Applicatives Methods (3B)

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The definition of Applicative

class (Functor f) => Applicative f where pure :: a -> f a (<*>) :: f (a -> b) -> f a -> f b

The class has a two methods :

pure brings arbitrary values into the functor

(<*>) takes a function wrapped in a functor f and a value wrapped in a functor f and returns the result of the application which is also wrapped in a functor f f (a -> b) :: a function <u>wrapped in f</u>

f a :: a value wrapped in f

The Maybe instance of Applicative

instance Applicative Maybe where	
pure	= Just
(Just f) <*> (Just x)	= Just (f x)
_ <*> _	= Nothing

pure wraps the value with Just;

(<*>) applies

the <u>function</u> wrapped in **Just** to the <u>value</u> wrapped in **Just** if both exist, and results in **Nothing** otherwise.

https://en.wikibooks.org/wiki/Haskell/Applicative_functors

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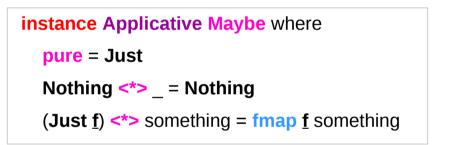
An Instance of the Applicative Typeclass

class (Functor f) => Applicative f where pure :: a -> f a (<*>) :: f (a -> b) -> f a -> f b

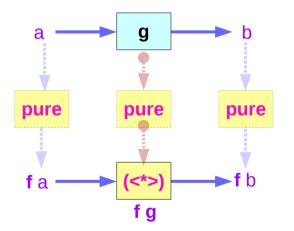


```
(Functor f) => Applicative f
```

f : Functor, Applicative



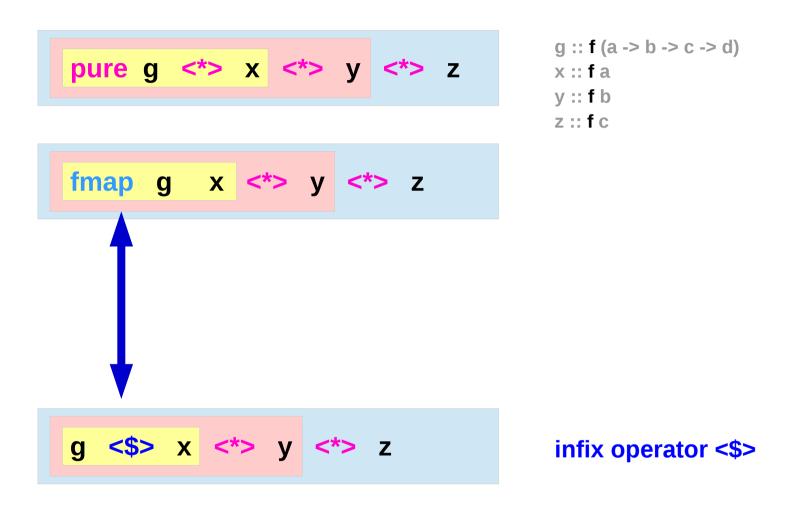
\underline{f} : function in a context



(Functor f) => Applicative f

http://learnyouahaskell.com/functors-applicative-functors-and-monoids

Left associative <*>, fmap, and <\$>

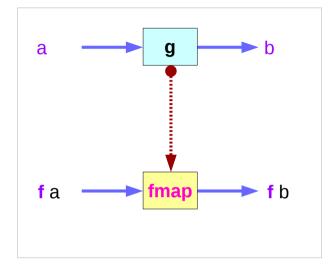


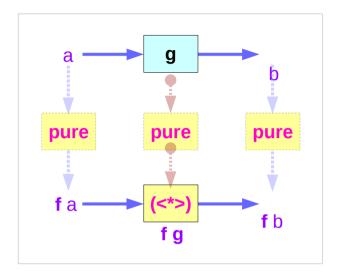
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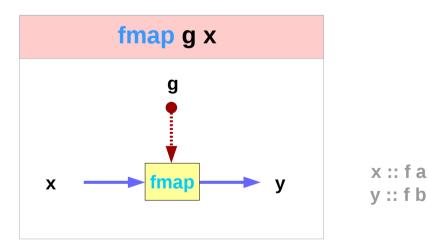
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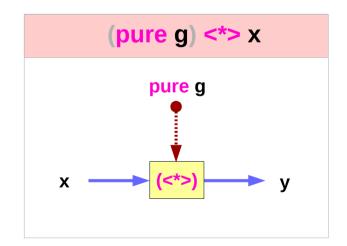
fmap g x = (pure g) < > x





pure = f





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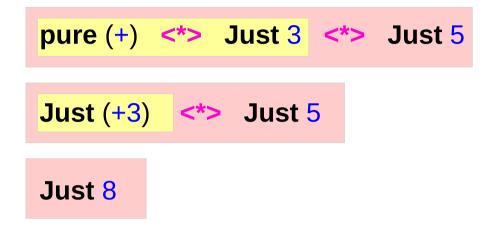
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Left Associative <*> examples

ghci> pure (+) <*> Just 3 <*> Just 5

Just 8



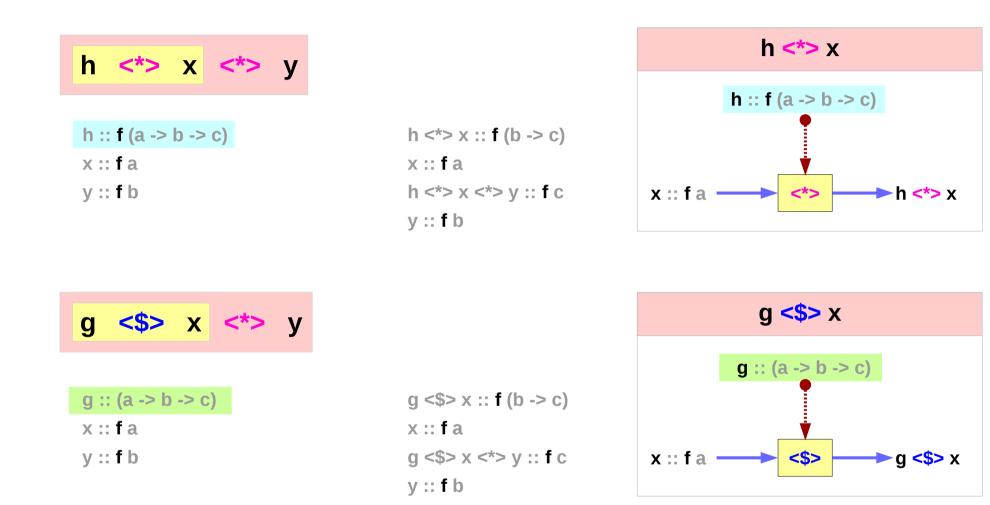
ghci> pure (+) <*> Just 3 <*> Nothing Nothing

```
ghci> pure (+) <*> Nothing <*> Just 5
Nothing
```

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Infix Operators <*> vs <\$>



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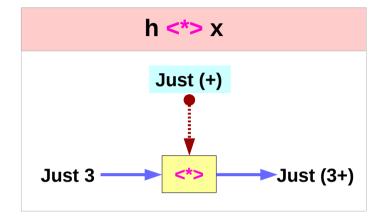
Applicatives Methods (3B)

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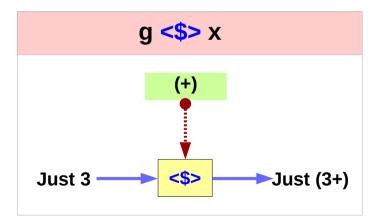
Infix Operators <*> vs <\$> examples



Just (+) <*> Just 3 <*> Just 2 Just (+3) <*> Just 2 Just 5



(+) <\$> Just 3 <*> Just 2 Just (+3) <*> Just 2 Just 5



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Applicatives Methods (3B)

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the minimal complete definition

class (Functor f) => Applicative f where
 pure :: a -> f a
 (<*>) :: f (a -> b) -> f a -> f b

the minimal complete definition

(<\$>) :: (Functor f) => (a -> b) -> f a -> f b

g **<\$>** x = fmap g x

Not in the minimal complete definition

g::a->b, x::fa

instance Applicative Maybe where
pure = Just
Nothing <*> _ = Nothing
(Just g) <*> something = fmap g something

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Applicatives Methods (3B)

The Applicative Typeclass

Applicative is a <u>superclass</u> of Monad. every Monad is also a Functor and an Applicative fmap, pure, (<*>) can all be used with monads.

a **Monad** instance

requires **Functor** and **Applicative** instances. defines the types and roles of **return** and (>>)

- fmap : defined in Functors
- pure, (<*>) : defined in Applicatives
- return, (>>) : defined in Monads

(<\$>) infix operator

(<\$>) :: (Functor f) => (a -> b) -> f a -> f b g <\$> x = fmap g x

The \$ operator is for avoiding parentheses

putStrLn (show (1 + 1))
putStrLn \$ show (1 + 1)
putStrLn \$ show \$ 1 + 1

(\$) calls the <u>function</u> which is its left-hand argument of \$ on the <u>value</u> which is its right-hand argument of \$

The Applicative Laws

The identity law:	pure id <*> v = v	id :: a -> a v :: f a
Homomorphism:	pure g <*> pure x = pure (g x)	g :: a -> b x :: a
Interchange:	u <*> pure y = pure (\$ y) <*> u	u :: f (a -> b) y :: a
Composition:	u <*> (v <*> w) = pure (.) <*> u <*> v <*> w	w :: f a v :: f (a -> b) u :: f (b -> c)
Left associative	u <*> v <*> w = (u <*> v) <*> w	u :: f (c -> b -> a) v :: f c
		u <*> v :: f (b -> a)
		w :: f b u <*> v <*> w = f c

Applicativ	/es
Methods ((3B)

The Identity Law

The identity law	pure id <*> v = v	id :: a -> a	v :: f a

pure to inject <u>values</u> into the <u>functor</u>in a default, featureless way,so that the result is as close as possibleto the <u>plain</u> value.

applying the **pure id** morphism does nothing, exactly like with the plain **id** function.

The Homomorphism Law

The homomorphism lawpure g <*> pure x = pure (g x)g :: a -> bx :: a	
---	--

applying a "**pure**" <u>function</u> to a "**pure**" <u>value</u> is the same as applying the <u>function</u> to the <u>value</u> in the normal way and then using **pure** on the result. means **pure** <u>preserves</u> function application.

applying a non-effectful function g
to a non-effectful argument x in an effectful context pure
is the same as just applying the function g to the argument x
and then injecting the result (f x)
into the effectual context with pure.

The Interchange Law

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u <*> pure y = pure (\$ y) <*> u

u :: f (a -> b) y :: a

(\$ y) is the *function* that <u>supplies</u> y

as argument to another function

– the higher order functions

Function \$ Argument \$ y

 (\mathbf{y}) as a single argument

applying a <u>morphism</u> u to a <u>"pure" value</u> pure y	Just (+3) <*> Just 2
is the same as applying pure (\$ y) to the <u>morphism</u> u	Just (\$ 2) <*> Just (+3)

when evaluating the application of	u :: f (a -> b)	y :: a
an <u>effectful function</u> (u) to a <u>pure argument</u> (pure y),	u <*> pure y :: f b	pure y :: f a
the <u>order</u> <u>doesn't</u> <u>matter</u> .	pure (\$ y) <*> u :: f b	pure (\$ y) :: f (a)

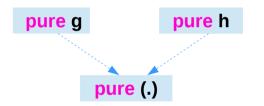
The Composition Law

The composition law pure (.) <*> u <*> v <*> w = u <*> (v <*> w) w :: f a v :: f (a -> b) u :: f (b -> c)

```
pure (.) composes morphisms similarly to how (.) composes functions:
```

```
pure (.) <*> pure f <*> pure g <*> pure x
= pure f <*> (pure g <*> pure x)
```

$$(g . h) x = g (h x)$$



applying the composed morphism pur	e (.) <*> u <*> v to w	u = <mark>pure</mark> g
gives the same result as applying u	u	v = <mark>pure</mark> h
to the result of applying ${f v}$ to ${f w}$	(v <*> w)	w = pure x

it is expressing a sort of <u>associativity</u> property of (<*>).

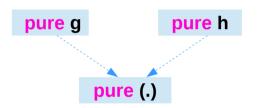
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The Composition Law – type view

pure (.) <*> u <*> v <*> w = u <*> (v <*> w)

w :: f a	value	
v :: f (a -> b)	func1	
u : f (b -> c)	func2	
v <*> w :: f b		application of func1
u <*> <mark>(</mark> v <*> w)::fc	application of func2

pure (.) <*> u <*> v :: f (a -> c) -- composite function pure (.) <*> u <*> v <*> w :: f c -- application (g.h) x = g(hx)



https://en.wikibooks.org/wiki/Haskell/Applicative_functors

Applicatives Methods (3B)

liftA2

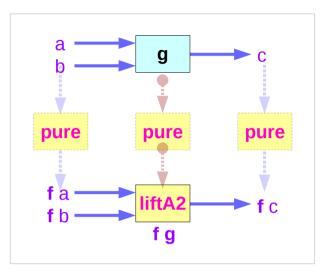
liftA2 :: (a -> b -> c) -> f a -> f b -> f c

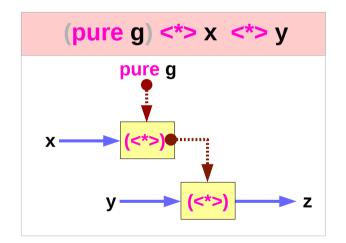
<u>lift</u> a <u>binary</u> <u>function</u> (**a->b->c**) to actions.

Some functors support an implementation of **liftA2** that is more efficient than the default one.

liftA2 may have an <u>efficient</u> implementation whereas **fmap** is an <u>expensive</u> operation,

sometimes better to use **liftA2** than to use **fmap** over the structure and then use **<*>**.

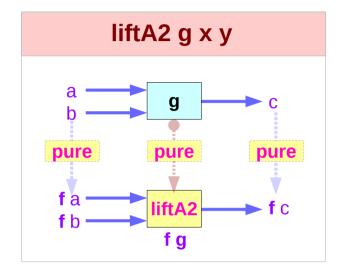




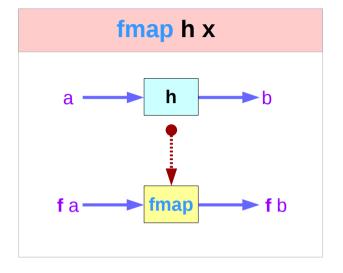
http://hackage.haskell.org/package/base-4.10.1.0/docs/Control-Applicative.html#v:liftA2

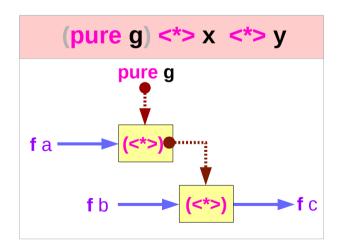
Applicatives Methods (3B)

liftA2, <*>, fmap, <\$>

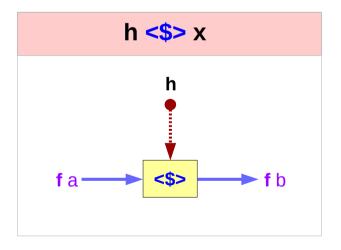


g :: a -> b -> c h :: a -> b



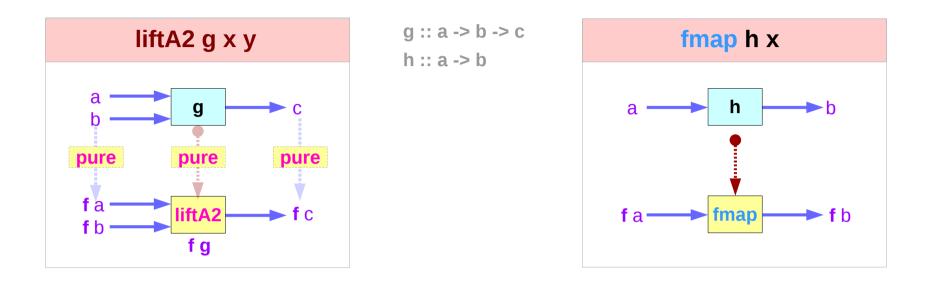


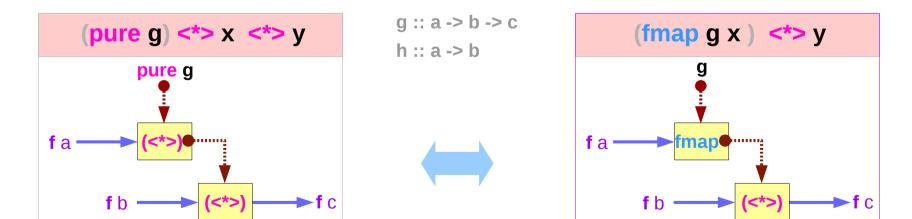
g :: a -> b -> c h :: a -> b



Applicatives Methods (3B)

Young Won Lim 4/10/18 pure g <*> x <*> y = (fmap g x) <*> y





Applicatives Methods (3B)

liftA2

variables

g :: a -> b -> c x :: f a y :: f b z :: f c pure g <*> x <*> y liftA2 g x y liftA2 :: (a -> b -> c) -> f a -> f b -> f c g :: a -> b -> c x :: f a

y :: f b liftA2 g x y :: f c

https://wiki.haskell.org/Applicative_functor

liftA2

Actually, using the **liftA** commands we can pull <u>results</u> of applicative functors into a <u>scope</u> where we can talk exclusively about <u>functor results</u> c and not about <u>effects</u>. f c Note that <u>functor results</u> can also be <u>functions</u>. c This scope is simply a <u>function</u>, which contains the code that we used in the non-functorial setting.

```
liftA3
(\x g h -> let y = g x in h y y)
fx fg fh
```

The order of effects is entirely determined by the order of arguments to liftA3

http://hackage.haskell.org/package/base-4.10.1.0/docs/Control-Applicative.html#v:liftA2

liftA2 :: (a -> b -> c) -> f a -> f b -> f c

A minimal complete definition must include implementations of **pure** and of either **<*>** or **liftA2**.

pure and <*>
pure and liftA2

If it defines both, then they must behave the same as their default definitions:

(<*>) = liftA2 id liftA2
f x y = f <\$> x <*> y

http://hackage.haskell.org/package/base-4.10.1.0/docs/Control-Applicative.html#v:liftA2

Applicatives Methods (3B)

liftA2

Consider the non-functorial expression:

x :: x g :: x -> y h :: y -> y -> z

let y = g x in h y y

Very simple. Now we like to generalize this to

fx :: f x fg :: f (x -> y) fh :: f (y -> y -> z)

https://wiki.haskell.org/Applicative_functor

liftA2

However, we note that let fy = fq <*> fx

in fh <*> fy <*> fy

runs the effect of fy

twice. E.g. if fy

writes something to the terminal then fh <*> fy <*> fy

writes twice. This could be intended, but how can we achieve, that the effect is run only once and the result is used twice? Actually, using the liftA

commands we can pull results of applicative functors into a scope where we can talk exclusively about functor results and not about effects. Note that functor results can also be functions. This scope is simply a function, which contains the code that we used in the non-functorial setting.

liftA3

```
(\x g h -> let y = g x in h y y)
fx fg fh
```

The order of effects is entirely determined by the order of arguments to liftA3

https://wiki.haskell.org/Applicative_functor

liftA2(<*>)

10 down vote accepted

The wiki article says that **liftA2 (<*>)** can be used to <u>compose applicative functors</u>. It's easy to see how to use it from its type:

o :: (Applicative f, Applicative f1) =>
 f (f1 (a -> b)) -> f (f1 a) -> f (f1 b)
o = liftA2 (<*>)

https://stackoverflow.com/questions/12587195/examples-of-haskell-applicative-transformers

liftA2(<*>) Examples

So to if f is Maybe and f1 is [] we get:	[(+1), (+6)]	[1, 6]
> Just [(+1),(+6)] `o` Just [1, 6] Just [2,7,7,12]	(+1) [1, 6] (+6) [1, 6]	

The other way around is:

> [Just (+1),Just (+6)] `o` [Just 1, Just 6]
[Just 2,Just 7,Just 7,Just 12]

[Just (+1),Just (+6)] [Ju

[Just 1, Just 6]

Just (+1) [Just 1, Just 6] Just (+6) [Just 1, Just 6]

https://stackoverflow.com/questions/12587195/examples-of-haskell-applicative-transformers

Applicatives Methods (3B)



LiftA2 (:)

your ex function is equivalent to liftA2 (:):

```
test1 = liftA2 (:) "abc" ["pqr", "xyz"]
```

To use (:) with deeper applicative stack you need multiple applications of liftA2:

```
*Main> (liftA2 . liftA2) (:) (Just "abc") (Just ["pqr", "xyz"])
Just ["apqr","axyz","bpqr","bxyz","cpqr","cxyz"]
```

However it only works when both operands are equally deep. So besides double liftA2 you should use pure to fix the level:

```
*Main> (liftA2 . liftA2) (:) (pure "abc") (Just ["pqr", "xyz"])
Just ["apqr","axyz","bpqr","bxyz","cpqr","cxyz"]
```

https://stackoverflow.com/questions/12587195/examples-of-haskell-applicative-transformers

<\$> related operators

Functor map <\$>

- (<\$>) :: Functor f => (a -> b) -> f a -> f b
- (<\$) :: Functor f => a -> f b -> f a
- (\$>) :: Functor f => f a -> b -> f b

The **<\$>** operator is just a synonym for the **fmap** function from the Functor typeclass.

This function generalizes the **map** function for lists to many other data types, such as **Maybe**, **IO**, and **Map**.

<\$> examples

#!/usr/bin/env stack

-- stack --resolver ghc-7.10.3 runghc

```
import Data.Monoid ((<>))
```

main :: IO ()

main = do

putStrLn "Enter your year of birth"

year <- read <\$> getLine

let age :: Int

age = 2020 - year

putStrLn \$ "Age in 2020: " <> show age

<\$, \$> operators

In addition, there are two additional operators provided which replace a value inside a Functor instead of applying a function. This can be both more convenient in some cases, as well as for some Functors be more efficient.

value <\$ functor = const value <\$> functor
functor \$> value = const value <\$> functor

x <\$ y = y \$> x x \$> y = y <\$ x

<*> related operators

Applicative function application <*>

(<*>) :: Applicative f => f (a -> b) -> f a -> f b
(*>) :: Applicative f => f a -> f b -> f b
(<*) :: Applicative f => f a -> f b -> f a

Commonly seen with **<\$>**, **<*>** is an operator that applies a wrapped function to a wrapped value. It is part of the Applicative typeclass, and is very often seen in code like the following:

foo <\$> bar <*> baz

<*> examples

For cases when you're dealing with a Monad, this is equivalent to:

do x <- bar

y <- baz

return (foo x y)

Other common examples including parsers and serialization libraries. Here's an example you might see using the aeson package:

data Person = Person { name :: Text, age :: Int } deriving Show

-- We expect a JSON object, so we fail at any non-Object value.

instance FromJSON Person where

parseJSON (Object v) = Person <\$> v .: "name" <*> v .: "age"
parseJSON _ = empty

***>** operator

To go along with this, we have two helper operators that are less frequently used:

***>** ignores the value from the first argument. It can be defined as:

a1 *> a2 = (id <\$ a1) <*> a2

Or in do-notation:

a1 *> a2 = do _ <- a1 a2

For Monads, this is completely equivalent to >>.

<* operator

<* is the same thing in reverse: perform the first action then the second, but only take the value from the first action. Again, definitions in terms of <*> and do-notation:

(<*) = liftA2 const

a1 <* a2 = do

res <- a1

_ <- a2

return res

(*> v.s. >>) and (pure v.s. return)

- (*>) :: Applicative f => fa -> fb -> fb
- (>>) :: Monad m => m a -> m b -> m b

pure :: Applicative f => a -> f a

return :: Monad m => a -> m a

the constraint changes from Applicative to Monad.

(*>) in Applicative(>>) in Monadpure in Applicativereturn in Monad

References

- [1] ftp://ftp.geoinfo.tuwien.ac.at/navratil/HaskellTutorial.pdf
- [2] https://www.umiacs.umd.edu/~hal/docs/daume02yaht.pdf