Applicatives Laws (2B)

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The definition of Applicative

class (Functor f) => Applicative f where pure :: a -> f a (<*>) :: f (a -> b) -> f a -> f b

The class has a two methods :

pure brings arbitrary values into the functor

(<*>) takes a <u>function</u> wrapped in a functor **f** and a <u>value</u> wrapped in a functor **f** and returns the <u>result</u> of the <u>application</u> which is also wrapped in a functor **f**

The Maybe instance of Applicative

instance Applicative Maybe where	
pure	= Just
(Just f) <*> (Just x)	= Just (f x)
_ <*> _	= Nothing

pure wraps the value with Just;

(<*>) applies

the <u>function</u> wrapped in **Just** to the <u>value</u> wrapped in **Just** if both exist, and results in **Nothing** otherwise.

https://en.wikibooks.org/wiki/Haskell/Applicative_functors

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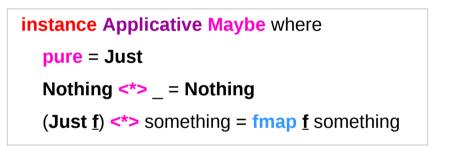
The Applicative Typeclass

class (Functor f) => Applicative f where pure :: a -> f a (<*>) :: f (a -> b) -> f a -> f b

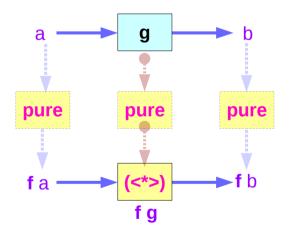


(Functor f) => Applicative f

f : Functor, Applicative



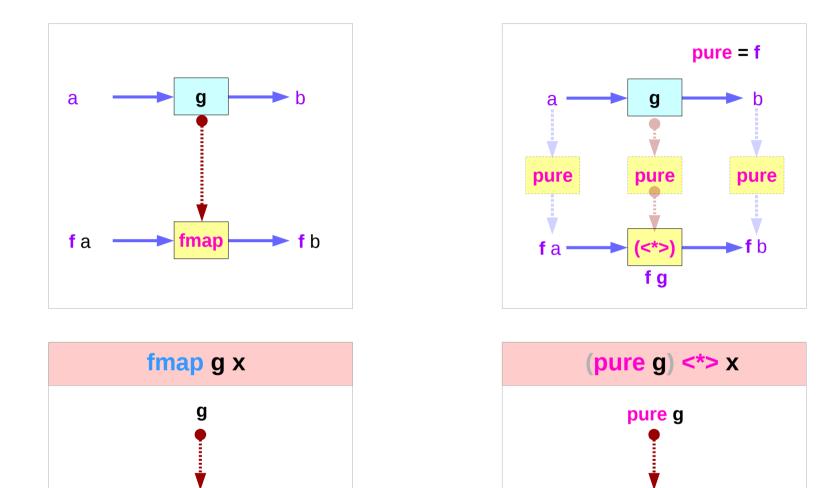
<u>f</u> : function in a context



(Functor f) => Applicative f

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fmap g x = (pure g) < > x



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fmap

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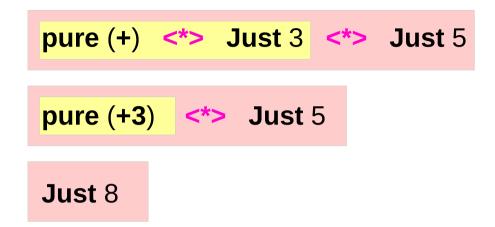
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Left Associative <*>

ghci> pure (+) <*> Just 3 <*> Just 5 Just 8



ghci> pure (+) <*> Just 3 <*> Nothing Nothing

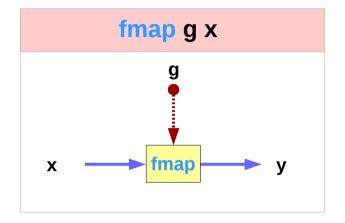
ghci> pure (+) <*> Nothing <*> Just 5 Nothing

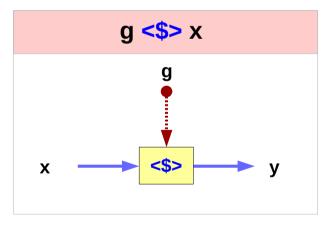
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Infix Operator <\$>





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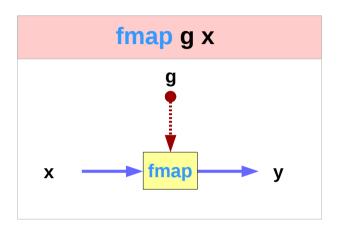
Infix Operator <\$> : not a class method

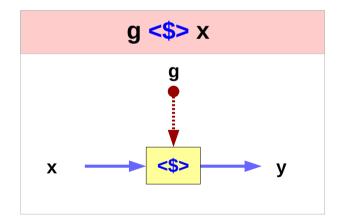
class (Functor f) => Applicative f where
 pure :: a -> f a
 (<*>) :: f (a -> b) -> f a -> f b

not a class method

(<\$>) :: (Functor f) => (a -> b) -> f a -> f b
f <\$> x = fmap f x

instance Applicative Maybe where		
pure = Just		
Nothing <*> _ = Nothing		
(Just f) <*> something = fmap f something		





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Applicatives Laws (2B)

The Applicative Typeclass

Applicative is a <u>superclass</u> of Monad. every Monad is also a Functor and an Applicative fmap, pure, (<*>) can all be used with monads.

a **Monad** instance also requires **Functor** and **Applicative** instances.

the types and roles of **return** and (>>)

(*> v.s. >>) and (pure v.s. return)

- (*>) :: Applicative f => fa -> fb -> fb
- (>>) :: Monad m => m a -> m b -> m b

pure :: Applicative f => a -> f a

return :: Monad m => a -> m a

the constraint changes from Applicative to Monad.

(*>) in Applicative	(>>) in Monad
pure in Applicative	return in Monad

The Applicative Laws

The identity law: pure id <*> v = v

Homomorphism: pure f <*> pure x = pure (f x)

Interchange: u <*> pure y = pure (\$ y) <*> u

Composition: u <*> (v <*> w) = pure (.) <*> u <*> v <*> w

The Identity Law

The identity law

pure id <*> v = v

pure to inject values into the functor

in a default, featureless way,

so that the result is as close as possible to the <u>plain</u> value.

applying the **pure id** morphism does nothing, exactly like with the plain **id** function.

The Homomorphism Law

The homomorphism law

pure f <*> pure x = pure (f x)

applying a "**pure**" <u>function</u> to a "**pure**" <u>value</u> is the same as applying the function to the <u>value</u> in the normal way and then using **pure** on the result. means **pure** preserves function application.

applying a <u>non-effectful</u> function f
to a <u>non-effectful</u> argument x in an <u>effectful</u> context pure
is the same as just applying the function f to the argument x
and then injecting the result (f x) into the <u>context</u> with pure.

The Interchange Law

The interchange law

u <*> pure y = pure (\$ y) <*> u

applying a morphism **u** to a "**pure**" value **pure y** is the same as applying **pure** (**\$ y**) to the morphism **u**

(**\$ y**) is the function that supplies **y** as <u>argument</u> to another function – the higher order functions

when evaluating the application of an <u>effectful function</u> **u** to a <u>pure argument</u> **pure y**, the <u>order</u> in which we evaluate the <u>function</u> **u** and its <u>argument</u> <u>pure</u> **y** <u>doesn't</u> <u>matter</u>.

The Composition Law

The composition law p	oure (.) <*> u <*> v <*> w = u <*> (v <*> w)
pure (.) composes morphisms similarly to how (.) composes functions:	(f.g) x = f (g x)
<pre>pure (.) <*> pure f <*> pure g <*> pure ;; = pure f <*> (pure g <*> pure x)</pre>	x u = pure f v = pure g w = pure x
applying the composed morphism pure (gives the same result as applying u to the result of applying v to w	.) <*> u <*> v to w u (v <*> w)

it is expressing a sort of associativity property of (<*>).

References

- [1] ftp://ftp.geoinfo.tuwien.ac.at/navratil/HaskellTutorial.pdf
- [2] https://www.umiacs.umd.edu/~hal/docs/daume02yaht.pdf