## Applicatives Methods (3B)

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## The definition of Applicative

```
class (Functor f) => Applicative f where
    pure :: a -> fa
    (<*>) :: f(a -> b) -> fa -> fb
```

$$
\begin{aligned}
& f(a->b):: \text { a function wrapped in } f \\
& f a:: \text { a value wrapped in } f
\end{aligned}
$$

The class has a two methods :
pure brings arbitrary values into the functor
(<*>) takes a function wrapped in a functor f and a value wrapped in a functor $f$
and returns the result of the application
which is also wrapped in a functor $f$

## The Maybe instance of Applicative

```
instance Applicative Maybe where
```

```
pure = Just
```

pure = Just
(Just f) <*> (Just x) = Just (f x)
(Just f) <*> (Just x) = Just (f x)
<*> _ = Nothing

```
    <*> _ = Nothing
```

pure wraps the value with Just;
(<*>) applies
the function wrapped in Just
to the value wrapped in Just if both exist, and results in Nothing otherwise.
https://en.wikibooks.org/wiki/Haskell/Applicative_functors

## An Instance of the Applicative Typeclass

```
class (Functor f) => Applicative f where
    pure :: a -> fa
    (<*>) :: f (a -> b) -> f a -> f b
```

f: Functor, Applicative
instance Applicative Maybe where
pure = Just

Nothing <*> $=$ Nothing
(Just $\underline{\mathbf{f}}$ ) <*> something $=$ fmap $\underline{\mathbf{f}}$ something
$\underline{f}$ : function in a context

(Functor f) => Applicative f

(Functor $f$ ) $=>$ Applicative $f$

## fmap $g x=($ pure $g)<*>x$


pure = f

http://learnyouahaskell.com/functors-applicative-functors-and-monoids

## Applicatives Methods (3B)

## Left Associative <*>

ghci> pure (+) <*> Just $3<*>$ Just 5 Just 8
ghci> pure (+) <*> Just 3 <*> Nothing
Nothing

## pure (+) <*> Just 3 <*> Just 5

## Just (+3) <*> Just 5

## Just 8

Nothing

## Infix Operator <\$>


http://learnyouahaskell.com/functors-applicative-functors-and-monoids

## Infix Operators <*> vs <\$>

```
h <*> x <*> y
h :: f (a -> b)
    Just (+) <*> Just 3 <*> Just 2
    Just (+3) <*> Just 2
    Just }
```


(+) <\$> Just 3 <*> Just 2
Just (+3) <*> Just 2
Just 5

$$
\mathbf{g}<\$>x \quad<*>y
$$

$$
\mathrm{g}:: \mathbf{a}->\mathbf{b}
$$

## Infix Operator <\$> : not in the minimal complete definition

```
class (Functor f) => Applicative f where
    pure :: a -> fa
    (<*>) :: f(a -> b) -> fa -> f b
```

Not in the minimal complete definition

```
(<\$>) :: (Functor f) => (a ->b) -> fa ->fb
\(\mathbf{g}\) <\$> \(\mathbf{x}=\) fmap \(\mathbf{g} \mathbf{x}\)
\[
\begin{aligned}
& \mathrm{g}:: \mathrm{a}->\mathrm{b} \\
& \mathrm{x}:: \mathrm{fa}
\end{aligned}
\]
```

```
instance Applicative Maybe where
    pure = Just
    Nothing <*> _ = Nothing
    (Just g) <*> something = fmap g something
```




## The Applicative Typeclass

Applicative is a superclass of Monad.
every Monad is also a Functor and an Applicative
fmap, pure, (<*>) can all be used with monads.
a Monad instance also requires
Functor and Applicative instances.
the types and roles of return and (>>)
https://en.wikibooks.org/wiki/Haskell/Applicative_functors

## The Applicative Laws

The identity law: pure id <*> v=v

Homomorphism: pure $\mathbf{f}$ <*> pure $\mathbf{x}=$ pure ( $\mathbf{f} \mathbf{x}$ )
f:: a -> b
$\mathrm{x}:$ : $\mathbf{a}$

Interchange: u <*> pure $\mathbf{y}=$ pure $(\$ \mathbf{y})<*>\mathbf{u}$
$u:: f(a->b) \quad y:: a$

Composition: u <*> (v <*> w) = pure (.) <*> u <*> v <*> w

## The Identity Law

## The identity law <br> pure id <*>v=v

pure to inject values into the functor
in a default, featureless way,
so that the result is as close as possible to the plain value.
applying the pure id morphism does nothing, exactly like with the plain id function.

## The Homomorphism Law

```
The homomorphism law
pure \(\mathbf{f}<\) *> pure \(\mathbf{x}=\) pure ( \(\mathbf{f} \mathbf{x}\) )
```

applying a "pure" function to a "pure" value is the same as applying the function to the value in the normal way and then using pure on the result. means pure preserves function application.
applying a non-effectful function $f$
to a non-effectful argument $\mathbf{x}$ in an effectful context pure is the same as just applying the function $f$ to the argument $\mathbf{x}$ and then injecting the result ( $\mathbf{f} \mathbf{x}$ ) into the context with pure.

## The Interchange Law

```
The interchange law
\(\mathbf{u}\) <*> pure \(\mathbf{y}=\) pure \((\$ \mathbf{y})\) <*> \(\mathbf{u}\)
```

applying a morphism $\mathbf{u}$ to a "pure" value pure $\mathbf{y}$ is the same as applying pure (\$y) to the morphism $\mathbf{u}$

```
Just (+3) <*> Just 2
Just ($ 2) <*> Just (+3)
```

(\$y) is the function that supplies $y$ as argument to another function

- the higher order functions
when evaluating the application of an effectful function $\mathbf{u}$ to a pure argument pure $\mathbf{y}$, the order in which we evaluate the function $\mathbf{u}$ and its argument pure $\mathbf{y}$ doesn't matter.


## The Composition Law

## The composition law

$$
\text { pure } \left.(.)<*>\mathbf{u}<{ }^{*}\right\rangle \text { v <*> w }=\mathbf{u}<*>(v<*>\text { w })
$$

pure (.) composes morphisms similarly to how (.) composes functions:
pure (.) <*> pure f <*> pure $\mathbf{g}$ <*> pure $x$
$=$ pure $\mathbf{f}<*>$ (pure $\mathbf{g}<*>$ pure $\boldsymbol{x}$ )
pure $h$
pure (.)
applying the composed morphism pure (.) <*> u <*> v to w
gives the same result as applying $\mathbf{u}$ to the result of applying $\mathbf{v}$ to $\mathbf{w}$
pure $\mathbf{g}$
$\square$

$$
(v<*>w)
$$

(g.h) $x=g(h x)$
it is expressing a sort of associativity property of (<*>).

## The Composition Law

The composition law

```
pure (.) <*> u <*> v <*> w = u <*> (v <*> w)
(g.h) x = g (h x)
w :: fa
v :: f(a -> b)
u : f(b -> c)
v <*> w :: f b
u <*> (v <*> w) :: f c
pure (.)<*> u <*> v :: f(a -> c )
pure (.) <*> u <*> v <*> w :: f c
```


## liftA2

## liftA2 :: (a -> b -> c) -> fa -> f b -> f c

Lift a binary function to actions.

Some functors support an implementation of liftA2 that is more efficient than the default one.

liftA2 may have an efficient implementation whereas fmap is an expensive operation, then better to use liftA2 than To use fmap over the structure and then use <*>.

$$
(\text { pure } \mathbf{g})<*>x<*>y
$$


http://hackage.haskell.org/package/base-4.10.1.0/docs/Control-Applicative.html\#v:liftA2
liftA2
variables

$$
\begin{aligned}
& \text { g :: a -> b -> c } \\
& x:=f a \\
& y \text { :: fb } \\
& \text { z: : f c } \\
& \text { pure } \mathbf{g} \text { <*> } \boldsymbol{x} \text { <*> } \boldsymbol{y} \\
& \text { liftA2 } \mathbf{g x y} \\
& \text { liftA2 :: (a -> b-> c) -> fa -> fb->fc } \\
& \text { g :: a -> b-> c } \\
& x: \text { fa } \\
& y:: f b \\
& \text { liftA2 gxy : f c }
\end{aligned}
$$

## liftA2

Actually, using the liftA commands
we can pull results of applicative functors
into a scope where we can talk
exclusively about functor results c
and not about effects. fec
Note that functor results can also be functions. c
This scope is simply a function,
which contains the code that we used in the non-functorial setting.

## liftA3

(lx gh h let $y=g x$ in $h y y)$
fx fg fh

The order of effects is entirely determined by the order of arguments to liftA3

## liftA2

liftA2 :: (a-> b-> c) -> fa->fb->fc

A minimal complete definition must include implementations of pure and of either <*> or liftA2.
pure and <*>
pure and liftA2

If it defines both, then they must behave the same as their default definitions:
(<*>) = liftA2 id liftA2
fxy=f<\$>>* y
http://hackage.haskell.org/package/base-4.10.1.0/docs/Control-Applicative.html\#v:liftA2
liftA2

10
down vote
accepted

The wiki article says that liftA2 (<*>)
can be used to compose applicative functors.
It's easy to see how to use it from its type:
$0::($ Applicative $f$, Applicative f1) $=>$

$$
f(f 1(a->b))->f(f 1 a)->f(f 1 b)
$$

$\left.0=\operatorname{liftA} 2\left(<^{*}\right\rangle\right)$

So to if $f$ is Maybe and $f 1$ is [] we get:
$>$ Just $[(+1),(+6)]$ `o` Just $[1,6]$
Just [2,7,7,12]
https://stackoverflow.com/questions/12587195/examples-of-haskell-applicative-transformers
liftA2

The other way around is:
> [Just (+1),Just (+6)] `o` [Just 1, Just 6] [Just 2,Just 7,Just 7,Just 12]
your ex function is equivalent to liftA2 (:):
test1 = liftA2 (:) "abc" ["pqr", "xyz"]
https://stackoverflow.com/questions/12587195/examples-of-haskell-applicative-transformers

## Applicatives

Methods (3B)

## liftA2

To use (:) with deeper applicative stack
you need multiple applications of liftA2:
*Main> (liftA2 . liftA2) (:) (Just "abc") (Just ["pqr", "xyz"])
Just ["apqr","axyz","bpqr","bxyz","cpqr","cxyz"]

However it only works when both operands are equally deep.
So besides double liftA2 you should use pure to fix the level:
*Main> (liftA2 . liftA2) (:) (pure "abc") (Just ["pqr", "xyz"])
Just ["apqr","axyz","bpqr","bxyz","cpqr","cxyz"]
liftA2

Consider the non-functorial expression:
x:: $x$
g :: x -> y
h :: y -> y -> z
let $\mathrm{y}=\mathrm{g} \mathrm{x}$
in hyy

Very simple. Now we like to generalize this to
fx : : f $x$
fg :: f(x -> y)
fh :: f(y -> y -> z)

## liftA2

However, we note that
let $\mathrm{fy}=\mathrm{fg}$ <*> fx
in fh <*> fy <*> fy
runs the effect of fy
twice. E.g. if fy
writes something to the terminal then fh <*> fy <*> fy
writes twice. This could be intended, but how can we achieve, that the effect is run only once and the result is used twice? Actually, using the liftA
commands we can pull results of applicative functors into a scope where we can talk exclusively about functor results and not about effects. Note that functor results can also be functions. This scope is simply a function, which contains the code that we used in the non-functorial setting.

## liftA3

(lx gh h let $\mathrm{y}=\mathrm{gx} \mathrm{x} \mathrm{n} \mathrm{h} y \mathrm{y}$ )
fx fg fh

The order of effects is entirely determined by the order of arguments to liftA3
https://wiki.haskell.org/Applicative_functor

## <\$> related operators

Functor map <\$>
(<\$>) :: Functor f => (a -> b) -> fa -> f b
(<\$) :: Functor f => a -> fb -> fa
(\$>) :: Functor f => fa->b -> f b

The <\$> operator is just a synonym
for the fmap function from the Functor typeclass.

This function generalizes the map function for lists to many other data types, such as Maybe, IO, and Map.

## <\$> examples

\#!/usr/bin/env stack
-- stack --resolver ghc-7.10.3 runghc
import Data.Monoid ((<>))
main :: 10 ()
main $=$ do
putStrLn "Enter your year of birth"
year <- read <\$> getLine
let age :: Int
age $=2020$ - year
putStrLn \$ "Age in 2020: " <> show age
https://haskell-lang.org/tutorial/operators

## <\$, \$> operators

In addition, there are two additional operators provided which replace a value inside a Functor
instead of applying a function.
This can be both more convenient in some cases, as well as for some Functors be more efficient.
value $<\$$ functor $=$ const value $<\$>$ functor
functor $\$>$ value $=$ const value $<\$>$ functor
$x<\$ y=y \$>x$
$x$ \$>y=y<\$x

## <*> related operators

Applicative function application <*>
(<*>) :: Applicative $f=>f(a->b)->f a->f b$
(*>) :: Applicative $f=>f a->f b->f b$
(<*) :: Applicative f => fa->fb->fa

Commonly seen with <\$>, <*> is an operator
that applies a wrapped function to a wrapped value.
It is part of the Applicative typeclass,
and is very often seen in code like the following:
foo <\$> bar <*> baz
https://haskell-lang.org/tutorial/operators

## <*> examples

For cases when you're dealing with a Monad, this is equivalent to:

```
do x <- bar
    y <- baz
    return (foo x y)
```

Other common examples including parsers and serialization libraries.
Here's an example you might see using the aeson package:
data Person = Person \{ name :: Text, age :: Int \} deriving Show
-- We expect a JSON object, so we fail at any non-Object value.
instance FromJSON Person where parseJSON (Object v) = Person <\$> v .: "name" <*> v .: "age" parseJSON _ = empty

## *> operator

To go along with this, we have two helper operators that are less frequently used:
*> ignores the value from the first argument. It can be defined as:

```
a1 *> a2 = (id <$ a1) <*> a2
```

Or in do-notation:

```
a1 *> a2 = do
```

_<- a1
a2

For Monads, this is completely equivalent to >>.

## <* operator

<* is the same thing in reverse: perform the first action then the second, but only take the value from the first action.

Again, definitions in terms of <*> and do-notation:
(<*) $=$ liftA2 const
a1 $<^{*}$ a2 = do
res <- a1
<- a2
return res

## (*> v.s. >>) and (pure v.s. return)

(*>) :: Applicative $f=>\quad$ fa $->$ fb $->\quad$ fb
$(\gg):$ Monad $m \quad=>\quad$ ma-> mb-> mb
pure :: Applicative f $=>\quad$ a $->\quad$ fa
return :: Monad m => a -> ma
the constraint changes from Applicative to Monad.
(*>) in Applicative pure in Applicative
( $\gg$ ) in Monad
return in Monad

## References

[1] ftp://ftp.geoinfo.tuwien.ac.at/navratil/HaskellTutorial.pdf
[2] https://www.umiacs.umd.edu/~hal/docs/daume02yaht.pdf

