Variable Block Adder (1C)

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Please send corrections (or suggestions) to youngwlim@hotmail.com.

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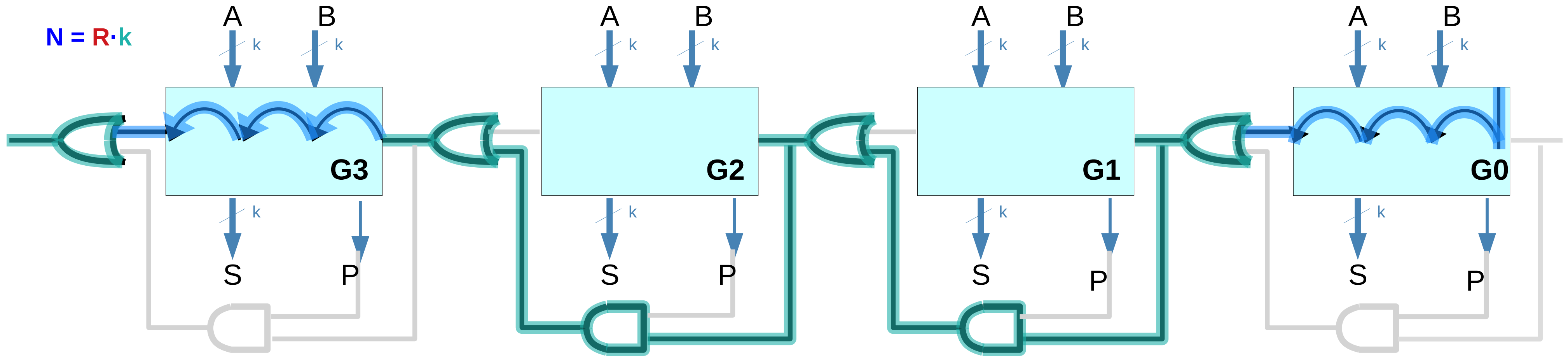
Carry Skip Adder



Fixed block size = k bits

$(k-1)\,\Delta_{rca}$	$(R-2)\,\Delta_{SKIP}$	$(k-1)\,\Delta_{rca}$
$(k-1)\,t$	$(m-2)\,T$	$(k-1)\,t$

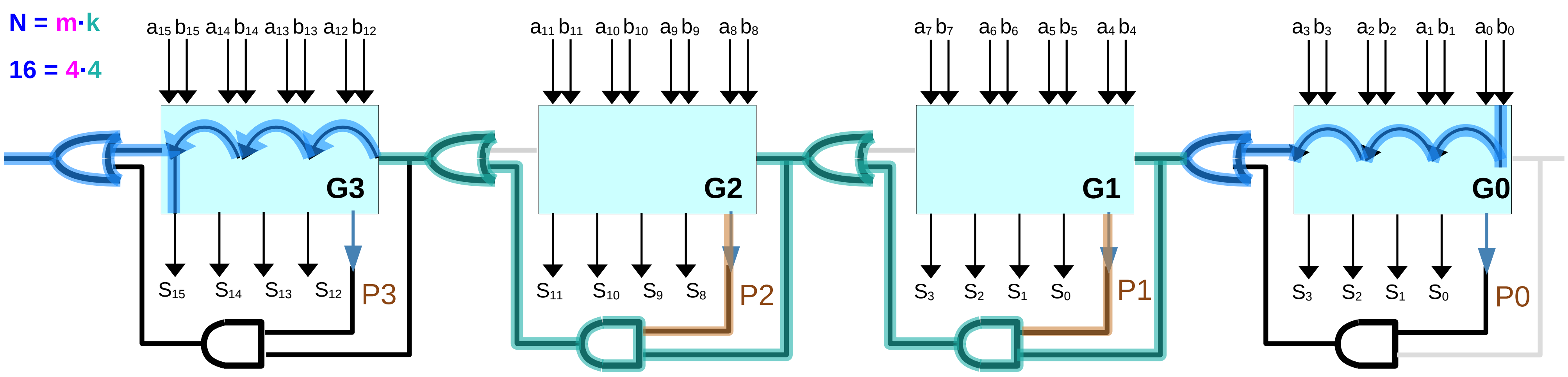
Variable block size = x_i bits for the i-th group

$(x_i-1)\,t$	$(m-2)\,T$	$(x_j-1)\,t$

t denote the time required for a carry signal to ripple across a bit
T denote the time required for the signal to skip over a group of bits
m denotes the optimal number of groups for an n-bit carry chain

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Carry Skip Adder – fixed block size



$P3 = p_{15} \cdot p_{14} \cdot p_{13} \cdot p_{12}$ $P2 = p_{11} \cdot p_{10} \cdot p_9 \cdot p_8$ $P1 = p_7 \cdot p_6 \cdot p_5 \cdot p_4$ $P0 = p_3 \cdot p_2 \cdot p_1 \cdot p_0$

$(k-1)\,\Delta_{rca}$ $(m-2)\,\Delta_{SKIP}$ $(k-1)\,\Delta_{rca}$

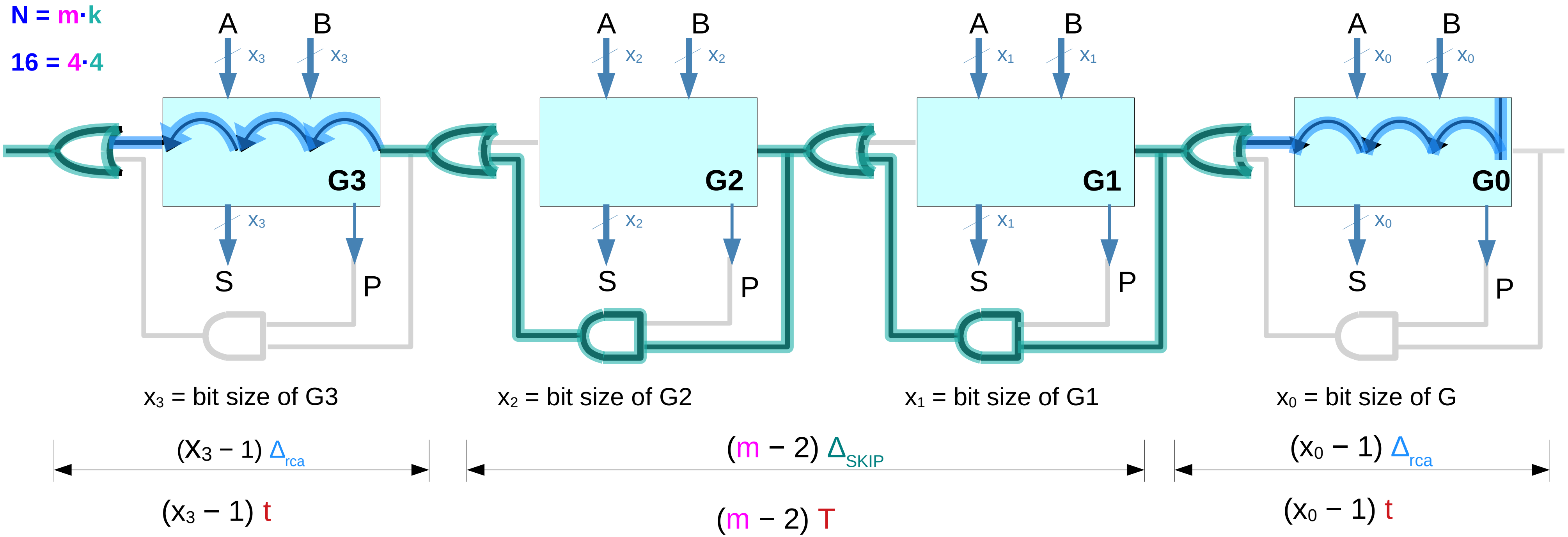
$(k-1)\,t$ $(m-2)\,T$ $(k-1)\,t$

t denote the time required for a carry signal to ripple across a bit
T denote the time required for the signal to skip over a group of bits
m denotes the optimal number of groups for an n-bit carry chain

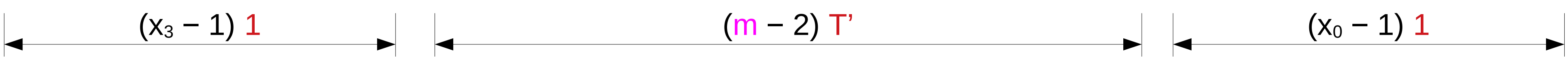
Fixed Block Size ➡ delay(P3) = delay(P2) = delay(P1) = delay(P0) = Fixed Delay

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Carry Skip Adder – maximum carry delay (3)



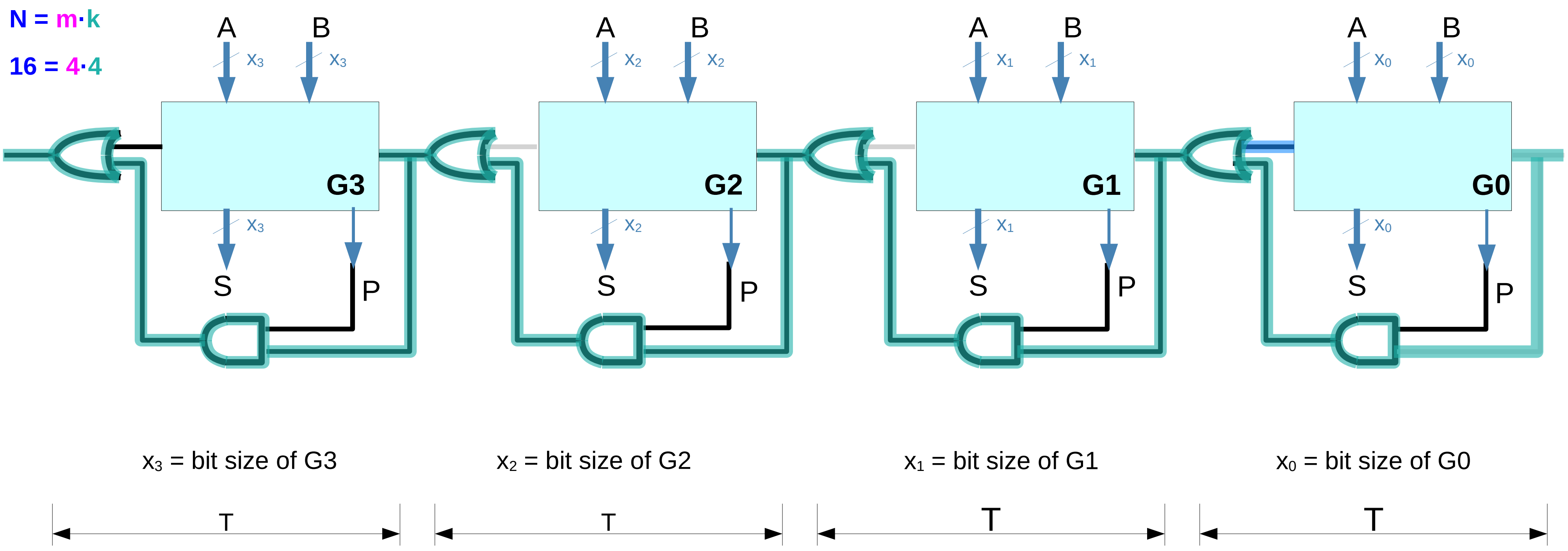
t denote the time required for a carry signal to ripple across a bit
T denote the time required for the signal to skip over a group of bits
m denotes the optimal number of groups for an n-bit carry chain



T' normalized delay of *T* over *t*

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Carry Skip Adder – maximum carry delay (3)



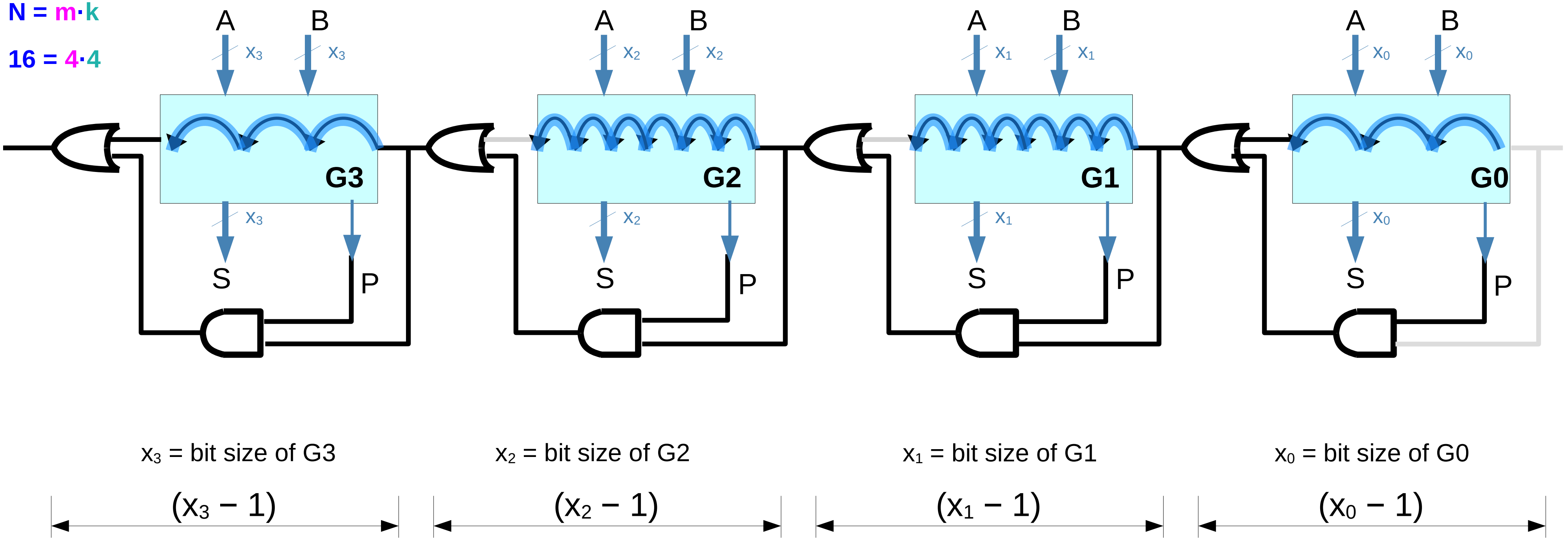
Carry Skip Delays

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Carry Skip Adder – maximum carry delay (3)

$N = m \cdot k$

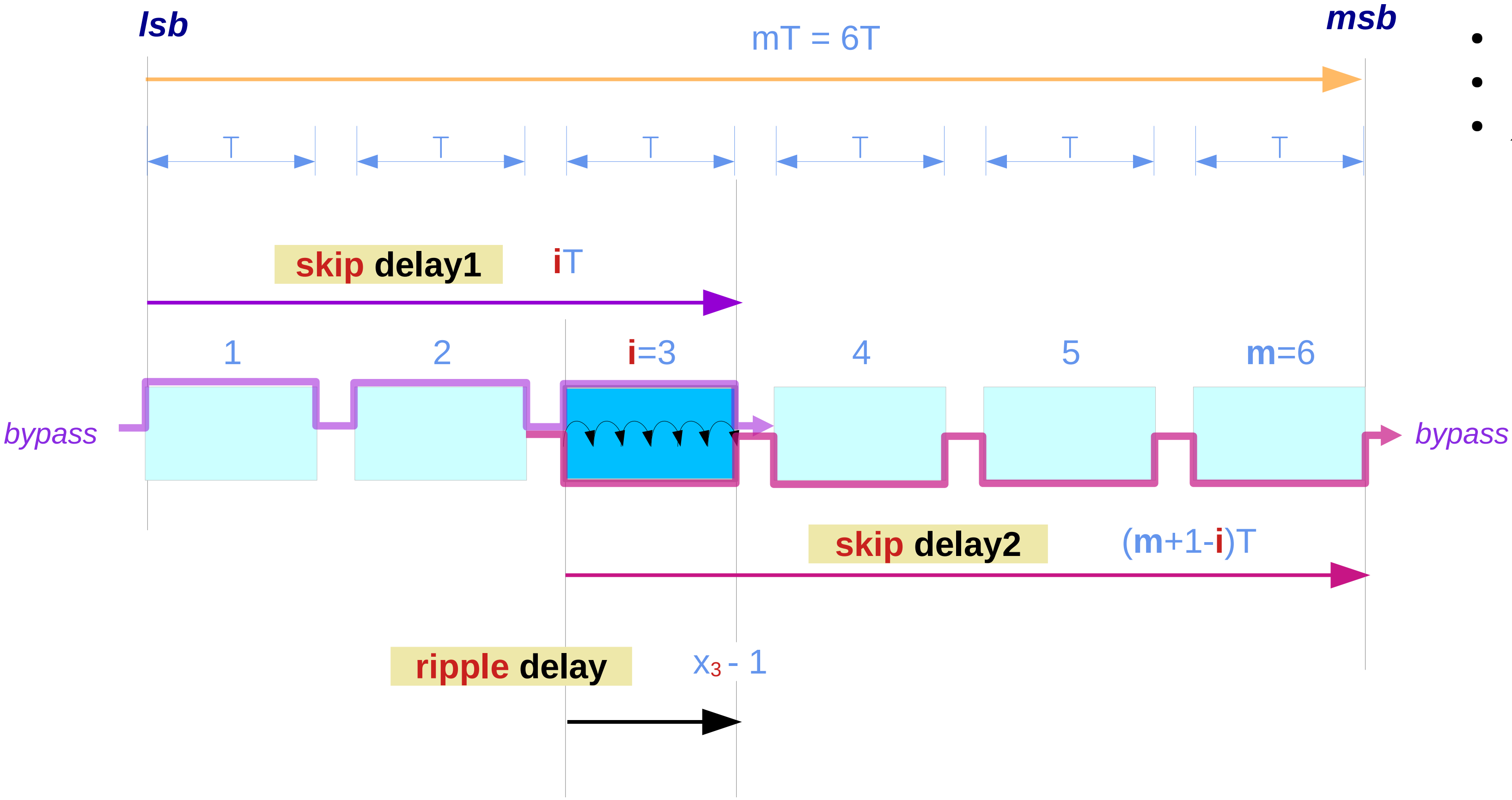
$16 = 4 \cdot 4$



Carry Ripple delays

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Skip path delays and ripple delays



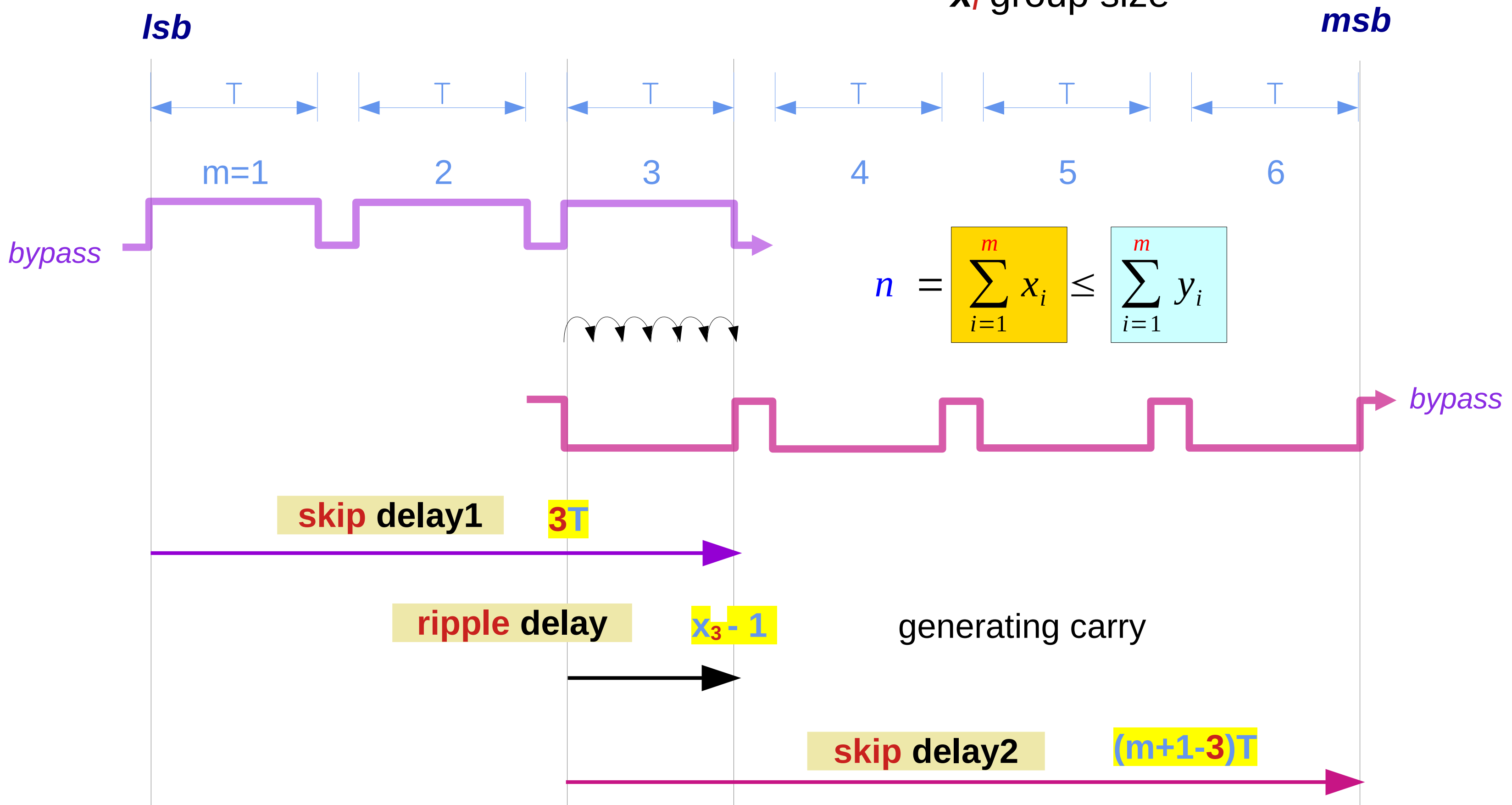
- n bits
- m groups
- x_i group size

ripple delay ≤ skip delay1
ripple delay ≤ skip delay2

ripple delay ≤ min (skip delay1, skip delay2)

Overlapping Delay Paths

- x_i group size



$$n = \sum_{i=1}^{m} x_i \le \sum_{i=1}^{m} y_i$$

ripple delay ≤ skip delay1
ripple delay ≤ skip delay2

m = 2;
y_1 = 1+T
y_2 = 1+T

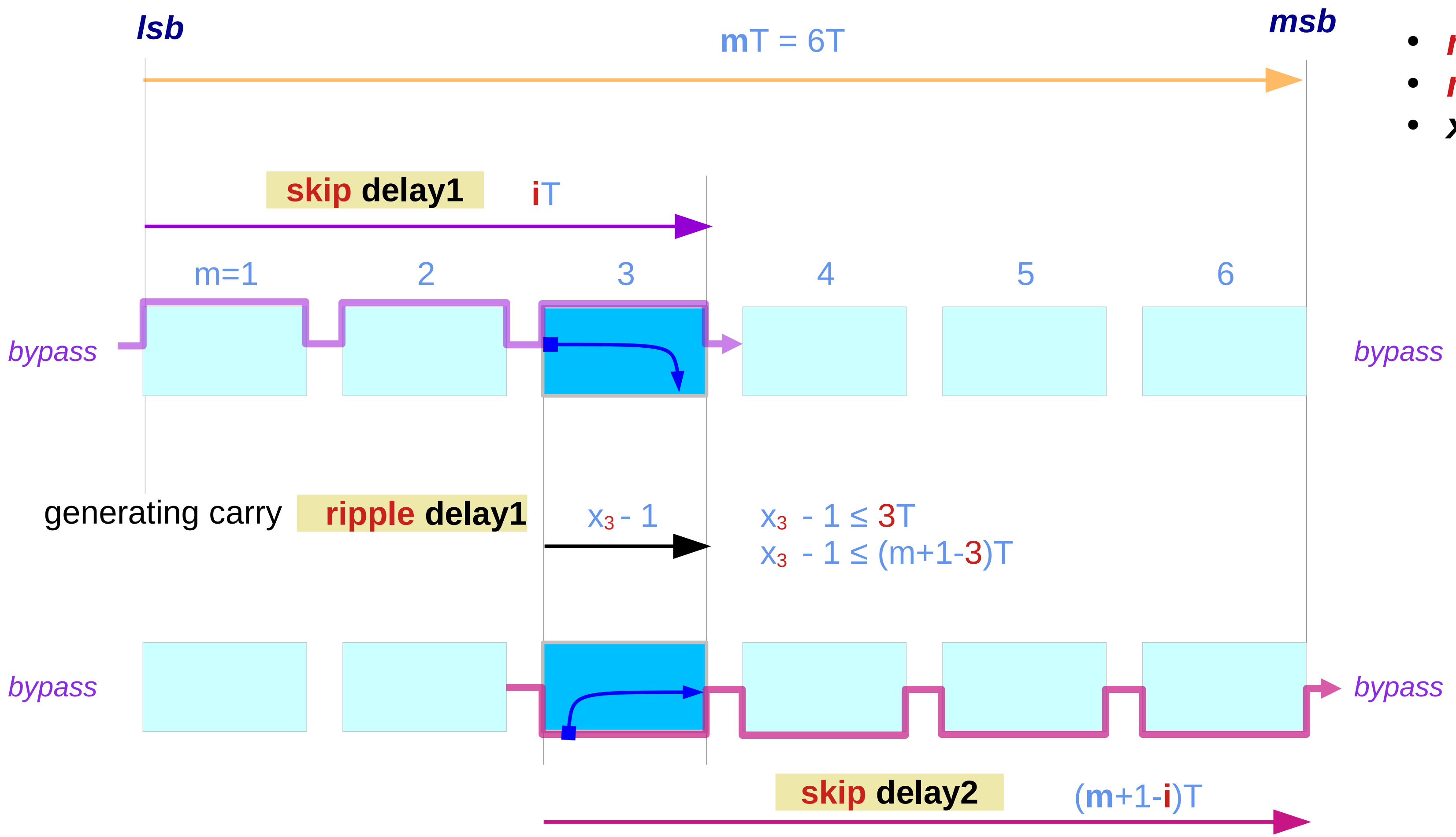
m = 3;
y_1 = 1+T
y_2 = 1+2T
y_3 = 1+T

m = 4;
y_1 = 1+T
y_2 = 1+2T
y_3 = 1+2T
y_4 = 1+T

m = 5;
y_1 = 1+T
y_2 = 1+2T
y_3 = 1+3T
y_4 = 1+2T
y_5 = 1+T

m = 6;
y_1 = 1+T
y_2 = 1+2T
y_3 = 1+3T
y_4 = 1+3T
y_6 = 1+T

Minimum skip path delay ($y_i - 1$) of the i^{th} group



- ***n*** bits
- ***m*** groups
- x_i group size

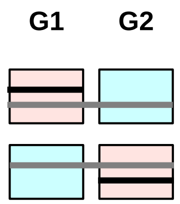
$$x_3 \le y_3 = \min(1 + iT, 1 + (m+1-i)T)$$

Minimum skip path delay of the i^{th} group (1)

m = 2;

min(T, 2T) = T

min(2T, T) = T

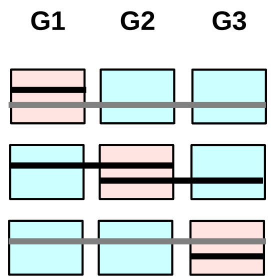


m = 3;

min(T, 3T) = T

min(2T, 2T) = 2T

min(3T, T) = T



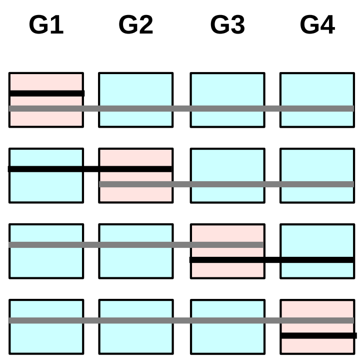
m = 4;

min(T, 4T) = T

min(2T,3T) = 2T

min(3T,2T) = 2T

min(4T,T) = T



m = 5;

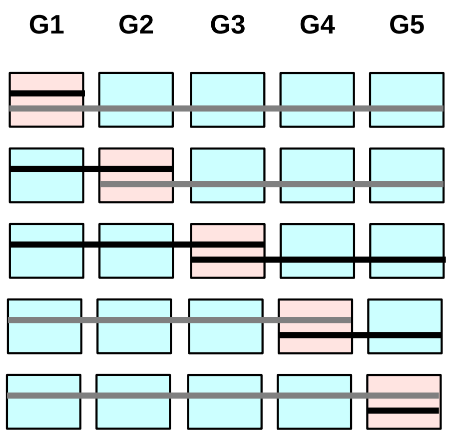
min(T, 5T) = T

min(2T, 4T) = 2T

min(3T, 3T) = 3T

min(4T, 2T) = 2T

min(5T, T) = T



$$y_i = 1 + min\{iT,\ (m+1-i)T\}$$

$$y_i = min\{1+iT, 1+(m+1-i)T\}, \quad i = 1, \ldots, m$$

Minimum skip path delay of the i^{th} group (2)

m = 6;

min(T, 6T) = T

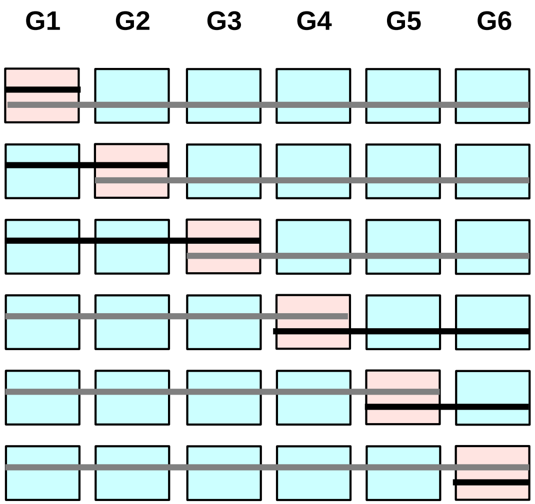
min(2T, 5T) = 2T

min(3T, 4T) = 3T

min(4T, 3T) = 3T

min(5T, 2T) = 2T

min(6T, T) = T



m = 7;

min(T, 7T) = T

min(2T, 6T) = 2T

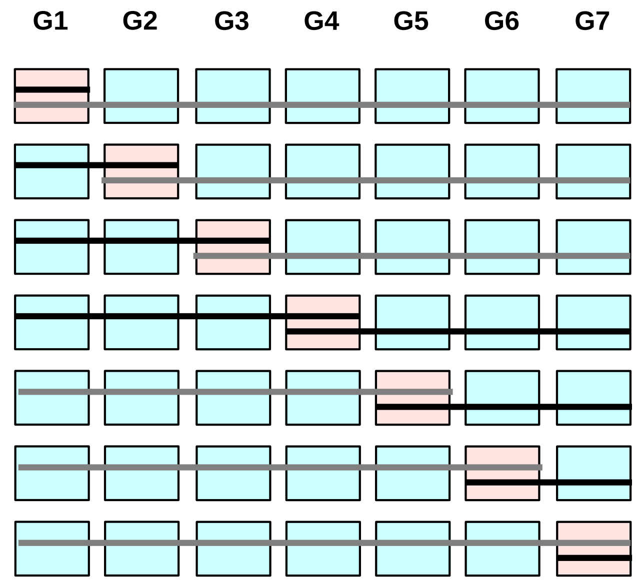
min(3T, 5T) = 3T

min(4T, 4T) = 4T

min(5T, 3T) = 3T

min(6T, 2T) = 2T

min(7T, T) = T



$$y_i = 1 + min\{iT,\ (m+1-i)T\}$$

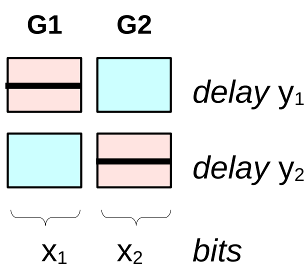
$$\Rightarrow \quad y_i = min\{1+iT,\ 1+(m+1-i)T\}, \quad i = 1, \dots, m$$

The i^{th} group has x_i bits for a given m (1)

m = 2;

$x_1 \le y_1 = 1+T$

$x_2 \le y_2 = 1+T$

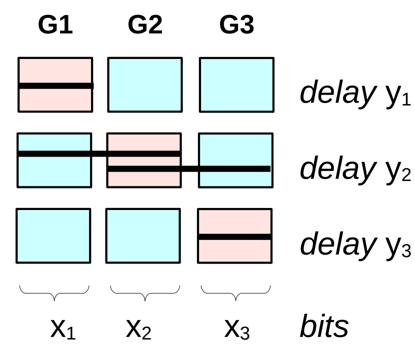


m = 3;

$x_1 \le y_1 = 1+T$

$x_2 \le y_2 = 1+2T$

$x_3 \le y_3 = 1+T$



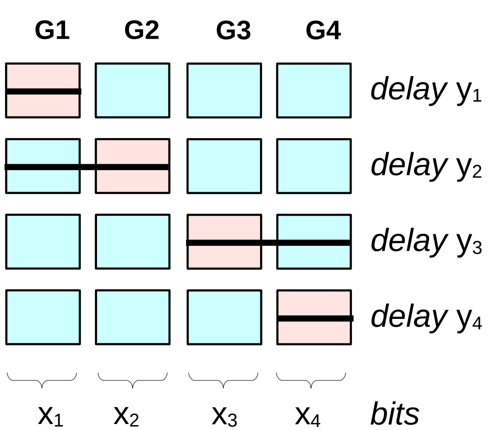
m = 4;

$x_1 \le y_1 = 1+T$

$x_2 \le y_2 = 1+2T$

$x_3 \le y_3 = 1+2T$

$x_4 \le y_4 = 1+T$



m = 5;

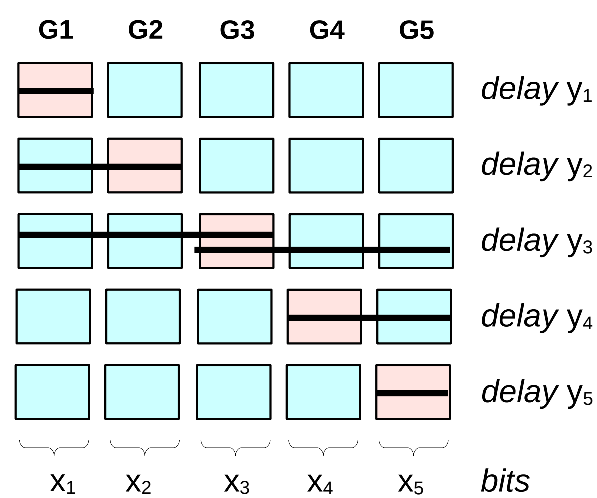
$x_1 \le y_1 = 1+T$

$x_2 \le y_2 = 1+2T$

$x_3 \le y_3 = 1+3T$

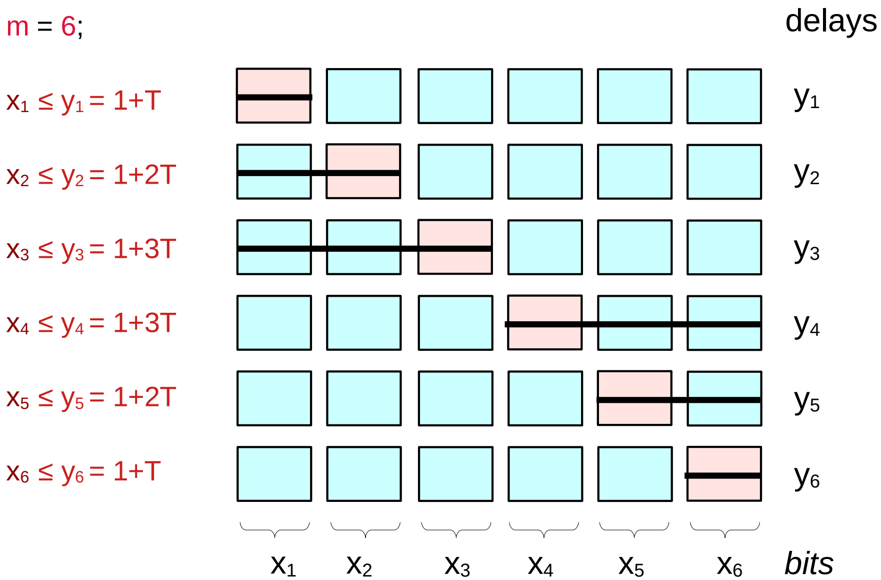
$x_4 \le y_4 = 1+2T$

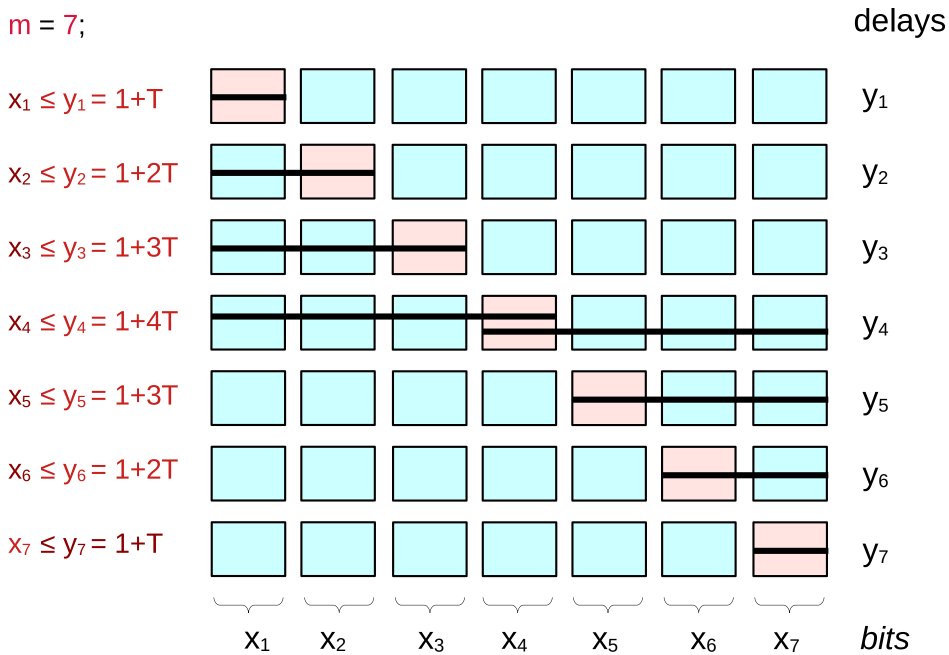
$x_5 \le y_5 = 1+T$



$x_i - 1 \le \min(iT, (m+1-i)T)$ → $x_i \le y_i = \min(1 + iT, 1 + (m+1-i)T)$

The i^{th} group has x_i bits for a given m (2)





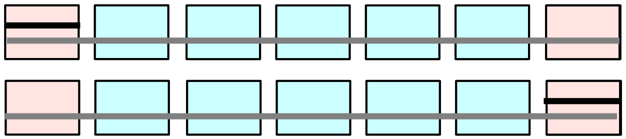
$x_i - 1 \le \min(\ iT, (m+1-i)T\ )$ → $x_i \le y_i = \min(\ 1 + iT, 1 + (m+1-i)T\ )$

Symmetric histograms of y_i's

$$y_1 = min\{1+1\cdot T, 1+(m+1-1)T\} = 1 + T$$

$$y_m = min\{1+m\cdot T, 1+(m+1-m)T\} = 1 + T$$

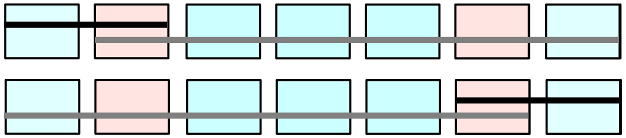
$$x_1 \le y_1 = 1+T \quad (bits)$$
$$x_m \le y_m = 1+T \quad (bits)$$



$$y_2 = min\{1+2\cdot T, 1+(m+1-2)T\} = 1 + 2T$$

$$y_{m-1} = min\{1+(m-1)\cdot T, 1+(m+1-(m-1))T\} = 1 + 2T$$

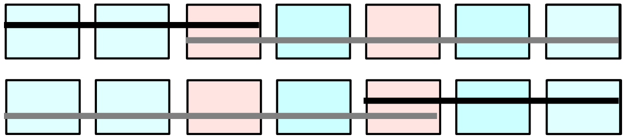
$$x_2 \le y_2 = 1+2T \quad (bits)$$
$$x_{m-1} \le y_{m-1} = 1+2T \quad (bits)$$



$$y_3 = min\{1+3\cdot T, 1+(m+1-3)T\} = 1 + 3T$$

$$y_{m-2} = min\{1+(m-2)\cdot T, 1+(m+1-(m-2))T\} = 1 + 3T$$

$$x_3 \le y_3 = 1+3T \quad (bits)$$
$$x_{m-2} \le y_{m-2} = 1+3T \quad (bits)$$



Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Procedure

(I) Let m be the smallest positive integer

$$n \le \sum_{i=1}^{m} y_i \qquad i = 1, \dots, m$$

(II) Let

$$y_i = min\{1+iT, 1+(m+1-i)T\}$$

(III) Let x_i , i = 1, …, m

starting with the first row, row by row

$$n = \sum_{i=1}^{m} x_i \le \sum_{i=1}^{m} y_i$$

Variable block size = x_i bits for the i-th group

find the smallest m

$$n \le \sum_{i=1}^{m} y_i = \sum_{i=1}^{m} min\{1+iT, 1+(m+1-i)T\}$$

m = 2;
while ($y_1+\cdots+y_m$ < n) m = m+1 ;

m = 2;	m = 4;	m = 5;	m = 6;	m = 7;
y_1 = 1+T	y_1 = 1+T	y_1 = 1+T	y_1 = 1+T	y_1 = 1+T
y_2 = 1+T	y_2 = 1+2T	y_2 = 1+2T	y_2 = 1+2T	y_2 = 1+2T
	y_3 = 1+2T	y_3 = 1+3T	y_3 = 1+3T	y_3 = 1+3T
m = 3;	y_4 = 1+T	y_4 = 1+2T	y_4 = 1+3T	y_4 = 1+4T
y_1 = 1+T		y_5 = 1+T	y_5 = 1+2T	y_5 = 1+3T
y_2 = 1+2T			y_6 = 1+T	y_6 = 1+2T
y_3 = 1+T				y_7 = 1+T

the scheme (i), (ii), (iii)
gives the **max propagation time** mT

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Determining *m* the number of groups (1)

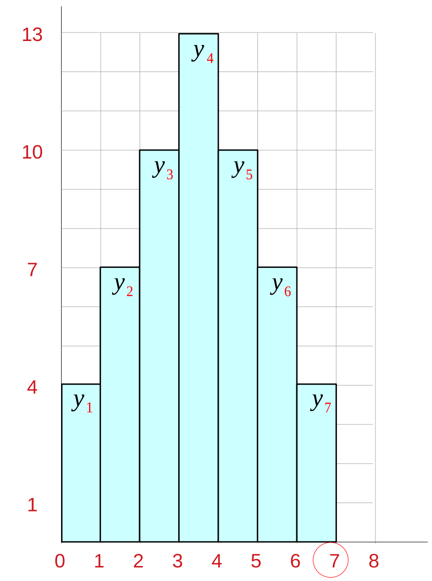
Method 1 – using a histogram

Let m be the <u>smallest</u> positive integer such that

$$n \leq \sum_{i=1}^{m} y_i$$

```
m = 2;
while (y1+···+ym < n)  m = m+1 ;
```

$$y_i = min\{1+iT, 1+(m+1-i)T\}, \quad i = 1, \ldots, m$$



m = 2; T = 3
y_1 = min {1+T, 1+2T} = 1+T = 4
y_2 = min {1+2T, 1+T} = 1+T = 4

m = 3; T = 3
y_1 = min {1+T, 1+3T} = 1+T = 4
y_2 = min {1+2T, 1+2T} = 1+2T = 7
y_3 = min {1+3T, 1+T} = 1+T = 4

m = 4; T = 3
y_1 = min {1+T, 1+4T} = 1+T = 4
y_2 = min {1+2T, 1+3T} = 1+2T = 7
y_3 = min {1+3T, 1+2T} = 1+2T = 7
y_4 = min {1+4T, 1+T} = 1+T = 4

m = 5; T = 3
y_1 = min {1+T, 1+5T} = 1+T = 4
y_2 = min {1+2T, 1+4T} = 1+2T = 7
y_3 = min {1+3T, 1+3T} = 1+3T = 10
y_4 = min {1+4T, 1+2T} = 1+2T = 7
y_5 = min {1+5T, 1+T} = 1+T = 4

m = 6; T = 3
y_1 = min {1+T, 1+6T} = 1+T = 4
y_2 = min {1+2T, 1+5T} = 1+2T = 7
y_3 = min {1+3T, 1+4T} = 1+3T = 10
y_4 = min {1+4T, 1+3T} = 1+3T = 10
y_5 = min {1+5T, 1+2T} = 1+2T = 7
y_6 = min {1+6T, 1+T} = 1+T = 4

m = 7; T = 3
y_1 = min {1+T, 1+7T} = 1+T = 4
y_2 = min {1+2T, 1+6T} = 1+2T= 7
y_3 = min {1+3T, 1+5T} = 1+3T= 10
y_4 = min {1+4T, 1+4T} = 1+4T= 13
y_5 = min {1+5T, 1+3T} = 1+3T= 10
y_6 = min {1+6T, 1+2T} = 1+2T= 7
y_7 = min {1+7T, 1+1T} = 1+T = 4

4+4 = 8
4+7+4 = 15
4+7+7+4 = 22
4+7+10+7+4 = 32
4+7+10+10+7+4 = 42
4+7+10+13+10+7+4 = 55

Determining *m* the number of groups (2)

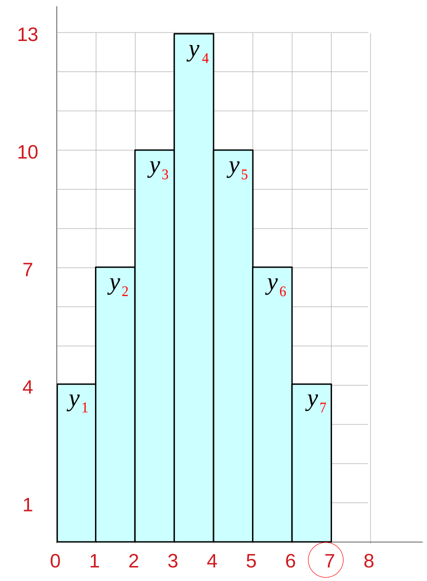
Method 1 – using a histogram

Let *m* be the smallest positive integer such that

$$n \le \sum_{i=1}^{m} y_i$$

m = 2;
while ($y_1+\cdots+y_m$ < n) m = m+1 ;

$$y_i = min\{1+iT, 1+(m+1-i)T\}, \quad i = 1, \ldots, m$$

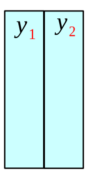


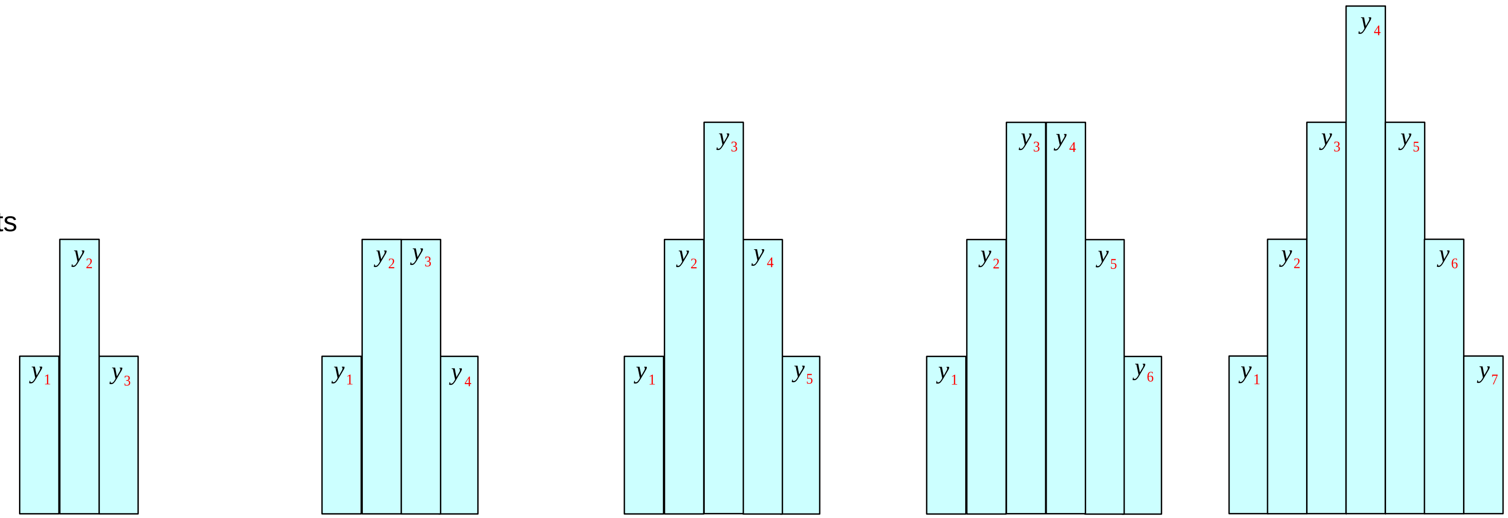
$m = 2$	$m = 3$	$m = 4$	$m = 5$	$m = 6$	$m = 7$
$\sum_{i=1}^{2} y_i$	$\sum_{i=1}^{3} y_i$	$\sum_{i=1}^{4} y_i$	$\sum_{i=1}^{5} y_i$	$\sum_{i=1}^{6} y_i$	$\sum_{i=1}^{7} y_i$
$2 + 2\cdot T$	$3 + 4\cdot T$	$4 + 6\cdot T$	$5 + 9\cdot T$	$6 + 12\cdot T$	$7 + 16\cdot T$
$2 + 2\cdot T < n$	$3 + 4\cdot T < n$	$4 + 6\cdot T < n$	$5 + 9\cdot T < n$	$6 + 12\cdot T < n$	$7 + 16\cdot T > n$
$2 + 2\cdot 3 < 48$	$3 + 4\cdot 3 < 48$	$4 + 6\cdot 3 < 48$	$5 + 9\cdot 3 < 48$	$6 + 12\cdot 3 < 48$	$7 + 16\cdot 3 > 48$
$8 < 48$	$15 < 48$	$22 < 48$	$32 < 48$	$42 < 48$	$55 > 48$
					$m = 7$

$n = 48$
$T = 3$
→ $m = 7$

n the number of bits, ***m*** the number of groups

4+4 = 8 bits
4+7+4 = 15 bits
4+7+7+4 = 22 bits
4+7+10+7+4 = 32 bits
4+7+10+10+7+4 = 42 bits
4+7+10+13+10+7+4 = 55 bits





no of groups
m = 2; T = 3
y_1 = 1+T = 4
y_2 = 1+T = 4
--
8

no of groups
m = 3; T = 3
y_1 = 1+T = 4
y_2 = 1+2T = 7
y_3 = 1+T = 4
--
15

no of groups
m = 4; T = 3
y_1 = 1+T = 4
y_2 = 1+2T= 7
y_3 = 1+2T= 7
y_4 = 1+T = 4
--
22

no of groups
m = 5; T = 3
y_1 = 1+T = 4
y_2 = 1+2T = 7
y_3 = 1+3T = 10
y_4 = 1+2T = 7
y_5 = 1+T = 4
--
32

no of groups
m = 6; T = 3
y_1 = 1+T = 4
y_2 = 1+2T = 7
y_3 = 1+3T = 10
y_4 = 1+3T = 10
y_5 = 1+2T = 7
y_6 = 1+T = 4
--
42

no of groups
m = 7; T = 3
y_1 = 1+T = 4
y_2 = 1+2T = 7
y_3 = 1+3T = 10
y_4 = 1+4T = 13
y_5 = 1+3T = 10
y_6 = 1+2T = 7
y_7 = 1+T = 4
--
55

no of bits
$n \le 8$

no of bits
$8 < n \le 15$

no of bits
$15 < n \le 22$

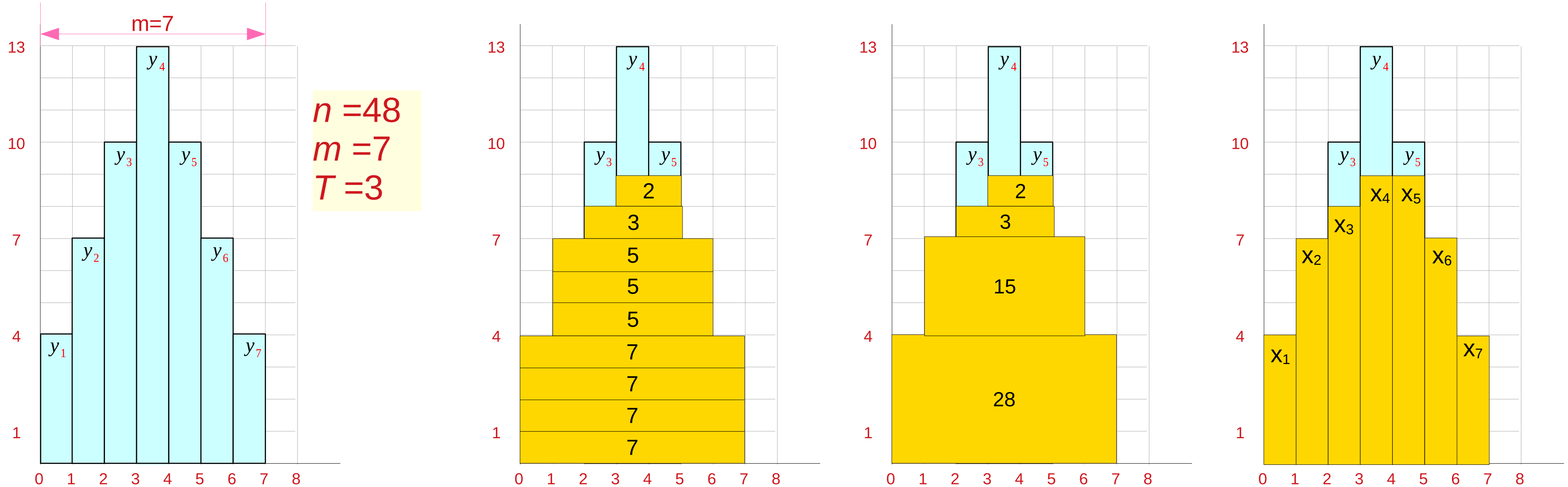
no of bits
$22 < n \le 32$

no of bits
$32 < n \le 42$

no of bits
$42 < n \le 55$

n =48
T =3
→ *m* =7

Determining x_i the group size of the i^{th} group



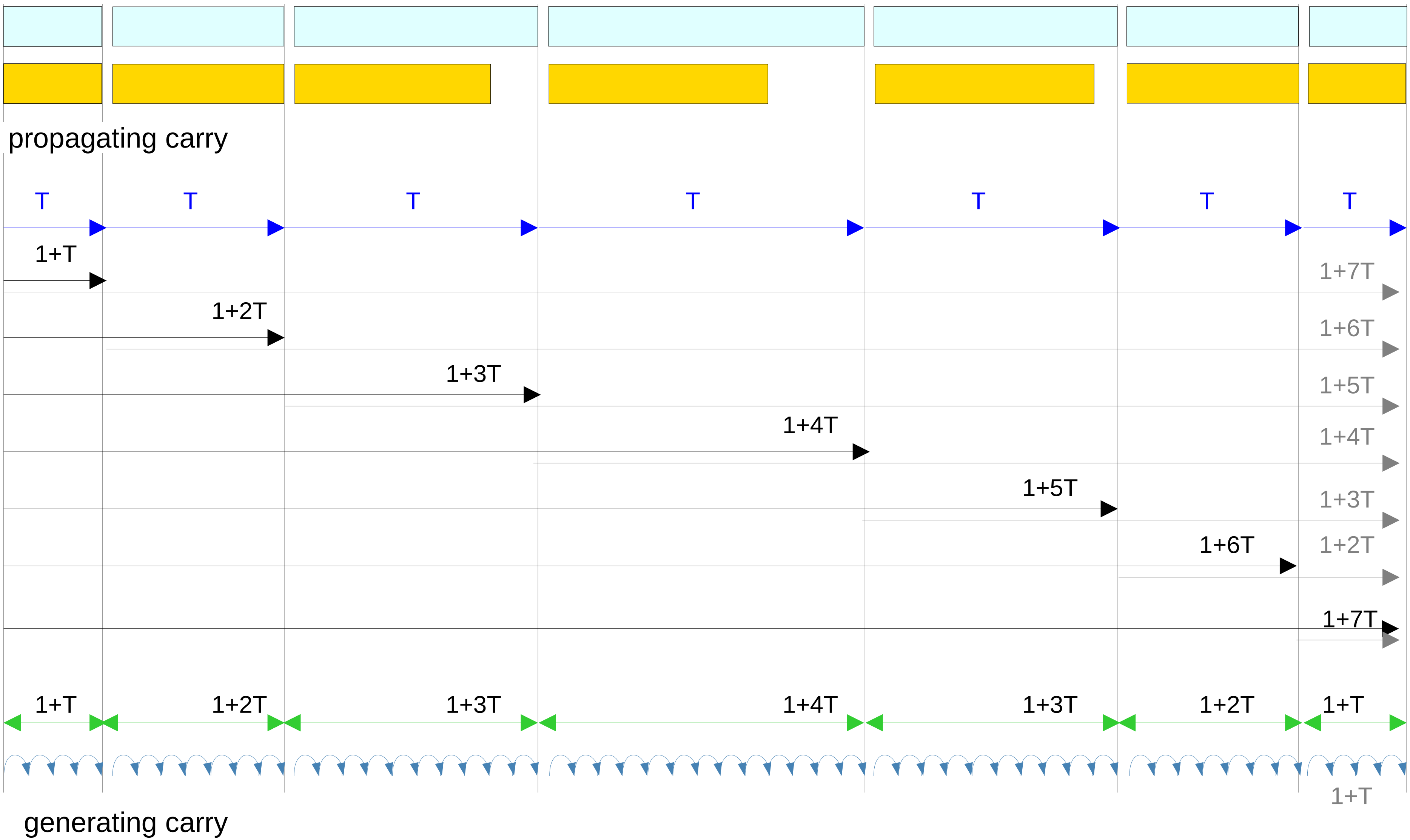
7	7
7	14
7	21
7	28
5	33
5	38
5	43
3	46
2	48

4 * 7 = 28	28
3 * 5 = 15	43 = 28 + 15
1 * 3 = 3	46 = 28 + 15 + 3
1 * 2 = 2	48 = 28 + 15 + 3 + 2

$x_1 = 4 \leq y_1 = 4$
$x_2 = 7 \leq y_2 = 7$
$x_3 = 8 < y_3 = 10$
$x_4 = 9 < y_4 = 13$
$x_5 = 9 < y_5 = 10$
$x_6 = 7 \leq y_6 = 7$
$x_7 = 4 \leq y_7 = 4$

m =7 groups

Determining x_i the group size of the i^{th} group



Procedure

(I) Let m be the smallest positive integer such that

$$n \le m + \frac{1}{2} m T + \frac{1}{4} m^2 T + \left(1 - (-1)^m\right) \frac{1}{8} T = \sum_{i=1}^{m} y_i$$

- total n =48 bits
- m =7 groups
- i-th group has x_i bits (size)
- constant skip delay $T = T(x_i) = 3$

(II) Let

$$y_i = \min\{1 + iT, 1 + (m + 1 - i)T\}, \quad i = 1, \ldots, m$$

and construct a histogram whose i-th column has height y_i

for example, for T=3, and n=48, we have m=7

(III) It is easily verified that the area of the histogram in (II) is

$$\sum_{i=1}^{m} y_i = m + \frac{1}{2} m T + \frac{1}{4} m^2 T + \left(1 - (-1)^m\right) \frac{1}{8} T \ge n$$

so these are at least n unit squares in the histogram
starting with the first row, shade in n of the squares, row by row
Let x_i denote the number of shaded squares in column i of the histogram,
$i = 1, \ldots, m$

$$n = \sum_{i=1}^{m} x_i \le \sum_{i=1}^{m} y_i$$

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Procedure - determining *m* the number of groups

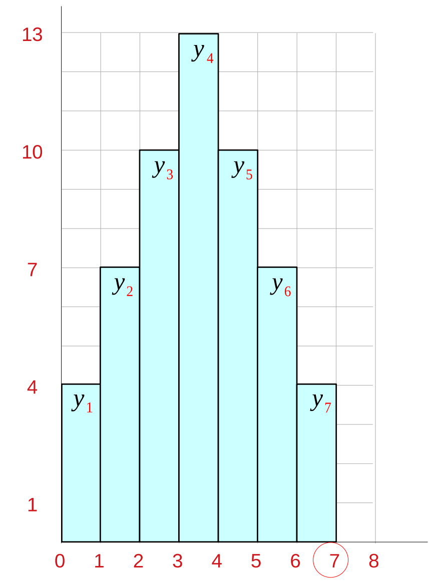
Method 1 – using a histogram

Let *m* be the smallest positive integer such that

$$n \le \sum_{i=1}^{m} y_i$$

```
m = 2;
while (y1+···+ym < n)  m = m+1 ;
```

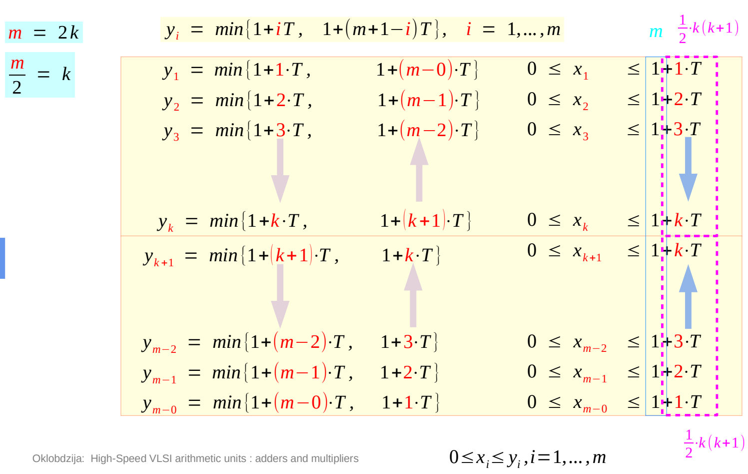
$$y_i = \min\{1+iT,\ 1+(m+1-i)T\}, \quad i = 1,\dots,m$$



Method 2 – using a *closed formula*

Let *m* be the smallest positive integer such that

$$n \le m + \frac{1}{2}mT + \frac{1}{4}m^2T + \left(1-(-1)^m\right)\frac{1}{8}T$$



Processing - determining x_i the group size of the i^{th} group

construct a histogram
whose *i-th* column has height y_i

so these y_i's are at least n unit squares
in the histogram, starting with the first row,
shade in n of the squares, row by row

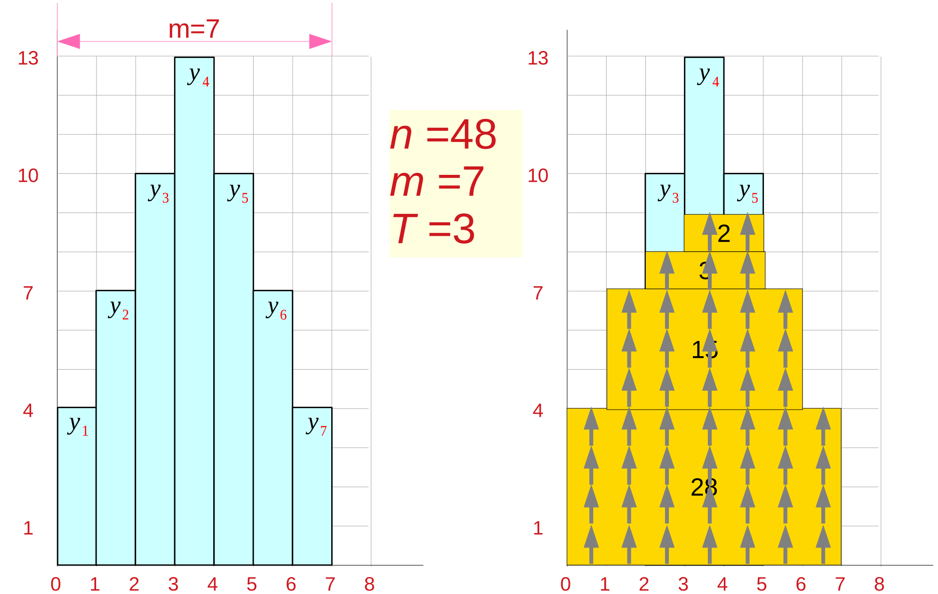
let x_i denote the number of shaded squares
in column i of the histogram,
$i = 1, \ldots, m$

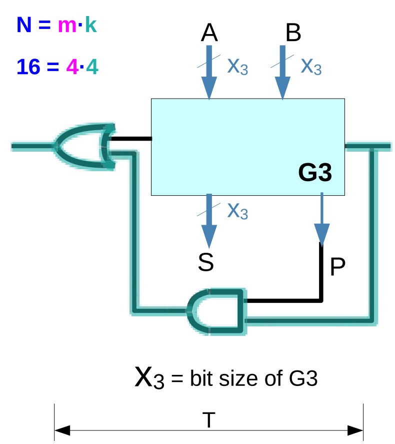
$$0 \le x_i \le y_i, \quad i=1,\ldots,m$$

$$n = \sum_{i=1}^{m} x_i \le \sum_{i=1}^{m} y_i$$

$$n = \sum_{i=1}^{7} x_i$$

$$n = 4+7+8+9+9+7+4=48$$
$$\le 4+7+10+13+10+7+4=55$$





m = 7; T = 3

$x_1 = 4 \le y_1 = 4$
$x_2 = 7 \le y_2 = 7$
$x_3 = 8 < y_3 = 10$
$x_4 = 9 < y_4 = 13$
$x_5 = 9 < y_5 = 10$
$x_6 = 7 \le y_6 = 7$
$x_7 = 4 \le y_7 = 4$

Maximum Propagation Time P_{max}

find the <u>smallest</u> *m*

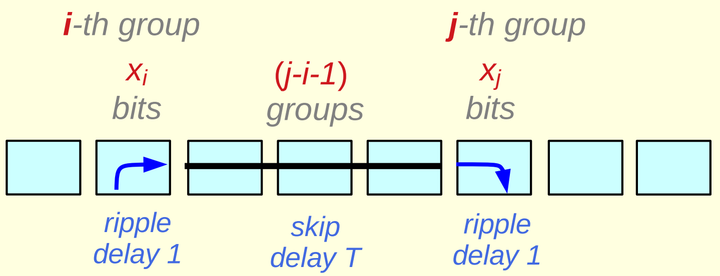
$$n \le \sum_{i=1}^{m} y_i = \sum_{i=1}^{m} min\{1+iT,\ 1+(m+1-i)T\}$$

m = 2;
while ($y_1+\cdots+y_m$ < n) m = m+1 ;

the scheme (i), (ii), (iii)
gives the max propagation time *mT*

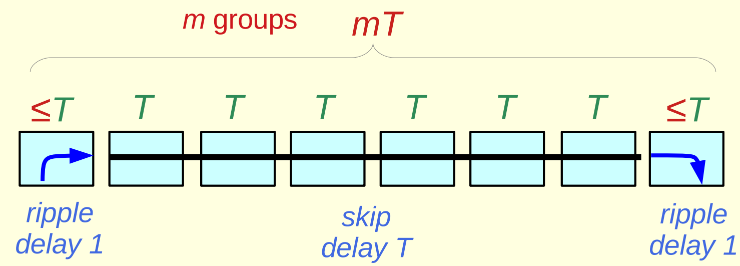
propagation time of a carry signal $\le mT$
the maximum propagation time = mT

propagation time $P = P_{i,j}$



$$P = P_{i,j} \qquad \forall i,\ \forall j \quad 1 \le i,\ j \le m$$

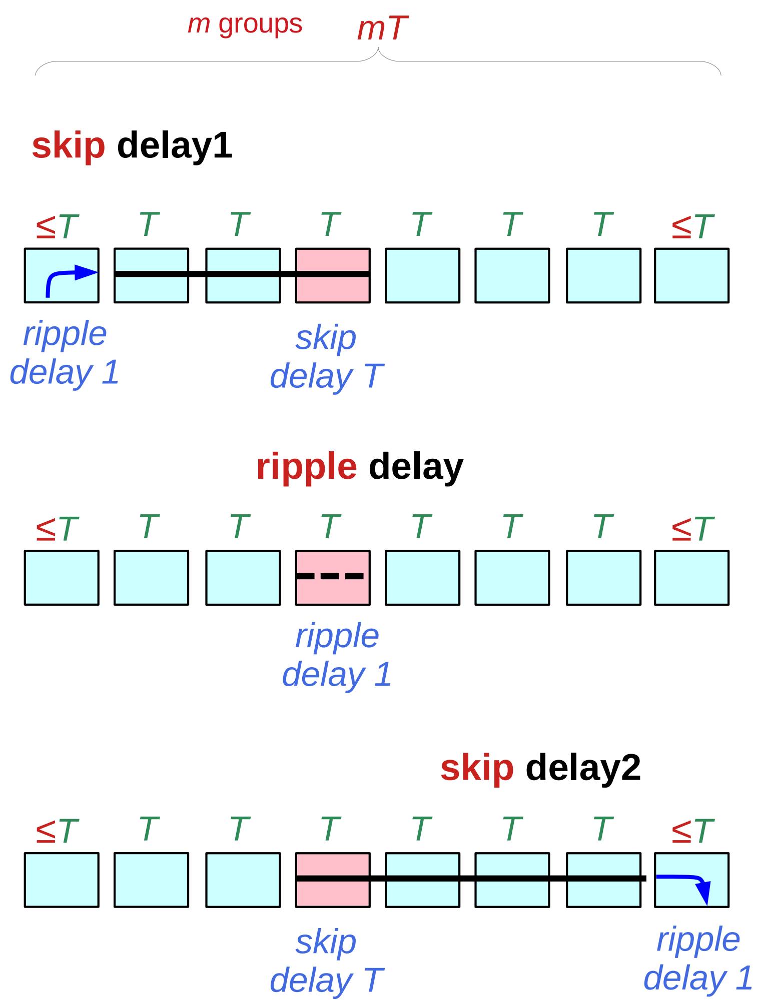
maximum propagation time $P_{max} = P_{1,m}$



$$P_{max} = P_{1,m} \le mT$$

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Maximum propagation time P



the scheme (i), (ii), (iii)
gives the max prop time *mT*

skip delay1	iT	generating carry
ripple delay	$x_i - 1$	
skip delay2	$(m+1-i)T$	terminating carry

$$\begin{cases} x_i - 1 \le iT \\ x_i - 1 \le (m+1-i)T \end{cases}$$

$$\Rightarrow \begin{cases} x_i \le 1 + iT \\ x_i \le 1 + (m+1-i)T \end{cases}$$

$$\Rightarrow x_i \le \min\{1 + iT,\ 1 + (m+1-i)T\} = y_i$$

$$\Rightarrow x_i \le y_i$$

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Maximum propagation time P

$$y_i = min\{1+iT, 1+(m+1-i)T\}, \quad i = 1, \ldots, m$$

the scheme (i), (ii), (iii)
gives the max prop time mT

$$y_1 = min\{1+1 \cdot T, 1+(m+1-1)T\} = 1 + T$$

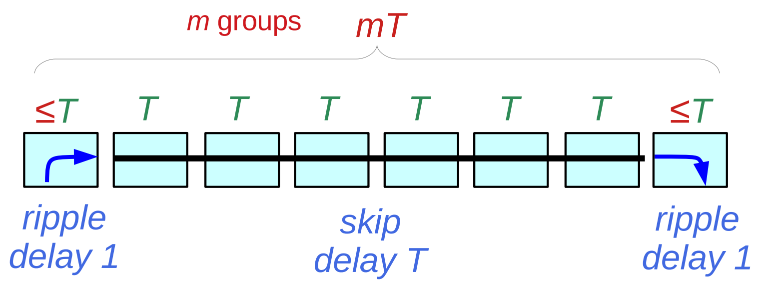
$$y_m = min\{1+m \cdot T, 1+(m+1-m)T\} = 1 + T$$

$$x_1 \le y_1 = 1+T$$

$$x_m \le y_m = 1+T$$

$x_1 - 1 \le 1T$ $x_1 - 1 \le (m+1-1)T$	$x_1 \le T+1$ $x_1 \le mT + 1$
$x_2 - 1 \le 2T$ $x_2 - 1 \le (m+1-2)T$	$x_2 \le 2T + 1$ $x_2 \le (m-1)T + 1$
• • •	• • •
$x_{m-1} - 1 \le (m-1)T$ $x_{m-1} - 1 \le (m+1-(m-1))T$	$x_{m-1} \le (m-1)T + 1$ $x_{m-1} \le 2T + 1$
$x_m - 1 \le mT$ $x_m - 1 \le (m+1-m)T$	$x_m - 1 \le mT$ $x_m - 1 \le T$

maximum propagation time



$$P_{max} = P_{1,m} \le mT$$

$$P = P_{i,j} \le mT$$

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Maximum propagation time P

Lemma 1 When the bits of a carry skip adder are grouped according to the scheme (i)-(iii), the maximum propagation time of a carry signal is mT

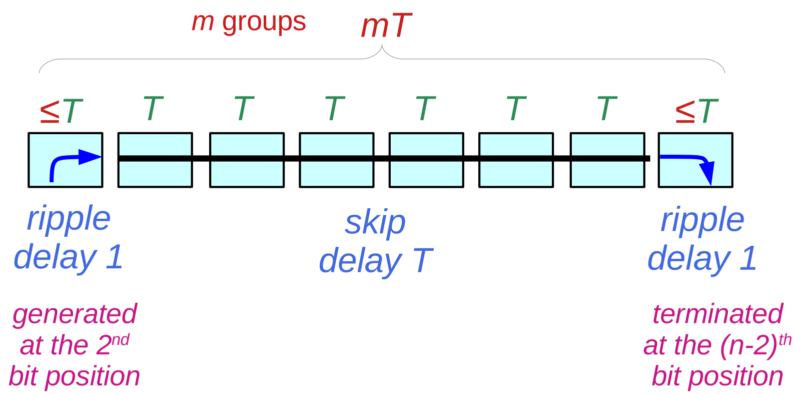
the scheme (i), (ii), (iii) gives the max prop time mT

The carry generated at the 2nd bit position and terminating at the $(n\text{-}1)^{th}$ bit position clearly has propagation time mT.

We must show that *any other* carry signal has propagation time smaller than or equal to mT

propagation time of a carry signal $\leq mT$
the maximum propagation time = mT

maximum propagation time



$$P_{max} = P_{1,m} \leq mT$$

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Delay range of a carry signal

Lemma 2 Let *D* denote the maximum delay of a carry signal in a *n* bit carry skip adder with group sizes chosen optimally. Then

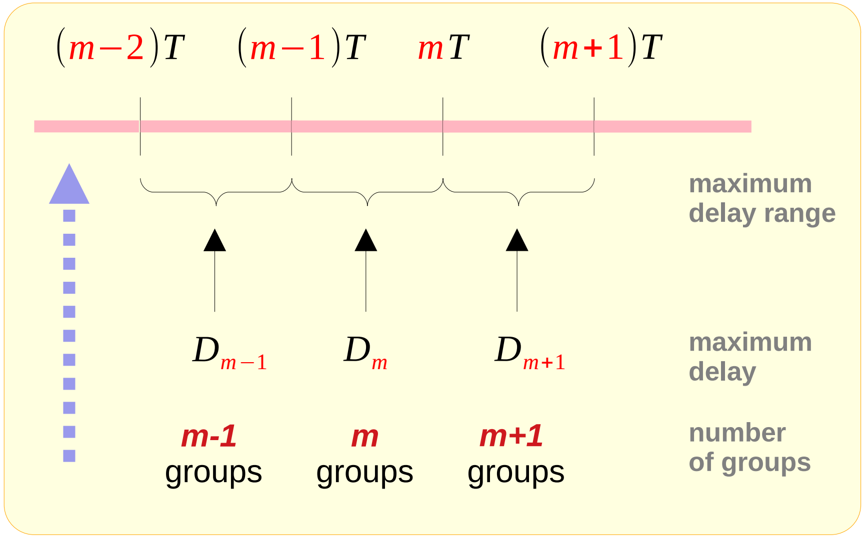
$$(m-1)T \;\le\; D \;\le\; mT$$

- ***n*** bits
- ***r*** groups

Since we have exhibited a division of the carry chain into groups
In such a way that the maximum delay of a carry signal is mT
We clearly have $D \le mT$

the **maximum delay** = D
the **optimal group size** = m

$$(m-1)T \;\le\; D \;\le\; mT$$



Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Maximum delay and optimal group size

the maximum propagation time
$\propto$ the number of groups

$$D \propto m$$

- *not an optimal optimal division*
  - *larger number of groups →*
  - *larger delays →*

- when group size m is not *optimal*

  then there is an optimal group size = r

  - the maximum delay with the group size m $D_m = mT$
  - the maximum delay with the group size r $D_r = rT$
  - *r must be smaller than m* $r \le m$

$$D_r < D_m$$

$$\Rightarrow rT < mT$$

$$\Rightarrow r < m$$

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Maximum delay of a carry signal

$$(m-1)T \le D \le mT$$

*Assume there are **r** groups*

*then 2 cases : **even r**, **odd r***

for each of these 2 cases

prove $mT - D < T + 1$

➡ $mT - D \le T$

➡ $(m-1)T \le D$

P: ***the propagation delay of any carry signal path ≤ mT***

D: ***the max of P***

the upper bound of D

diff (** mT**, D) ≤ T

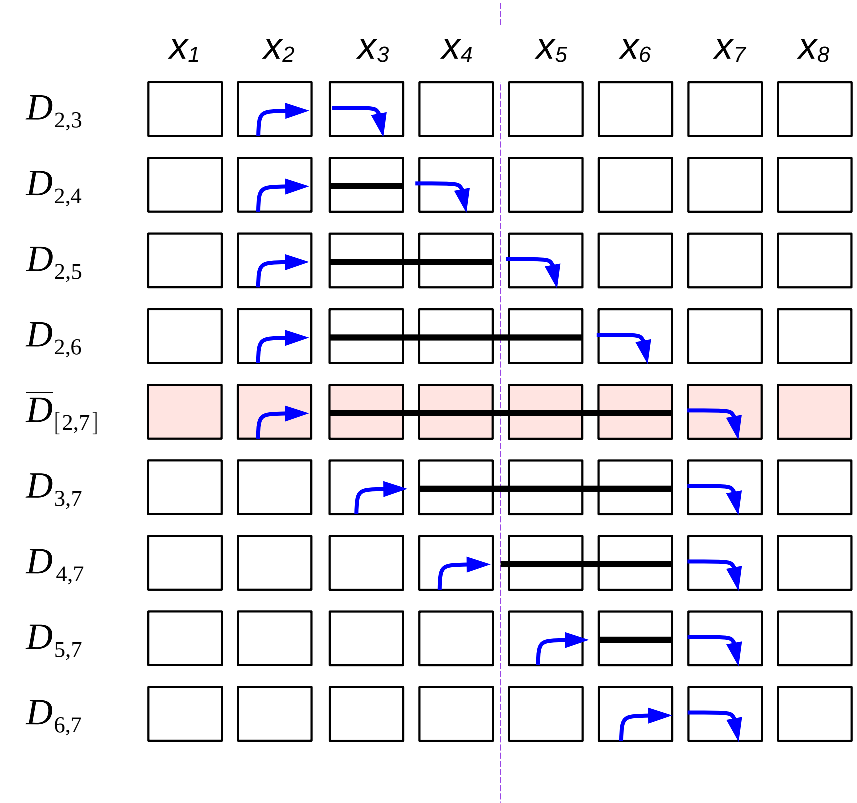
➡ ***diff (** mT**,** max **P) ≤ T***

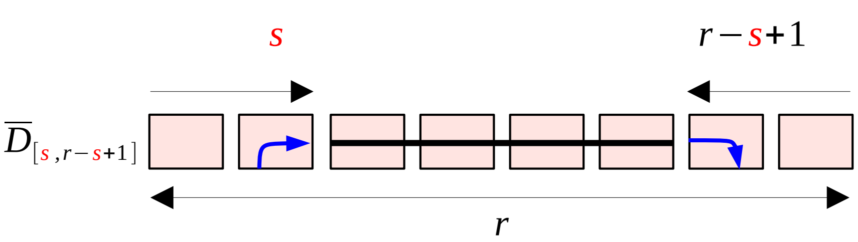
the lower bound of D

$(m-1)T \le D$

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Maximum delays of carry signals (**r = 2k**)





$\overline{D}_{[2,7]}$ = *the maximum delay of carry signals generated in the i-th group and terminated in the j-th group such that* $2 \le i, j \le 7$ $\;\le\; D$

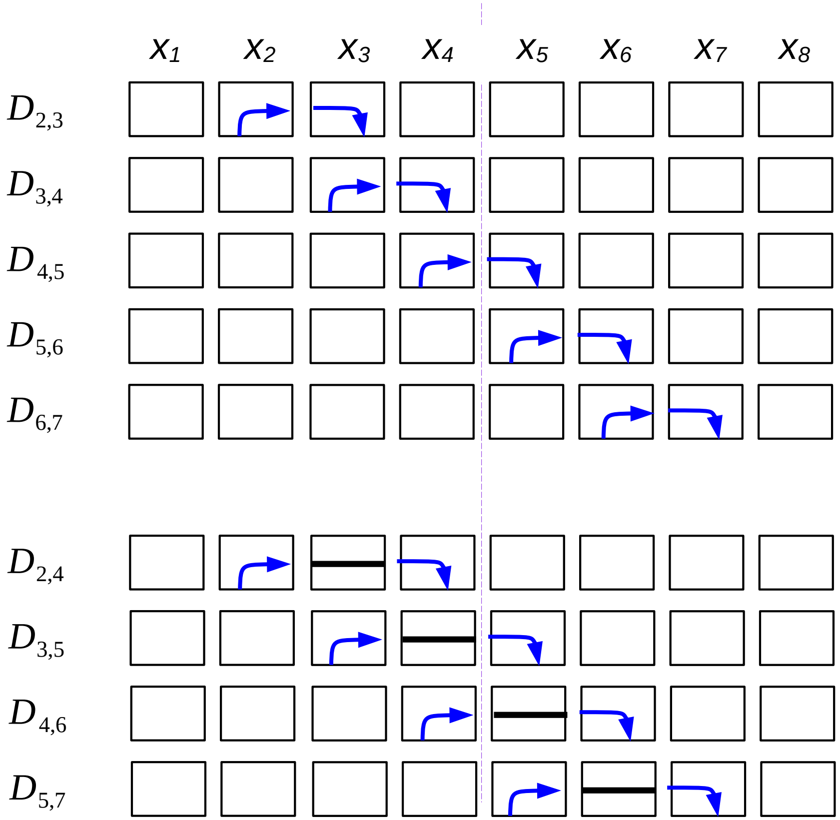
$$\overline{D}_{[2,7]} = \max \left\{ \begin{matrix} D_{2,3}, & D_{2,4}, & D_{2,5}, & D_{2,6}, \\ & D_{2,7}, & & \\ D_{3,7}, & D_{4,7}, & D_{5,7}, & D_{6,7} \end{matrix} \right\}$$

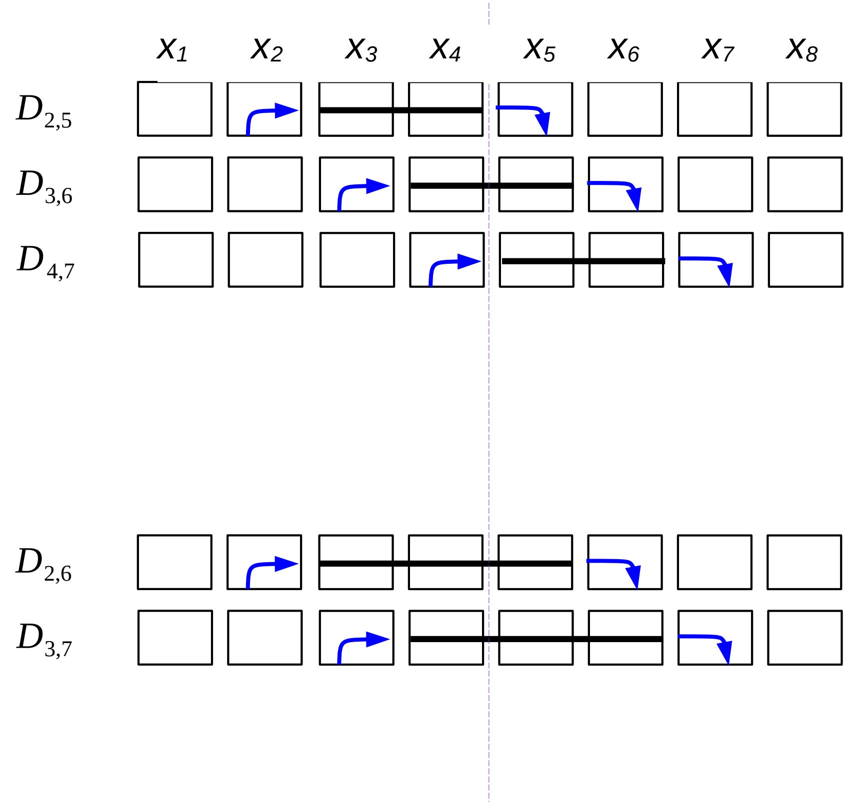
$\overline{D}_{[s,r-s+1]}$ = *the maximum delay of carry signals generated in the i-th group and terminated in the j-th group such that* $s \le i, j \le r-s+1$

$$\overline{D}_{[2,7]} = \overline{D}_{[2,8-2+1]} = \overline{D}_{[s,8-s+1]} \quad \leftarrow s = 2$$

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Maximum delays of carry signals (**r = 2k**)

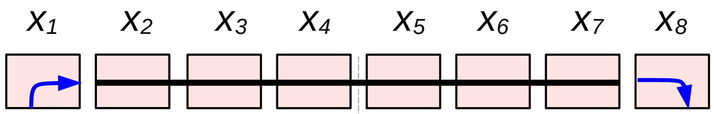




Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

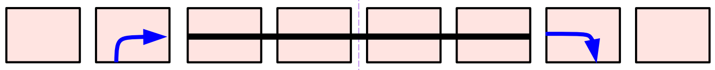
Maximum delays of carry signals (r = 2k)

$\overline{D}_{[1,8]} =$



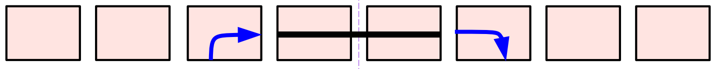
$\overline{D}_{[s,r-s+1]}\Big|_{s=1}$

$\overline{D}_{[2,7]} =$



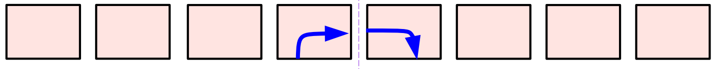
$\overline{D}_{[s,r-s+1]}\Big|_{s=2}$

$\overline{D}_{[3,6]} =$



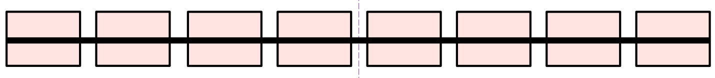
$\overline{D}_{[s,r-s+1]}\Big|_{s=3}$

$\overline{D}_{[4,5]} =$



$\overline{D}_{[s,r-s+1]}\Big|_{s=4}$

$\overline{D}_{all} =$



$\overline{D}_{[1,8]} =$ *The maximum delay of carry signals generated in the i-th group or terminated in the j-th group such that 1 ≤ i, j ≤ 8* $\leq D$

$\overline{D}_{[2,7]} =$ *The maximum delay of carry signals generated in the i-th group or terminated in the j-th group such that 2 ≤ i, j ≤ 7* $\leq D$

$\overline{D}_{[3,6]} =$ *The maximum delay of carry signals generated in the i-th group or terminated in the j-th group such that 3 ≤ i, j ≤ 6* $\leq D$

$\overline{D}_{[4,5]} =$ *The maximum delay of carry signals generated in the i-th group or terminated in the j-th group such that 4 ≤ i, j ≤ 5* $\leq D$

$\overline{D}_{all} =$ *All skip delay* $\leq D$

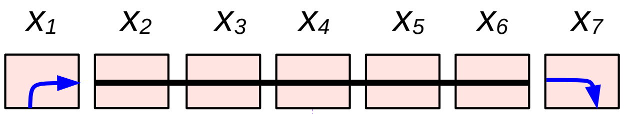
$$D = max\{\overline{D}_{[1,8]},\ \overline{D}_{[2,7]},\ \overline{D}_{[3,6]},\ \overline{D}_{[4,5]}\}$$

Max delay of all carry signals

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

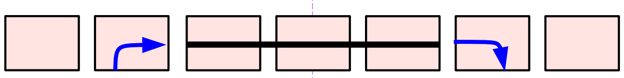
Maximum delays of carry signals (**r = 2k+1**)

$\overline{D}_{[1,7]} =$



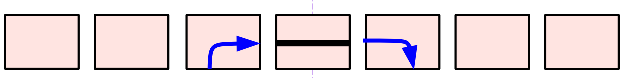
$\overline{D}_{[s,r-s+1]} \Big|_{s=1}$

$\overline{D}_{[2,6]} =$



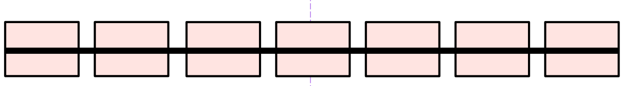
$\overline{D}_{[s,r-s+1]} \Big|_{s=2}$

$\overline{D}_{[3,5]} =$

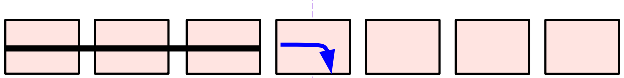


$\overline{D}_{[s,r-s+1]} \Big|_{s=3}$

$\overline{D}_{all} =$



$\overline{D}_{all} =$



$\overline{D}_{[1,7]} =$ *The maximum delay of carry signals generated in the i-th group or terminated in the j-th group such that* 1 ≤ *i, j* ≤ 8 $\leq$ D

$\overline{D}_{[2,6]} =$ *The maximum delay of carry signals generated in the i-th group or terminated in the j-th group such that* 2 ≤ *i, j* ≤ 7 $\leq$ D

$\overline{D}_{[3,65]} =$ *The maximum delay of carry signals generated in the i-th group or terminated in the j-th group such that* 3 ≤ *i, j* ≤ 6 $\leq$ D

$\overline{D}_{all} =$ *All skip delay*

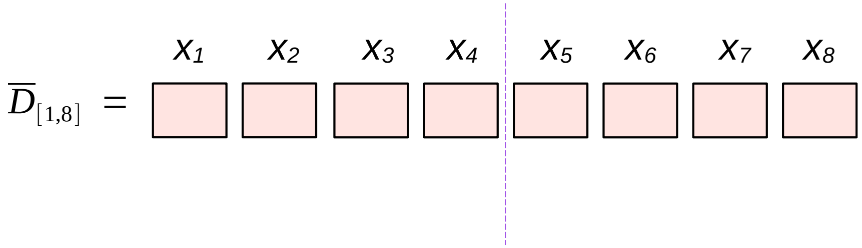
$\tilde{D}_{all} =$ *Comparable to all skip delay* $\leq$ D

$$D = max\{\overline{D}_{[1,8]}, \overline{D}_{[2,7]}, \overline{D}_{[3,6]}, \overline{D}_{[4,5]}\}$$

Max delay of all carry signals

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Maximum delays of carry signals (**r = 2k**)



$$D = \max_{s=1}^{r/2} \overline{D}_{[s,r-s+1]}$$

$$= \max_{s=1}^{k} \overline{D}_{[s,2k+1-s]}$$

$$= \max_{s=1}^{4} \overline{D}_{[s,9-s]}$$

Max delay of all carry signals

$$\overline{D}_{[1,r]} \le D$$
$$\overline{D}_{[2,r-1]} \le D$$
$$\vdots \qquad \vdots$$
$$\overline{D}_{[k,k+1]} \le D$$

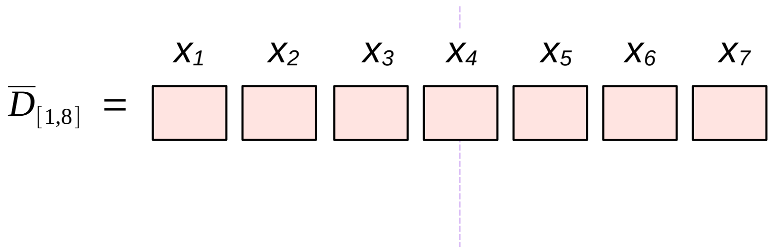
$$\overline{D}_{all} \le D$$

$$(m-1)T \le D$$

Lower bound of D

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Maximum delays of carry signals (**r = 2k+1**)



$$D = \max_{s=1}^{floor(r/2)} \overline{D}_{[s,r-s+1]}$$

$$= \max_{s=1}^{k} \overline{D}_{[s,2k+2-s]}$$

$$= \max_{s=1}^{3} \overline{D}_{[s,8-s]}$$

Max delay of all carry signals

$$\overline{D}_{[1,r]} \leq D$$
$$\overline{D}_{[2,r-1]} \leq D$$
$$\vdots \qquad \vdots$$
$$\overline{D}_{[k,k+1]} \leq D$$

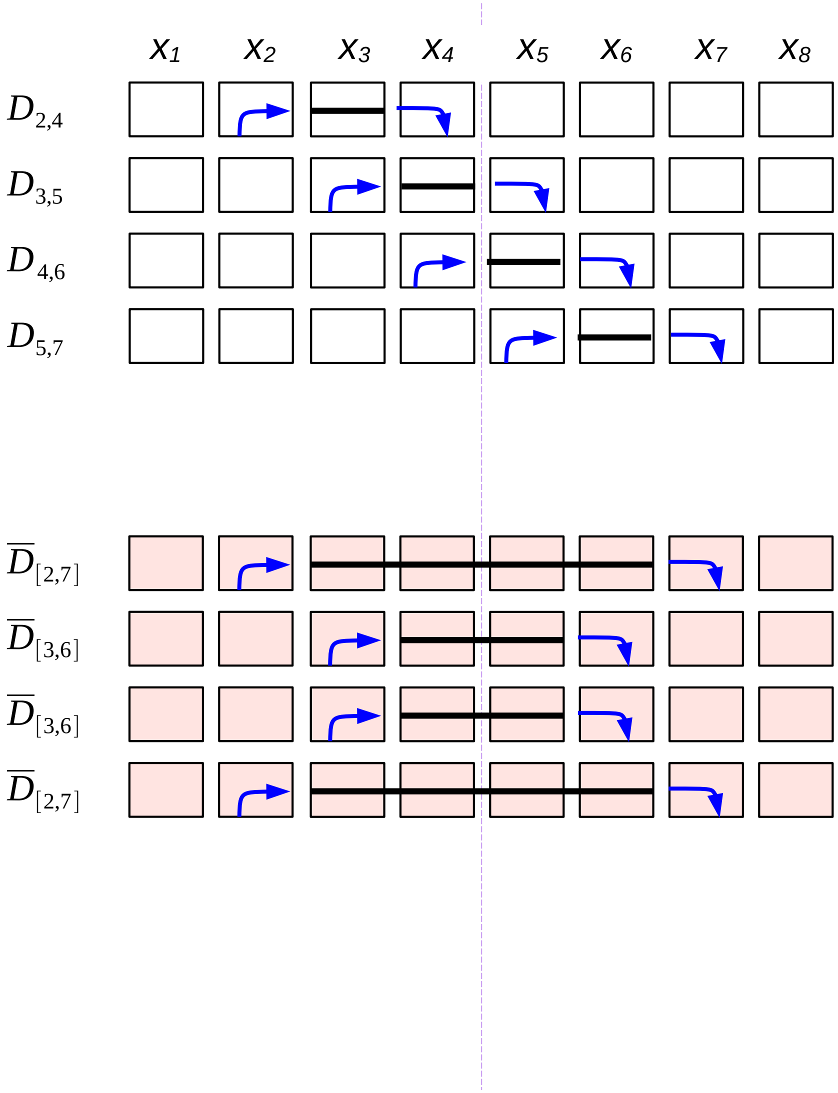
$$\widetilde{D}_{all} \leq D$$

$$(m-1)T \leq D$$

Lower bound of D

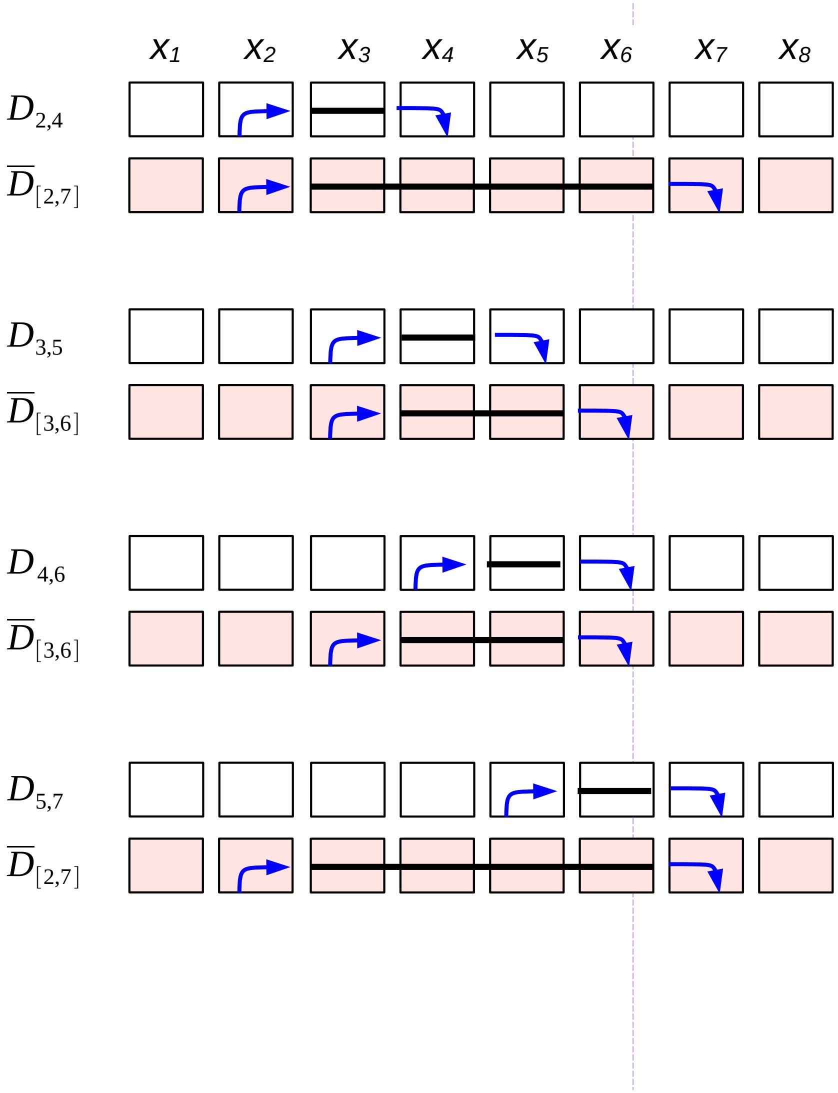
Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Example delays of carry signals (**r = 2k**) (1)



Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Example delays of carry signals (**r = 2k**) (2)



$$D_{2,4} \leq \overline{D}_{[2,7]}$$

$$D_{3,5} \leq \overline{D}_{[3,6]}$$

$$D_{4,6} \leq \overline{D}_{[3,6]}$$

$$D_{5,7} \leq \overline{D}_{[2,7]}$$

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Optimal division into groups (1-1)

Theorem 1

The scheme 2(i) – 2(iii) given above
for dividing the bits of a carry skip adder into groups
is optimal for $2 \le T \le 7$

dividing the bits into groups by the scheme 2(i) – 2(iii)
gives *m* groups

propagation time of a carry signal $\le mT$
the maximum propagation time = mT

the maximum delay = D
the optimal group size = m

$$(m-1)T \le D \le mT$$

(I) Let m be the smallest positive integer such that

$$n \le m + \frac{1}{2}mT + \frac{1}{4}m^2T + (1-(-1)^m)\frac{1}{8}T$$

(II) Let $y_i = min\{1+iT, 1+(m+1-i)T\}$,
$i = 1, \dots, m$

and construct a histogram
whose i-*th* column has height y_i

(III) the area of the histogram in (II) is

$$m + \frac{1}{2}mT + \frac{1}{4}m^2T + (1-(-1)^m)\frac{1}{8}T \ge n$$

so these are at least n unit squares
in the histogram starting with the first row,
shade in n of the squares, row by row
Let x_i denote the number of shaded squares
in column i of the histogram,
$i = 1, \dots, m$

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Optimal division into groups (1-2)

Assume

- the scheme by 2(i) – 2(iii) (*m* groups) is not optimal
- let ***D*** be the maximum delay
  corresponding to an optimal division of the bits into groups
- there are ***r*** groups in the optimal division.

m groups
- *not optimal division*
- ***D*** = maximum delay
- ***m****T* skip delay

r groups
- optimal division
- ***r****T* skip delay

Since a carry in signal to the least significant bit group
can skip over each group

we have $rT \le D \le mT$ so $r \le m$

skip delay
$rT \le D \le mT$

$r \le m$

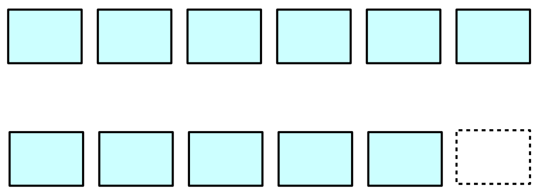
if m is not optimal, but r is *then* $mT \ge rT$ (smaller delay rT) *thus* $m \ge r$ (smaller *r* exists)

D *= max delay is assumed*
To be greater than all skip
delay ***rT*** *of the optimal division*

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Optimal division into groups (1-2)

If the optimal division gives *m* groups

m groups  mT

(*m-1*) groups $(m-1)T$

$$D \le mT$$

$$(m-1)T \le D$$

- when optimal group size = m
  the maximum delay $D_m \le mT$
- when optimal group size = (*m-1*)
  the maximum delay $D_{m-1} \le (m-1)T$

Normally, by 2(i) – 2(iii) (*m* groups) is optimal
and its maximum delay *D* is less than all skip delay mT

$D \le mT$

D = maximum delay

To prove this, first, negate that

- *m is <u>not</u> by the optimal division, but r is*
- *D is <u>greater</u> than all skip delay of the optimal division*

$rT \le D \le mT$

$r \le m$ ➡ $r < (m-1)$

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Optimal division into groups (1-2)

$rT \leq D \leq mT$ so $r \leq m$

Optimal division : r groups

$D' \leq$ all skip delay rT (r groups)

Non-optimal division : m groups

$D \leq$ all skip delay mT (m groups)

too many partitions m $r \leq m$

Assume max delay D is greater than all skip delay rT of the optimal division

if m is <u>not</u> optimal, <u>but</u> r is
then $mT \geq rT$ (smaller delay rT)
thus $m \geq r$ (smaller r exists)

D is max delay for m groups
D' is max delay for r groups
then $D' \leq rT \leq D \leq mT$

D = maximum delay

$rT \leq D \leq mT$

$r \leq m$ ➡ $r < (m-1)$

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Optimal division into groups (2)

we have $rT \leq D \leq mT$ so $r \leq m$

If $r = m$

then $D = mT$ ➡ $D = rT$ $rT = D$

If $r = m-1$, ($r < m$)

$D \geq (m-1)T$ ➡ $D \geq (m-1)T = rT$ $rT \leq D$

if $r < m-1$, ($r < m$) $r = m-2, m-3, m-3, \ ...$

$D \geq (m-1)T$ ➡ $D \geq (m-1)T > rT$ $rT < D$

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Optimal division into groups (3)

we have $rT \leq D \leq mT$ so $r \leq m$

If $r = m$ then $D = mT$
and the **theorem** holds by **lemma** 1

When the bits of a carry skip adder are grouped according to the scheme (i)-(iii), the maximum propagation time of a carry signal is mT

$(m\text{-}1)T \leq D \leq mT$

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Lemma 1 *When the bits of a carry skip adder are grouped according to the scheme (i)-(iii), the maximum propagation time of a carry signal is mT*

Theorem 1 *The scheme 2(i) – 2(iii) given above for dividing the bits of a carry skip adder into groups is optimal for* $2 \leq T \leq 7$

(5) $r = 2k$ $\quad X = 4 - T^2$

$$mT - D \leq T + \frac{-8(T/n) + 4}{\sqrt{4(T/n) + 8(T/n^2)} + \sqrt{4(T/n) + 4/n^2}}$$

$r = 2k+1$ $\quad X = 4$

$$mT - D \leq T + \frac{(T-2)^2/n}{\sqrt{4(T/n) + 4(T/n^2)} + \sqrt{4(T/n) + (T/n)^2 + 4/n^2}}$$

m groups – not optimal division
r groups – optimal division

D = maximum delay

$rT \leq D \leq mT$

$r \leq m$

Optimal division into groups (3)

If $r = m-1$, ($r < m$)
m and r have different parities and
it follows from (5)
that $mT - D \le T$ for $2 \le T \le 7$

so that $D \ge (m-1)T$
since $r = m-1$,
$D \ge (m-1)T = rT$ $\quad rT \le D$

This means that a signal which
skips over each of the r groups (rT)
has delay less than the maximum D.

$rT \le D \le mT$

m is <u>not</u> optimal division
r is optimal division

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Lemma 1 *When the bits of a carry skip adder are* <u>*grouped*</u> *according to the scheme (i)-(iii), the maximum propagation time of a carry signal is mT*

Theorem 1 *The scheme 2(i) – 2(iii) given above for dividing the bits of a carry skip adder into groups is optimal for 2 ≤ T ≤ 7*

(5) $r = 2k$ $\quad X = 4 - T^2$

$$mT - D \le T + \frac{-8(T/n) + 4}{\sqrt{4(T/n) + 8(T/n^2)} + \sqrt{4(T/n) + 4/n^2}}$$

$r = 2k+1$ $\quad X = 4$

$$mT - D \le T + \frac{(T-2)^2/n}{\sqrt{4(T/n) + 4(T/n^2)} + \sqrt{4(T/n) + (T/n)^2 + 4/n^2}}$$

m groups – not optimal division
r groups – optimal division

D = maximum delay

$rT \le D \le mT$

$r \le m$

Optimal division into groups (4)

Similarly,
if $r < m-1$, ($r < m$)

$(m-1)T \leq D$

since $r < m-1$,
$rT < (m-1)T \leq D$

so that a signal which skips over each group
has delay $rT < D$.

$rT < D \leq mT$

m is not optimal division
r is optimal division

m groups – not optimal division
r groups – optimal division

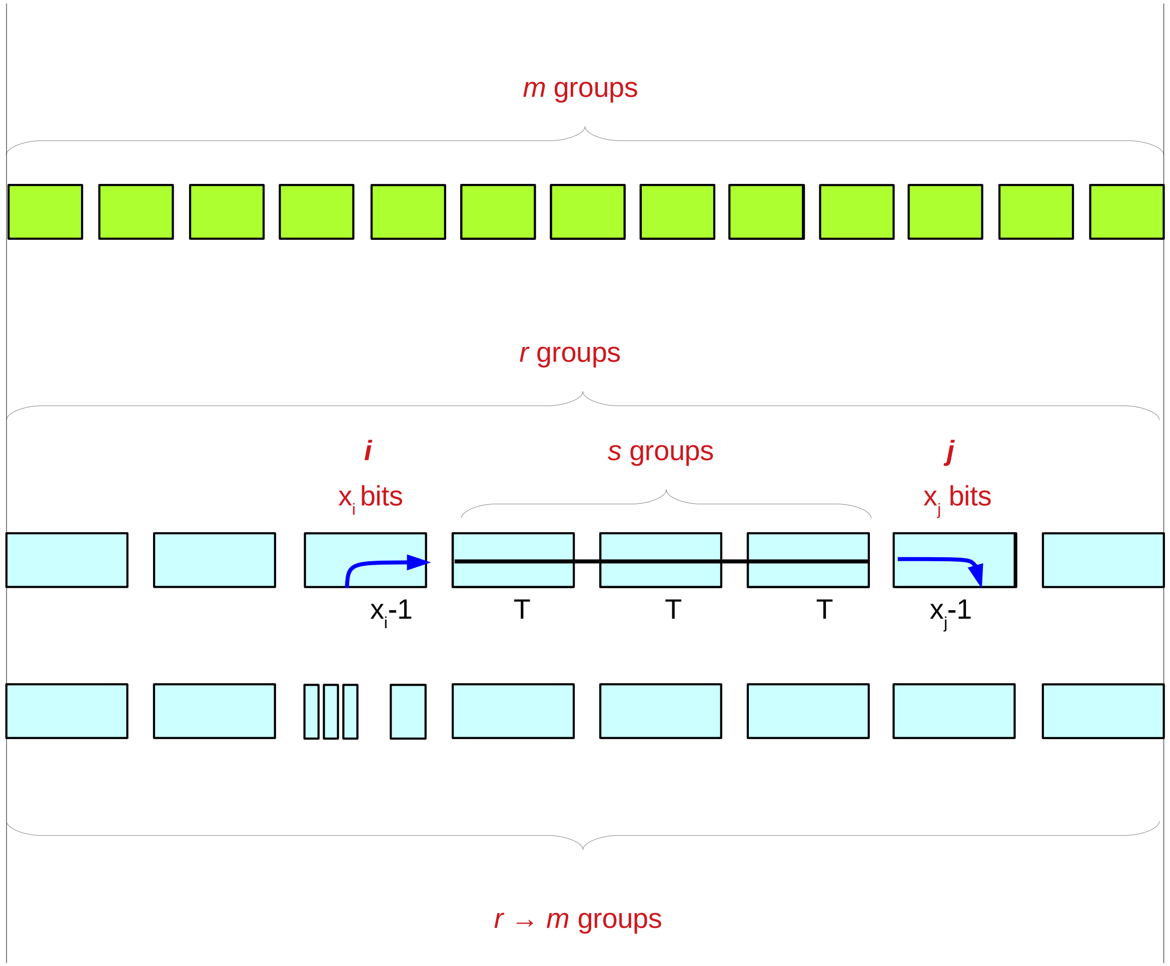
D = maximum delay

$rT \leq D \leq mT$

$r \leq m$

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Optimal division into groups (5)



m groups – not optimal division
r groups – optimal division

D = maximum delay

$rT \leq D \leq mT$

$r \leq m$

if m is not optimal, but r is

(((r +1) +1) +1) … → *m*

contradiction! *r must be m*

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

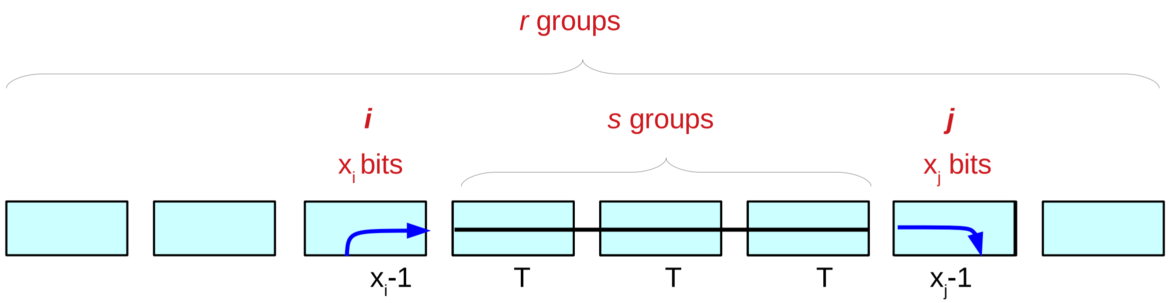
Optimal division into groups (5)

It follows that a signal with delay D

must start in a group i,
ripple to the end of group i,

then skip over $s < r$ groups and

either terminate, or ripple through
the first few bits of a group $j > i$.



m groups – not optimal division
r groups – optimal division

D = maximum delay

$rT \le D \le mT$

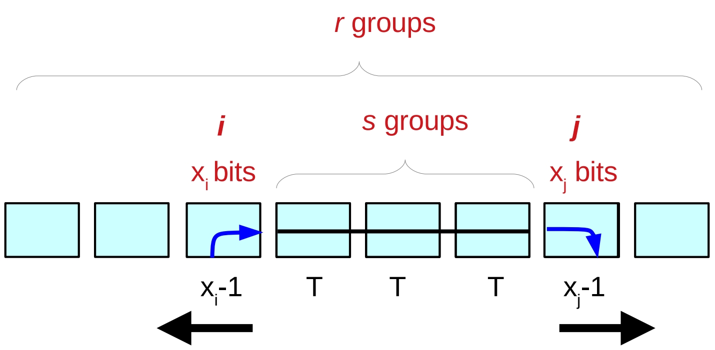
$r \le m$

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Optimal division into groups (6-1)

Let x_i and x_j denote
the lengths of the *i-th* and *j-th* groups respectively.

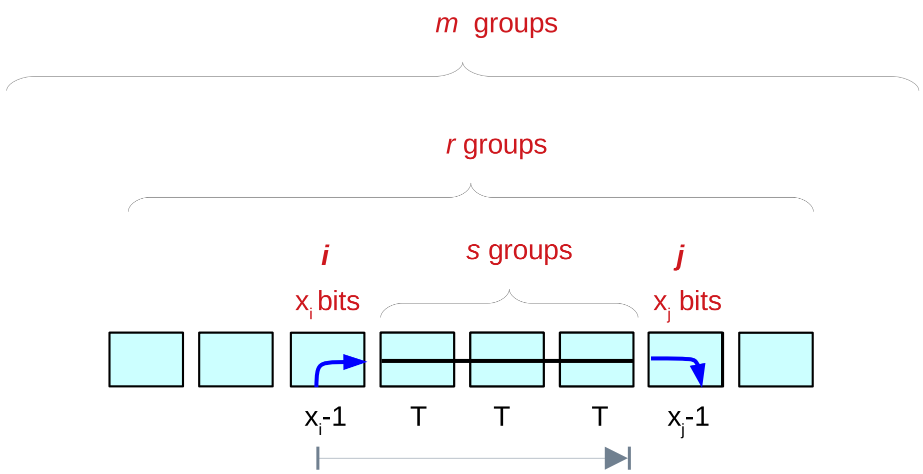
Assume that *i* is chosen as small as possible
and *j* as large as possible. (longer path)



Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Optimal division into groups (6-2)

A signal originating in group i, rippling to the end of this group i and then skipping over the next s group has delay $(x_i - 1) + sT$



$$D \le (x_i - 1) + sT$$
$$\le (x_i - 1) + (r-1)T$$
$$\le (x_i - 1) + (m-2)T.$$

$s < r$ groups ➡ $s \le (r-1)$

$r < m$ groups ➡ $r \le (m-1)$

$$D \le (x_i - 1) + sT$$
$$< (x_i - 1) + rT$$
$$< (x_i - 1) + mT.$$

if m is not optimal, but r is

$$s < r < m$$
$$s \le (r-1) < (m-1)$$
$$s \le (r-1) \le (m-2)$$

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Optimal division into groups (6-3)

$$D \le (x_i - 1) + sT$$
$$\le (x_i - 1) + (r - 1)T$$
$$\le (x_i - 1) + (m - 2)T$$

$$(m-1)T \le D$$
$$D \le (x_i - 1) + (m - 2)T$$

$$(m-1)T \le D \le (x_i - 1) + (m - 2)T$$

$$(m-1)T \le (x_i - 1) + (m - 2)T$$

Since $D \ge (m-1)T$
this implies that $x_i \ge T+1$

$$T \le (x_i - 1)$$

$$T + 1 \le x_i$$

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Optimal division into groups (7)

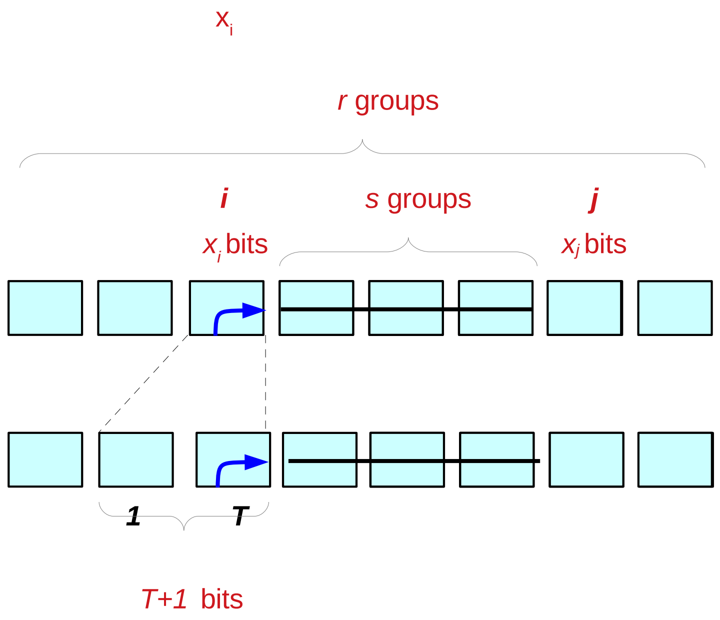
Divide group *i* into two groups
such that the group containing the msb has size *T*.

Since the *i*-th group is the first group
in which a signal having maximum delay can originate,

this subdivision does not increase
the delay of any carry signal of maximum delay

However, it increases the number of groups by **1**

$x_i \geq T+1$ *bits size*



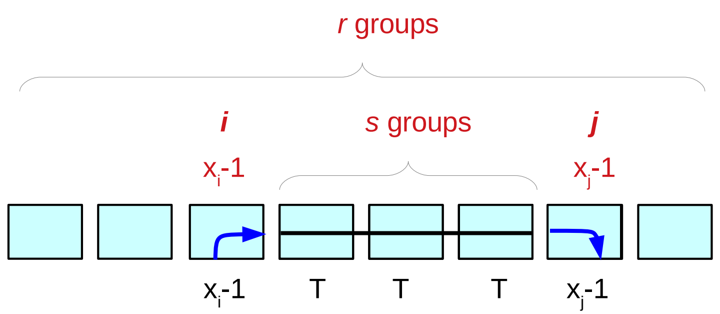
$$D \leq (x_i - 1) + sT$$
$$\leq (x_i - 1) + (r-1)T$$
$$\leq (x_i - 1) + (m-2)T.$$

$$(m-1)T < D$$
$$D < (x_i - 1) + (m-2)T$$
$$x_i \geq (T+1)$$

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Optimal division into groups (8-1)

Suppose now that a carry signal originates in a group *i*, ripples to its end, skips over $s \le r\text{-}2$ groups and finally ripples through the first few bits of a group *j* and terminates.



We then have

$$D \le (x_i - 1) + sT + (x_j - 1)$$
$$\le x_i + x_j - 2 + (m-3)T$$

So that either $x_i \ge T+1$ or $x_j \ge T+1$

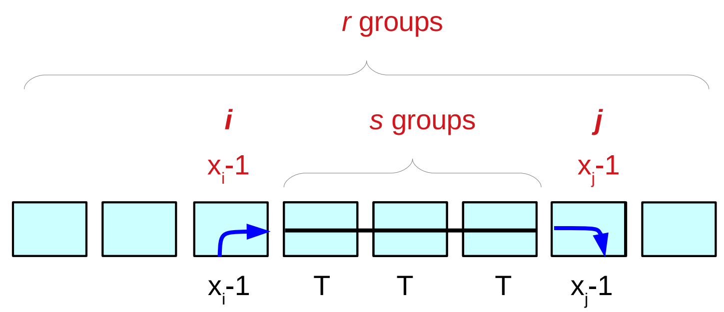
$s < r$ groups	$s < r$ groups
$s \le (r\text{-}1)$ groups	$s \le (r\text{-}2)$ groups
$r < m$ groups	$r < m$ groups
$r \le (m\text{-}1)$ groups	$r \le (m\text{-}2)$ groups

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Optimal division into groups (8-2)

$$s < r < m$$
$$s \le (r-1) < (m-1)$$
$$s \le (r-1) \le (m-2)$$

$$s \le (r-2) \le (m-3)$$



$$s \le (r-2)$$

$$D \le (x_i - 1) + sT + (x_j - 1)$$
$$\le x_i + x_j - 2 + (r-2)T$$
$$\le x_i + x_j - 2 + (m-3)T$$

$$(m-1)T < D$$
$$D < (x_i - 1) + (m-2)T \iff x_i \ge (T+1)$$
$$D < (x_j - 1) + (m-2)T \iff x_j \ge (T+1)$$

So that either $x_i \ge T+1$ or $x_j \ge T+1$

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

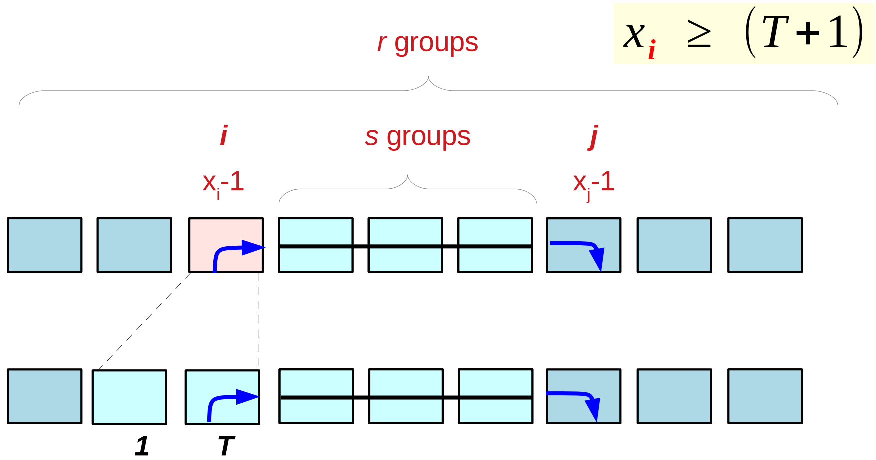
Optimal division into groups (9)

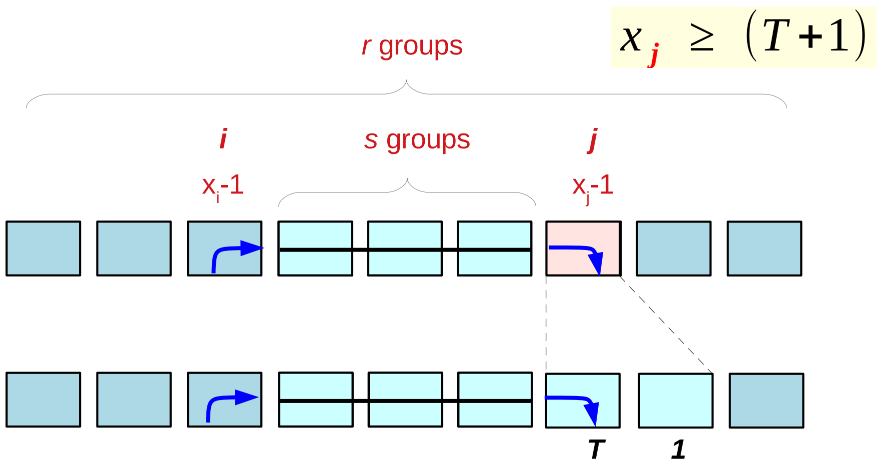
So that either $x_i \geq T+1$ or $x_j \geq T+1$

This means that we can subdivide
one of the groups i, j
without increasing D
not both of them

Continuing in this way, we can always increase
the number r of group in an optimal division
of a carry chain by 1
without increasing D if $r < m$

This means that we can arrive
at an optimal division of the carry chain
into m groups.





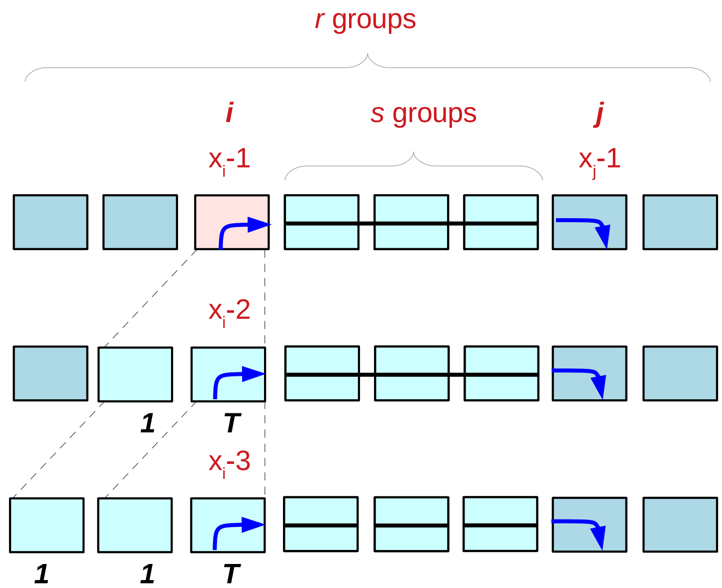
Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Optimal division into groups (9)

$x_i \geq (T+1) > (T+2) \cdots$

$x_i \geq (T+1)$

$x_i \geq (T+2)$



$$D \leq (x_i - 1) + sT$$
$$\leq (x_i - 1) + (r-1)T$$
$$\leq (x_i - 1) + (m-2)T$$

$$(m-1)T \leq D \leq (x_i - 1) + (m-2)T$$

if m is not optimal, but r is

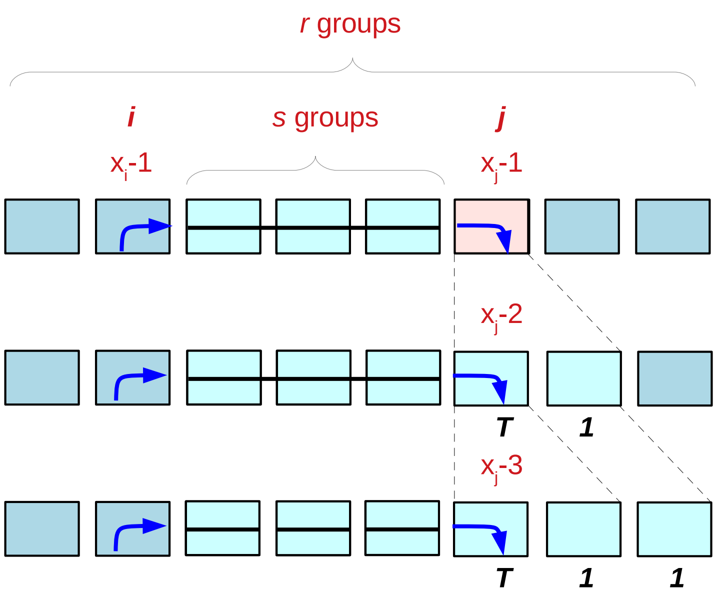
(((r +1) +1) +1) … → m

contradiction! m must be r

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Optimal division into groups (9)

$$x_j \geq (T+1) > (T+2) \cdots$$



$$x_j \geq (T+1)$$

$$x_j \geq (T+2)$$

$$D \leq (x_i - 1) + sT$$
$$\leq (x_i - 1) + (r-1)T$$
$$\leq (x_i - 1) + (m-2)T$$

if m is not optimal, but r is

$(((r\ \ +1)\ \ +1)\ \ +1)\ \ \ldots \longrightarrow m$

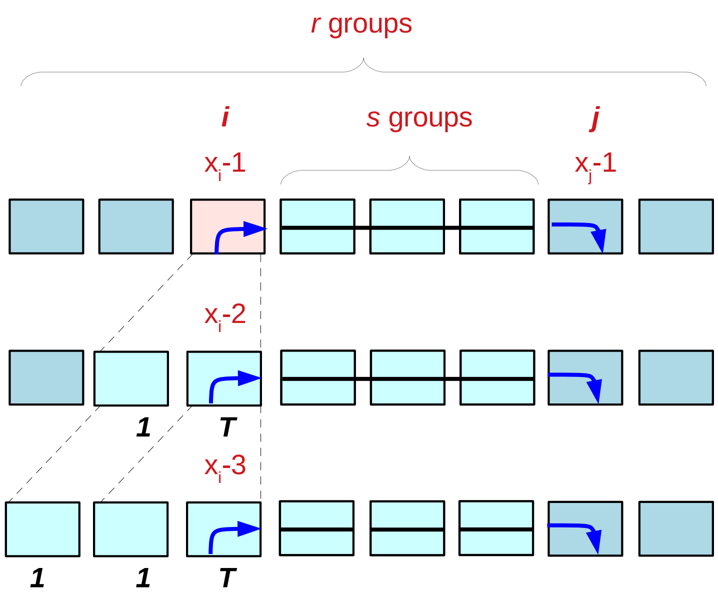
contradiction! m must be r

$$(m-1)T \leq D \leq (x_i - 1) + (m-2)T$$

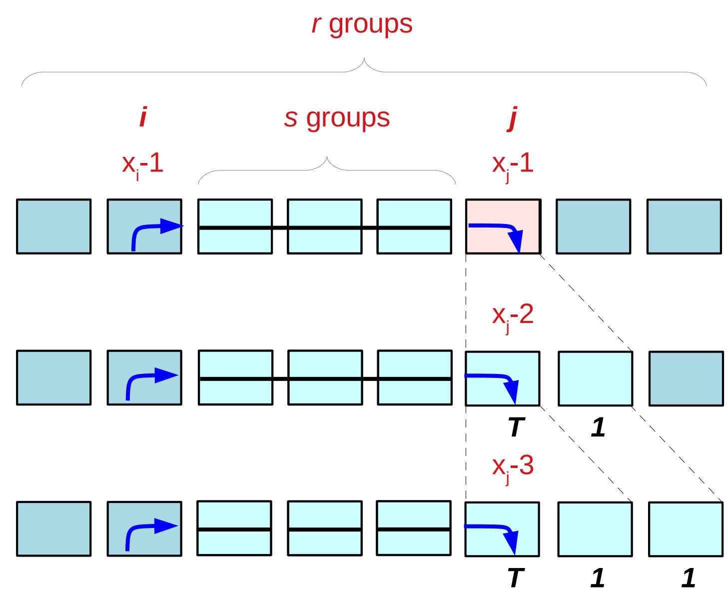
Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Optimal division into groups (9)

$x_i \geq (T+1) > (T+2) \cdots$



$x_j \geq (T+1) > (T+2) \cdots$



if m is not optimal, but r is

$(((r +1) +1) +1) \ldots \rightarrow m$

contradiction! m must be r

if m is not optimal, but r is

$(((r +1) +1) +1) \ldots \rightarrow m$

contradiction! m must be r

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Optimal division into groups (9)

Normally, by 2(i) – 2(iii) (m groups) is optimal
and its maximum delay D is less than all skip delay mT

$$D \le mT$$

To prove this, first, negate that

- *m is not by the optimal division, but r is*
- *D is greater than all skip delay of the optimal division*

Assume

- the scheme by 2(i) – 2(iii) (m groups) is not optimal
- let D be the maximum delay corresponding to an optimal division
- there are r groups in the optimal division.

(...(((r+1)+1)+1) … +1) → m : optimal

if m is not optimal, but r is

$(((r\ +1)\ +1)\ +1)\ \ldots \rightarrow m$

contradiction! m must be r

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers

Optimal division into groups (11)

We must then have $D \geq mT$
which, together with **Lemma 2**,
Implies $D = mT$

This completes the proof of the theorem

m groups – not optimal division
r groups – optimal division

D = maximum delay

$rT \leq D \leq mT$

$r \leq m$

Lemma 2
Let D denote the maximum delay
of a carry signal in a n bit carry skip adder
with group sizes chosen optimally.

$$(m-1)T \leq D \leq mT$$

Theorem 1

The scheme 2(i) – 2(iii) given above
for dividing the bits of a carry skip adder into groups
is optimal for $2 \leq T \leq 7$

Oklobdzija: High-Speed VLSI arithmetic units : adders and multipliers