Carry and Overflow

Young W. Lim

2024-04-27 Sat

1/112

Young W. Lim Carry and Overflow 2024-04-27 Sat

Outline

- Based on
- Overview
 - Overview
- Carry flag
 - TOC: Carry flag
 - Examples of signed and unsigned integer arithmetic
 - Carry flag in unsigned and signed computations
 - Rules for the carry flag
 - Method for computing the carry flag
 - More examples of the carry flag
- Overflow flag
 - TOC: Overflow flag
 - Overflow flag in unsigned and signed computations
 - Rules for the overflow flag
 - Method 1 for computing the overflow flag
 - Method 2 for computing the overflow flag
 - More examples of the overflow flag



Based on

"Self-service Linux: Mastering the Art of Problem Determination",

Mark Wilding

"Computer Architecture: A Programmer's Perspective", Bryant & O'Hallaron

I, the copyright holder of this work, hereby publish it under the following licenses: GNU head Permission is granted to copy, distribute and/or modify this document under the terms of the GNU Free Documentation License, Version 1.2 or any later version published by the Free Software Foundation; with no Invariant Sections, no Front-Cover Texts, and no Back-Cover Texts. A copy of the license is included in the section entitled GNU Free Documentation License.

CC BY SA This file is licensed under the Creative Commons Attribution ShareAlike 3.0 Unported License. In short: you are free to share and make derivative works of the file under the conditions that you appropriately attribute it, and that you distribute it only under a license compatible with this one.

3 / 112

Compling 32-bit program on 64-bit gcc

- gcc -v
- gcc -m32 t.c
- sudo apt-get install gcc-multilib
- sudo apt-get install g++-multilib
- gcc-multilib
- g++-multilib
- gcc -m32
- objdump -m i386

TOC: Overview

- Carry flag and overflow flag
- Signed and unsigned computations
- Flags for an unsigned number
- Flags for a signed number
- Detecting errors in usigned and signed arithmetic
- The verb to overflow v.s. the overflow flag

Carry flag and overflow flag

- considering carry and overflow flags in x86
- do not confuse the carry flag with the overflow flag in integer arithmetic.
- the ALU always sets these flags appropriately when doing any integer math.
- these flags can occur on its own, or both together.

Signed and unsigned computations

- the CPU's ALU <u>doesn't</u> care or know whether <u>signed</u> or <u>unsigned</u> computations are performed;
- the <u>ALU</u> just performs integer arithmetic and sets the flags appropriately.
- It's up to the <u>programmer</u> to know which flag to check after the arithmetic is done.

Flags for an unsigned number

- if a word is treated as an unsigned number,
 - the carry flag must be used to check if the result is fit into n-bit or (n+1)-bit number
 - the overflow flag is irrelevant to an unsigned number arithmetic

Flags for a signed number

- if a word is treated as an signed number,
 - the carry flag is *irrelevant* to an signed number arithmetic
 - the overflow flag must be used to check if the result is wrong or not

Detecting errors in usigned and signed arithmetic (1)

	unsigned integer arithmetic	signed integer arithmetic
CF Carry Flag	detects overflows extends an n -bit result into an $(n+1)$ -bit result	
OF Overflow Flag		detects <i>overflows</i> errors the result cannot be used

Detecting errors in usigned and signed arithmetic (2)

- unsigned integer arithmetic overflow is indicated by the carry flag
 - P + P CF=1 \rightarrow carry out the result is too large for an *n-bit* integer
 - P-P CF=1 \rightarrow borrow in the result is too small for an *n-bit* integer
- signed integer arithmetic overflow is indicated by the overflow flag
 - $P + P \rightarrow N$ OF=1 \rightarrow overflow the result is not correct
 - $N + N \rightarrow P$ OF=1 \rightarrow overflow the result is $\overline{\text{not}}$ correct
- P (positive), N (negative)

Detecting errors in usigned and signed arithmetic (3)

- unsigned integer arithmetic overflow is indicated by the carry flag
 - the *overflowed n*-bit result can be extended into (n+1)-bit result by using the carry flag
- signed integer arithmetic overflow is indicated by the overflow flag
 - the overflowed n-bit result cannot be used

https://stackoverflow.com/questions/47333458/assembly-x86-64-setting-carry-flag-f

The verb to overflow v.s. the overflow flag (1)

- Do not confuse the <u>English verb</u> to overflow with the overflow flag in the ALU.
- The <u>verb</u> to overflow is used casually to indicate that some math result doesn't fit in the number of bits available;
- it could be integer math, or floating-point math, or whatever.
- The overflow flag is set specifically by the ALU
 it isn't the same as the casual English verb "to overflow"

The verb to overflow v.s. the overflow flag (2)

- In English, we may say
 "the binary/integer math overflowed the number of bits available for the result, causing the carry flag to come on".
- Note how this English usage of the verb "to overflow" is not the same as saying the overflow flag is on".
- A math result can <u>overflow</u> (the <u>verb</u>) the number of bits available <u>without</u> turning on the ALU <u>overflow flag</u>

Computing Carry and Overflow Flags

CF (carry flag) and OF (overflow flag) computation

ADD (addition)	SUB (subtraction)
$CF = C_n$	$CF = \overline{C_n}$
OF = $C_n \bigoplus C_{n-1}$	OF = $C_n \bigoplus C_{n-1}$
a 2's complement addition $A + B = A + B + 0$	a transformed addition $A - B = A + \overline{B} + 1$
$\{C_n, S_{n-1}\} = a_{n-1} + b_{n-1} + c_{n-1}$	$\{C_n, S_{n-1}\} = a_{n-1} + \overline{b_{n-1}} + c_{n-1}$
${C_{n-1}, S_{n-2}} = a_{n-2} + b_{n-2} + c_{n-2}$	${C_{n-1}, S_{n-2}} = a_{n-2} + \overline{b_{n-2}} + c_{n-2}$

https://www.csie.ntu.edu.tw/~cyy/courses/assembly/12fall/lectures/handouts/lec14_

TOC: Carry flag

- Examples of signed and unsigned integer arithmetic
- Carry flag in unsigned and signed computations
- Rules for the carry flag
- Method for computing the carry flag
- More examples of the carry flag

TOC: Examples of signed and unsigned integer arithmetic

- Examples of interpreting signed and unsigned numbers
- Examples of signed and unsigned integer arithmetic
- 2's complements
- Unsigned subtraction
- Signed subtraction
- Interpreting the result as a signed or an unsigned integer
- Summary of signed and unsigned subtractions
- Examples of unsigned integer overflows
- Examples of signed integer overflows

Examples of interpreting signed and unsigned numbers (1)

• interpreting 0xFFFFBDC3

```
as an unsigned (positive) number +0xFFFFBDC3 +4294950339_{10} as a signed (negative) number -0x0000423D -16957_{10}
```

https://stackoverflow.com/questions/47333458/assembly-x86-64-setting-carry-flag-fe

18 / 112

Young W. Lim Carry and Overflow 2024-04-27 Sat

Examples of interpreting signed and unsigned numbers (2)

- interpreting 0xFFFFBDC3
 - ullet as an unsigned (positive) number \mid +0xFFFFBDC3 \mid +4294950339 $_{10}\mid$

$$\begin{array}{l} 15*16^7 + 15*16^6 + 15*16^5 + 15*16^4 \\ + 11*16^3 + 13*16^2 + 12*16^1 + 3*16^0 \end{array}$$

 \bullet as a signed (negative) number \mid -0x0000423D \mid -16957 $_{10}\mid$

$$0*16^7 + 0*16^6 + 0*16^5 + 0*16^4 + 4*16^3 + 2*16^2 + 3*16^1 + 13*16^0$$

Examples of interpreting signed and unsigned numbers (3)

 \bullet the 2's complement of 0xFFFFBDC3 : 0x0000423D (= +16957₁₀)

• the 2's complement of 0x0000423D : 0xFFFFBDC3 (= -16957₁₀)

```
0 0 0 0 4 2 3 D

0x0000423D 0x0000_0000_0000_0100_0010_0011_1101

0x0000BDC2 0x1111_1111_1111_1111_1011_1101_00010 (1's complement)

0xFFFFBDC3 0x1111_1111_1111_1111_1011_1100_0011 (2's complement)

F F F F B D C 3
```

https://stackoverflow.com/questions/47333458/assembly-x86-64-setting-carry-flag-fe

Examples of signed and unsigned integer arithmetic

• subtracting 0x0000618D from 0x0000195D

0x0000195D - 0x0000618D

unsigned subtraction

subtraction by hand

0x0000195D + (-0x0000618D)

signed subtraction

the *transformed addition* using the 2's complement of <u>subtrahend</u>

https://stackoverflow.com/questions/47333458/assembly-x86-64-setting-carry-flag-f

Young W. Lim Carry and Overflow 2024-04-27 Sat 21 / 112

2's complements

• the 2's complement of 0x0000618D: 0xFFFF8E73 (= -24973₁₀)

 \bullet the 2's complement of <code>0xFFFF8E73</code> : <code>0x0000618D</code> (= +24973₁₀)

```
0x0000618D 0x0000_0000_0000_0110_0001_1000_1101

0xFFFF9E72 0x1111_1111_1111_1111_1001_1110_0111_0010 (1's complement)

0xFFFF9E73 0x1111_1111_1111_1111_1001_1110_0111_0011 (2's complement)

F F F F S E 7 3
```

https://stackoverflow.com/questions/47333458/assembly-x86-64-setting-carry-flag-fe

Young W. Lim Carry and Overflow 2024-04-27 Sat 22 / 112

= 900 €

Unsigned subtraction

 0x0000195D - 0x0000618D : unsigned subtraction subtraction by hand

```
0 0 0 0 1 9 5 D
0x0000195D 0x0000_0000_0000_0001_1001_0101_1101
0x0000618D 0x0000_0000_0000_00110_0001_1000_1101
0 0 0 0 6 1 8 D
0xFFFFB7D0 1 0x1111_1111_1111_1111_1011_0111_1101_0000 (hand subtraction)
1 F F F F B 7 D 0

V borrow (CF=1) : unsigned integer overflow
```

- A borrow is indicated by the carry flag (CF=1)
 - whenever an unsigned integer overflow happened
 - A B, when A < B, for non-negative integers A, B

Signed subtraction

• 0x0000195D + (-0x0000618D) : signed subtraction the *transformed addition* using the 2's complement of <u>subtrahend</u>

- signed integer overflow is indicated by the overflow flag (OF)
 - the carry flag is set by the inverted carry of a transformed addition

4 D > 4 A > 4 E > 4 E > 9 Q P

Interpreting the result as a signed or an unsigned integer

- subtracting 0x0000618D from 0x0000195D
 the results of unsigned and signed subtractions have
 the same bit pattern 0xFFFFB7D0
- the 2's complement of 0x00004830: 0xFFFFB7D0 (= -18480₁₀) 0 0 0 0 4 8 3 0 0x00004830 0x0000_0000_0000_0100_1000_0011_0000 0xFFFFB7CF 0x1111_1111_1111_1111_0111_1110_11111 (1's complement) 0xFFFFB7D0 0x1111_1111_1111_1111_1011_0111_1101_0000 (2's complement) F F F F B 7 D 0

https://stackoverflow.com/questions/47333458/assembly-x86-64-setting-carry-flag-fe

Summary of signed and unsigned subtractions (1)

- subtracting 0x0000618D from 0x0000195D
 - 0x0000195D 0x0000618D : unsigned integer subtraction hand subtraction
 - 0x0000195D + (-0x0000618D) : signed integer subtraction the *transformed addition* using the 2's complement of the subtrahend
 - the same result : 0xFFFFB7D0 (the same bit pattern)
 - interpreting as a unsigned integer 4294948816₁₀
 0xFFFFB7D0 with a borrow (CF=1)
 - interpreting as a signed integer -18480₁₀
 -0x00004830 (meaningless CF=1)

https://stackoverflow.com/questions/47333458/assembly-x86-64-setting-carry-flag-fe

Summary of signed and unsigned subtractions (2)

0xFFFFB7D0 with CF=1	the result of unsigned subtraction with unsigned integer overflow	4294948816 ₁₀
-0x00004830	the result of signed subtraction	-18480 ₁₀

Young W. Lim Carry and Overflow 2024-04-27 Sat 27 / 112

Examples of unsigned integer overflows

- 0x0000195D 0x0000618D : unsigned subtraction
 - there is an unsigned integer overflow so the carry flag will be set (CF=1) to indicate a borrow
 - A B, when A < B, for non-negative integers A, B (unsigned integers can't be negative),

Examples of signed integer overflows

- 0x0000195D + (-0x0000618D) : signed subtraction
 - there is no signed integer overflow the overflow flag won't be set (OF=0)
 - signed overflw occurrs, in the transformed addition,
 - two *positive* numbers are added and the result is a *negative*, $(P + P \rightarrow N)$, or
 - two *negative* numbers are added and the result is a *positive*, $(N + N \rightarrow P)$

https://stackoverflow.com/questions/47333458/assembly-x86-64-setting-carry-flag-f

TOC Carry flag in unsigned and signed computations

- 2's complement numbers : 4-bit
- Addend and augend in a n-bit addition
- Full adder operation in each bit position
- Internal and external carry bits
- Addition and Subtraction
- Using the Carry Flag as a borrow

2's complement numbers : 4-bit

0111	(+7)	1000	(-8)
0110	(+6)	1001	(-7)
0101	(+5)	1010	(-6)
0100	(+4)	1011	(-5)
0011	(+3)	1100	(-4)
0010	(+2)	1101	(-3)
0001	(+1)	1110	(-2)
0000	(0)	1111	(-1)

Addend and augend in a *n*-bit addition

n	bits	addened	Α	$\{a_{n-1}, a_{n-2}, \cdots, a_1, a_0\}$
n	bits	augend	В	$\{b_{n-1},b_{n-2},\cdots,b_1,b_0\}$
(n+1)	bits	carry bits	C	$\{C_n, C_{n-1}, C_{n-2}, \cdots, C_1, C_0\}$
n	bits	sum bits	S	$\{S_{n-1}, S_{n-2}, \cdots, S_1, S_0\}$

external carry bits : C_n carry out, C_0 carry in

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt

- **◆ロト ◆御 ト ◆ 恵 ト ◆ 恵 ・ か**へで

Full adder operation in each bit position

full adder operation in the i^{th} bit position

$$\{C_{i+1},S_i\}=a_i+b_i+C_i$$

$$\begin{array}{ccc}
a_i \\
b_i \\
C_i
\end{array}$$

$$C_{i+1} S_i$$

Internal and external carry bits

$$\begin{array}{lll} \text{external carrys} & \textit{C_n output, C_0 input} \\ \text{internal carrys} & \{\textit{$C_{n-1}, C_{n-2}, \cdots \cdots, C_2, C_1$}\} \\ \text{sum bits} & \{\textit{$S_{n-1}, S_{n-2}, \cdots \cdots, S_1, S_0$}\} & \text{output} \ \end{array}$$

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt

◆□▶ ◆□▶ ◆필▶ ◆필▶ · 필 · 虳९○·

Addition and Subtraction

addition

$$\{C_n, S\} = A + B = A + B + 0$$

	a_{n-1}	a_{n-2}	 a_1	a_0
	b_{n-1}	b_{n-2}	 b_1	b_0
	C_{n-1}	C_{n-2}	 C_1	0
Cn	S_{n-1}	S_{n-2}	 S_1	<i>S</i> ₀

subtraction - transformed addition

$$\{C_n,S\}=A-B=A+\overline{B}+\mathbf{1}$$

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt

35 / 112

Using the Carry Flag as a borrow (1)

- a borrow (CF=1) occurs in the subtraction A - Bwhen b is larger than a (A < B)as unsigned numbers
- Computer hardware can detect a borrow (CF=1) in subtraction by looking at whether a carry out (Cn) occurred in the transformed addition

https://stackoverflow.com/questions/47333458/assembly-x86-64-setting-carry-flag-f-

36 / 112

Young W. Lim Carry and Overflow

Using the Carry flag as a borrow (2)

- a borrow (CF=1) occurs
 in the subtraction A B (A < B)
 as unsigned numbers
- a carry out (Cn) in the transformed addition
 - If there is no carry (Cn=0) then there is a borrow (CF=1)
 - If there is a carry (Cn=1) then there is no borrow (CF=0)
 - CF = !Cn

https://stackoverflow.com/questions/47333458/assembly-x86-64-setting-carry-flag-fe

Using the Carry Flag as a borrow (3)

- the same addition and subtraction instructions are used for both unsigned and signed integer arithmetic.
 - no special addition and subtraction instructions for unsigned and signed integer arithmetic
- the only difference is
 - which flags you test afterwards and
 - how you interpret the result

TOC Rules for the carry flag

- 2's complement numbers : 4-bit
- The 1st rule for setting the carry flag
- The 2nd rule for setting the carry flag
- Cases for <u>clearing</u> the carry flag
- Computing CF in unsigned additions and subtractions

2's complement numbers : 4-bit

0111	(+7)	1000	(-8)
0110	(+6)	1001	(-7)
0101	(+5)	1010	(-6)
0100	(+4)	1011	(-5)
0011	(+3)	1100	(-4)
0010	(+2)	1101	(-3)
0001	(+1)	1110	(-2)
0000	(0)	1111	(-1)

The 1st rule for setting the carry flag

- ① CF = 1 : carry in unsigned addition
 - the carry flag is set
 if the addition of two unsigned numbers causes
 a carry out of the most significant bits added.
 - unsigned integer overflow in unsigned addition
 - hand addition rule

The 2nd rule for setting the carry flag

- 2 CF = 1 : borrow in unsigned subtraction
 - the carry flag is <u>also</u> set
 if the <u>subtraction</u> of two <u>unsigned</u> numbers requires
 a <u>borrow</u> into the most significant bits subtracted.
 - unsigned integer overflow in unsigned subtraction
 - hand subtraction rule

Cases for clearing the carry flag (1)

- Otherwise, the carry flag is turned off (zero).
 - all three interpretations have the same CF=1, the same S=0000

unsigned addition	signed addition	signed subtraction
0111 (7) +1001 +(9)	0111 (+7) +1001 +(-7)	0111 (+7) -0111 -(+7)
10000 (16)	 10000 (0)	10000 (0)
CF=1	Cn=1 -> CF=1	Cn=1 -> CF=1
CF means 16 S = 0000	CF meaningless S = 0000	CF meaningless S = 0000
* think hand addition	* think Cn of the co CF <- Cn	rresponding addition

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt

Young W. Lim Carry and Overflow 2024-04-27 Sat 43 / 112

Cases for clearing the carry flag (2)

- Otherwise, the carry flag is turned off (zero).
 - all three interpretations have the same CF=0, the same S=1111

- 1	signed addition	signed subtraction
	0111 (+7) +1001 +(-7) 10000 (0)	0111 (+7) -0111 -(+7) 10000 (0)
	Cn=1 -> CF=1	Cn=1 -> CF=1
i	CF meaningless S = 0000	CF meaningless S = 0000
	* think Cn of the cor CF <- Cn	responding addition
		0111 (+7) +1001 +(-7) 10000 (0) Cn=1 -> CF=1 CF meaningless S = 0000

Computing CF in unsigned additions and subtractions

- Computing CF in an unsigned addition
 - do the signed addition
 - Cn is the carry out
 - $CF \leftarrow Cn$
- Computing CF in an unsigned subraction
 - do the transformed signed addition
 - do the signed addition
 - Cn is the carry out
 - $\bullet \ \texttt{CF} \leftarrow \, \texttt{!Cn}$

TOC: Method for computing the carry flag

Carry flag computation

Young W. Lim Carry and Overflow 2024-04-27 Sat 46 / 112

Carry flag computation (1)

ADD (addition)	SUB (subtraction)
$CF = C_n$	$CF = \overline{C_n}$
normal carry of a 2's complement addition $A + B = A + B + 0$	inverted carry of a transformed addition $A - B = A + \overline{B} + 1$
$\{C_n, S_{n-1}\}\$ = $a_{n-1} + b_{n-1} + c_{n-1}$	$ \{C_n, S_{n-1}\} = a_{n-1} + \overline{b_{n-1}} + c_{n-1} $

 $\verb|https://www.csie.ntu.edu.tw/^cyy/courses/assembly/12fall/lectures/handouts/lec14_information and the control of the course o$

◄□▶
4□▶
4□▶
4□▶
4□▶
4□▶
4□▶
4□▶
4□▶

Carry flag computation (2)

- In unsigned arithmetic,
 - the carry flag is used to detect overflow
 - the carry flag is used to extend n-bit result into (n+1)-bit result
 - for addition, the carry flag is a carry out
 - for subtraction, the carry flag is a borrow in
- In signed arithmetic,
 - the carry flag is useless
 - the carry flag neither detects overflow nor extends n-bit result

Carry flag computation (3)

In unsigned arithmetic,

```
Addition CF = 1 means carry out when Cn = 1
Subtraction CF = 1 means borrow in when Cn = 0
```

- CF Carry Flag in x86
- Cn the normal carry out
 - the <u>carry out</u> of a 2's complement addition for <u>ADD</u>
 - the carry out of a transformed addition for SUB
- In signed arithmetic,
 - the carry flag is useless

TOC: More examples of the carry flag

- Summary I
- Summary II
- Cases for <u>setting</u> the carry flag
- Cases for <u>clearing</u> the carry flag

Summary I

unsign	ed add,	/sub		signed	additio	n 	signed	subtrac	tion	CF	OF
1101	(13)		1	1101	(-3)		1101	(-3)		1	
+1110	+(14)	ADD	-1	+1110	+(-2)	ADD	-0010	-(+2)			
			-								
11011	(11)	(+16)	- [11011	(-5)		11011	(-5)		1	0
0011	(3)			0011	(+3)		0011	(+3)			
-1110	-(14)	SUB	i	+0010	+(+2)		-1110	-(-2)	SUB	i	
			i							i	
10101	(5)	(-16)	į	00101	(+5)		00101	(+5)		1	0
0011	(3)		i	0011	(+3)		0011	(+3)		i	
+0010	+(2)	ADD	- 1	+0010	+(+2)	ADD	-1110	-(-2)		1	
										1	
00101	(5)	(+ 0)	1	00101	(+5)		00101	(+5)		0	0
1101	(13)			1101	(-3)		1101	(-3)		1	
-0010	, , ,	SUB			+(-2)		-0010	, . ,	SUB	1	
-0010	-(2)	DUD	-	+1110	+(-2)		-0010	-(+2)	DOD	-	
44044	(44)	(10)	- !	44044	(5)		44044	(5)		1 ^	^
11011	(11)	(-16)	ı	11011	(-5)		11011	(-5)		1 0	0

Summary II

		 /b							OF.	
unsign	ea aaa,	sub	signed	additio	n 	signed	subtrac	tion	CF	UF
1011	(11)		1011	(-5)		1011	(-5)			
+1100	+(12)	ADD	+1100	+(-4)	ADD	-0100	-(+4)			
		(, , 0)								
10111	(7)	(+16)	10111	(+7)		10111	(+7)		1	1
0101	(5)		0101	(+5)		0101	(+5)			
-1100	-(12)	SUB	+0100	+(+4)		-1100	-(-4)	SUB		
			l							
11001	(9)	(-16)	01001	(-7)		01001	(-7)		1	1
0101	(5)		 0101	(+5)		0101	(+5)			
+0100	, ,	ADD	l +0100		ADD		-(-4)			
01001	(9)	(+ 0)	01001	(-7)		01001	(-7)		0	1
			1							
1011	(11)		1011	(-5)		1011	(-5)			
-0100	-(4)	SUB	+1100	+(-4)		-0100	-(+4)	SUB		
			l							
00111	(7)	(0)	10111	(+7)		10111	(+7)		0	1

◆□▶ ◆□▶ ◆■▶ ◆■▶ ● からぐ

Cases for setting the carry flag (1) CF=1, OF=0

• unsigned integer overflow (CF=1 means +16)

* unsigned addition	* signed addition	signed subtraction
1101 (13) +1110 +(14) ADD	1101 (-3) +1110 +(-2) ADD	1101 (-3) -0010 -(+2)
11011 (11) (+16)	 11011 (-5)	11011 (-5)
CF=1	Cn=1 -> CF=1	Cn=1 -> CF=1
CF means 16 S = 0000	CF meaningless S = 0000	CF meaningless S = 0000
* think hand addition	* think Cn of the corre CF <- Cn (for unsign	

^{*} CF=1, S=1011, OF=0 for all three interpretations

Cases for setting the carry flag (2) CF=1, OF=0

• unsigned integer overflow (CF=1 means -16)

unsigned subtraction	signed addition	* signed subtraction
0011 (3) -1110 -(14) SUB	0011 (+3) +0010 +(+2)	0011 (+3) -1110 -(-2) SUB
10101 (5) (-16)	 00101 (+5)	00101 (+5)
CF=1	Cn=0 -> CF=1	Cn=0 -> CF=1
CF means -16 S = 0101	CF meaningless S = 0101	CF meaningless S = 0101
think hand subtraction	* think Cn of the tra CF <- !Cn (for uns	

^{*} CF=1, S=0101, OF=0 for all three interpretations

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt

54 / 112

Young W. Lim Carry and Overflow 2024-04-27 Sat

Cases for setting the carry flag (3) CF=1, OF=1

• unsinged integer overflow (CF=1 means +16)

* unsigned addition	* signed addition	signed subtraction
1011 (11) +1100 +(12) ADD	1011 (-5) +1100 +(-4) ADD	1011 (-5) -0100 -(+4)
10111 (7) (+16)	10111 (+7)	10111 (+7)
CF=1	Cn=1 -> CF=1	Cn=1 -> CF=1
CF means +16 S = 0111	CF meaningless S = 0111	CF meaningless S = 0111
* think hand addition	* think Cn of the corres CF <- Cn (for unsigne	. 0

^{*} CF=1, S=0111, OF=1 for all three interpretations

Cases for setting the carry flag (4) CF=1, OF=1

• unsinged integer overflow (CF=1 means -16)

* unsigned subtraction	I	signed addition	* signed subtraction
0101 (5) -1100 -(12) SUB		0101 (+5) +0100 +(+4)	0101 (+5) -1100 -(-4) SUB
11001 (9) (-16)		01001 (-7)	01001 (-7)
CF=1		Cn=0 -> CF=1	Cn=0 -> CF=1
CF means -16 S = 1001	i I	CF meaningless S = 1001	CF meaningless S = 1001
* think hand subtraction	: :	* think Cn of the transfo CF <- !Cn (for unsigned	

^{*} CF=1, S=1001, OF=1 for all three interpretations

Cases for clearing the carry flag (1) CF=0, OF=0

• no unsigned integer overflow (CF=0)

* unsigned addition	* signed addition	signed subtraction
0011 (3) +0010 +(2) ADD	0011 (+3) +0010 +(+2) ADD	0011 (+3) -1110 -(-2)
00101 (5) (+ 0)	00101 (+5)	00101 (+5)
CF=0	Cn=0 -> CF=0	Cn=0 -> CF=0
CF means 0 S = 0101	CF meaningless S = 0101	CF meaningless S = 0101
* think hand addition	* think Cn of the corres	• 0

^{*} CF=0, S=0101, OF=0 for all three interpretations

Cases for clearing the carry flag (2) CF=0, OF=0

• no unsigned integer overflow (CF=0)

* unsigned addition	* signed addition	signed subtraction
1101 (13) -0010 -(2) SUB	1101 (-3) +1110 +(-2)	1101 (-3) -0010 -(+2) SUB
11011 (11) (-16)	 11011 (-5)	11011 (-5)
CF=0	Cn=0 -> CF=0	Cn=0 -> CF=0
CF means 0 S = 1011	CF meaningless S = 1011	CF meaningless S = 1011
* think hand subtraction	* think Cn of the corre CF <- Cn (for unsign	

^{*} CF=0, S=1011, OF=0 for all three interpretations

Cases for clearing the carry flag (3) CF=0, OF=1

• no unsinged integer overflow (CF=0)

* unsigned addition	* signed addition	signed subtraction
0101 (5) +0100 +(4) ADD	0101 (+5) +0100 +(+4) ADD	0101 (+5) -1100 -(-4)
01001 (9) (+ 0)	01001 (-7)	01001 (-7)
CF=0	Cn=0 -> CF=0	Cn=0 -> CF=0
CF means +0 S = 1001	CF meaningless S = 1001	CF meaningless S = 1001
* think hand addition	* think Cn of the corre CF <- Cn (for unsign	

^{*} CF=0, S=1001, OF=1 for all three interpretations

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt

59 / 112

Cases for clearing the carry flag (4) CF=0, OF=1

• no unsinged integer overflow (CF=0)

unsigned subtraction		signed addition	* signed subtraction
1011 (11) -0100 -(4) SUB	 	1011 (-5) +1100 +(-4)	1011 (-5) -0100 -(+4) SUB
00111 (7) (0)	 	10111 (+7)	10111 (+7)
CF=0	į	Cn=1 -> CF=0	Cn=1 -> CF=0
CF means 0 S = 0111	i I	CF meaningless S = 0111	CF meaningless S = 0111
think hand subtraction		* think Cn of the tra CF <- !Cn (for unsi	

^{*} CF=0, S=0111, OF=1 for all three interpretations

TOC: Overflow flag

- Overflow flag in unsigned and signed computations
- Rules for the overflow flag
- Method 1 for computing the overflow flag
- Method 2 for computing the overflow flag
- More examples of the overflow flag

TOC Overflow flag in unsigned and signed computations

Overflow flag

Young W. Lim Carry and Overflow 2024-04-27 Sat 62 / 112

Overflow flag (1)

 only need to look at the sign bits (leftmost) of the three numbers

```
augend + addend = sum
minuend - subrahend = difference
```

to decide if the overflow flag is turned on or off.

overflow flag is based on signed arithmetic

Overflow flag (2)

- in signed arithmetic,
 - watch the overflow flag to detect errors
 - overflow flag on means the result is wrong
 - errors can be detected by examining the <u>sign</u> of the result, in the 2's complement arithmetic
- in unsigned arithmetic,
 - the overflow flag tells you <u>nothing</u> interesting

Overflow flag (3)

- when two positive numbers are added
 - ullet if the result is a negative, (P + P ightarrow N), then $\underline{\text{overflow}}$
 - ullet if the result is a positive, (P + P ightarrow P), then \underline{no} overflow
- when two negative numbers are added
 - ullet the result is a positive, (N + N ightarrow P), then overflow
 - the result is a negative, (N + N \rightarrow N), then <u>no</u> <u>overflow</u>

Overflow flag (4)

- adding <u>negative</u> and <u>positive</u> numbers <u>cannot</u> be <u>wrong</u>, because the sum is between the addends.
 - mixed-sign addition never turns on the overflow flag.
 - opposite signed numbers are added, then no overflow
 - both of the <u>addends</u> lies in the <u>allowable range</u> of numbers, and their <u>sum</u> is between the addends, therefore the <u>sum</u> lies also in the allowable range
 - $(P + N \rightarrow P)$ no overflow
 - $(P + N \rightarrow N)$ no overflow
 - $(N + P \rightarrow P)$ no overflow
 - $(N + P \rightarrow N)$ no overflow

TOC Rules for the overflow flag

- the 1st rule for setting OF
- the 2nd rule for setting OF
- ullet cases for clearing OF $(1\sim6)$

Young W. Lim

Overflow flag setting and clearing conditions

ADD conditions SUB conditions

0F=1	P+P o N	$P - N \rightarrow N$	$C_n \bigoplus C_{n-1} = 1$
0F=1	$N + N \rightarrow P$	N-P o P	$C_n \bigoplus C_{n-1} = 1$
0F=0	P+P o P	$P - N \rightarrow P$	$C_n \bigoplus C_{n-1} = 0$
0F=0	$N + N \rightarrow N$	N-P o N	$C_n \bigoplus C_{n-1} = 0$
0F=0	$P + N \rightarrow P$	$P - N \rightarrow P$	$C_n \bigoplus C_{n-1} = 0$
0F=0	$P + N \rightarrow N$	P-P o N	$C_n \bigoplus C_{n-1} = 0$
0F=0	$N + P \rightarrow P$	$N - N \rightarrow P$	$C_n \bigoplus C_{n-1} = 0$
0F=0	$N + P \rightarrow N$	$N - P \rightarrow N$	$C_n \bigoplus C_{n-1} = 0$

$$+P = -(-P) = -N$$

 $+N = -(-N) = -P$

The 1st rule for setting the overflow flag

If the sum of two signed numbers with the sign bits off (0, 0) yields a result number with the sign bit on (1) the overflow flag is turned on (OF =1 : P + P → N)

signed	addition	signed sub	traction	unsign	ed addition
0100	carries				
0100	(+4)	0100 (+	4)	0100	(4)
+0100	+(+4)	-1100 -(-	4)	+0100	+(4)
01000	(-8)	01000 (-	8)	01000	(8)

• OF =
$$C_n \bigoplus C_{n-1} = C_4 \bigoplus C_3 = 0 \bigoplus 1 = 1$$



The 2nd rule for setting the overflow flag

② If the sum of two numbers with the sign bits on (1, 1) yields a result number with the sign bit off (0) the overflow flag is turned on. (OF =1: N + N → P)

signed addition	signed subtraction	unsigned addition
1001 carries 1001 (-7) +1001 +(-7)	1001 (-7) -0111 -(+7)	1001 (9) +1001 +(9)
10010 (2)	10010 (2)	10010 (18)

• OF = $C_n \bigoplus C_{n-1} = C_4 \bigoplus C_3 = 1 \bigoplus 0 = 1$

Cases for clearing the overflow flag (1)

• overflow flag is turned off. (OF = 0 : $P + P \rightarrow P$)

signed	addition	signed	subtraction	unsign	ed additi	on
0011	carries					
0011	(+3)	0011	(+3)	0011	(3)	
+0011	+(+3)	-1101	-(-3)	+0011	+(3)	
00110	(+6)	00110	(+6)	00110	(6)	

• OF = $C_n \bigoplus C_{n-1} = C_4 \bigoplus C_3 = 0 \bigoplus 0 = 0$

Cases for clearing the overflow flag (2)

• overflow flag is turned off. (OF = 0 : $N + N \rightarrow N$)

signed	l addition	signed subtraction	unsigned addition
1101	carries		
1101	(-3)	1101 (-3)	1101 (13)
+1101	+(-3)	-0011 -(+3)	+1101 +(13)
11010	(-6)	11010 (-6)	11010 (26)

• OF = $C_n \bigoplus C_{n-1} = C_4 \bigoplus C_3 = 1 \bigoplus 1 = 0$

Cases for clearing the overflow flag (3)

• overflow flag is turned off. (OF = 0 : $P + N \rightarrow P$)

signed	addition	signed	subtraction	unsigne	ed addition
1100	carries				
0100	(+4)	0100	(+4)	0100	(4)
+1101	+(-3)	-0011	-(+3)	+1101	+(13)
10001	(+1)	10001	(+1)	10001	(17)

• OF = $C_n \bigoplus C_{n-1} = C_4 \bigoplus C_3 = 1 \bigoplus 1 = 0$

Cases for clearing the overflow flag (4)

• overflow flag is turned off. (OF = 0 : $P + N \rightarrow N$)

signed	addition	signed	subtraction	unsigne	ed addition	1
0000	carries					
0011	(+3)	0011	(+3)	0011	(3)	
+1100	+(-4)	-0100	-(+4)	+1100	+(12)	
01111	(-1)	01111	(-1)	01111	(15)	

• OF = $C_n \bigoplus C_{n-1} = C_4 \bigoplus C_3 = 0 \bigoplus 0 = 0$

Cases for clearing the overflow flag (5)

• overflow flag is turned off. (OF = 0 : $N + P \rightarrow P$)

signed addition	signed subtraction	unsigned addition
1100 carries		
1101 (-3)	0011 (-3)	1101 (13)
+0100 (+4)	-1100 -(-4)	+0100 +(4)
10001 (+1)	10001 (+1)	10001 (17)

• OF =
$$C_n \bigoplus C_{n-1} = C_4 \bigoplus C_3 = 1 \bigoplus 1 = 0$$

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt

Young W. Lim

Cases for clearing the overflow flag (6)

• overflow flag is turned off. (OF = 0 : $N + P \rightarrow N$)

signed	d addition	signed	d subtraction	unsign	ed addition
0000	carries				
1100	(-4)	0100	(-4)	1100	(12)
+0011	+(+3)	-1101	-(-3)	+0011	+(3)
01111	(-1)	01111	(-1)	01111	(15)

• OF =
$$C_n \bigoplus C_{n-1} = C_4 \bigoplus C_3 = 0 \bigoplus 0 = 0$$

TOC Method 1 for computing the overflow flag

- Adding two numbers with the same sign
- Overflow conditions for additions and subtractions
- Overflow condition for an addition
- Overflow conditions for a subtraction
- Overflow in signed computations

Adding two numbers with the same sign

- overflow can only happen when adding two numbers of the same sign results in a different sign $(P + P \rightarrow N, N + N \rightarrow P)$
- *n*-bit signed binary arithmetic A + B = C

$$A = (a_{n-1}, \dots, a_1, a_0)$$

$$B = (b_{n-1}, \dots, b_1, b_0)$$

$$C = (c_{n-1}, \dots, c_1, c_0)$$

- to detect overflow
 - only the sign bits are considered
 - msb (most significant bit) $a_{n-1}, b_{n-1}, c_{n-1}$
 - the other bits are ignored



Overflow conditions for additions and subtractions

- with two operands (A and B) and one result (C), three sign bits $(a_{n-1}, b_{n-1}, c_{n-1})$ are considered $\rightarrow 2^3 = 8$ possible combinations
- only two cases result in overflow for an addition
 - 0 0 1 $(p+p \rightarrow n)$
 - 1 1 0 $(n+n \to p)$
- only two cases are considered as overflow for an subtraction
 - 0 1 1 $(p-n \to n)$
 - 100 $(n-p \to p)$



Overflow condition for an addition

• Overflow in an addition (A + B)

	a_{n-1}	b_{n-1}	c_{n-1}	
	0	0	0	$\overline{p+p o p}$
OVER	0	0	1	p+p ightarrow n
	0	1	0	$p + n \rightarrow p$
	0	1	1	$p + n \rightarrow n$
	1	0	0	n+p o p
	1	0	1	n+p o n
OVER	1	1	0	$n + n \rightarrow p$
	1	1	1	$n + n \rightarrow n$

- adding two positives should be positive
- adding two negatives should be negative



Overflow conditions for a subtraction

• Overflow in a subtraction (A - B)

	a_{n-1}	b_{n-1}	c_{n-1}	
	0	0	0	p-p o p
	0	0	1	p-p ightarrow n
	0	1	0	p-n o p
OVER	0	1	1	p-n ightarrow n
OVER	1	0	0	n-p o p
	1	0	1	n-p o n
	1	1	0	n-n o p
	1	1	1	$n-n \rightarrow n$

- subtracting a negative is the same as adding a positive
- subtracting a positive is the same as adding a negative

Overflow in signed computations

- ALU might contain a small logic that <u>sets</u> the <u>overflow</u> flag to "1" if and only if any one of the above four <u>OV</u> conditions is met.
- in signed computations, <u>adding</u> two numbers of the <u>same sign</u> must produce a <u>result</u> of the <u>same sign</u>, otherwise overflow happened.

TOC Method 2 for computing the overflow flag

- Carry into and carry out of the sign bit
- Overflow in 2's complement arithmetic
- Overflow flag = $C_n \bigoplus C_{n-1}$
- Examples of 4-bit signed additions
- C_n and C_{n-1} in a *n*-bit addition
- Overflow flag computation
- Examples of computing overflow flag
- Hexadecimal carry, octal carry, decimal carry
- No carry into the sign bit

Carry into and carry out of the sign bit

- When adding two n-bit binary values, consider
 - the carry coming into the most significant bit (msb)
 C_{n-1}: carry into the sign bit
 - the carry going out of the most significant bit (msb)
 C_n: carry out of the sign bit
 this is the carry flag (CF) in the processor

Overflow in 2's complement arithmetic

- overflow in 2's complement happens (OF=1) when
 - there is a carry into the sign bit $(C_{n-1} = 1)$ but no carry out of the sign bit $(C_n = 0)$
 - there is no carry into the sign bit $(C_{n-1} = 0)$ but a carry out of the sign bit $(C_n = 1)$

Overflow flag = $C_n \bigoplus C_{n-1}$

- the overflow flag is the XOR $(C_n \bigoplus C_{n-1})$ of
 - of the carry coming into the sign bit (C_{n-1})
 - with the carry going out of the sign bit (C_n)
- overflow happens when the carry in (C_{n-1}) does not equal to the carry out (C_n)

Examples of 4-bit signed additions (1)

4-bit 2's complement addition examples

```
0000
                                0100
0100 (+4) (pos sign 0)
                                0100 (+4) (pos sign 0)
+1000 (-8) (neg sign 1)
                                +0100 (+4) (pos sign 0)
01100 (-4) (neg sign 1)
                                01000 (-8) (neg sign 1)
C4 carry out 0 (1+0+0)
                               C4 carry out 0 (0+0+1)
C3 carry in 0 (0+1+0)
                                C3 carry in 1 (1+1+0)
O XOR O = NO OVERFI.OW
                                O XOR 1 = OVERFI.OW!
1100
                                1000
1100 (-4) (neg sign 1)
                               1100 (-4) (neg sign 1)
+0100 (+4) (pos sign 0)
                                +1000 (-8) (neg sign 1)
10000 (0) (pos sign 0)
                                10100 (+4) (pos sign 0)
C4 carry out 1 (1+0+1)
                                C4 carry out 1 (1+1+0)
C3 carry in 1 (1+1+0)
                                C3 carry in 0 (1+0+0)
1 XOR 1 = NO OVERFI.OW
                                1 XOR O = OVERFI.OW!
```

Examples of 4-bit signed additions (2)

ullet same sign addition o possible overflow

+ +, -	, +	+ +, +	, -
+5	-5	+5	-5
+5	-5	+1	-1
-6(0	F) +6(0	F) +6	-6
0101	1011	0001	1111
0101	1011	0101	1011
0101	1011	0001	1111
01010	10110	00110	11010
C4 = 0	C4 = 1	C4 = 0	C4 = 1
C3 = 1	C3 = 0	C3 = 0	C3 = 1
OF = 1	OF = 1	OF = O	OF = O

Examples of 4-bit signed additions (3)

ullet mixed sign addition o no overflow

+ -, +	+ -, -	- +, +	- +, -
+5	+5	-5	-5
-1	-6	+6	+1
+4	-1	+1	-4
1111	0000	1110	0011
0101	0101	1011	1011
1111	1010	0110	0001
10100	01111	10001	01100
C4 = 1	C4 = 0	C4 = 1	C4 = 0
C3 = 1	C3 = 0	C3 = 1	C3 = 0
OF = O	OF = O	OF = O	OF = O

$\overline{C_n}$ and $\overline{C_{n-1}}$ in a *n*-bit addition

$(n-1)^{th}$ bit – MSB

- adding operations at the (n-1) bit position
- $\{C_n, S_{n-1}\} = a_{n-1} + b_{n-1} + c_{n-1}$

C_n:carry coming out of the msb

$$(n-2)^{th}$$
 bit

- adding operations at the (n-2) bit position
- $\{C_{n-1}, S_{n-2}\} = a_{n-2} + b_{n-2} + c_{n-2}$

$$\begin{array}{ccc}
 & a_{n-2} \\
 & b_{n-2} \\
 & C_{n-2}
\end{array}$$
 $\begin{array}{ccc}
 & C_{n-1} & S_{n-2}
\end{array}$

● C_{n-1}: carry coming *into* the msb

Overflow flag computation

ADD (addition)	SUB (subtraction)
OF = $C_n \bigoplus C_{n-1}$	OF = $C_n \bigoplus C_{n-1}$
a 2's complement addition $A + B = A + B + 0$ ($C_0 = 0$)	the transformed addition $A - B = A + \overline{B} + 1 \ (C_0 = 1)$
${C_n, S_{n-1}}$ = $a_{n-1} + b_{n-1} + c_{n-1}$	${C_n, S_{n-1}} = a_{n-1} + \overline{b_{n-1}} + c_{n-1}$
${C_{n-1}, S_{n-2}}$ = $a_{n-2} + b_{n-2} + c_{n-2}$	$ \{C_{n-1}, S_{n-2}\} $ $= a_{n-2} + \overline{b_{n-2}} + c_{n-2} $



Hexadecimal carry, octal carry, decimal carry

- Note that this XOR method only works with the binary carry that goes into the sign bit.
- not works with hexadecimal carry decimal carry, octal carry
 - the carry doesn't go into the sign bit
 - can't XOR that non-binary carry with the outgoing carry.

No carry into the sign bit

Hexadecimal addition example
 (showing that XOR doesn't work for hex carry):
 ^{8Ah}
 +8Ah
 ====
 114h

- The hexadecimal carry of 1 resulting from A+A does not affect the sign bit.
- If you do the math in binary, you'll see
 that there is no carry into the sign bit;
 but, there is carry out of the sign bit.
 Therefore, the above example sets OVERFLOW on.
 (The example adds two negative numbers and gets a positive number.)

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt

Young W. Lim Carry and Overflow 2024-04-27 Sat 93 / 112

Summary I

unsigned ac	ld/sub	signed	additio	n 	signed	subtrac	tion	CF	OF
1101 (13	3)	1101	(-3)		1101	(-3)		l	
+1110 +(14	1) ADD	+1110	+(-2)	ADD	-0010	-(+2)			
11011 (1:	 l) (+16)	1 11011	(-5)		11011	(-5)		 1	0
0011	2)		(+2)		0011	(.2)			
0011 (3	-	0011				(+3)	~	!	
-1110 -(14	1) SUB	+0010	+(+2)		-1110	-(-2)	SUB	<u> </u>	
								l	
10101 (5	5) (-16)	00101 	(+5)		00101	(+5)		1 	0
0011 (3	3)	0011	(+3)		0011	(+3)		İ	
+0010 +(2	2) ADD	+0010	+(+2)	ADD	-1110	-(-2)		l	
		l						l	
00101 (9	5) (+ 0)	00101	(+5)		00101	(+5)		0	0
1101 (13	3)	l l 1101	(-3)		1101	(-3)		 	
•	2) SUB		+(-2)		-0010		SUB	l	
								l	
11011 (1:	1) (-16)	11011	(-5)		11011	(-5)		I 0	0

Summary II

		/ l								
unsign	iea aaa,	/sub 	signed	additio	n 	signea	subtrac	tion	CF	10
1011	(11)		1011	(-5)		1011	(-5)			
+1100	+(12)	ADD	+1100	+(-4)	ADD	-0100	-(+4)			
10111	(7)	(116)		(+7)		10111	(+7)		 1	1
10111	(1)	(+16)	10111	(+1)		10111	(+1)		 1	1
0101	(5)		0101	(+5)		0101	(+5)		İ	
-1100	-(12)	SUB	+0100	+(+4)		-1100	-(-4)	SUB		
11001	(9)	(-16)	01001	(-7)		01001	(-7)		1 	1
0101	(5)		0101	(+5)		0101	(+5)			
+0100	+(4)	ADD	+0100	+(+4)	ADD	-1100	-(-4)			
			l						l	
01001	(9)	(+ 0)	01001 	(-7)		01001	(-7)		0 	1
1011	(11)		1011	(-5)		1011	(-5)		İ	
-0100	-(4)	SUB	+1100	+(-4)		-0100	-(+4)	SUB		
	(7)	(0)		(.7)		40444	(.7)			
00111	(7)	(0)	10111	(+7)		10111	(+7)		0	1

◆□▶ ◆□▶ ◆■▶ ◆■▶ ● からぐ

95 / 112

Cases for setting the overflow flag (1) CF=1, OF=1

singed integer overflow (OF=1 means incorrect S)

* unsigned addition	* signed addition	signed subtraction	
1011 (11) +1100 +(12) ADD	1000 1011 (-5) +1100 +(-4) ADD	1011 (-5) -0100 -(+4)	
10111 (7) (+16)	10111 (+7)	10111 (+7)	
0F=1	n + n -> p (OF=1)	n - p -> p (OF=1)	
OF meaningless	-> incorrect S	-> incorrect S	
S = 0111	S = 0111	S = 0111	
* think hand addition	* 0F <- C4 XOR C3 = 1 XO		

 $\boldsymbol{*}$ CF=1, S=0111, OF=1 for all three interpretations

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt = > = > > = > > <

Young W. Lim Carry and Overflow 2024-04-27 Sat 96 / 112

Cases for setting the overflow flag (2) CF=1, OF=1

singed integer overflow (OF=1 means incorrect S)

* unsigned subtraction	Ι	signed addition	* signed subtraction
0101 (5) -1100 -(12) SUB 11001 (9) (-16)	 	0100 0101 (+5) +0100 +(+4) 01001 (-7)	0101 (+5) -1100 -(-4) SUB 01001 (-7)
0F=1		p + p -> n (0F=1)	p - n -> n (0F=1)
OF meaningless S = 1001	 	-> incorrect S S = 1001	-> incorrect S S = 1001
* think hand subtraction	* 	of signed additio	

* CF=1, S=1001, OF=1 for all three interpretations

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt = > = > > = > > <

Young W. Lim Carry and Overflow 2024-04-27 Sat 97 / 112

Cases for setting the overflow flag (3) CF=0, OF=1

• singed integer overflow (OF=1 means incorrect S)

* unsigned addition	* signed addition	signed subtraction
0101 (5) +0100 +(4) ADD	0100 0101 (+5) +0100 +(+4) ADD	0101 (+5) -1100 -(-4)
01001 (9) (+ 0)	 01001 (-7)	01001 (-7)
0F=1	p + p -> n (0F=1)	p - n -> n (0F=1)
OF meaningless S = 1001	-> incorrect S S = 1001	-> incorrect S S = 1001
* think hand addition	* OF <- C4 XOR C3 = O XOR of signed additio	

 $\boldsymbol{*}$ CF=0, S=1001, OF=1 for all three interpretations

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt = > = > > = > > <

Young W. Lim Carry and Overflow 2024-04-27 Sat 98 / 112

Cases for setting the overflow flag (4) CF=0, OF=1

• singed integer overflow (OF=1 means incorrect S)

* unsigned subtraction	 I	signed addition	* signed subtraction
1011 (11) -0100 -(4) SUB 00111 (7) (0)	 	1000 1011 (-5) +1100 +(-4) 10111 (+7)	1011 (-5) -0100 -(+4) SUB 10111 (+7)
0F=1	 	n + n -> p (OF=1)	n - p -> p (0F=1)
OF meaningless S = 0111	 	-> incorrect S S = 0111	-> incorrect S S = 0111
* think hand subtraction	* 	OF <- C4 XOR C3 = 1 XOR of signed addition	· -

 \ast CF=0, S=0111, OF=1 for all three interpretations

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt

99 / 112

Cases for clearing the overflow flag (1) CF=1, OF=0

• no signed integer overflow (CF=0 means correct S)

unsigned addition	* signed addition	signed subtraction	
	1100		
1101 (13)	1101 (-3)	1101 (-3)	
+1110 +(14) ADD	+1110 +(-2) ADD	-0010 -(+2)	
11011 (11) (+16)	11011 (-5) 	11011 (-5)	
OF=0	n + n -> n (OF=0)	n - p -> n (OF=0)	
OF meaningless	-> correct S	-> correct S	
S = 0000	S = 0000	S = 0000	
think hand addition	* OF <- C4 XOR C3 = 1 XO of signed additi		

* CF=1, S=1011, OF=0 for all three interpretations

Cases for clearing the overflow flag (2) CF=1, OF=0

• no signed integer overflow (CF=0 means correct S)

unsigned subtraction			* signed subtractio
		0010	
0011 (3)	1	0011 (+3)	0011 (+3)
-1110 -(14) SUB		+0010 +(+2)	-1110 -(-2) SUE
10101 (5) (-16)	 	00101 (+5)	00101 (+5)
CF=1	 	p + p -> p (0F=0)	p - n -> p (OF=0)
OF meaningless	ĺ	-> correct S	-> correct S
S = 0101	İ	S = 0101	S = 0101
think hand	*	OF <- C4 XOR C3 = 0 XO	R 0 = 0
subtraction	1	of signed additi	on

 $\boldsymbol{*}$ CF=1, S=0101, OF=0 for all three interpretations

Cases for clearing the overflow flag (3) CF=0, OF=0

• no signed integer overflow (CF=0 means correct S)

unsigned addition	* signed addition	signed subtractio
	0010	
0011 (3)	0011 (+3)	0011 (+3)
+0010 +(2) ADD	+0010 +(+2) ADD	-1110 -(-2)
00101 (5) (+ 0)	00101 (+5)	00101 (+5)
0F=0	p + p -> p (OF=0)	p - n -> p (OF=0)
OF meaningless	-> correct S	-> correct S
S = 0101	S = 0101	S = 0101
think hand	* OF <- C4 XOR C3 = 0 X	OR 0 = 0
addition	of signed addit	ion

 $\boldsymbol{*}$ CF=0, S=0101, OF=0 for all three interpretations

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt

Young W. Lim Carry and Overflow 2024-04-27 Sat 102 / 112

Cases for clearing the overflow flag (4) CF=0, OF=0

• no signed integer overflow (CF=0 means correct S)

unsigned addition	* signed addition	signed subtraction
1101 (13) -0010 -(2) SUB	1100 1101 (-3) +1110 +(-2)	1101 (-3) -0010 -(+2) SUB
11011 (11) (-16)	 11011 (-5)	11011 (-5)
0F=0	n + n -> n (OF=0)	n - p -> n (0F=0)
OF meaningless S = 1011	-> correct S S = 1011	-> correct S S = 1011
think hand subtraction	* OF <- C4 XOR C3 = 1 X	

 $\boldsymbol{*}$ CF=0, S=1011, OF=0 for all three interpretations

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt

Young W. Lim Carry and Overflow 2024-04-27 Sat 103 / 112

2's complement numbers : 4-bit

0111	(+7)	1000	(-8)
0110	(+6)	1001	(-7)
0101	(+5)	1010	(-6)
0100	(+4)	1011	(-5)
0011	(+3)	1100	(-4)
0010	(+2)	1101	(-3)
0001	(+1)	1110	(-2)
0000	(0)	1111	(-1)

Unsigned 4-bit addition table (1)

	0000	0001	0010	0011	0100	0101	0110	0111
	(0)	(1)	(2)	(3)	(4)	(5)	(6)	(7)
0000	0000	0001	0010	0011	0100	0101	0110	0111
(0)	(0)	(1)	(2)	(3)	(4)	(5)	(6)	(7)
0001	0001	0010	0011	0100	0101	0110	0111	1000
(1)	(1)	(2)	(3)	(4)	(5)	(6)	(7)	(8)
0010	0010	0011	0100	0101	0110	0111	1000	1001
(2)	(2)	(3)	(4)	(5)	(6)	(7)	(8)	(9)
0011	0011	0100	0101	0110	0111	1000	1001	1010
(3)	(3)	(4)	(5)	(6)	(7)	(8)	(9)	(10)
0100	0100	0101	0110	0111	1000	1001	1010	1011
(4)	(4)	(5)	(6)	(7)	(8)	(9)	(10)	(11)
0101	0101	0110	0111	1000	1001	1010	1011	1100
(5)	(5)	(6)	(7)	(8)	(9)	(10)	(11)	(12)
0110	0110	0111	1000	1001	1010	1011	1100	1101
(6)	(6)	(7)	(8)	(9)	(10)	(11)	(12)	(13)
0111	0111	1000	1001	1010	1011	1100	1101	1110
(7)	(7)	(8)	(9)	(10)	(11)	(12)	(13)	(14)

Unsigned 4-bit addition table (2)

	0000	0001	0010	0011	0100	0101	0110	0111
	(0)	(1)	(2)	(3)	(4)	(5)	(6)	(7)
1000	1000	1001	1010	1011	1100	1101	1110	1111
(8)	(8)	(9)	(10)	(11)	(12)	(13)	(14)	(15)
1001	1001	1010	1011	1100	1101	1110	1111	0000 CF
(9)	(9)	(10)	(11)	(12)	(13)	(14)	(15)	(0)
1010	1010	1011	1100	1101	1110	1111	0000 CF	0001 CF
(10)	(10)	(11)	(12)	(13)	(14)	(15)	(16+0)	(16+1)
1011	1011	1100	1101	1110	1111	0000 CF	0001 CF	0010 CF
(11)	(11)	(12)	(13)	(14)	(15)	(16+0)	(16+1)	(16+2)
1100	1100	1101	1110	1111	0000 CF	0001 CF	0010 CF	0011 CF
(12)	(12)	(13)	(14)	(15)	(16+0)	(16+1)	(16+2)	(16+3)
1101	1101	1110	1111	0000 CF	0001 CF	0010 CF	0011 CF	0100 CF
(13)	(13)	(14)	(15)	(16+0)	(16+1)	(16+2)	(16+3)	(16+4)
1110	1110	1111	0000 CF	0001 CF	0010 CF	0011 CF	0100 CF	0101 CF
(14)	(14)	(15)	(16+0)	(16+1)	(16+2)	(16+3)	(16+4)	(16+5)
1111	1111	0000 CF	0001 CF	0010 CF	0011 CF	0100 CF	0101 CF	0110 CF
(15)	(15)	(16+0)	(16+1)	(16+2)	(16+3)	(16+4)	(16+5)	(16+6)

Unsigned 4-bit addition table (3)

100 (8) 0000 100 (0) (8) 0001 100 (1) (9) 0010 10) (9) 00 1001) (9) 01 1010	(10) 1010 (10) 1011 (11)	1011 (11) 1011 (11) 1100 (12)	1100 (12) 1100 (12) 1101	1101 (13) 1101 (13) 1110	1110 (14) 1110 (14) 1111	1111 (15) 1111 (15) 0000 CF
0000 100 (0) (8) 0001 100 (1) (9)	00 1001) (9) 01 1010) (10)	1010 (10) 1011 (11)	1011 (11) 1100	1100 (12) 1101	1101 (13)	1110 (14)	1111 (15)
(0) (8) 0001 10 (1) (9)) (9) 01 1010) (10)	(10) 1011 (11)	(11) 1100	(12) 1101	(13)	(14)	(15)
0001 10 (1) (9)	01 1010) (10)	1011 (11)	1100	1101			
(1) (9)) (10)	(11)			1110	1111	nnnn CE
			(12)	(10)			0000 CF
0010 10	10 1011	1100		(13)	(14)	(15)	(16+0)
0010 10		1100	1101	1110	1111	0000 CF	0001 CF
(2) (10	0) (11)	(12)	(13)	(14)	(15)	(16+0)	(16+1)
0011 10	11 1100	1101	1110	1111	0000 CF	0001 CF	0010 CF
(3) (11	1) (12)	(13)	(14)	(15)	(16+0)	(16+1)	(16+2)
0100 11	00 1101	1110	1111	0000 CF	0001 CF	0010 CF	0011 CF
(4) (12	2) (13)	(14)	(15)	(16+0)	(16+1)	(16+2)	(16+3)
0101 11	01 1110	1111	0000 CF	0001 CF	0010 CF	0011 CF	0100 CF
(5) (13	3) (14)	(15)	(16+0)	(16+1)	(16+2)	(16+3)	(16+4)
0110 11	10 1111	0000 C	F 0001 CF	0010 CF	0011 CF	0100 CF	0101 CF
(6) (14	4) (15)	(16+0)	(16+1)	(16+2)	(16+3)	(16+4)	(16+5)
0111 11	11 0000	CF 0001 C	F 0010 CF	0011 CF	0100 CF	0101 CF	0110 CF
(7) (15	5) (16+	0) (16+1)) (16+2)	(16+3)	(16+4)	(16+5)	(16+6)

Unsigned 4-bit addition table (4)

	1000	1001	1010	1011	1100	1101	1110	1111
	(8)	(9)	(10)	(11)	(12)	(13)	(14)	(15)
1000	0000 CF	0001 CF	0010 CF	0011 CF	0100 CF	0101 CF	0110 CF	0111 CF
(8)	(16+0)	(16+1)	(16+2)	(16+3)	(16+4)	(16+5)	(16+6)	(16+7)
1001	0001 CF	0010 CF	0011 CF	0100 CF	0101 CF	0110 CF	0111 CF	1000 CF
(9)	(16+1)	(16+2)	(16+3)	(16+4)	(16+5)	(16+6)	(16+7)	(16+8)
1010	0010 CF	0011 CF	0100 CF	0101 CF	0110 CF	0111 CF	1000 CF	1001 CF
(10)	(16+2)	(16+3)	(16+4)	(16+5)	(16+6)	(16+7)	(16+8)	(16+9)
1011	0011 CF	0100 CF	0101 CF	0110 CF	0111 CF	1000 CF	1001 CF	1010 CF
(11)	(16+3)	(16+4)	(16+5)	(16+6)	(16+7)	(16+8)	(16+9)	(16+10)
1100	0100 CF	0101 CF	0110 CF	0111 CF	1000 CF	1001 CF	1010 CF	1011 CF
(12)	(16+4)	(16+5)	(16+6)	(16+7)	(16+8)	(16+9)	(16+10)	(16+11)
1101	0101 CF	0110 CF	0111 CF	1000 CF	1001 CF	1010 CF	1011 CF	1100 CF
(13)	(16+5)	(16+6)	(16+7)	(16+8)	(16+9)	(16+10)	(16+11)	(16+12)
1110	0110 CF	0111 CF	1000 CF	1001 CF	1010 CF	1011 CF	1100 CF	1101 CF
(14)	(16+6)	(16+7)	(16+8)	(16+9)	(16+10)	(16+11)	(16+12)	(12+13)
1111	0111 CF	1000 CF	1001 CF	1010 CF	1011 CF	1100 CF	1101 CF	1110 CF
(15)	(16+7)	(16+8)	(16+9)	(16+10)	(16+11)	(16+12)	(16+13)	(16+14)

http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt

- 4 ロ ト 4 個 ト 4 恵 ト 4 恵 ト - 恵 - か Q (C)

Signed 4-bit addition table (1)

	0000	0001	0010	0011	0100	0101	0110	0111
	(0)	(+1)	(+2)	(+3)	(+4)	(+5)	(+6)	(+7)
0000	0000	0001	0010	0011	0100	0101	0110	0111
(0)	(0)	(+1)	(+2)	(+3)	(+4)	(+5)	(+6)	(+7)
0001	0001	0010	0011	0100	0101	0110	0111	1000
(+1)	(+1)	(+2)	(+3)	(+4)	(+5)	(+6)	(+7)	(-8) OF
0010	0010	0011	0100	0101	0110	0111	1000	1001
(+2)	(+2)	(+3)	(+4)	(+5)	(+6)	(+7)	(-8) OF	(-7) OI
0011	0011	0100	0101	0110	0111	1000	1001	1010
(+3)	(+3)	(+4)	(+5)	(+6)	(+7)	(-8) OF	(-7) OF	(-6) OF
0100	0100	0101	0110	0111	1000	1001	1010	1011
(+4)	(+4)	(+5)	(+6)	(+7)	(-8) OF	(-7) OF	(-6) OF	(-5) OF
0101	0101	0110	0111	1000	1001	1010	1011	1100
(+5)	(+5)	(+6)	(+7)	(-8) OF	(-7) OF	(-6) OF	(-5) OF	(-4) OF
0110	0110	0111	1000	1001	1010	1011	1100	1101
(+6)	(+6)	(+7)	(-8) OF	(-7) OF	(-6) OF	(-5) OF	(-4) OF	(-3) OI
0111	0111	1000	1001	1010	1011	1100	1101	1110
(+7)	(+7)	(-8) OF	(-7) OF	(-6) OF	(-5) OF	(-4) OF	(-3) OF	(-2) OI

Signed 4-bit addition table (2)

	0000	0001	0010	0011	0100	0101	0110	0111
	(0)	(+1)	(+2)	(+3)	(+4)	(+5)	(+6)	(+7)
1000	1000	1001	1010	1011	1100	1101	1110	1111
(-8)	(-8)	(-7)	(-6)	(-5)	(-4)	(-3)	(-2)	(-1)
1001	1001	1010	1011	1100	1101	1110	1111	0000
(-7)	(-7)	(-6)	(-5)	(-4)	(-3)	(-2)	(-1)	(0)
1010	1010	1011	1100	1101	1110	1111	0000	0001
(-6)	(-6)	(-5)	(-4)	(-3)	(-2)	(-1)	(0)	(+1)
1011	1011	1100	1101	1110	1111	0000	0001	0010
(-5)	(-5)	(-4)	(-3)	(-2)	(-1)	(0)	(+1)	(+2)
1100	1100	1101	1110	1111	0000	0001	0010	0011
(-4)	(-4)	(-3)	(-2)	(-1)	(0)	(+1)	(+2)	(+3)
1101	1101	1110	1111	0000	0001	0010	0011	0100
(-3)	(-3)	(-2)	(-1)	(0)	(+1)	(+2)	(+3)	(+4)
1110	1110	1111	0000	0001	0010	0011	0100	0101
(-2)	(-2)	(-1)	(0)	(+1)	(+2)	(+3)	(+4)	(+5)
1111	1111	0000	0001	0010	0011	0100	0101	0110
(-1)	(-1)	(0)	(+1)	(+2)	(+3)	(+4)	(+5)	(+6)

Signed 4-bit addition table (3)

	1000	1001	1010	1011	1100	1101	1110	1111
	(-8)	(-7)	(-6)	(-5)	(-4)	(-3)	(-2)	(-1)
0000	1000	1001	1010	1011	1100	1101	1110	1111
(0)	(-8)	(-7)	(-6)	(-5)	(-4)	(-3)	(-2)	(-1)
0001	1001	1010	1011	1100	1101	1110	1111	0000
(+1)	(-7)	(-6)	(-5)	(-4)	(-3)	(-2)	(-1)	(0)
0010	1010	1011	1100	1101	1110	1111	0000	0001
(+2)	(-6)	(-5)	(-4)	(-3)	(-2)	(-1)	(0)	(+1)
0011	1011	1100	1101	1110	1111	0000	0001	0010
(+3)	(-5)	(-4)	(-3)	(-2)	(-1)	(0)	(+1)	(+2)
0100	1100	1101	1110	1111	0000	0001	0010	0011
(+4)	(-4)	(-3)	(-2)	(-1)	(0)	(+1)	(+2)	(+3)
0101	1101	1110	1111	0000	0001	0010	0011	0100
(+5)	(-3)	(-2)	(-1)	(0)	(+1)	(+2)	(+3)	(+4)
0110	1110	1111	0000	0001	0010	0011	0100	0101
(+6)	(-2)	(-1)	(0)	(+1)	(+2)	(+3)	(+4)	(+5)
0111	1111	0000	0001	0010	0011	0100	0101	0110
(+7)	(-1)	(0)	(+1)	(+2)	(+3)	(+4)	(+5)	(+6)

.....

Signed 4-bit addition table (4)

	1000	1001	1010	1011	1100	1101	1110	1111
	(-8)	(-7)	(-6)	(-5)	(-4)	(-3)	(-2)	(-1)
1000	0000	0001	0010	0011	0100	0101	0110	0111
(-8)	(0) OF	(+1) OF	(+2) OF	(+3) OF	(+4) OF	(+5) OF	(+6) OF	(+7) OF
1001	0001	0010	0011	0100	0101	0110	0111	1000
(-7)	(+1) OF	(+2) OF	(+3) OF	(+4) OF	(+5) OF	(+6) OF	(+7) OF	(-8)
1010	0010	0011	0100	0101	0110	0111	1000	1001
(-6)	(+2) OF	(+3) OF	(+4) OF	(+5) OF	(+6) OF	(+7) OF	(-8)	(-7)
1011	0011	0100	0101	0110	0111	1000	1001	1010
(-5)	(+3) OF	(+4) OF	(+5) OF	(+6) OF	(+7) OF	(-8)	(-7)	(-6)
1100	0100	0101	0110	0111	1000	1001	1010	1011
(-4)	(+4) OF	(+5) OF	(+6) OF	(+7) OF	(-8)	(-7)	(-6)	(-5)
1101	0101	0110	0111	1000	1001	1010	1011	1100
(-3)	(+5) OF	(+6) OF	(+7) OF	(-8)	(-7)	(-6)	(-5)	(-4)
1110	0110	0111	1000	1001	1010	1011	1100	1101
(-2)	(+6) OF	(+7) OF	(-8)	(-7)	(-6)	(-5)	(-4)	(-3)
1111	0111	1000	1001	1010	1011	1100	1101	1110
(-1)	(+7) OF	(-8)	(-7)	(-6)	(-5)	(-4)	(-3)	(-2)

 $\verb|http://teaching.idallen.com/dat2343/10f/notes/040_overflow.txt|$

2024-04-27 Sat