## Graph Search (6A)

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## Graph Traversal

graph traversal (graph search) refers to the process of visiting (checking and/or updating) each vertex in a graph.

Such traversals are classified
by the order in which the vertices are visited.
Tree traversal is a special case of graph traversal.

A depth-first search (DFS)
is an algorithm for traversing a finite graph.
DFS visits the child vertices
before visiting the sibling vertices;
that is, it traverses the depth of any particular path before exploring its breadth.

A stack (often the program's call stack via recursion) is generally used when implementing the algorithm.

## DFS Backtrack

The algorithm begins with a chosen "root" vertex;
it then iteratively transitions from the current vertex to an adjacent, unvisited vertex, until it can no longer find an unexplored vertex to transition to from its current location.

The algorithm then backtracks along previously visited vertices, until it finds a vertex connected to yet more uncharted territory.

It will then proceed down the new path as it had before, backtracking as it encounters dead-ends, and ending only when the algorithm has backtracked past the original "root" vertex from the very first step.

## BFS

A breadth-first search (BFS) is another technique for traversing a finite graph.

BFS visits the neighbor vertices before visiting the child vertices
a queue is used in the search process
This algorithm is often used to find the shortest path from one vertex to another.

## Depth First Search Example


https://en.wikipedia.org/wiki/Graph_traversal

## Breadth First Search Example


https://en.wikipedia.org/wiki/Graph_traversal

## General Graph Search Algorithm - 1

```
Search(Start, isGoal, Criteria)
    insert(Start, Open);
    repeat
        if (empty(Open)) then return fail;
        select node from Open using Criteria;
        mark node as visited;
        if (isGoal(node)) then return node;
        for each child of node do
        if (child not already visited)
                then insert(child, Open);
```

DFS

## OPEN

## CLOSED "visited"


unvisited children : 1, 3


## BFS

## OPEN

CLOSED marked "visited"

unvisited children : 1, 3


## Possible duplication

## DFS Stack


possible duplication - not yet expanded

## BFS Queue


possible duplication

- not yet expanded


## Must check before expansion

## DFS Stack

must check if the selected node is already "visited"

possible duplication

## BFS Queue

must check if the selected node is already "visited"

possible duplication

## General Graph Search Algorithm - 1

```
Search(Start, isGoal, Criteria)
    insert(Start, Open);
    repeat
        if (empty(Open)) then return fail;
        select node from Open using Criteria;
        mark node as visited;
        if (isGoal(node)) then return node;
        for each child of node do
        if (child not already visited)
                then insert(child, Open);
```


## DFS-1 (Depth First Search)

Open list - use a stack
Select with Criteria - pop
DFS(Start, isGoal)
push(Start, Open); // push
repeat
if (empty(Open)) then return fail; node := pop(Open); // pop mark node as visited; if (isGoal(node)) then return node; for each child of node do if (child not already visited) then push(child, Open); // push

## DFS-1 Example (1)


pop
a


pop



## DFS-1 Example (2)



## BFS-1 (Breadth First Search)

Open list - use a FIFO
Select with Criteria - dequeue

```
BFS(Start, isGoal)
    enqueue(Start, Open); // enqueue
    repeat
    if (empty(Open)) then return fail;
    node := dequeue(Open); // dequeue
    mark node as visited;
    if (isGoal(node)) then return node;
    for each child of node do
                if (child not already visited) then
                    enqueue(child, Open); // enqueue
```

BFS-1 Example (1)


## BFS-1 Example (2)



## General Graph Search Algorithm - 2

```
Initialize as follows:
    unmark all nodes in N;
    mark node s;
    pred(s) = 0; {that is, it has no predecessor}
    LIST = {s}
while LIST f ø do
    select a node i in LIST;
    if node j is incident to an admissible arc (i,j) then
        mark node j;
        pred(j) := i;
        add node j to the end of LIST;
    else
        delete node i from LIST
```

DFS : select the last node i in LIST;


## Admissible arc

$\operatorname{pred}(\mathrm{j})$ is a node that precedes j on some path from s ;
A node is either marked or unmarked.
Initially only node s is marked.
If a node is marked, it is reachable from node $s$.
An arc $(\mathrm{i}, \mathrm{j}) \in \mathrm{A}$ is admissible

if node $i$ is marked and $j$ is not.

Before a node is added into LIST, the node is marked

LIST contains only the marked nodes thus, the selected node $\mathbf{i}$ is marked already

The node $\mathbf{j}$ incident to the admissible $\operatorname{arc}(\mathbf{i}, \mathbf{j})$ must be unmarked


This node $\mathbf{j}$ is marked and added into LIST
In this way, LIST contains only marked and non-repeating nodes

Check before inserting

## DFS-2

```
Initialize as follows:
    unmark all nodes in N;
    mark node s;
    pred(s) = 0; {that is, it has no predecessor}
    push s onto LIST
while LIST = ø do
    pop a node i from LIST;
    if node j is incident to an admissible arc (i,j) then
        mark node j;
        pred(j) := i;
        push(node j) onto LIST;
    else
        delete node i from LIST
```


## DFS-2 Example (1)



## DFS-2 Example (2)



## BFS-2

```
Initialize as follows:
    unmark all nodes in N;
    mark node s;
    pred(s) = 0; {that is, it has no predecessor}
    enqueue s onto LIST
while LIST = ø do
    dequeue node i from LIST;
    if node j is incident to an admissible arc (i,j) then
        mark node j;
        pred(j) := i;
        enqueue node j onto LIST;
    else
    delete node i from LIST
```

BFS-2 Example (1)

https://en.wikipedia.org/wiki/Graph_traversal

BFS-2 Example (2)


## DFS Pseudocode

1 procedure DFS(G, v):
2 label v as explored
3 for all edges e in G.incidentEdges(v) do
4 if edge $e$ is unexplored then
$5 \quad \mathrm{w} \leftarrow \mathrm{G} . \operatorname{adjacentVertex}(\mathrm{v}, \mathrm{e})$
6 if vertex $w$ is unexplored then
8 recursively call DFS(G, w)
9 else
10 label e as a back edge

## BFS Pseudocode

1 procedure BFS(G, v):
2 create a queue Q
3 enqueue $v$ onto Q
4 mark v
5 while Q is not empty:
$6 \quad \mathrm{t} \leftarrow \mathrm{Q}$.dequeue()
7 if t is what we are looking for:
8 return t
9 for all edges e in G.adjacentEdges(t) do

12
13
14
15 enqueue o onto Q
16 return null

## References

[1] http://en.wikipedia.org/
[2]

