Background – Type Classes (1B)

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Based on

http://learnyouahaskell.com/making-our-own-types-and-typeclasses#the-functor-typeclass

http://learnyouahaskell.com/functors-applicative-functors-and-monoids

Haskell in 5 steps https://wiki.haskell.org/Haskell_in_5_steps

Polymorphism in Haskell

The **polymorphism** features of Haskell

- Purity
- higher order functions
- parameterized algebraic data types
- typeclasses

- -- side effects
- -- function passing and returning

Typeclasses

Types in Haskell

- no explicit <u>hierarchy</u> of <u>types</u>
- similar types can act like each other
- connect such similar types with the appropriate typeclasses

Example:

An Int can act like many things

• like an <u>equatable</u> thing, **Eq**

• like an <u>ordered</u> thing, **Ord**

• like an <u>enumerable</u> thing, etc. **Enum**

Open Typeclasses

Typeclasses are **open**:

• can define our own data type,

can think about what it can act like

 can connect it with the typeclasses that define its behaviors.

typeclasses define <u>behaviors</u> that is very <u>general</u> and <u>abstract</u>

action

behavior

define behaviors operation of a functions

Defining behavior

defining behaviors:

define behaviors

the type declarations of functions

operation of a functions

general and abstract:

A **typeclass defintion** include

the type declarations of functions,

define behavior

which give a lot of informations

connect

about functions

Examples of defining behavior

Example:

typeclasses that define operations

to see if two things are <u>equal</u> to <u>compare</u> two things by some <u>ordering</u>.

- very abstract and elegant <u>behaviors</u>,
- not anything very special
 because these operations are most common

Typeclasses and Instances

typeclasses are like interfaces

defines some behavior

- comparing for *equality*
- comparing for ordering
- enumeration

instances of that typeclass

types possessing such behavior

such behavior is defined by

- function definition
- function **type declaration** only

a function definition

a function type declaration

A function definition can be overloaded

Typeclasses and Type

typeclasses are like interfaces

defines some behavior

- comparing for equality
- comparing for ordering
- enumeration

instances of that typeclass

types possessing such behavior

class AAA bbb where

func1 :: a -> b -> c

func2 :: b -> c -> a

instance **AAA BBB** where

func1 definition func2 definition

a type is an instance of a typeclass implies

the function types <u>declared</u> by the **typeclass** are <u>defined</u> (implemented) in the **instance** so that the functions can be used, which the **typeclass** defines with that **type**

Instance Example

the Eq typeclass

defines the functions == and /=

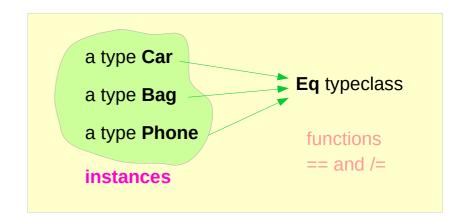
a type Car

comparing two cars c1 and c2 with the equality function ==

The Car type is an **instance** of Eq **typeclass**

Instances: various types

Typeclass: a group or a class of these similar types



Instance of a typeclass (1)

(State a) is an instance of Show a should be an instance of Show

```
State { runState = (\c -> (c, c)) }
```

https://stackoverflow.com/questions/7966956/instance-show-state-where-doesnt-compile

Instance of a typeclass (2)

```
*Main> getState
["0 => (0,0)","1 => (1,1)","2 => (2,2)","3 => (3,3)"]

*Main> putState 1
["0 => ((),1)","1 => ((),1)","2 => ((),1)","3 => ((),1)"]
```

https://stackoverflow.com/questions/7966956/instance-show-state-where-doesnt-compile

TrafficLight Type Example (2)

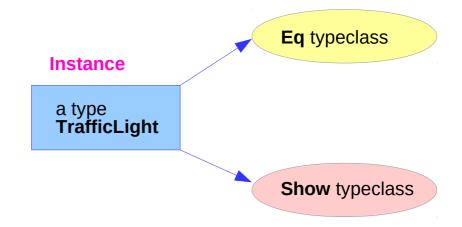
```
class Show a where
show :: a -> String
* * *
```

- a type declaration

data TrafficLight = Red | Yellow | Green

instance Show TrafficLight where

show Red = "Red light"
show Yellow = "Yellow light"
show Green = "Green light"



ghci> [Red, Yellow, Green]
[Red light, Yellow light, Green light]

Class Constraints

```
class (Eq a) => Num a where
...

class Num a where
...
```

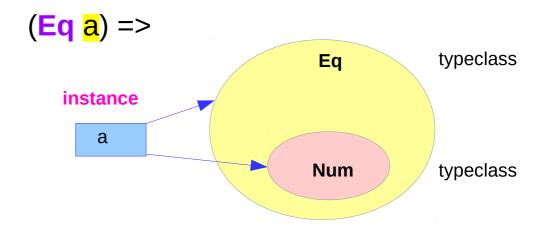
class constraint on a class declaration

an instance of **Eq**<u>before</u> being an instance of **Num**

the required function bodies can be defined in

- · the class declaration
- an instance declarations,

we can safely use == because a is a part of Eq



Num: a subclass of Eq

Class Constraints: class & instance declarations

class constraints in class declarations

to make a typeclass a **<u>subclass</u>** of another typeclass

subclass

```
class (Eq a) => Num a where ...
```

class constraints in instance declarations

to express <u>requirements</u> about the contents of some type.

<u>requirements</u>

```
instance (Eq x, Eq y) => Eq (Pair x y) where
Pair x0 y0 == Pair x1 y1 = x0 == x1 && y0 == y1
```

http://cmsc-16100.cs.uchicago.edu/2016/Lectures/07-type-classes.php

Class constraints in instance declaration examples

```
instance (Eq x, Eq y) => Eq (Pair x y) where

Pair x0 y0 == Pair x1 y1 = x0 == x1 & y0 == y1

Eq (Pair x y) Eq x Eq y
```

Derived instance

Class constraints and Overloading

```
class Eq a where
(==) :: a -> a -> Bool
```

```
instance Eq Integer where
x == y = x integerEq y
```

```
instance Eq Float where
x == y = x `floatEq` y
```

```
== of Eq (Tree a) == of Eq a
```

A Concrete Type and a Type Constructor

a : a concrete type

Maybe : not a concrete type

: a type constructor that takes one parameter

produces a concrete type.

Maybe a : a concrete type

Instance of **Eq**

```
data TrafficLight = Red | Yellow | Green
```

```
class Eq a where

(==) :: a -> a -> Bool

(/=) :: a -> a -> Bool

x == y = not (x /= y)

x /= y = not (x == y)

instance Eq TrafficLight where

Red == Red = True

Green == Green = True

Yellow == Yellow = True

_ == _ = False
```

to define our own type (defining a new data type) allowed values are Red, Yellow, and Green no class (type) instances

class:

defining new typeclasses

instance

making types instances of a typeclasses

http://learnyouahaskell.com/making-our-own-types-and-typeclasses

Instance of **Show**

instance Show TrafficLight where

show Red = "Red light"

show Yellow = "Yellow light"

show Green = "Green light"

ghci> Red == Red

True

ghci> Red == Yellow

False

ghci> Red 'elem' [Red, Yellow, Green]

True

ghci> [Red, Yellow, Green]

[Red light, Yellow light, Green light]

◄ instance Eq TrafficLight

◄ instance Eq TrafficLight

◄ instance Eq TrafficLight

◄ instance Show TrafficLight

http://learnyouahaskell.com/making-our-own-types-and-typeclasses

Instance Maybe m

```
instance Eq Maybe where ...
```

```
instance Eq (Maybe m) where
Just x == Just y = x == y
Nothing == Nothing = True
_ == _ = False
```

```
instance (Eq m) => Eq (Maybe m) where
Just x == Just y = x == y
Nothing == Nothing = True
_ == _ = False
```

Maybe is not a concrete type

Maybe m is a concrete type

all types of the form **Maybe m** to be part of the **Eq** typeclass,

but only those types where the **m** (what's contained inside the **Maybe**) is also a part of **Eq**.

http://learnyouahaskell.com/making-our-own-types-and-typeclasses

Eq, Ord, Show classes

Since <u>equality tests</u> between values are frequently used most of your own data types should be <u>members</u> of **Eq**.

Prelude classes

- Eq
- Ord
- Show

for the convenience, Haskell has a way to declare such "obvious" **instance definitions** using the keyword **deriving**.

Deriving instance example

```
data Foo = Foo {x :: Integer, str :: String}
```

deriving (Eq, Ord, Show)

This makes **Foo** an **instance** of **Eq** ith an <u>automatically generated</u> <u>definition</u> of **==**

also an instance of Ord and Show

deriving (Eq, Ord, Show)

```
*Main> Foo 3 "orange" == Foo 6 "apple"
False
*Main> Foo 3 "orange" /= Foo 6 "apple"
True
```

Deriving instance pros and cons

The **types** of **elements** inside the **data** type must also be **instances** of the **class** you are <u>deriving</u>.

Deriving instances

- synthesis of functions for a <u>limited</u> set of predefined classes
- against the general Haskell philosophy : "built in things are not special",
- induces compact codes
- often reduces errors in coding

 (an example: an instance of Eq such that x == y
 would not be equal to y == x would be flat out wrong).

Derivable Classes

Eq

Equality operators == and /=

Ord

Comparison operators < <= > >=; min, max, and compare.

Enum

For enumerations only. Allows the use of list syntax such as [Blue .. Green].

Bounded

Also for enumerations, but can also be used on types that have only one constructor.

Provides minBound and maxBound as the lowest and highest values that the type can take.

Show

Defines the function show, which converts a value into a string, and other related functions.

Read

Defines the function **read**, which <u>parses</u> a <u>string</u> into a <u>value</u> of the type, and other related functions.

Functors, Applicatives, Monads

functors: you apply a <u>function</u> to a <u>wrapped</u> <u>value</u>

applicatives: you apply a <u>wrapped</u> <u>function</u> to a <u>wrapped</u> <u>value</u>

monads: you apply a <u>function</u> that <u>returns</u> a <u>wrapped</u> <u>value</u>, to a <u>wrapped</u> <u>value</u>

functors: using fmap or <\$>

applicatives: using <*> or liftA

monads: using >>= or liftM

Functors

Functors use the **fmap** or <\$> functions

```
fmap :: Functor f => (a -> b) -> f a -> f b

< :: Functor f => (a -> b) -> f a -> f b
```

This takes a function and applies to to the wrapped elements

```
fmap (x \rightarrow x + 1) (Just 1)
-- Applies (+1) to the inner value, returning (Just 2)
-- Applies (+1) to an empty wrapper, returning Nothing

fmap (x \rightarrow x + 1) [1, 2, 3]
-- Applies (+1) to all inner values, returning [2, 3, 4]

(x \rightarrow x + 1) <$> [1, 2, 3]
-- Same as above [2, 3, 4]
```

Applicatives

Applicatives use the <*> function:

"> :: Applicative
$$f => f(a -> b) -> fa -> fb$$

This takes a wrapped function and applies it to the wrapped elements

Monads – return

There are two relevant functions in the **Monad typeclass**:

```
return :: Monad m => a -> m a
(>>=) :: Monad m => m a -> (a -> m b) -> m b
```

The return function takes a raw, <u>unwrapped</u> value, and <u>wraps</u> it up in the desired monadic type.

```
makeJust :: a -> Maybe a
makeJust x = return x

let foo = makeJust 10 -- returns (Just 10)
```

Monads - bind

The bind function lets you temporarily unwrap the inner elements of a **Monad** and pass them to a <u>function</u> that performs some action that <u>wraps</u> them back UP in the same monad.

This can be used with the return function in trivial cases:

Monads – a binding operand

functions to chain together that don't require to use **return**.

```
getLine :: IO String -- return String type value as a result
putStrLn :: String -> IO ()
```

function call examples

```
getLine >>= (\x -> putStrLn x) -- gets a line from IO and prints it to the console
getLine >>= putStrLn -- with currying, this is the same as above
```

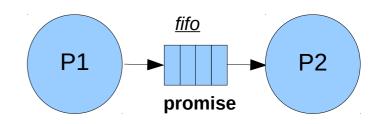
Monads – a chain of functions

functions to chain together that don't require to use **return**.

Promises and Mediators

the concept of **promises** (particularly in Javascript)

A **promise** is an **object** that <u>acts</u> as a <u>placeholder</u> for the **result value** of an <u>asynchronous</u>, <u>background</u> **computation**, like fetching some data from a remote service.



it serves as a mediator

between the <u>asynchronous computation</u> and <u>functions</u> that need to **operate** on its <u>anticipated</u> **result**.

Act

Behavior

Operation

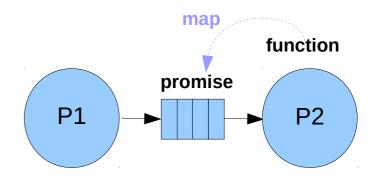
Define

Connect

Map a function over a promise

A mediator allows us to say what <u>function</u> should apply to the <u>result</u> of a <u>background</u> task, <u>before</u> that task has <u>completed</u>.

When you **map** a **function** over a **promise**, the <u>value</u> that your function should apply to may <u>not</u> have been <u>computed</u> yet and in fact, if there is an error somewhere it may never be computed.

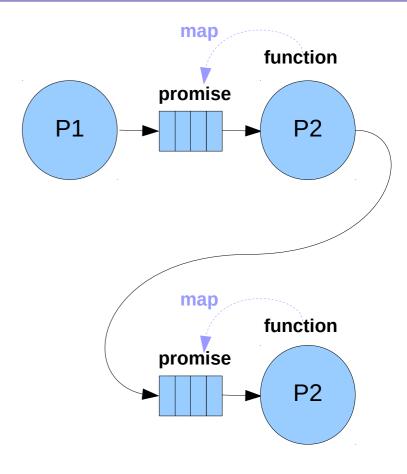


Chaining a function onto a promise

Promise libraries usually support
a functorial/monadic API
where you can chain a function onto a promise,
which produces another promise
that produces the result of applying that function
to the original promise's result.

the value of the functor/monad interface

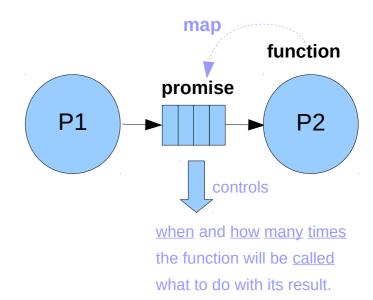
Promises allow you to say what **function** should apply to the **result** of a background task, before that task has completed.



Interfaces

think functor/applicative/monad
as interfaces for mediator objects
that sit in between functions and arguments,
and connect them indirectly according to some policy.

The <u>simplest</u> way to use a function is just to <u>call</u> it with some <u>arguments</u>;



First-class functions

A **higher-order function** is a function that <u>takes</u> other functions <u>as arguments</u> or returns a function as result.

https://softwareengineering.stackexchange.com/questions/303472/what-is-the-purpose-of-wrapped-values-in-haskell

Interfaces with first-class functions

if you have <u>first-class functions</u>, you have other, indirect options—

you can <u>supply</u> the function to a <u>mediator</u> object that will <u>control</u> <u>when</u> and <u>how many times</u> the function will be <u>called</u>, and what to do with its result.

function promise P1 P2 controls

when and how many times the function will be called what to do with its result.

First-class functions

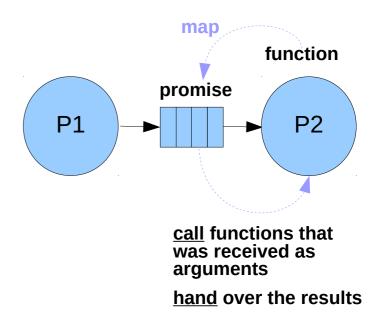
A **higher-order function** is a function that <u>takes</u> other functions <u>as arguments</u> or <u>returns</u> a function as result.

https://softwareengineering.stackexchange.com/questions/303472/what-is-the-purpose-of-wrapped-values-in-haskell

Promises and Mediators

Promises <u>call</u> the <u>functions</u> supplied to them when the <u>result</u> of some background task is <u>completed</u>

The <u>results</u> of those functions are then <u>handed</u> <u>over</u> to <u>other promises</u> that are waiting for them.



https://softwareengineering.stackexchange.com/questions/303472/what-is-the-purpose-of-wrapped-values-in-haskell

General Monad - MonadPlus

Haskell's **Control.Monad** module defines a typeclass, **MonadPlus**, that enables abstract the common pattern eliminating **case** expressions.

```
class Monad m => MonadPlus m where
mzero :: m a
mplus :: m a -> m a -> m a
```

```
class (Monad m) => MonadPlus m where
```

```
instance MonadPlus [] where
mzero = []
mplus = (++)
```

```
instance MonadPlus Maybe where
  mzero = nothing
  nothing `mplus` ys = ys
  xs   `mplus` _ = xs
```

http://book.realworldhaskell.org/read/programming-with-monads.html

General Monad - MonadPlus Laws

The class **MonadPlus** is used for monads that have a <u>zero</u> <u>element</u> and a <u>plus</u> operation:

```
class (Monad m) => MonadPlus m where
```

mzero :: m a

mplus :: m a -> m a -> m a

For lists, the zero value is [], the empty list.

The I/O monad has no zero element and

is not a member of this class.

```
m >>= \x -> mzero = mzero
mzero >>= m = mzero
```

The zero element laws:

```
m `mplus` mplus = m
mplus `mplus` m = m
```

The laws governing the mplus operator

The mplus operator is ordinary list concatenation in the list monad.

http://book.realworldhaskell.org/read/programming-with-monads.html

Functional Dependency (fundep)

class Mult a b c | a b -> c where
(*) :: a -> b -> c

a b -> c means

c is <u>uniquely</u> <u>determined</u> from **a** and **b**

fundeps are not standard Haskell 98.

(Nor are multi-parameter type classes, for that matter.) They are, however, supported at least in **GHC** and **Hugs** and will almost certainly end up in Haskell'.

class Mult a b c where

(*) :: a -> b -> c

https://wiki.haskell.org/Functional_dependencies

Functional Dependency – a type inferencer

In a multiparameter typeclass, by default, the **type variables** are considered <u>independently</u>.

The **type inferencer** has to determine **a** and **b** <u>independently</u>, then check to see if the **instance** exists.

class Foo a b

Functional dependencies narrow down possible choices.

effective, useful

class Foo a b | a -> b

Look, if you determine what **a** is, then there is a unique **b** so that **Foo a b** exists, so don't bother trying to infer **b**, just go look up the instance and typecheck that.

https://stackoverflow.com/questions/20040224/functional-dependencies-in-haskell

Functional Dependency – return type polymorphism

Fundep is useful with return type polymorphism

class Foo a b c where

bar :: a -> b -> c

there's no way to infer

bar (bar "foo" 'c') 1

Because we have no way of determining c of $a \rightarrow b \rightarrow c$.

Even if we only wrote one instance for **String** and **Char**, we have to assume that someone might/will come along and add another instance later on.

https://stackoverflow.com/questions/20040224/functional-dependencies-in-haskell

Functional Dependency – determining the return type

With **fundeps** we <u>don't</u> have to specify the **return type**, which is annoying.

And now it's easy to see that the return type **c** of **bar "foo" 'c'** is <u>unique</u> and thus <u>inferable</u>.

class Foo a b c | a b -> c where bar :: a -> b -> c

https://stackoverflow.com/guestions/20040224/functional-dependencies-in-haskell

Type Constructors with parameters

Type constructors take other types as parameters to eventually produce concrete types. – like a function

type constructors can be <u>partially applied</u> just like functions can Either String is a type that <u>takes</u> one type and <u>produces</u> a concrete type,

by using type declarations

like Either String Int

formally defining how **types** are <u>applied</u> to **type constructors**, formally defining how **values** are <u>applied</u> to **functions**

Kind of a type

values like

3 – Int

"YEAH" – String

takeWhile – a function value

each have their own type.

types are little labels that values carry

so that we can reason about the values.

types have their own another little labels, called kinds.

A **kind** can be considered as the **type of a type**.

Examining the kind of a type

To examine the kind of a type

using the :k command in GHCI.

ghci>:k Int

Int :: *

A * means that the type is a **concrete type**.

A **concrete type** is a type that doesn't take any type **parameters** and **values** can only have types that are **concrete types**.

Kind of a type constructor

```
ghci>:k Maybe
  Maybe :: * -> *
The Maybe type constructor
     takes one concrete type
                                      (like Int)
     and returns a concrete type
                                      (like Maybe Int)
Int -> Int represents a function
     taking an Int and returning an Int,
* -> * represents a type constructor
     taking an concrete type and returning a concrete type
```

Kind of a type constructor applied with a type parameter

```
apply the type parameter to Maybe
  ghci>:k Maybe Int
  Maybe Int :: *
the type parameter Int is applied to Maybe
The kind of Maybe Int is a concrete type
  :t isUpper
                                :k isUpper
  Char -> Bool
  :t isUpper 'A' (True)
                                :k isUpper 'A'
 Bool
```

References

- [1] ftp://ftp.geoinfo.tuwien.ac.at/navratil/HaskellTutorial.pdf
- [2] https://www.umiacs.umd.edu/~hal/docs/daume02yaht.pdf