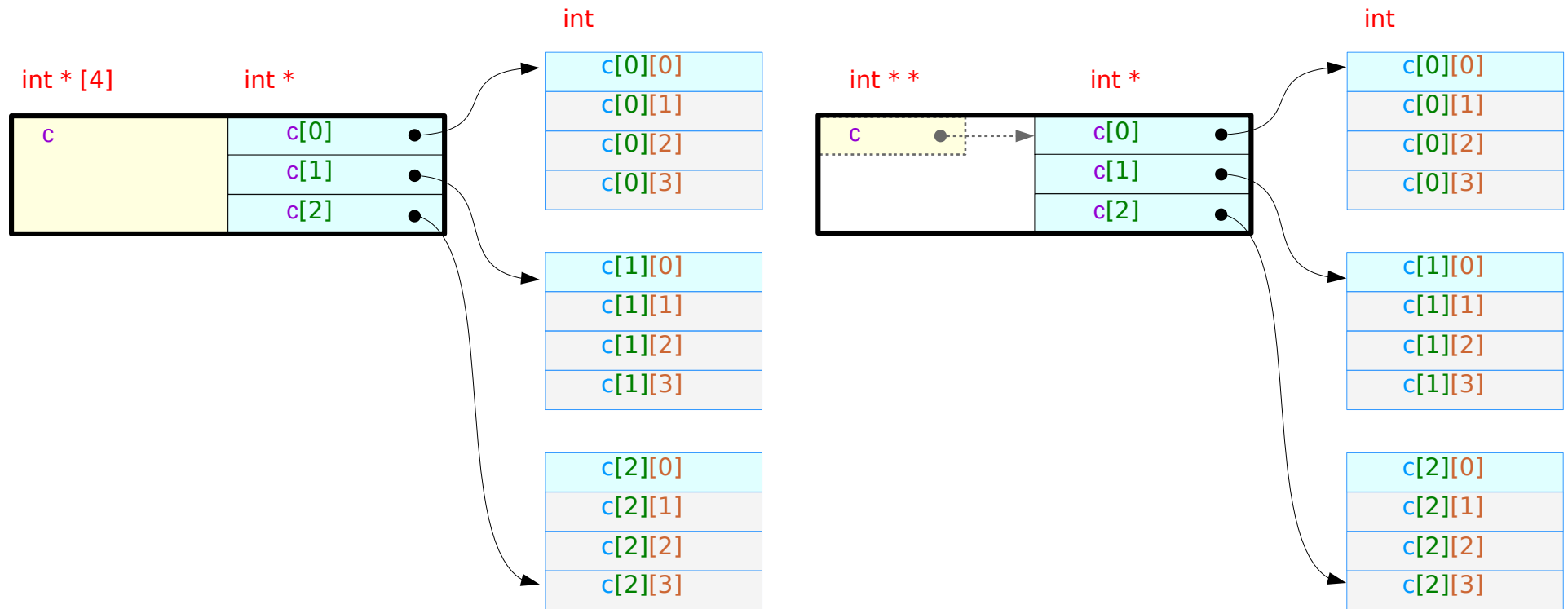


`(* (c+i))[j]`

nested structure

2-d array `c`

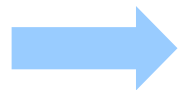
Nested structure – abstract data, virtual pointer



Case 3

Case 4

`*(c[i]+j)`



`*(*(c+i)+j)`

nested structure

1-d array `c`

-
- Pointer conversions in array types
 - **Simulating array accesses by real pointers**
 - Dynamic memory allocation

c is a double pointer and a **1-d** array pointer

$*(*(\mathbf{c}+\mathbf{0})+\mathbf{0})$



****c**

a double pointer

$(*(\mathbf{c}+\mathbf{0}))[\mathbf{0}]$



(*c)[0]

a **1-d** array pointer

c is a double pointer and a 1-d array pointer

Case 1

int [3][4]



Incrementing the 1st dimension pointer

Case 2

array pointer

int (*)[4]



Incrementing the 2nd dimension pointer

Case 4

double pointer

int **

Case 1

int [3][4]



Incrementing the 2nd dimension pointer

Case 3

pointer array

int *[3]



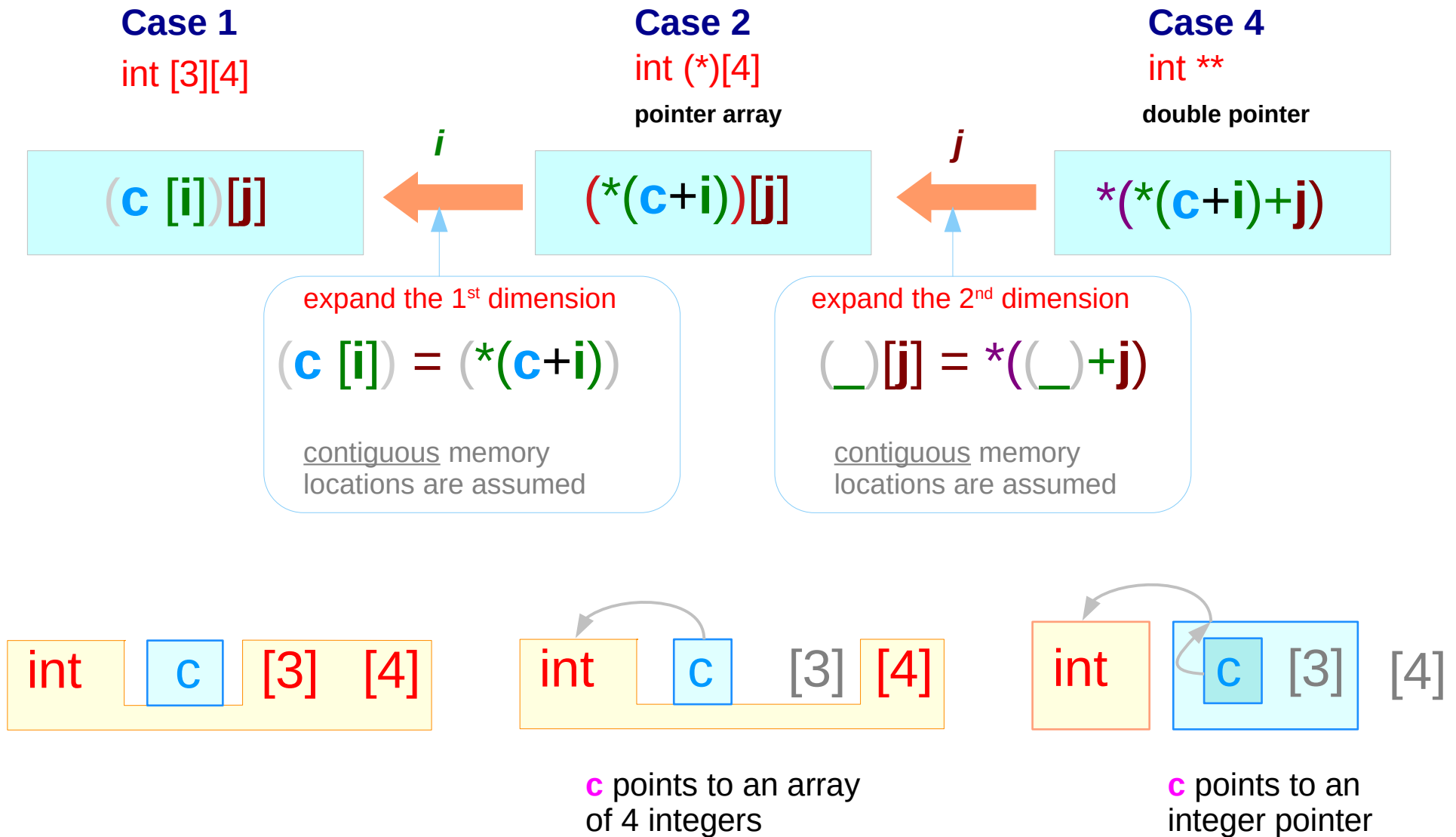
Incrementing the 1st dimension pointer

Case 4

double pointer

int **

2-d array access via a double indirection



2-d array access via a double indirection

Case 1

`int [3][4]`

`(c [i])[j]`

expand the 2nd dimension

`(_) [j] = *((_) + j)`

contiguous memory locations are assumed

Case 3

`int * [3]`

array pointer

`*((c [i]) + j)`

expand the 1st dimension

`(c [i]) = (*(c + i))`

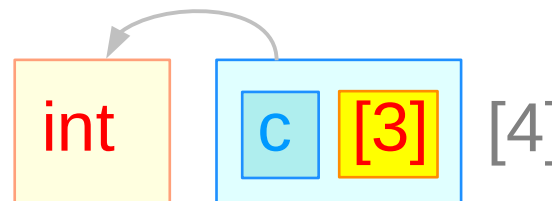
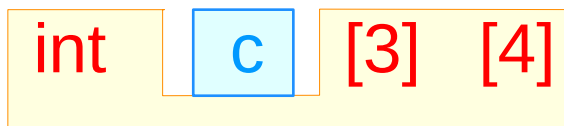
contiguous memory locations are assumed

Case 4

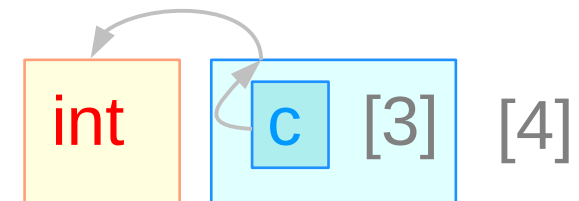
`int **`

double pointer

`*(*(c + i) + j)`



c is an array of 3 integer pointers



c points to an integer pointer

Cases 1, 2, 4

```
int c [3] [4];
```

```
int (*p) [4];
```

Case 1

int [3][4]

(c [i]) [j]

p = c

(p [i]) [j]

Case 2

int (*) [4]

(*(c+i)) [j]

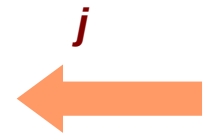
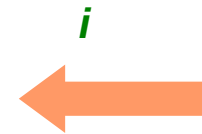
(*(p+i)) [j]

Case 4

int **

((c+i)+j)

((p+i)+j)



p[0]=c[0],
p[1]=c[1],
p[2]=c[2];

equivalence

Cases 1, 3, 4

```
int c [3] [4];
```

```
int **p, *q[3];
```

Case 1

int [3][4]

(c [i])[j]

p = q;

(p [i])[j]

q[0]=c[0],
q[1]=c[1],
q[2]=c[2];

must be allocated
and initialized

Case 3

int * [3]

*((c [i])+j)

*((p [i])+j)

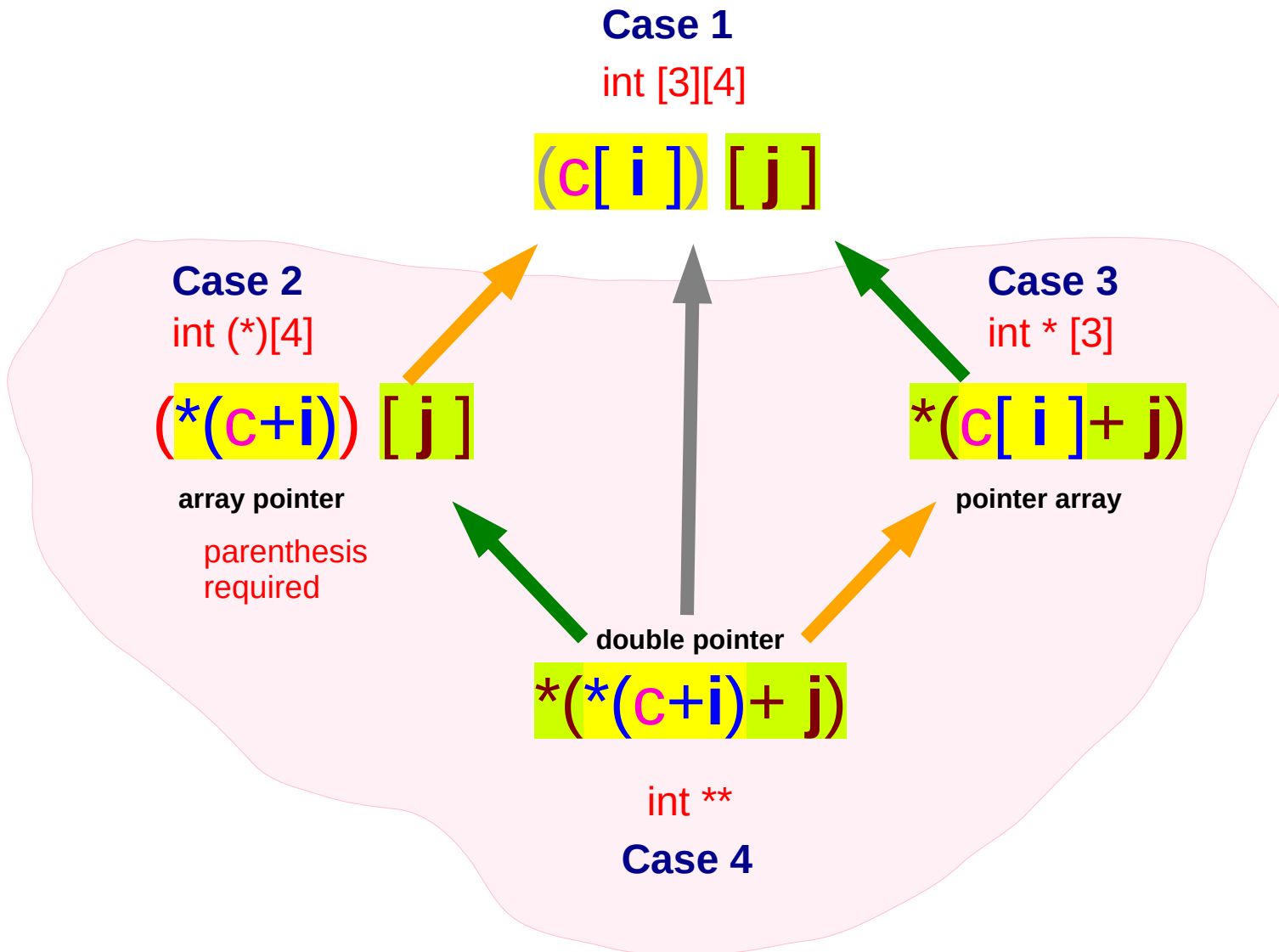
Case 4

int **

((c+i)+j)

((p+i)+j)

Simulating 2-d array accesses by real pointers



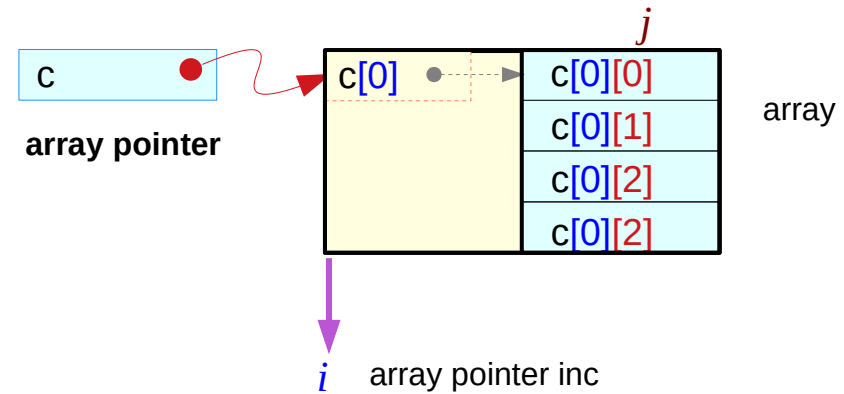
Incrementing pointers

Case 2

$(*(c+i))[j]$

$\text{int } (*)[4]$

c points to a **1-d** array with 4 elements

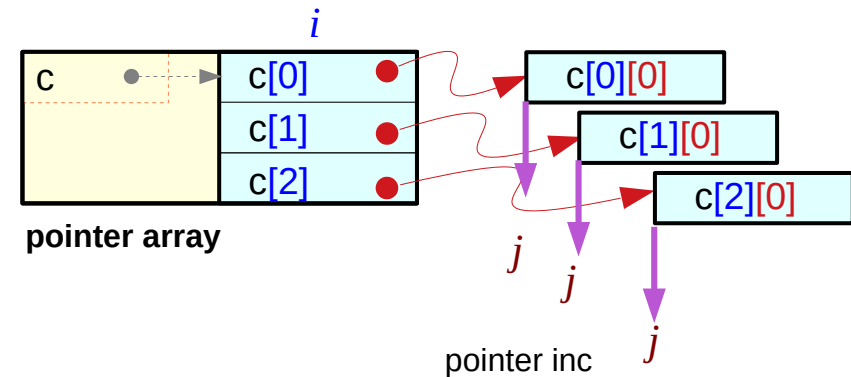


Case 3

$*(c[i]+j)$

$\text{int } * [3]$

c is a **1-d** array of integer pointers

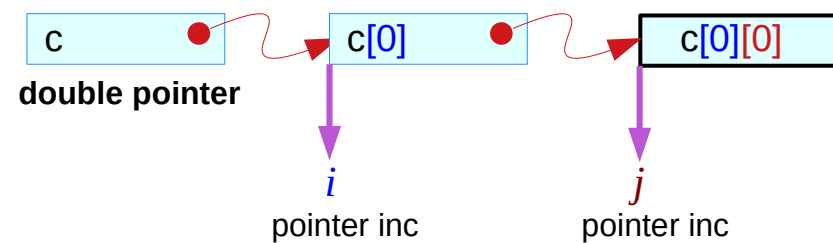


Case 4

$*(*(c+i)+j)$

$\text{int } **$

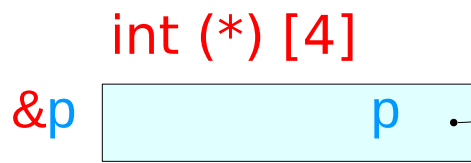
c points to an integer pointer



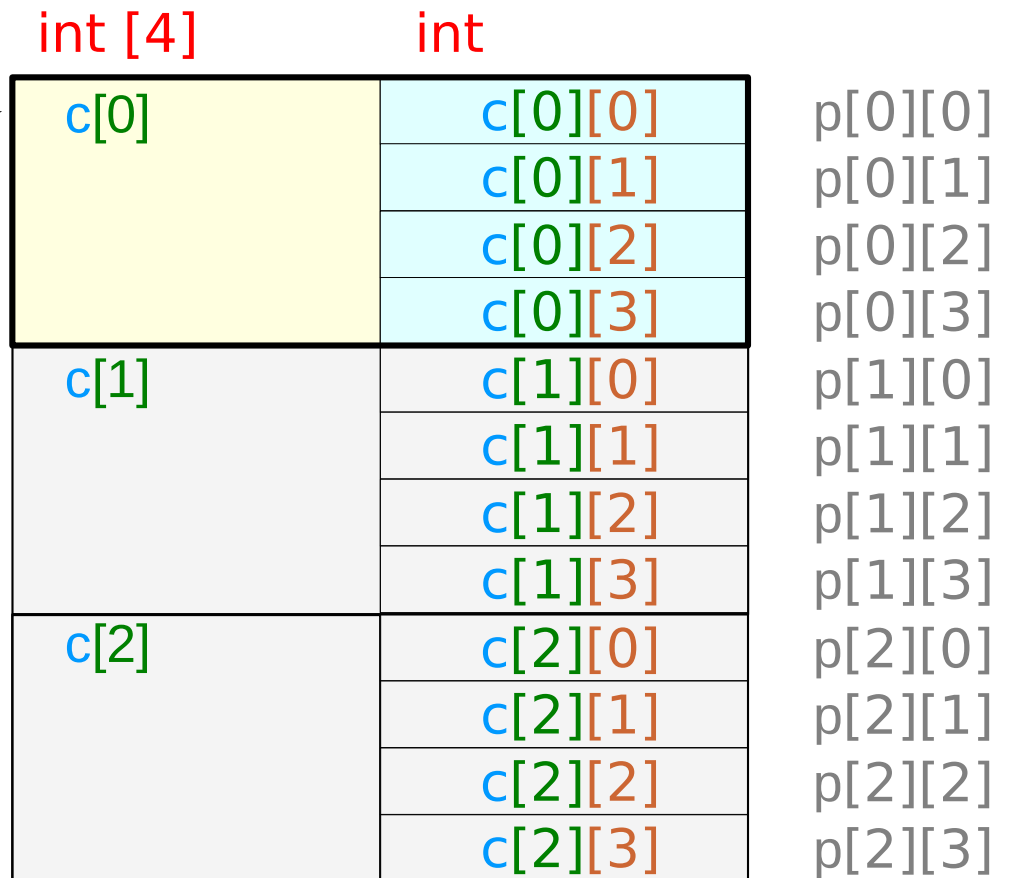
2-d array access using an array pointer **p**

```
int c [3] [4];
```

```
int (*p) [4];
```



```
p = &c[0];
```



Case 2

```
int (*)[4]
```

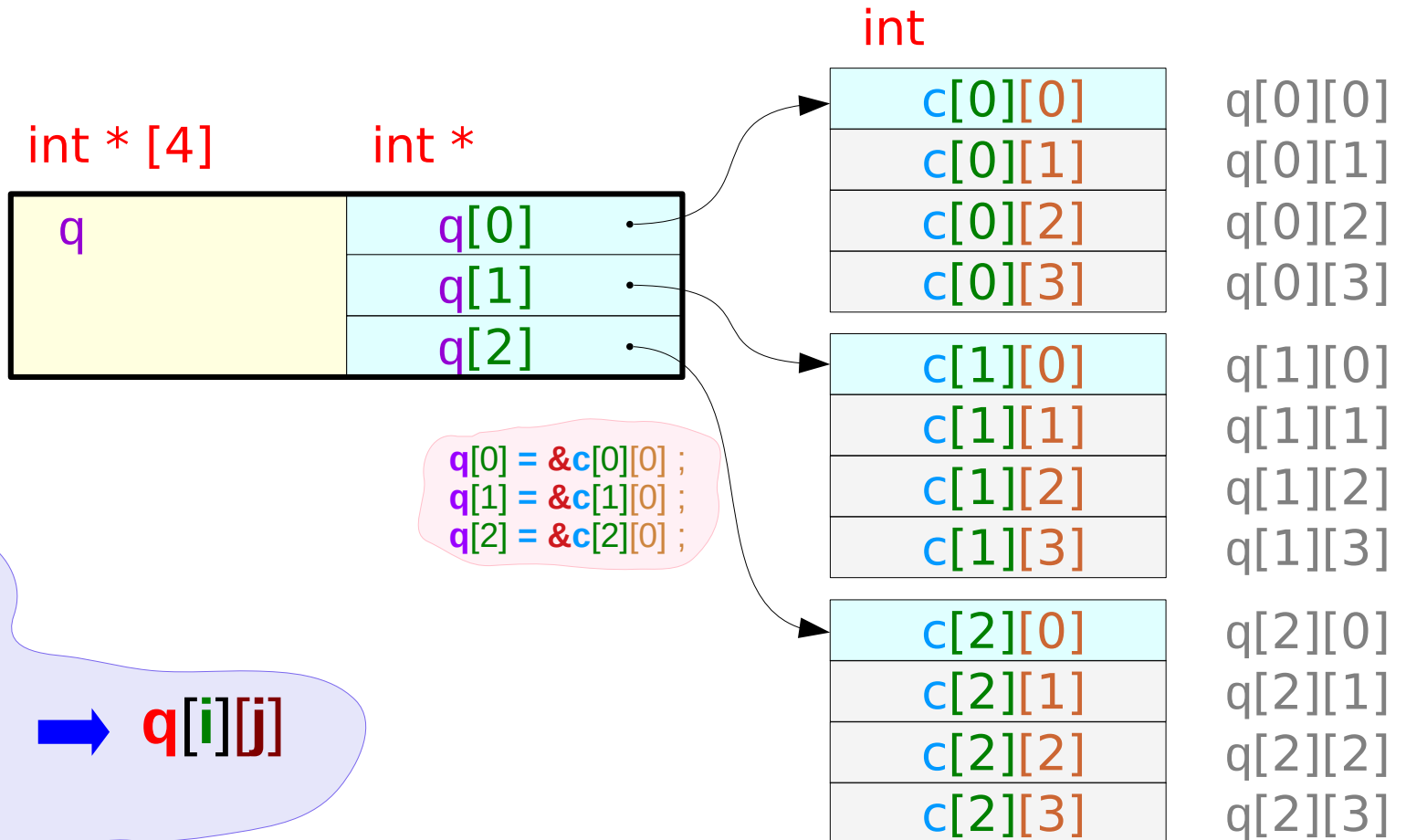
```
(*(c+i))[j]
```

```
(*(p+i))[j] → p[i][j]
```

2-d array access using a pointer array q

```
int c [3] [4];
```

```
int *q[3];
```



Case 3

```
int * [3]
```

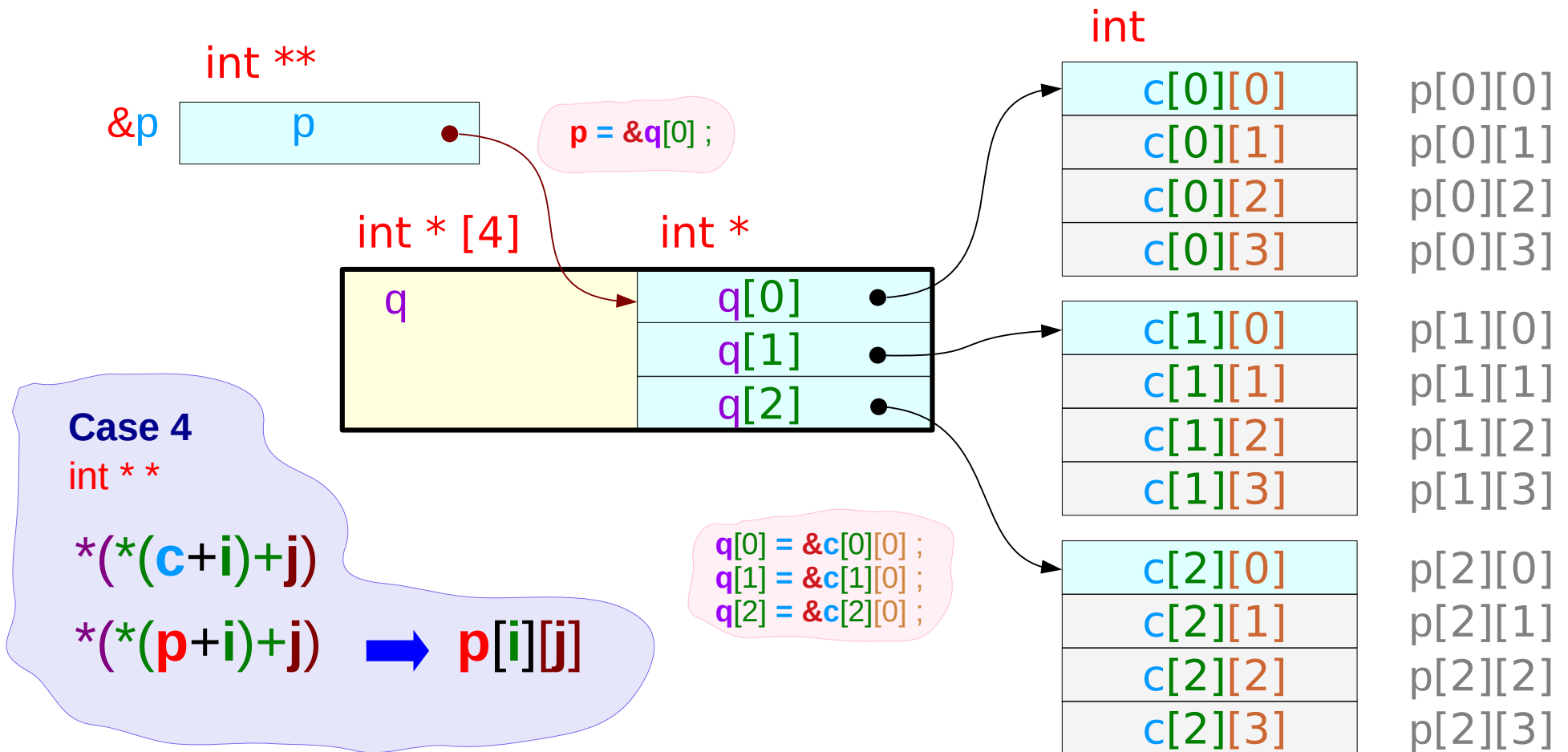
```
*((c [i]) + j)
```

```
*((q [i]) + j) → q[i][j]
```

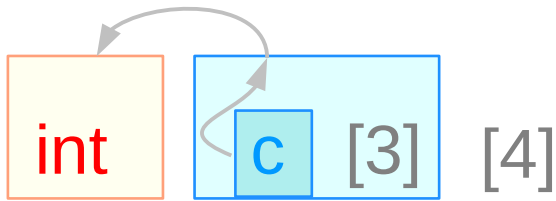

2-d array access using double pointers q

```
int c [3] [4];
```

```
int **p, *q[4];
```



Case 4) double pointer **c**, pointer **c[i]**



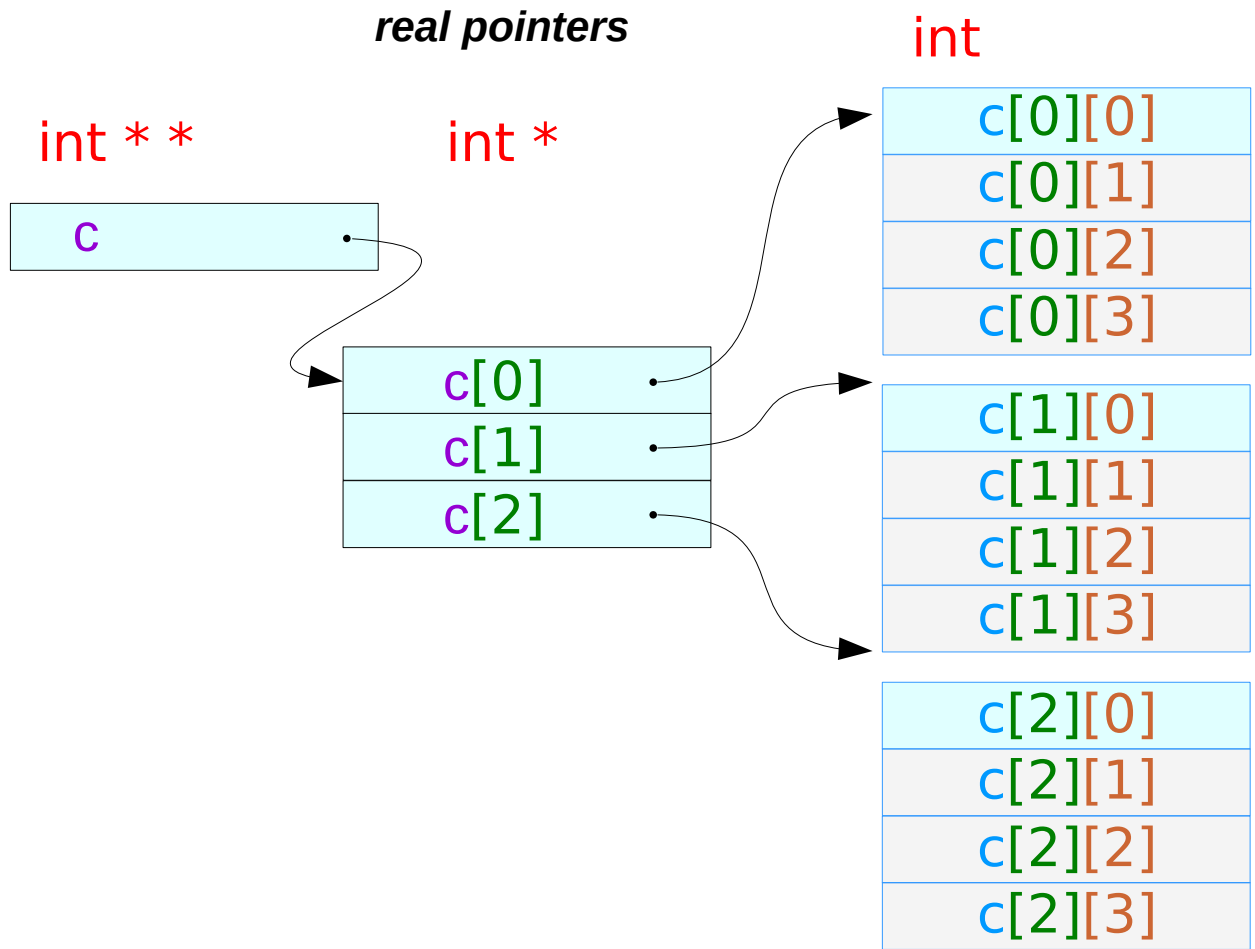
c double pointer
type : **int ****



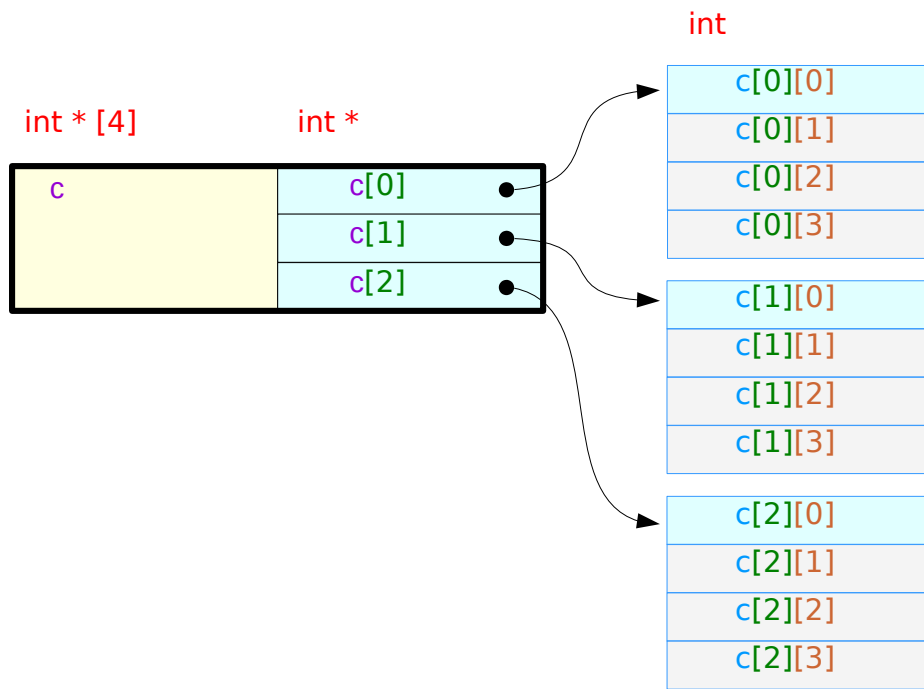
c[i] pointer
type : **int ***

Double pointer

$*(*(\mathbf{c} + \mathbf{i}) + \mathbf{j})$

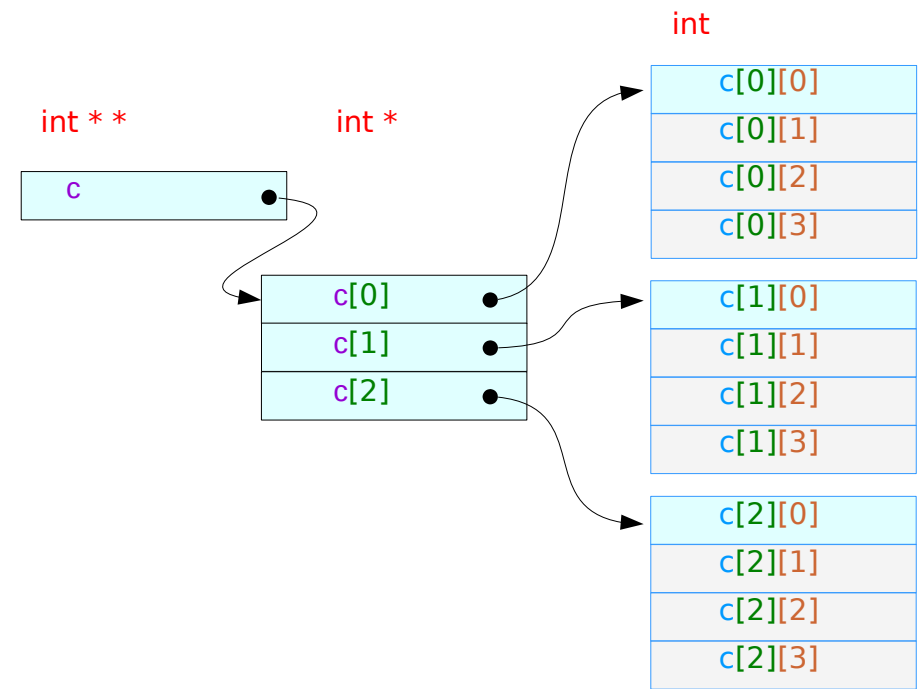


Nested structure - Pointers



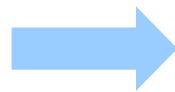
Case 3

`*(c[i]+j)`



Case 4

`*(*(c+i)+j)`



nested structure

-
- Pointer conversions in array types
 - Simulating array accesses by real pointers
 - **Dynamic memory allocation**

Dynamic Memory Allocation of 2-d Arrays

method 1

```
int ** c ;
c = (int **) malloc(3 * sizeof (int *) ) ;
c[0] = (int *) malloc(4 * sizeof (int)) ;
c[1] = (int *) malloc(4 * sizeof (int)) ;
c[2] = (int *) malloc(4 * sizeof (int)) ;
```

Case 4

```
int **
```

method 2

```
int ** c ;
int * p ;
c = (int **) malloc( 3 * sizeof(int *) ) ;
p = (int *) malloc( 4 * 4 * sizeof(int) ) ;
for (i=0; i<M; i++) c[i] = p + i*N;
```

Case 4

```
int **
```

method 3

```
int (*p) [3] ;
p = (int (*) [3]) malloc(3 * 4 * sizeof (int)) ;
```

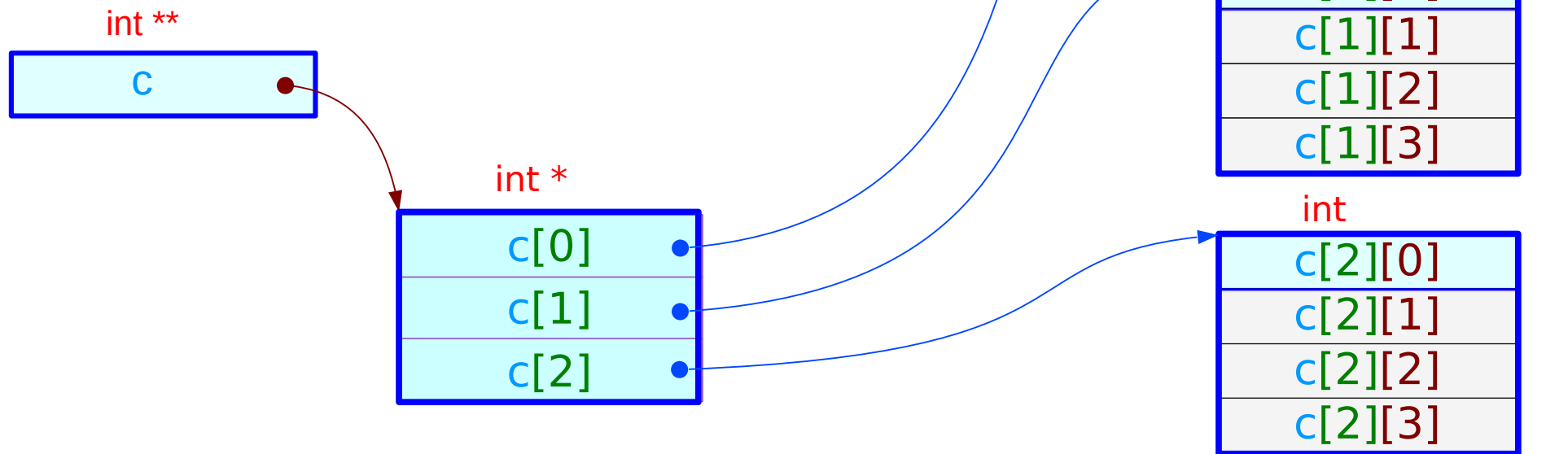
Case 2

```
int (*)[4]
```

2-d array dynamic allocation : method 1

method 1

```
int ** c ;  
c = (int **) malloc(3 * sizeof (int *) ) ;  
c[0] = (int *) malloc(4 * sizeof (int)) ;  
c[1] = (int *) malloc(4 * sizeof (int)) ;  
c[2] = (int *) malloc(4 * sizeof (int)) ;
```



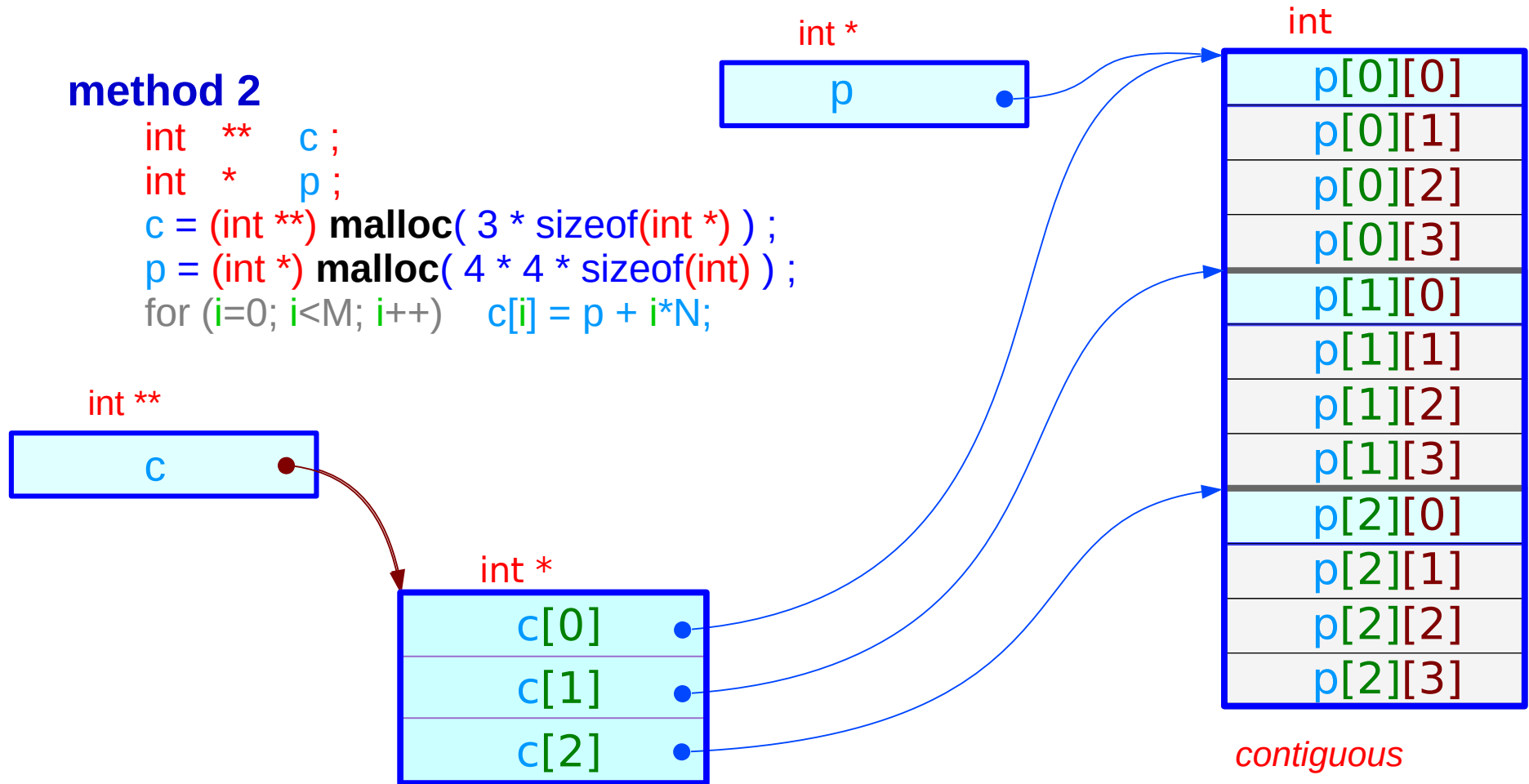
c: an array of integer pointers

may not be contiguous because of memory alignments

2-d array dynamic allocation : method 2

method 2

```
int ** c ;  
int * p ;  
c = (int **) malloc( 3 * sizeof(int *) ) ;  
p = (int *) malloc( 4 * 4 * sizeof(int) ) ;  
for (i=0; i<M; i++) c[i] = p + i*N;
```

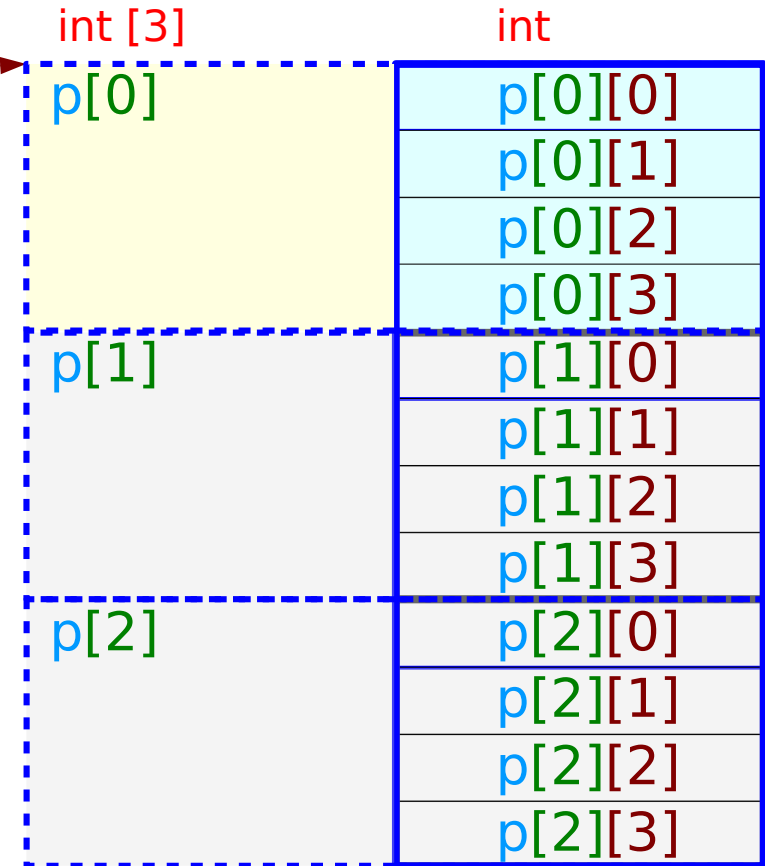
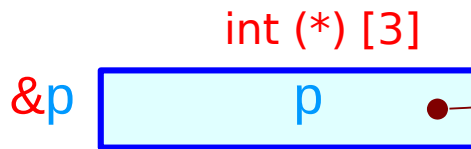


array of pointers
allocated physically in memory

2-d array dynamic allocation : method 3

method 3

```
int (*p) [3];  
p = (int (*) [3]) malloc(3 * 4 * sizeof (int));
```



utilize pointer
addition property

Pointer to Arrays :
No physical allocation

C scalar data types

Arithmetic types and **pointer** types

- collectively called **scalar types**
- hold single data item
- size and format

Array and **structure** / **union** types

- collectively called **aggregate types**
- hold more than one data items

C scalar data types provide their size and format
The alignment of a scalar data types is equal to its size

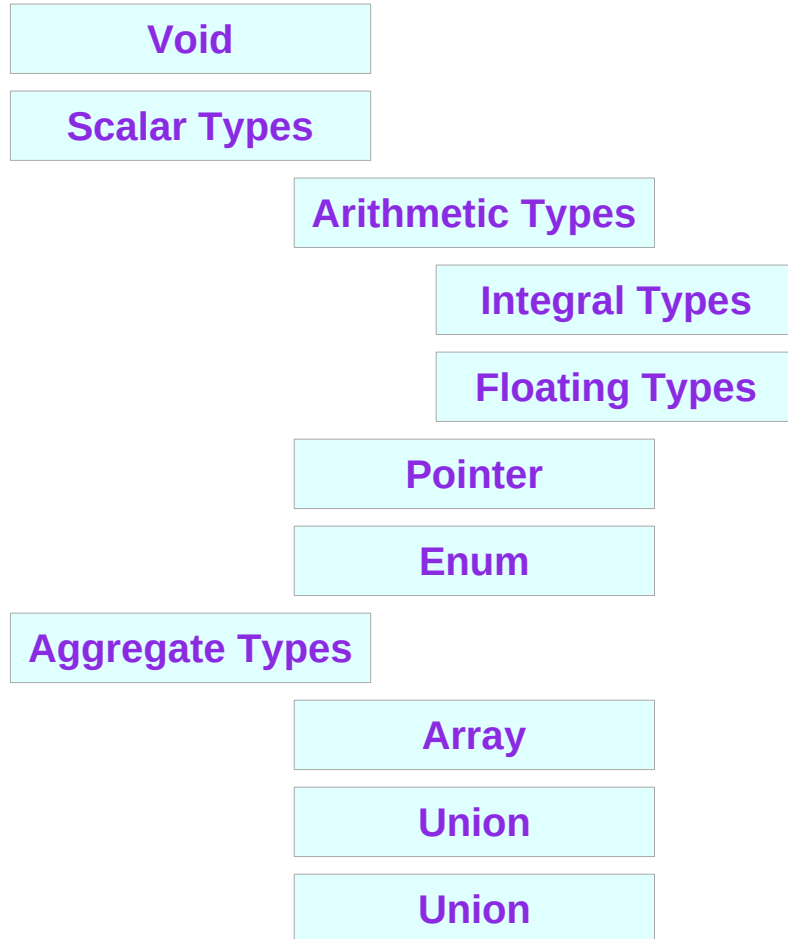
Scalar alignment shows scalar alignments that apply to individual scalars and to scalars that are elements of an array or members of a structure or union.

Wide characters are supported (character constants prefixed with an L)

The size of each wide character is 4 bytes

<https://stackoverflow.com/questions/35722514>

Data Types



<https://stackoverflow.com/questions/35722514>

C abstract data types

Array as an Abstract Data Type and as a Data Structure

Abstract Data Types, ADTs are a way of classifying data structures based on how they are used and the behaviors they provide.

They do not specify how the data structure must be implemented but simply provide a minimal expected interface and set of behaviors.

Data structure is a concrete implementation of a data type.

It's possible to analyze the time and memory complexity of a Data Structure but not from a data type.

The Data Structure can be implemented in several ways and its implementation may vary from language to language

[Lucasmagnum.edium.com/sidenotes-array-sbstract-data-type-data-structure...](https://lucasmagnum.edium.com/sidenotes-array-sbstract-data-type-data-structure...)

References

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- [3] C A Reference Manual, Samuel P. Harbison & Guy L. Steele Jr.
- [4] C Language Express, I. K. Chun