## DLTI Convolution (1B)

- Circular Convolution
- Numerical Convolution
- Moving Average Filter

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## Based on

Introduction to Signal Processing
S. J. Ofranidis

The necessities in DSP C Programming
FIR Filter (A.pdf) 20191114

## Linear Convolution using the DFT



## Circular Convolution using the DFT

$$
y[n]=\sum_{k=-\infty}^{+\infty} h[k] x[n-k]
$$



## Linear Convolution

$$
\begin{array}{ll}
x[n] & 0 \leq n \leq L-1 \\
h[n] & 0 \leq n \leq M-1 \\
y[n] & 0 \leq n \leq L+M-2
\end{array}
$$

$$
Y\left(e^{j \hat{\omega}}\right)=H\left(e^{j \hat{\omega}}\right) X\left(e^{j \hat{\omega}}\right)
$$

DTFT
$\omega$ continuous frequency variable


$$
y[n]=\sum_{m=0}^{N-1} h[m] x\left[\langle n-k\rangle_{N}\right] \quad 0 \leq n \leq N-1
$$

Circular Convolution
$y[n]=h[n] \times[n]$

## Circular Convolution using the DFT

Linear Convolution

$$
y[n]=h[n] * x[n]
$$

$$
Y\left(e^{j \hat{\omega}}\right)=H\left(e^{j \hat{\omega}}\right) X\left(e^{j \hat{\omega}}\right)
$$

$$
y[n]=\sum_{k=-\infty}^{+\infty} h[k] x[n-k]
$$

$\omega$ continuous frequency variable

## Circular Convolution



## Circular Convolution (1)

$$
\begin{aligned}
& y[n]=\frac{1}{N} \sum_{k=0}^{N-1} H[k] X[k] W_{N}^{-k n} \\
&=\frac{1}{N} \sum_{k=0}^{N-1}\left\{\sum_{m=0}^{N-1} h[m] W_{N}^{+k m}\right\}\left\{\sum_{l=0}^{N-1} x[l] W_{N}^{+k l}\right\} W_{N}^{-k n} \\
&=\sum_{m=0}^{N-1} h[m] \sum_{l=0}^{N-1} x[l]\left\{\frac{1}{N} \sum_{k=0}^{N-1} W_{N}^{+k(m+l-n)}\right\} \\
&\left\{\frac{1}{N} \sum_{k=0}^{N-1} W_{N}^{+k(m+l-n)}\right\}=1, \quad m+l-n=r N \\
& l=n-m+r N=(n-m) \bmod N=\langle n-m\rangle_{N}
\end{aligned}
$$

## Circular Convolution (1)

$$
\begin{aligned}
& X_{1}[k]=\sum_{n=0}^{N-1} x_{1}[n] e^{-j 2 \pi n k / N} \quad k=0,1, \cdots, N-1 \\
& X_{2}[k]=\sum_{n=0}^{N-1} X_{2}[n] e^{-j 2 \pi n k / N} \quad k=0,1, \cdots, N-1 \\
& X_{3}[k]=X_{1}[k] X_{2}[k] \quad k=0,1, \cdots, N-1 \\
& x_{3}[n]=\frac{1}{N} \sum_{k=0}^{N-1} X_{3}[k] e^{+j 2 \pi k m / N} \\
& =\frac{1}{N} \sum_{k=0}^{N-1} X_{1}[k] X_{2}[k] e^{+j 2 \pi k m / N} \\
& =\frac{1}{N} \sum_{k=0}^{N-1}\left\{\sum_{n=0}^{N-1} x_{1}[n] e^{-j 2 \pi k n / N}\right\}\left\{\sum_{l=0}^{N-1} \chi_{2}[l] e^{-j 2 \pi k l / N}\right\} e^{+j 2 \pi k m / N} \\
& =\frac{1}{N} \sum_{n=0}^{N-1} x_{1}[n] \sum_{l=0}^{N-1} x_{2}[l] \sum_{k=0}^{N-1} e^{-j 2 \pi k n / N} e^{-j 2 \pi k l / N} e^{+j 2 \pi k m / N} \\
& =\frac{1}{N} \sum_{n=0}^{N-1} x_{1}[n] \sum_{l=0}^{N-1} x_{2}[l] \sum_{k=0}^{N-1} e^{+j 2 \pi k(m-n-l) / N}
\end{aligned}
$$

## Circular Convolution (2)

$$
\begin{aligned}
& X_{1}[k]=\sum_{n=0}^{N-1} x_{1}[n] e^{-j 2 \pi n k / N} \\
& X_{2}[k]=\sum_{n=0}^{N-1} x_{2}[n] e^{-j 2 \pi n k / N} \quad k=0,1, \cdots, N-1 \\
& X_{3}[k]=X_{1}[k] X_{2}[k] \quad k=0,1, \cdots, N-1 \\
& x[k]=\frac{1}{N} \sum_{n=0}^{N-1} x_{1}[n] \sum_{l=0}^{N-1} x_{2}[l] \sum_{k=0}^{N-1} e^{+j 2 \pi k(m-n-l) / N} \\
& a=e^{+j 2 \pi(m-n-l) / N}
\end{aligned}
$$

$$
\begin{array}{rlr}
\sum_{k=0}^{N-1} a^{k} & =N \quad(a=1) \\
& =\frac{1-a^{N}}{1-a} \quad(a \neq 1)
\end{array}
$$

## Circular Convolution (2)

$$
\begin{aligned}
& \frac{(m-n-l)}{N} \in\left\{\begin{array}{cccc}
\frac{0}{N}, & \frac{1}{N}, & \cdots, & \frac{N-1}{N}, \\
\frac{0}{N}+1, & \frac{1}{N}+1, & \cdots, & \frac{N-1}{N}+1, \\
\frac{0}{N}+2, & \frac{1}{N}+2, & \cdots, & \frac{N-1}{N}+2, \\
\vdots & \vdots & & \vdots
\end{array}\right\} \\
& e^{j \frac{2 \pi(m-n-l)}{N}} \in\left\{e^{j 2 \pi \cdot 0 / N}, e^{j 2 \pi \cdot 1 / N}, \cdots, e^{+j 2 \pi(N-1) / N}\right\} \quad e^{j \frac{2 \pi(m-n-l)}{N}}=e^{j \frac{2 \pi((m-n-l))_{N}}{N}} \\
& ((m-n-l))_{N}=0 \quad e^{j \frac{2 \pi(m-n-l)}{N}}=e^{j \frac{2 \pi((m-n-l))_{N}}{N}}=1 \\
& \square \sum_{k=0}^{N-1} e^{j \frac{2 \pi k(m-n-l)}{N}}=\sum_{k=0}^{N-1} e^{j \frac{2 \pi k((m-n-l))_{N}}{N}}=N \\
& ((m-n))_{N}=l
\end{aligned}
$$

## Circular Convolution

| $x[2]$ | $x[1]$ | $x[0]$ | $x[3]$ |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $x[3]$ | $x[1]$ | $x[2]$ |  | $x[0]$ | $x[1]$ |
| $x[0]$ | $x[3]$ | $x[3]$ | $x[0]$ | $x[2]$ |  |
| $x\left[\langle n-0\rangle_{4}\right]$ | $x\left[\langle n-1\rangle_{4}\right]$ | $x[2]$ | $x[1]$ |  |  |
|  |  |  |  |  |  |
|  |  | $x\left[\langle n-2\rangle_{4}\right]$ | $x\left[\langle n-3\rangle_{4}\right]$ |  |  |

## Circular Convolution

| $x[3]$ | $\frac{x[2]}{h[2]}$ | $x[1]$ |  |  | $x[0]$ |  |  | $x[3]$ |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $h[2]$ |  |  | $x[0] \quad x[1]$ | $h[3]$ | $h[1]$ | $x[3] x[0]$ | $h[3] \quad h[1] \times[2]$ |  |  |
|  | $h[3] \quad h[1]$ | $x[1] x[2]$ | $h[3]$ | $h[1]$ |  |  |  |  |  |  |  |
|  | $h[0]$ |  | $h[$ |  |  |  |  |  |  |  |  |
|  | $x[0]$ |  | $x$ |  |  |  |  |  |  |  |  |
|  | $x\left[\langle n-0\rangle_{4}\right]$ |  | $x[<$ | $-1\rangle_{4}$ ] |  | $x[\langle n$ |  |  |  | $\left.-3\rangle_{4}\right]$ |  |

## Circular Convolution

## Octave conv()

conv ( $a, b$ )
conv (a, b, shape)
Convolve two vectors $a$ and $b$.

The length of the output vector : length (a) + length (b) - 1
When $a$ and $b$ are the coefficient vectors of two polynomials, the convolution represents the coefficient vector of the product polynomial.

The optional shape argument may be
shape $=$ "full"
Return the full convolution. (default)
shape = "same"
Return the central part of the convolution with the same size as a.
See also: deconv conv2 convn fftconv

## Octave fftconv()

fftconv ( $x, y$ )
fftconv ( $\mathrm{x}, \mathrm{y}, \mathrm{n}$ )
Convolve two vectors using the FFT for computation.
$\mathrm{c}=\mathrm{fftconv}(\mathrm{x}, \mathrm{y})$ returns
a vector of length equal to length $(x)+$ length $(y)-1$.
If $x$ and $y$ are the coefficient vectors of two polynomials, the returned value is the coefficient vector of the product polynomial.

The computation uses the FFT by calling the function fftfilt.
If the optional argument n is specified, an N -point FFT is used.
fftfilt (b, x, n)
With two arguments, fftfilt filters x with the FIR filter b using the FFT.
Given the optional third argument, n ,
fftfilt uses the overlap-add method to filter x with b using an N -point FFT.
If $x$ is a matrix, filter each column of the matrix.

## Linear delay-line buffer

At each time instant, the data in the delay line are shifted one memory location ahead.
shifting the data forward
while holding the buffer addresses fixed


## Wrapped linear delay-line buffer


the data are shifted one location forward the buffer addresses are fixed

the starting address $p$ is fixed
index pointer at time $\mathbf{n}$ pointer at time $\mathbf{n + 1}$ data at time $\mathbf{n}$ data at time $\mathbf{n + 1}$

## Circular delay-line buffer


the data are kept fixed the addresses are shifted backwards

the starting address $p$ is shifted backward

| 0 | 1 | 2 | 3 | location index |
| :--- | :--- | :--- | :--- | :--- |
| $p$ | $p+1$ | $p+2$ | $p+3$ | pointer at time $\mathbf{n}$ |
| $p+1$ | $p+2$ | $p+3$ | $p$ | pointer at time $\mathbf{n + 1}$ |
| $x[n]$ | $x[n-1]$, | $x[n-2]$, | $x[n-3]$ | data at time $\mathbf{n}$ |
| $x[n]$, | $x[n-1]$, | $x[n-2]$, | $x[n+1]$ | data at time $\mathbf{n + 1}$ |

## Pointer variable p

A pointer variable $\mathbf{p}$ always points to the current input

A buffer address w (or w register)

At each time instant, the buffer (register) pointed to by $\mathbf{p}$ gets loaded with the current input sample
$* p=x \quad p[0]=x$

## Offset integer q

The pointer $\mathbf{p}$ is restricted to lie within the pointer range of the linear buffer $\mathbf{w}$,

$$
w \leq p \leq w+M
$$

always point at somewhere in the w buffer, w[q]

$$
\begin{aligned}
& p \longrightarrow \begin{array}{c}
w \\
\\
\\
\\
\\
{ }^{*}(w+0) \\
{ }^{*}(w+1) \\
{ }^{*}(w+2) \\
{ }^{*}(w+3)
\end{array} \\
& p=w+3 \\
& q=3
\end{aligned}
$$

$$
\begin{aligned}
& p=w+q \quad \Rightarrow \\
& * p=p[0]=w[q]
\end{aligned}
$$

$\mathbf{q}$ is an integer that gives the offset of $\mathbf{p}$ from the fixed beginning of the w buffer.

$$
0 \leq q \leq M
$$

## Address pointer $p$ and index $q$



$$
\begin{gathered}
\mathrm{n}=\quad 0 \\
4, \\
8, \\
12 \\
\cdots \\
\mathrm{w}[0]=*(w+0) \\
\mathrm{p} \text { when } \mathrm{q}=0
\end{gathered}
$$


$M=3$
modulo $(\mathrm{M}+1)=$ modulo 4

## Address pointer p (moving) and w (fixed)

modulo ( $\mathrm{M}+1$ )
$w \leq p \leq w+M$
$p=w+q \quad * p=*(w+q)=w[q]$

| $\Delta$ | w | *(w+0) |
| :---: | :---: | :---: |
|  | w+1 | *(w+1) |
|  | w+2 | *(w+2) |
| $\mathrm{p} \longrightarrow$ | w+3 | *(w+3) |
| $\triangle$ |  |  |

$0 \leq q \leq M$

$$
\begin{array}{ll}
p[0] & *(p+0)=*(w+3+0)=w[3] \\
p[1] & *(p+1)=*(w+3+1)=w[0] \\
p[2] & *(p+2)=*(w+3+2)=w[1] \\
p[3] & *(p+3)=*(w+3+3)=w[2]
\end{array}
$$

## Accessing the buffer w[k] by $\mathrm{p}[\mathrm{i}]$

modulo (3+1)

| $\mathrm{p} \longrightarrow \mathrm{w}$ | *(w+0) |
| :---: | :---: |
|  | *(w+1) |
|  | *(w+2) |
|  | *(w+3) |


| $w$ |
| :--- |
| $*(w+0)$ |
| ${ }^{*}(w+1)$ |
| *(w+2) |
| $*(w+3)$ |



$$
\begin{array}{ccll}
p[0]=w[0] & p[0]=w[3] & p[0]=w[2] & p[0]=w[1] \\
p[1]=w[1] & p[1]=w[0] & p[1]=w[3] & p[1]=w[2] \\
p[2]=w[2] & p[2]=w[1] & p[2]=w[0] & p[2]=w[3] \\
p[3]=w[3] & p[3]=w[2] & p[3]=w[1] & p[3]=w[0] \\
\text { when } q=0 & \text { when } q=3 & \text { when } q=2 & \text { when } q=1
\end{array}
$$

## Address pointer p moves backward

modulo (3+1)


$$
\begin{aligned}
& p[0] \\
& p[1] \\
& p[2] \\
& p[3]
\end{aligned}
$$

$$
*(p+0)=*(w+q+0)=w[0], w[3], w[2], w[1]
$$

$$
p[1] \quad *(p+1)=*(w+q+1)=w[1], w[0], w[3], w[2]
$$

$$
p[2] \quad *(p+2)=*(w+q+2)=w[2], w[1], w[0], w[3]
$$

$$
*(p+3)=*(w+q+3)=w[3], w[2], w[1], w[0]
$$

modulo $4 \quad q=0(q=3, q=2(q=1)$

## Address pointer $p$

offset index
modulo ( $\mathrm{M}+1$ )

$$
w \leq p \leq w+M \quad 0 \leq q \leq M
$$01

$$
p=w+q
$$

$$
\begin{aligned}
& * p=*(w+q)=w[q] \\
& *(p+i)=*(w+q+i)=w[q+i]
\end{aligned}
$$

$$
p+i=w+(p-w+i)
$$

$$
*(p+i)=*(w+(p-w+i))
$$

$$
=*(w+(q+i))
$$

$$
w[(p-w+i) \%(M+1)]=w[(q+i) \%(M+1)]
$$

## Accessing w buffer by the pointer $\mathbf{p}$

modulo ( $\mathrm{M}+1$ )

$$
w[(p-w+i) \%(M+1)]=w[(q+i) \%(M+1)]
$$

offset index


## Address pointer p moves backward

modulo ( $\mathrm{M}+1$ )

$$
w[(p-w+i) \%(M+1)]=w[(q+i) \%(M+1)]
$$



## Increasing and decreasing the pointer $\mathbf{p}$



## Internal States



## Address pointer $\mathbf{p}$ and incoming inputs


$n=4 \quad x 4$

$n=1 \quad x 1$

$n=5 \quad x 5$

$n=2 \quad x 2$

$n=6 \quad x 6$

$n=3 \quad x 3$

$n=7 \quad x 8$


## Address pointer p and incoming inputs - linear array view



## Internal states over the first 8 steps of n

| n | q | [w0] | [w1] | [w2] | [w3] | s0 s1 s2 s3 | s0 | s1 | s2 | s3 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | x0 | 0 | 0 | 0 | [w0] [w1] [w2] [w3] | x 0 | 0 | 0 | 0 |
| 1 | 3 | x0 | 0 | 0 | x1 | [w3] [w0] [w1] [w2] | x 1 | x0 | 0 | 0 |
| 2 | 2 | x0 | 0 | x2 | x1 | [w2] [w3] [w0] [w1] | x 2 | x1 | x0 | 0 |
| 3 | 1 | x0 | x3 | x2 | x1 | [w1] [w2] [w3] [w0] | $x 3$ | x2 | x1 | x0 |
| 4 | 0 | $\times 4$ | x3 | x2 | x1 | [w0] [w1] [w2] [w3] | x 4 | x3 | x2 | x1 |
| 5 | 3 | x4 | x3 | x2 | x5 | [w3] [w0] [w1] [w2] | x 5 | x4 | x3 | x2 |
| 6 | 2 | x4 | x3 | x6 | x5 | [w2] [w3] [w0] [w1] | x 6 | x5 | x4 | x3 |
| 7 | 1 | x4 | x7 | x6 | x5 | [w1] [w2] [w3] [w0] | x7 | x6 | x5 | x4 |

[w0] value at the buffer wo [w1] value at the buffer w1 [w2] value at the buffer w2 [w3] value at the buffer w3

$$
\left(\begin{array}{l}
s_{0} \\
s_{1} \\
s_{2} \\
s_{3}
\end{array}\right) \quad\left(\begin{array}{l}
p[0] \\
p[1] \\
p[2] \\
p[3]
\end{array}\right) \quad\left(\begin{array}{l}
w[2] \\
w[3] \\
w[0] \\
w[1]
\end{array}\right)
$$

## References

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